

# Bitmagic

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**theory** *WordInterval*

**imports** *Main*

*NumberWang*

*~~/src/HOL/Word/Word*

*~~/src/HOL/Library/Code-Target-Nat*

**begin**

Intervals of consecutive words

**value**  $(2::nat) < 2^{32}$

**datatype**  $('a::len)$  *wordinterval* = *WordInterval*

$('a::len)$  *word* — start (inclusive)

$('a::len)$  *word* — end (inclusive)

| *RangeUnion*  $'a$  *wordinterval*  $'a$  *wordinterval*

**fun** *wordinterval-to-set* ::  $'a::len$  *wordinterval*  $\Rightarrow$   $('a::len$  *word*) *set* **where**

*wordinterval-to-set* (*WordInterval* *start end*) =  $\{start .. end\}$  |

*wordinterval-to-set* (*RangeUnion* *r1 r2*) = (*wordinterval-to-set* *r1*)  $\cup$  (*wordinterval-to-set* *r2*)

**fun** *wordinterval-element* ::  $'a::len$  *word*  $\Rightarrow$   $'a::len$  *wordinterval*  $\Rightarrow$  *bool* **where**

*wordinterval-element* *el* (*WordInterval* *s e*) =  $(s \leq el \wedge el \leq e)$  |

*wordinterval-element* *el* (*RangeUnion* *r1 r2*) = (*wordinterval-element* *el* *r1*  $\vee$  *wordinterval-element* *el* *r2*)

```

lemma wordinterval-element-set-eq[simp]: wordinterval-element el rg = (el ∈
wordinterval-to-set rg)
  by(induction rg rule: wordinterval-element.induct) simp-all

fun wordinterval-union :: 'a::len wordinterval ⇒ 'a::len wordinterval ⇒ 'a::len
wordinterval where
  wordinterval-union r1 r2 = RangeUnion r1 r2
lemma wordinterval-union-set-eq[simp]: wordinterval-to-set (wordinterval-union
r1 r2) = wordinterval-to-set r1 ∪ wordinterval-to-set r2 by simp

fun wordinterval-empty :: 'a::len wordinterval ⇒ bool where
  wordinterval-empty (WordInterval s e) = (e < s) |
  wordinterval-empty (RangeUnion r1 r2) = (wordinterval-empty r1 ∧ wordinterval-empty
r2)
lemma wordinterval-empty-set-eq[simp]: wordinterval-empty r ⟷ wordinterval-to-set
r = {}
  by(induction r) auto

fun wordinterval-optimize-empty where
  wordinterval-optimize-empty (RangeUnion r1 r2) = (let r1o = wordinterval-optimize-empty
r1 in (let r2o = wordinterval-optimize-empty r2 in (
    if wordinterval-empty r1o then r2o else (if wordinterval-empty r2o then r1o
else (RangeUnion r1o r2o)))))) |
  wordinterval-optimize-empty r = r
lemma wordinterval-optimize-empty-set-eq[simp]: wordinterval-to-set (wordinterval-optimize-empty
r) = wordinterval-to-set r
  by(induction r) (simp-all add: Let-def)
lemma wordinterval-optimize-empty-double[simp]: wordinterval-optimize-empty
(wordinterval-optimize-empty r) = wordinterval-optimize-empty r
  apply(induction r)
  by(simp-all add: Let-def)
fun wordinterval-empty-shallow where
  wordinterval-empty-shallow (WordInterval s e) = (e < s) |
  wordinterval-empty-shallow (RangeUnion - -) = False
lemma helper-optimize-shallow: wordinterval-empty (wordinterval-optimize-empty
r) = wordinterval-empty-shallow (wordinterval-optimize-empty r)
  by(induction r) fastforce+
fun wordinterval-optimize-empty2 where
  wordinterval-optimize-empty2 (RangeUnion r1 r2) = (let r1o = wordinterval-optimize-empty
r1 in (let r2o = wordinterval-optimize-empty r2 in (
    if wordinterval-empty-shallow r1o then r2o else (if wordinterval-empty-shallow
r2o then r1o else (RangeUnion r1o r2o)))))) |
  wordinterval-optimize-empty2 r = r
lemma wordinterval-optimize-empty-code[code-unfold]: wordinterval-optimize-empty
= wordinterval-optimize-empty2
  by (subst fun-eq-iff, clarify, rename-tac r, induct-tac r)
  (unfold wordinterval-optimize-empty.simps wordinterval-optimize-empty2.simps
Let-def helper-optimize-shallow[symmetric], simp-all)

```

**definition** *Empty-WordInterval* :: 'a::len wordinterval **where** *Empty-WordInterval*  
 $\equiv$  *WordInterval* 1 0

**lemma** *wordinterval-empty-Empty-WordInterval*: *wordinterval-empty Empty-WordInterval*  
**by**(*simp add: Empty-WordInterval-def*)

**lemma** *Empty-WordInterval-set-eq[simp]*: *wordinterval-to-set Empty-WordInterval*  
 $= \{\}$  **by**(*simp add: Empty-WordInterval-def*)

**fun** *wordinterval-to-list* :: 'a::len wordinterval  $\Rightarrow$  ('a::len wordinterval) list **where**  
*wordinterval-to-list* (*RangeUnion* r1 r2) = *wordinterval-to-list* r1 @ *wordinterval-to-list* r2 |  
*wordinterval-to-list* r = (if *wordinterval-empty* r then [] else [r])

**lemma** *wordinterval-to-list-set-eq*: ( $\bigcup$  set (map *wordinterval-to-set* (*wordinterval-to-list* rs))) = *wordinterval-to-set* rs  
**by**(*induction rs*) *simp-all*

**fun** *list-to-wordinterval* **where**  
*list-to-wordinterval* [r] = r |  
*list-to-wordinterval* (r#rs) = (*RangeUnion* r (*list-to-wordinterval* rs)) |  
*list-to-wordinterval* [] = *Empty-WordInterval*

**lemma** *list-to-wordinterval-set-eq*: ( $\bigcup$  set (map *wordinterval-to-set* rs)) = *wordinterval-to-set* (*list-to-wordinterval* rs)  
**by**(*induction rs rule: list-to-wordinterval.induct*) *simp-all*

**lemma** *list-to-wordinterval-set-eq-simp[simp]*: *wordinterval-to-set* (*list-to-wordinterval* (a # as)) = *wordinterval-to-set* (*wordinterval-union* a (*list-to-wordinterval* as))  
**by**(*cases as*) *auto*

**fun** *wordinterval-linearize* **where** *wordinterval-linearize* rs = *list-to-wordinterval* (*wordinterval-to-list* rs)

**lemma** *wordinterval-to-set* (*wordinterval-linearize* r) = *wordinterval-to-set* r  
**by**(*simp, metis list-to-wordinterval-set-eq wordinterval-to-list-set-eq*)

**fun** *wordinterval-optimize-same* **where** *wordinterval-optimize-same* rs = *list-to-wordinterval* (*remdups* (*wordinterval-to-list* rs))

**lemma** *wordinterval-optimize-same-set-eq[simp]*: *wordinterval-to-set* (*wordinterval-optimize-same* rs) = *wordinterval-to-set* rs

**by**(*simp, subst list-to-wordinterval-set-eq[symmetric]*) (*metis image-set wordinterval-to-list-set-eq set-remdups*)

**fun** *wordinterval-is-simple* **where** *wordinterval-is-simple* (*WordInterval* - -) = *True* | *wordinterval-is-simple* (*RangeUnion* - -) = *False*

**fun** *wordintervalist-union-free* **where**  
*wordintervalist-union-free* (r#rs) = (*wordinterval-is-simple* r  $\wedge$  *wordintervalist-union-free* rs) |  
*wordintervalist-union-free* [] = *True*

**lemma** *wordintervalist-union-freeX*: *wordintervalist-union-free* (r # rs)  $\implies \exists$  s

```

e. r = WordInterval s e
  by (induction rs) (cases r, simp, simp)+
lemma wordintervalist-union-free-append: wordintervalist-union-free (a@b) =
(wordintervalist-union-free a  $\wedge$  wordintervalist-union-free b)
  by (induction a) (auto)
lemma wordinterval-to-list-union-free: l = wordinterval-to-list r  $\implies$  wordintervalist-union-free
l
  by (induction r arbitrary: l) (simp-all add: wordintervalist-union-free-append)

```

previous and next words addresses, without wrap around

```

definition word-next :: 'a::len word  $\Rightarrow$  'a::len word where
  word-next a  $\equiv$  if a = max-word then max-word else a + 1
definition word-prev :: 'a::len word  $\Rightarrow$  'a::len word where
  word-prev a  $\equiv$  if a = 0 then 0 else a - 1

```

```

lemma word-next (2:: 8 word) = 3 by eval
lemma word-prev (2:: 8 word) = 1 by eval
lemma word-prev (0:: 8 word) = 0 by eval

```

```

fun wordinterval-setminus :: 'a::len wordinterval  $\Rightarrow$  'a::len wordinterval  $\Rightarrow$  'a::len
wordinterval where
  wordinterval-setminus (WordInterval s e) (WordInterval ms me) = (
    if s > e  $\vee$  ms > me then WordInterval s e else
    if me  $\geq$  e
    then
      WordInterval (if ms = 0 then 1 else s) (min e (word-prev ms))
    else if ms  $\leq$  s
    then
      WordInterval (max s (word-next me)) (if me = max-word then 0 else e)
    else
      RangeUnion (WordInterval (if ms = 0 then 1 else s) (word-prev ms))
    (WordInterval (word-next me) (if me = max-word then 0 else e))
  ) |
  wordinterval-setminus (RangeUnion r1 r2) t = RangeUnion (wordinterval-setminus
r1 t) (wordinterval-setminus r2 t) |
  wordinterval-setminus t (RangeUnion r1 r2) = wordinterval-setminus (wordinterval-setminus
t r1) r2

```

```

lemma wordinterval-setminus-rr-set-eq[simp]: wordinterval-to-set(wordinterval-setminus
(WordInterval s e) (WordInterval ms me)) =
  wordinterval-to-set (WordInterval s e) - wordinterval-to-set (WordInterval ms
me)
  apply(simp only: wordinterval-setminus.simps)
  apply(case-tac e < s)
  apply simp
  apply(case-tac me < ms)
  apply simp
  apply(case-tac [!] e  $\leq$  me)

```



**lemma** *wordinterval-invert-set-eq*[simp]: *wordinterval-to-set* (*wordinterval-invert* *r*) = *UNIV* - *wordinterval-to-set* *r* **by** (*auto*)

**lemma** *wordinterval-invert-UNIV-empty*: *wordinterval-empty* (*wordinterval-invert* *wordinterval-UNIV*) **by** *simp*

**fun** *wordinterval-intersection* :: 'a::len *wordinterval*  $\Rightarrow$  'a::len *wordinterval*  $\Rightarrow$  'a::len *wordinterval* **where**  
*wordinterval-intersection* *r1* *r2* =  
*wordinterval-optimize-same* (*wordinterval-setminus* (*wordinterval-union* *r1* *r2*)  
(*wordinterval-union* (*wordinterval-invert* *r1*) (*wordinterval-invert* *r2*)))  
**lemma** *wordinterval-intersection-set-eq*[simp]: *wordinterval-to-set* (*wordinterval-intersection* *r1* *r2*) = *wordinterval-to-set* *r1*  $\cap$  *wordinterval-to-set* *r2*  
**unfolding** *wordinterval-intersection.simps* *wordinterval-optimize-same-set-eq*  
**by** *auto*

**lemma** *wordinterval-setminus-intersection-empty-struct-rr*:  
*wordinterval-empty* (*wordinterval-intersection* (*WordInterval* *r1s* *r1e*) (*WordInterval* *r2s* *r2e*))  $\implies$   
*wordinterval-setminus* (*WordInterval* *r1s* *r1e*) (*WordInterval* *r2s* *r2e*) = (*WordInterval* *r1s* *r1e*)  
**apply** (*subst* (*asm*) *wordinterval-empty-set-eq*)  
**apply** (*subst* (*asm*) *wordinterval-intersection-set-eq*)  
**apply** (*unfold* *wordinterval-to-set.simps* (1))  
**apply** (*cases* *wordinterval-empty* (*WordInterval* *r1s* *r1e*), *case-tac* [!]  
*wordinterval-empty* (*WordInterval* *r2s* *r2e*))  
**apply** (*unfold* *wordinterval-empty.simps* (1))  
**apply** (*force*, *force*, *force*)  
**apply** (*cases* *r1e* < *r2s*)  
**defer**  
**apply** (*subgoal-tac* *r2e* < *r1s*)  
**defer**  
**apply** *force*  
**apply** (*simp* *only*: *wordinterval-setminus.simps*)  
**apply** (*case-tac* [!] *r1e*  $\leq$  *r2e*, *case-tac* [!] *r2s*  $\leq$  *r1s*)  
**apply** (*auto*)  
**apply** (*metis* *add.commute* *inc-i* *le-minus* *min-absorb1* *word-le-sub1* *word-prev-def* *word-zero-le*)  
**apply** (*metis* *inc-le* *word-next-def* *max.order-iff*)  
**done**

