

# Nonnegative Matrix Theory and its Applications to Communication Network Problems

Chee Wei Tan

California Institute of Technology

University of Illinois, Chicago

26th Aug 2009

# Acknowledgement

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- Mung Chiang (Princeton)
- R. Srikant (UIUC)
- Shmuel Friedland (UIC)
- Steven Low (Caltech)

# Outline

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- Model
- Sum Rate Maximization
- Friedland-Karlin Inequalities & Minimax Theorem
- Global Optimization Algorithm
- Max-min Weighted SIR & Nonlinear Perron-Frobenius Theory
- Fast Polynomial-time Algorithms
- Conclusion

# What makes a problem easy or hard

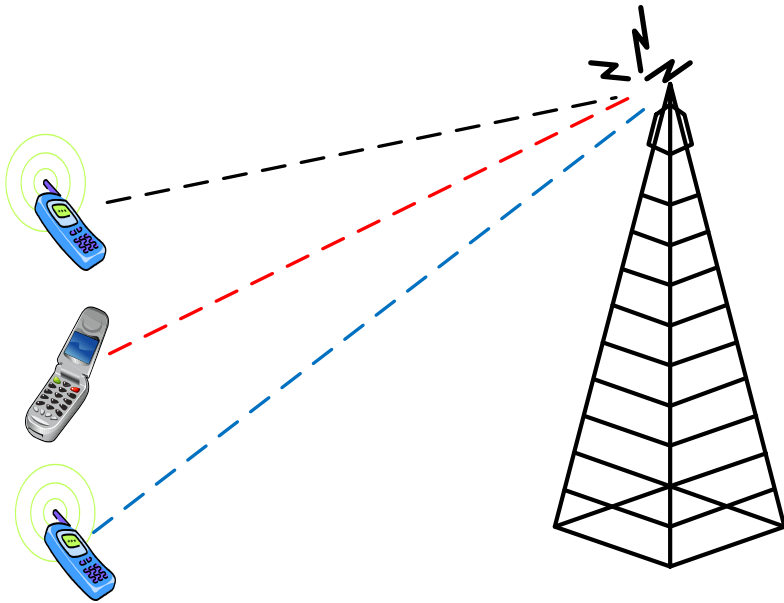
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... the great watershed in optimization isn't between linearity and nonlinearity, but **convexity** and **nonconvexity**.

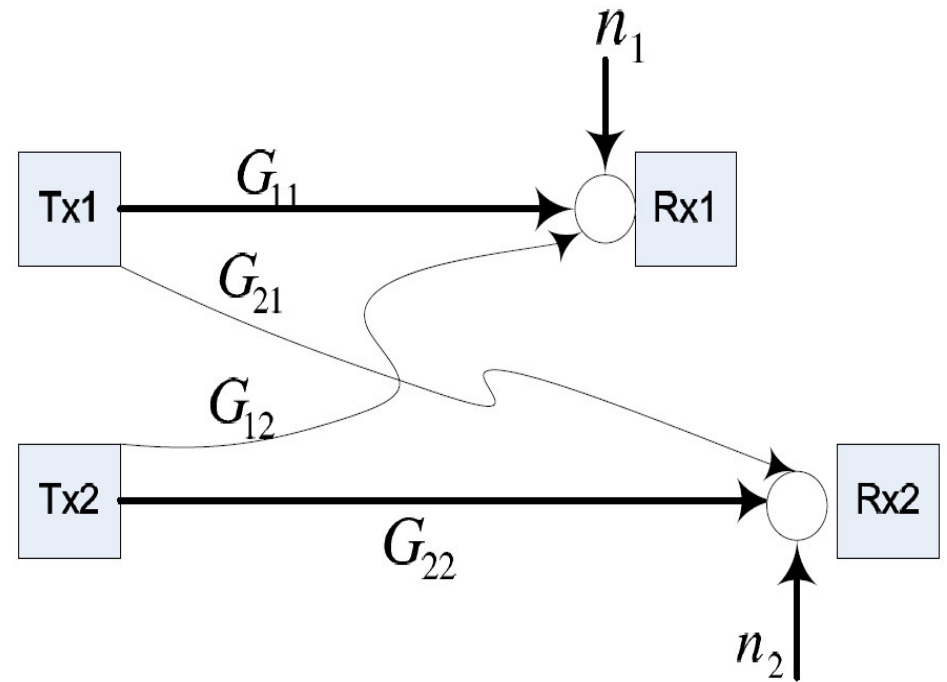
– *SIAM Review* 1993, R. Rockafellar

- Linear inequality theory & nonconvex integer programming (1947)
- Semidefinite matrix theory & nonconvex quadratic programming (1995)
- Nonnegative matrix theory & nonconvex cone programming ([this talk](#))

# Model



(a) Cellular wireless network



(b) Interference channel

# Performance Metric

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- Signal-to-Interference Ratio:

$$\text{SIR}_l(\mathbf{p}) = \frac{G_{ll}p_l}{\sum_{j \neq l} G_{lj}p_j + n_l}$$

with  $G_{lj}$  the channel gains from transmitter  $j$  to receiver  $l$  and  $n_l$  the additive white Gaussian noise (AWGN) power at receiver  $l$

- Attainable data rate (nats per channel use) is a function of  $\text{SIR}$ , e.g., Shannon capacity formula  $r_l = \log(1 + \text{SIR}_l)$
- Total power constraints  $\mathbf{1}^\top \mathbf{p} \leq \bar{P}$

# Interference Parameters

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- Let  $\mathbf{F}$  be a nonnegative matrix with entries:

