

Bitmagic

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```
theory IPv4Addr
imports Main
  NumberWang
  ~/src/HOL/Word/Word
  ~/src/HOL/Library/Code-Target-Nat
begin

value (2::nat) < 2^32
```

1 Modelling IPv4 Addresses

An IPv4 address is basically a 32 bit unsigned integer

```
type-synonym ipv4addr = 32 word
```

```
value 42 :: ipv4addr
```

```
value (42 :: ipv4addr) ≤ 45
```

Conversion between natural numbers and IPv4 addresses

```
definition nat-of-ipv4addr :: ipv4addr ⇒ nat where
  nat-of-ipv4addr a = unat a
```

```
definition ipv4addr-of-nat :: nat ⇒ ipv4addr where
  ipv4addr-of-nat n = of-nat n
```

```
lemma ((nat-of-ipv4addr (42::ipv4addr))::nat) = 42 by eval
```

```
lemma ((ipv4addr-of-nat (42::nat))::ipv4addr) = 42 by eval
```

The maximum IPv4 address

definition *max-ipv4-addr* :: *ipv4addr* **where**
max-ipv4-addr \equiv *ipv4addr-of-nat* ($(2^{32}) - 1$)

lemma *max-ipv4-addr-number*: *max-ipv4-addr* = 4294967295
by *eval*

lemma *max-ipv4-addr* = 0b11111111111111111111111111111111
by(*fact max-ipv4-addr-number*)

lemma *max-ipv4-addr-max-word*: *max-ipv4-addr* = *max-word*
by(*simp add: max-ipv4-addr-number max-word-def*)

lemma *max-ipv4-addr-max*[*simp*]: $\forall a. a \leq \text{max-ipv4-addr}$
by(*simp add: max-ipv4-addr-max-word*)

lemma *range-0-max-UNIV*[*simp*]: $\{0 \dots \text{max-ipv4-addr}\} = \text{UNIV}$
by(*simp add: max-ipv4-addr-max-word*) *fastforce*

identity functions

lemma *nat-of-ipv4addr-ipv4addr-of-nat*: $\llbracket n \leq \text{nat-of-ipv4addr max-ipv4-addr} \rrbracket$
 $\implies \text{nat-of-ipv4addr (ipv4addr-of-nat } n) = n$

by (*metis ipv4addr-of-nat-def le-unat-uoI nat-of-ipv4addr-def*)

lemma *nat-of-ipv4addr-ipv4addr-of-nat-mod*: *nat-of-ipv4addr (ipv4addr-of-nat* *n*)
 $= n \bmod 2^{32}$

by(*simp add: ipv4addr-of-nat-def nat-of-ipv4addr-def unat-of-nat*)

lemma *ipv4addr-of-nat-nat-of-ipv4addr*: *ipv4addr-of-nat (nat-of-ipv4addr* *addr*)
 $= \text{addr}$

by(*simp add: ipv4addr-of-nat-def nat-of-ipv4addr-def*)

Equality of IPv4 addresses

lemma $\llbracket n \leq \text{nat-of-ipv4addr max-ipv4-addr} \rrbracket \implies \text{nat-of-ipv4addr (ipv4addr-of-nat } n) = n$

apply(*simp add: nat-of-ipv4addr-def ipv4addr-of-nat-def*)

apply(*induction* *n*)

apply(*simp-all*)

by(*unat-arith*)

lemma *ipv4addr-of-nat-eq*: $x = y \implies \text{ipv4addr-of-nat } x = \text{ipv4addr-of-nat } y$
by(*simp add: ipv4addr-of-nat-def*)

1.1 Representing IPv4 Addresses

fun *ipv4addr-of-dotteddecimal* :: *nat* \times *nat* \times *nat* \times *nat* \Rightarrow *ipv4addr* **where**
ipv4addr-of-dotteddecimal (*a,b,c,d*) = *ipv4addr-of-nat* (*d* + 256 * *c* + 65536 * *b* + 16777216 * *a*)

fun *dotteddecimal-of-ipv4addr* :: *ipv4addr* \Rightarrow *nat* \times *nat* \times *nat* \times *nat* **where**
dotteddecimal-of-ipv4addr *a* = (*nat-of-ipv4addr ((a >> 24) AND 0xFF)*,
nat-of-ipv4addr ((a >> 16) AND 0xFF),
nat-of-ipv4addr ((a >> 8) AND 0xFF),
nat-of-ipv4addr (a AND 0xFF))

```

declare ipv4addr-of-dotteddecimal.simps[simp del]
declare dotteddecimal-of-ipv4addr.simps[simp del]

lemma ipv4addr-of-dotteddecimal (192, 168, 0, 1) = 3232235521 by eval
lemma dotteddecimal-of-ipv4addr 3232235521 = (192, 168, 0, 1) by eval

a different notation for ipv4addr-of-dotteddecimal

lemma ipv4addr-of-dotteddecimal-bit:
  ipv4addr-of-dotteddecimal (a,b,c,d) = (ipv4addr-of-nat a << 24) + (ipv4addr-of-nat
b << 16) + (ipv4addr-of-nat c << 8) + ipv4addr-of-nat d
proof -
  have a: (ipv4addr-of-nat a) << 24 = ipv4addr-of-nat (a * 16777216)
  by(simp add: ipv4addr-of-nat-def shiftl-t2n, metis Abs-fnat-hom-mult comm-semiring-1-class.normalizing-s
of-nat-numeral)
  have b: (ipv4addr-of-nat b) << 16 = ipv4addr-of-nat (b * 65536)
  by(simp add: ipv4addr-of-nat-def shiftl-t2n, metis Abs-fnat-hom-mult comm-semiring-1-class.normalizing-s
of-nat-numeral)
  have c: (ipv4addr-of-nat c) << 8 = ipv4addr-of-nat (c * 256)
  by(simp add: ipv4addr-of-nat-def shiftl-t2n, metis Abs-fnat-hom-mult comm-semiring-1-class.normalizing-s
of-nat-numeral)
  have ipv4addr-of-nat-suc:  $\bigwedge x. \text{ipv4addr-of-nat } (\text{Suc } x) = \text{word-succ } (\text{ipv4addr-of-nat } (x))$ 
  by(simp add: ipv4addr-of-nat-def, metis Abs-fnat-hom-Suc of-nat-Suc)
  { fix x y
  have ipv4addr-of-nat x + ipv4addr-of-nat y = ipv4addr-of-nat (x+y)
  apply(induction x arbitrary: y)
  apply(simp add: ipv4addr-of-nat-def)
  apply(simp add: ipv4addr-of-nat-suc)
  by (metis (hide-lams, no-types) comm-semiring-1-class.normalizing-semiring-rules(22)
comm-semiring-1-class.normalizing-semiring-rules(24) word-succ-p1)
  } from this a b c
  show ?thesis
  apply(simp add: ipv4addr-of-dotteddecimal.simps)
  apply(thin-tac ?x)+
  apply(rule ipv4addr-of-nat-eq)
  by presburger
qed

lemma size-ipv4addr: size (x::ipv4addr) = 32 by(simp add:word-size)
lemma ipv4addr-of-nat-shiftr-slice: ipv4addr-of-nat a >> x = slice x (ipv4addr-of-nat
a)
  by(simp add: ipv4addr-of-nat-def shiftr-slice)
value (4294967296::ipv4addr) = 2^32

lemma nat-of-ipv4addr-slice-ipv4addr-of-nat:
  nat-of-ipv4addr (slice x (ipv4addr-of-nat a)) = (nat-of-ipv4addr (ipv4addr-of-nat
a)) div 2^x

```

```

proof –
  have mod4294967296: int a mod 4294967296 = int (a mod 4294967296)
    using zmod-int by auto
  have int-pullin: int (a mod 4294967296) div 2 ^ x = int (a mod 4294967296
div 2 ^ x)
    using zpower-int zdiv-int by (metis of-nat-numeral )
  show ?thesis
    apply(simp add: shiftr-slice[symmetric])
    apply(simp add: ipv4addr-of-nat-def word-of-nat)
    apply(simp add: nat-of-ipv4addr-def unat-def)
    apply(simp add: shiftr-div-2n)
    apply(simp add: uint-word-of-int)
    apply(simp add: mod4294967296 int-pullin)
  done
qed

lemma ipv4addr-and-255: (x::ipv4addr) AND 255 = x AND mask 8
  apply(subst pow2-mask[of 8, simplified, symmetric])
  by simp

lemma ipv4addr-of-nat-AND-mask8: (ipv4addr-of-nat a) AND mask 8 = (ipv4addr-of-nat
(a mod 256))
  apply(simp add: ipv4addr-of-nat-def and-mask-mod-2p)
  apply(simp add: word-of-nat)
  apply(simp add: uint-word-of-int)
  apply(subst mod-mod-cancel)
  apply simp
  apply(simp add: zmod-int)
  done

lemma dotteddecimal-of-ipv4addr-ipv4addr-of-dotteddecimal:
  [ a < 256; b < 256; c < 256; d < 256 ] ==> dotteddecimal-of-ipv4addr (ipv4addr-of-dotteddecimal
(a,b,c,d)) = (a,b,c,d)
proof –
  assume a < 256 and b < 256 and c < 256 and d < 256
  note assms = ⟨a < 256⟩ ⟨b < 256⟩ ⟨c < 256⟩ ⟨d < 256⟩
  hence a: nat-of-ipv4addr ((ipv4addr-of-nat (d + 256 * c + 65536 * b +
16777216 * a) >> 24) AND mask 8) = a
    apply(simp add: ipv4addr-of-nat-def word-of-nat)
    apply(simp add: nat-of-ipv4addr-def unat-def)
    apply(simp add: and-mask-mod-2p)
    apply(simp add: shiftr-div-2n)
    apply(simp add: uint-word-of-int)
    apply(subst mod-pos-pos-trivial)
    apply simp-all
    apply(subst mod-pos-pos-trivial)
    apply simp-all
    apply(subst mod-pos-pos-trivial)
    apply simp-all
  done
from assms have b: nat-of-ipv4addr ((ipv4addr-of-nat (d + 256 * c + 65536

```

