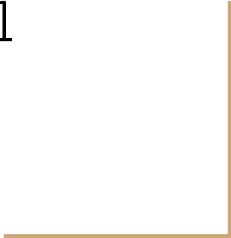


Programming, Problem Solving, and Algorithms

CPSC203, 2019 W1



Announcements

Project 3 released soon. Due 11:59p, Nov 29.

“Problem of the Day” continues!

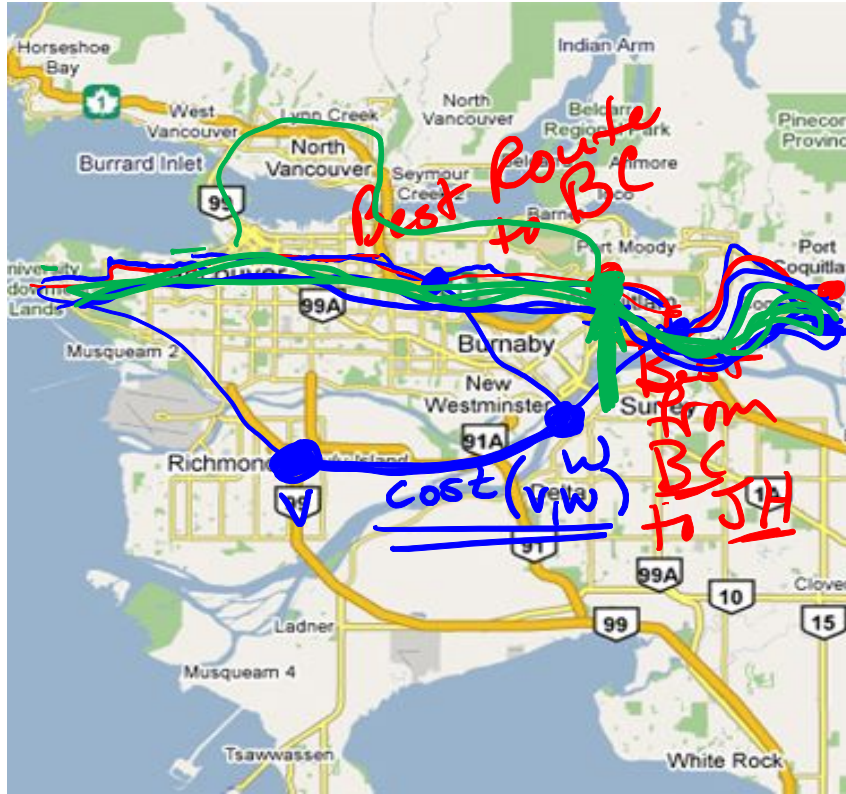
Today:

Shortest Path

Maps!

How many Starbucks are in Vancouver?

Single Source Shortest Path



Given a start vertex (source) s , find the path of least total cost from s to every vertex in the graph.

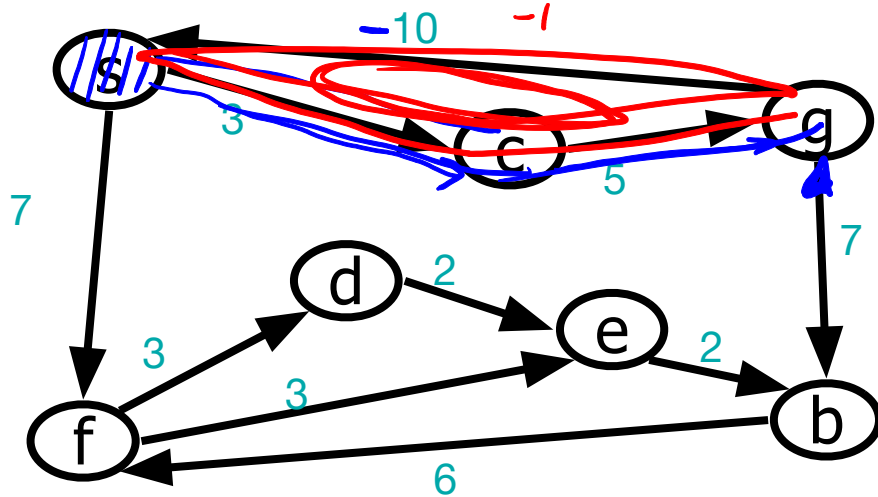
cost metric is time
but we could choose
distance, steps
traffic, transit switches.
of Starbucks on rt.
Need to elucidate "Cost".

Single Source Shortest Path

Input: directed graph G with non-negative edge weights, and a start vertex s. "source"

Output: A subgraph G' consisting of the shortest (minimum total cost) paths from s to every other vertex in the graph.

