

Graph Theory: Basics

Introduction to network Science

Instructor: Michele Starnini — <https://github.com/chatox/networks-science-course>

Network Science or Graph Theory?

	Field	When	What	How
Graphs	mathematics (computer science)	1960-70	Theory	Structural properties
Networks	physics	2000-present	Applications	Complex systems, Dynamical processes

Nobody cares! Completely equivalent

What is (modern) Network Science

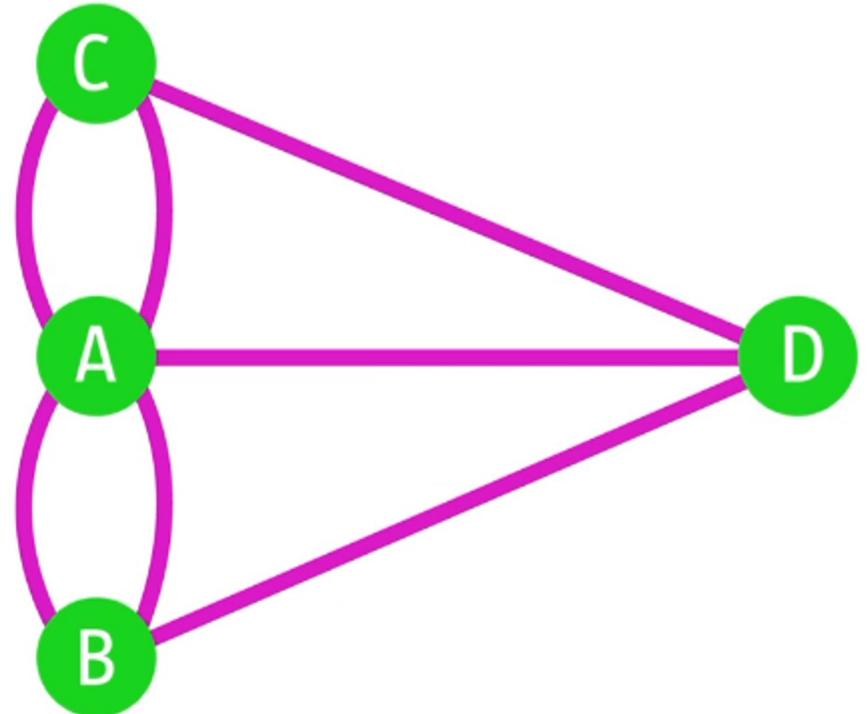
- **Mathematical formalism** from graph theory
- New insights from **statistical physics**
- Recent availability of very **large databases** on networked systems
- **Computational resources** to analyze them

Contents

- Directed, weighted, ... many kind of graphs!
- Density & Sparsity
- Degrees
- Adjacency matrices

Notation for a graph

- $G = (V, E)$
 - V : nodes or vertices
 - E : links or edges
- $|V| = N$ size of graph
- $|E| = L$ (or E) number of links



Subgraph

- Given $G = (V, E)$
- A **subgraph** induced by a nodeset S is the graph $G'=(S,F)$ defined by:
 - nodes in S
 - edges in $F = \{ (u,v) \in E \text{ s.t. } u \in S \text{ and } v \in S \}$

Directed vs undirected graphs

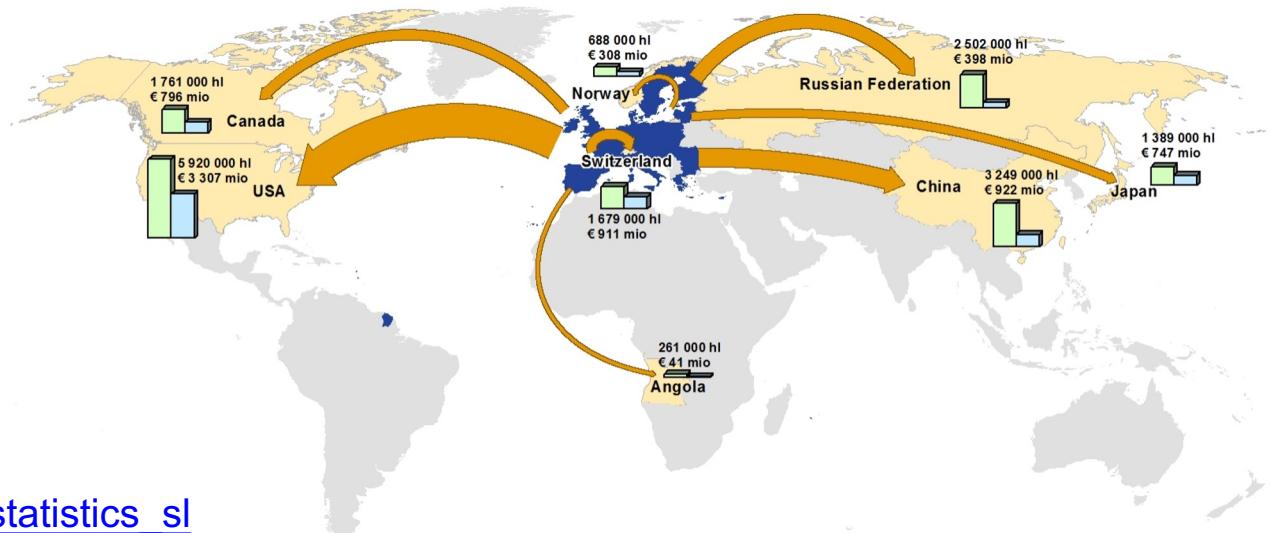
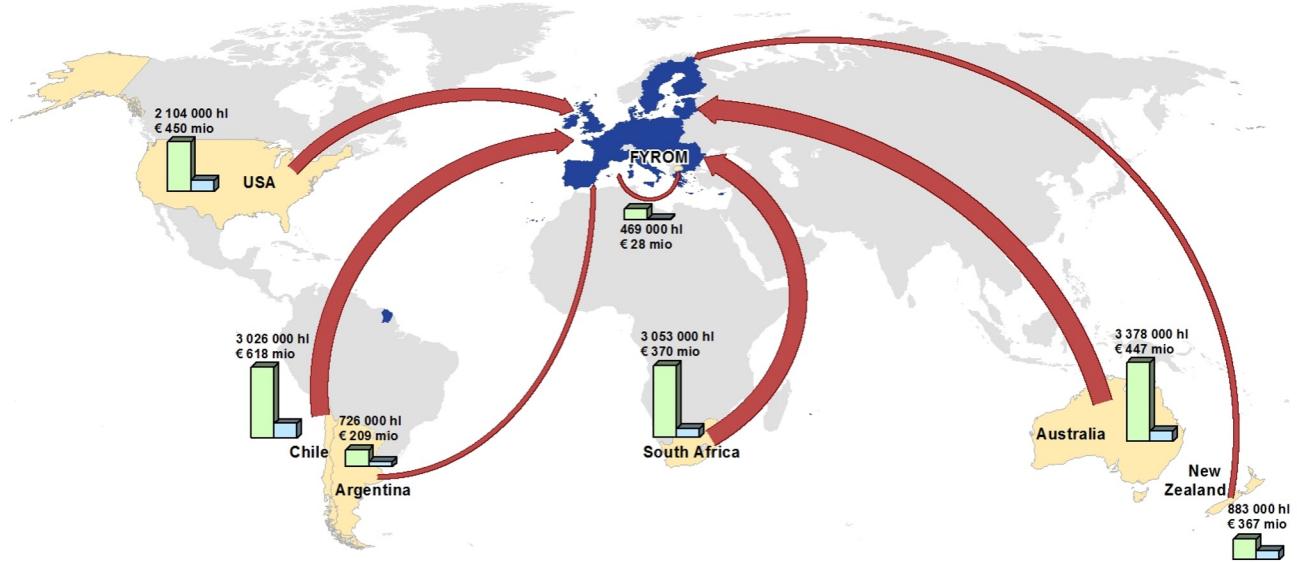
- In an undirected graph
 - E is a symmetric relation
$$(u, v) \in E \Rightarrow (v, u) \in E$$
- In a directed graph, also known as “digraph”
 - E is not a symmetric relation
$$(u, v) \in E \not\Rightarrow (v, u) \in E$$

Weighted vs unweighted graphs

- In a weighted graph edges have **weights** denoting the strength or importance of a connection
- When drawing, typically larger weights are drawn with thicker lines