```

* b + 16777216 * a) >> 16) AND mask 8) = b
  apply(simp add: ipv4addr-of-nat-def word-of-nat)
  apply(simp add: nat-of-ipv4addr-def unat-def)
  apply(simp add: and-mask-mod-2p)
  apply(simp add: shiftr-div-2n)
  apply(simp add: uint-word-of-int)
  apply(subst mod-pos-pos-trivial)
  apply simp-all
  apply(subst mod-pos-pos-trivial[where b=4294967296])
  apply simp-all
  apply(simp add: NumberWang.div65536)
done
from assms have c: nat-of-ipv4addr ((ipv4addr-of-nat (d + 256 * c + 65536
* b + 16777216 * a) >> 8) AND mask 8) = c
  apply(simp add: ipv4addr-of-nat-def word-of-nat)
  apply(simp add: nat-of-ipv4addr-def unat-def)
  apply(simp add: and-mask-mod-2p)
  apply(simp add: shiftr-div-2n)
  apply(simp add: uint-word-of-int)
  apply(subst mod-pos-pos-trivial)
  apply simp-all
  apply(subst mod-pos-pos-trivial[where b=4294967296])
  apply simp-all
  apply(simp add: NumberWang.div256)
done
from (d < 256) have d: nat-of-ipv4addr (ipv4addr-of-nat (d + 256 * c +
65536 * b + 16777216 * a) AND mask 8) = d
  apply(simp add: ipv4addr-of-nat-AND-mask8)
  apply(simp add: ipv4addr-of-nat-def word-of-nat)
  apply(simp add: nat-of-ipv4addr-def)
  apply(subgoal-tac (d + 256 * c + 65536 * b + 16777216 * a) mod 256 = d)
  prefer 2
  apply(simp add: NumberWang.mod256)
  apply(simp)
  apply(simp add: unat-def)
  apply(simp add: uint-word-of-int)
  apply(simp add: mod-pos-pos-trivial)
done
from a b c d show ?thesis
apply(simp add: ipv4addr-of-dotteddecimal.simps dotteddecimal-of-ipv4addr.simps)
  apply(simp add: ipv4addr-and-255)
done
qed

```

lemma *ipv4addr-of-dotteddecimal-eqE*: $\llbracket \text{ipv4addr-of-dotteddecimal } (a,b,c,d) = \text{ipv4addr-of-dotteddecimal } (e,f,g,h); a < 256; b < 256; c < 256; d < 256; e < 256; f < 256; g < 256; h < 256 \rrbracket \implies$
 $a = e \wedge b = f \wedge c = g \wedge d = h$

by (*metis Pair-inject dotteddecimal-of-ipv4addr-ipv4addr-of-dotteddecimal*)

previous and next ip addresses, without wrap around

definition *ip-next* :: *ipv4addr* \Rightarrow *ipv4addr* **where**
ip-next *a* \equiv if *a* = *max-ipv4-addr* then *max-ipv4-addr* else *a* + 1
definition *ip-prev* :: *ipv4addr* \Rightarrow *ipv4addr* **where**
ip-prev *a* \equiv if *a* = 0 then 0 else *a* - 1

lemma *ip-next* 2 = 3 **by** *eval*

lemma *ip-prev* 2 = 1 **by** *eval*

lemma *ip-prev* 0 = 0 **by** *eval*

1.2 IP ranges

lemma *UNIV-ipv4addrset*: (*UNIV* :: *ipv4addr* set) = {0 .. *max-ipv4-addr*}
by(*auto*)

lemma (42::*ipv4addr*) \in *UNIV* **by** *eval*

definition *ipv4range-set-from-netmask*::*ipv4addr* \Rightarrow *ipv4addr* \Rightarrow *ipv4addr* set **where**

ipv4range-set-from-netmask *addr* *netmask* \equiv let *network-prefix* = (*addr* AND *netmask*) in {*network-prefix* .. *network-prefix* OR (*NOT* *netmask*)}

lemma *ipv4range-set-from-netmask* (*ipv4addr-of-dotteddecimal* (192,168,0,42))
(*ipv4addr-of-dotteddecimal* (255,255,0,0)) =
{*ipv4addr-of-dotteddecimal* (192,168,0,0) .. *ipv4addr-of-dotteddecimal*
(192,168,255,255)}

by(*simp* add: *ipv4range-set-from-netmask-def* *ipv4addr-of-dotteddecimal.simps*
ipv4addr-of-nat-def)

lemma *ipv4range-set-from-netmask* (*ipv4addr-of-dotteddecimal* (192,168,0,42))
(*ipv4addr-of-dotteddecimal* (0,0,0,0)) = *UNIV*

by(*simp* add: *UNIV-ipv4addrset* *ipv4addr-of-dotteddecimal.simps* *ipv4addr-of-nat-def*
ipv4range-set-from-netmask-def *max-ipv4-addr-max-word*)

192.168.0.0/24

definition *ipv4range-set-from-bitmask*::*ipv4addr* \Rightarrow *nat* \Rightarrow *ipv4addr* set **where**
ipv4range-set-from-bitmask *addr* *bitmask* \equiv *ipv4range-set-from-netmask* *addr*
(*of-bl* ((*replicate* *bitmask* *True*) @ (*replicate* (32 - *bitmask*) *False*)))

lemma (*replicate* 3 *True*) = [*True*, *True*, *True*] **by** *eval*

lemma *of-bl* (*replicate* 3 *True*) = (7::*ipv4addr*) **by** *eval*

lemma *ipv4range-set-from-bitmask-alt1*:

```

    ipv4range-set-from-bitmask addr bitmask = ipv4range-set-from-netmask addr
((mask bitmask) << (32 - bitmask))
    apply(simp add: ipv4range-set-from-bitmask-def mask-bl)
    apply(simp add: Word.shiftl-of-bl)
done

```

```

lemma ipv4range-set-from-bitmask (ipv4addr-of-dotteddecimal (192,168,0,42))
16 =

```

```

    {ipv4addr-of-dotteddecimal (192,168,0,0) .. ipv4addr-of-dotteddecimal
(192,168,255,255)}

```

```

    by(simp add: ipv4range-set-from-bitmask-def ipv4range-set-from-netmask-def ipv4addr-of-dotteddecimal.simps
ipv4addr-of-nat-def)

```

```

lemma ipv4range-set-from-bitmask-UNIV: ipv4range-set-from-bitmask 0 0 =
UNIV

```

```

    apply(simp add: ipv4range-set-from-bitmask-def ipv4range-set-from-netmask-def)
    by (metis max-ipv4-addr-max-word range-0-max-UNIV)

```

```

lemma ip-in-ipv4range-set-from-bitmask-UNIV: ip ∈ (ipv4range-set-from-bitmask
(ipv4addr-of-dotteddecimal (0, 0, 0, 0)) 0)

```

```

    by(simp add: ipv4addr-of-dotteddecimal.simps ipv4addr-of-nat-def ipv4range-set-from-bitmask-UNIV)

```

```

lemma ipv4range-set-from-bitmask-0: ipv4range-set-from-bitmask foo 0 = UNIV

```

```

    apply(rule)

```

```

    apply(simp-all)

```

```

    apply(simp add: ipv4range-set-from-bitmask-alt1 ipv4range-set-from-netmask-def
Let-def)

```

```

    apply(simp add: range-0-max-UNIV[symmetric] del: range-0-max-UNIV)

```

```

    apply(simp add: mask-def)

```

```

done

```

```

lemma ipv4range-set-from-bitmask-32: ipv4range-set-from-bitmask foo 32 = {foo}

```

```

    apply(rule)

```

```

    apply(simp-all)

```

```

    apply(simp-all add: ipv4range-set-from-bitmask-alt1 ipv4range-set-from-netmask-def
Let-def)

```

```

    apply(simp-all add: mask-def)

```

```

    apply(simp-all only: max-ipv4-addr-number[symmetric] max-ipv4-addr-max-word
Word.word-and-max)

```