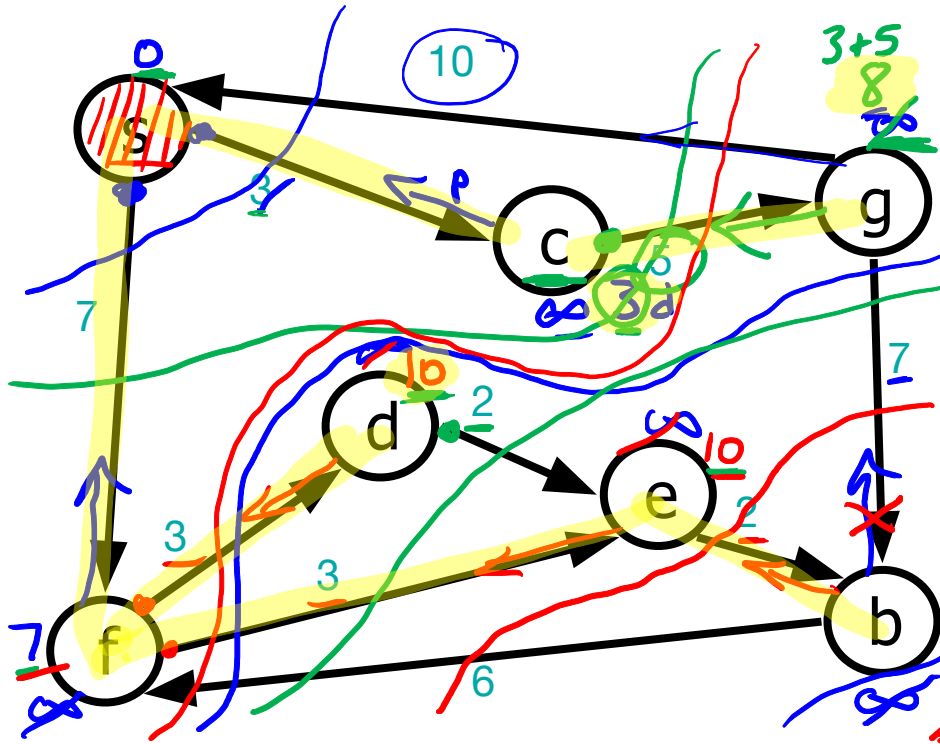
neg cycle no alg works
neg edge (no neg cycles)
Dijkstra's doesn't work.



Dijkstra's Algorithm (1959)

Shortest path doesn't exist if there is a negative wt cycle.

Single Source Shortest Path



Given a source vertex s , we wish to find the shortest path from s to every other vertex in the graph.

Initialize structure:

$dist[v]$: current best from s to v .
 Priority Queue Q of $d[v]$
 $pred[v] = w$ if (w, v) is in soln

Repeat these steps:

1. Label a new (unlabelled) vertex v , whose shortest distance has been found **s c f g d e b**
2. Update v 's neighbors with an improved distance

Single Source Shortest Path

Initialize structure:

1. For all v , $d[v] = \text{"infinity"}$, $p[v] = \text{null}$
current best distance
2. Initialize source: $d[s] = 0$
3. Initialize priority (min) queue of the $d[v]$
predecessor

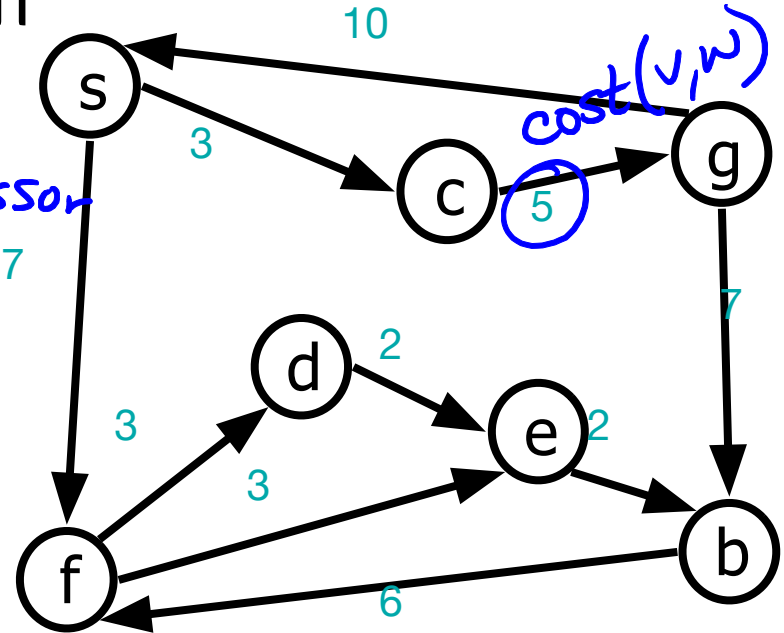
→ Repeat these steps n times:

- Find minimum $d[]$ unlabelled vertex: v
PQ operation
- Label vertex v *done*
- For all unlabelled neighbors w of v ,
"is better than"

If ($d[v] + \text{cost}(v,w) < d[w]$)

