

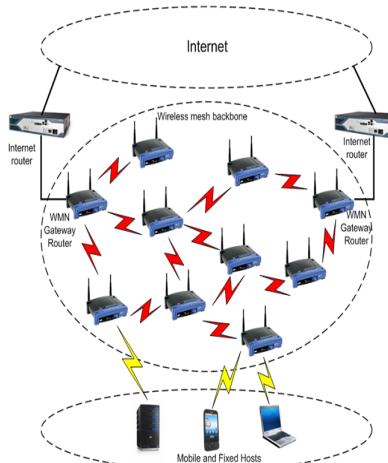
# BackLog Based Random Access in Wireless Networks : (In)Stability Issues and Fluid Limits

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Hong Kong Talk  
Hong Kong

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# Background and Motivation



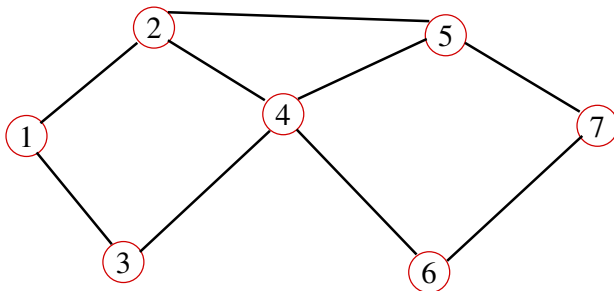
## Wireless Mesh Networks - Applications and Issues:

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- Applications
  - Electric Meters for Billing
  - One lap top per child program
  - Iridium satellite constellation
- Future Wireless Networks
  - Very Large Number of Users
  - No centralised control and no infrastructure
  - Autonomous radio nodes share resources
- Internet of things, sensor networks, ...

Challenge: Efficient use of resources in such networks

# WirelessMesh Network: Interference Model



- Users cannot Transmit Simultaneously because of Interference
- Interference relationship represented by **Conflict Graph**
- Feasible schedules correspond to Independent Sets e.g  $\{1, 5, 6\}$

# Wireless Mesh Network: Traffic Model

- Packets arrivals Poisson
- Packet Transmission Times are unit exponential
- $U_i(t)$  activity r.v. for node  $i$ , time  $t$
- $Q_i(t)$  backlog for node  $i$ , time  $t$
- Interference Constraints Respected
- Performance: Throughput and Delay

# Medium Access Algorithms

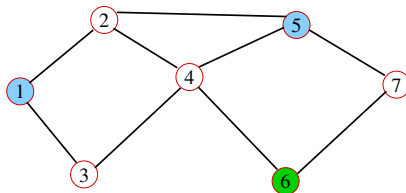
- **Capacity region:** Convex Hull of Feasible Schedules
- **Throughput-optimal algorithm:** Stabilises the queues, for any traffic intensity inside the capacity region (positive recurrence)
- **Centralised algorithms:** High Complexity, throughput optimal
  - e.g Max. Weight - selects a schedule to maximise  $\mathbf{Q}^T \mathbf{R}$
- **Distributed Algorithms:** (low complexity, not clear if optimal)
  - **Random Access Algorithms:** e.g. CSMA in 802.11
  - User listens to channel and transmits with some probability when idle

# Random Access Algorithm: Description

- When node  $i$  ends activity at time  $t$  it then backs-off w.p.  $g(Q_i(t))$ 
  - Back-off periods are unit exponential
- Back-off clock is **suspended** if any neighbour of  $i$  is active
  - Resumes when channel becomes idle
- At end of back-off period, node  $i$  begins transmission w.p.  $f(Q_i(t))$