```

    apply(simp-all add: word32-or-NOT4294967296)

```

```

done

```

```

lemma ipv4range-set-from-bitmask-alt: ipv4range-set-from-bitmask pre len = {(pre
AND ((mask len) << (32 - len))) .. pre OR (mask (32 - len))}

```

```

    apply(simp only: ipv4range-set-from-bitmask-alt1 ipv4range-set-from-netmask-def
Let-def)

```

```

    apply(subst Word.word-aa-dist)

```

```

    apply(simp only: word-or-not)

```

```

    apply(simp only: Word.word-and-max)

```

```

    apply(simp only: NOT-mask-len32)

```

```

done

```

making element check executable

```

lemma addr-in-ipv4range-set-from-netmask-code[code-unfold]:
  addr ∈ (ipv4range-set-from-netmask base netmask)  $\longleftrightarrow$  (base AND netmask)
 $\leq$  addr  $\wedge$  addr  $\leq$  (base AND netmask) OR (NOT netmask)
  by(simp add: ipv4range-set-from-netmask-def Let-def)
lemma addr-in-ipv4range-set-from-bitmask-code[code-unfold]: addr ∈ (ipv4range-set-from-bitmask
pre len)  $\longleftrightarrow$ 
  (pre AND ((mask len) << (32 - len)))  $\leq$  addr  $\wedge$  addr  $\leq$  pre OR
  (mask (32 - len))
unfolding ipv4range-set-from-bitmask-alt by simp

value ipv4addr-of-dotteddecimal (192,168,4,8) ∈ (ipv4range-set-from-bitmask
(ipv4addr-of-dotteddecimal (192,168,0,42)) 16)

```

```

datatype ipv4range = IPv4Range
  ipv4addr — start (inclusive)
  ipv4addr — end (inclusive)
  | IPv4Union ipv4range ipv4range

```

```

fun ipv4range-to-set :: ipv4range  $\Rightarrow$  ipv4addr set where
  ipv4range-to-set (IPv4Range start end) = {start .. end} |
  ipv4range-to-set (IPv4Union r1 r2) = (ipv4range-to-set r1)  $\cup$  (ipv4range-to-set
r2)

```

```

fun ipv4range-element where
  ipv4range-element el (IPv4Range s e) = (s  $\leq$  el  $\wedge$  el  $\leq$  e) |
  ipv4range-element el (IPv4Union r1 r2) = (ipv4range-element el r1  $\vee$  ipv4range-element
el r2)

```

```

lemma ipv4range-element-set-eq[simp]: ipv4range-element el rg = (el ∈ ipv4range-to-set
rg)
by(induction rg rule: ipv4range-element.induct) simp-all

```

```

fun ipv4range-union where ipv4range-union r1 r2 = IPv4Union r1 r2
lemma ipv4range-union-set-eq[simp]: ipv4range-to-set (ipv4range-union r1 r2)
= ipv4range-to-set r1  $\cup$  ipv4range-to-set r2 by simp

```

```

fun ipv4range-empty where
  ipv4range-empty (IPv4Range s e) = (e < s) |
  ipv4range-empty (IPv4Union r1 r2) = (ipv4range-empty r1  $\wedge$  ipv4range-empty
r2)
lemma ipv4range-empty-set-eq[simp]: ipv4range-empty r  $\longleftrightarrow$  ipv4range-to-set r
= {}
by(induction r) auto

```

```

fun ipv4range-optimize-empty where
  ipv4range-optimize-empty (IPv4Union r1 r2) = (let r1o = ipv4range-optimize-empty
r1 in (let r2o = ipv4range-optimize-empty r2 in (
    if ipv4range-empty r1o then r2o else (if ipv4range-empty r2o then r1o else

```



```

(IPv4Union r1o r2o)))) |
  ipv4range-optimize-empty r = r
lemma ipv4range-optimize-empty-set-eq[simp]: ipv4range-to-set (ipv4range-optimize-empty
r) = ipv4range-to-set r
  by(induction r) (simp-all add: Let-def)
lemma ipv4range-optimize-empty-double[simp]: ipv4range-optimize-empty (ipv4range-optimize-empty
r) = ipv4range-optimize-empty r
  apply(induction r)
  by(simp-all add: Let-def)
fun ipv4range-empty-shallow where
  ipv4range-empty-shallow (IPv4Range s e) = (e < s) |
  ipv4range-empty-shallow (IPv4Union - -) = False
lemma helper-optimize-shallow: ipv4range-empty (ipv4range-optimize-empty r)
= ipv4range-empty-shallow (ipv4range-optimize-empty r)
  by(induction r) fastforce+
fun ipv4range-optimize-empty2 where
  ipv4range-optimize-empty2 (IPv4Union r1 r2) = (let r1o = ipv4range-optimize-empty
r1 in (let r2o = ipv4range-optimize-empty r2 in (
    if ipv4range-empty-shallow r1o then r2o else (if ipv4range-empty-shallow r2o
then r1o else (IPv4Union r1o r2o)))))) |
  ipv4range-optimize-empty2 r = r
lemma ipv4range-optimize-empty-code[code-unfold]: ipv4range-optimize-empty =
ipv4range-optimize-empty2
  by (subst fun-eq-iff, clarify, rename-tac r, induct-tac r)
    (unfold ipv4range-optimize-empty.simps ipv4range-optimize-empty2.simps
Let-def helper-optimize-shallow[symmetric], simp-all)

fun ipv4range-to-list where
  ipv4range-to-list (IPv4Union r1 r2) = ipv4range-to-list r1 @ ipv4range-to-list
r2 |
  ipv4range-to-list r = (if ipv4range-empty r then [] else [r])

lemma fold ( $\lambda r s. s \cup \text{ipv4range-to-set } r$ ) (ipv4range-to-list rs) {} = ipv4range-to-set
rs
  apply(induction rs, simp)
  apply(subst ipv4range-to-list.simps(1))
  apply simp
  thm fold-append-concat-rev
  oops

lemma ipv4range-to-list-set-eq: ( $\bigcup \text{set } (\text{map } \text{ipv4range-to-set } (\text{ipv4range-to-list }
rs))$ ) = ipv4range-to-set rs
  by(induction rs) simp-all

fun list-to-ipv4range where
  list-to-ipv4range [r] = r |
  list-to-ipv4range (r#rs) = (IPv4Union r (list-to-ipv4range rs)) |
  list-to-ipv4range [] = IPv4Range 2 1

```

lemma *list-to-ipv4range-set-eq*: $(\bigcup \text{set } (\text{map } \text{ipv4range-to-set } rs)) = \text{ipv4range-to-set } (\text{list-to-ipv4range } rs)$

by(*induction* *rs* *rule*: *list-to-ipv4range.induct*) *simp-all*

fun *ipv4range-linearize* **where** *ipv4range-linearize* *rs* = *list-to-ipv4range* (*ipv4range-to-list* *rs*)

lemma *ipv4range-to-set* (*ipv4range-linearize* *r*) = *ipv4range-to-set* *r*

by(*simp*, *metis* *list-to-ipv4range-set-eq* *ipv4range-to-list-set-eq*)

fun *ipv4range-optimize-same* **where** *ipv4range-optimize-same* *rs* = *list-to-ipv4range* (*remdups* (*ipv4range-to-list* *rs*))

lemma *ipv4range-optimize-same-set-eq*[*simp*]: *ipv4range-to-set* (*ipv4range-optimize-same* *rs*) = *ipv4range-to-set* *rs*

by(*simp*, *subst* *list-to-ipv4range-set-eq*[*symmetric*]) (*metis* *image-set* *ipv4range-to-list-set-eq* *set-remdups*)

fun *ipv4range-is-simple* **where** *ipv4range-is-simple* (*IPv4Range* -) = *True* | *ipv4range-is-simple* (*IPv4Union* -) = *False*

fun *ipv4rangelist-union-free* **where**

ipv4rangelist-union-free (*r* # *rs*) = (*ipv4range-is-simple* *r* \wedge *ipv4rangelist-union-free* *rs*) |

ipv4rangelist-union-free [] = *True*

lemma *ipv4rangelist-union-freeX*: *ipv4rangelist-union-free* (*r* # *rs*) $\implies \exists s e. r = \text{IPv4Range } s e$

by (*induction* *rs*) (*cases* *r*, *simp*, *simp*) +

lemma *ipv4rangelist-union-free-append*: *ipv4rangelist-union-free* (*a* @ *b*) = (*ipv4rangelist-union-free* *a* \wedge *ipv4rangelist-union-free* *b*)

by (*induction* *a*) (*auto*)

lemma *ipv4range-to-list-union-free*: *l* = *ipv4range-to-list* *r* $\implies \text{ipv4rangelist-union-free } l$

by(*induction* *r* *arbitrary*: *l*) (*simp-all* *add*: *ipv4rangelist-union-free-append*)

fun *ipv4range-setminus* :: *ipv4range* \Rightarrow *ipv4range* \Rightarrow *ipv4range* **where**

ipv4range-setminus (*IPv4Range* *s* *e*) (*IPv4Range* *ms* *me*) = (

if *s* > *e* \vee *ms* > *me* *then* *IPv4Range* *s* *e* *else*

if *me* \geq *e*

then

IPv4Range (*if* *ms* = 0 *then* 1 *else* *s*) (*min* *e* (*ip-prev* *ms*))

else if *ms* \leq *s*

then

IPv4Range (*max* *s* (*ip-next* *me*)) (*if* *me* = *max-ipv4-addr* *then* 0 *else* *e*)

else

IPv4Union (*IPv4Range* (*if* *ms* = 0 *then* 1 *else* *s*) (*ip-prev* *ms*)) (*IPv4Range*

(*ip-next* *me*) (*if* *me* = *max-ipv4-addr* *then* 0 *else* *e*))

) |

ipv4range-setminus (*IPv4Union* *r1* *r2*) *t* = *IPv4Union* (*ipv4range-setminus* *r1*

t) (*ipv4range-setminus* *r2* *t*) |

ipv4range-setminus *t* (*IPv4Union* *r1* *r2*) = *ipv4range-setminus* (*ipv4range-setminus* *t* *r1*) *r2*


