Comparative Analysis of Data Structures

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kd-Trees Implementation

- Each node contains:
 - Coordinates
 - Pointer to right subtree
 - Pointer to left subtree
 - discriminator (index of coordinate that splits right and left subtrees)
- kd_tree class contains:
 - pointer to root
 - d (dimension of coordinates in tree)
 - Methods:
 - void insert_node(int* coordinates)
 - void delete_node(int* coordinates)
 - Node* find_node(int* coordinates)



Random Datasets

Random Datasets

- 100 cases for n = 100, 1000 and 10 cases for n = 100K, 1M
- sampled from uniform[0, n*10]
- Goal: Generate histogram of run times for find, insert, delete and compute average runtime for each n. (Varying data)

Random Dataset Permuted

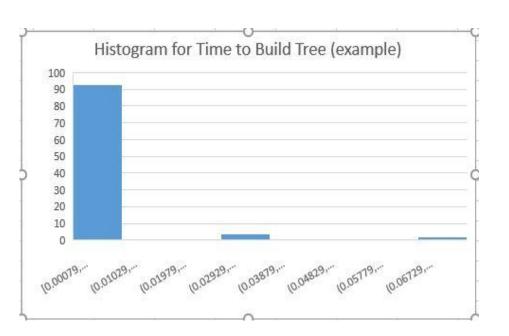
- 100 cases for n = 100, 1000 and 10 cases for n = 100K, 1M
- Permutations of 0th random dataset for each n-case
- Goal: Generate histogram of run times for find, insert, delete and compute average runtime for each n. (Varying shapes)

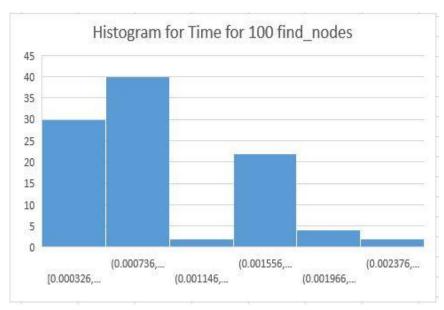


Results

- find_node, insert_node, delete_node work for Handwritten_Test_Cases
- delete_node fails on random datasets
 - difficulty: finding the replacement node for the deleted node is not trivial
- The histograms I generated for insert/find random datasets were not binned correctly (to be completed by final report).
 - difficulty: (finding empirical parameters) either I should zoom in on sections of the histogram and plot that histogram or I should increase the number of insertions significantly









If only I had more time . . .

- Finish implementing delete_node for random datasets
- Finish time-to-run comparisons between Quadtree and kd-tree.
- Implement nearest neighbor query and compare time-to-run with Quadtree's nearest neighbor query.



Quadtree

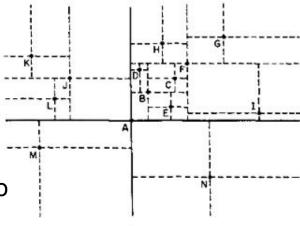
- Tree data structure that contains exactly four children
- used to represented 2D as a data structure
- They decompose into adaptable cells
- Have a tree pyramid but are unbalanced
- Different types:
 - Point quadtree
 - Point-region quadtree
 - edge quadtree
 - Matrix region Quadtree
- Applications
 - Image manipulation
 - Maps
 - Data compression

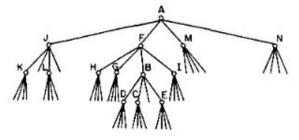


Meshing

Point Quadtree

- Multidimensional generalization of BST
- Non-uniform-sized cells containing one element apiece
- Can reduce storage requirements
- Insertion
 - O(log4N)
- Deletion:
 - Average number or nodes needed to be reinserted is 67%
 - O(Nlogk+1(N))
- Search
 - O(2N^(1-1/k)





Samet, H. (1984). "The quadtree and related hierarchical data structures" (PDF). ACM Computing Surveys. ACM. **16** (2): 187–260. doi:10.1145/356924.356930.