# Random Access Algorithm: Example I



Active

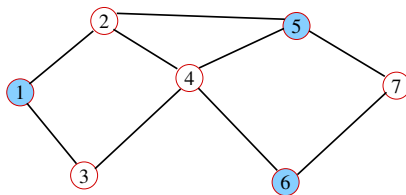


Suspended



Running Back-Off Clock

## Random Access Algorithm: Example II



Active

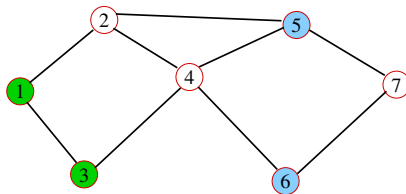


Suspended



Running Back-Off Clock

## Random Access Algorithm: Example III



Active

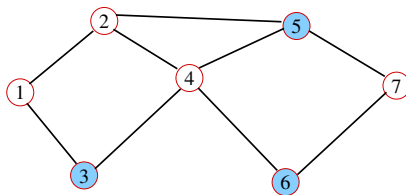


Suspended



Running Back-Off Clock

## Random Access Algorithm: Example IV



Active



Suspended



Running Back-Off Clock

# Throughput-Optimal Random Access

- $f = 1$  and deactivation function  $g = g(Q_i(t))$
- Rajogopalan and Shah, Shin (2009)
  - $g(x) = 1/\log(1+x)$  throughput optimal
  - Proof requires knowledge of maximum queue size in network
- Ghaderi and Srikant (2010)
  - $g(x) = (1+x)^{-1/h(1+x)}$
  - $h(x)$  slowly increasing function e.g.  $\log \log(e+x)$
- However excessive backlogs and delays are produced!!

# Remainder of Talk

- More aggressive/persistent schemes can improve delay
- No maximum stability guarantees
- Examine stability/instability via fluid limits
- Demonstrate instability in certain topologies

# Fluid Limits

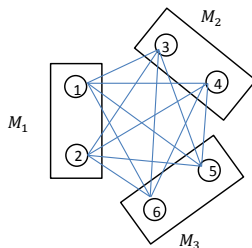
- $R$  scale parameter,  $(\frac{1}{R}\mathbf{Q}(Rt), \mathbf{U}(Rt))_{t>0}$
- Continuous queue paths obtained by linear interpolation have "skip free property" (see Meyn)
- $\|\mathbf{Q}(0)\| = R$ , take limits as  $R \rightarrow \infty$
- Sample paths of queues **Lipschitz continuous**, sequence is **tight**
- Fluid limits of queue process exist in  $C[0, \infty)$

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- Fast time scale with coarse granularity
- Fluid limits can be used to infer (in)stability

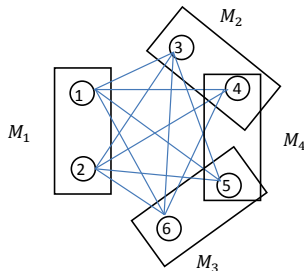


## Example: Complete Partite Network



- Backoff  $g(x) \doteq \frac{1}{(x+1)^\gamma}$
- Traffic -  $(0.4, 0.4, 0.3, 0.2, 0.2, 0.3) \rho$
- Stable  $\rho < 1, \gamma > 1$
- Nodes hold on to channel until backlog cleared
- Activity process does not leave component until all queues drained
- Random switching to alternate component
- Fluid limit reaches 0 in finite time  $\Rightarrow$  **Stability**

## Example: Network with Reduced Interference



- Removing one edge causes **instability**,  $\forall \gamma > 1$
- Traffic -  $(0.4, 0.4, 0.3, 0.2, 0.2, 0.3) \rho$
- $M_4$  visits occur i.o. use positive fraction of time
- $L(t)$  total queue length - explodes as  $\rho \uparrow 1^-$

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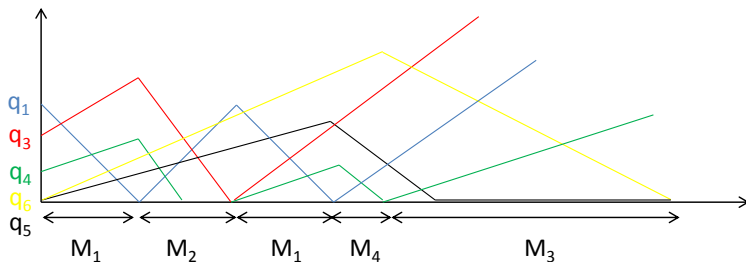
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- Switching out of  $M_4$  :  $\mathbb{P}\{M_1\} = 0$



# Sample Path Illustrating $M_1$ Cycles



# Fluid Limit Instability

$L(t)$  grows faster than any sub linear function of  $t$ ,

Theorem (Ghaderi, Borst, W.)