**declare** *wordinterval-intersection.simps*[*simp del*]  
**declare** *wordinterval-setminus.simps*(1)[*simp del*]

**lemma** *wordinterval-setminus-intersection-empty-struct*:  
*wordinterval-empty* (*wordinterval-intersection* *r1* *r2*)  $\implies$   
*wordinterval-setminus* *r1* *r2* = *r1*  
**by** (*induction* *r1* *r2* *rule*: *wordinterval-setminus.induct*, *auto* *simp* *add*: *wordinterval-setminus-intersection-empty-struct*)  
*fastforce*

```

definition wordinterval-subset :: 'a::len wordinterval  $\Rightarrow$  'a::len wordinterval  $\Rightarrow$ 
bool where
  wordinterval-subset r1 r2  $\equiv$  wordinterval-empty (wordinterval-setminus r1 r2)
lemma wordinterval-subset-set-eq[simp]: wordinterval-subset r1 r2 = (wordinterval-to-set
r1  $\subseteq$  wordinterval-to-set r2)
  unfolding wordinterval-subset-def by simp

definition wordinterval-eq :: 'a::len wordinterval  $\Rightarrow$  'a::len wordinterval  $\Rightarrow$  bool
where
  wordinterval-eq r1 r2 = (wordinterval-subset r1 r2  $\wedge$  wordinterval-subset r2 r1)
lemma wordinterval-eq-set-eq: wordinterval-eq r1 r2  $\longleftrightarrow$  wordinterval-to-set r1
= wordinterval-to-set r2
  unfolding wordinterval-eq-def by auto
thm iffD1[OF wordinterval-eq-set-eq]
declare iffD1[OF wordinterval-eq-set-eq, simp]
lemma wordinterval-eq-comm: wordinterval-eq r1 r2  $\longleftrightarrow$  wordinterval-eq r2 r1
  unfolding wordinterval-eq-def by fast
lemma wordinterval-to-set-alt: wordinterval-to-set r = {x. wordinterval-element
x r}
  unfolding wordinterval-element-set-eq by blast

lemma wordinterval-un-empty: wordinterval-empty r1  $\implies$  wordinterval-eq (wordinterval-union
r1 r2) r2
  by (subst wordinterval-eq-set-eq, simp)
lemma wordinterval-un-empt-b: wordinterval-empty r2  $\implies$  wordinterval-eq (wordinterval-union
r1 r2) r1
  by (subst wordinterval-eq-set-eq, simp)

lemma wordinterval-Diff-triv:
  assumes wordinterval-empty (wordinterval-intersection a b) shows wordinterval-eq
(wordinterval-setminus a b) a
  using wordinterval-setminus-intersection-empty-struct[OF assms] wordinterval-eq-set-eq[of
a a] by simp

fun wordinterval-size where
  wordinterval-size (RangeUnion a b) = wordinterval-size a + wordinterval-size
b |
  wordinterval-size (WordInterval s e) = (if s  $\leq$  e then 1 else 0)
lemma wordinterval-size r = length (wordinterval-to-list r)
  by (induction r, simp-all)

lemma [simp]:  $\exists x::('a::len wordinterval). y \in \text{wordinterval-to-set } x$ 
proof show  $y \in \text{wordinterval-to-set wordinterval-UNIV}$  by simp qed

```

```

lemma wordinterval-eq-reflp:
  reflp wordinterval-eq
  apply(rule reflpI)
  by(simp only: wordinterval-eq-set-eq)
lemma wordinterval-eq-symp:
  symp wordinterval-eq
  apply(rule sympI)
  by(simp add: wordinterval-eq-comm)
lemma wordinterval-eq-transp:
  transp wordinterval-eq
  apply(rule transpI)
  by(simp only: wordinterval-eq-set-eq)

lemma wordinterval-eq-equivp:
  equivp wordinterval-eq
  by (auto intro: equivpI wordinterval-eq-reflp wordinterval-eq-symp wordinterval-eq-transp)

end

theory IPv4Addr
imports Main
  NumberWang
  WordInterval
  ~~ /src/HOL/Word/Word
  ~~ /src/HOL/Library/Code-Target-Nat
begin

value (2::nat) < 2^32

```

## 1 Modelling IPv4 Addresses

An IPv4 address is basically a 32 bit unsigned integer

```
type-synonym ipv4addr = 32 word
```

```
value 42 :: ipv4addr
```

```
value (42 :: ipv4addr) ≤ 45
```

Conversion between natural numbers and IPv4 addresses

```
definition nat-of-ipv4addr :: ipv4addr ⇒ nat where
  nat-of-ipv4addr a = unat a
```

```
definition ipv4addr-of-nat :: nat ⇒ ipv4addr where
  ipv4addr-of-nat n = of-nat n
```

```
lemma ((nat-of-ipv4addr (42::ipv4addr))::nat) = 42 by eval
```

```
lemma ((ipv4addr-of-nat (42::nat))::ipv4addr) = 42 by eval
```



The maximum IPv4 address

**definition** *max-ipv4-addr* :: *ipv4addr* **where**  
*max-ipv4-addr*  $\equiv$  *ipv4addr-of-nat* ( $(2^{32}) - 1$ )

**lemma** *max-ipv4-addr-number*: *max-ipv4-addr* = 4294967295  
**by** *eval*

**lemma** *max-ipv4-addr* = 0b11111111111111111111111111111111  
**by**(*fact max-ipv4-addr-number*)

**lemma** *max-ipv4-addr-max-word*: *max-ipv4-addr* = *max-word*  
**by**(*simp add: max-ipv4-addr-number max-word-def*)

**lemma** *max-ipv4-addr-max*[*simp*]:  $\forall a. a \leq \text{max-ipv4-addr}$   
**by**(*simp add: max-ipv4-addr-max-word*)

**lemma** *range-0-max-UNIV*:  $\text{UNIV} = \{0 .. \text{max-ipv4-addr}\}$   
**by**(*simp add: max-ipv4-addr-max-word*) *fastforce*

identity functions

**lemma** *nat-of-ipv4addr-ipv4addr-of-nat*:  $\llbracket n \leq \text{nat-of-ipv4addr max-ipv4-addr} \rrbracket$   
 $\implies \text{nat-of-ipv4addr (ipv4addr-of-nat } n) = n$

**by** (*metis ipv4addr-of-nat-def le-unat-uoI nat-of-ipv4addr-def*)

**lemma** *nat-of-ipv4addr-ipv4addr-of-nat-mod*: *nat-of-ipv4addr (ipv4addr-of-nat n)*  
 $= n \bmod 2^{32}$

**by**(*simp add: ipv4addr-of-nat-def nat-of-ipv4addr-def unat-of-nat*)

**lemma** *ipv4addr-of-nat-nat-of-ipv4addr*: *ipv4addr-of-nat (nat-of-ipv4addr addr)*  
 $= \text{addr}$

**by**(*simp add: ipv4addr-of-nat-def nat-of-ipv4addr-def*)

Equality of IPv4 addresses

**lemma**  $\llbracket n \leq \text{nat-of-ipv4addr max-ipv4-addr} \rrbracket \implies \text{nat-of-ipv4addr (ipv4addr-of-nat } n) = n$

**apply**(*simp add: nat-of-ipv4addr-def ipv4addr-of-nat-def*)

**apply**(*induction n*)

**apply**(*simp-all*)

**by**(*unat-arith*)

**lemma** *ipv4addr-of-nat-eq*:  $x = y \implies \text{ipv4addr-of-nat } x = \text{ipv4addr-of-nat } y$   
**by**(*simp add: ipv4addr-of-nat-def*)

## 1.1 Representing IPv4 Addresses

**fun** *ipv4addr-of-dotdecimal* ::  $\text{nat} \times \text{nat} \times \text{nat} \times \text{nat} \Rightarrow \text{ipv4addr}$  **where**  
*ipv4addr-of-dotdecimal* (*a,b,c,d*) = *ipv4addr-of-nat* ( $d + 256 * c + 65536 * b + 16777216 * a$ )

**fun** *dotdecimal-of-ipv4addr* :: *ipv4addr*  $\Rightarrow \text{nat} \times \text{nat} \times \text{nat} \times \text{nat}$  **where**  
*dotdecimal-of-ipv4addr* *a* = (*nat-of-ipv4addr* ( $((a \gg 24) \text{ AND } 0xFF)$ ),  
*nat-of-ipv4addr* ( $((a \gg 16) \text{ AND } 0xFF)$ ),  
*nat-of-ipv4addr* ( $((a \gg 8) \text{ AND } 0xFF)$ ),  
*nat-of-ipv4addr* ( $(a \text{ AND } 0xFF)$ ))

```

declare ipv4addr-of-dotdecimal.simps[simp del]
declare dotdecimal-of-ipv4addr.simps[simp del]

```

```

lemma ipv4addr-of-dotdecimal (192, 168, 0, 1) = 3232235521 by eval

```

```

lemma dotdecimal-of-ipv4addr 3232235521 = (192, 168, 0, 1) by eval

```

a different notation for *ipv4addr-of-dotdecimal*

```

lemma ipv4addr-of-dotdecimal-bit:

```

```

  ipv4addr-of-dotdecimal (a,b,c,d) = (ipv4addr-of-nat a << 24) + (ipv4addr-of-nat
b << 16) + (ipv4addr-of-nat c << 8) + ipv4addr-of-nat d

```

```

proof -

```

```

  have a: (ipv4addr-of-nat a) << 24 = ipv4addr-of-nat (a * 16777216)

```

```

    by(simp add: ipv4addr-of-nat-def shiftl-t2n of-nat-mult)

```

```

  have b: (ipv4addr-of-nat b) << 16 = ipv4addr-of-nat (b * 65536)

```

```

    by(simp add: ipv4addr-of-nat-def shiftl-t2n of-nat-mult)

```

```

  have c: (ipv4addr-of-nat c) << 8 = ipv4addr-of-nat (c * 256)

```

```

    by(simp add: ipv4addr-of-nat-def shiftl-t2n of-nat-mult)

```

```

  have ipv4addr-of-nat-suc:  $\bigwedge x. \text{ipv4addr-of-nat } (\text{Suc } x) = \text{word-succ } (\text{ipv4addr-of-nat } (x))$ 

```

```

    by(simp add: ipv4addr-of-nat-def, metis Abs-fnat-hom-Suc of-nat-Suc)

```

```

  { fix x y

```

```

    have ipv4addr-of-nat x + ipv4addr-of-nat y = ipv4addr-of-nat (x+y)

```

```

    apply(induction x arbitrary: y)

```

```

    apply(simp add: ipv4addr-of-nat-def)

```

```

    apply(simp add: ipv4addr-of-nat-suc word-succ-p1)

```

```

    done

```

```

  } from this a b c

```

```

show ?thesis

```

```

apply(simp add: ipv4addr-of-dotdecimal.simps)

```

```

apply(rule ipv4addr-of-nat-eq)

```

```

by presburger

```

```

qed

```

```

lemma size-ipv4addr: size (x::ipv4addr) = 32 by(simp add:word-size)

```

```

lemma ipv4addr-of-nat-shiftr-slice: ipv4addr-of-nat a >> x = slice x (ipv4addr-of-nat
a)

```

```

  by(simp add: ipv4addr-of-nat-def shiftr-slice)

```

```

value (4294967296::ipv4addr) = 2^32

```

```

lemma nat-of-ipv4addr-slice-ipv4addr-of-nat:

```

```

  nat-of-ipv4addr (slice x (ipv4addr-of-nat a)) = (nat-of-ipv4addr (ipv4addr-of-nat
a)) div 2^x

```

```

proof -

```

```

  have mod4294967296: int a mod 4294967296 = int (a mod 4294967296)

```

```

    using zmod-int by auto

```

```

  have int-pullin: int (a mod 4294967296) div 2 ^ x = int (a mod 4294967296
div 2 ^ x)

