CHC-COMP 2018

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1st Competition on Solving Constrained Horn Clauses



CHC-COMP: CHC Solving Competition

First edition on July 13, 2018 at HVCS@FLOC

Constrained Horn Clauses (CHC) is a fragment of First Order Logic (FOL) that is sufficiently expressive to describe many verification, inference, and synthesis problems including inductive invariant inference, model checking of safety properties, inference of procedure summaries, regression verification, and sequential equivalence. The CHC competition (CHC-COMP) will compare state-of-the-art tools for CHC solving with respect to performance and effectiveness on a set of publicly available benchmarks. The winners among participating solvers are recognized by measuring the number of correctly solved benchmarks as well as the runtime.

Web: https://chc-comp.github.io/

Gitter: https://gitter.im/chc-comp/Lobby

GitHub: https://github.com/chc-comp

Format: https://chc-comp.github.io/2018/format.html





Constrained Horn Clauses (CHC)

A Constrained Horn Clause (CHC) is a FOL formula of the form

$$\forall V \cdot (\varphi \wedge p_1[X_1] \wedge \cdots \wedge p_n[X_n] \to h[X]$$

where

- \mathcal{T} is a background theory (e.g., Linear Arithmetic, Arrays, Bit-Vectors, or combinations of the above)
- $ullet \varphi$ is a constraint in the background theory ${\mathcal T}$
- $p_1, ..., p_n, h$ are n-ary predicates
- p_i[X] is an application of a predicate to first-order terms



CHC Satisfiability

A \mathcal{T} -model of a set of a CHCs Π is an extension of the model M of \mathcal{T} with a first-order interpretation of each predicate p_i that makes all clauses in Π true in M

A set of clauses is **satisfiable** if and only if it has a model

This is the usual FOL satisfiability

A \mathcal{T} -solution of a set of CHCs Π is a substitution σ from predicates p_i to \mathcal{T} -formulas such that $\Pi \sigma$ is \mathcal{T} -valid

In the context of program verification

- a program satisfies a property iff corresponding CHCs are satisfiable
- solutions are inductive invariants
- refutation proofs are counterexample traces



Format 1/2

```
benchmark ::=
       logic
       fun_decl+
       (assert chc_assert)*
       (assert chc query)
       (check-sat)
logic ::= (set-logic HORN)
fun decl ::= (declare-fun symbol ( sort* ) Bool)
chc assert ::= ;; a well-formed first-order sentence of the form
       | (forall (var decl+) (=> chc tail chc head))
       chc head
```



Format 2/2

```
var_decl ::= (symbol sort)
chc head ::=
        (u predicate var*), where all variables are
DTSTTNCT
chc tail ::=
       (u predicate var*)
       | i formula
       (and (u predicate var*)+ i formula*)
chc query ::= ;; a well-formed first-order sentence of the form
       | (forall ( var decl+ ) (=> chc tail false)
u predicate ::= uninterperted predicate (i.e., Boolean function)
i formula ::= an SMT-LIB formula over variables, and
              interpreted functions and predicates
```



Example

```
(set-logic HORN)
(declare-fun Inv (Int) Bool)
;; fact
(assert (forall ((x Int)) (=> (= x 0) (Inv x))))
;; chc
(assert (forall ((x Int) (y Int))
       (=> (and (Inv x) (<= x 10) (= y (+ x 1))) (Inv y))))
;; query
(assert (forall ((x Int))
              (=> (and (Inv x) (> x 15)) false)))
(check-sat)
```



Benchmark Selection

chc-comp.github.io as collection for CHC benchmarks

- Collected from participants
- Generated with SeaHorn from SV-COMP Device Driver problems

Converted to competition format using format script (Adriene)

https://github.com/chc-comp/scripts

Out of successfully formatted benchmarks selected

- 339 CHC-LINEAR(LIA)
- 132 CHC-LINEAR(LRA)
- Random selection, favoring hard-for-spacer benchmarks

Benchmarks released after competition:

https://github.com/chc-comp/chc-comp18-benchmarks



Participants

Eldarica Philipp Rümmer

Hoice Adriene Champion

PECOS John Gallagher

TransfHORNer Fabio Fioravanti

Aultimate TreeAutomizer Alexander Nutz

Aultimate Unihorn Alexander Nutz

Hors Concours

spacer Arie Gurfinkel

rebus unnamed solver not entered

in the competition



Hoice [Champion et al., 2018]