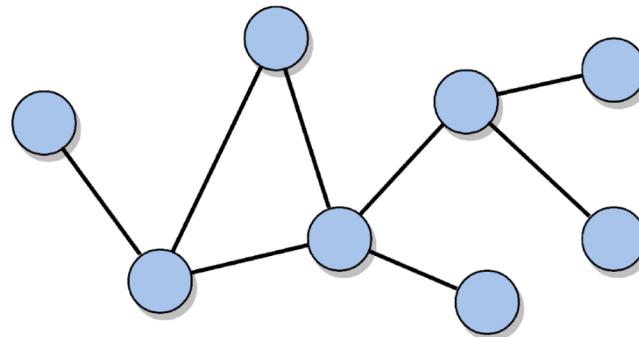
Weighted networks

EU imports (top)
and exports (bottom)
of wine

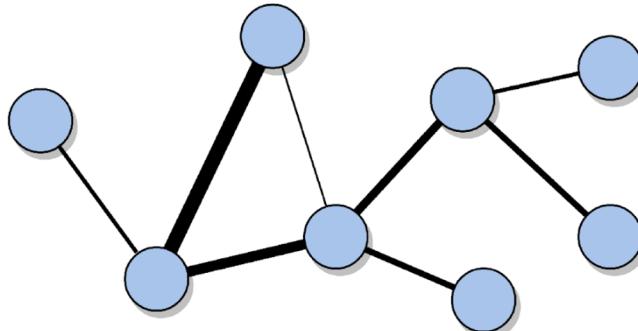


Undirected

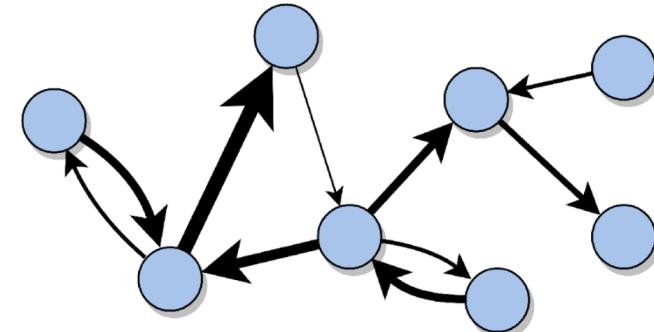
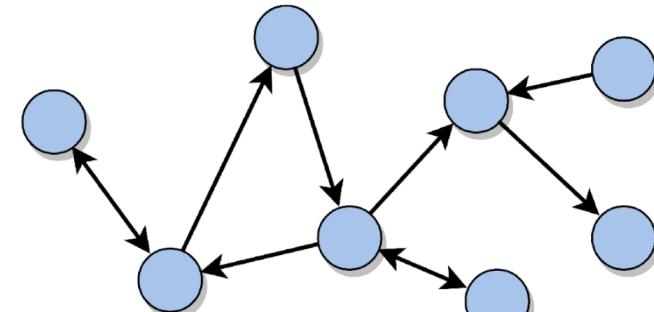
Unweighted



Weighted



Directed



Example graphs we will use

Network	N	E
Zachary's Karate Club (karate.gml)	34	78
Game of Thrones (got-relationships.csv)	84	216
US companies ownership	1351	6721
Marvel comics (hero-network.csv)	6K	167K

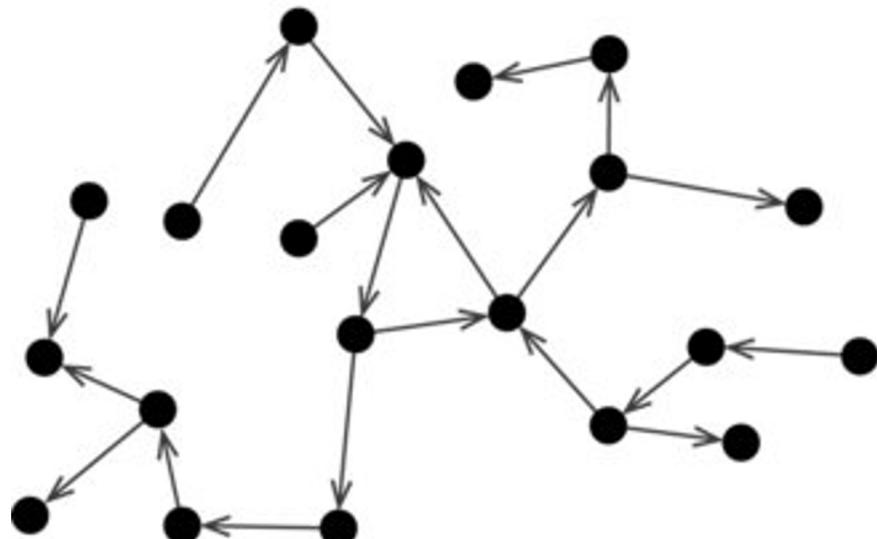
<https://github.com/chatox/networks-science-course/tree/master/practicum/data>

Density and sparsity

- Network size N = number of nodes
- L = number of links
- Maximum possible number of links:
- Density: $d = \frac{L}{L_{max}} = \frac{2L}{N(N - 1)}$
- The network is **sparse** if $d \ll 1$

$$L_{max} = \binom{N}{2} = \frac{N(N - 1)}{2}$$

Sparse network



Dense network



[[Source](#)]

Example: Facebook

Rough orders-of-magnitude approximations:

- $N \approx 10$
- $L \approx 10^3 \times N$
- $d \approx L / N^2 \approx 10^3 N / N^2 \approx 10^3 / 10^9 = 10^{-6}$
- Most (but not all) real-world networks are similarly sparse because the number of links scales proportionally to N , whereas the maximum scales with N^2

Complex networks are sparse

- Theoretically $L_{\max} = \binom{N}{2} = \frac{N(N - 1)}{2}$
- Most real networks are sparse, i.e.,
$$L \ll L_{\max}$$

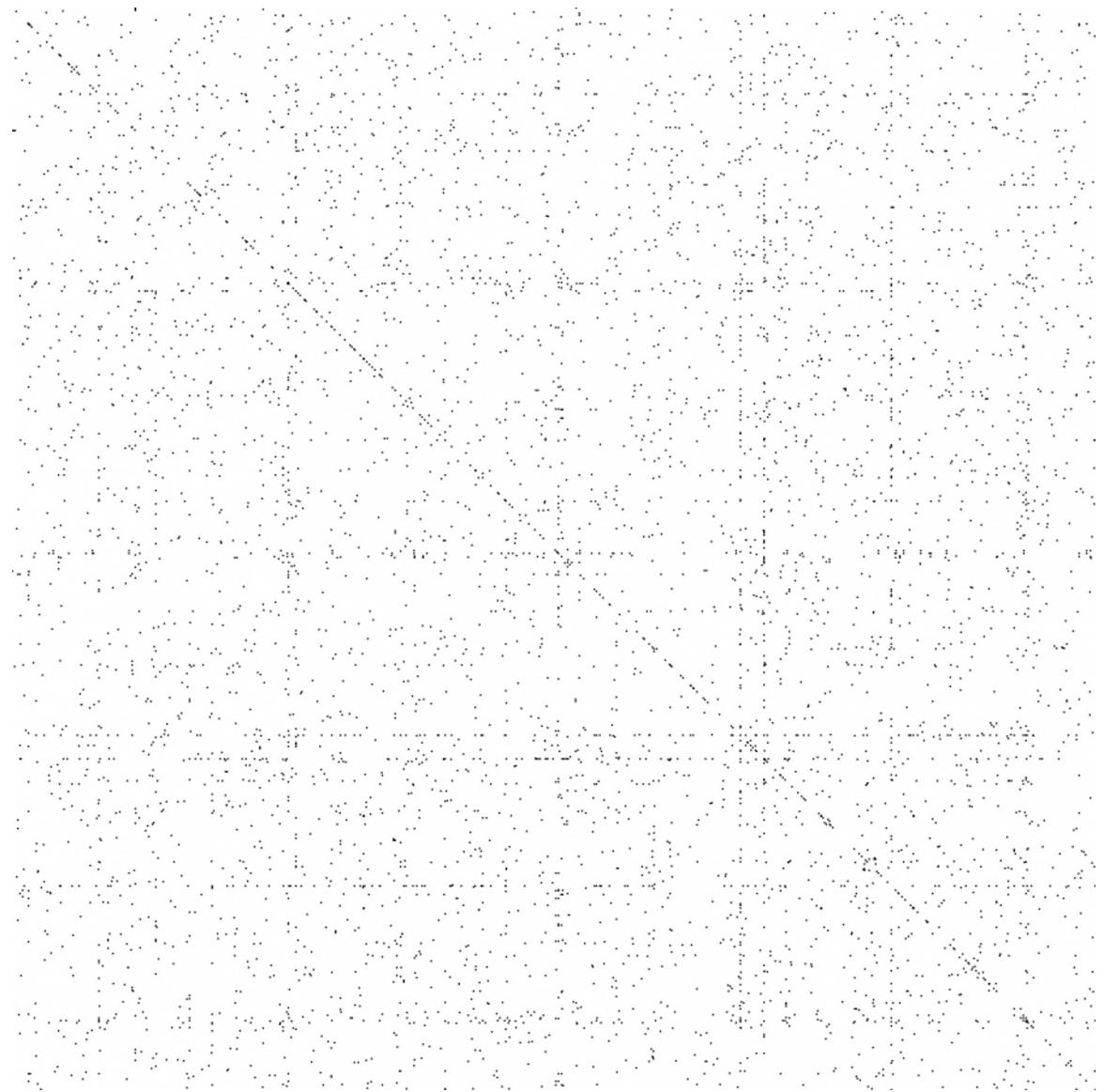
L is the number of links in the network, N is the number of nodes on it

How sparse are some networks?