```

```

    using zpower-int zdiv-int by (metis of-nat-numeral )
show ?thesis
  apply(simp add: shiftr-slice[symmetric])
  apply(simp add: ipv4addr-of-nat-def word-of-nat)
  apply(simp add: nat-of-ipv4addr-def unat-def)
  apply(simp add: shiftr-div-2n)
  apply(simp add: uint-word-of-int)
  apply(simp add: mod4294967296 int-pullin)
done
qed
lemma ipv4addr-and-255: (x::ipv4addr) AND 255 = x AND mask 8
  apply(subst pow2-mask[of 8, simplified, symmetric])
  by simp
lemma ipv4addr-of-nat-AND-mask8: (ipv4addr-of-nat a) AND mask 8 = (ipv4addr-of-nat
(a mod 256))
  apply(simp add: ipv4addr-of-nat-def and-mask-mod-2p)
  apply(simp add: word-of-nat)
  apply(simp add: uint-word-of-int)
  apply(subst mod-mod-cancel)
  apply simp
  apply(simp add: zmod-int)
done

lemma dotdecimal-of-ipv4addr-ipv4addr-of-dotdecimal:
   $\llbracket a < 256; b < 256; c < 256; d < 256 \rrbracket \implies \text{dotdecimal-of-ipv4addr } (\text{ipv4addr-of-dotdecimal } (a,b,c,d)) = (a,b,c,d)$ 
proof -
  assume a < 256 and b < 256 and c < 256 and d < 256
  note assms =  $\langle a < 256 \rangle \langle b < 256 \rangle \langle c < 256 \rangle \langle d < 256 \rangle$ 
  hence a: nat-of-ipv4addr ((ipv4addr-of-nat (d + 256 * c + 65536 * b +
16777216 * a) >> 24) AND mask 8) = a
    apply(simp add: ipv4addr-of-nat-def word-of-nat)
    apply(simp add: nat-of-ipv4addr-def unat-def)
    apply(simp add: and-mask-mod-2p)
    apply(simp add: shiftr-div-2n)
    apply(simp add: uint-word-of-int)
    apply(subst mod-pos-pos-trivial)
    apply simp-all
    apply(subst mod-pos-pos-trivial)
    apply simp-all
    apply(subst mod-pos-pos-trivial)
    apply simp-all
  done
  from assms have b: nat-of-ipv4addr ((ipv4addr-of-nat (d + 256 * c + 65536
* b + 16777216 * a) >> 16) AND mask 8) = b
    apply(simp add: ipv4addr-of-nat-def word-of-nat)
    apply(simp add: nat-of-ipv4addr-def unat-def)
    apply(simp add: and-mask-mod-2p)
    apply(simp add: shiftr-div-2n)

```

```

    apply(simp add: uint-word-of-int)
    apply(subst mod-pos-pos-trivial)
    apply simp-all
    apply(subst mod-pos-pos-trivial[where b=4294967296])
    apply simp-all
    apply(simp add: NumberWang.div65536)
    done
  from assms have c: nat-of-ipv4addr ((ipv4addr-of-nat (d + 256 * c + 65536
* b + 16777216 * a) >> 8) AND mask 8) = c
    apply(simp add: ipv4addr-of-nat-def word-of-nat)
    apply(simp add: nat-of-ipv4addr-def unat-def)
    apply(simp add: and-mask-mod-2p)
    apply(simp add: shiftr-div-2n)
    apply(simp add: uint-word-of-int)
    apply(subst mod-pos-pos-trivial)
    apply simp-all
    apply(subst mod-pos-pos-trivial[where b=4294967296])
    apply simp-all
    apply(simp add: NumberWang.div256)
    done
  from (d < 256) have d: nat-of-ipv4addr (ipv4addr-of-nat (d + 256 * c +
65536 * b + 16777216 * a) AND mask 8) = d
    apply(simp add: ipv4addr-of-nat-AND-mask8)
    apply(simp add: ipv4addr-of-nat-def word-of-nat)
    apply(simp add: nat-of-ipv4addr-def)
    apply(subgoal-tac (d + 256 * c + 65536 * b + 16777216 * a) mod 256 = d)
    prefer 2
    apply(simp add: NumberWang.mod256)
    apply(simp)
    apply(simp add: unat-def)
    apply(simp add: uint-word-of-int)
    apply(simp add: mod-pos-pos-trivial)
    done
  from a b c d show ?thesis
    apply(simp add: ipv4addr-of-dotdecimal.simps dotdecimal-of-ipv4addr.simps)
    apply(simp add: ipv4addr-and-255)
    done
qed

```

**lemma** *ipv4addr-of-dotdecimal-eqE*:  $\llbracket \text{ipv4addr-of-dotdecimal } (a,b,c,d) = \text{ipv4addr-of-dotdecimal } (e,f,g,h); a < 256; b < 256; c < 256; d < 256; e < 256; f < 256; g < 256; h < 256 \rrbracket \implies$

$$a = e \wedge b = f \wedge c = g \wedge d = h$$

**by** (*metis* *Pair-inject dotdecimal-of-ipv4addr-ipv4addr-of-dotdecimal*)

previous and next ip addresses, without wrap around

**definition** *ip-next* :: *ipv4addr*  $\Rightarrow$  *ipv4addr* **where**  
*ip-next* a  $\equiv$  if a = max-ipv4-addr then max-ipv4-addr else a + 1

**definition** *ip-prev* :: *ipv4addr*  $\Rightarrow$  *ipv4addr* **where**  
*ip-prev* *a*  $\equiv$  if *a* = 0 then 0 else *a* - 1

**lemma** *ip-next* 2 = 3 **by** *eval*  
**lemma** *ip-prev* 2 = 1 **by** *eval*  
**lemma** *ip-prev* 0 = 0 **by** *eval*

## 1.2 IP ranges

**lemma** *UNIV-ipv4addrset*: (*UNIV* :: *ipv4addr* set) = {0 .. *max-ipv4-addr*}  
**by** (*auto*)  
**lemma** (42::*ipv4addr*)  $\in$  *UNIV* **by** *eval*

**definition** *ipv4range-set-from-netmask*::*ipv4addr*  $\Rightarrow$  *ipv4addr*  $\Rightarrow$  *ipv4addr* set  
**where**  
*ipv4range-set-from-netmask* *addr netmask*  $\equiv$  let *network-prefix* = (*addr* AND  
*netmask*) in {*network-prefix* .. *network-prefix* OR (*NOT netmask*)}

**lemma** *ipv4range-set-from-netmask* (*ipv4addr-of-dotdecimal* (192,168,0,42)) (*ipv4addr-of-dotdecimal*  
(255,255,0,0)) =  
{*ipv4addr-of-dotdecimal* (192,168,0,0) .. *ipv4addr-of-dotdecimal* (192,168,255,255)}  
**by** (*simp add: ipv4range-set-from-netmask-def ipv4addr-of-dotdecimal.simps ipv4addr-of-nat-def*)

**lemma** *ipv4range-set-from-netmask* (*ipv4addr-of-dotdecimal* (192,168,0,42)) (*ipv4addr-of-dotdecimal*  
(0,0,0,0)) = *UNIV*  
**by** (*simp add: UNIV-ipv4addrset ipv4addr-of-dotdecimal.simps ipv4addr-of-nat-def*  
*ipv4range-set-from-netmask-def max-ipv4-addr-max-word*)

192.168.0.0/24

**definition** *ipv4range-set-from-bitmask*::*ipv4addr*  $\Rightarrow$  *nat*  $\Rightarrow$  *ipv4addr* set **where**  
*ipv4range-set-from-bitmask* *addr bitmask*  $\equiv$  *ipv4range-set-from-netmask* *addr*  
(*of-bl* ((*replicate* *bitmask* *True*) @ (*replicate* (32 - *bitmask*) *False*)))

**lemma** (*replicate* 3 *True*) = [*True*, *True*, *True*] **by** *eval*  
**lemma** *of-bl* (*replicate* 3 *True*) = (7::*ipv4addr*) **by** *eval*

**lemma** *ipv4range-set-from-bitmask-alt1*:  
*ipv4range-set-from-bitmask* *addr bitmask* = *ipv4range-set-from-netmask* *addr*  
((*mask* *bitmask*) << (32 - *bitmask*))  
**apply** (*simp add: ipv4range-set-from-bitmask-def mask-bl*)  
**apply** (*simp add: Word.shiftl-of-bl*)  
**done**

**lemma** *ipv4range-set-from-bitmask* (*ipv4addr-of-dotdecimal* (192,168,0,42)) 16  
=

```

    {ipv4addr-of-dotdecimal (192,168,0,0) .. ipv4addr-of-dotdecimal (192,168,255,255)}
  by (simp add: ipv4range-set-from-bitmask-def ipv4range-set-from-netmask-def ipv4addr-of-dotdecimal.simps
    ipv4addr-of-nat-def)
  lemma ipv4range-set-from-bitmask-UNIV: ipv4range-set-from-bitmask 0 0 =
    UNIV
  apply (simp add: ipv4range-set-from-bitmask-def ipv4range-set-from-netmask-def)
)
  by (simp add: UNIV-ipv4addrset max-ipv4-addr-max-word)
  lemma ip-in-ipv4range-set-from-bitmask-UNIV: ip ∈ (ipv4range-set-from-bitmask
    (ipv4addr-of-dotdecimal (0, 0, 0, 0)) 0)
  by (simp add: ipv4addr-of-dotdecimal.simps ipv4addr-of-nat-def ipv4range-set-from-bitmask-UNIV)

  lemma ipv4range-set-from-bitmask-0: ipv4range-set-from-bitmask foo 0 = UNIV
  apply (rule)
  apply (simp-all)
  apply (simp add: ipv4range-set-from-bitmask-alt1 ipv4range-set-from-netmask-def
    Let-def)
  apply (simp add: range-0-max-UNIV)
  apply (simp add: mask-def)
  done

  lemma ipv4range-set-from-bitmask-32: ipv4range-set-from-bitmask foo 32 = {foo}
  apply (simp add: ipv4range-set-from-bitmask-alt1 ipv4range-set-from-netmask-def
    Let-def)
  apply (simp add: mask-def)
  apply (simp only: max-ipv4-addr-number[symmetric] max-ipv4-addr-max-word
    Word.word-and-max)
  apply (simp add: word32-or-NOT4294967296)
  done

  lemma ipv4range-set-from-bitmask-alt: ipv4range-set-from-bitmask pre len = {(pre
    AND ((mask len) << (32 - len))) .. pre OR (mask (32 - len))}
  apply (simp only: ipv4range-set-from-bitmask-alt1 ipv4range-set-from-netmask-def
    Let-def)
  apply (subst Word.word-aa-dist)
  apply (simp only: word-or-not)
  apply (simp only: Word.word-and-max)
  apply (simp only: NOT-mask-len32)
  done

  making element check executable

  lemma addr-in-ipv4range-set-from-netmask-code[code-unfold]:
    addr ∈ (ipv4range-set-from-netmask base netmask) ↔ (base AND netmask)
    ≤ addr ∧ addr ≤ (base AND netmask) OR (NOT netmask)
  by (simp add: ipv4range-set-from-netmask-def Let-def)
  lemma addr-in-ipv4range-set-from-bitmask-code[code-unfold]: addr ∈ (ipv4range-set-from-bitmask
    pre len) ↔
    (pre AND ((mask len) << (32 - len))) ≤ addr ∧ addr ≤ pre OR
    (mask (32 - len))

```

**unfolding** *ipv4range-set-from-bitmask-alt* **by** *simp*

**value** *ipv4addr-of-dotdecimal* (192,168,4,8) ∈ (*ipv4range-set-from-bitmask* (*ipv4addr-of-dotdecimal* (192,168,0,42)) 16)

**definition** *ipv4range-single* :: *ipv4addr* ⇒ 32 *wordinterval* **where**  
*ipv4range-single ip* ≡ *WordInterval ip ip*

**fun** *ipv4range-range* :: (*ipv4addr* × *ipv4addr*) ⇒ 32 *wordinterval* **where**  
*ipv4range-range (ip-start, ip-end)* = *WordInterval ip-start ip-end*  
**declare** *ipv4range-range.simps*[*simp del*]

**definition** *ipv4range-to-set* :: 32 *wordinterval* ⇒ (*ipv4addr*) *set* **where**  
*ipv4range-to-set rg* = *wordinterval-to-set rg*

**definition** *ipv4range-element* :: 'a::len *word* ⇒ 'a::len *wordinterval* ⇒ *bool* **where**  
*ipv4range-element el rg* = *wordinterval-element el rg*

**definition** *ipv4range-union* :: 32 *wordinterval* ⇒ 32 *wordinterval* ⇒ 32 *wordinterval* **where**  
*ipv4range-union r1 r2* = *wordinterval-union r1 r2*

**definition** *ipv4range-empty* :: 32 *wordinterval* ⇒ *bool* **where**  
*ipv4range-empty rg* = *wordinterval-empty rg*

**definition** *ipv4range-setminus* :: 32 *wordinterval* ⇒ 32 *wordinterval* ⇒ 32 *wordinterval* **where**  
*ipv4range-setminus r1 r2* = *wordinterval-setminus r1 r2*

**definition** *ipv4range-UNIV* :: 32 *wordinterval* **where** *ipv4range-UNIV* ≡ *wordinterval-UNIV*

**definition** *ipv4range-invert* :: 32 *wordinterval* ⇒ 32 *wordinterval* **where**  
*ipv4range-invert r* = *ipv4range-setminus ipv4range-UNIV r*

