

# Proposal

## Blue-Collar Node : A Sharding Approach for Optimizing Storage on Blockchains

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**Abstract**—The scalability of blockchain is a primary and urgent concern. The current popular blockchains have fundamental bottlenecks which limit their ability to have a higher throughput and a lower latency. One of the bottlenecks is the storage requirements of the current blockchains. All the full nodes and miners in a blockchain are required to store the complete blockchains and it grows everytime more transactions are added. The network relies mainly on commodity hardware for block propagation and transaction verification. With the increase in the size of blockchains, storing the entire blockchain on voluntary full nodes becomes impractical. Our project focuses at sharding this blockchain and storing it in several nodes. There is a two-fold advantage of this, transactions can be verified in parallel by different nodes and consensus can be achieved in a faster way increasing the throughput of the network, and the storage requirement of the full node will be decreased substantially. That directly will prevent the blockchain from getting centralized towards supercomputers. However, in this project we will focus only on the storage part of the problem. Executing transaction in parallel will be a part of the future work.

**Index Terms**—blockchain, storage, sharding, blue-collar, decentralized

### I. INTRODUCTION

In 2009, Bitcoin [7] introduced the first blockchain which is an append only database that stores transactions. Blockchain maintains the distributed database in a decentralized network aiming to solve the Byzantine Generals Problem [15]. Byzantine Generals Problem when extrapolated to the banking system i.e to digital money translates to double spending problem, where since there is no material money a malicious node can spend the money twice, and since there is no central leader to check the money can be spent. Bitcoin solves this problem by distributing the ledger to every node in the network and cryptographically securing it making it immutable. Special nodes called miners try to include the transactions into blocks by solving a puzzle. This mechanism is called Proof of Work. Each miner then tries to add this block to the distributed replicated ledger and after the block is added each node on the network updates their ledger.

#### A. Problems

Satoshi Nakamoto randomly decided a block interval of 10 minutes between adding 2 blocks to the chain. The block size

limited to 1MB. Since, the block size is capped and the time interval is constant, the number of transactions that can be dealt with are limited to 3-7 transactions per second. Also, since the ledger is duplicated on each node on the network, as the ledger grows the storage requirement of the node increases.

### II. BACKGROUND AND RELATED WORK

#### A. Background

Blockchain is a decentralized, tamper-proof digital ledger technology that is transforming transaction prospects for many industries. In addition to being the foundation of cryptocurrencies like Ethereum, the blockchain structure, which consists of linked blocks of information, allows for direct transactions across a network of computers without need for a central authority. The Ethereum network consists of three kind of nodes: **full nodes**, **light-weight nodes** and **miners**. Nodes save the complete blockchain and have two main functions:

- 1) to validate transactions
- 2) to propagate transactions and blocks in the network

Although full nodes validate transactions by checking their inputs and outputs, light clients verify transactions by SPV, or by requesting random chunks of a block and using the merkle root to verify each chunk. This verification technique relies on trust on a full node or the node sending data. This can be dangerous, because a malicious node can withhold small amount of data and can get unnoticed by many light clients. The network becomes much more trustworthy if the number of full nodes are a majority since data is always verified. As the blockchain technology is being adopted more, the storage requirement is increasing for each individual node. This increase in infrastructure requirement forces volunteer nodes to move to light-clients. This harms the network in 2 ways :

- 1) The network gets more centralized to people who can afford supercomputers
- 2) The trust on the network weakens as most transactions are not completely validated by full nodes

In this project, we want to introduce a new kind of node that saves only a part of the blockchain thus reducing the storage

requirement and also validates each transaction that is a part of the block. This will introduce the following advantages to the network:

- 1) The trust on the network will be high as the nodes will validate each transaction fully.
- 2) Randomizing block storage in such a way that transactions are fully validated by a small part of the network leading to multiple transactions being validated in parallel.

We will be focusing only on the first advantage as a part of this project, the second would remain a future work.

### B. Related Work

CRUSH [1] and RUSH [6] are used in a number of distributed storage systems to handle the problem of data distribution by replication without causing overhead on the network. The former uses a better approach but since the replication factor is high in that algorithm, the blockchain might ultimately store the same amount of data that it was storing earlier thus causing the blue-collar nodes to store the same amount of data as a full node. We want to change the algorithm to redirect the blocks to different nodes, but we want to leverage the property of equal or balanced distribution form [1].

"Sharding FAQ [3] gives a good glimpse of the problem we are trying to solve with techniques that will come in handy for us to make our solution better. The difference in the approach is that [3] tries to scale the throughput of the blockchain while we focus on the storage optimization of the system.

