

# DBMS PROJECT

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## # Topic - Bank Management System

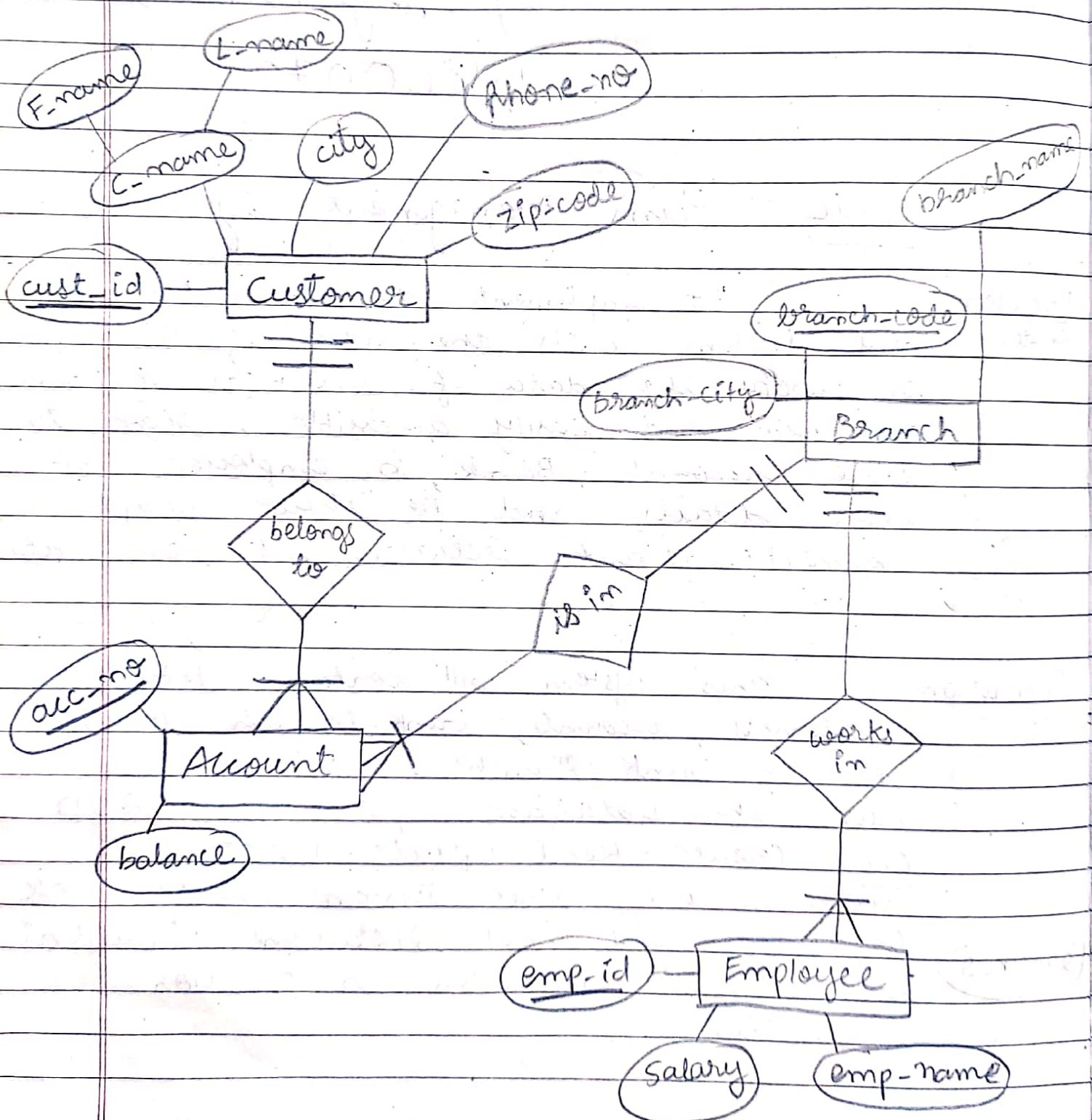
### Problem Statement

To implement the Bank Management system with the help of DBMS, to keep the data of customers of Bank organised and easily accessible. Also to store account, Bank, & employees of bank details and to show proper connection / link between all these data.

### Definition

- This system will contain details of customers, accounts, bank branch & employees of bank branches. Admin can run the Database operations CRUD i.e. Create - Read - Update - Delete, in such a way that system should work properly and this operations shouldn't create any updation or Deletion anomalies.