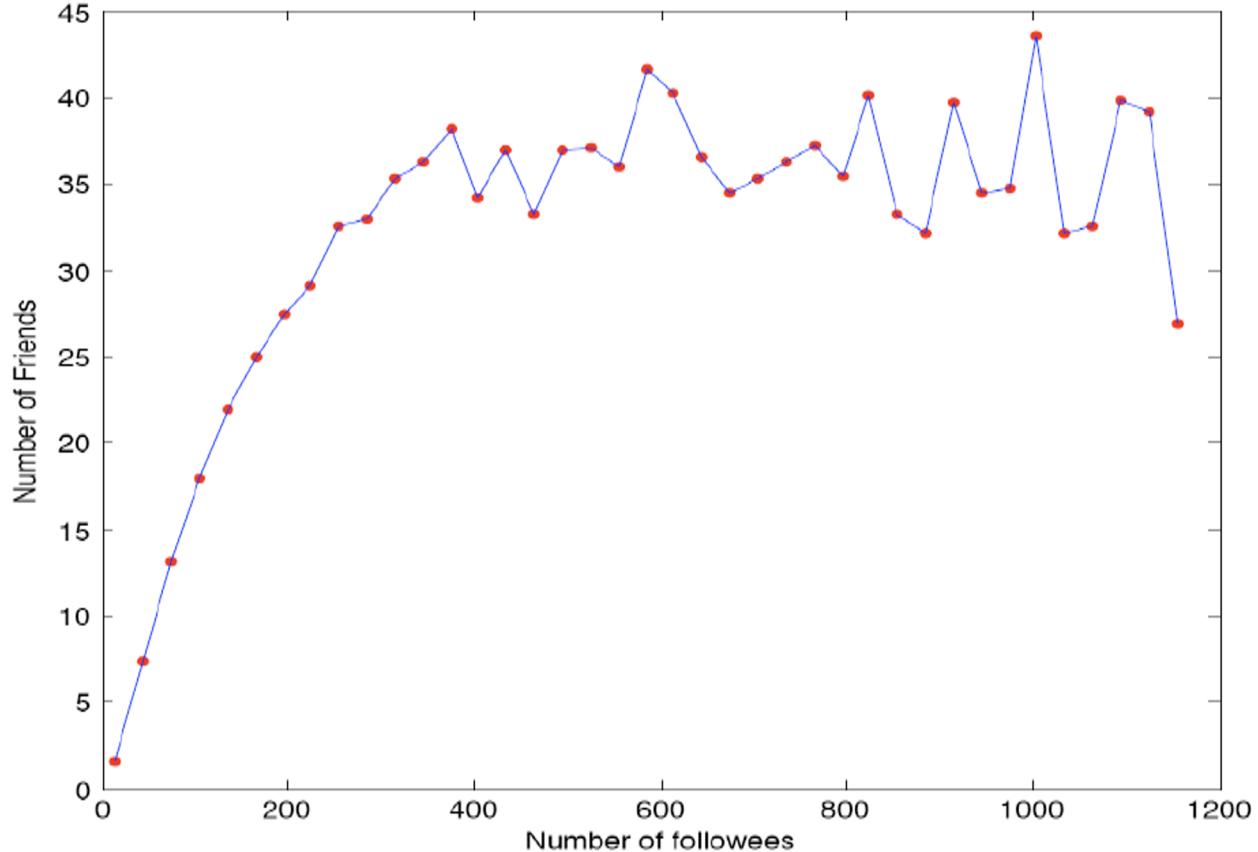
Network	$ V $	$ E $	Max $ E $
Zachary's Karate Club	34	78	561
Game of Thrones	84	216	3496
US companies ownership	1351	6721	911K
Marvel comics	6K	167K	17M

Example:
protein
interaction
network

($N=2K$, $L=3K$)



Example: people you follow on Twitter (followees) vs people you have sent at least two messages to (“friends”)



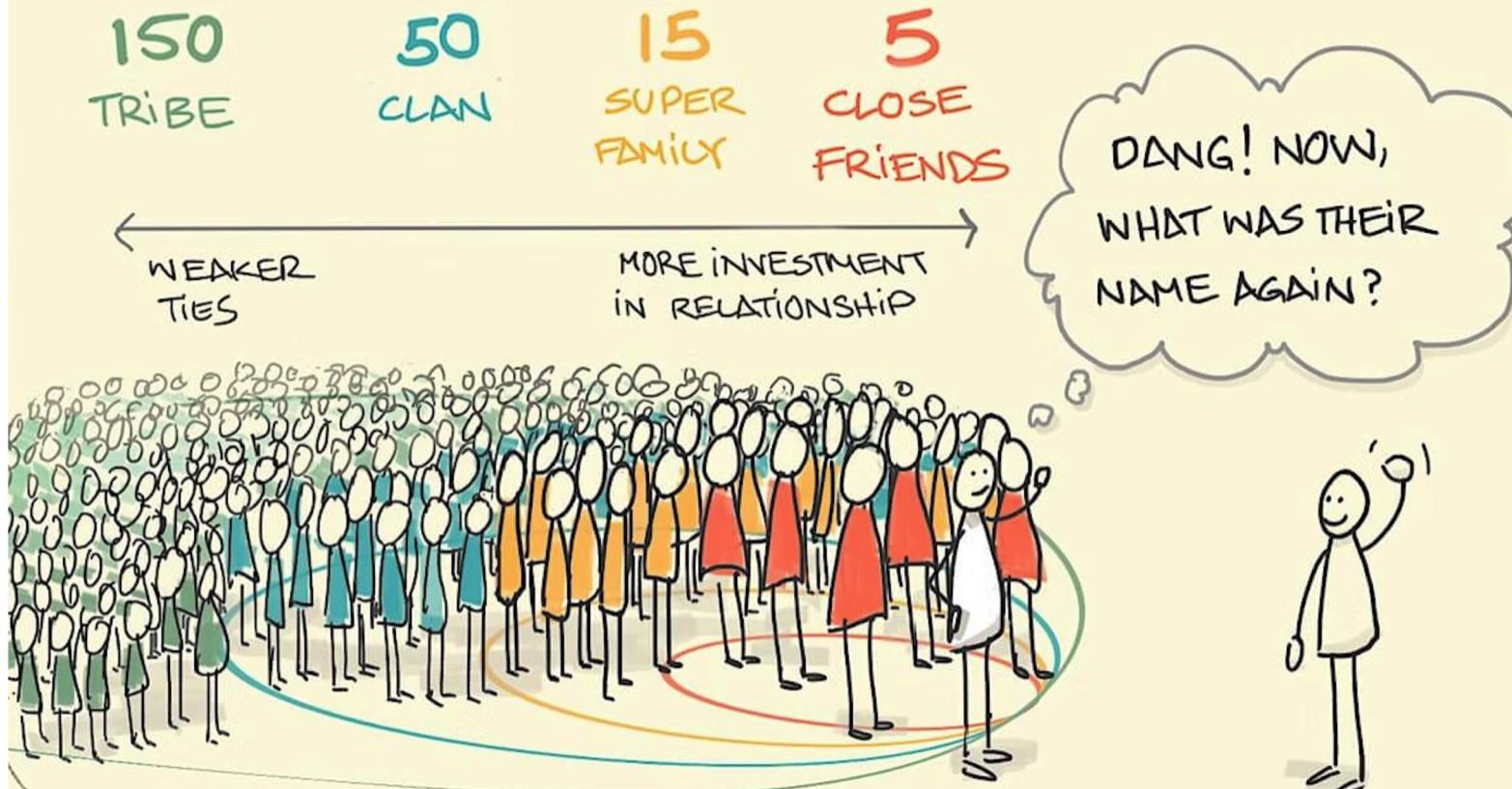
Why are networks sparse?

