



Functional Programming

Week 9 – Calendar Application, Scope, Modules

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Last Lecture – Library Functions

`take, drop :: Int -> [a] -> [a]`

`splitAt :: Int -> [a] -> ([a], [a])`

`takeWhile, dropWhile :: (a -> Bool) -> [a] -> [a]`

`span :: (a -> Bool) -> [a] -> ([a], [a])`

`zipWith :: (a -> b -> c) -> [a] -> [b] -> [c]`

`zip :: [a] -> [b] -> [(a, b)]`

`unzip :: [(a, b)] -> ([a], [b])`

`words, lines :: String -> [String]`

`unwords, unlines :: [String] -> String`

`concatMap :: (a -> [b]) -> [a] -> [b]`

`($) :: (a -> b) -> a -> b`

Last Lecture – List Comprehension

- list comprehension
 - shape: `[(x,y,z) | x <- [1..n], let y = x ^ 2, y > 100, Just z <- f y]`
 - consists of guards, generators, local declarations
 - translated via `concatMap`
- examples

```
prime n = n >= 2 && null [ x | x <- [2 .. n - 1], n `mod` x == 0]
```

```
ptriples n = [ (x,y,z) |  
  x <- [1..n], y <- [x..n], z <- [y..n], x^2 + y^2 == z^2]
```

Last Lecture – Printing a Calendar

- given a month and a year, print the corresponding calendar

- example: December 2021

Mo Tu We Th Fr Sa Su

1 2 3 4 5

6 7 8 9 10 11 12

...

- we concentrate on printing, assuming machinery for construction

```
type Month = Int
```

```
type Year = Int
```

```
type Dayname = Int -- Mo = 0, Tu = 1, ..., So = 6
```

```
-- monthInfo returns name of 1st day in m. and number of days in m.
```

```
monthInfo :: Month -> Year -> (Dayname, Int)
```

Design and Functionality for Representing Character-Pictures

```
type Height = Int
type Width = Int
type Picture = (Height, Width, [[Char]])
```

```
above :: Picture -> Picture -> Picture
stack :: [Picture] -> Picture
```

```
beside :: Picture -> Picture -> Picture
spread :: [Picture] -> Picture
```

```
tile :: [[Picture]] -> Picture
tile = stack . map spread
```

Finalizing the Calendar

Creating Pictures

- single 'pixels'

```
pixel :: Char -> Picture  
pixel c = (1, 1, [[c]])
```

- rows

```
row :: String -> Picture  
row r = (1, length r, [r])
```

- blank

```
blank :: Height -> Width -> Picture  
blank h w = (h, w, blanks)  
  where  
    blanks = replicate h (replicate w ' ')
```

Constructing a Month

- as indicated, assume function

```
monthInfo :: Month -> Year -> (Dayname, Int)
```

where daynames are 0 (Monday), 1 (Tuesday), ...

```
daysOfMonth :: Month -> Year -> [Picture]
```

```
daysOfMonth m y =
```

```
  map (row . rjustify 3 . pic) [1 - d .. numSlots - d]
```

```
  where
```

```
    (d, t) = monthInfo m y
```

```
    numSlots = 6 * 7 -- max 6 weeks * 7 days per week
```

```
    pic n = if 1 <= n && n <= t then show n else ""
```

```
rjustify :: Int -> String -> String
```

```
rjustify n xs
```

```
  | 1 <= n = replicate (n - 1) ' ' ++ xs
```

```
  | otherwise = error ("text (" ++ xs ++ ") too long")
```

```
  where l = length xs
```


Tiling the Days

- `daysOfMonth` delivers list of 42 single pictures (of size 1×3)
- missing: layout + header for final picture (of size 7×21)

```
month :: Month -> Year -> Picture
```

```
month m y = above weekdays . tile . groupsOfSize 7 $ daysOfMonth m y
  where weekdays = row " Mo Tu We Th Fr Sa Su"
```

```
-- groupsOfSize splits list into sublists of given length
```

```
groupsOfSize :: Int -> [a] -> [[a]]
```

```
groupsOfSize n [] = []
```

```
groupsOfSize n xs = ys : groupsOfSize n zs
```

```
  where (ys, zs) = splitAt n xs
```

Printing a Month

- transform a `Picture` into a `String`
`showPic :: Picture -> String`
`showPic (_, _, css) = unlines css`
- show result of `month m y` as `String`
`showMonth :: Month -> Year -> String`
`showMonth m y = showPic $ month m y`
- display final string via `putStr :: String -> IO ()` to properly print newlines and drop double quotes

```
> showMonth 12 2021
```

```
" Mo Tu We Th Fr Sa Su\n      1  2  3  4  5\n  6  7  8  9 10 11 12\n 13 14 15 16 17 18 19\n 20 21 22 23 24 25 26\n 27 28 29 30 31"
```

```
> putStr $ showMonth 12 2021
```

```
Mo Tu We Th Fr Sa Su
      1  2  3  4  5
  6  7  8  9 10 11 12
 13 14 15 16 17 18 19
 20 21 22 23 24 25 26
 27 28 29 30 31
```