```

definition ipv4range-UNIV  $\equiv$  IPv4Range 0 max-ipv4-addr
lemma ipv4range-UNIV-set-eq[simp]: ipv4range-to-set ipv4range-UNIV = UNIV
  unfolding ipv4range-UNIV-def by simp

fun ipv4range-invert where ipv4range-invert r = ipv4range-setminus ipv4range-UNIV
r
lemma ipv4range-invert-set-eq[simp]: ipv4range-to-set (ipv4range-invert r) =
UNIV - ipv4range-to-set r by(auto)

lemma ipv4range-invert-UNIV-empty: ipv4range-empty (ipv4range-invert ipv4range-UNIV)
by simp

fun ipv4range-intersection where ipv4range-intersection r1 r2 =
  ipv4range-optimize-same (ipv4range-setminus (ipv4range-union r1 r2) (ipv4range-union
(ipv4range-invert r1) (ipv4range-invert r2)))
lemma ipv4range-intersection-set-eq[simp]: ipv4range-to-set (ipv4range-intersection
r1 r2) = ipv4range-to-set r1  $\cap$  ipv4range-to-set r2
  unfolding ipv4range-intersection.simps ipv4range-optimize-same-set-eq by auto

lemma ipv4range-setminus-intersection-empty-struct-rr:
  ipv4range-empty (ipv4range-intersection (IPv4Range r1s r1e) (IPv4Range r2s
r2e))  $\implies$ 
  ipv4range-setminus (IPv4Range r1s r1e) (IPv4Range r2s r2e) = (IPv4Range
r1s r1e)
  apply(subst(asm) ipv4range-empty-set-eq)
  apply(subst(asm) ipv4range-intersection-set-eq)
  apply(unfold ipv4range-to-set.simps(1))
  apply(cases ipv4range-empty (IPv4Range r1s r1e), case-tac [!]  
ipv4range-empty (IPv4Range r2s r2e))
  apply(unfold ipv4range-empty.simps(1))
  apply(force, force, force)
  apply(cases r1e < r2s)
  defer
  apply(subgoal-tac r2e < r1s)
  defer
  apply force
  apply(simp only: ipv4range-setminus.simps)
  apply(case-tac [!] r1e  $\leq$  r2e, case-tac [!] r2s  $\leq$  r1s)
  apply(auto)
  apply(metis (hide-lams, no-types) comm-semiring-1-class.normalizing-semiring-rules(24)
inc-i ip-prev-def le-minus min.absorb-iff1 word-le-sub1 word-zero-le)
  apply(metis inc-le ip-next-def max.order-iff)
done

declare ipv4range-intersection.simps[simp del]
declare ipv4range-setminus.simps(1)[simp del]

lemma ipv4range-setminus-intersection-empty-struct:
  ipv4range-empty (ipv4range-intersection r1 r2)  $\implies$ 

```

$ipv4range\text{-}setminus r1\ r2 = r1$
by (induction r1 r2 rule: $ipv4range\text{-}setminus.induct$, auto simp add: $ipv4range\text{-}setminus\text{-}intersection\text{-}empty$ -s fastforce

definition $ipv4range\text{-}subset\ r1\ r2 \equiv ipv4range\text{-}empty\ (ipv4range\text{-}setminus r1\ r2)$
lemma $ipv4range\text{-}subset\text{-}set\text{-}eq[simp]$: $ipv4range\text{-}subset\ r1\ r2 = (ipv4range\text{-}to\text{-}set\ r1 \subseteq ipv4range\text{-}to\text{-}set\ r2)$
unfolding $ipv4range\text{-}subset\text{-}def$ **by** simp

definition $ipv4range\text{-}eq$ **where**
 $ipv4range\text{-}eq\ r1\ r2 = (ipv4range\text{-}subset\ r1\ r2 \wedge ipv4range\text{-}subset\ r2\ r1)$
lemma $ipv4range\text{-}eq\text{-}set\text{-}eq$: $ipv4range\text{-}eq\ r1\ r2 \longleftrightarrow ipv4range\text{-}to\text{-}set\ r1 = ipv4range\text{-}to\text{-}set\ r2$
unfolding $ipv4range\text{-}eq\text{-}def$ **by** auto
thm $iffD1[OF\ ipv4range\text{-}eq\text{-}set\text{-}eq]$
declare $iffD1[OF\ ipv4range\text{-}eq\text{-}set\text{-}eq, simp]$
lemma $ipv4range\text{-}eq\text{-}comm$: $ipv4range\text{-}eq\ r1\ r2 \longleftrightarrow ipv4range\text{-}eq\ r2\ r1$
unfolding $ipv4range\text{-}eq\text{-}def$ **by** fast
lemma $ipv4range\text{-}to\text{-}set\text{-}alt$: $ipv4range\text{-}to\text{-}set\ r = \{x.\ ipv4range\text{-}element\ x\ r\}$
unfolding $ipv4range\text{-}element\text{-}set\text{-}eq$ **by** blast

lemma $ipv4range\text{-}un\text{-}empty$: $ipv4range\text{-}empty\ r1 \implies ipv4range\text{-}eq\ (ipv4range\text{-}union\ r1\ r2)\ r2$
by (subst $ipv4range\text{-}eq\text{-}set\text{-}eq$, simp)
lemma $ipv4range\text{-}un\text{-}empty\text{-}b$: $ipv4range\text{-}empty\ r2 \implies ipv4range\text{-}eq\ (ipv4range\text{-}union\ r1\ r2)\ r1$
by (subst $ipv4range\text{-}eq\text{-}set\text{-}eq$, simp)

lemma $ipv4range\text{-}Diff\text{-}triv$:
assumes $ipv4range\text{-}empty\ (ipv4range\text{-}intersection\ a\ b)$ **shows** $ipv4range\text{-}eq\ (ipv4range\text{-}setminus a\ b)\ a$
using $ipv4range\text{-}setminus\text{-}intersection\text{-}empty\text{-}struct[OF\ assms]$ $ipv4range\text{-}eq\text{-}set\text{-}eq[of\ a\ a]$ **by** simp

fun $ipv4range\text{-}size$ **where**
 $ipv4range\text{-}size\ (IPv4Union\ a\ b) = ipv4range\text{-}size\ a + ipv4range\text{-}size\ b$ |
 $ipv4range\text{-}size\ (IPv4Range\ s\ e) = (if\ s \leq e\ then\ 1\ else\ 0)$
lemma $ipv4range\text{-}size\ r = length\ (ipv4range\text{-}to\text{-}list\ r)$
by (induction r, simp-all)

lemma [simp]: $\exists x::ipv4range.\ y \in ipv4range\text{-}to\text{-}set\ x$
proof show $y \in ipv4range\text{-}to\text{-}set\ ipv4range\text{-}UNIV$ **by** simp **qed**

quotient-type $ipv4rq = ipv4range / ipv4range\text{-}eq$
by (unfold equivp-def, simp only: fun-eq-iff, unfold $ipv4range\text{-}eq\text{-}set\text{-}eq$) auto

lift-definition $ipv4rq\text{-}union :: ipv4rq \Rightarrow ipv4rq \Rightarrow ipv4rq$ **is** $IPv4Union$ **unfolding** $ipv4range\text{-}eq\text{-}set\text{-}eq$ **by** simp
lift-definition $ipv4rq\text{-}setminus :: ipv4rq \Rightarrow ipv4rq \Rightarrow ipv4rq$ **is** $ipv4range\text{-}setminus$

unfolding *ipv4range-eq-set-eq* **by** *simp*
lift-definition *ipv4rq-intersection* :: *ipv4rq* \Rightarrow *ipv4rq* \Rightarrow *ipv4rq* **is** *ipv4range-intersection*
unfolding *ipv4range-eq-set-eq* **by** *simp*
lift-definition *ipv4rq-empty* :: *ipv4rq* \Rightarrow *bool* **is** *ipv4range-empty* **unfolding**
ipv4range-eq-set-eq **by** *simp*
lift-definition *ipv4rq-element* :: *ipv4addr* \Rightarrow *ipv4rq* \Rightarrow *bool* **is** *ipv4range-element*
unfolding *ipv4range-eq-set-eq* **by** *simp*
lift-definition *ipv4rq-to-set* :: *ipv4rq* \Rightarrow *ipv4addr* *set* **is** *ipv4range-to-set* **unfolding**
ipv4range-eq-set-eq **by** *simp*
lift-definition *ipv4rq-UNIV* :: *ipv4rq* **is** *ipv4range-UNIV* .
lift-definition *ipv4rq-eq* :: *ipv4rq* \Rightarrow *ipv4rq* \Rightarrow *bool* **is** *ipv4range-eq* **unfolding**
ipv4range-eq-set-eq **by** *simp*
lemma *ipv4rq-setminus-set-eq*: *ipv4rq-to-set* (*ipv4rq-setminus* *r1* *r2*) = *ipv4rq-to-set*
r1 - *ipv4rq-to-set* *r2* **by** *transfer simp*
lemma *ipv4rq-intersection-set-eq*: *ipv4rq-to-set* (*ipv4rq-intersection* *r1* *r2*) =
ipv4rq-to-set *r1* \cap *ipv4rq-to-set* *r2* **by** *transfer simp*
lemma *ipv4rq-empty-set-eq*: *ipv4rq-empty* *r* = (*ipv4rq-to-set* *r* = {}) **by** *transfer*
simp
lemma *ipv4rq-element-set-eq*: *ipv4rq-element* *x* *r* = (*x* \in *ipv4rq-to-set* *r*) **by**
transfer simp
lemma *ipv4rq-UNIV-set-eq*: *ipv4rq-to-set* *ipv4rq-UNIV* = *UNIV* **by** *transfer simp*
lemmas *ipv4rq-eqs*[*simp*] = *ipv4rq-intersection-set-eq* *ipv4rq-setminus-set-eq* *ipv4rq-empty-set-eq*
ipv4rq-UNIV-set-eq

instantiation *ipv4rq* :: *equal*
begin
definition *equal-ipv4rq* *r1* *r2* = *ipv4rq-eq* *r1* *r2*
instance
proof
case *goal1* **thus** ?*case* **unfolding** *equal-ipv4rq-def* **by** *transfer simp*
qed
end

abbreviation *ipv4rq-abbr* :: *ipv4addr* \Rightarrow *ipv4addr* \Rightarrow *ipv4rq* ([*-*; *-*]) **where**
[*s*; *e*] \equiv *abs-ipv4rq* (*IPv4Range* *s* *e*)
abbreviation *ipv4un-abbr* :: *ipv4rq* \Rightarrow *ipv4rq* \Rightarrow *ipv4rq* (*-* \cup_{rg} *-*) **where**
r1 \cup_{rg} *r2* \equiv *ipv4rq-union* *r1* *r2*

lemma *rq-on-set*: *ipv4rq-to-set* *x* = *ipv4rq-to-set* *y* \longleftrightarrow *x* = *y*
by (*metis* *Quotient-ipv4rq* *Quotient-rel-rep* *ipv4range-eq-set-eq* *ipv4rq-to-set.rep-eq*)

fun *ipv4range-intersection2* :: *ipv4range* \Rightarrow *ipv4range* \Rightarrow *ipv4range* **where**
ipv4range-intersection2 (*IPv4Union* *a1* *a2*) (*IPv4Union* *b1* *b2*) = *IPv4Union*
(*IPv4Union* (*ipv4range-intersection2* *a1* *b1*) (*ipv4range-intersection2* *a1* *b2*))
(*IPv4Union* (*ipv4range-intersection2* *a2* *b1*) (*ipv4range-intersection2* *a2* *b2*)) |
ipv4range-intersection2 *r* (*IPv4Union* *r1* *r2*) = (*IPv4Union* (*ipv4range-intersection2*
r *r1*) (*ipv4range-intersection2* *r* *r2*)) |
ipv4range-intersection2 (*IPv4Union* *r1* *r2*) *r* = (*IPv4Union* (*ipv4range-intersection2*
r1 *r*) (*ipv4range-intersection2* *r2* *r*)) |