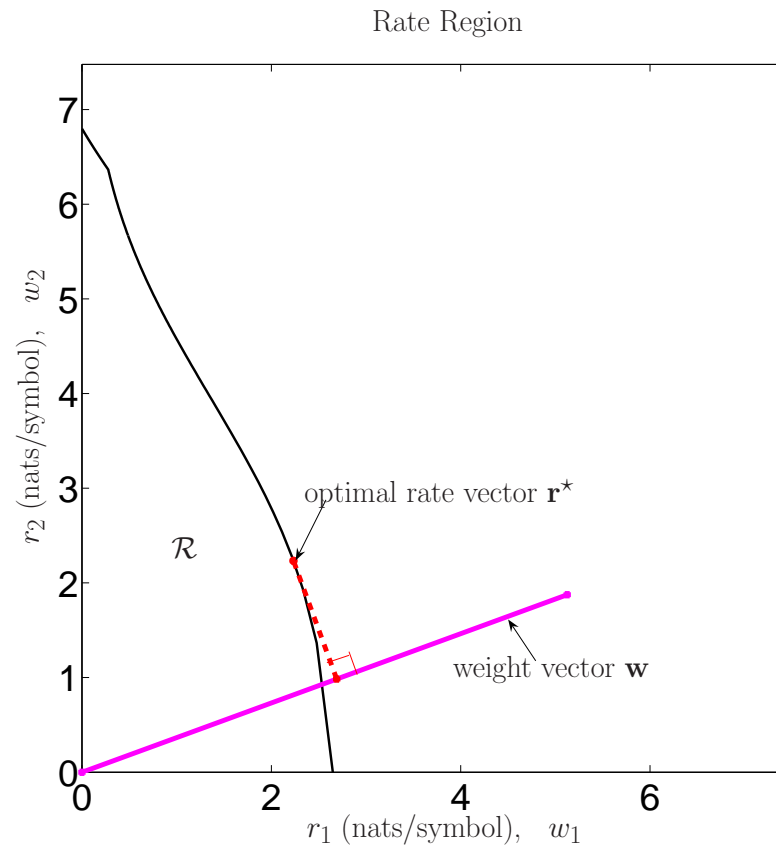
$$F_{ij} = \begin{cases} 0, & \text{if } i = j \\ \frac{G_{ij}}{G_{ii}}, & \text{if } i \neq j \end{cases}$$

and

$$\mathbf{v} = \left( \frac{n_1}{G_{11}}, \frac{n_2}{G_{22}}, \dots, \frac{n_L}{G_{LL}} \right)^\top.$$

# Sum Rate Geometry Illustration

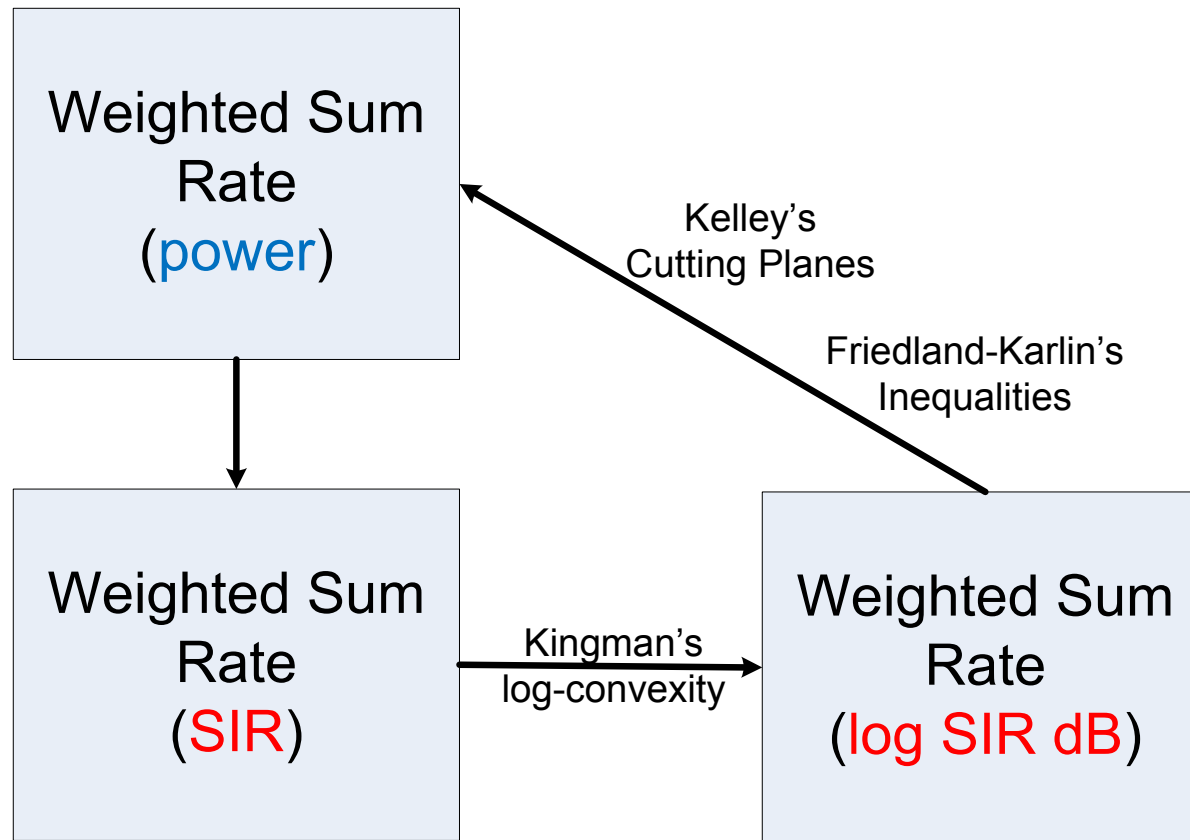
$$\begin{array}{ll}\text{maximize} & \sum_l w_l \log(1 + \text{SIR}_l(\mathbf{p})) = \sum_l w_l r_l \\ \text{subject to} & 0 \leq p_l \leq \bar{p}_l \quad \forall l, \\ \text{variables:} & p_l \quad \forall l.\end{array}$$





