



# **Functional Programming**

Week 2 – Tree Shaped Data and Datatypes

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#### Last Lecture

- algorithm (can be informal) vs. program (concrete prog.-language)
- Haskell script (code, program, ...), e.g., program.hs
   fahrenheitToCelsius f = (f 32) \* 5 / 9
   consists of function definitions that describe input-output behaviour
- function- and parameter-names have to start with lowercase letters
- read-eval-print loop:
   load script, enter expressions and let these be evaluated

```
$ ghci progam.hs
... welcome message ...
Main> fahrenheitToCelsius (3 + 20) - 7
-12.0
Main> ... further expressions ...
...
Main> :q
```

# Structured Data

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# Different Representations of Data

- some (abstract) element can be represented in various ways
- example: numbers roman:
  - decimal:
  - binary:
  - English: tally list:
  - fact: algorithms depend on concrete representation
  - example: addition
  - decimal + binary: process digits of both numbers from right to left

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English: not well-suited

roman: algorithm?

- - tally list: just write the two numbers side-by-side

in Haskell: numbers are built-in, representation not revealed to user

Week 2

7823

- + 9098732

(||| + || = ||||)

(IV + IX = XIII)

(twentvnine + two = thirtvone)

ΧI

11

4/24

1011

eleven

# Different Representations of Data - Continued

- representation must be chosen appropriately
- example: person
  - photographer:



• social analysis: Paul Tom
Sue — Alice — Bol

- advertizing:
- genealogist:
- Bob (bob@foo.com, employee, hobbies: photography, jazz music, . . . )

  Carmen 

   John Suzan 

   — Jack

Alice — @ — Bo

Jack

## Tree Shaped Data

- in functional programming most of the data is tree shaped
- a tree
  - has exactly one root node
  - can have several subtrees: nodes without subtrees are leaves
  - nodes and edges can be labeled
- in computer science, trees are usually displayed upside-down
- examples from previous slide

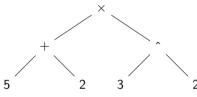


Bob Alice father father mother mother Carmen John Suzan Jack Week 2

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#### Expressions = Trees

- mathematical expressions can be represented as trees
- example
  - expression in textual form:  $(5+2) \times 3^2$
  - expression as tree



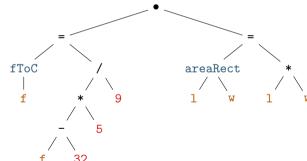
- remarks
  - the process of converting text into tree form is called parsing
  - ullet operator precedences ( $\hat{}$  binds stronger than  $\times$  b.s.t. +) and parentheses are only required for parsing
    - parsing  $(5+2) \times (3^2)$  results in tree above
    - $5+2\times3^2$  and  $((5+2)\times3)^2$  represent other trees
  - algorithm of calculator
    - convert textual input into tree
    - evaluate the tree bottom-up, i.e., start at leaves and end at root

## **Programs** = Trees

- programs can be represented as trees, too: abstract syntax tree
- example
  - program in textual form

```
-- some comment
fToC f = (f - 32) * 5 / 9
areaRect 1 w = 1 * w
```

abstract syntax tree (draft)



comments and parentheses are no longer present in syntax tree

## **Tree Shaped Data**

- many programs deal with tree shaped data
- examples
  - calculator evaluates expression tree
  - compiler translates abstract syntax tree into machine code
  - search engine translates query into HTML (tree shaped)
  - contact application manages tree shaped personal data
  - file systems are organised as trees
- trees as mental model or representation of data is often suitable
- good news: processing tree shaped data is well-supported in functional programming
- next lecture: define functions on trees
- this lecture: restriction of trees via types