- Different mechanisms, think about it from the node perspective:
 - How many items **could** the node be connected to
 - Would it be **realistic** to connect to a large fraction of them?
- In social networks, Dunbar's number ($\simeq 150$)

DUNBAR'S NUMBER : 150

[Sketchplanations]

TYPICAL NUMBER OF PEOPLE WE CAN KEEP TRACK OF AND
CONSIDER PART OF OUR ONGOING SOCIAL NETWORK



Degree

Node i has degree k_i

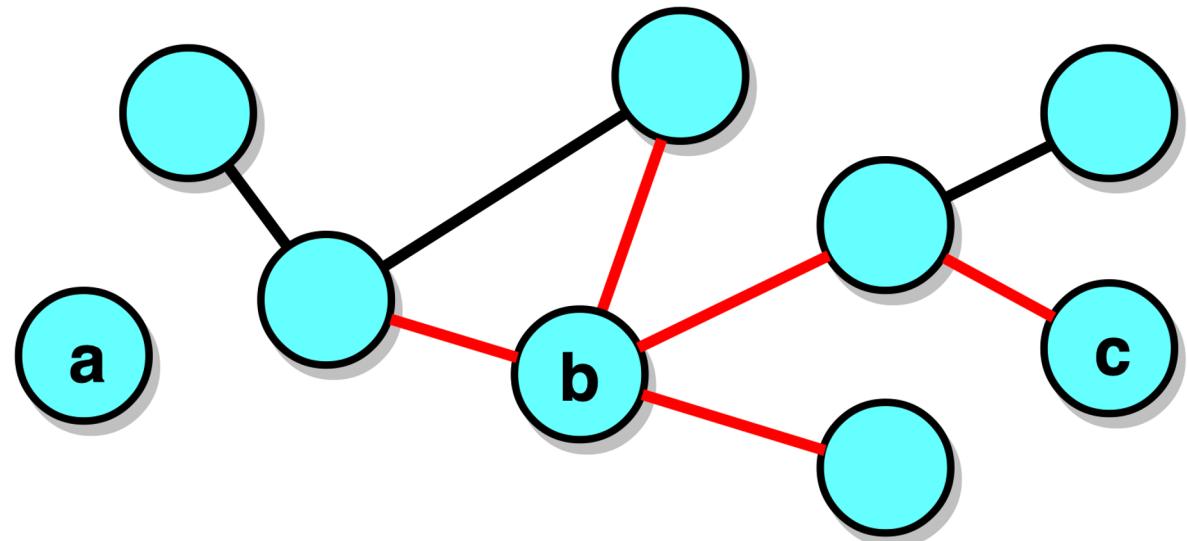
This is the number of links incident on this node

The total number of links L is given by

$$L = \frac{1}{2} \sum_{i=1}^N k_i$$

Average degree

$$\langle k \rangle = \frac{1}{N} \sum_{i=1}^N k_i = \frac{2L}{N}$$



Average degree & density

$$d = \frac{\langle k \rangle}{N - 1} = \frac{\langle k \rangle}{k_{max}}$$

- Let us revisit the Facebook example
- Since N is very large, $N - 1$ can be approximated by N

$$d = \frac{\langle k \rangle}{N - 1} \approx \frac{\langle k \rangle}{N} \approx \frac{10^3}{10^9} = 10^{-6}$$

In directed graphs

We distinguish **in-degree** from **out-degree**

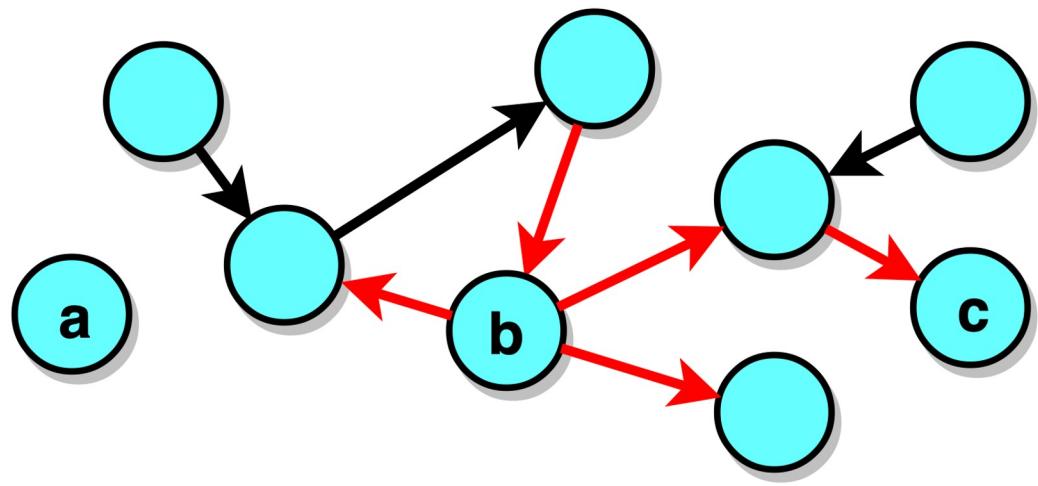
Incoming and outgoing links, respectively

Degree is the sum of both

$$k_i = k_i^{\text{in}} + k_i^{\text{out}}$$

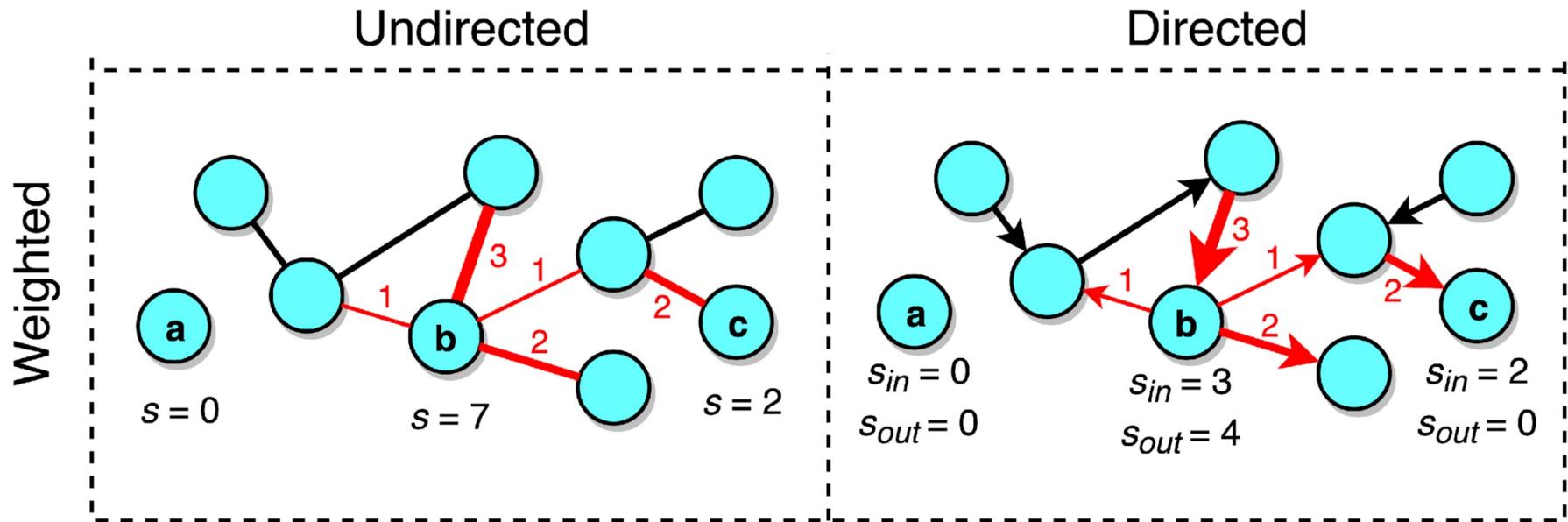
Counting total number of links:

$$L = \sum_{i=1}^N k_i^{\text{in}} = \sum_{i=1}^N k_i^{\text{out}}$$



In weighted graphs

- We speak of “weighted degree” or “strength”

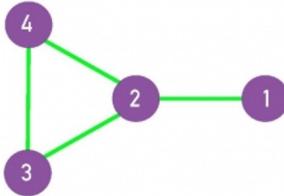


Degree distribution

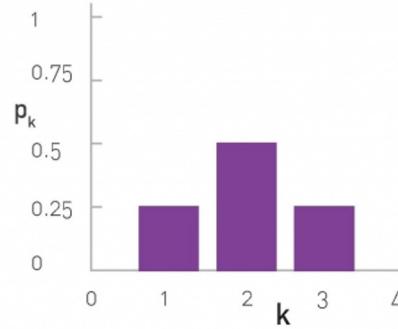
- If there are N_k nodes with degree k
- The **degree distribution** is given by $p_k = \frac{N_k}{N}$
- The average degree is then $\langle k \rangle = \sum_{k=0}^{\infty} kp_k$

Degree distribution; two toy graphs

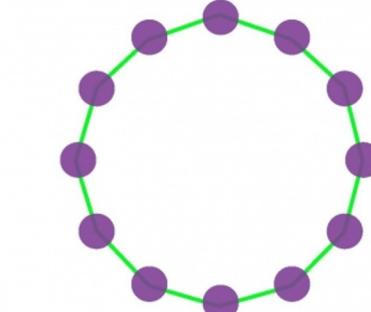
a.