```

    ipv4range-intersection2 (IPv4Range s1 e1) (IPv4Range s2 e2) = IPv4Range (max
s1 s2) (min e1 e2)
  lemma ipv4range-intersection2-set-eq[simp]: ipv4range-to-set (ipv4range-intersection2
r1 r2) =
    ipv4range-to-set r1 ∩ ipv4range-to-set r2
  by (induction rule: ipv4range-intersection2.induct) auto

  lift-definition ipv4rq-intersection2 :: ipv4rq ⇒ ipv4rq ⇒ ipv4rq
  is ipv4range-intersection2 unfolding ipv4range-eq-set-eq by simp
  lemma ipv4rq-intersection2-set-eq: ipv4rq-to-set (ipv4rq-intersection2 r1 r2) =
    ipv4rq-to-set r1 ∩ ipv4rq-to-set r2
  by transfer simp

  lift-definition ipv4rq-optimize-empty :: ipv4rq ⇒ ipv4rq
  is ipv4range-optimize-empty unfolding ipv4range-eq-set-eq by simp
  lemma ipv4rq-optimize-empty-type-id: (ipv4rq-optimize-empty r) = r
  by (transfer, rename-tac rt, induct-tac rt)
    (unfold ipv4range-eq-set-eq, simp add: Let-def)+

  lemma ipv4rq-int-int2-code[code-unfold]: ipv4rq-intersection = (λx y. ipv4rq-optimize-empty
(ipv4rq-intersection2 x y))
  unfolding fun-eq-iff
  unfolding ipv4rq-optimize-empty-type-id rq-on-set[symmetric]
  unfolding ipv4range-intersection2-set-eq ipv4rq-intersection-set-eq
  by clarify

  lemma rule-ipv4rq-eq-set: ipv4rq-to-set x = ipv4rq-to-set y ⟹ x = y
  using ipv4range-eq-set-eq by transfer blast

  definition is-lowest-element x S = (x ∈ S ∧ (∀ y ∈ S. y ≤ x ⟶ y = x))
  lemma is-lowest-element-alt: (x ∈ S ∧ (∀ y ∈ S. x ≤ y)) = is-lowest-element x S
  unfolding is-lowest-element-def
  oops

  fun ipv4range-lowest-element where
    ipv4range-lowest-element (IPv4Range s e) = (if s ≤ e then Some s else None)
  |
    ipv4range-lowest-element (IPv4Union A B) = (case (ipv4range-lowest-element
A, ipv4range-lowest-element B) of
    (Some a, Some b) ⇒ Some (if a < b then a else b) |
    (None, Some b) ⇒ Some b |
    (Some a, None) ⇒ Some a |
    (None, None) ⇒ None)

  lemma ipv4range-lowest-none-empty: ipv4range-lowest-element r = None ⟷
    ipv4range-empty r
  by (induction r, simp-all, fastforce)

```

```

lemma ipv4range-lowest-element-correct-A: ipv4range-lowest-element  $r = \text{Some } x \implies \text{ipv4range-element } x \ r \wedge (\forall y \in \text{ipv4range-to-set } r. (y \leq x \longrightarrow y = x))$ 
  apply (induction  $r$  arbitrary:  $x$  rule: ipv4range-lowest-element.induct)
  apply (rename-tac  $rs$   $re$   $x$ , case-tac  $rs \leq re$ , auto)[1]
  apply (subst (asm) ipv4range-lowest-element.simps(2))
  apply (rename-tac  $A$   $B$   $x$ )
  apply (case-tac ipv4range-lowest-element  $B$ )
  apply (case-tac [!] ipv4range-lowest-element  $A$ )
  apply (simp-all add: ipv4range-lowest-none-empty)[3]
  apply fastforce
done

lemma smallerequalgreater:  $((y :: \text{ipv4addr}) \leq s \longrightarrow y = s) = (y \geq s)$  by fastforce

lemma somemcase:  $x = \text{Some } y \implies \text{case } x \text{ of } \text{None} \Rightarrow a \mid \text{Some } z \Rightarrow b \ z = b \ y$ 
by simp