# Solution Map

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Friedland and Tan, *Nonnegative Matrix Inequalities and Applications to Multiuser Communication Problems*,  
submitted to SIAM Journ. on Matrix Analysis & Applications, 2009

# Sum Shannon Rate Global Optimization

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- Nonlinear map between power  $\mathbf{p}$  and SIR  $\gamma = \exp(\tilde{\gamma})$ :

$$\mathbf{p}^* = (\mathbf{I} - \text{diag}(\exp(\tilde{\gamma}^*))\mathbf{F})^{-1} \text{diag}(\exp(\tilde{\gamma}^*))\mathbf{v} \quad (1)$$

- Constraints of  $\mathbf{p} \leq \bar{\mathbf{p}}$  into spectral radius constraints  $\tilde{\gamma}$
- Convert into **convex maximization** (dB domain):

$$\begin{aligned} & \text{maximize} && \sum_l w_l \log(1 + \exp(\tilde{\gamma}_l)) \\ & \text{subject to} && \log \rho(\text{diag}(\exp(\tilde{\gamma}))(\mathbf{F} + (1/\bar{p}_l)\mathbf{v}\mathbf{e}_l^\top)) \leq 0 \quad \forall l, \\ & \text{variables:} && \tilde{\gamma}_l, \quad \forall l. \end{aligned}$$

# Nonnegative Matrix Theory: Minimax Theorem

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- **Theorem 1.** *Friedland-Karlin inequality [FriedlandKarlin'75]: For any irreducible nonnegative matrix  $\mathbf{A}$ ,*

$$\prod_l ((\mathbf{A}\mathbf{z})_l / z_l)^{x_l y_l} \geq \rho(\mathbf{A})$$

*for all strictly positive  $\mathbf{z}$ , where  $\mathbf{x}$  and  $\mathbf{y}$  are the Perron and left eigenvectors of  $\mathbf{A}$  respectively. Equality holds in (2) if and only if  $\mathbf{z} = a\mathbf{x}$  for some positive  $a$ .*

- Donsker-Varadhan's variational principle (1975):

$$\max_{\boldsymbol{\lambda} \geq 0, \mathbf{1}^\top \boldsymbol{\lambda} = 1} \min_{\mathbf{p} \geq 0} \sum_l \lambda_l \frac{(\mathbf{A}\mathbf{p})_l}{p_l} = \min_{\mathbf{p} \geq 0} \max_{\boldsymbol{\lambda} \geq 0, \mathbf{1}^\top \boldsymbol{\lambda} = 1} \sum_l \lambda_l \frac{(\mathbf{A}\mathbf{p})_l}{p_l}$$

- Extensions (see [FriedlandTan'09])

# Sum Shannon Rate Global Optimization

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- **Convex Maximization** (dB domain):

$$\begin{aligned} & \text{maximize} && \sum_l w_l \log(1 + \exp(\tilde{\gamma}_l)) \\ & \text{subject to} && \log \rho(\text{diag}(\exp(\tilde{\gamma}))(\mathbf{F} + (1/\bar{p}_l)\mathbf{v}\mathbf{e}_l^\top)) \leq 0 \quad \forall l, \\ & \text{variables:} && \tilde{\gamma}_l, \quad \forall l. \end{aligned}$$

- Relaxation of the constraint set by the **Friedland-Karlin Inequalities**:

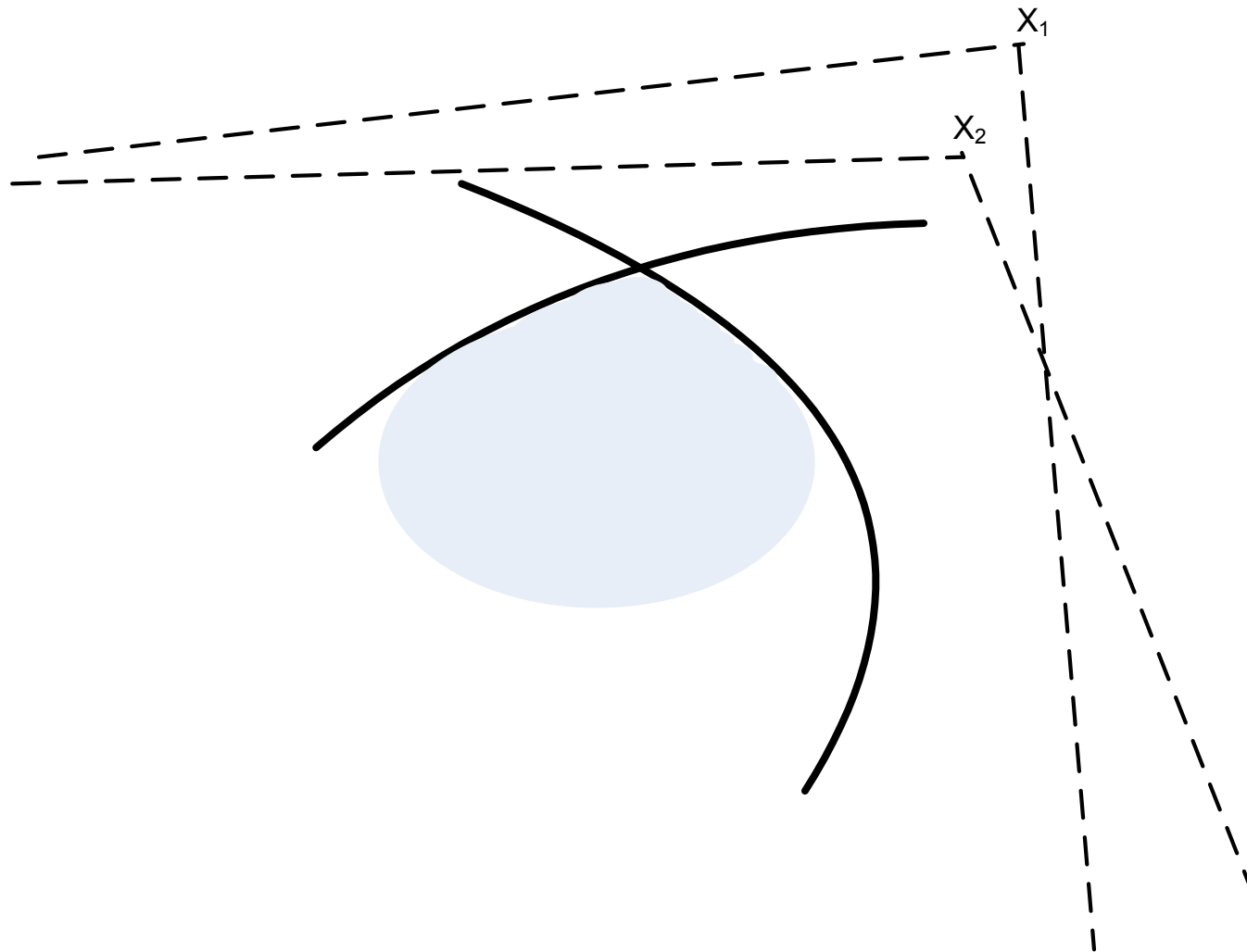
$$\prod_l \gamma_l^{x_l(\mathbf{A})y_l(\mathbf{A})} \rho(\mathbf{A}) \leq \rho(\text{diag}(\gamma)\mathbf{A})$$