Scope

Scope

- consider program (1 compile error)

```
radius = 15
```

```
area radius = pi^2 * radius
```

```
squares x = [ x^2 | x <- [0 .. x]]
```

```
length [] = 0
```

```
length (_,xs) = 1 + length xs
```

```
data Rat = Rat Integer Integer
```

```
createRat n d = normalize $ Rat n d where normalize ... = ...
```

- **scope**
 - resolve ambiguities
 - defines which names of variables, functions, types, ... are visible at a given program position
 - controlling scope to structure larger programs (imports / exports)

Scope of Names

```
radius = 15
```

```
area radius = pi^2 * radius
```

- in the following we assume that `name_i` in the real code is always just `name` and the `_i` is used for addressing the different occurrences of `name`

- renamed Haskell program

```
radius_1 = 15
```

```
area_1 radius_2 = pi_1^2 * radius_3
```

- scope of names in right-hand sides of equations
 - is `radius_3` referring to `radius_2` or `radius_1`?
 - what is `pi_1` referring to?
- rule of thumb for searching `name`: search **inside-out**
 - think of abstract syntax tree of expression
 - whenever you pass a **let**, **where**, **case**, or function definition where `name` is **bound**, then refer to that **local** name
 - if nothing is found, then search **global** function `name`, also in Prelude
- `radius_3` refers to `radius_2`, `pi_1` to `Prelude.pi`

Local Names in Case-Expressions

- general case: `case expr of { pat1 -> expr1; ...; patN -> exprN }`
 - each `patI` binds the variables that occur in `patI`
 - these variables can be used in `exprI`
 - the newly bound variables of `patI` bind stronger than any previously bound variables
- example Haskell expression

```
case xs_1 of                -- renamed Haskell expression
  [] -> xs_2
  (x_1 : xs_3) -> case xs_4 ++ ys_1 of
    [] -> ys_2
    (x_2 : xs_5) -> x_3 : xs_6 ++ ys_3
```

- `x_3` refers to `x_2` (since `x_2` is further inside than `x_1`)
- `xs_6` refers to `xs_5` (since `xs_5` is further inside than `xs_3`)
- `xs_4` refers to `xs_3`
- `xs_1`, `xs_2`, `ys_1`, `ys_2`, and `ys_3` are not bound in this expression (the proper references need to be determined further outside)

Local Names in Let-Expressions

let {

pat1 = expr1; ...; patN = exprN;

f1 pats1 = fexpr1; ...; fM patsM = fexprM

} **in** expr

- all variables in pat1 ... patN and all names f1 ... fM are bound
 - these can be used in expr, in each exprI and in each fexprJ
 - variables of patsJ bind strongest, but only in fexprJ
- **let** (x_1, y_1) = (y_2 + 1, 5) -- renamed Haskell expression
f_1 x_2 = x_3 + g_1 y_3 id_1
g_2 y_4 f_2 = f_3 \$ g_3 x_4 f_4
in (f_5, g_4, x_5, y_5)
 - y_2, y_3 and y_5 refer to y_1
 - x_3 refers to x_2 since x_2 binds stronger than x_1
 - x_4 and x_5 refer to x_1
 - f_3 and f_4 refer to f_2 since f_2 binds stronger than f_1
 - g_1, g_3 and g_4 refer to g_2
 - f_5 refers to f_1
 - id_1 is not bound in this expression

Global Function Definitions

- general case:

`fname pats = expr`

- all variables in `pats` are bound locally and can be used in `expr`
- `fname` is **not** locally bound, but added to global lookup table
- all variables/names in `expr` without local reference will be looked up in global lookup table
- lookup in global table does not permit ambiguities

- `radius_1 = 15` -- renamed Haskell program

`area_2 radius_2 = pi_1^2 * radius_3`

`length_1 [] = 0`

`length_2 (_,xs_1) = 1 + length_3 xs_2`

- `radius_1`, `area_2` and `length_1/2` are stored in global lookup table
- global lookup table has ambiguity: `length_1/2` vs. `Prelude.length`
- `pi_1` is not locally bound and therefore refers to `Prelude.pi`
- `radius_3` refers to local `radius_2` and not to global `radius_1`
- `xs_2` refers to `xs_1`
- `length_3` is not locally bound and because of mentioned ambiguity, this leads to a compile error

