## Models of Computation

## **Tutorial Exercises 3**

- 1. Use the Pumping Lemma to prove that the following languages over the alphabet  $\Sigma = \{0, 1\}$  are not regular:
  - (i)  $L_1 = \{0^n 1^m 0^n : m, n > 0\}.$
  - (ii)  $L_2 = \Sigma^* \setminus \{ 0^n 1^n : n \ge 0 \}.$
  - (iii)  $L_3 = \{ 0^m 1^n : m \neq n \}.$
- 2. Consider  $L = \{ w \in \{0, 1\} : w \text{ is not a palindrome } \}$ . A palindrome is a string that reads the same forward or backward. Show that L is not regular in three different ways:
  - (i) using the fact that the set of palindroms is not regular;
  - (ii) using the Myhill-Nerode Theorem;
  - (iii) using the Pumping Lemma.
- 3. Let  $L_1 = \{01^i 01^j 01^j \mid i, j > 0\}$  and  $L_2 = \{w \in \{0, 1\}^* \mid w \text{ contains } 00\}.$ 
  - (i) Show that  $L = L_1 \cup L_2$  satisfies the Pumping Lemma with pumping length 3.
  - (ii) Prove that  $L = L_1 \cup L_2$  is not regular using the Myhill-Nerode theorem.
  - (iii) Does (ii) contradict (i)?
- 4. The alphabet is  $\Sigma = \{0, 1\}$ . Give context-free grammars that generate the following languages.
  - (i) The empty set.
  - (ii) The language of palindromes.
  - (iii) The set of strings containing at least three 1's.
  - (iv) The set of odd-length strings whose middle symbol is 0.
- 5. Find context-free grammars for the following languages (where  $m, n \ge 0$ ):
  - (a)  $L = \{ a^n b^m : n \le m+3 \}$
  - (b)  $L = \{ a^n b^m : n \neq 2m \}$
- 6. (a) Describe what is meant by *well-bracketing* in the case where there are two kinds of parentheses, namely, ( ) and [ ]. For example ( [ ] ) and ( [ ] ( ( ) ) ) [ ] are well-bracketed, but ( [ ) ] is not.
  - Give a CFG over  $\Sigma = \{ [, ], (, ) \}$  that generates well-bracketed parentheses.
  - (b) Say that ( and ] are positive and ) and [ are negative. A string over  $\Sigma$  is said to be *alternating* if any two adjacent symbols have opposite polarities.
    - Give a CFG that generates the set L of alternating and well-bracketed strings over  $\Sigma$ . For example  $[\ (\ )\ ]\ [\ ]$  and  $(\ [\ (\ [\ ]\ )\ ]\ )\ (\ )$  are in L but  $[\ ]\ (\ )$  and  $[\ [\ ]\ ]$  are not.

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Adapted from materials by Hanno Nickau and Luke Ong