

MODELS OF COMPUTATION

Tutorial Exercises 2

- Give regular expressions for each of the following subsets of $\{a, b\}^*$.
 - $\{x : x \text{ contains an even number of } a\text{'s}\}$
 - $\{x : x \text{ contains an odd number of } b\text{'s}\}$
 - $\{x : x \text{ contains an even number of } a\text{'s or an odd number of } b\text{'s}\}$
- Give NFAs accepting the sets of strings denoted by the following regular expressions:
 - $(000^* + 111^*)^*$
 - $(01 + 10)(01 + 10)(01 + 10)$
 - $(0 + 1(01^*0)^*1)^*$

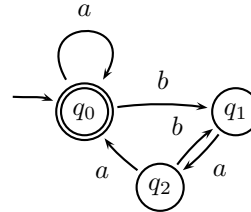
Try to simplify as much as possible.

- Use the procedure given in the lectures to convert the following NFA

$$N = (\{q_0, q_1, q_2\}, \{a, b\}, \delta, \{q_0\}, \{q_0\})$$

to a regular expression where δ is given by:

δ	a	b
q_0	$\{q_0\}$	$\{q_1\}$
q_1	$\{q_2\}$	\emptyset
q_2	$\{q_0\}$	$\{q_1\}$



- Prove or disprove each of the following:
 - $(E + F)^* \equiv E^* + F^*$
 - $(EF + E)^*E \equiv E(FE + E)^*$
 - $E(FE + E)^*F \equiv EE^*F(EE^*F)^*$
- Prove that if L is a regular language, so is $L^R = \{w \mid \text{the reversal of } w \text{ is in } L\}$.
- Show that the regular languages are closed under the following operations:
 - $\min(L) = \{w \mid w \text{ is in } L \text{ but no proper prefix of } w \text{ is in } L\}$.
 - $\max(L) = \{w \mid w \text{ is in } L \text{ and for no } x \text{ other than } \epsilon \text{ is } wx \text{ in } L\}$
 - $\text{init}(L) = \{w \mid \text{for some } x, wx \text{ is in } L\}$.

[Hint: Start with a DFA for L and perform an appropriate construction.]

Tom Melham, Trinity 2009
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