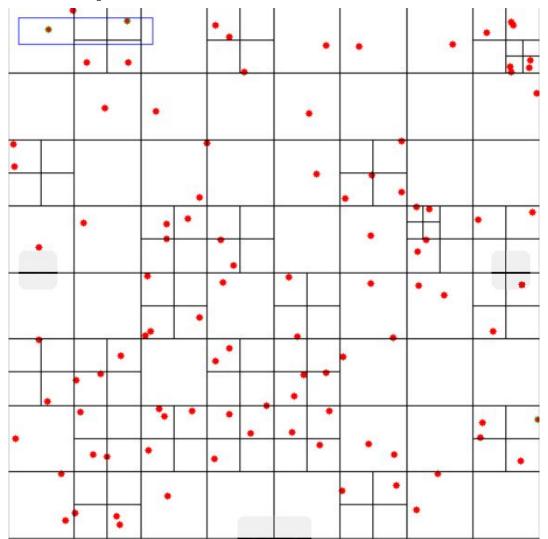


Region Based Quadtree

- Similar to Point Quadtree
- Regions are subvided equally
- Each square contains a point
- Data points are inserted into an MX quadtree by searching for them
- Insertion O(Nlog(N))

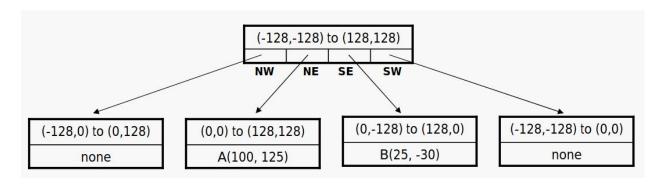


Sample implementation



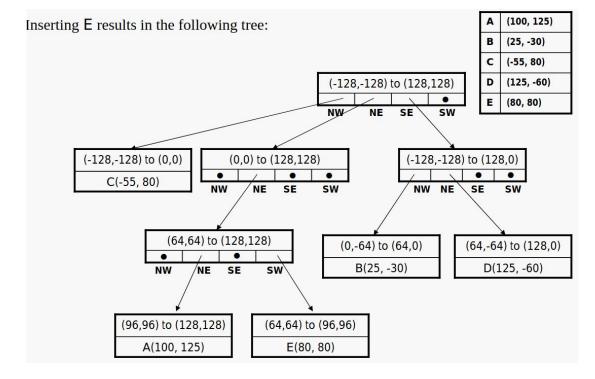


Point Region Quadtree



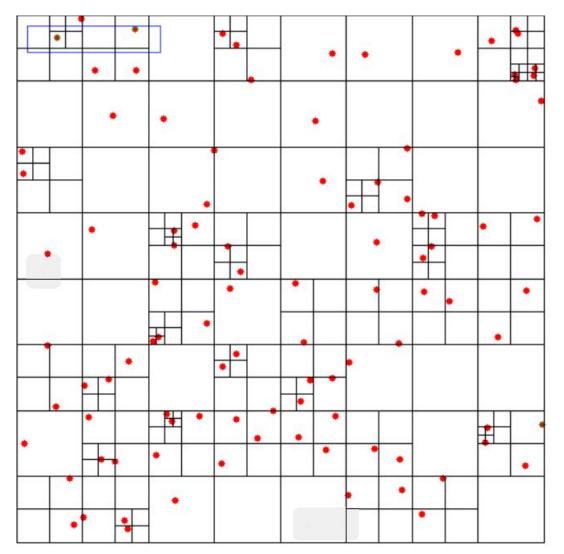
source:

http://courses.cs.vt.edu/~c s3114/Summer15/Notes/T 04_PRQuadTrees.pdf





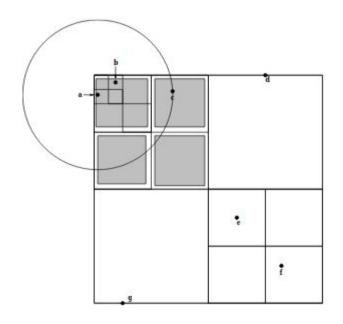
Sample Implementation





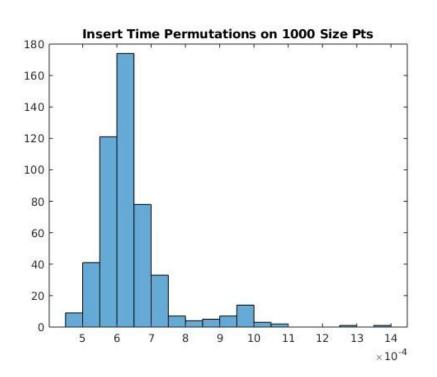
Region Find on Quadtree

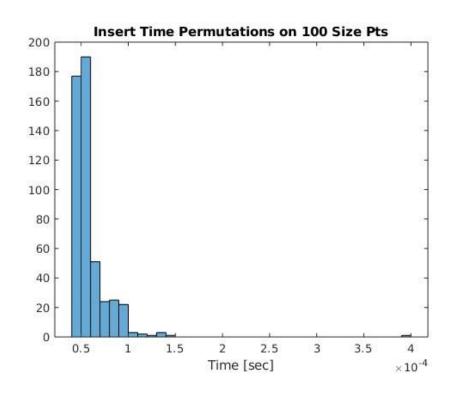
- Same as Point Finding
- Also used for nearest neighbour finder
- Finds closest point in neighbourhood or region based on region





Sample Insertion Time Histogram







KD-TREE vs QUADTREE

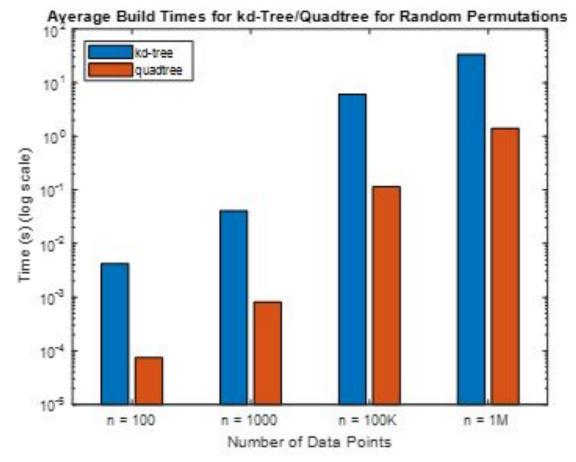
AVG Time insertion Sample 1000

kd-tree : 0.0054

sec

quadtree: 0.0006

sec





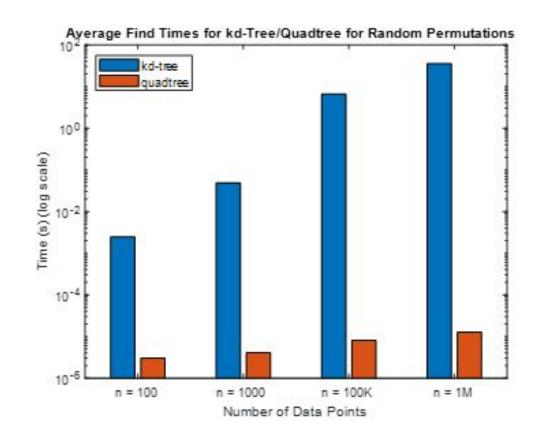
15

KD-TREE vs QUADTREE Avg. Find Time

AVG Finding Time Sample 1000

kd-tree: 0.0013 sec

quadtree: 0.0014 sec





Works cited:

kd-trees: Jon Louis Bentley. 1975. Multidimensional binary search trees used for associative searching. Commun. ACM 18, 9 (September 1975), 509-517. DOI=http://dx.doi.org/10.1145/361002.361007

Quadtree:

Samet, H. (1984). "The quadtree and related hierarchical data structures" (PDF). ACM Computing Surveys. ACM. 16
(2): 187–260. doi:10.1145/356924.356930.