lemma ipv4range-lowest-element-set-eq:
   $\neg \text{ipv4range-empty } r \implies$ 
   $(\text{ipv4range-lowest-element } r = \text{Some } x) = (\text{is-lowest-element } x \ (\text{ipv4range-to-set } r))$ 
  unfolding is-lowest-element-def
  apply (rule iffI)
  using ipv4range-lowest-element-correct-A ipv4range-lowest-none-empty apply simp
  apply (induction  $r$  arbitrary:  $x$  rule: ipv4range-lowest-element.induct)
  apply simp
  apply (rename-tac  $A$   $B$   $x$ )
  apply (case-tac ipv4range-lowest-element  $B$ )
  apply (case-tac [!] ipv4range-lowest-element  $A$ )
  apply (auto)[3]
  apply (subgoal-tac  $\neg \text{ipv4range-empty } A \wedge \neg \text{ipv4range-empty } B$ )
  prefer 2
  using arg-cong[where  $f = \text{Not}$ , OF ipv4range-lowest-none-empty] apply (simp, metis)
  apply (clarsimp simp add: ipv4range-lowest-none-empty)
  proof –
    fix  $A :: \text{ipv4range}$  and  $B :: \text{ipv4range}$  and  $xa :: 32 \text{ word}$  and  $a :: 32 \text{ word}$ 
and  $aa :: 32 \text{ word}$ 
    assume  $a1: \bigwedge x. x \in \text{ipv4range-to-set } B \wedge (\forall y \in \text{ipv4range-to-set } B. y \leq x \longrightarrow y = x) \implies a = x$ 
    assume  $a2: \text{ipv4range-lowest-element } B = \text{Some } a$ 
    assume  $a3: \text{ipv4range-lowest-element } A = \text{Some } aa$ 
    assume  $a4: xa \in \text{ipv4range-to-set } A \vee xa \in \text{ipv4range-to-set } B$ 
    assume  $a5: \forall y \in \text{ipv4range-to-set } A \cup \text{ipv4range-to-set } B. y \leq xa \longrightarrow y = xa$ 
    obtain  $sk_0 :: 32 \text{ word} \Rightarrow 32 \text{ word}$  where  $f1: \forall x_0. x_0 \notin \text{ipv4range-to-set } B \vee sk_0 \ x_0 \in \text{ipv4range-to-set } B \wedge sk_0 \ x_0 \leq x_0 \wedge sk_0 \ x_0 \neq x_0 \vee a = x_0$ 
    using  $a1$  by (metis (lifting))
    have  $\forall x_0. x_0 \notin \{uub. uub \in \text{ipv4range-to-set } A \vee uub \in \text{ipv4range-to-set } B\}$ 

```


$\vee \neg x_0 \leq xa \vee xa = x_0$
using *a5* **by** *blast*
hence *f2*: $\forall x_0. \neg (x_0 \in \text{ipv4range-to-set } A \vee x_0 \in \text{ipv4range-to-set } B) \vee xa$
 $= x_0 \vee \neg x_0 \leq xa$
by *blast*
hence $xa \notin \text{ipv4range-to-set } B \vee a = xa$
using *f1* **by** (*metis* (*lifting*))
hence $aa = xa \vee a = xa$
using *f2 a3 a4* **by** (*metis* (*lifting*) *ipv4range-element-set-eq ipv4range-lowest-element-correct-A le-less-linear less-asym'*)
thus $(aa < a \longrightarrow aa = xa) \wedge (\neg aa < a \longrightarrow a = xa)$
using *a2 f2 a3* **by** (*metis* (*lifting*) *ipv4range-element-set-eq ipv4range-lowest-element-correct-A le-less-linear less-asym'*)
qed

lift-definition *ipv4rq-lowest-element* :: *ipv4rq* \Rightarrow *ipv4addr option* **is** *ipv4range-lowest-element*
unfolding *ipv4range-eq-set-eq*

proof –
fix *r1 r2*
assume *eq*: *ipv4range-to-set r1* = *ipv4range-to-set r2*
show *ipv4range-lowest-element r1* = *ipv4range-lowest-element r2*
proof(*cases ipv4range-empty r1*)
case *True*
moreover
with *eq* **have** *ipv4range-empty r2* **by** *simp*
ultimately
have *ipv4range-lowest-element r1* = *None* *ipv4range-lowest-element r2* =
None
using *ipv4range-lowest-none-empty[symmetric]* **by** *simp-all*
then show *?thesis* ..
next
case *False*
with *eq* **have** *False2*: $\neg \text{ipv4range-empty } r2$ **by** *simp*
note *ipv4range-lowest-element-set-eq[OF False] ipv4range-lowest-element-set-eq[OF False2]*
with *eq* **show** *?thesis*
by (*metis not-Some-eq*)
qed
qed

lemma *ipv4rq-lowest-element-set-eq*:

$\neg \text{ipv4rq-empty } r \Longrightarrow$
 $(\text{ipv4rq-lowest-element } r = \text{Some } x) = (\text{is-lowest-element } x (\text{ipv4rq-to-set } r))$
by(*transfer, simp add: ipv4range-lowest-element-set-eq*)

lemma *ipv4rq-lowest-in*:

assumes $\neg \text{ipv4rq-empty } r$
shows *ipv4rq-element* (*the* (*ipv4rq-lowest-element r*)) *r*
using *assms* **by**(*transfer, metis ipv4range-lowest-element-correct-A ipv4range-lowest-none-empty*)

option.exhaust option.sel)

```

fun list-to-ipv4rq :: ipv4rq list  $\Rightarrow$  ipv4rq where
  list-to-ipv4rq [] = ipv4rq-setminus ipv4rq-UNIV ipv4rq-UNIV |
  list-to-ipv4rq [x] = x |
  list-to-ipv4rq (x#xs) = ipv4rq-union x (list-to-ipv4rq xs)
lemma list-to-ipv4rq-set-eq[simp]: ipv4rq-to-set (list-to-ipv4rq rs) = ( $\bigcup$  set (map
  ipv4rq-to-set rs))
  apply(induction rs rule: list-to-ipv4rq.induct)
  apply(simp-all)
oops

```

```

end
theory NumberWangCaesar
imports IPv4Addr
  ./autocorres-0.98/lib/WordLemmaBucket
begin

```

```

type-synonym prefix-match = (ipv4addr  $\times$  nat)
abbreviation pfxm-prefix p  $\equiv$  fst p
abbreviation pfxm-length p  $\equiv$  snd p
abbreviation pfxm-mask x  $\equiv$  mask (32 - pfxm-length x)

```

```

definition valid-prefix where
  valid-prefix pf = ((pfxm-mask pf) AND pfxm-prefix pf = 0)
lemma valid-prefix-E: valid-prefix pf  $\implies$  ((pfxm-mask pf) AND pfxm-prefix pf =
  0)
  unfolding valid-prefix-def .
lemma valid-prefix-alt-def: valid-prefix p = (pfxm-prefix p AND ( $2^{\wedge}(32 - \text{pfxm-length } p) - 1$ ) = 0)
  unfolding valid-prefix-def
  unfolding mask-def
  using word-bw-comms(1)
  arg-cong[where f =  $\lambda x. (\text{pfxm-prefix } p \text{ AND } x - 1 = 0)$ ]
  shiftl-1
  by metis

```

1.3 Address Semantics

```

definition prefix-match-semantics where
  prefix-match-semantics m a = (pfxm-prefix m = (NOT pfxm-mask m) AND a)

```

```

lemma mask-32-max-word: mask 32 = (max-word :: 32 word) by eval

```

1.4 Set Semantics

```

definition prefix-to-ipset :: prefix-match  $\Rightarrow$  ipv4addr set where

```

$\text{prefix-to-ipset } pfx = \{pfxm\text{-prefix } pfx \text{ .. } pfxm\text{-prefix } pfx \text{ OR } pfxm\text{-mask } pfx\}$

lemma *pfx-not-empty*: $\text{valid-prefix } pfx \implies \text{prefix-to-ipset } pfx \neq \{\}$
unfolding *valid-prefix-def prefix-to-ipset-def* **by** (*simp add: le-word-or2*)

definition *ipset-prefix-match* **where**

$\text{ipset-prefix-match } pfx \text{ } rg = (\text{let } pfxrg = \text{prefix-to-ipset } pfx \text{ in } (rg \cap pfxrg, rg - pfxrg))$

lemma *ipset-prefix-match-m*[*simp*]: $\text{fst } (\text{ipset-prefix-match } pfx \text{ } rg) = rg \cap (\text{prefix-to-ipset } pfx)$ **by** (*simp only: Let-def ipset-prefix-match-def, simp*)

lemma *ipset-prefix-match-nm*[*simp*]: $\text{snd } (\text{ipset-prefix-match } pfx \text{ } rg) = rg - (\text{prefix-to-ipset } pfx)$ **by** (*simp only: Let-def ipset-prefix-match-def, simp*)

lemma *ipset-prefix-match-distinct*: $\text{rpm} = \text{ipset-prefix-match } pfx \text{ } rg \implies (\text{fst } \text{rpm}) \cap (\text{snd } \text{rpm}) = \{\}$ **by** *force*

lemma *ipset-prefix-match-complete*: $\text{rpm} = \text{ipset-prefix-match } pfx \text{ } rg \implies (\text{fst } \text{rpm}) \cup (\text{snd } \text{rpm}) = rg$ **by** *force*

lemma *rpm-m-dup-simp*: $rg \cap \text{fst } (\text{ipset-prefix-match } (\text{routing-match } r) \text{ } rg) = \text{fst } (\text{ipset-prefix-match } (\text{routing-match } r) \text{ } rg)$
by *simp*

1.5 Equivalence Proofs

lemma *helper1*: $\text{NOT } (0::32 \text{ word}) = x_{19} \text{ OR } \text{NOT } x_{19}$ **using** *word-bool-alg.double-compl*
by *simp*

lemma *helper2*: $(x_0::32 \text{ word}) \text{ AND } \text{NOT } 0 = x_0$ **by** *simp*

lemma *helper3*: $(x_{48}::32 \text{ word}) \text{ OR } x_{49} = x_{48} \text{ OR } x_{49} \text{ AND } \text{NOT } x_{48}$ **using** *helper1 helper2* **by** (*metis word-oa-dist2*)

lemma *packet-ipset-prefix-eq1*:

assumes $\text{addr} \in \text{addrrg}$

assumes *valid-prefix match*

assumes $\neg \text{prefix-match-antics match addr}$

shows $\text{addr} \in (\text{snd } (\text{ipset-prefix-match match addrrg}))$

using *assms*

proof —

have $\text{pfxm-prefix match} \leq \text{addr} \implies \neg \text{addr} \leq \text{pfxm-prefix match OR pfxm-mask match}$

proof —

case *goal1*

have *a1*: $\text{pfxm-mask match AND pfxm-prefix match} = 0$

using *assms(2)* **unfolding** *valid-prefix-def* .

have *a2*: $\text{pfxm-prefix match} \neq \text{NOT pfxm-mask match AND addr}$

using *assms(3)* **unfolding** *prefix-match-antics-def* .

have *f1*: $\text{pfxm-prefix match} = \text{pfxm-prefix match AND NOT pfxm-mask match}$

using *a1* **by** (*metis mask-eq-0-eq-x word-bw-comms(1)*)

hence *f2*: $\forall x_{11}. (\text{pfxm-prefix match OR } x_{11}) \text{ AND NOT pfxm-mask match} = \text{pfxm-prefix match OR } x_{11} \text{ AND NOT pfxm-mask match}$

by (*metis word-bool-alg.conj-disj-distrib2*)

moreover