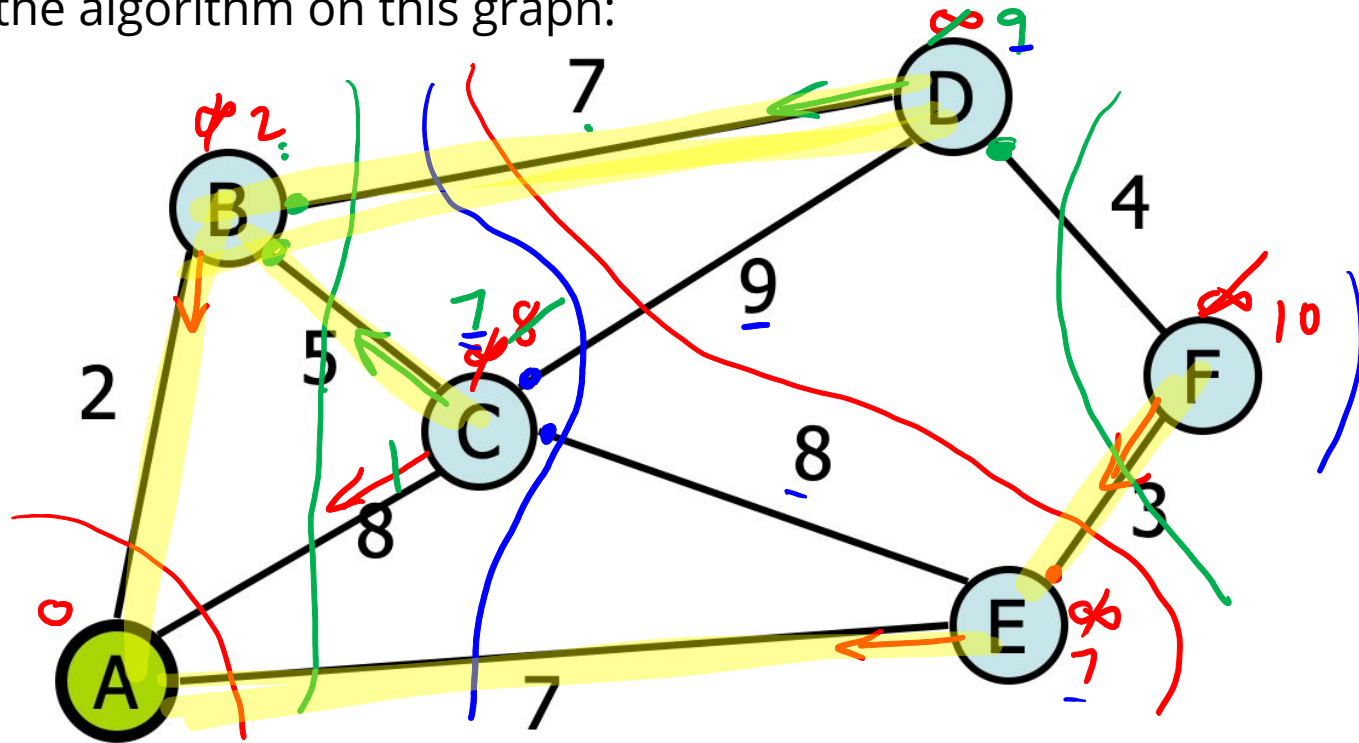
$d[w] = d[v] + \text{cost}(v,w)$ *update step.*

$p[w] = v$



Your Turn...

Execute the algorithm on this graph:



Dijkstra's Algorithm ^{PQ}

How is this algorithm similar to BFS/DFS?

Structure of final graph provided by

How is this algorithm different from BFS/DFS?

*no enqueueing of children
'cu'z vertices are already there*

Initialize structure:

1. For all v , $d[v] = \text{"infinity"}$, $p[v] = \text{null}$
2. Initialize source: $d[s] = 0$
3. Initialize priority (min) queue
4. Initialize set of labeled vertices to \emptyset .

While PQ is not empty:

Repeat these steps n times:

- Find & remove minimum $d[]$ unlabelled vertex v *@ degree*

- Label vertex v ✓

- *Same*
For all unlabelled neighbors w of v ,

If $\text{cost}(v, w) < d[w]$
 $d[w] = \text{cost}(v, w)$
 $p[w] = v$

update

Map applications

Three parts:

1. Assembling the data - OSM, local data stores, statsCan, etc. This is mostly the art of assembling geodataframes.
2. Computing on the data - osmnx simplifies graph algorithms and computation, but also supports other spatial computation.
3. Visualizing the data - matplotlib for static maps, folium for interactive maps.

POTD #36 Thu

<https://github.students.cs.ubc.ca/cpsc203-2019w-t1/potd36>

Describe any snags you run into:

1. Line ____: _____
2. Line ____: _____
3. Line ____: _____
4. Line ____: _____
5. Line ____: _____

ToDo for next class...

POTD: Continue every weekday! Submit to repo.

Reading: TLACS Ch 10 & 12 (lists and dictionaries)

References:

<https://www.youtube.com/watch?v=wsSEKm-rU6U>

<https://github.com/gboeing/osmnx-examples/tree/master/notebooks>

<https://gist.github.com/psychemedia/b49c49da365666ba9199d2e27d002d07>