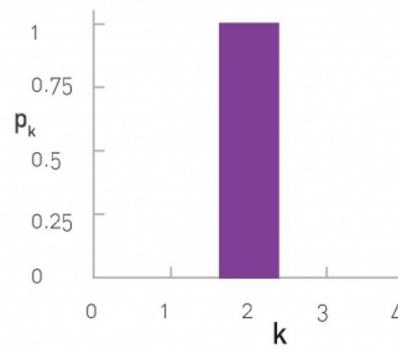
b.



c.

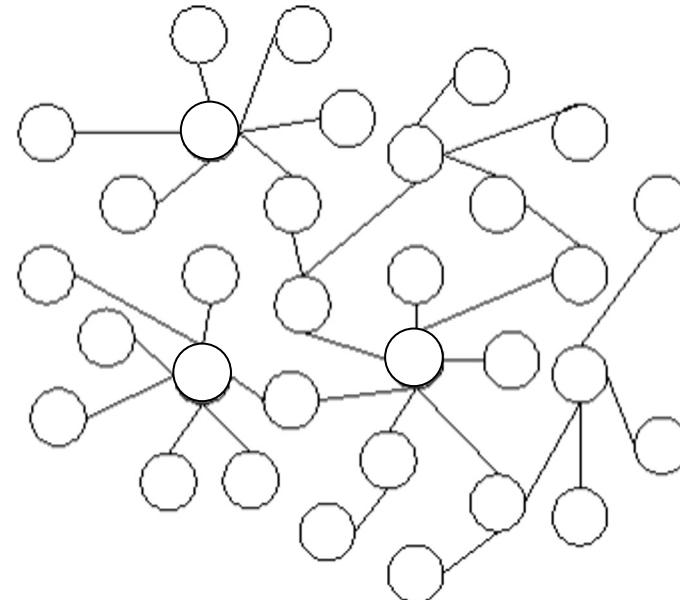
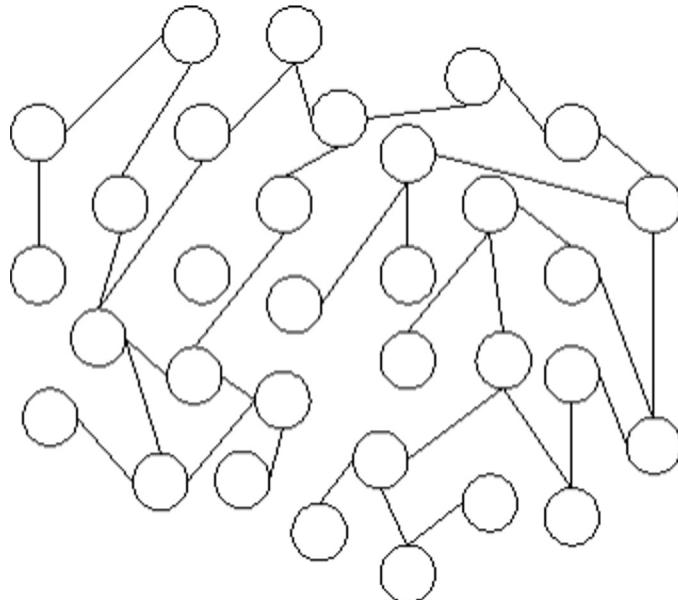


d.



Exercise

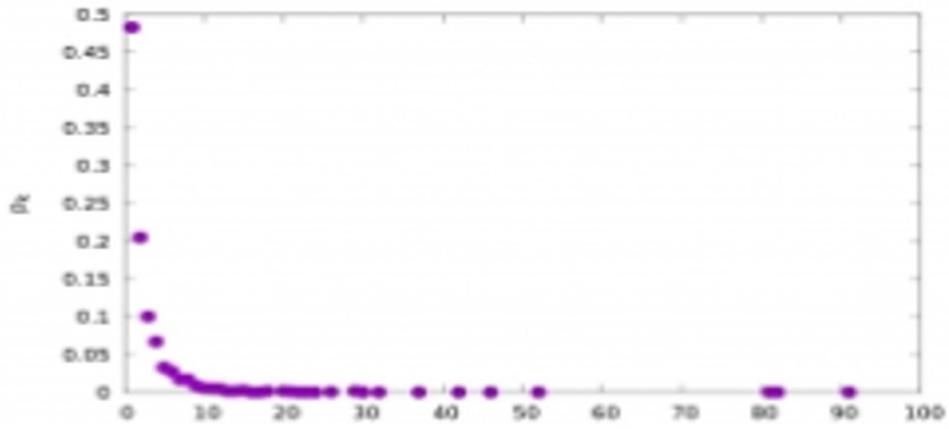
Draw the degree distribution of these graphs



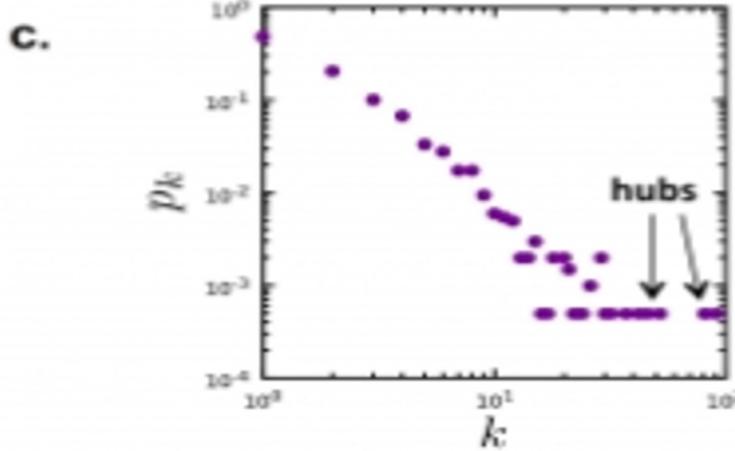
Spreadsheet links: <https://upfbarcelona.padlet.org/chato/shyq9m6f2g2dh1bw>



Real graphs



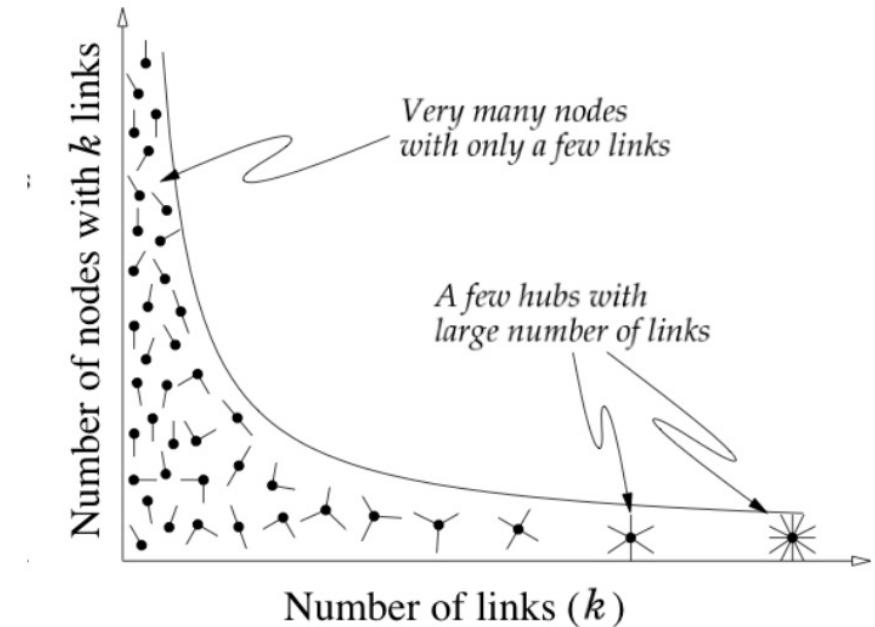
Linear
scale



Log-log
scale

Degree distribution $P(k)$

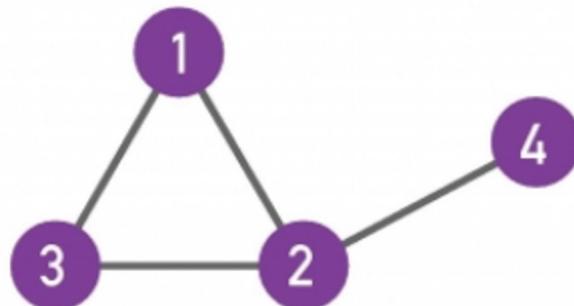
- Most important statistical property of networks
- Nodes with large degree = hubs,
nodes with small degree = leaves
- When hubs are present,
the $P(k)$ is heterogeneous
- Most real (complex) networks
are heterogeneous