```

    { assume  $\neg \text{pf}\bar{x}\text{m-prefix match} \leq \text{addr AND NOT pf}\bar{x}\text{m-mask match}$ 
      hence  $\neg (\text{pf}\bar{x}\text{m-prefix match} \leq \text{addr} \wedge \text{addr} \leq \text{pf}\bar{x}\text{m-prefix match OR pf}\bar{x}\text{m-mask match})$ 
    }
    using f1 neg-mask-mono-le by metis }
  moreover
    { assume  $\text{pf}\bar{x}\text{m-prefix match} \leq \text{addr AND NOT pf}\bar{x}\text{m-mask match} \wedge \text{addr AND NOT pf}\bar{x}\text{m-mask match} \neq (\text{pf}\bar{x}\text{m-prefix match OR pf}\bar{x}\text{m-mask match}) \text{ AND NOT pf}\bar{x}\text{m-mask match}$ 
      hence  $\exists x_0. \neg \text{addr AND NOT mask } x_0 \leq (\text{pf}\bar{x}\text{m-prefix match OR pf}\bar{x}\text{m-mask match}) \text{ AND NOT mask } x_0$ 
      using f2 by (metis dual-order.antisym word-bool-alg.conj-cancel-right word-log-esimps(3))
      hence  $\neg (\text{pf}\bar{x}\text{m-prefix match} \leq \text{addr} \wedge \text{addr} \leq \text{pf}\bar{x}\text{m-prefix match OR pf}\bar{x}\text{m-mask match})$ 
      using neg-mask-mono-le by auto }
    ultimately show ?case
    using a2 by (metis goal1 word-bool-alg.conj-cancel-right word-bool-alg.conj-commute word-log-esimps(3))
  qed
  from this show ?thesis using assms(1)
  unfolding ipset-prefix-match-def Let-def snd-conv prefix-to-ipset-def
  by simp
qed

```

```

lemma packet-ipset-prefix-eq2:
  assumes  $\text{addr} \in \text{addr}\bar{r}\bar{g}$ 
  assumes valid-prefix match
  assumes prefix-match-semantics match addr
  shows  $\text{addr} \in (\text{fst } (\text{ipset-prefix-match match addr}\bar{r}\bar{g}))$ 
using assms
  apply(subst ipset-prefix-match-def)
  apply(simp only: Let-def fst-def Case-def)
  apply(simp add: prefix-to-ipset-def)
  apply(transfer)
  apply(simp only: prefix-match-semantics-def valid-prefix-def)
  apply(simp add: word-and-le1)
  apply(metis helper3 le-word-or2 word-bw-comms(1) word-bw-comms(2))
done

```

```

lemma packet-ipset-prefix-eq3:
  assumes  $\text{addr} \in \text{addr}\bar{r}\bar{g}$ 
  assumes valid-prefix match
  assumes  $\text{addr} \in (\text{snd } (\text{ipset-prefix-match match addr}\bar{r}\bar{g}))$ 
  shows  $\neg \text{prefix-match-semantics match addr}$ 
using assms
  apply(subst(asm) ipset-prefix-match-def)
  apply(simp only: Let-def fst-def Case-def)
  apply(simp)
  apply(subst(asm) prefix-to-ipset-def)

```

```

    apply(transfer)
    apply(simp only: prefix-match-semantics-def valid-prefix-def Set-Interval.ord-class.atLeastAtMost-iff
prefix-to-ipset-def)
    apply(simp)
    apply(metis helper3 le-word-or2 word-and-le2 word-bw-comms(1) word-bw-comms(2))
done

```

```

lemma packet-ipset-prefix-eq4:
  assumes addr ∈ addrrg
  assumes valid-prefix match
  assumes addr ∈ (fst (ipset-prefix-match match addrrg))
  shows prefix-match-semantics match addr
using assms
proof -
  have pfxm-prefix match = NOT pfxm-mask match AND addr
  proof -
    have a1: pfxm-mask match AND pfxm-prefix match = 0 using assms(2)
  unfolding valid-prefix-def .
    have a2: pfxm-prefix match ≤ addr ∧ addr ≤ pfxm-prefix match OR pfxm-mask
match
    using assms(3) unfolding ipset-prefix-match-def Let-def fst-conv prefix-to-ipset-def
  by simp
    have f2: ∀ x0. pfxm-prefix match AND NOT mask x0 ≤ addr AND NOT mask
x0
    using a2 neg-mask-mono-le by blast
    have f3: ∀ x0. addr AND NOT mask x0 ≤ (pfxm-prefix match OR pfxm-mask
match) AND NOT mask x0
    using a2 neg-mask-mono-le by blast
    have f4: pfxm-prefix match = pfxm-prefix match AND NOT pfxm-mask match
    using a1 by (metis mask-eq-0-eq-x word-bw-comms(1))
    hence f5: ∀ x6. (pfxm-prefix match OR x6) AND NOT pfxm-mask match =
pfxm-prefix match OR x6 AND NOT pfxm-mask match
    using word-ao-dist by (metis)
    have f6: ∀ x2 x3. addr AND NOT mask x2 ≤ x3 ∨ ¬ (pfxm-prefix match OR
pfxm-mask match) AND NOT mask x2 ≤ x3
    using f3 dual-order.trans by auto
    have pfxm-prefix match = (pfxm-prefix match OR pfxm-mask match) AND
NOT pfxm-mask match
    using f5 by auto
    hence pfxm-prefix match = addr AND NOT pfxm-mask match
    using f2 f4 f6 by (metis eq-iff)
    thus pfxm-prefix match = NOT pfxm-mask match AND addr
    by (metis word-bw-comms(1))
  qed
  from this show ?thesis unfolding prefix-match-semantics-def .
qed

```

```

lemma packet-ipset-prefix-eq24:
  assumes addr ∈ addrrg

```

```

assumes valid-prefix match
shows prefix-match-semantics match addr = (addr ∈ (fst (ipset-prefix-match
match addrrg)))
using packet-ipset-prefix-eq2[OF assms] packet-ipset-prefix-eq4[OF assms] by fast

lemma packet-ipset-prefix-eq13:
assumes addr ∈ addrrg
assumes valid-prefix match
shows  $\neg$ prefix-match-semantics match addr = (addr ∈ (snd (ipset-prefix-match
match addrrg)))
using packet-ipset-prefix-eq1[OF assms] packet-ipset-prefix-eq3[OF assms] by fast

lemma prefix-match-if-in-my-set: assumes valid-prefix pfx
shows prefix-match-semantics pfx (a :: ipv4addr)  $\longleftrightarrow$  a ∈ prefix-to-ipset pfx
using packet-ipset-prefix-eq24[OF - assms]
by (metis (erased, hide-lams) Int-iff UNIV-I fst-conv ipset-prefix-match-def)

lemma prefix-match-if-in-corny-set:
assumes valid-prefix pfx
shows prefix-match-semantics pfx (a :: ipv4addr)  $\longleftrightarrow$  a ∈ ipv4range-set-from-netmask
(pfxm-prefix pfx) (NOT pfxm-mask pfx)
unfolding prefix-match-if-in-my-set[OF assms]
unfolding prefix-to-ipset-def ipv4range-set-from-netmask-def Let-def
unfolding word-bool-alg.double-compl
proof –
  case goal1
  have *: pfxm-prefix pfx AND NOT pfxm-mask pfx = pfxm-prefix pfx
  unfolding mask-eq-0-eq-x[symmetric] using valid-prefix-E[OF assms] word-bw-comms(1)[of
pfxm-prefix pfx] by simp
  hence **: pfxm-prefix pfx AND NOT pfxm-mask pfx OR pfxm-mask pfx =
pfxm-prefix pfx OR pfxm-mask pfx
  by simp
  show ?case unfolding * ** ..
qed

lemma ipv4addr-and-maskshift-eq-and-not-mask: (base::32 word) AND (mask m
<< 32 – m) = base AND NOT mask (32 – m)
apply word-bitwise
apply (subgoal-tac m > 32 ∨ m ∈ set (map nat (upto 0 32)))
apply (simp add: upto-code upto-aux-rec, elim disjE)
apply (simp add: size-mask-32word)
apply (simp-all add: size-mask-32word) [33]
apply (simp add: upto-code upto-aux-rec, presburger)
done

lemma maskshift-eq-not-mask: ((mask m << 32 – m) :: 32 word) = NOT mask
(32 – m)

```

[illegible]