**definition** *ipv4range-intersection* :: 32 *wordinterval* ⇒ 32 *wordinterval* ⇒ 32 *wordinterval* **where**  
*ipv4range-intersection r1 r2* = *wordinterval-intersection r1 r2*

**definition** *ipv4range-subset* :: 32 *wordinterval* ⇒ 32 *wordinterval* ⇒ *bool* **where**  
*ipv4range-subset r1 r2* ≡ *wordinterval-subset r1 r2*

**definition** *ipv4range-eq* :: 32 *wordinterval* ⇒ 32 *wordinterval* ⇒ *bool* **where**  
*ipv4range-eq r1 r2* = *wordinterval-eq r1 r2*

```

lemma ipv4range-single-set-eq: ipv4range-to-set (ipv4range-single ip) = {ip}
  by(simp add: ipv4range-single-def ipv4range-to-set-def)
lemma ipv4range-range-set-eq: ipv4range-to-set (ipv4range-range (ip1, ip2)) =
{ip1 .. ip2}
  by(simp add: ipv4range-range.simps ipv4range-to-set-def)

lemma ipv4range-element-set-eq[simp]: ipv4range-element el rg = (el ∈ ipv4range-to-set
rg)
  by(simp add: ipv4range-element-def ipv4range-to-set-def)
lemma ipv4range-union-set-eq[simp]: ipv4range-to-set (ipv4range-union r1 r2)
= ipv4range-to-set r1 ∪ ipv4range-to-set r2
  by(simp add: ipv4range-to-set-def ipv4range-union-def)
lemma ipv4range-empty-set-eq[simp]: ipv4range-empty r  $\longleftrightarrow$  ipv4range-to-set r
= {}
  by(simp add: ipv4range-to-set-def ipv4range-empty-def)
lemma ipv4range-setminus-set-eq[simp]: ipv4range-to-set (ipv4range-setminus r1
r2) = ipv4range-to-set r1 - ipv4range-to-set r2
  by(simp add: ipv4range-setminus-def ipv4range-to-set-def)
lemma ipv4range-UNIV-set-eq[simp]: ipv4range-to-set ipv4range-UNIV = UNIV
  by(simp only: ipv4range-UNIV-def ipv4range-to-set-def wordinterval-UNIV-set-eq)
lemma ipv4range-invert-set-eq[simp]: ipv4range-to-set (ipv4range-invert r) =
UNIV - ipv4range-to-set r
  by(simp add: ipv4range-invert-def)
lemma ipv4range-intersection-set-eq[simp]: ipv4range-to-set (ipv4range-intersection
r1 r2) = ipv4range-to-set r1 ∩ ipv4range-to-set r2
  by(simp add: ipv4range-intersection-def ipv4range-to-set-def)
lemma ipv4range-subset-set-eq[simp]: ipv4range-subset r1 r2 = (ipv4range-to-set
r1 ⊆ ipv4range-to-set r2)
  by(simp add: ipv4range-subset-def ipv4range-to-set-def)
lemma ipv4range-eq-set-eq: ipv4range-eq r1 r2  $\longleftrightarrow$  ipv4range-to-set r1 = ipv4range-to-set
r2
  unfolding ipv4range-eq-def ipv4range-to-set-def using wordinterval-eq-set-eq
by blast

declare ipv4range-range-set-eq[unfolded ipv4range-range.simps, simp]
declare ipv4range-union-set-eq[unfolded ipv4range-union-def wordinterval-union.simps,
simp]

thm iffD1[OF ipv4range-eq-set-eq]
declare iffD1[OF ipv4range-eq-set-eq, cong]
lemma ipv4range-eq-comm: ipv4range-eq r1 r2  $\longleftrightarrow$  ipv4range-eq r2 r1
  unfolding ipv4range-eq-def wordinterval-eq-set-eq by blast
lemma ipv4range-to-set-alt: ipv4range-to-set r = {x. ipv4range-element x r}
  unfolding ipv4range-element-set-eq by blast

lemma ipv4range-un-empty: ipv4range-empty r1  $\implies$  ipv4range-eq (ipv4range-union
r1 r2) r2

```



by(subst ipv4range-eq-set-eq, simp)  
**lemma** *ipv4range-un-empt-b*: *ipv4range-empty* *r2*  $\implies$  *ipv4range-eq* (*ipv4range-union* *r1* *r2*) *r1*  
 by(subst ipv4range-eq-set-eq, simp)

**lemma** *ipv4range-Diff-triv*: *ipv4range-empty* (*ipv4range-intersection* *a* *b*)  $\implies$  *ipv4range-eq* (*ipv4range-setminus* *a* *b*) *a*  
 by(simp only: wordinterval-Diff-triv ipv4range-eq-def ipv4range-setminus-def ipv4range-intersection-def ipv4range-empty-def)

**definition** *is-lowest-element* *x S* = (*x*  $\in$  *S*  $\wedge$  ( $\forall y \in S. y \leq x \longrightarrow y = x$ ))

**fun** *ipv4range-lowest-element* :: 32 wordinterval  $\Rightarrow$  ipv4addr option **where**  
*ipv4range-lowest-element* (*WordInterval* *s e*) = (if *s*  $\leq$  *e* then *Some s* else *None*)  
 |  
*ipv4range-lowest-element* (*RangeUnion* *A B*) = (case (*ipv4range-lowest-element* *A*, *ipv4range-lowest-element* *B*) of  
 (*Some a*, *Some b*)  $\Rightarrow$  *Some* (if *a* < *b* then *a* else *b*) |  
 (*None*, *Some b*)  $\Rightarrow$  *Some b* |  
 (*Some a*, *None*)  $\Rightarrow$  *Some a* |  
 (*None*, *None*)  $\Rightarrow$  *None*)

**lemma** *ipv4range-lowest-none-empty*: *ipv4range-lowest-element* *r* = *None*  $\longleftrightarrow$  *ipv4range-empty* *r*  
**proof**(induction *r*)  
 case *WordInterval* **thus** ?case **by** simp  
 next  
 case *RangeUnion* **thus** ?case **by** fastforce  
 qed

**lemma** *ipv4range-lowest-element-correct-A*: *ipv4range-lowest-element* *r* = *Some* *x*  $\implies$  *is-lowest-element* *x* (*ipv4range-to-set* *r*)  
**unfolding** *is-lowest-element-def*  
**apply**(induction *r* arbitrary: *x* rule: *ipv4range-lowest-element.induct*)  
**apply**(rename-tac *rs re x*, case-tac *rs*  $\leq$  *re*, auto)[1]  
**apply**(subst(asm) *ipv4range-lowest-element.simps*(2))  
**apply**(rename-tac *A B x*)  
**apply**(case-tac *ipv4range-lowest-element B*)  
**apply**(case-tac[!] *ipv4range-lowest-element A*)  
**apply**(simp-all add: *ipv4range-lowest-none-empty*)[3]  
**apply** fastforce  
**done**

**lemma** *ipv4range-lowest-element-set-eq*: **assumes**  $\neg$  *ipv4range-empty* *r*  
**shows** (*ipv4range-lowest-element* *r* = *Some x*) = (*is-lowest-element* *x* (*ipv4range-to-set* *r*))

```

proof(rule iffI)
  assume ipv4range-lowest-element r = Some x
  thus is-lowest-element x (ipv4range-to-set r)
  using ipv4range-lowest-element-correct-A ipv4range-lowest-none-empty by
simp
next
  assume is-lowest-element x (ipv4range-to-set r)
  with assms show (ipv4range-lowest-element r = Some x)
  proof(induction r arbitrary: x rule: ipv4range-lowest-element.induct)
  case 1 thus ?case by(simp add: is-lowest-element-def)
  next
  case (2 A B x)

    have is-lowest-RangeUnion: is-lowest-element x (ipv4range-to-set A  $\cup$ 
ipv4range-to-set B)  $\implies$ 
      is-lowest-element x (ipv4range-to-set A)  $\vee$  is-lowest-element x (ipv4range-to-set
B)
    by(simp add: is-lowest-element-def)

    have ipv4range-lowest-element-RangeUnion:  $\bigwedge a b A B.$  ipv4range-lowest-element
A = Some a  $\implies$  ipv4range-lowest-element B = Some b  $\implies$ 
      ipv4range-lowest-element (RangeUnion A B) = Some (min a b)
    by(auto dest!: ipv4range-lowest-element-correct-A simp add: is-lowest-element-def
min-def)

    from 2 show ?case
    apply(case-tac ipv4range-lowest-element B)
    apply(case-tac[!] ipv4range-lowest-element A)
    apply(auto simp add: is-lowest-element-def)[3]
    apply(subgoal-tac  $\neg$  ipv4range-empty A  $\wedge$   $\neg$  ipv4range-empty B)
    prefer 2
    using arg-cong[where f = Not, OF ipv4range-lowest-none-empty] ap-
ply(simp, metis)
    apply(drule(1) ipv4range-lowest-element-RangeUnion)
    apply(simp split: option.split-asm add: min-def)
    apply(drule is-lowest-RangeUnion)
    apply(elim disjE)
    apply(simp add: is-lowest-element-def)
    apply(clarsimp simp add: ipv4range-lowest-none-empty)

    apply(simp add: is-lowest-element-def)
    apply(clarsimp simp add: ipv4range-lowest-none-empty)
    using ipv4range-lowest-element-correct-A[simplified is-lowest-element-def]
    by (metis Un-iff not-le)
  qed
qed

```

```

end
theory NumberWangCaesar
imports ./IPv4Addr
        ./l4v/lib/WordLemmaBucket
begin

context
begin

type-synonym prefix-match = (ipv4addr × nat)
definition pfxm-prefix p ≡ fst p
definition pfxm-length p ≡ snd p
definition pfxm-mask x ≡ mask (32 - pfxm-length x)
lemmas pfxm-defs = pfxm-prefix-def pfxm-mask-def pfxm-length-def

definition valid-prefix where
  valid-prefix pf = ((pfxm-mask pf) AND pfxm-prefix pf = 0)
private lemma valid-prefix-E: valid-prefix pf ⇒ ((pfxm-mask pf) AND pfxm-prefix
pf = 0)
  unfolding valid-prefix-def .
private lemma valid-prefix-alt-def: valid-prefix p = (pfxm-prefix p AND (2 ^ (32
- pfxm-length p) - 1) = 0)
  unfolding valid-prefix-def
  unfolding mask-def
  using word-bw-comms(1)
  arg-cong[where f = λx. (pfxm-prefix p AND x - 1 = 0)]
  shiftl-1
  unfolding pfxm-prefix-def pfxm-mask-def mask-def
  by metis

```

### 1.3 Address Semantics

```

definition prefix-match-semantics where
  prefix-match-semantics m a = (pfxm-prefix m = (NOT pfxm-mask m) AND a)

private lemma mask-32-max-word: mask 32 = (max-word :: 32 word) using
WordLemmaBucket.mask-32-max-word by simp

```

### 1.4 Set Semantics

```

definition prefix-to-ipset :: prefix-match ⇒ ipv4addr set where
  prefix-to-ipset pfx = {pfxm-prefix pfx .. pfxm-prefix pfx OR pfxm-mask pfx}

private lemma pfx-not-empty: valid-prefix pfx ⇒ prefix-to-ipset pfx ≠ {}
  unfolding valid-prefix-def prefix-to-ipset-def by (simp add: le-word-or2)

definition ipset-prefix-match where

```

$ipset\text{-}prefix\text{-}match\ pfx\ rg = (let\ pfxrg = prefix\text{-}to\text{-}ipset\ pfx\ in\ (rg \cap pfxrg, rg - pfxrg))$   
**private lemma** *ipset-prefix-match-m[simp]*:  $fst\ (ipset\text{-}prefix\text{-}match\ pfx\ rg) = rg \cap (prefix\text{-}to\text{-}ipset\ pfx)$  **by** (*simp only: Let-def ipset-prefix-match-def, simp*)  
**private lemma** *ipset-prefix-match-nm[simp]*:  $snd\ (ipset\text{-}prefix\text{-}match\ pfx\ rg) = rg - (prefix\text{-}to\text{-}ipset\ pfx)$  **by** (*simp only: Let-def ipset-prefix-match-def, simp*)  
**private lemma** *ipset-prefix-match-distinct*:  $rpm = ipset\text{-}prefix\text{-}match\ pfx\ rg \implies (fst\ rpm) \cap (snd\ rpm) = \{\}$  **by** *force*  
**private lemma** *ipset-prefix-match-complete*:  $rpm = ipset\text{-}prefix\text{-}match\ pfx\ rg \implies (fst\ rpm) \cup (snd\ rpm) = rg$  **by** *force*  
**private lemma** *rpm-m-dup-simp*:  $rg \cap fst\ (ipset\text{-}prefix\text{-}match\ (routing\text{-}match\ r)\ rg) = fst\ (ipset\text{-}prefix\text{-}match\ (routing\text{-}match\ r)\ rg)$  **by** *simp*

