

# New stochastic sketching methods for Big Data Ridge Regression

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## **Abstract**

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# 1. Randomized Newton Method

## 1.1 Algorithm

## 1.2 Convergence rate (draft)

### 1.2.1 General case

$$Z = AI_C^T(I_C AI_C^T)^{-1}I_C A$$

$$\rho = 1 - \lambda_{\min}(A^{-\frac{1}{2}}E[Z]A^{-\frac{1}{2}}).$$

$$A^{-\frac{1}{2}}E[Z]A^{-\frac{1}{2}} = \sum_i p_i A^{\frac{1}{2}} I_{C_i}^T (I_{C_i} A I_{C_i}^T)^{-1} I_{C_i} A^{\frac{1}{2}}$$

for any  $i \in \{1, \dots, n\}$ ,  $A^{\frac{1}{2}} I_{C_i}^T (I_{C_i} A I_{C_i}^T)^{-1} I_{C_i} A^{\frac{1}{2}}$  is a projection matrix and then its eigenvalues are a nonempty subset of  $\{0, 1\}$ .

Since  $\lambda_{\max}$  is convex, we obtain that :

$$0 \leq \lambda_{\min}(A^{-\frac{1}{2}}E[Z]A^{-\frac{1}{2}}) \leq \lambda_{\max}(A^{-\frac{1}{2}}E[Z]A^{-\frac{1}{2}}) \leq \sum_i p_i \lambda_{\max}(A^{\frac{1}{2}} I_{C_i}^T (I_{C_i} A I_{C_i}^T)^{-1} I_{C_i} A^{\frac{1}{2}}) \leq 1.$$

$$C = (I_{C_1}^T, \dots, I_{C_r}^T).$$

$$A^{-\frac{1}{2}}E[Z]A^{-\frac{1}{2}} = (A^{\frac{1}{2}}CD)(DC^T A^{\frac{1}{2}}) \text{ where } D = \text{diag}(\sqrt{p_1}(I_{C_1} A I_{C_1}^T)^{-\frac{1}{2}}, \dots, \sqrt{p_r}(I_{C_r} A I_{C_r}^T)^{-\frac{1}{2}})$$

#### Proposition 1.2.1

$$\lambda_{\min}(A^{-\frac{1}{2}}E[Z]A^{-\frac{1}{2}}) \geq \binom{n-1}{s-1} \frac{\lambda_{\min}(A)}{\lambda_{\max}(A)} \min_i p_i$$

**Proof :**

$$\lambda_{\min}(A^{-\frac{1}{2}}E[Z]A^{-\frac{1}{2}}) \geq \lambda_{\min}(C^T A C) \lambda_{\min}(D^2)$$

$$\lambda_{\min}(D^2) = \min_i \frac{p_i}{\lambda_{\max}(I_{C_i} A I_{C_i}^T)} \geq \min_i \frac{p_i}{\lambda_{\max}(I_{C_i}^T I_{C_i}) \lambda_{\max}(A)} \geq \min_i \frac{p_i}{\lambda_{\max}(A)}, \text{ since for any } i \in \{1, \dots, n\}, \text{ for any } x \text{ in } \mathbb{R}^n \langle I_{C_i}^T I_{C_i} x | x \rangle = \|I_{C_i} x\|^2 \leq \|x\|^2 \text{ and then } \lambda_{\max}(I_{C_i}^T I_{C_i}) \leq 1.$$

$$\text{Therefore, } \lambda_{\min}(A^{-\frac{1}{2}}E[Z]A^{-\frac{1}{2}}) \geq \min_i p_i \frac{\lambda_{\min}(C^T A C)}{\lambda_{\max}(A)} = \min_i p_i \frac{\lambda_{\min}(A) \lambda_{\min}(C C^T)}{\lambda_{\max}(A)}.$$

$CC^T = \sum_i I_{C_i}^T I_{C_i} = \binom{n-1}{s-1} I_n$  and then we obtain that :

$$\lambda_{\min}(A^{-\frac{1}{2}} E[Z] A^{-\frac{1}{2}}) \geq \binom{n-1}{s-1} \frac{\lambda_{\min}(A)}{\lambda_{\max}(A)} \min_i p_i$$

### 1.2.2 Uniform case

For any  $i$ ,  $p_i = \frac{1}{\binom{n}{s}}$  is the uniform probability of choosing  $s$  rows uniformly on  $\{1, \dots, n\}$ , knowing that  $s$  is the sketchsize.

