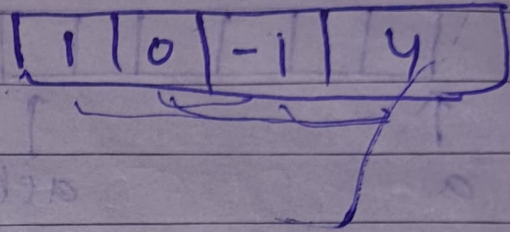


Arrays

→ Generating subarrays → contiguous



```
for (int i = 0; i < n; i++)
```

```
{
  for (int j = i; j < n; j++)
```

```
{
  for (int k = i; k <= j; k++)
```

```
{
```

```
    cout << a[k];
```

```
}
```

```
    cout << "\n";
```

```
}
```

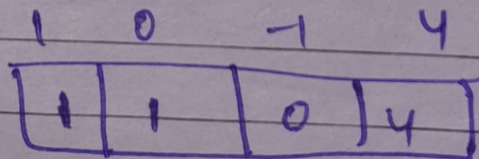
```
}
```

One line

→ Maximum subarrays sum

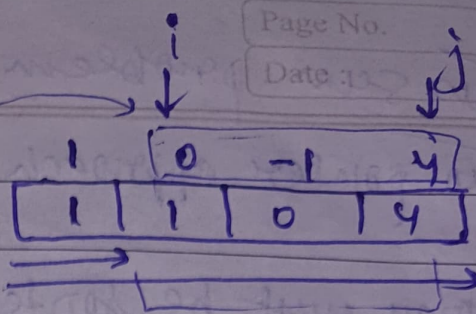
①st way brute force :→ add in above loop \nearrow sum, temp.
 $O(n^3)$

②nd way cumulative sum :→ make a cumulative sum array.
 $O(n^2)$



$$CS[i] = CS[i-1] + a[i]$$

also for



sum from i to j

$$\text{Sum} = \text{cs}[j] - \text{cs}[i-1]$$

③rd way: Kadane's Algo \Rightarrow

1st maintain 2 variable curr and max Till now now initialize, $\text{curr} = 0$
 $\text{Max Till now} = 0$

2nd Iterate from $i=0$ to $i=n-1$, and do

$\text{curr} = \text{curr} + a[i]$
 $\text{Max Till now} = \max[\text{Max Till now}, \text{curr}]$
 but if $\text{curr} > 0$ then only

otherwise,

③rd if $\text{curr} < 0$, then do $\text{curr} = 0$

$\text{max Till now} = \max(\text{max Till now}, a[i])$ use element ko le hi mat, phir se array shuru kar

④th when loop ends, output max Till now .

Two sum problem!!

→ Two pointer approach

① array must be sorted.

② if $(a[i] + a[j] > \text{sum})$ $\{1, 3, 5, 7, 10, 11, 12, 13\}$
 $j--$; → for its kam pane h method

else if $(a[i] + a[j] < \text{sum})$
 $i++$; → for its badhane h method.