ICE-based Refinement Type Discovery for Higher-Order Functional Programs

Adrien Champion¹, Tomoya Chiba¹, Naoki Kobayashi¹, Ryosuke Sato²

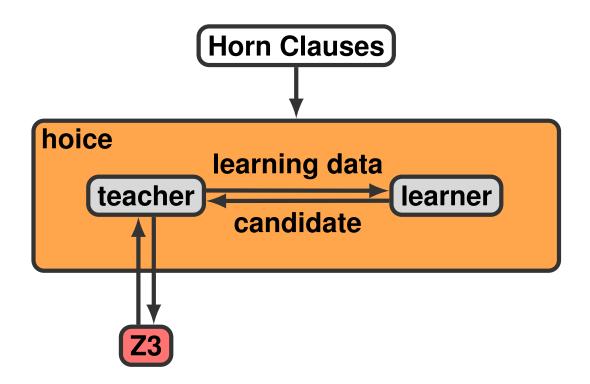
¹ The University of Tokyo
² Kyushu University

https://github.com/hopv/hoice

- machine-learning-based Horn clause solver: generalized ICE framework [Garg et al., 2014]
- context: higher-order program verification
- supports Int, Real, and Array
- support for datatypes coming soon

Hoice [Champion et al., 2018]

- learner produces candidates for the predicates
- teacher checks each clause is respected
- • ⇒ each check is a quantifier-free (non-Horn) formula
- using Z3 [de Moura and Bjørner, 2008] (separate process)



TransfHORNer

```
E. De Angelis <sup>(1)</sup>, F. Fioravanti <sup>(1)</sup>, A. Pettorossi <sup>(2)</sup>, M. Proietti <sup>(3)</sup>
```

(1) DEC, University "G. d'Annunzio" of Chieti-Pescara, Italy(2) DICII, University of Rome Tor Vergata, Rome, Italy(3) CNR-IASI, Rome, Italy

HCVS 2018

Oxford - July 13, 2018

TransfHORNer distinctive features

Apply satisfiability-preserving CHC transformations (Unfold/Define/Fold)

Discover invariants and try to increase the effectiveness of backend CHC solvers

Related work

- Verification of C programs with integers and arrays (PEPM13, VMCAI14, SCP14,...)
- VeriMAP (TACAS14) https://fmlab.unich.it/VeriMAP/
- Semantics-based VC generation for C programs (PPDP15, SCP17)
- Relational program verification (SAS16, TPLP18)
- Verification of BPMN business processes (LOPSTR16, RULEML+RR17)
- Satisfiability checking of CHCs with inductively defined data structures (lists, trees, ...) without induction



PECOS: Partial evaluation and Constraint Specialisation

Developed by John P. Gallagher, Bishoksan Kafle and José F. Morales (Roskilde, IMDEA, Melbourne). Based partly on previous experiments with RAHFT (Kafle, Gallagher & Morales, CAV'2016).

Oriented towards pure LRA Horn clauses (challenging to modify for mixed Boolean/numeric constraints). Iteratively performs the following transformations.

- Constraint specialisation (PEPM'2015, Sci. Comput. Program. 2017) for invariant discovery. Backwards/forwards propagation, convex polyhedral global analysis with query-answer (magic-set) transformation.
- Partial evaluation of the computation trees for false. Achieves control-flow refinement and predicate specialisation. Employs a property-based finite abstract domain. Described in PEPM'1993, ICLP'2018, WST'2018.

Prototype software on github.com/jpgallagher/pecos. Implemented in Ciao Prolog with interfaces to PPL and Yices.

Ultimate TreeAutomizer

 $Tree\ Automizer = Trace\ Abstraction\ +\ Tree\ Automata$

Trace Abstraction (Automizer) approach:

Program is correct iff each error trace (word) is infeasible.

Tree Automizer approach:

Set of CHCs is sat iff the constraints in each derivation of false (tree) are unsat.