## 1.5 Equivalence Proofs

**private lemma** *helper3*:  $(x::32\ word)\ OR\ y = x\ OR\ y\ AND\ NOT\ x$  **by** (*simp add: word-oa-dist2*)

**private lemma** *packet-ipset-prefix-eq1*:  
**assumes**  $addr \in addrrg$   
**assumes** *valid-prefix match*  
**assumes**  $\neg prefix\text{-}match\text{-}semantics\ match\ addr$   
**shows**  $addr \in (snd\ (ipset\text{-}prefix\text{-}match\ match\ addrrg))$   
**using** *assms*  
**proof** –  
**have**  $pfxm\text{-}prefix\ match \leq addr \implies \neg addr \leq pfxm\text{-}prefix\ match\ OR\ pfxm\text{-}mask\ match$   
**proof** –  
**case** *goal1*  
**have**  $a1: pfxm\text{-}mask\ match\ AND\ pfxm\text{-}prefix\ match = 0$   
**using** *assms(2) unfolding valid-prefix-def .*  
**have**  $a2: pfxm\text{-}prefix\ match \neq NOT\ pfxm\text{-}mask\ match\ AND\ addr$   
**using** *assms(3) unfolding prefix-match-semantics-def .*  
**have**  $f1: pfxm\text{-}prefix\ match = pfxm\text{-}prefix\ match\ AND\ NOT\ pfxm\text{-}mask\ match$   
**using**  $a1$  **by** (*metis mask-eq-0-eq-x word-bw-comms(1)*)  
**hence**  $f2: \forall x_{11}. (pfxm\text{-}prefix\ match\ OR\ x_{11})\ AND\ NOT\ pfxm\text{-}mask\ match = pfxm\text{-}prefix\ match\ OR\ x_{11}\ AND\ NOT\ pfxm\text{-}mask\ match$   
**by** (*metis word-bool-alg.conj-disj-distrib2*)  
**moreover**  
**{ assume**  $\neg pfxm\text{-}prefix\ match \leq addr\ AND\ NOT\ pfxm\text{-}mask\ match$   
**hence**  $\neg (pfxm\text{-}prefix\ match \leq addr \wedge addr \leq pfxm\text{-}prefix\ match\ OR\ pfxm\text{-}mask\ match)$   
**using**  $f1$  *neg-mask-mono-le* **unfolding** *pfxm-prefix-def pfxm-mask-def* **by** *metis* **}**  
**moreover**  
**{ assume**  $pfxm\text{-}prefix\ match \leq addr\ AND\ NOT\ pfxm\text{-}mask\ match \wedge addr\ AND$

```

NOT pfxm-mask match  $\neq$  (pfxm-prefix match OR pfxm-mask match) AND NOT
pfxm-mask match
  hence  $\exists x_0. \neg \text{addr AND NOT mask } x_0 \leq (\text{pfxm-prefix match OR pfxm-mask}$ 
match) AND NOT mask  $x_0$ 
  using f2 unfolding pfxm-prefix-def pfxm-mask-def by (metis dual-order.antisym
word-bool-alg.conj-cancel-right word-log-esimps(3))
  hence  $\neg (\text{pfxm-prefix match} \leq \text{addr} \wedge \text{addr} \leq \text{pfxm-prefix match OR pfxm-mask}$ 
match)
  using neg-mask-mono-le by auto }
ultimately show ?case
  using a2 by (metis goal1 word-bool-alg.conj-cancel-right word-bool-alg.conj-commute
word-log-esimps(3))
qed
from this show ?thesis using assms(1)
  unfolding ipset-prefix-match-def Let-def snd-conv prefix-to-ipset-def
  by simp
qed

private lemma packet-ipset-prefix-eq2:
  assumes addr  $\in$  addrrg
  assumes valid-prefix match
  assumes prefix-match-semantics match addr
  shows addr  $\in$  (fst (ipset-prefix-match match addrrg))
using assms
  apply(subst ipset-prefix-match-def)
  apply(simp only: Let-def fst-def)
  apply(simp add: prefix-to-ipset-def)
  apply(transfer)
  apply(simp only: prefix-match-semantics-def valid-prefix-def)
  apply(simp add: word-and-le1)
  apply(metis helper3 le-word-or2 word-bw-comms(1) word-bw-comms(2))
done

private lemma packet-ipset-prefix-eq3:
  assumes addr  $\in$  addrrg
  assumes valid-prefix match
  assumes addr  $\in$  (snd (ipset-prefix-match match addrrg))
  shows  $\neg$ prefix-match-semantics match addr
using assms
  apply(subst(asm) ipset-prefix-match-def)
  apply(simp only: Let-def fst-def)
  apply(simp)
  apply(subst(asm) prefix-to-ipset-def)
  apply(transfer)
  apply(simp only: prefix-match-semantics-def valid-prefix-def Set-Interval.ord-class.atLeastAtMost-iff
prefix-to-ipset-def)
  apply(simp)
  apply(metis helper3 le-word-or2 word-and-le2 word-bw-comms(1) word-bw-comms(2))
done

```

```

private lemma packet-ipset-prefix-eq4:
  assumes addr ∈ addrrg
  assumes valid-prefix match
  assumes addr ∈ (fst (ipset-prefix-match match addrrg))
  shows prefix-match-semantics match addr
using assms
proof -
  have pfxm-prefix match = NOT pfxm-mask match AND addr
  proof -
    have a1: pfxm-mask match AND pfxm-prefix match = 0 using assms(2)
  unfolding valid-prefix-def .
    have a2: pfxm-prefix match ≤ addr ∧ addr ≤ pfxm-prefix match OR pfxm-mask
    match
    using assms(3) unfolding ipset-prefix-match-def Let-def fst-conv prefix-to-ipset-def
  by simp
    have f2: ∀ x₀. pfxm-prefix match AND NOT mask x₀ ≤ addr AND NOT mask
    x₀
    using a2 neg-mask-mono-le by blast
    have f3: ∀ x₀. addr AND NOT mask x₀ ≤ (pfxm-prefix match OR pfxm-mask
    match) AND NOT mask x₀
    using a2 neg-mask-mono-le by blast
    have f4: pfxm-prefix match = pfxm-prefix match AND NOT pfxm-mask match
    using a1 by (metis mask-eq-0-eq-x word-bw-comms(1))
    hence f5: ∀ x₆. (pfxm-prefix match OR x₆) AND NOT pfxm-mask match =
    pfxm-prefix match OR x₆ AND NOT pfxm-mask match
    using word-ao-dist by (metis)
    have f6: ∀ x₂ x₃. addr AND NOT mask x₂ ≤ x₃ ∨ ¬ (pfxm-prefix match OR
    pfxm-mask match) AND NOT mask x₂ ≤ x₃
    using f3 dual-order.trans by auto
    have pfxm-prefix match = (pfxm-prefix match OR pfxm-mask match) AND
    NOT pfxm-mask match
    using f5 by auto
    hence pfxm-prefix match = addr AND NOT pfxm-mask match
    using f2 f4 f6 unfolding pfxm-prefix-def pfxm-mask-def by (metis eq-iff)
    thus pfxm-prefix match = NOT pfxm-mask match AND addr
    by (metis word-bw-comms(1))
  qed
  from this show ?thesis unfolding prefix-match-semantics-def .
qed

```

```

private lemma packet-ipset-prefix-eq24:
  assumes addr ∈ addrrg
  assumes valid-prefix match
  shows prefix-match-semantics match addr = (addr ∈ (fst (ipset-prefix-match
    match addrrg)))
using packet-ipset-prefix-eq2[OF assms] packet-ipset-prefix-eq4[OF assms] by fast

```

```

private lemma packet-ipset-prefix-eq13:

```

```

assumes  $addr \in addrrg$ 
assumes  $valid\_prefix\ match$ 
shows  $\neg prefix\_match\_semantics\ match\ addr = (addr \in (snd\ (ipset\_prefix\_match\ match\ addrrg)))$ 
using  $packet\_ipset\_prefix\_eq1\ [OF\ assms]\ packet\_ipset\_prefix\_eq3\ [OF\ assms]$  by  $fast$ 

private lemma  $prefix\_match\_if\_in\_my\_set$ : assumes  $valid\_prefix\ pfx$ 
shows  $prefix\_match\_semantics\ pfx\ (a :: ipv4addr) \longleftrightarrow a \in prefix\_to\_ipset\ pfx$ 
using  $packet\_ipset\_prefix\_eq24\ [OF\ -\ assms]$ 
by  $(metis\ (erased,\ hide\_lams)\ Int\_iff\ UNIV\_I\ fst\_conv\ ipset\_prefix\_match\_def)$ 

lemma  $prefix\_match\_if\_in\_corny\_set$ :
assumes  $valid\_prefix\ pfx$ 
shows  $prefix\_match\_semantics\ pfx\ (a :: ipv4addr) \longleftrightarrow a \in ipv4range\_set\_from\_netmask\ (pfxm\_prefix\ pfx)\ (NOT\ pfxm\_mask\ pfx)$ 
unfolding  $prefix\_match\_if\_in\_my\_set\ [OF\ assms]$ 
unfolding  $prefix\_to\_ipset\_def\ ipv4range\_set\_from\_netmask\_def\ Let\_def$ 
unfolding  $word\_bool\_alg.double\_compl$ 
proof  $-$ 
  case  $goal1$ 
  have  $*$ :  $pfxm\_prefix\ pfx\ AND\ NOT\ pfxm\_mask\ pfx = pfxm\_prefix\ pfx$ 
  unfolding  $mask\_eq\_0\_eq\_x\ [symmetric]$  using  $valid\_prefix\_E\ [OF\ assms]\ word\_bw\_comms(1)\ [of\ pfxm\_prefix\ pfx]$  by  $simp$ 
  hence  $**$ :  $pfxm\_prefix\ pfx\ AND\ NOT\ pfxm\_mask\ pfx\ OR\ pfxm\_mask\ pfx = pfxm\_prefix\ pfx\ OR\ pfxm\_mask\ pfx$ 
  by  $simp$ 
  show  $?case$  unfolding  $*\ **\ ..$ 
qed

private lemma  $ipv4addr\_and\_maskshift\_eq\_and\_not\_mask$ :  $(base :: 32\ word)\ AND\ (mask\ m << 32 - m) = base\ AND\ NOT\ mask\ (32 - m)$ 
apply  $word\_bitwise$ 
apply  $(subgoal\_tac\ m > 32 \vee m \in set\ (map\ nat\ (upto\ 0\ 32)))$ 
apply  $(simp\ add: upto\_code\ upto\_aux\_rec,\ elim\ disjE)$ 
  apply  $(simp\ add: size\_mask\_32word)$ 
  apply  $(simp\_all\ add: size\_mask\_32word)\ [33]$ 
apply  $(simp\ add: upto\_code\ upto\_aux\_rec,\ presburger)$ 
done

lemma  $maskshift\_eq\_not\_mask$ :  $((mask\ m << 32 - m) :: 32\ word) = NOT\ mask\ (32 - m)$ 
apply  $word\_bitwise$ 
apply  $(subgoal\_tac\ m > 32 \vee m \in set\ (map\ nat\ (upto\ 0\ 32)))$ 
apply  $(simp\ add: upto\_code\ upto\_aux\_rec,\ elim\ disjE)$ 
  apply  $(simp\ add: size\_mask\_32word)$ 
  apply  $(simp\_all\ add: size\_mask\_32word)\ [33]$ 

```

**apply** (simp add: upto-code upto-aux-rec, presburger)  
**done**

**private lemma** *ipv4addr-andnotmask-eq-ormaskandnot*:  $((base::32\ word)\ AND\ NOT\ mask\ (32 - m)) = ((base\ OR\ mask\ (32 - m))\ AND\ NOT\ mask\ (32 - m))$   
**apply** word-bitwise  
**apply** (subgoal-tac  $m > 32 \vee m \in set\ (map\ nat\ (upto\ 0\ 32))$ )  
**apply** (simp add: upto-code upto-aux-rec, elim disjE)  
**apply** (simp add: size-mask-32word)  
**apply** (simp-all add: size-mask-32word) [33]  
**apply** (simp add: upto-code upto-aux-rec, presburger)  
**done**