```

apply satx
apply satx
apply satx
apply satx
apply satx
apply satx
apply satx
apply satx
apply satx
apply (simp add: upto-code upto-aux-rec, presburger)
done

```

lemma *size-mask-32word'*: $\text{size } ((\text{mask } (32 - m)) :: 32 \text{ word}) = 32$ **by** (simp add: word-size)

```

lemma helper-32-case-split:  $32 < m \vee m \in \text{set } (\text{map nat } [0..32])$ 
  by (simp add: upto-code upto-aux-rec, presburger)
lemma ipv4addr-andnot-impl-takem:  $(a :: 32 \text{ word}) \text{ AND NOT mask } (32 - m) =$ 
 $b \implies (\text{take } (m) (\text{to-bl } a)) = (\text{take } (m) (\text{to-bl } b))$ 
  apply word-bitwise
  apply (subgoal-tac  $m > 32 \vee m \in \text{set } (\text{map nat } (\text{upto } 0 \ 32)))$ 
  prefer 2
  apply (simp only: helper-32-case-split)
  apply (simp add: upto-code upto-aux-rec, elim disjE)
  apply (simp add: size-mask-32word size-mask-32word')
  apply (simp-all add: size-mask-32word size-mask-32word') [33]
done

```

definition *ip-set* :: $32 \text{ word} \Rightarrow \text{nat} \Rightarrow 32 \text{ word set}$ **where** $\text{ip-set } i \ r = \{j \mid i \text{ AND NOT mask } (32 - r) = j \text{ AND NOT mask } (32 - r)\}$

lemma $(m1 \vee m2) \wedge (m3 \vee m4) \longleftrightarrow (m1 \wedge m3) \vee (m1 \wedge m4) \vee (m2 \wedge m3) \vee (m2 \wedge m4)$
by blast

lemmas *caesar-proof-unfolded* = *prefix-match-if-in-corny-set*[*unfolded valid-prefix-def*
prefix-match-semantics-def *Let-def*, *symmetric*]

lemma *caesar-proof-without-structures*: $\text{mask } (32 - l) \text{ AND pfxm-p} = 0 \implies$
 $(a \in \text{ipv4range-set-from-netmask } (\text{pfxm-p}) (\text{NOT mask } (32 - l))) =$
 $(\text{pfxm-p} = \text{NOT mask } (32 - l) \text{ AND } a)$
using *caesar-proof-unfolded* **by** force

lemma *mask-and-not-mask-helper*: $\text{mask } (32 - m) \text{ AND base AND NOT mask } (32 - m) = 0$
by (simp add: word-bw-lcs)

lemma *ipv4range-set-from-netmask-base-mask-consume*:
 $\text{ipv4range-set-from-netmask } (\text{base AND NOT mask } (32 - m)) (\text{NOT mask } (32 - m)) =$


```

    ipv4range-set-from-netmask base (NOT mask (32 - m))
  unfolding ipv4range-set-from-netmask-def
  by (simp add: AND-twice)

lemma ipv4range-set-from-bitmask-eq-ip-set: ipv4range-set-from-bitmask base m =
ip-set base m
  unfolding ip-set-def
  unfolding set-eq-iff
  unfolding mem-Collect-eq
  unfolding ipv4range-set-from-bitmask-alt1
  unfolding maskshift-eq-not-mask
  using caesar-proof-without-structures[OF mask-and-not-mask-helper, of - base m]
  unfolding ipv4range-set-from-netmask-base-mask-consume
  unfolding word-bw-comms(1)[of - ~ mask (32 - m)]
  ..

end
theory NumberWangCebewee
imports
  ./autocorres-0.98/lib/WordLemmaBucket
  NumberWangCaesar
begin

lemma and-not-mask-twice:
  (w && ~ mask n) && ~ mask m = w && ~ mask (max m n)
  apply (simp add: and-not-mask)
  apply (case-tac n < m)
  apply (simp-all add: shiftl-shiftr2 shiftl-shiftr1 not-less max-def
    shiftr-shiftr shiftl-shiftl)
  apply (cut-tac and-mask-shiftr-comm
    [where w=w and m=size w and n=m, simplified,symmetric])
  apply (simp add: word-size mask-def)
  apply (cut-tac and-mask-shiftr-comm
    [where w=w and m=size w and n=n, simplified,symmetric])
  apply (simp add: word-size mask-def)
  done

lemma X: j ∈ ip-set i r ⇒ ip-set j r = ip-set i r
  by (auto simp: ip-set-def)

lemma Z:
  fixes i :: ('a :: len) word
  assumes r2 ≤ r1 i && ~ mask r2 = x && ~ mask r2
  shows i && ~ mask r1 = x && ~ mask r1
  proof -
    have i AND NOT mask r1 = (i && ~ mask r2) && ~ mask r1 (is - = ?w

```

```

&& -)
  using ⟨r2 ≤ r1⟩ by (simp add: and-not-mask-twice max-def)
  also have ?w = x && ~ mask r2 by fact
  also have ... && ~ mask r1 = x && ~ mask r1
  using ⟨r2 ≤ r1⟩ by (simp add: and-not-mask-twice max-def)
  finally show ?thesis .
qed

lemma Y: r1 ≤ r2 ⟹ ip-set i r2 ⊆ ip-set i r1
  unfolding ip-set-def
  apply auto
  apply (rule Z[where ?r2.0=32 - r2])
  apply auto
  done

lemma ip-set-intersect-subset-helper:
  fixes i1 r1 i2 r2
  assumes disj: ip-set i1 r1 ∩ ip-set i2 r2 ≠ {} and r1 ≤ r2
  shows ip-set i2 r2 ⊆ ip-set i1 r1
  proof -
    from disj obtain j where j ∈ ip-set i1 r1 j ∈ ip-set i2 r2 by auto
    with ⟨r1 ≤ r2⟩ have j ∈ ip-set j r1 j ∈ ip-set j r1 using X Y by blast+

    show ip-set i2 r2 ⊆ ip-set i1 r1
    proof
      fix i assume i ∈ ip-set i2 r2
      with ⟨j ∈ ip-set i2 r2⟩ have i ∈ ip-set j r2 using X by auto
      also have ip-set j r2 ⊆ ip-set j r1 using ⟨r1 ≤ r2⟩ Y by blast
      also have ... = ip-set i1 r1 using ⟨j ∈ ip-set i1 r1⟩ X by blast
      finally show i ∈ ip-set i1 r1 .
    qed
  qed

lemma ip-set-notsubset-empty-inter: ¬ ip-set i1 r1 ⊆ ip-set i2 r2 ⟹ ¬ ip-set i2
r2 ⊆ ip-set i1 r1 ⟹ ip-set i1 r1 ∩ ip-set i2 r2 = {}
  apply(cases r1 ≤ r2)
  using ip-set-intersect-subset-helper apply blast
  apply(cases r2 ≤ r1)
  using ip-set-intersect-subset-helper apply blast
  apply(simp)
  done

end
theory Numberwang-Ln
imports NumberWangCebewee
begin

```

```

lemma ipv4range-bitmask-intersect:  $\neg \text{ipv4range-set-from-bitmask } b2 \ m2 \subseteq \text{ipv4range-set-from-bitmask } b1 \ m1 \implies$ 
 $\neg \text{ipv4range-set-from-bitmask } b1 \ m1 \subseteq \text{ipv4range-set-from-bitmask } b2 \ m2 \implies$ 
 $\text{ipv4range-set-from-bitmask } b1 \ m1 \cap \text{ipv4range-set-from-bitmask } b2 \ m2 = \{\}$ 
apply (simp add: ipv4range-set-from-bitmask-eq-ip-set)
using ip-set-notsubset-empty-inter
by presburger

```

end