# Types

# **Types**

- functions are often annotated by their domain and codomain, e.g.,
  - $(!): \mathbb{N} \to \mathbb{N}$
  - $(/): \mathbb{R} \times (\mathbb{R} \setminus \{0\}) \to \mathbb{R}$
  - $log_2: \mathbb{R}_{\geq 0} \to \mathbb{R}$
- domain and codomain provide useful information
  - domain: what are allowed inputs to a function
  - codomain: what are potential outputs of the function
- aim: specify domains and codomains of (Haskell-)functions
- notions
  - elements or values
    - maths: 5, 8,  $\pi$ ,  $-\frac{3}{4}$ , ...
    - Haskell: 5, 8, 3.141592653589793, -0.75, ..., "hello", 'c', ...
  - sets of elements to specify domain or codomain, in Haskell: types
    - maths:  $\mathbb{N}$ ,  $\mathbb{Z}$ ,  $\mathbb{Q}$ ,  $\mathbb{R}$ ,  $\mathbb{Q} \setminus \{0\}$ , ...
    - Haskell: Integer, Double, String, Char, ...

## **Typing Judgements**

- in maths, we write statements like  $7 \in \mathbb{Z}$ ,  $7 \in \mathbb{R}$ ,  $0.75 \notin \mathbb{Z}$
- in Haskell we can also express that a value or expression has a certain type via typing judgements
  - format: expression :: type
    - examples
      - 7 :: Integer or 7 :: Double
      - 'c' :: Char
- that an expression indeed has the specified type is checked by the Haskell compiler
  - if an expression has not the given type, a type error is displayed
    - examples
      - 7 :: String or 0.75 :: Integer or 'c' :: String
      - (7 :: Integer) :: Double
    - remarks
      - unlike in maths where  $\mathbb{N}\subseteq\mathbb{Z}\subseteq\mathbb{Q}$ , in Haskell the types Integer and Double are not subtypes of each other
      - although some expressions can have both types (e.g., 7 + 5),
         in general numbers of different types have to be converted explicitly
        - once a typing judgement is applied, the type of that expressions is fixed

# Typing of Haskell Expressions

- not only values but also functions have a type, e.g.,
  - (/) :: Double -> Double -> Double • (+) :: Integer -> Integer -> Integer
  - (+) :: Double -> Double -> Double
  - head :: String -> Char

- remarks
- a function can have multiple types, e.g., (+)
- limited expressivity, e.g. (/) :: Double -> Double \ {0} -> Double not allowed
- type of arguments matches input-types of function

type checking enforces that in all function applications,

- example: consider expressions expr1 / expr2 • recall: (/) :: Double -> Double -> Double
- it will be checked that both expr1 and expr2 have type Double
- type of the overall expression expr1 / expr2 will then be Double
- examples
  - $\bullet$  5 + 3 / 2 • 5 + '3' or 5.2 + 0.8 :: Integer

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# **Static Typing**

- Haskell performs static typing
- static typing: types will be checked before evaluation (by contrast, dynamic typing checks types during evaluation)
- when loading Haskell script
  - check types of all function definitions someFun x ... z = expr:
     check that lhs someFun x ... z has same type as rhs expr
  - consequence: expressions cannot change their type during evaluation
- when entering expression in REPL: type check expression before evaluation
- benefits
  - no type checking required during evaluation
  - no type errors during evaluation

# Built-In Types – A First Overview

- numbers
  - Integer arbitrary-precision integers
  - Int fixed-precision integers with at least 29 bits (-100, 0, 999)
  - Float single-precision floating-point numbers (-12.34, 5.78e36)
  - Double double-precision floating-point numbers
- characters and text
  - Char a single character ('a', 'Z', ' ')
  - String text of arbitrary length ("", "a", "The answer is 42.")
  - some characters have to be escaped via the backslash-symbol \:
    - '\t' and '\n' tabulator and new-line
    - '\"' and '\'' double- and single quote
    - '\\' the backslash character
    - example: in the program

text = "Please say \"hello\"\nwhenever you enter the room"
the string text corresponds to the following two lines:

Please say "hello"

whenever you enter the room

• Bool - yes/no-decisions or truth-values (True, False)