$$\sum_l x_l(\mathbf{A})y_l(\mathbf{A})\tilde{\gamma}_l + \log \rho(\mathbf{A}) \leq \log \rho(\text{diag}(\exp(\tilde{\gamma}))\mathbf{A}) \quad (\text{dB domain}).$$

- Outer approximation algorithm (**Kelley's cutting planes**)

# Outer Approximation Algorithm Illustration

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# Sum Rate Global Optimization: Algorithm

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## Algorithm 1. [Sum Rate Outer Approximation Algorithm]

1. *Compute the vertices of the enclosing linear polyhedron  $D^{(0)}$ , described by the set of constraints:*

$$\sum_j (\mathbf{x}(\mathbf{F} + (1/\bar{p}_l)\mathbf{v}\mathbf{e}_l^\top) \circ \mathbf{y}(\mathbf{F} + (1/\bar{p}_l)\mathbf{v}\mathbf{e}_l^\top))_j \tilde{\gamma}_j + \log \rho(\mathbf{F} + (1/\bar{p}_l)\mathbf{v}\mathbf{e}_l^\top) \leq 0,$$

*and  $\tilde{\gamma}_l \geq -K$  for all  $l$ . Let  $V^{(0)}$  be the set of vertices of  $D^{(0)}$ . Set  $k = 1$  and go to Step 2.*

2. *Iteration  $k$ : Solve the problem:*

$$\begin{aligned} & \text{maximize} && \sum_l w_l \log(1 + e^{\tilde{\gamma}_l}) \\ & \text{subject to} && \tilde{\gamma}_l \in D^{(k-1)} \end{aligned}$$

*by selecting  $\max \{ \sum_l w_l \log(1 + e^{\tilde{\gamma}_l}) : v \in V^{(k-1)} \}$ . Let  $\tilde{\gamma}^k$  be the optimizer to (2).*

3. Compute

$$\mathbf{p}^k = \left( \mathbf{I} - \text{diag}(\exp(\tilde{\gamma}^k)) \mathbf{F} \right)^{-1} \text{diag}(\exp(\tilde{\gamma}^k)) \mathbf{v}.$$

4. If  $\mathbf{p}^k \leq \bar{\mathbf{p}}$ , stop:  $\tilde{\gamma}^k$  is the solution to (2) and  $\mathbf{p}^k$  is the solution to (1). Otherwise, let

$$\begin{aligned} J^k &= \{l : \log \rho(\text{diag}(\exp(\tilde{\gamma}^k))(\mathbf{F} + (1/\bar{p}_l) \mathbf{v} \mathbf{e}_l^\top)) \\ &= \max_{1 \leq j \leq L} \log \rho(\text{diag}(\exp(\tilde{\gamma}^k))(\mathbf{F} + (1/\bar{p}_j) \mathbf{v} \mathbf{e}_j^\top))\} \end{aligned} \quad (2)$$

and choose any  $j^k \in J^k$ .

5. Compute the left eigenvector  $\mathbf{y}_{j^k}$  and right (Perron) eigenvector  $\mathbf{x}_{j^k}$  of  $\text{diag}(\exp(\tilde{\gamma}^k))(\mathbf{F} + (1/\bar{p}_{j^k}) \mathbf{v} \mathbf{e}_{j^k}^\top)$ . Set

$$\begin{aligned} G_{j^k}^k(\tilde{\gamma}) &= \log \rho(\text{diag}(\exp(\tilde{\gamma}^k))(\mathbf{F} + (1/\bar{p}_{j^k}) \mathbf{v} \mathbf{e}_{j^k}^\top)) + \\ &\quad \frac{[\exp(\tilde{\gamma}^k) \circ \mathbf{x}_{j^k} \circ \mathbf{y}_{j^k}]^\top (\tilde{\gamma} - \tilde{\gamma}^k)}{\rho(\text{diag}(\exp(\tilde{\gamma}^k))(\mathbf{F} + (1/\bar{p}_{j^k}) \mathbf{v} \mathbf{e}_{j^k}^\top))}. \end{aligned} \quad (3)$$

6. Set  $D^{(k)} = D^{(k-1)} \cap \{\tilde{\gamma} : G_{j^k}^k(\tilde{\gamma}) \leq 0\}$ ,  $V^{(k)} = \{\text{extreme points of } D^{(k)}\}$ .