*For any  $m > 1$ , there exists a constant  $\rho^* < 1$  such that for all  $\rho \in (\rho^*, 1)$ ,*

$$\limsup_{T \rightarrow \infty} \mathbb{E} \left[ \frac{T}{L^m(T)} \right] = 0$$

*for any initial condition.*

## Two Cycle Argument

### Prop

$\exists C > 0$  such that  $\theta = 1 - (1 - \rho)C$ ,  $p = 1/12$ . Over cycle pairs  $D_k$ ,  $k = 1, 2, \dots$ ,

- (i)  $\Delta T_k \leq CL_k$ ;
- (ii)  $L(t) \geq \theta L_k$  for all  $t \in [T_k, T_{k+1}]$ ;
- (iii)  $\mathbb{P}(L_{k+1} - \theta L_k \geq \delta(\rho)\theta L_k | L_k) \geq p$ ,

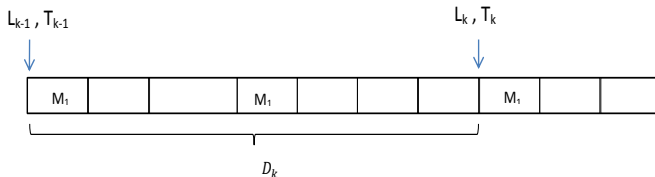
with  $\delta(\rho)$  a constant, depending on  $\rho$ , and

$\delta(\rho) \uparrow \delta = \frac{1}{\beta^{\max}(1+\beta^{\max})(1+\alpha-\min\{\kappa_3, \kappa_6\})}$ , as  $\rho \uparrow 1$ .

This proposition shows that for  $\rho$  close to 1 the load cannot significantly decrease over a pair of cycles and will increase by a substantial amount with non-zero probability.

# Proof Sketch I: Preliminaries

$M_1$  cycles,



Definitions and 2-cycle arguments,

$$T_{N_t-1} < t \leq T_{N_t} \quad (1)$$

$$L(t) \geq \theta L_{N_t-1}, \quad \theta \in (0, 1] \quad (2)$$

$$T_{N_t} \leq T_{N_t-1} + CL_{N_t-1} \quad (3)$$

$$L_{N_t-1} \leq 2T_{N_t-1} \quad (4)$$

# Proof Sketch II: Enough to bound $T_{N_t-1} L_{N_t-1}^{-m}$

$$\begin{aligned}
 \limsup_{t \rightarrow \infty} \mathbb{E} [t L^{-m}(t)] &\stackrel{(2)}{\leq} \limsup_{t \rightarrow \infty} \mathbb{E} [T_{N_t} \theta^{-m} L_{N_t-1}^{-m}] \\
 &\stackrel{(3)}{\leq} \theta^{-m} \left( \limsup_{t \rightarrow \infty} \mathbb{E} [T_{N_t-1} L_{N_t-1}^{-m}] + C \limsup_{t \rightarrow \infty} \mathbb{E} [L_{N_t-1}^{-m+1}] \right) \\
 &\stackrel{(4)}{\leq} \theta^{-m} (1 + 2C) \limsup_{t \rightarrow \infty} \mathbb{E} [T_{N_t-1} L_{N_t-1}^{-m}].
 \end{aligned}$$

## Proof Sketch III: $L_k^{-m}$ is a Geometric Supermartingale

$$\begin{aligned}\mathbb{E} [L_{k+1}^{-m} | \mathcal{F}_k] &\leq (1-p)(\theta L_k)^{-m} + p((\theta + \delta)L_k)^{-m} \\ &= \alpha_m L_k^{-m},\end{aligned}$$

where,

$$\alpha_m := (1-p)\theta^{-m} + p(\theta + \delta)^{-m}, p = \frac{1}{12}$$

by two cycle argument,  $\alpha_m < 1$ ,  $\rho$  close to 1

Hence a.s. convergence to 0 and .....

$$T_k \rightarrow \infty, \text{ a.s.}$$

because  $L_k \leq \rho T_k + 1 \leq T_k + 1$ , so stopping times well defined ...

# Proof Sketch IV: $T_k L_k^{-m}$ *almost* a Geom. Supermart.

$$\begin{aligned}\mathbb{E} [T_k L_k^{-m} | \mathcal{F}_{k-1}] &\leq (T_{k-1} + CL_{k-1}) \mathbb{E} [L_k^{-m} | \mathcal{F}_{k-1}] \\ &\leq (T_{k-1} + CL_{k-1}) \alpha_m L_{k-1}^{-m} \\ &= \alpha_m T_{k-1} L_{k-1}^{-m} + \alpha_m CL_{k-1}^{-m+1}.\end{aligned}$$

$\epsilon_k \doteq C\alpha_m L_k^{-m+1}$ , so that  $\sum_{k=1}^{\infty} \mathbb{E} [\epsilon_k] \leq C\alpha_m \sum_{k=1}^{\infty} \alpha_m^k < \infty$  and hence  $\lim_{k \rightarrow \infty} T_k L_k^{-m} = 0$  almost surely.