#### Proposition 1.2.2

$$\lambda_{\min}(A^{-\frac{1}{2}} E[Z] A^{-\frac{1}{2}}) \geq \frac{s}{n} \frac{\lambda_{\min}(A)}{\lambda_{\max}(A)}$$

## 2. Hadamard Sketches

### 2.1 Algorithm

### 2.2 Convergence rate (draft)

$$Z = AS^T(SAS^T)^{-1}SA$$

$$\rho = 1 - \lambda_{\min}(A^{-\frac{1}{2}}E[Z]A^{-\frac{1}{2}})$$

$$S_i = I_{C_i}H.$$

#### Proposition 2.2.1

$$\lambda_{\min}(A^{-\frac{1}{2}}E[Z]A^{-\frac{1}{2}}) \geq \frac{s}{n} \frac{\lambda_{\min}(A)}{\lambda_{\max}(A)}$$

**Proof :**

$$\tilde{A} = \frac{1}{n}HAH^T = \frac{H}{\sqrt{n}}A\frac{H^T}{\sqrt{n}}.$$

$$A^{-\frac{1}{2}}E[Z]A^{-\frac{1}{2}} = E[A^{\frac{1}{2}}S^T(SAS^T)^{-1}SA^{\frac{1}{2}}] = \sum_i p_i A^{\frac{1}{2}}H^T I_{C_i}^T (I_{C_i}HAH^T I_{C_i}^T)^{-1} I_{C_i}HA^{\frac{1}{2}}$$

$$= A^{\frac{1}{2}}H^T \frac{1}{n} E[I_C^T (I_C \tilde{A} I_C^T)^{-1} I_C] HA^{\frac{1}{2}} = H^{-1} \tilde{A}^{\frac{1}{2}} E[I_C^T (I_C \tilde{A} I_C^T)^{-1} I_C] \tilde{A}^{\frac{1}{2}} \left(\frac{H^T}{n}\right)^{-1} = H^{-1} \tilde{A}^{\frac{1}{2}} E[I_C^T (I_C \tilde{A} I_C^T)^{-1} I_C] \tilde{A}^{\frac{1}{2}} H.$$

Hence :

$$\rho = 1 - \lambda_{\min}(A^{-\frac{1}{2}}E[Z]A^{-\frac{1}{2}}) = 1 - \lambda_{\min}(\tilde{A}^{\frac{1}{2}}E[I_C^T (I_C \tilde{A} I_C^T)^{-1} I_C] \tilde{A}^{\frac{1}{2}})$$

We recognize the convergence rate in the Randomized Newton Method and then, denoting by  $\rho_{Newton}(M)$  the convergence rate of the Newton method associated with the definite positive matrix  $M$ , we obtain that :

$$\rho = 1 - \lambda_{\min}(\tilde{A}^{\frac{1}{2}}E[I_C^T (I_C \tilde{A} I_C^T)^{-1} I_C] \tilde{A}^{\frac{1}{2}}) = 1 - (1 - \rho_{Newton}(\tilde{A})) = \rho_{Newton}(\tilde{A}) = \rho_{Newton}(A), \text{ since } A \text{ and } \tilde{A} \text{ have the same eigenvalues.}$$

## 3. *Count-min Sketches*

### 3.1 Algorithm

### 3.2 Convergence rate

$S$  is constructed as follows :

For every  $i \in \{1, \dots, n\}$ ,  $l$  is chosen uniformly on  $\{1, \dots, n\}$  and  $\epsilon$  uniformly on  $\{-1, 1\}$ , then  $S$  is updated in his  $l^{th}$  row as :

$S(l, :) := S(l, :) + \epsilon e_i^T$ , where  $e_i^T$  is the  $i^{th}$  coloumn of the identity matrix.

## 4. *Conclusion*



## *References*

- [1] ROBERT GOWER AND PETER RICHTARIK, Randomized iterative methods for linear systems, SIAM, (2015).