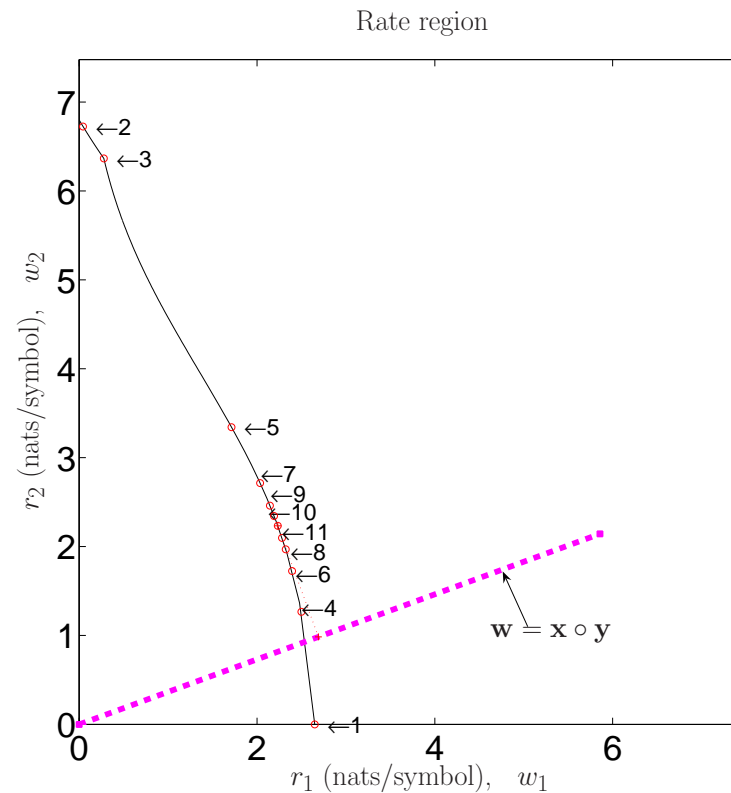
7. Set  $k \leftarrow k + 1$ . Go to Step 2.

- Step 3 yields a **feasible** power vector  $\hat{\mathbf{p}}^k$ :  $\hat{p}_l^k = \min\{p_l^k, \bar{p}_l\}$  for all  $l$ .



# Global Optimizing Sum Rate: Examples

- $\lim_{k \rightarrow \infty} \min \left\{ \left( \mathbf{I} - \text{diag}(\exp(\tilde{\gamma}^k)) \mathbf{F} \right)^{-1} \text{diag}(\exp(\tilde{\gamma}^k)) \mathbf{v}, \bar{p}_l \right\} = \mathbf{p}^*$
- Fast convergence in numerical examples



# Global Optimizing Sum Rate: Examples

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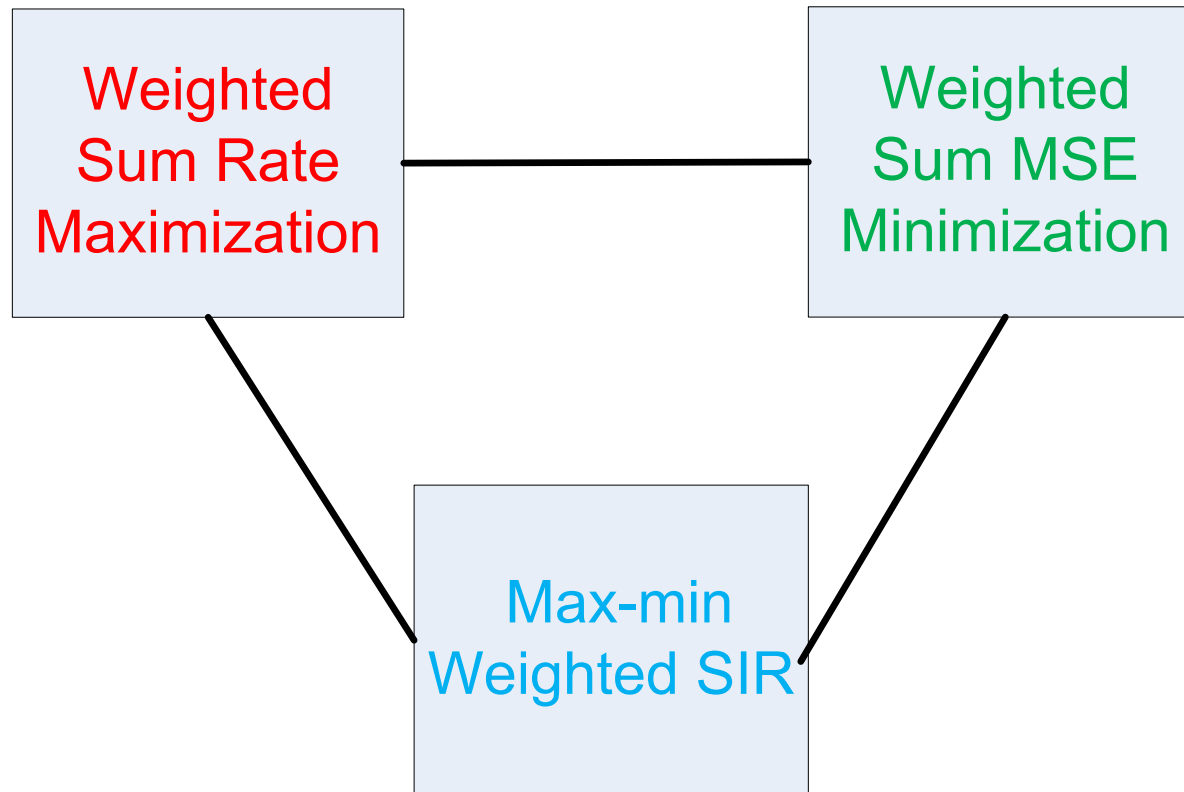
- Efficient and fast for **small to medium-sized** networks