**private lemma** *size-mask-32word'*:  $size\ ((mask\ (32 - m))::32\ word) = 32$  **by** (simp add: word-size)

**lemma** *wordinterval-to-set-ipv4range-set-from-bitmask*: **assumes** *valid-prefix pfx*  
**shows** *prefix-to-ipset pfx = ipv4range-set-from-bitmask (pfxm-prefix pfx)*  
*(pfxm-length pfx)*  
**proof** –  
**have** *prefix-match-if-in-corny-set*:  $(prefix-to-ipset\ pfx) = ipv4range-set-from-netmask\ (pfxm-prefix\ pfx)\ (NOT\ pfxm-mask\ pfx)$   
**unfolding** *prefix-to-ipset-def ipv4range-set-from-netmask-def Let-def*  
**unfolding** *word-bool-alg.double-compl*  
**proof** –  
**case** *goal1*  
**have** \*:  $pfxm-prefix\ pfx\ AND\ NOT\ pfxm-mask\ pfx = pfxm-prefix\ pfx$   
**unfolding** *mask-eq-0-eq-x[symmetric]* **using** *valid-prefix-E[OF assms]*  
*word-bw-comms(1)[of pfxm-prefix pfx]* **by** *simp*  
**hence** \*\*:  $pfxm-prefix\ pfx\ AND\ NOT\ pfxm-mask\ pfx\ OR\ pfxm-mask\ pfx = pfxm-prefix\ pfx\ OR\ pfxm-mask\ pfx$   
**by** *simp*  
**show** ?case **unfolding** \* \*\* ..  
**qed**

**have**  $\bigwedge len. ((mask\ len)::ipv4addr) << 32 - len = \sim\sim\ mask\ (32 - len)$   
**using** *maskshift-eq-not-mask* **by** *simp*  
**from** *this[of (pfxm-length pfx)]* **have** *mask-def2-symmetric*:  $((mask\ (pfxm-length\ pfx))::ipv4addr) << 32 - pfxm-length\ pfx = NOT\ pfxm-mask\ pfx$   
**unfolding** *pfxm-mask-def* **by** *simp*

**have** *ipv4range-set-from-netmask-bitmask*:  
 $ipv4range-set-from-netmask\ (pfxm-prefix\ pfx)\ (NOT\ pfxm-mask\ pfx) = ipv4range-set-from-bitmask\ (pfxm-prefix\ pfx)\ (pfxm-length\ pfx)$   
**unfolding** *ipv4range-set-from-netmask-def ipv4range-set-from-bitmask-alt*



```

unfolding pfxm-mask-def[symmetric]
unfolding mask-def2-symmetric
apply(simp)
unfolding Let-def
using assms[unfolded valid-prefix-def] by (metis helper3 word-bw-comms(2))

show ?thesis by (metis ipv4range-set-from-netmask-bitmask local.prefix-match-if-in-corny-set)

qed

private lemma helper-32-case-split:  $32 < m \vee m \in \text{set } (\text{map nat } [0..32])$ 
by (simp add: upto-code upto-aux-rec, presburger)
private lemma ipv4addr-andnot-impl-takem:  $(a::32 \text{ word}) \text{ AND NOT mask } (32 - m) = b \implies (\text{take } (m) (\text{to-bl } a)) = (\text{take } (m) (\text{to-bl } b))$ 
apply word-bitwise
apply (subgoal-tac  $m > 32 \vee m \in \text{set } (\text{map nat } (\text{upto } 0 \ 32)))$ 
prefer 2
apply(simp only: helper-32-case-split)
apply (simp add: upto-code upto-aux-rec, elim disjE)
apply (simp add: size-mask-32word size-mask-32word')
apply (simp-all add: size-mask-32word size-mask-32word')
done

definition ip-set ::  $32 \text{ word} \Rightarrow \text{nat} \Rightarrow 32 \text{ word set}$  where  $\text{ip-set } i \ r = \{j . i \text{ AND NOT mask } (32 - r) = j \text{ AND NOT mask } (32 - r)\}$ 

private lemma  $(m1 \vee m2) \wedge (m3 \vee m4) \longleftrightarrow (m1 \wedge m3) \vee (m1 \wedge m4) \vee (m2 \wedge m3) \vee (m2 \wedge m4)$ 
by blast

private lemmas caesar-proof-unfolded = prefix-match-if-in-corny-set[unfolded valid-prefix-def
prefix-match-semantics-def Let-def, symmetric]
private lemma caesar-proof-without-structures:  $\text{mask } (32 - l) \text{ AND pfxm-p} = 0 \implies$ 
 $(a \in \text{ipv4range-set-from-netmask } (\text{pfxm-p}) (\text{NOT mask } (32 - l))) =$ 
 $(\text{pfxm-p} = \text{NOT mask } (32 - l) \text{ AND } a)$ 
using caesar-proof-unfolded unfolding pfxm-defs by force

private lemma mask-and-not-mask-helper:  $\text{mask } (32 - m) \text{ AND base AND NOT mask } (32 - m) = 0$ 
by(simp add: word-bw-lcs)

lemma ipv4range-set-from-netmask-base-mask-consume:
 $\text{ipv4range-set-from-netmask } (\text{base AND NOT mask } (32 - m)) (\text{NOT mask } (32 - m)) =$ 
 $\text{ipv4range-set-from-netmask base } (\text{NOT mask } (32 - m))$ 
unfolding ipv4range-set-from-netmask-def

```

```

by(simp add: AND-twice)

lemma ipv4range-set-from-bitmask-eq-ip-set: ipv4range-set-from-bitmask base m =
ip-set base m
  unfolding ip-set-def
  unfolding set-eq-iff
  unfolding mem-Collect-eq
  unfolding ipv4range-set-from-bitmask-alt1
  unfolding maskshift-eq-not-mask
  using caesar-proof-without-structures[OF mask-and-not-mask-helper, of - base m]
  unfolding ipv4range-set-from-netmask-base-mask-consume
  unfolding word-bw-comms(1)[of -  $\sim\sim$  mask (32 - m)]
  ..

lemma cornys-hacky-call-to-prefix-to-range-to-start-with-a-valid-prefix: valid-prefix
(base AND NOT mask (32 - len), len)
  apply(simp add: valid-prefix-def pfxm-mask-def pfxm-length-def pfxm-prefix-def)
  by (metis mask-and-not-mask-helper)
end
end
theory NumberWangCebewee
imports
  ./l4v/lib/WordLemmaBucket
  NumberWangCaesar
begin

lemma and-not-mask-twice:
  (w &&  $\sim\sim$  mask n) &&  $\sim\sim$  mask m = w &&  $\sim\sim$  mask (max m n)
  apply (simp add: and-not-mask)
  apply (case-tac n<m)
  apply (simp-all add: shiftl-shiftr2 shiftl-shiftr1 not-less max-def
    shiftr-shiftr shiftl-shiftl)
  apply (cut-tac and-mask-shiftr-comm
    [where w=w and m=size w and n=m, simplified,symmetric])
  apply (simp add: word-size mask-def)
  apply (cut-tac and-mask-shiftr-comm
    [where w=w and m=size w and n=n, simplified,symmetric])
  apply (simp add: word-size mask-def)
  done

lemma X: j ∈ ip-set i r  $\implies$  ip-set j r = ip-set i r
  by (auto simp: ip-set-def)

lemma Z:

```

```

fixes  $i :: ('a :: len) \text{ word}$ 
assumes  $r2 \leq r1 \ i \ \&\& \sim\sim \text{ mask } r2 = x \ \&\& \sim\sim \text{ mask } r2$ 
shows  $i \ \&\& \sim\sim \text{ mask } r1 = x \ \&\& \sim\sim \text{ mask } r1$ 
proof -
  have  $i \text{ AND NOT mask } r1 = (i \ \&\& \sim\sim \text{ mask } r2) \ \&\& \sim\sim \text{ mask } r1$  (is - = ?w
&& -)
    using  $\langle r2 \leq r1 \rangle$  by (simp add: and-not-mask-twice max-def)
  also have  $?w = x \ \&\& \sim\sim \text{ mask } r2$  by fact
  also have  $\dots \ \&\& \sim\sim \text{ mask } r1 = x \ \&\& \sim\sim \text{ mask } r1$ 
    using  $\langle r2 \leq r1 \rangle$  by (simp add: and-not-mask-twice max-def)
  finally show ?thesis .
qed

```

```

lemma  $Y: r1 \leq r2 \implies \text{ip-set } i \ r2 \subseteq \text{ip-set } i \ r1$ 
  unfolding ip-set-def
  apply auto
  apply (rule Z[where ?r2.0=32 - r2])
  apply auto
  done

```

```

lemma ip-set-intersect-subset-helper:
  fixes  $i1 \ r1 \ i2 \ r2$ 
  assumes  $\text{disj}: \text{ip-set } i1 \ r1 \cap \text{ip-set } i2 \ r2 \neq \{\}$  and  $r1 \leq r2$ 
  shows  $\text{ip-set } i2 \ r2 \subseteq \text{ip-set } i1 \ r1$ 
  proof -
    from  $\text{disj}$  obtain  $j$  where  $j \in \text{ip-set } i1 \ r1 \ j \in \text{ip-set } i2 \ r2$  by auto
    with  $\langle r1 \leq r2 \rangle$  have  $j \in \text{ip-set } j \ r1 \ j \in \text{ip-set } j \ r1$  using  $X \ Y$  by blast+

    show  $\text{ip-set } i2 \ r2 \subseteq \text{ip-set } i1 \ r1$ 
    proof
      fix  $i$  assume  $i \in \text{ip-set } i2 \ r2$ 
      with  $\langle j \in \text{ip-set } i2 \ r2 \rangle$  have  $i \in \text{ip-set } j \ r2$  using  $X$  by auto
      also have  $\text{ip-set } j \ r2 \subseteq \text{ip-set } j \ r1$  using  $\langle r1 \leq r2 \rangle \ Y$  by blast
      also have  $\dots = \text{ip-set } i1 \ r1$  using  $\langle j \in \text{ip-set } i1 \ r1 \rangle \ X$  by blast
      finally show  $i \in \text{ip-set } i1 \ r1$  .
    qed
  qed

```

```

lemma ip-set-notsubset-empty-inter:  $\neg \text{ip-set } i1 \ r1 \subseteq \text{ip-set } i2 \ r2 \implies \neg \text{ip-set } i2$ 
 $r2 \subseteq \text{ip-set } i1 \ r1 \implies \text{ip-set } i1 \ r1 \cap \text{ip-set } i2 \ r2 = \{\}$ 
  apply(cases  $r1 \leq r2$ )
  using ip-set-intersect-subset-helper apply blast
  apply(cases  $r2 \leq r1$ )
  using ip-set-intersect-subset-helper apply blast
  apply(simp)
  done

```

```

end
theory Numberwang-Ln
imports NumberWangCebewee
begin

lemma ipv4range-bitmask-intersect:  $\neg \text{ipv4range-set-from-bitmask } b2 \ m2 \subseteq \text{ipv4range-set-from-bitmask } b1 \ m1 \implies$ 
 $\neg \text{ipv4range-set-from-bitmask } b1 \ m1 \subseteq \text{ipv4range-set-from-bitmask } b2 \ m2 \implies$ 
 $\text{ipv4range-set-from-bitmask } b1 \ m1 \cap \text{ipv4range-set-from-bitmask } b2 \ m2 = \{\}$ 
apply(simp add: ipv4range-set-from-bitmask-eq-ip-set)
using ip-set-notsubset-empty-inter by presburger

```

```

lemma ipv4addr-of-dotdecimal-dotdecimal-of-ipv4addr:
  (ipv4addr-of-dotdecimal (dotdecimal-of-ipv4addr ip)) = ip
proof -
  have ip-and-mask8-bl-drop24: (ip::ipv4addr) AND mask 8 = of-bl (drop 24 (to-bl ip))
  by(simp add: WordLemmaBucket.of-drop-to-bl size-ipv4addr)

  have List-rev-drop-geqn:  $\bigwedge x \ n. \text{length } x \geq n \implies (\text{take } n \ (\text{rev } x)) = \text{rev } (\text{drop } (\text{length } x - n) \ x)$ 
  by(simp add: List.rev-drop)

  have and-mask-bl-take:  $\bigwedge x \ n. \text{length } x \geq n \implies ((\text{of-bl } x) \text{ AND mask } n) = (\text{of-bl } (\text{rev } (\text{take } n \ (\text{rev } (x)))))$ 
  apply(simp add: List-rev-drop-geqn)
  apply(simp add: WordLib.of-bl-drop)
  done

  have bit-equality:  $((ip >> 24) \text{ AND } 0xFF << 24) + ((ip >> 16) \text{ AND } 0xFF << 16) + ((ip >> 8) \text{ AND } 0xFF << 8) + (ip \text{ AND } 0xFF) =$ 
 $\text{of-bl } (\text{take } 8 \ (\text{to-bl } ip)) @ \text{take } 8 \ (\text{drop } 8 \ (\text{to-bl } ip)) @ \text{take } 8 \ (\text{drop } 16 \ (\text{to-bl } ip)) @ \text{drop } 24 \ (\text{to-bl } ip)$ 
  apply(simp add: ipv4addr-and-255)
  apply(simp add: shiftr-slice)
  apply(simp add: Word.slice-take' size-ipv4addr)
  apply(simp add: and-mask-bl-take)
  apply(simp add: List-rev-drop-geqn)
  apply(simp add: drop-take)
  apply(simp add: Word.shiffl-of-bl)
  apply(simp add: of-bl-append)
  apply(simp add: ip-and-mask8-bl-drop24)
  done

  have blip-split:  $\bigwedge \text{blip}. \text{length } \text{blip} = 32 \implies \text{blip} = (\text{take } 8 \ \text{blip}) @ (\text{take } 8 \ (\text{drop } 8 \ \text{blip})) @ (\text{take } 8 \ (\text{drop } 16 \ \text{blip})) @ (\text{drop } 24 \ \text{blip})$ 