"A note on data availability and erasure coding" [10] is another blog by the founder of Ethereum [8] [9] which touches our areas of concerns. It talks about light clients and how their lack of validation techniques might be harmful for the network. It exposes a loop-hole in the system which we want to remediate by the creation of "blue-collar" nodes.

Elastico [4] also attempted sharding but they also try to solve the problem as discussed above i.e to have a higher throughput system. [4] also discusses centralized approaches of sharding in distributed databases, some of them are Google's Spanner [12], Scatter [13], and RSCoin [14]. These sharding protocols however do not handle byzantine failures.

## III. PROPOSED SOLUTION

### A. Components

The key component of our system will be what we are calling a "**Blue-collar node**". This node will be a combination of a light and full-node. This node will be required to store the sharded blockchain and the parity bits such that if needed it can re-create the data from a connected full node. **Data-distribution algorithm** : Ceph utilizes CRUSH [1] (Controlled Replication Under Scalable Hashing) which is a data distribution algorithm that distributes object replicas across devices that belong to a heterogeneous storage cluster. In this project, we want to focus in creating a variant of this algorithm, which doesn't store replicas of objects (blocks in a blockchain) but stores only a part of the blockchain and

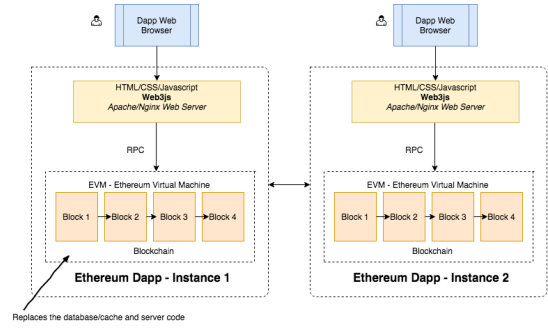


Fig. 1. Ethereum Architecture [5]

redirects the newly created block to only a part of nodes. The algorithm should then create parity bits, that is different for each node and helps each node create the blocks that are stored in other nodes when any of the other nodes fail. To recreate this data we will be utilizing Erasure Code [10] and Reed-Solomon algorithm [11]. **Monitor** : The monitor here will become a part of the ethereum protocol that stores a map or DHT (Distributed Hash Table) of the data in its peer nodes and health of its peer nodes.

### B. Architecture

The current architecture of ethereum blockchain is shown in 1. This architecture shows how in order to create Decentralized Application (dapp) on Ethereum platform one needs to store the entire blockchain that is replicated on both the instances (or all the instances). We intend on developing an algorithm could be activated as a parameter to the ethereum protocol in a similar fashion as light, fast or full node is activated. The new node type that we will be adding is "blue-collar". This node will store part of the data and parity bits calculated using the Reed-Solomon algorithm in order to re-create data on another node when the node associated with the parity goes down.

### C. Challenges

In the methodology described above there are certain challenges that are not yet handled. Some of them are listed below:

- 1) The data distribution has to be randomized in such a way that a malicious intentions can't get hold of data that can be completely lost. In such a case, the contingency would be to recover the data from a full node and distribute it again over the network among peers but would require a lot of data movement.
- 2) The data has to be distributed in such a way that in case some nodes fail, the re-distribution of data does not cause much data movement and doesn't occupy bandwidth that could be used for block propagation.

### D. Implementation Plan

We intend on using a container service to build a decentralized ditributed network of Ethereum. There are multiple ready to use configurations that can be utilized to run a private production ready Ethereum network. These available

container configurations automatically connect to the ethereum mainchain to run a full node by running the default wallet. This can be used as a base system for us to implement our data-distribution algorithm. This algorithm will have three variants suited for the three different kinds of nodes on the network.

#### E. Evaluation Plan

The proposed system would then be evaluated as to how much maximum storage is required by the blue-collar node after we have sharded the blockchain. Metrics that will be used to evaluate the success of the project:

- 1) Ratio of storage utilized by a full node to a blue-collar node
- 2) Ratio of the bandwidth used by a full node to a blue-collar node
- 3) Ratio of the bandwidth used by a light-client to a blue collar node
- 4) Distribution of data among all the peer blue-collar nodes
- 5) Data movement on blue-collar node failure

#### IV. PROPOSED SCHEDULE

Task	Begin date	End date
More literature survey	February 1st, 2018	February 13th, 2018
Setting up the Ethereum system on Docker or Kubernetes	February 14th, 2018	February 28th, 2018
Testing scenarios for Ethereum Blockchain	March 1st, 2018	March 7th, 2018
Creating the data distribution algorithm	February 15th, 2018	March 15th, 2018
Monitor	March 16th, 2018	March 31st, 2018
Testing	April 1st, 2018	End Sem

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