Problem size	Maximal number of generated vertices	Number of iterations	CPU time (minutes)
2	15	12	0.062
4	139	760	4.1
6	14022	1238	83
8	283681	1968	468

Table 1: A comparison of the typical convergence and complexity statistics of Algorithm 1 with the problem size. The CPU time is computed based on an implementation on a 64-bit Sun/Solaris 10 (SunOS 5.10) computer.

# Nonconvex Power Control Problems

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Tan, Chiang and Srikant, IEEE INFOCOM 2009 & IEEE ISIT 2009

# Max-min Weighted SIR

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$$\text{SIR}_l(\mathbf{p}) = \frac{G_{ll}p_l}{\sum_{j \neq l} G_{lj}p_j + n_l} \quad \mathbf{B} = \mathbf{F} + (1/\bar{P})\mathbf{v}\mathbf{1}^\top$$

•

$$\begin{aligned} & \text{maximize} \quad \min_l \frac{\text{SIR}_l(\mathbf{p})}{\beta_l} \\ & \text{subject to} \quad \mathbf{1}^\top \mathbf{p} \leq \bar{P}, \quad \mathbf{p} \geq \mathbf{0}, \\ & \text{variables:} \quad \mathbf{p}. \end{aligned}$$

- Optimal value:  $1/\rho(\text{diag}(\beta)\mathbf{B})$  and optimal solution:  $(\bar{P}/\mathbf{1}^\top \mathbf{x}(\text{diag}(\beta)\mathbf{B}))\mathbf{x}(\text{diag}(\beta)\mathbf{B})$ .

# Nonlinear Perron-Frobenius Theory

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- Find  $(\check{\lambda}, \check{\mathbf{s}})$  in

$$\lambda \mathbf{s} = \mathbf{A}\mathbf{s} + \mathbf{b}, \quad \lambda \in \mathbb{R}, \quad \mathbf{s} \geq \mathbf{0}, \quad \|\mathbf{s}\| = 1,$$

where  $\mathbf{A}$  and  $\mathbf{b}$  is a square irreducible nonnegative matrix and nonnegative vector, respectively and  $\|\cdot\|$  a monotone vector norm.

- $(\check{\lambda}, \check{\mathbf{s}})$  is Perron-Frobenius eigenvalue-vector pair of  $\mathbf{A} + \mathbf{b}\mathbf{c}_*^\top$ , where

$$\mathbf{c}_* = \arg \max_{\|\mathbf{c}\|_* = 1} \rho(\mathbf{A} + \mathbf{b}\mathbf{c}^\top),$$

where  $\|\cdot\|_*$  is the dual norm of  $\|\cdot\|$ , and  $\check{\mathbf{s}} = (\mathbf{A}\check{\mathbf{s}} + \mathbf{b})/\|\mathbf{A}\check{\mathbf{s}} + \mathbf{b}\|$ .

V. D. Blondel, L. Ninove and P. Van Dooren, *An affine eigenvalue problem on the nonnegative orthant*, Linear Algebra & its Applications, 2005

# Nonlinear Perron-Frobenius Theory: Max-min SIR

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$$\text{SIR}_l(\mathbf{p}^*) = \tau^* \beta_l \Rightarrow \frac{(p_l^*/\bar{P})}{\sum_{j \neq l} F_{lj}(p_l^*/\bar{P}) + (v_l/\bar{P})} = \tau^* \beta_l$$

Let  $\mathbf{s}^* = (1/\bar{P})\mathbf{p}^*$ :

$$(1/\tau^*)\mathbf{s}^* = \text{diag}(\boldsymbol{\beta})\mathbf{F}\mathbf{s}^* + (1/\bar{P})\text{diag}(\boldsymbol{\beta})\mathbf{v}, \quad \|\mathbf{s}\|_1 = 1$$

- ■  $s_l = p_l/\bar{p}_l$ ,  $\mathbf{A} = \text{diag}(\boldsymbol{\beta})\mathbf{F}$ ,  $\mathbf{b} = (1/\bar{P})\text{diag}(\boldsymbol{\beta})\mathbf{v}$  and  $\lambda = 1/\tau^*$
- $\|\cdot\| = \|\cdot\|_1 \longleftrightarrow \|\cdot\|_* = \|\cdot\|_\infty \quad \& \quad \mathbf{c}_* = \mathbf{1}$
- $(\check{\lambda}, \check{\mathbf{s}})$  is the Perron-Frobenius eigenvalue and vector pair of  $\text{diag}(\boldsymbol{\beta})(\mathbf{F} + (1/\bar{P})\mathbf{v}\mathbf{1}^\top) = \text{diag}(\boldsymbol{\beta})\mathbf{B}$

# Max-min Weighted SIR: Algorithm

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- **Algorithm 2. [Max-min Weighted SIR]**

1. *Update power  $\mathbf{p}(k + 1)$ :*

$$p_l(k + 1) = \left( \frac{\beta_l}{\text{SIR}(\mathbf{p}(k))} \right) p_l(k) \quad \forall l$$

2. *Normalize  $\mathbf{p}(k + 1)$ :*      $\mathbf{p}(k + 1) = \mathbf{p}(k + 1) / \mathbf{1}^\top \mathbf{p}(k + 1) \cdot \bar{P}$