```

```

8 blip)) @ (take 8 (drop 16 blip)) @ (take 8 (drop 24 blip))
  apply(case-tac blip)
  apply(simp-all)

  apply(rename-tac blip,case-tac blip,simp-all)+
done

have ipv4addr-of-dotdecimal (dotdecimal-of-ipv4addr ip) = of-bl (to-bl ip)
  apply(subst blip-split)
  apply(simp)
  apply(simp add: ipv4addr-of-dotdecimal-bit dotdecimal-of-ipv4addr.simps)
  apply(simp add: ipv4addr-of-nat-nat-of-ipv4addr)
  apply(simp add: bit-equality)
done

```

```

thus ?thesis using Word.word-bl.Rep-inverse[symmetric] by simp
qed

```

```

end
theory CIDRSplit
imports IPv4Addr NumberWangCaesar
begin

context
begin

```

## 2 CIDR Split Motivation

When talking about ranges of IP addresses, we can make the ranges explicit by listing them.

```

value map (ipv4addr-of-nat ∘ nat) [1 .. 4]
definition ipv4addr-upto :: ipv4addr ⇒ ipv4addr ⇒ ipv4addr list where
  ipv4addr-upto i j ≡ map (ipv4addr-of-nat ∘ nat) [int (nat-of-ipv4addr i) .. int
(nat-of-ipv4addr j)]
lemma ipv4addr-upto: set (ipv4addr-upto i j) = {i .. j}
proof –
  have helpX: ⋀f (i::nat) (j::nat). (f ∘ nat) ‘ {int i..int j} = f ‘ {i .. j}
    apply(intro set-eqI)
    apply(safe)
    apply(force)
  by (metis Set-Interval.transfer-nat-int-set-functions(2) image-comp image-eqI)
  have ipv4addr-of-nat-def': ipv4addr-of-nat = of-nat using ipv4addr-of-nat-def
fun-eq-iff by presburger
show ?thesis
  unfolding ipv4addr-upto-def

```

```

apply(intro set-eqI)
apply(simp add: ipv4addr-of-nat-def ' nat-of-ipv4addr-def)
apply(safe)
  apply(simp-all)
  apply (metis (no-types, hide-lams) le-unat-uoI nat-mono uint-nat unat-def
word-le-nat-alt)
  apply (metis (no-types, hide-lams) le-unat-uoI nat-mono uint-nat unat-def
word-le-nat-alt)
  apply(simp add: helpX)
  by (metis atLeastAtMost-iff image-eqI word-le-nat-alt word-unat.Rep-inverse)
qed

```

The function *ipv4addr-upto* gives back a list of all the ips in the list. This list can be pretty huge! In the following, we will use CIDR notation (e.g. 192.168.0.0/24) to describe the list more compactly.

## 2.1 Prefix Match Range stuff

**definition** *prefix-to-range* :: *prefix-match*  $\Rightarrow$  32 *wordinterval* **where**  
*prefix-to-range pfx* = *WordInterval (pfxm-prefix pfx) (pfxm-prefix pfx OR pfxm-mask pfx)*

**lemma** *prefix-to-range-set-eq*: *wordinterval-to-set (prefix-to-range pfx) = prefix-to-ipset pfx*

**unfolding** *prefix-to-range-def prefix-to-ipset-def* **by** *simp*

**lemma** *prefix-to-range-ipv4range-range*: *prefix-to-range pfx = ipv4range-range ((pfxm-prefix pfx), (pfxm-prefix pfx OR pfxm-mask pfx))*

**unfolding** *ipv4range-range.simps prefix-to-range-def* **by** *simp*

**corollary** *valid-prefix pfx*  $\Rightarrow$  *wordinterval-to-set (prefix-to-range pfx) = ipv4range-set-from-bitmask (pfxm-prefix pfx) (pfxm-length pfx)*

**using** *wordinterval-to-set-ipv4range-set-from-bitmask prefix-to-range-set-eq* **by** *simp*

**lemma** *prefix-bitrang-list-union*:  $\forall pfx \in \text{set cidrlist}. (\text{valid-prefix } pfx) \Rightarrow$   
*wordinterval-to-set (list-to-wordinterval (map prefix-to-range cidrlist)) =*  
 $\bigcup ((\lambda(\text{base}, \text{len}). \text{ipv4range-set-from-bitmask base len}) \text{ ' set (cidrlist)})$

```

apply(induction cidrlist)
  apply(simp)
  apply(simp)
  apply(subst prefix-to-range-set-eq)
  apply(subst wordinterval-to-set-ipv4range-set-from-bitmask)
  apply(simp)
  apply(simp add: pfxm-prefix-def pfxm-length-def)
  apply(clarify)
  apply(simp)
done

```

**lemma** *prefix-to-ipset-subset-ipv4range-set-from-bitmask*:  
*prefix-to-ipset pfx*  $\subseteq$  *ipv4range-set-from-bitmask (pfxm-prefix pfx) (pfxm-length pfx)*  
**apply**(*rule*)  
**apply**(*simp add: prefix-to-ipset-def addr-in-ipv4range-set-from-bitmask-code*)  
**apply**(*intro impI conjI*)  
**apply** (*metis (erased, hide-lams) order-trans word-and-le2*)  
**by** (*metis pfxm-mask-def*)

**private definition** *pfxes* :: *nat list* **where** *pfxes*  $\equiv$  *map nat [0..32]*

**definition** *ipv4range-split1* *r*  $\equiv$  (  
*let ma = ipv4range-lowest-element r in*  
*case ma of None*  $\Rightarrow$  (*None, r*) |  
*Some a*  $\Rightarrow$  *let cs = (map ( $\lambda s. (a,s)$ ) pfxes) in*  
*let cfs = filter ( $\lambda s. \text{valid-prefix } s \wedge \text{ipv4range-subset (prefix-to-range s) r}$ ) cs in* (\* anything that is a valid prefix should also be a subset. but try proving that.\*)  
*let mc = find (const True) cfs in*  
*(case mc of None*  $\Rightarrow$  (*None, r*) |  
*Some m*  $\Rightarrow$  (*mc, ipv4range-setminus r (prefix-to-range m)*)))

**private lemma** *flipnot*: *a=b*  $\Rightarrow$  ( $\neg a$ )= $(\neg b)$  **by** *simp*

**private lemma** *find-const-True*: *find (const True) l = None*  $\longleftrightarrow$  *l = []*

**by**(*cases l, simp-all add: const-def*)

**private lemma** *ipv4range-split-innards-helper*: *ipv4range-lowest-element r = Some a*  $\Rightarrow$

[*s*  $\leftarrow$  *map (Pair a) pfxes . valid-prefix s*  $\wedge$  *ipv4range-to-set (prefix-to-range s)*  $\subseteq$  *ipv4range-to-set r*]  $\neq$  []

**proof** –

**assume** *a: ipv4range-lowest-element r = Some a*

**have** *b: (a,32)  $\in$  set (map (Pair a) pfxes)*

**unfolding** *pfxes-def*

**unfolding** *set-map set-upto*

**using** *Set.image-iff atLeastAtMost-iff int-eq-iff of-nat-numeral order-refl*

**by** (*metis (erased, hide-lams)*)

**have** *c: valid-prefix (a,32)* **unfolding** *valid-prefix-def pfxm-defs* **by** *simp*

**have** *ipv4range-to-set (prefix-to-range (a,32)) = {a}* **unfolding** *prefix-to-range-def pfxm-defs* **by** *simp*

**moreover** **have** *a  $\in$  ipv4range-to-set r* **using** *a ipv4range-lowest-element-set-eq ipv4range-lowest-none-empty*

**by** (*metis is-lowest-element-def option.distinct(1)*)

**ultimately** **have** *d: ipv4range-to-set (prefix-to-range (a,32))  $\subseteq$  ipv4range-to-set r* **by** *simp*

**show** *?thesis*

```

    unfolding flipnot[OF set-empty[symmetric]]
    unfolding set-filter
    using b c d by blast
qed
private lemma r-split1-not-none:  $\neg \text{ipv4range-empty } r \implies \text{fst } (\text{ipv4range-split1 } r) \neq \text{None}$ 
  unfolding ipv4range-split1-def Let-def
  apply(cases ipv4range-lowest-element r)
  apply(simp add: ipv4range-lowest-none-empty)
  apply(simp only:)
  apply(case-tac find (const True) [s ← map (Pair a) pfxes . valid-prefix s ∧
    ipv4range-subset (prefix-to-range s) r])
  apply(simp add: find-const-True ipv4range-split-innard-helper)
  apply(simp)
done
private lemma find-in:  $\text{Some } a = \text{find } f \ s \implies a \in \{x \in \text{set } s. f \ x\}$ 
  by (metis findSomeD mem-Collect-eq)
theorem ipv4range-split1-preserve:  $(\text{Some } s, u) = \text{ipv4range-split1 } r \implies \text{ipv4range-eq}$ 
 $(\text{ipv4range-union } (\text{prefix-to-range } s) \ u) \ r$ 
proof(unfold ipv4range-eq-set-eq)
  assume as:  $(\text{Some } s, u) = \text{ipv4range-split1 } r$ 
  have nn:  $\text{ipv4range-lowest-element } r \neq \text{None}$ 
    using as unfolding ipv4range-split1-def Let-def
    by (metis (erased, lifting) Pair-inject option.distinct(2) option.simps(4))
  then obtain a where a:  $\text{Some } a = (\text{ipv4range-lowest-element } r)$  unfolding
not-None-eq by force
  then have cpf:  $\text{find } (\text{const True}) [s \leftarrow \text{map } (\text{Pair } a) \text{ pfxes } . \text{valid-prefix } s \wedge$ 
 $\text{ipv4range-subset } (\text{prefix-to-range } s) \ r] \neq \text{None}$  (is ?cpf  $\neq \text{None}$ )
    unfolding flipnot[OF find-const-True]
    using ipv4range-split-innard-helper
    by simp
  then obtain m where m:  $m = \text{the } ?\text{cpf}$  by blast
  have s-def:  $\text{ipv4range-split1 } r =$ 
 $(\text{find } (\text{const True}) [s \leftarrow \text{map } (\text{Pair } a) \text{ pfxes } . \text{valid-prefix } s \wedge \text{ipv4range-subset}$ 
 $(\text{prefix-to-range } s) \ r], \text{ipv4range-setminus } r \ (\text{prefix-to-range } m))$ 
    unfolding m ipv4range-split1-def Let-def using cpf
    unfolding a[symmetric]
    unfolding option.simps(5)
    using option.collapse
    by force
  have u =  $\text{ipv4range-setminus } r \ (\text{prefix-to-range } s)$ 
    using as unfolding s-def using m by (metis (erased, lifting) Pair-inject
handy-lemma)
  moreover have  $\text{ipv4range-subset } (\text{prefix-to-range } s) \ r$ 
    using as unfolding s-def
    apply(rule Pair-inject)
    apply(unfold const-def)
    apply(drule find-in)
    apply(unfold set-filter)

```



```

    by blast
  ultimately show ipv4range-to-set (ipv4range-union (prefix-to-range s) u) =
    ipv4range-to-set r by auto
qed

```

```

private lemma ((a,b),(c,d)) = ((a,b),c,d) by simp

```

```

private lemma prefix-never-empty: ¬ipv4range-empty (prefix-to-range d)
proof -
  have ie: pfxm-prefix d ≤ pfxm-prefix d || pfxm-mask d by (metis le-word-or2)
  have ipv4range-element (fst d) (prefix-to-range d)
    unfolding ipv4range-element-set-eq
    unfolding ipv4range-to-set-def
    unfolding prefix-to-range-set-eq
    unfolding prefix-to-ipset-def
    using first-in-uptoD[OF ie]
    unfolding pfxm-defs
    .
  thus ?thesis
    unfolding ipv4range-empty-set-eq
    unfolding ipv4range-element-set-eq
    by blast
qed