By induction,

$$\mathbb{E} [T_k L_k^{-m}] \leq \alpha^{k-1} C(1 + T^{(0)} + (k-1)\alpha),$$

for  $\rho$  sufficiently close to 1  $\rightarrow$  convergence in  $\mathcal{L}_1$

# Proof Sketch V: Stopped Sequence UI too

By Theorems of Rogers & Williams, (slightly adapted)

$T_k L_k^{-m}$  U.I sequence and then  $T_{N_t} L_{N_t}^{-m}$  U.I family

Whence, given  $\varepsilon > 0, \exists K_\varepsilon > 0$

$$\begin{aligned} \mathbb{E} [T_{N_t} L_{N_t}^{-m}] &\leq \sum_{k=1}^{\infty} \mathbb{E} \left[ T_k L_k^{-m} \mathbb{1}_{\{N_t = k\}} \mathbb{1}_{\{T_{N_t} L_{N_t}^{-m} \leq K_\varepsilon\}} \right] + \varepsilon \\ &\leq K_\varepsilon \mathbb{P} \{N_t \leq D\} + \sum_{k=D+1}^{\infty} Ck \alpha^{k-1} + \varepsilon, \quad D \in \mathbb{N} \end{aligned}$$

Fixing  $\varepsilon$  and  $D$ , we find that

$$\limsup_{t \rightarrow \infty} \mathbb{E} [T_{N_t} L_{N_t}^{-m}] \leq (D+1) \frac{\alpha^D}{1-\alpha} + \varepsilon$$

so that  $\limsup_{t \rightarrow \infty} \mathbb{E} [T_{N_t} L_{N_t}^{-m}] = 0$  for  $\rho$  sufficiently close to 1



- A theorem of Meyn (Ann. Prob. 1995), establishes a criteria for transience based on a fluid limit
- The idea is that the fluid limit may be used to determine a Lypuanov function
- Strict growth of the fluid limit is required
- Only growth in expectation our case
- Nevertheless . . . . .

# Instability Theorem

## Theorem (Meyn)

*Suppose that for a Markov chain  $\{X(n); n = 1, 2, \dots\}$  with state space  $S$ , there exist positive functions  $W(\cdot)$  and  $\Delta(\cdot)$  on  $S$ , and a positive constant  $c_0$ , such that*

$$\mathbb{E}[W(X(n+1)) | \mathcal{F}_n] \leq W(X(n)) - \Delta(X(n)),$$

*whenever  $X(n) \in S_{c_0} = \{x \in S : W(x) \leq c_0\}$ , with  $\mathcal{F}_n := \sigma(X(0), X(1), \dots, X(n))$ . Then for all  $x \in S$ ,*

$$\mathbb{P}_x \left\{ \sum_{n=0}^{\infty} \Delta(X(n)) < \infty \right\} \geq 1 - W(x)/c_0.$$

# Brief Sketch I: Definition of $\mathcal{W}$

Define,  $W(x) \doteq \mathbb{E} [\mathcal{W} | Q(0) = x]$ , where  $\mathcal{W}$  is defined as

$$\mathcal{W} := \sum_{n=0}^{\|Q(0)\|T} [1 + \|Q(0)\| + a\|Q(n)\|]^{-m}$$

for some positive constants  $a$  and  $T$  to be determined later and  $m > 1$ .

Note that,  $W(Q(0) = x, U(0) = u) = W(x)$ ,

Interpret this as an approximation to,

$$\|x\|^{m+1} W(x) \approx \mathbb{E}_{\hat{x}} \left\{ \int_0^T (1 + aL_{\hat{x}}(t))^{-m} dt \right\} = V(q_{\hat{x}}(t))$$

equality when  $\|x\| \rightarrow \infty$ , and  $\hat{x} = \frac{x}{\|x\|}$  is initial state

## Brief Sketch II: Drift Break Up

$\theta^1$  backward shift operator

$$\theta^1 \mathcal{W} - \mathcal{W} = A + B + C$$

where

$$\begin{aligned} A &\doteq -[1 + \|Q(0)\| + a\|Q(n)\|]^{-m}, \\ B &\doteq \sum_{n=1}^{\|Q(0)\| T} \{[1 + \|Q(1)\| + a\|Q(n)\|]^{-m} \\ &\quad - [1 + \|Q(0)\| + a\|Q(n)\|]^{-m}\}, \\ C &= \sum_{n=\|Q(0)\| T+1}^{\|Q(1)\| T} [1 + \|Q(1)\| + a\|Q(n)\|]^{-m}. \end{aligned}$$