Under the hood:

- Ultimate Automata Library
- SMTInterpol

Contributors:

Daniel Dietsch, Alexander Nutz, Mostafa M. Mohamed, Daniel Tischner, Jochen Hoenicke, Matthias Heizmann, Andreas Podelski



Ultimate Unihorn Automizer

Input: Set of constrained Horn clauses Φ Approach:

- Construct (possibly recursive) program P_{Φ} such that: P_{Φ} is safe iff Φ is sat
- Apply off-the-shelf program verifier

Under the hood:

- Program verifier: Ultimate Automizer
- Predicate providers: Newton-style interpolation, SMTInterpol
- ► SMT Solvers: CVC4, MathSAT5, SMTInterpol, Z3
- Ultimate Automata Library

Contributors:

Daniel Dietsch, Matthias Heizmann, Jochen Hoenicke, Alexander Nutz, Andreas Podelski



The ELDARICA Horn Solver

Hossein Hojjat¹ Philipp Rümmer ²

¹Rochester Institute of Technology ²Uppsala University

HCVS 2018: 5th Workshop on Horn Clauses for Verification and Synthesis 13 July 2018

ELDARICA Overview

- Horn solver developed since 2011
- Open-source, implemented in Scala, running in JVM
- Input formats:

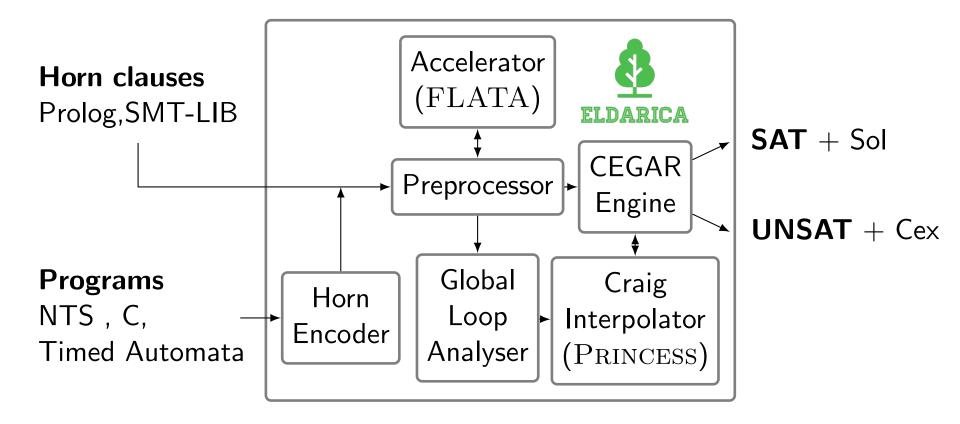
SMT-LIB, Prolog, C, timed automata

• Theories:

LIA, NIA, arrays, algebraic data-types, bit-vectors

- Scala/Java API
- Support for linear + non-linear clauses
- https://github.com/uuverifiers/eldarica

ELDARICA Architecture



Competition Setup

StarExec cluster environment

Dedicated Queue of 20 nodes

2 jobs per node

64GB per job (more than promised)

900s timeout enforced by runsolver

About 12 hours for one complete run of all tools and categories

Detailed results (and benchmarks) will be publicly available on StarExec



Results: LRA

solver	cnt	ok	sat	uns	fld	to	mo	time	real	space	uniq
rebus	132	96	82	14	36	36	0	42838	42842	59	15
spacer	132	82	73	9	50	50	0	54070	54032	93	3
tree-aut	132	25	24	1	107	107	0	101241	92580	7488	0
eldarica	132	20	18	2	112	77	0	72526	35825	2120	0
TransfHORNer	132	1 5	2	1 3	117	105	1	103490	103619	1294	2
uni-aut	132	9	9	0	123	90	0	82877	70345	7536	0
hoice	132	1	0	1	131	131	0	118793	118459	31	0
pecos	132	0	0	0	132	57	1	54743	54811	893	0

cnt – number of benchmarks fld – failed

ok – solved

sat – solved sat

uns – solved unsat

to – timeout

Space – sum of mem

uniq – unique solved

mo -- memory out

time – sum of time

unsound



real – sum of wall

Results: LIA

solver	cnt	ok	sat	uns	fld	to	mo	time	real	space	uniq
spacer	339	225	153	72	114	112	1	122323	122337	509	56
rebus	339	185	139	46	154	154	0	147539	147553	182	12
hoice	339	107	75	32	232	231	0	219810	217971	93	2
eldarica	339	105	97	8	234	234	0	216152	133402	5856	5
TransfHORNer	339	99	63	36	240	160	8	196339	196471	1642	0
pecos	339	61	61	0	278	152	1	171955	173970	3819	3
tree-aut	339	53	34	19	286	247	0	241323	220198	19234	0
uni-aut	339	51	43	8	288	211	0	199278	140978	19378	0



Big Thanks to





Discussion

CHC-COMP 2019

• Dates, format, co-location, tracks

Ranking

• Should we decide on 1st three places?

Non-Linear CHC as a category?

Benchmark storage

Github is limited to 1-2GB per repo

Benchmark selection / availability