```

```

lemma ipv4range-split1-never-empty: (Some s, u) = ipv4range-split1 r ⇒ ¬ipv4range-empty
(prefix-to-range s)
  unfolding ipv4range-split1-def Let-def
  using prefix-never-empty
  by simp

```

```

lemma ipv4range-split1-some-r-ne: (Some s, u) = ipv4range-split1 r ⇒ ¬ipv4range-empty
r
proof(rule ccontr)
  case goal1
  have ipv4range-lowest-element r = None unfolding ipv4range-lowest-none-empty
  using goal1(2) unfolding not-not .
  then have ipv4range-split1 r = (None, r) unfolding ipv4range-split1-def Let-def
  by simp
  then show False using goal1(1) by simp
qed

```

```

lemma ipv4range-split1-distinct: (Some s, u) = ipv4range-split1 r ⇒ ipv4range-empty
(ipv4range-intersection (prefix-to-range s) u)
proof -
  case goal1
  note ne = ipv4range-split1-never-empty[OF goal1]
  have nn: ipv4range-lowest-element r ≠ None using ipv4range-split1-some-r-ne[OF
    goal1, unfolded ipv4range-lowest-none-empty[symmetric]] .
  obtain a where ad: Some a = ipv4range-lowest-element r using nn by force

```

```

{
  fix rr :: 32 word × nat ⇒ 'a option × 32 wordinterval
  have (case find (const True) [s←map (Pair a) pfxes . valid-prefix s ∧ ipv4range-subset
(prefix-to-range s) r] of None ⇒ (None, r)
        | Some m ⇒ rr m) = rr (the (find (const True) [s←map (Pair a)
pfxes . valid-prefix s ∧ ipv4range-subset (prefix-to-range s) r]))
    using ipv4range-split-innard-helper[OF ad[symmetric]] find-const-True
by fastforce
} note uf2 = this
from goal1 have u = ipv4range-setminus r (prefix-to-range s)
  unfolding ipv4range-split1-def Let-def
  unfolding ad[symmetric] option.cases
  unfolding uf2
  unfolding Pair-eq
  by (metis option.sel)
then show ?thesis by force
qed

```

```

function ipv4range-split :: 32 wordinterval ⇒ (ipv4addr × nat) list where
  ipv4range-split rs = (if ¬ipv4range-empty rs then case ipv4range-split1 rs of
(Some s, u) ⇒ s # ipv4range-split u | - ⇒ [] else [])
  by(simp, blast)

```

**termination** *ipv4range-split*

**proof**(*relation measure (card ∘ ipv4range-to-set), rule wf-measure, unfold in-measure comp-def*)

note vernichter = *ipv4range-empty-set-eq ipv4range-intersection-set-eq ipv4range-union-set-eq ipv4range-eq-set-eq*

case goal1

note some = *goal1(2)[unfolded goal1(3)]*

from *ipv4range-split1-never-empty*[OF this] have ¬ *ipv4range-empty (prefix-to-range x2)* .

thus ?case

unfolding vernichter

unfolding *ipv4range-split1-preserve*[OF some, unfolded vernichter, symmetric]

unfolding *card-Un-disjoint*[OF finite finite *ipv4range-split1-distinct*[OF some, unfolded vernichter]]

by (metis *add.commute add-left-cancel card-0-eq finite linorder-neqE-nat monoid-add-class.add.right-neutral not-add-less1*)

qed

**lemma** *unfold-rsplit-case*:

assumes *su*: (*Some s, u*) = *ipv4range-split1 rs*

shows (case *ipv4range-split1 rs* of (*None, u*) ⇒ [] | (*Some s, u*) ⇒ *s # ipv4range-split u*) = *s # ipv4range-split u*

using *su* by (metis *option.simps(5) split-conv*)

**lemma** *ipv4range-split-union*: *ipv4range-eq (list-to-wordinterval (map prefix-to-range (ipv4range-split r))) r*

```

proof(induction r rule: ipv4range-split.induct, subst ipv4range-split.simps, case-tac
ipv4range-empty rs)
  case goal1
    thm Empty-WordInterval-set-eq ipv4range-eq-set-eq[of Empty-WordInterval rs,
unfolded ipv4range-to-set-def Empty-WordInterval-set-eq]
    show ?case using goal1(2)
    by(simp add: ipv4range-eq-set-eq[of Empty-WordInterval rs, unfolded ipv4range-to-set-def
Empty-WordInterval-set-eq] ipv4range-to-set-def)
  next
    case goal2
    obtain u s where su: (Some s, u) = ipv4range-split1 rs using r-split1-not-none[OF
goal2(2)] by (metis option.collapse surjective-pairing)
    note mIH = goal2(1)[OF goal2(2) su, of s]
    show ?case
      unfolding eqTrueI[OF goal2(2)]
      unfolding if-True
      unfolding unfold-rsplit-case[OF su]
      unfolding ipv4range-eq-set-eq
      unfolding ipv4range-to-set-def
      unfolding list.map
      unfolding list-to-wordinterval-set-eq-simp
      using mIH[unfolded ipv4range-eq-set-eq ipv4range-to-set-def]
      using ipv4range-split1-preserve[OF su, unfolded ipv4range-eq-set-eq ipv4range-to-set-def
ipv4range-union-def]
      unfolding wordinterval-union-set-eq
      by presburger
qed

```

```

value ipv4range-split (RangeUnion (WordInterval (ipv4addr-of-dotdecimal (64,0,0,0))
0x5FEFBBCC) (WordInterval 0x5FEEBB1C (ipv4addr-of-dotdecimal (127,255,255,255))))
value ipv4range-split (WordInterval 0 (ipv4addr-of-dotdecimal (255,255,255,254)))

```

10.0.0.0/8 – 10.8.0.0/16

```

lemma map ( $\lambda(ip,n). (dotdecimal-of-ipv4addr\ ip, n)) (ipv4range-split (ipv4range-setminus
(ipv4range-range ((ipv4addr-of-dotdecimal (10,0,0,0)), (ipv4addr-of-dotdecimal
(10,255,255,255))))
(ipv4range-range ((ipv4addr-of-dotdecimal (10,8,0,0)), (ipv4addr-of-dotdecimal
(10,8,255,255)))))) =
[[((10, 0, 0, 0), 13), ((10, 9, 0, 0), 16), ((10, 10, 0, 0), 15), ((10, 12, 0, 0),
14), ((10, 16, 0, 0), 12), ((10, 32, 0, 0), 11), ((10, 64, 0, 0), 10),
((10, 128, 0, 0), 9)] by eval$ 
```

```

declare ipv4range-split.simps[simp del]

```

```

corollary ipv4range-split: ( $\bigcup (prefix-to-ipset\ ' (set (ipv4range-split\ r)))$ ) = wordinterval-to-set
r

```

**proof** –

```

have prefix-to-range-set-eq-fun: prefix-to-ipset = (wordinterval-to-set  $\circ$  prefix-to-range)

```

```

by(simp add: prefix-to-range-set-eq fun-eq-iff)

{ fix r
  have  $\bigcup ((\text{wordinterval-to-set} \circ \text{prefix-to-range}) \text{ ` set } (\text{ipv4range-split } r)) =$ 
    ( $\text{wordinterval-to-set } (\text{list-to-wordinterval } (\text{map prefix-to-range } (\text{ipv4range-split } r))))$ )
  by (metis (erased, lifting) list.map-comp list-to-wordinterval-set-eq set-map)
  also have ... = ( $\text{wordinterval-to-set } r$ )
  by (metis ipv4range-eq-set-eq ipv4range-split-union ipv4range-to-set-def)
  finally have  $\bigcup ((\text{wordinterval-to-set} \circ \text{prefix-to-range}) \text{ ` set } (\text{ipv4range-split } r))$ 
    =  $\text{wordinterval-to-set } r$  .
} note ipv4range-eq-eliminator=this[of r]

show ?thesis
unfolding prefix-to-range-set-eq-fun
using ipv4range-eq-eliminator by auto
qed
corollary ipv4range-split-single: ( $\bigcup (\text{prefix-to-ipset } \text{ ` (set } (\text{ipv4range-split } (\text{WordInterval } \text{start } \text{end})))) = \{\text{start} \dots \text{end}\}$ )
using ipv4range-split by simp

lemma all-valid-Ball: Ball (set (ipv4range-split r)) valid-prefix
proof(induction r rule: ipv4range-split.induct, subst ipv4range-split.simps, case-tac
  ipv4range-empty rs)
  case goal1 thus ?case
  by(simp only: not-True-eq-False if-False Ball-def set-simps empty-iff) clarify
next
  case goal2
  obtain u s where su: (Some s, u) = ipv4range-split1 rs using r-split1-not-none[OF
    goal2(2)] by (metis option.collapse surjective-pairing)
  note mIH = goal2(1)[OF goal2(2) su refl]
  have vpfx: valid-prefix s
  proof -
    obtain a where a: ipv4range-lowest-element rs = Some a
    using goal2(2)[unfolded flipnot[OF ipv4range-lowest-none-empty, symmetric]]
    by force
    obtain m where m: find (const True) [s ← map (Pair a) pfxes . valid-prefix s
       $\wedge$  ipv4range-subset (prefix-to-range s) rs] = Some m
    using ipv4range-split-innard-helper[OF a, unfolded flipnot[OF find-const-True,
      symmetric]]
    by force
    note su[unfolded ipv4range-split1-def Let-def]
    then have (Some s, u) =
      (case find (const True) [s ← map (Pair a) pfxes . valid-prefix s  $\wedge$ 
        ipv4range-subset (prefix-to-range s) rs] of None  $\Rightarrow$  (None, rs)
        | Some m  $\Rightarrow$  (find (const True) [s ← map (Pair a) pfxes . valid-prefix
          s  $\wedge$  ipv4range-subset (prefix-to-range s) rs], ipv4range-setminus rs (prefix-to-range
          m)))
    unfolding a by simp

```

```

then have (Some s, u) =
  (Some m, ipv4range-setminus rs (prefix-to-range m))
  unfolding m by simp
moreover
note find-in[OF m[symmetric]]
ultimately
show valid-prefix s by simp
qed
show ?case
  unfolding eqTrueI[OF goal2(2)]
  unfolding if-True
  unfolding unfold-rsplit-case[OF su]
  unfolding list.set
  using mIH vpfx
  by blast
qed

```

**corollary** *ipv4range-split-bitmask:*

$(\bigcup ((\lambda (base, len). \text{ipv4range-set-from-bitmask } base \ len) \text{ ' (set (ipv4range-split } r))) ) = \text{wordinterval-to-set } r$

**proof** –

— without valid prefix assumption

**have** *prefix-to-ipset-subset-ipv4range-set-from-bitmask-helper:*

$\bigwedge X. (\bigcup_{x \in X. \text{prefix-to-ipset } x} \subseteq (\bigcup_{x \in X. \text{case } x \text{ of } (x, xa) \Rightarrow \text{ipv4range-set-from-bitmask } x \text{ } xa)$

**apply**(*rule*)

**using** *prefix-to-ipset-subset-ipv4range-set-from-bitmask[simplified pfxm-prefix-def pfxm-length-def]* **by** *fastforce*

**have** *ipv4range-set-from-bitmask-subseteq-prefix-to-ipset-helper:*

$\bigwedge X. \forall x \in X. \text{valid-prefix } x \implies (\bigcup_{x \in X. \text{case } x \text{ of } (x, xa) \Rightarrow \text{ipv4range-set-from-bitmask } x \text{ } xa) \subseteq (\bigcup_{x \in X. \text{prefix-to-ipset } x}$

**apply**(*rule*)

**apply**(*rename-tac x*)

**apply**(*safe*)

**apply**(*rename-tac a b*)

**apply**(*erule-tac x=(a,b) in ballE*)

**apply**(*simp-all*)

**apply**(*drule wordinterval-to-set-ipv4range-set-from-bitmask*)

**apply**(*rule-tac x=(a, b) in bexI*)

**apply**(*simp-all add: pfxm-prefix-def pfxm-length-def*)

**done**

**show** *?thesis*

**unfolding** *ipv4range-split[symmetric]*

**apply**(*simp add: ipv4range-range-def*)

**apply**(*rule*)

```

    apply(simp add: ipv4range-set-from-bitmask-subseteq-prefix-to-ipset-helper all-valid-Ball)
    apply(simp add: prefix-to-ipset-subset-ipv4range-set-from-bitmask-helper)
  done
qed
corollary ipv4range-split-bitmask-single:
  
$$\left( \bigcup ((\lambda (base, len). \text{ipv4range-set-from-bitmask } base \ len) \text{ ` (set (ipv4range-split (ipv4range-range (start, end)))))) \right) = \{start .. end\}$$

  using ipv4range-split-bitmask ipv4range-range.simps by simp

end

end

```