# Datatypes

#### **Current State**

- each value and function in Haskell has a type
- types are used to define input and output of function
- example: fahrenheitToCelsius :: Double -> Double
- built-in types for numbers, strings, and truth values
- missing: how to define types that describe tree shaped data?
- solution: definition of (algebraic) datatypes

## **Datatype Definitions**

- recall: a tree consists of a (labelled) root and 0 or more subtrees
- a datatype definition defines a set of trees by specifying all possible labelled roots together with a list of allowed subtrees
- Haskell scripts can contain many datatype definitions of the form data TName =

```
CName1 type1_1 ... type1_N1
| ...
| CNameM typeM_1 ... typeM_NM
deriving Show
```

where

- data is a Haskell keyword to define a new datatype
- TName is the name of the new type; type-names always start with capital letters
- CName1,...,CNameM are the labels of the permitted roots;
   these are called constructors and have to start with capital letters
  - typeI\_J can be any Haskell type, including TName itself
  - I is used as separator between different constructors
  - deriving Show is required for displaying values of type TName

# **Example Datatype Definition** – Date

# deriving Show

- here, there is only one constructor: DMY
  for day and month the precision of Int is sufficient
- the values of the type Date are exactly trees of the form



- in Haskell, these trees are built via the constructor DMY; DMY is a function of type
   Int -> Int -> Integer -> Date that is not evaluated
- example value of type Date: DMY 11 10 2021

# **Example Datatype Definition** – Person

```
data Person = -- name of type
Person -- constructor name can be same as type name
String -- first name
String -- last name
Bool -- married
Date -- birthday
```

#### deriving Show

- reuse of previously defined types is permitted, in particular Date
- this leads to trees with more than one level of subtrees
- example program that defines a person (and an auxiliary date)

```
today = DMY 11 10 2021
myself = Person "Rene" "Thiemann" True today
-- is the same as
myself = Person "Rene" "Thiemann" True (DMY 11 10 2021)
```

```
data Brand = Audi | BMW | Fiat | Opel deriving Show
data Vehicle =
    Car Brand Double -- horsepower
    | Bicycle
    | Truck Int -- number of wheels
    deriving Show
```

- Brand just defines 4 car brands; all "trees" of type Brand consist of a single node; such datatypes are called enumerations
- there are three kinds of Vehicles, each having a different list of types
- example expressions of type Vehicle:

```
Car Fiat (60 + 1)
Car Audi 149.5
Bicycle
```

Truck (-7) -- types don't enforce all sanity checks

# **Example Datatype Definition** – Expr

```
data Expr =
    Number Integer
```

- Plus Expr Expr
- Negate Expr

#### deriving Show

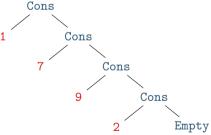
- type Expr models arithmetic expressions with addition and negation
- Expr ia a recursive datatype: Expr is defined via Expr itself
- recursive datatypes contain values (trees) of arbitrary large height
  - expression (-(5+2)) + 3 in Haskell (as value of type Expr):
    - Plus (Negate (Plus (Number 5) (Number 2))) (Number 3)
  - expression as tree Plus Negate Number Plus Number Number

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# **Example Datatype Definition – Lists**

```
    lists are just a special kind of trees, e.g., lists of integers
    data List =
        Empty
        | Cons Integer List
        deriving Show
```

- example representation of list [1, 7, 9, 2]
  - in Haskell: Cons 1 (Cons 7 (Cons 9 (Cons 2 Empty)))
  - as tree:



## **Summary**

- mental model: data = tree shaped data
- type = set of values; restricts shape of trees
- built-in types for numbers and strings
- user-definable datatypes, e.g., for expressions, lists, persons
   data TName =

```
CName1 type1_1 ... type1_N1
| ...
| CNameM typeM_1 ... typeM_NM
deriving Show
```

• missing: function definitions on trees