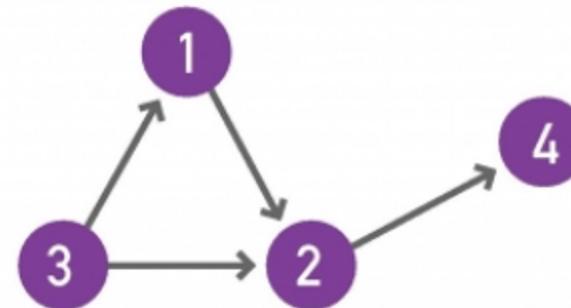
Adjacency matrix

- A is the **adjacency matrix** of $G = (V, E)$ iff:
 - A has $|V|$ rows and $|V|$ columns
 - $A_{ij} = 1$ if $(i,j) \in E$
 - $A_{ij} = 0$ if $(i,j) \notin E$
- **A_{ij} always means row i, column j**
 - Sometimes Barabási's book has this wrong

Examples



Undirected graph

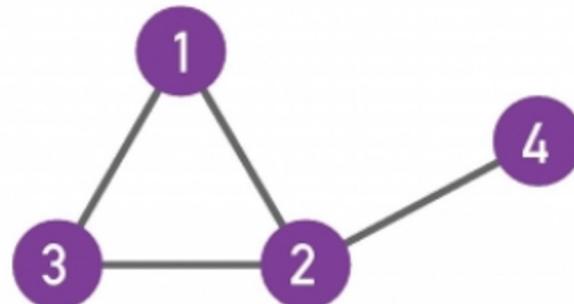


Directed graph

$$A_{ij} = \begin{matrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{matrix}$$

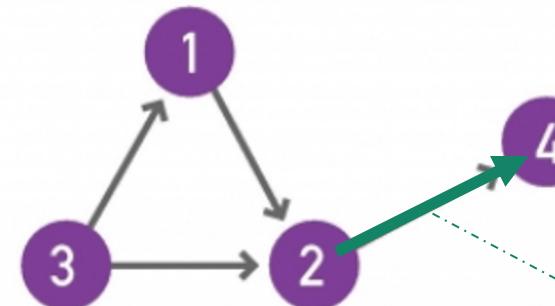
$$A_{ij} = \begin{matrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{matrix}$$

A_{ij} always means row i , column j



Undirected graph

$$A_{ij} = \begin{matrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{matrix}$$



Directed graph

$$A_{ij} = \begin{matrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{matrix}$$

Row 2
Column 4

Properties of adjacency matrices

- G is undirected $\Leftrightarrow A$ is symmetric
- G has a **self-loop**
 $\Leftrightarrow A$ has a non-zero element in the diagonal
- G is **complete** $\Leftrightarrow A_{ij} \neq 0$ (except if $i=j$)

Quick Exercise

- In terms of A, what is the expression for:

$$k_i^{\text{in}} =$$

$$k_i^{\text{out}} =$$

Weighted networks

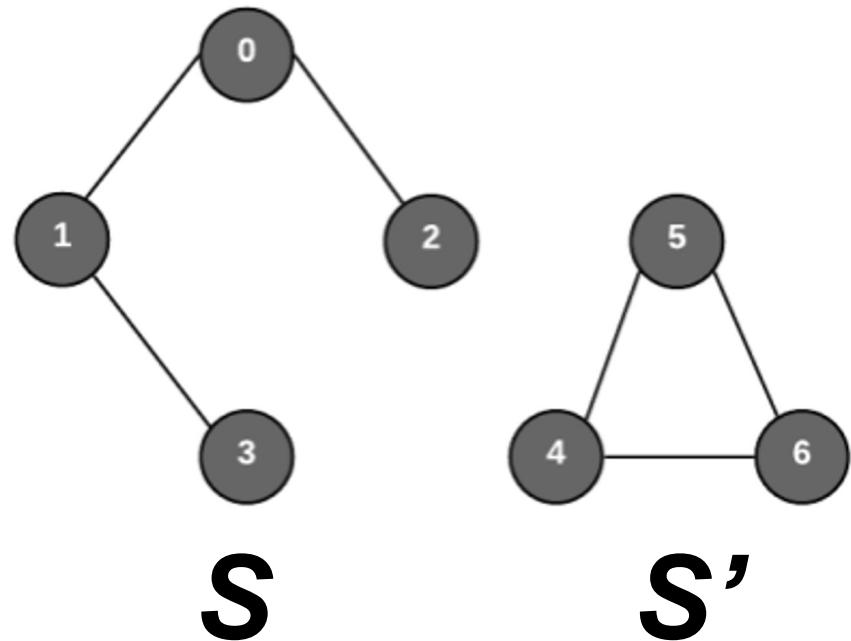
- Element w_{ij} indicates the **weight of the link** from node i to node j
- The link from node B to node C has a weight of 2
- Out-strength (weighted out-degree): $s_i^{out} = \sum_j A_{ij}$
- In-strength (weighted in-degree): $s_i^{in} = \sum_j A_{ji}$
- Total weight: $W = \sum_i s_i^{in} = \sum_i s_i^{out} = \sum_{ij} A_{ij}$

$$\begin{array}{cccccc} & A & B & C & D & E & F \\ \begin{matrix} A \\ B \\ C \\ D \\ E \\ F \end{matrix} & \left(\begin{matrix} 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \\ 2 & 1 & 3 & 1 & 1 & 0 \end{matrix} \right) \end{array}.$$

If a graph is disconnected

Disconnected graphs
have adjacency
matrices with **block
structure**

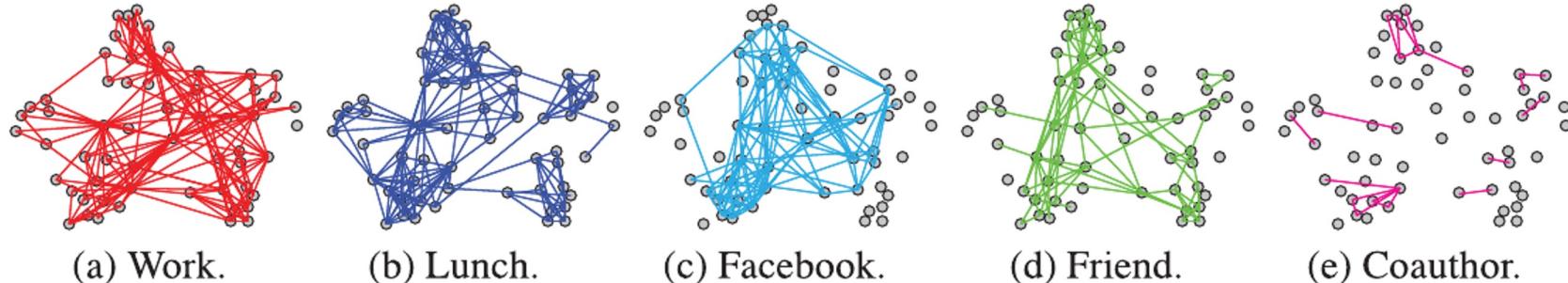
$$A = \begin{bmatrix} S & 0 \\ 0 & S' \end{bmatrix}$$