# Approximation to Weighted Sum Rate Problem

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- maximize  $\sum_{l=1}^L w_l \log(1 + \text{SIR}_l(\mathbf{p}))$   
subject to  $\sum_{l=1}^L p_l \leq \bar{P}, \quad p_l \geq 0 \quad \forall l,$   
variables:  $p_l \quad \forall l.$  (4)
- Quasi-inverse [Wong54]:  $\mathbf{B}$  is a quasi-inverse of  $\tilde{\mathbf{B}} \geq \mathbf{0}$  if  
 $\mathbf{B} - \tilde{\mathbf{B}} = \mathbf{B}\tilde{\mathbf{B}} = \tilde{\mathbf{B}}\mathbf{B} \geq \mathbf{0}$
- $\rho(\tilde{\mathbf{B}}) = \frac{\rho(\mathbf{B})}{1+\rho(\mathbf{B})}$
- $\mathbf{x}(\tilde{\mathbf{B}}) = \mathbf{x}(\mathbf{B})$  &  $\mathbf{y}(\tilde{\mathbf{B}}) = \mathbf{y}(\mathbf{B})$



# Interference & SNR Regime

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- Recall the matrix

$$\mathbf{B} = \mathbf{F} + (1/\bar{P})\mathbf{v}\mathbf{1}^\top$$

- (High SNR regime)  $\tilde{\mathbf{B}}$  does not exist

or any nonnegative matrix with a zero trace & positive off-diagonals

- (Low SNR regime)  $\tilde{\mathbf{B}}$  always exists

or any nonnegative matrix that is a dyad

- (Low interference/moderate SNR regime)  $\tilde{\mathbf{B}}$  almost always exists

# Approximation to Weighted Sum Rate Problem

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**Theorem 2.** If  $\tilde{\mathbf{B}} \geq \mathbf{0}$ ,

$$\sum_{l=1}^L w_l \log(1 + \text{SIR}_l(\mathbf{p})) \leq \|\mathbf{w}\|_{\infty}^{\mathbf{x}(\mathbf{B}) \circ \mathbf{y}(\mathbf{B})} \log(1 + 1/\rho(\mathbf{B})) \quad (5)$$

for all feasible  $\mathbf{p}$ .

*Equality is achieved if and only if  $\mathbf{w} = \mathbf{x}(\mathbf{B}) \circ \mathbf{y}(\mathbf{B})$ , and  $\text{SIR}_l(\mathbf{p}^*) = (1/\rho(\mathbf{B}))\mathbf{1} \ \forall l$ . In this case,  $\mathbf{p}^* = \mathbf{x}(\mathbf{B})$ .*

- Max-min SIR ( $\beta = 1$ ) as an “approximation” algorithm

Simple Characterization via Nonnegative Matrix Theory

# “Charnes-Cooper” Trick plus Friedland-Karlin

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- Maximizing objective (linear-fractional) function:

$$\max_{\mathbf{p}} \prod_l (((\mathbf{I} + \mathbf{F})\mathbf{p} + \mathbf{v})_l / (\mathbf{F}\mathbf{p} + \mathbf{v})_l)^{w_l}. \quad (6)$$

- Change of variable:  $\mathbf{z} = (\mathbf{I} + \mathbf{B})\mathbf{p}$
- Transform maximization over  $\mathbf{p}$  to  $\mathbf{z}$  (when  $\tilde{\mathbf{B}} \geq \mathbf{0}$ ):

$$\min_{\mathbf{z}} \prod_l \left( \frac{(\tilde{\mathbf{B}}\mathbf{z})_l}{z_l} \right)^{w_l} \quad (7)$$

- Friedland-Karlin inequalities on spectrum of  $\tilde{\mathbf{B}}$  ( $\sim \mathbf{B}$ )

# Weighted Sum Rate Maximization: Exact Solution

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**Theorem 3.** *If  $\tilde{\mathbf{B}} \succeq \mathbf{0}$ , then the optimal solution to Sum Rate problem is given by  $\mathbf{p}^* = (\mathbf{I} - \tilde{\mathbf{B}})\mathbf{z}^* \succeq \mathbf{0}$ , where  $\mathbf{z}^*$  is given by*

$$z_l^* = \frac{w_l}{\sum_j w_j \tilde{B}_{jl} / (\tilde{\mathbf{B}}\mathbf{z}^*)_j} \quad (8)$$

*for all  $l$  and satisfies  $\mathbf{1}^\top \mathbf{z}^* - \mathbf{1}^\top \tilde{\mathbf{B}}\mathbf{z}^* = \bar{P}$ .*

# Weighted Sum Rate Maximization: Algorithm

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## Algorithm 3. [Sum Rate Maximization Algorithm]

1. Initialize an arbitrarily small  $\epsilon > 0$ .
2. Update auxiliary variable  $\mathbf{z}(k+1)$ :

$$z_l(k+1) = \frac{w_l}{\sum_j w_j \tilde{B}_{jl} / (\tilde{\mathbf{B}} \mathbf{z}(k))_j} + \epsilon \quad \forall l. \quad (9)$$