The term 'A' provides the negative drift, see Meyn

## Brief Sketch III: Final Bound

After some manipulation . . . . .,

$$\limsup_{\|x\| \rightarrow \infty} \|x\|^m \mathbb{E}_{\hat{x}} \{ \theta^1 \mathcal{W} - \mathcal{W} \} \leq -(1+a)^{-m} + m \mathbb{E}_{\hat{x}} \left\{ \int_0^\infty (1 + a\tilde{L}(s))^{-m-1} ds \right\} + \mathbb{E}_{\hat{x}} \left\{ T(1 + a\tilde{L}(T))^{-m} \right\}$$

which bounds the drift in terms of integrals of the fluid process

Using earlier bounds,

$$m \mathbb{E}_{\hat{x}} \left\{ \int_0^\infty (1 + a\tilde{L}(s))^{-m-1} ds \right\} \leq mBC(a\theta)^{-m-1} \frac{1}{1-\alpha}$$

and take  $a$  large enough

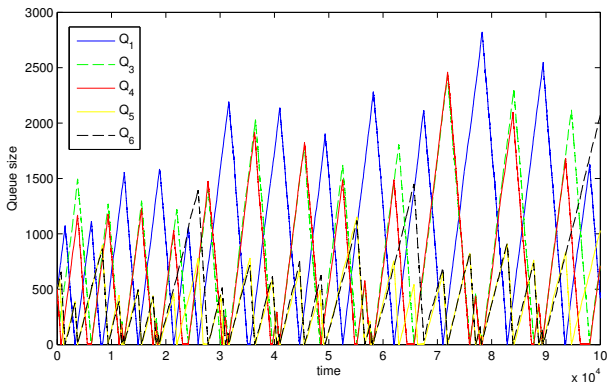
Also,

$$\mathbb{E}_{\hat{x}} \left\{ T[1 + a\tilde{L}(T)]^{-m} \right\} \leq a^{-m} \mathbb{E}_{\hat{x}} \left\{ TL^{-m}(T/b) \right\} \quad (5)$$

but RHS bounded as  $\limsup_{T \rightarrow \infty} \mathbb{E}_{\hat{x}} \{ TL^{-m}(T) \} = 0$ , for  $\rho \in (\rho^*, 1]$

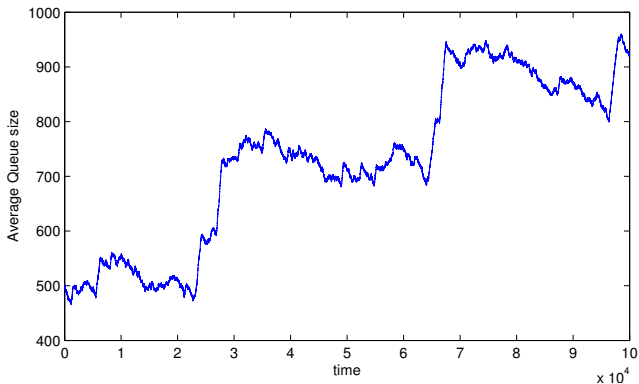
## Sample Paths for PreLimit Queues

- Initial queue size  $Q_\ell(0) = 500, \gamma = 2, \rho = 0.97$
- Backlogs cleared - queues exhibit qualitative behaviour of the fluid limit



## Sample Path for Mean Queue Length

- Saw Tooth Growth pattern in node average queue size
- Large Increments immediately after  $M_4$  visits (one per  $M_1$  cycle)
- Slight downward trend otherwise



# Conclusions








- Aggressive activation functions have piecewise linear (oscillatory) limits
- Instability for  $\gamma > 1$  in a broad class of interference networks
- Same applies for random capture algorithms



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- Aggressive activation functions have piecewise linear (oscillatory) limits
- Instability for  $\gamma > 1$  in a broad class of interference networks
- Same applies for random capture algorithms
- Under Investigation
- Instability for any  $\gamma > 1/K$  for network sizes of order  $K$
- Sharpen fluid limit instability/stability criteria

# Thanks!

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