Global vs. Local Definitions

```
length :: [a] -> Int

-- choose definition 1,
length = foldr (\ _-> (1 +)) 0

-- definition 2,
length =
  let { length [] = 0; length (x : xs) = 1 + length xs }
  in length

-- or definition 3
length [] = 0
length (_ : xs) = 1 + length xs
```

- definitions 1 and 2 compile since there is no `length` in the rhs that needs a global lookup
- in contrast, definition 3 does not compile
- still definitions 1 and 2 result in ambiguities in global lookup table
→ study Haskell's module system

Modules

Modules

- so far
 - Haskell program is a **single** file, consisting of several definitions
 - all global definitions are visible to user

```
-- functions on rational numbers
```

```
data Rat = Rat Integer Integer      -- internal definition of datatype
```

```
normalize (Rat n d) = ...           -- internal function
```

```
createRat n d = normalize $ Rat n d -- function for external usage
```

```
...
```

```
-- application: approximate pi to a certain precision
```

```
piApprox :: Integer -> Rat
```

```
piApprox p = ...
```

- motivation for modules
 - structure programs into smaller **reusable** parts without copying
 - distinguish between **internal** and **external** definitions
 - clear interface for users of modules
 - maintain invariants
 - improve maintainability

Modules in Haskell

```
-- first line of file ModuleName.hs
module ModuleName(exportList) where
-- standard Haskell type and function definitions
```

- each `ModuleName` has to start with uppercase letter
- each module is usually stored in separate file `ModuleName.hs`
- if Haskell file contains no `module` declaration, ghci inserts module name `Main`
- `exportList` is comma-separated list of function-names and type-names, these functions and types will be accessible for users of the module
- if (`exportList`) is omitted, then everything is exported
- for types there are different export possibilities
 - `module Name(Type)` exports `Type`, but no constructors of `Type`
 - `module Name(Type(..))` exports `Type` and its constructors

Example: Rational Numbers

```
module Rat(Rat, createRat, numerator, denominator) where
data Rat = Rat Integer Integer
normalize = ...
createRat n d = normalize $ Rat n d
numerator (Rat n d) = n
...
instance Num Rat where ...
instance Show Rat where ...
```

- external users know that a type `Rat` exists
- they only see functions `createRat`, `numerator` and `denominator`
- they don't have access to constructor `Rat` and therefore cannot form expressions like `Rat 2 4` which break invariant of cancelled fractions
- they can perform calculations with rational numbers since they have access to `(+)` of class `Num`, etc., in particular for the instance `Rat`
- for the same reason, they can display rational numbers via `show`

Example: Application

```
module PiApprox(piApprox, Rat) where
-- Prelude is implicitly imported

-- import everything that is exported by module Rat
import Rat

-- or only import certain parts
import Rat(Rat, createRat)

-- import declarations must be before other definitions
piApprox :: Integer -> Rat
piApprox n = let initApprox = createRat 314 100 in ...
```

- there can be multiple `import` declarations
- what is imported is not automatically exported
 - when importing `PiApprox`, type `Rat` is visible, but `createRat` is not
 - if application requires both `Rat` and `PiApprox`, import both modules:

```
import PiApprox
import Rat
```

Resolving Ambiguities

```
-- Foo.hs
```

```
module Foo where pi = 3.1415
```

```
-- Problem.hs
```

```
module Problem where
```

```
import Foo
```

```
pi = 3.1415
```

```
area r = pi * r^2
```

- problem: what is `pi` in definition of `area`? (global name)
- lookup map is ambiguous: `pi` defined in `Prelude`, `Foo`, and `Problem`
- ambiguity persists, even if definition is identical
- solution via **qualifier**: disambiguate by using `ModuleName.name` instead of `name`
 - write `area r = Problem.pi * r^2` in `Problem.hs`
(or `area r = Prelude.pi * r^2`)

Qualified Imports

```
module Foo where pi = 3.1415
module SomeLongModuleName where fun x = x + x

module ExampleQualifiedImports where

-- all imports of Foo have to use qualifier
import qualified Foo

-- result: no ambiguity on unqualified "pi"

import qualified SomeLongModuleName as S
-- "as"-syntax changes name of qualifier

area r = pi * r^2
myfun x = S.fun (x * x)
```


Summary

Summary

- calendar application
- scoping rules determine visibility of function names and variable names
- larger programs can be structured in **modules**
 - explicit **export-lists** to distinguish internal and external parts
 - advantage: changes of internal parts of module **M** are possible without having to change code that imports **M**, as long as exported functions of **M** have same names and types
 - if no module name is given: **Main** is used as module name
 - further information on modules

<https://www.haskell.org/onlinereport/modules.html>