# \* Entity - Relationship Diagram :-





## \* About ER Diagram

⇒ Entities - There are total 4 entities shown in diagram, all 4 are strong entities as it have its own attribute which acts as primary key.

- ① Customer - stores customer details
- ② Account - stores account details
- ③ Branch - stores branch details
- ④ Employee - stores employee details

⇒ Cardinalities - Total 3 cardinalities or say relationships are present.

- ① belongs-to - Many to one from account entity to customer entity.
- ② is-in - Many to one from account entity to branch entity.
- ③ works-in - Many to one from employee entity to branch entity.

## \* Constraints

- ① An account is owned by one and only one customer, but one customer can have multiple accounts.
- ② A customer can have only one phone number.
- ③ A particular account can be in only one branch only i.e. 2 branches can't have same account's number. So each account number is not null & unique.
- ④ An employee can work in one branch only, he can't have job at multiple branches.



## \* Conversion of ER Model to Relational Model

⇒ As customer is a strong Entity, here cust-id acts as a primary key which is Unique and not Null.

∴ Customer (cust-id, f-name, l-name, city, zip-code, phone-no)

→ Here we have not included c-name as shown in ER diagram because c-name is a Composite Attribute of first name (f-name) and last name (l-name).

⇒ Each account is different, so we have acc-no (Account number) as a primary key for strong Entity Account.

∴ Account (acc-no, balance)

⇒ Strong Entity Branch have branch-code as a primary key.

∴ Branch (branch-code, branch-name, branch-city)

⇒ Each employee is different, he/she have its own emp-id which is not null & Unique i.e. Primary key.

∴ Employee (emp-id, emp-name, salary)

→ A customer can have multiple accounts but an account is opened by only one customer so due to this many to one relationship we will update Account relation by adding cust-id as foreign key.

∴ Account ( acc-no, balance, cust-id )

→ Similarly many to one relationship exists between account & branch. Updation in Account relation by adding branch code as foreign key.

∴ ~~Star~~ Account ( acc-no, balance, cust-id, branch-code )

→ Also adding branch-code as foreign key to employee relation due to works-in many to one relationship bet<sup>n</sup> employee & branch.

∴ employee ( emp-id, salary, emp-name, branch-code )

Listing all the relations / tables

① customer ( cust-id, f-name, l-name, city, zip-code, phone-no )

② account ( acc-no, balance, cust-id, branch-code )

③ branch ( branch-code, branch-name, branch-city )

④ employee ( emp-id, salary, emp-name, branch-code )



## \* Normalization

⇒ List of all non-trivial functional dependencies are :-

(1)  $\text{cust-id} \rightarrow \text{f-name, l-name, city, phone-no, zip-code}$

(2)  $\text{zip-code} \rightarrow \text{city}$

(3)  $\text{branch-code} \rightarrow \text{branch-name, branch-city}$

(4)  $\text{acc-no} \rightarrow \text{balance, cust-id, branch-code}$

(5)  $\text{emp-id} \rightarrow \text{emp-name, salary, branch-code}$

⇒ Normal Forms

(i) First Normal form (1NF)

→ All attributes have atomic values and there are no composite and multivalued attributes. Hence relation is in 1NF.

(ii) Second Normal form (2NF)

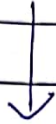
→ As there are no partial dependency in relations & relation is already in 1NF. Hence relation is in 2NF.

(iii) Third Normal form (3NF)

→ There exists 1 transitive dependency  
 $\text{zip-code} \rightarrow \text{city}$  i.e nonprime  $\rightarrow$  nonprime attribute

→ So we need to update customer relation and make a new relation.

⇒ customer ( cust-id, f-name, l-name, city, phone-no, zip-code )



- ① customer ( cust-id, f-name, l-name, phone-no, zip-code )
- ② city-code ( zip-code, city )

⇒ Now the relation is in 3NF.

(iv) Boyce Codd Normal Form (BCNF)

→ For all possible FDs, LHS of FDs is a super key & relation is in 3NF.  
 ∴ Relation is in BCNF.

(v) Fourth Normal Form (4NF)

→ As there is no multi valued dependency and relation is already in BCNF. Hence relation is in 4NF.

(vi) Fifth Normal form (5NF)

→ As there are no join dependencies & join is lossless, also relation is in 4NF. So relation is in 5NF.