Beyond simple graphs

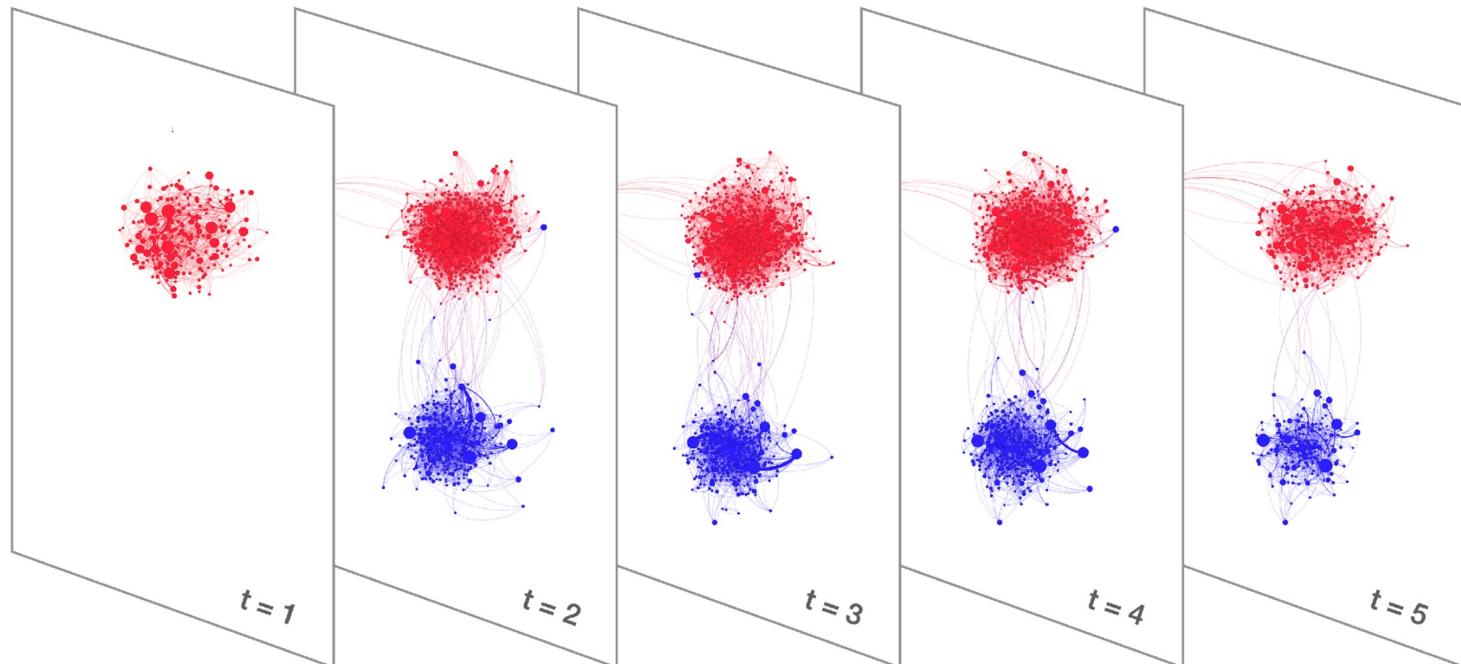
Some networks are multi-layer

- Multi-layer graphs have different edges **over the same nodes**
- Example graph of 61 employees of a university department, with links indicating: coworking, having lunch together, being connected on Facebook, being actual friends, or being co-authors



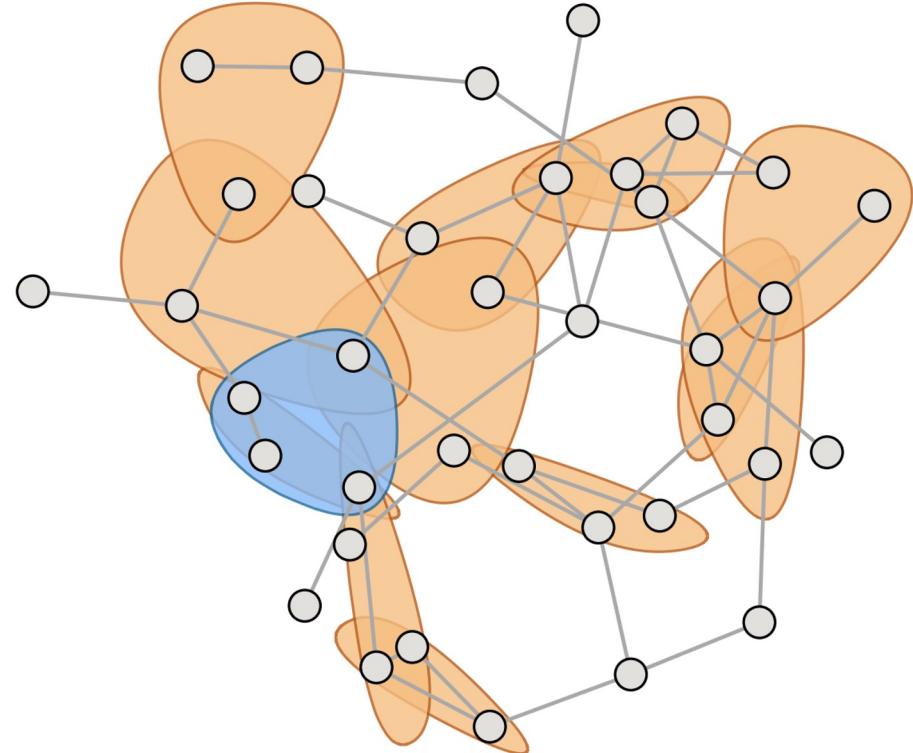
Some networks are time-evolving

Temporal, or
“time-evolving”
networks
At each timestep
there are new
nodes and/or
edges (and/or
deletions)



Some networks are higher-order

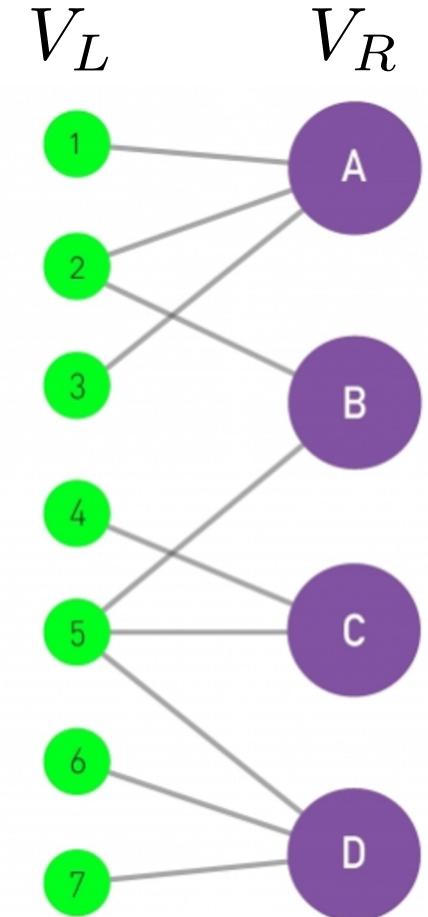
**Higher-order networks, or
“hypergraphs”:**
Hyper-links involve more than
two nodes



Some networks are **bi-partite**

- A **bipartite** graph is a graph
- $G = (V, E)$ such that

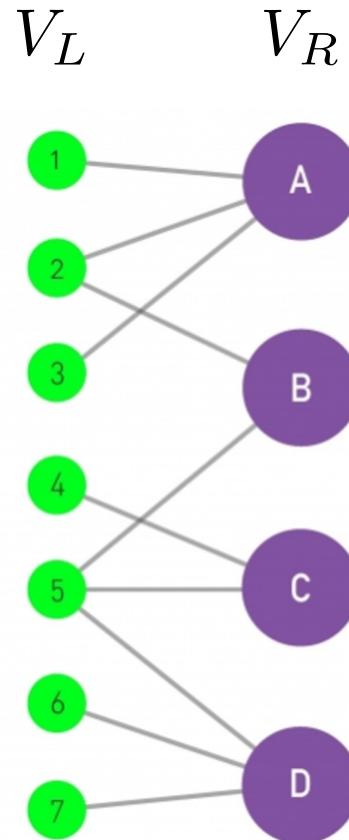
$$V = V_L \cup V_R, V_L \cap V_R = \emptyset, E \subseteq V_L \times V_R$$



Exercise: project a bipartite network

?

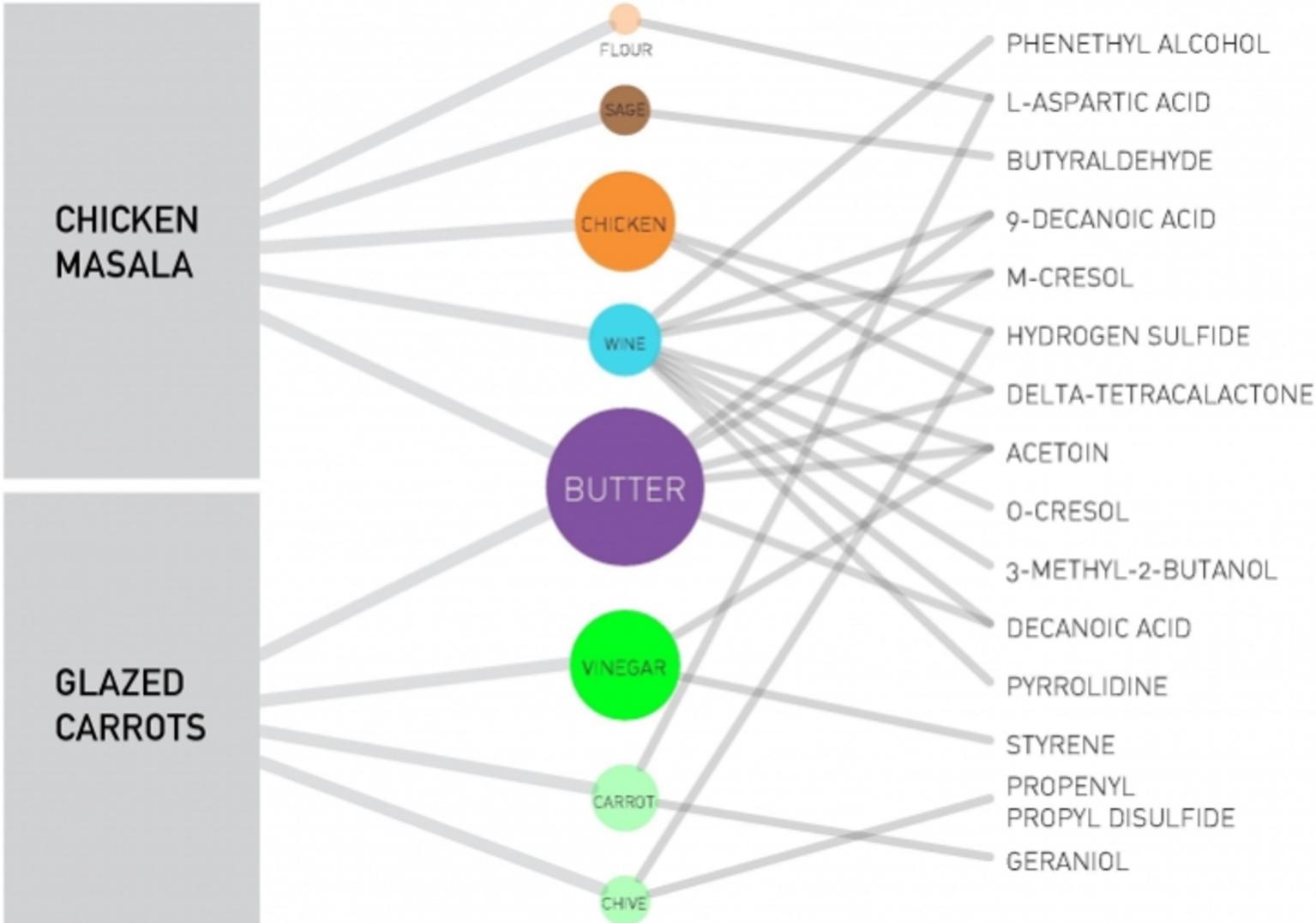
Left projection:
graph where
nodes
are 1, 2, ..., 7
and
nodes are
connected
if they share a
neighbor