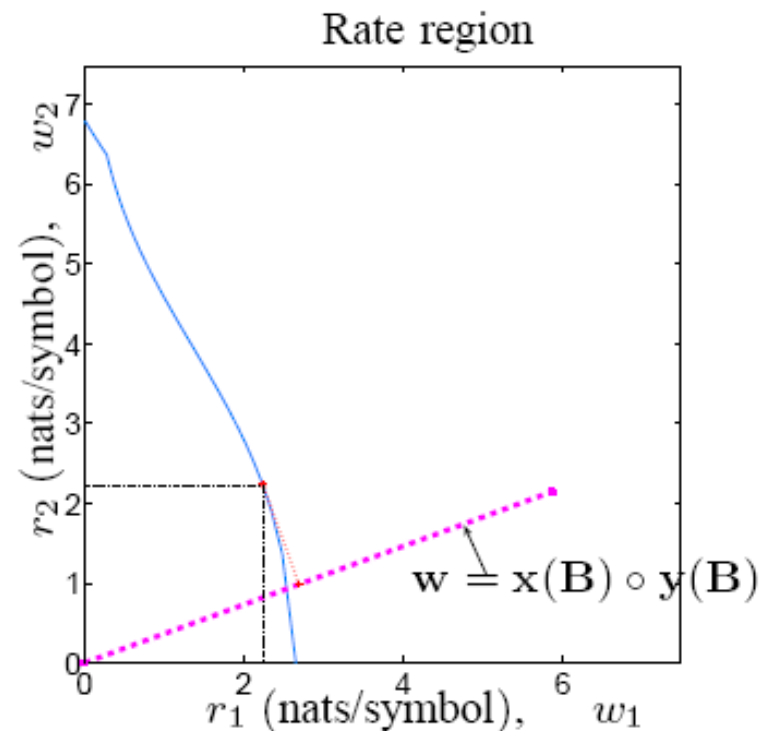
3. Update  $\mathbf{p}(k+1)$ :

$$p_l(k+1) = \frac{\text{SIR}_l(\mathbf{p}(k))}{1 + \text{SIR}_l(\mathbf{p}(k))} z_l(k+1) \quad \forall l. \quad (10)$$

4. Normalize  $\mathbf{p}(k+1)$ :  $\mathbf{p}(k+1) \leftarrow \mathbf{p}(k+1) \cdot \bar{P} / (\mathbf{1}^\top \mathbf{p}(k+1))$ .

# Link Sum Rate Maximization & Max-min SIR

- Friedland-Karlin inequalities & Arithmetic-geometric mean inequality &  $\mathbf{w} = \mathbf{x}(\mathbf{B}) \circ \mathbf{y}(\mathbf{B})$



New Link in Nonlinear Perron-Frobenius Theory & Optimization Theory

# Nonlinear Perron-Frobenius Theory: Friedland-Karlin Minimax

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- Minimax Theorem **[FriedlandKarlin'75]**:

$$\log \rho(\mathbf{A}) = \max_{\boldsymbol{\lambda} \geq 0, \mathbf{1}^\top \boldsymbol{\lambda} = 1} \min_{\mathbf{p} \geq 0} \sum_l \lambda_l \log \frac{(\mathbf{A}\mathbf{p})_l}{p_l} = \min_{\mathbf{p} \geq 0} \max_{\boldsymbol{\lambda} \geq 0, \mathbf{1}^\top \boldsymbol{\lambda} = 1} \sum_l \lambda_l \log \frac{(\mathbf{A}\mathbf{p})_l}{p_l}$$

- Nonlinear version ( $f(\mathbf{p}) = \mathbf{A}\mathbf{p} + \mathbf{b}$ ):

$$\begin{aligned} & \max_{\|\mathbf{c}\|_* = 1} \log \rho(\mathbf{A} + \mathbf{b}\mathbf{c}^\top) \\ &= \max_{\boldsymbol{\lambda} \geq 0, \mathbf{1}^\top \boldsymbol{\lambda} = 1} \min_{\|\mathbf{p}\| = 1} \sum_l \lambda_l \log \frac{(\mathbf{A}\mathbf{p} + \mathbf{b})_l}{p_l} \\ &= \min_{\|\mathbf{p}\| = 1} \max_{\boldsymbol{\lambda} \geq 0, \mathbf{1}^\top \boldsymbol{\lambda} = 1} \sum_l \lambda_l \log \frac{(\mathbf{A}\mathbf{p} + \mathbf{b})_l}{p_l}, \end{aligned}$$

where optimal  $\mathbf{p} = \mathbf{x}(\mathbf{A} + \mathbf{b}\mathbf{c}_*^\top)$  and  $\boldsymbol{\lambda} = \mathbf{x}(\mathbf{A} + \mathbf{b}\mathbf{c}_*^\top) \circ \mathbf{y}(\mathbf{A} + \mathbf{b}\mathbf{c}_*^\top)$

# Extension: Ideas for Fractional Programming

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- Sum-of-ratios fractional program (positive  $f(\mathbf{p})$ ,  $g(\mathbf{p})$ )  
**[FrenkSchaible05]:**

$$\min_{\mathbf{p}} \sum_l \frac{f_l(\mathbf{p})}{g_l(\mathbf{p})}. \quad (11)$$

- Domain transformation:  $\mathbf{z} = g(\mathbf{p})$
- Transform optimization over  $\mathbf{p}$  to  $\mathbf{z}$  (when  $g^{-1}(\mathbf{z}) \geq \mathbf{0}$ ):

$$\min_{\mathbf{z}} \sum_l \frac{(f(g^{-1}(\mathbf{z})))_l}{z_l} \quad (12)$$

- Link to minimax theorems in [Nonlinear Perron-Frobenius theory](#) when  $f(g^{-1}(\mathbf{z}))$  is concave and monotone



# Conclusion

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- Nonconvex power control problem in CDMA networks and OFDM/ADSL channels
- Eigenvalue characterization enables efficient global optimization of Weighted Sum Rate Maximization
- Link nonconvex nonnegative cone programming and nonnegative matrix theory
  - Friedland-Karlin inequalities
  - Log-convexity of spectral radius
- Solving nonconvex problems is an interesting art!

# Thank You

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<http://www.its.caltech.edu/~cheetan>

`cheetan@caltech.edu`