?

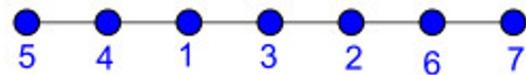
Right
projection:
graph where
nodes
are A, B, ..., D
and
nodes are
connected
if they share a
neighbor

Tripartite network

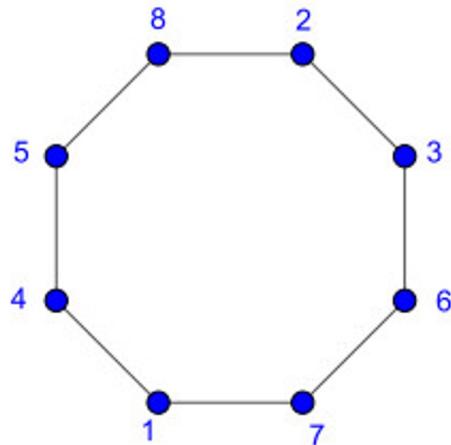


Some graphs have a name

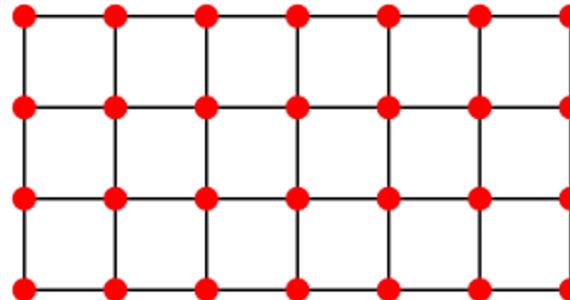
Line



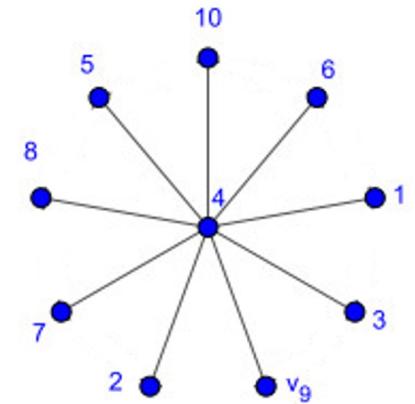
Cycle



Lattice



Star



These are not complex networks!

Clique and Bi-partite clique

- A **clique** is a complete (sub)graph $E = \{ (V \times V) \}$
- An **n-clique** is a complete graph of n nodes
- A **bi-partite clique** is such that

$$V = V_1 \cup V_2, V_1 \cap V_2 = \emptyset, E = (V_1 \times V_2)$$

- A **(n_1, n_2)-clique** is a bipartite clique such that
 $|V_1| = n_1, |V_2| = n_2$

Summary

Things to remember

- . Definitions
 - degree, in-degree, out-degree, strength
 - time-evolving graph, multi-layer graph
 - line graph, cycle graph, star graph, lattice, bi-partite graph, clique
- . Writing the adjacency matrix of a graph, and drawing a graph given its adjacency matrix
- . Plotting the degree distribution of a graph
- . Projecting a bi-partite graph

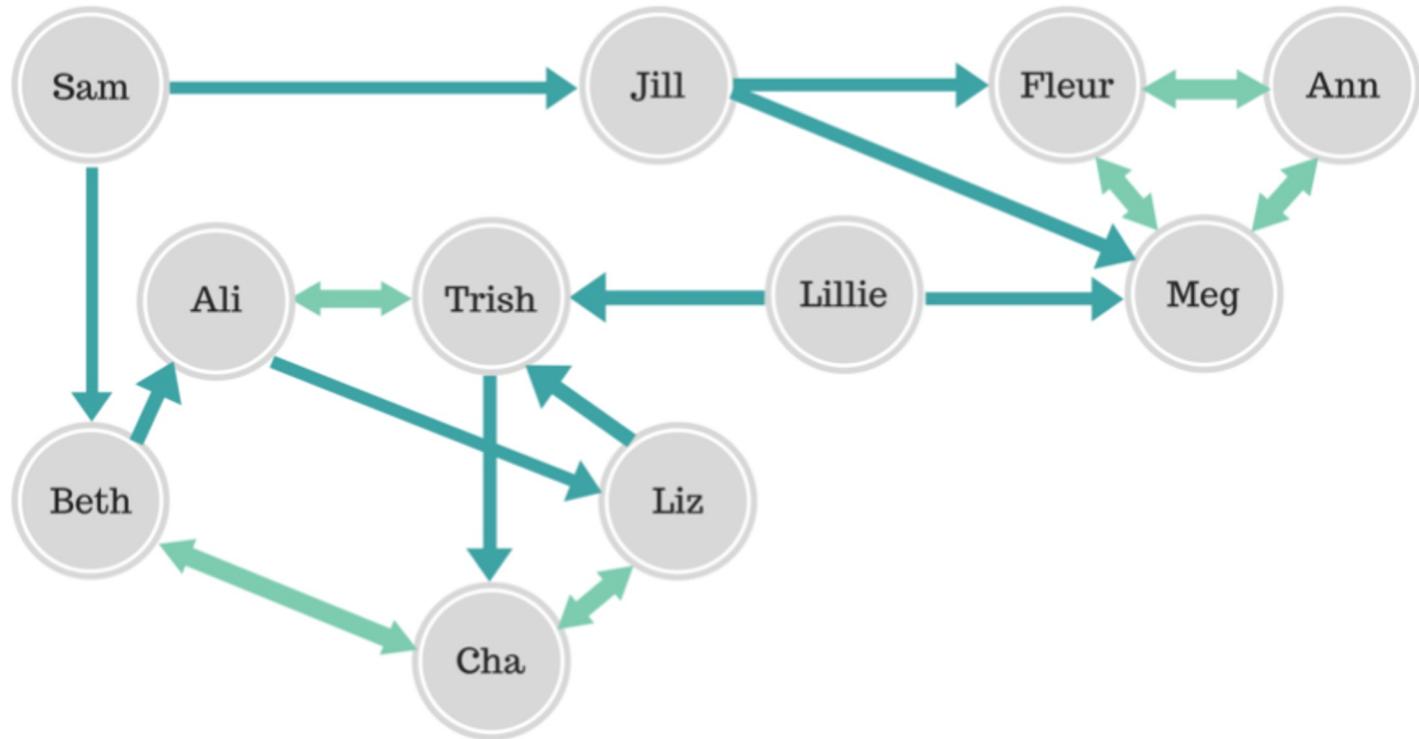
Sources

- A. L. Barabási (2016). Network Science –
[Chapter 02](#)
- URLs cited in the footer of specific slides

Practice on your own

Draw the
indegree,
outdegree,
degree
distribution

Write the
adjacency
matrix



Practice on your own

How do you call the sub-graph induced by nodesets:

- $\{H, A, B\}$
- $\{G, H, D\}$
- $\{B, D, E, G\}$
- $\{A, B, D, E\}$

