

# Theoretical limitations of multi-layer Transformer

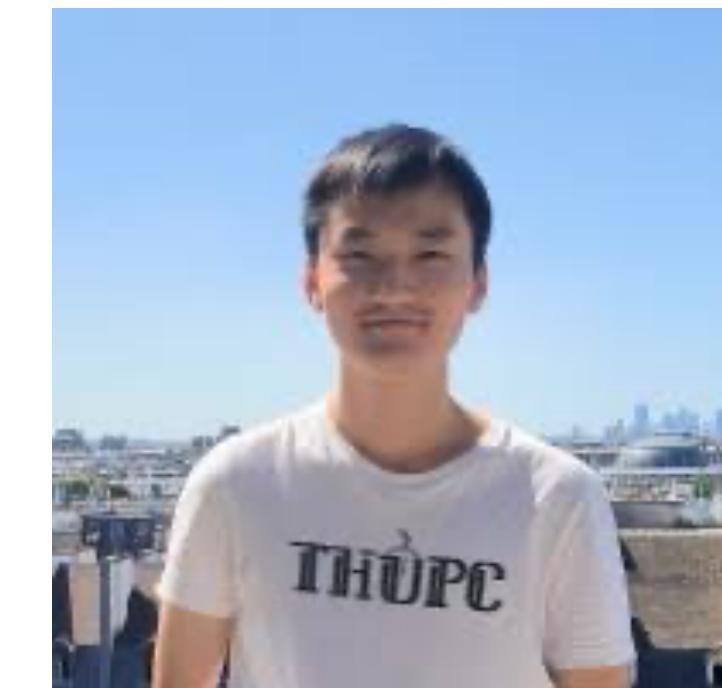
Warning: Theory paper, not meant for practicality.



**Lijie Chen**  
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**Binghui Peng**  
Stanford

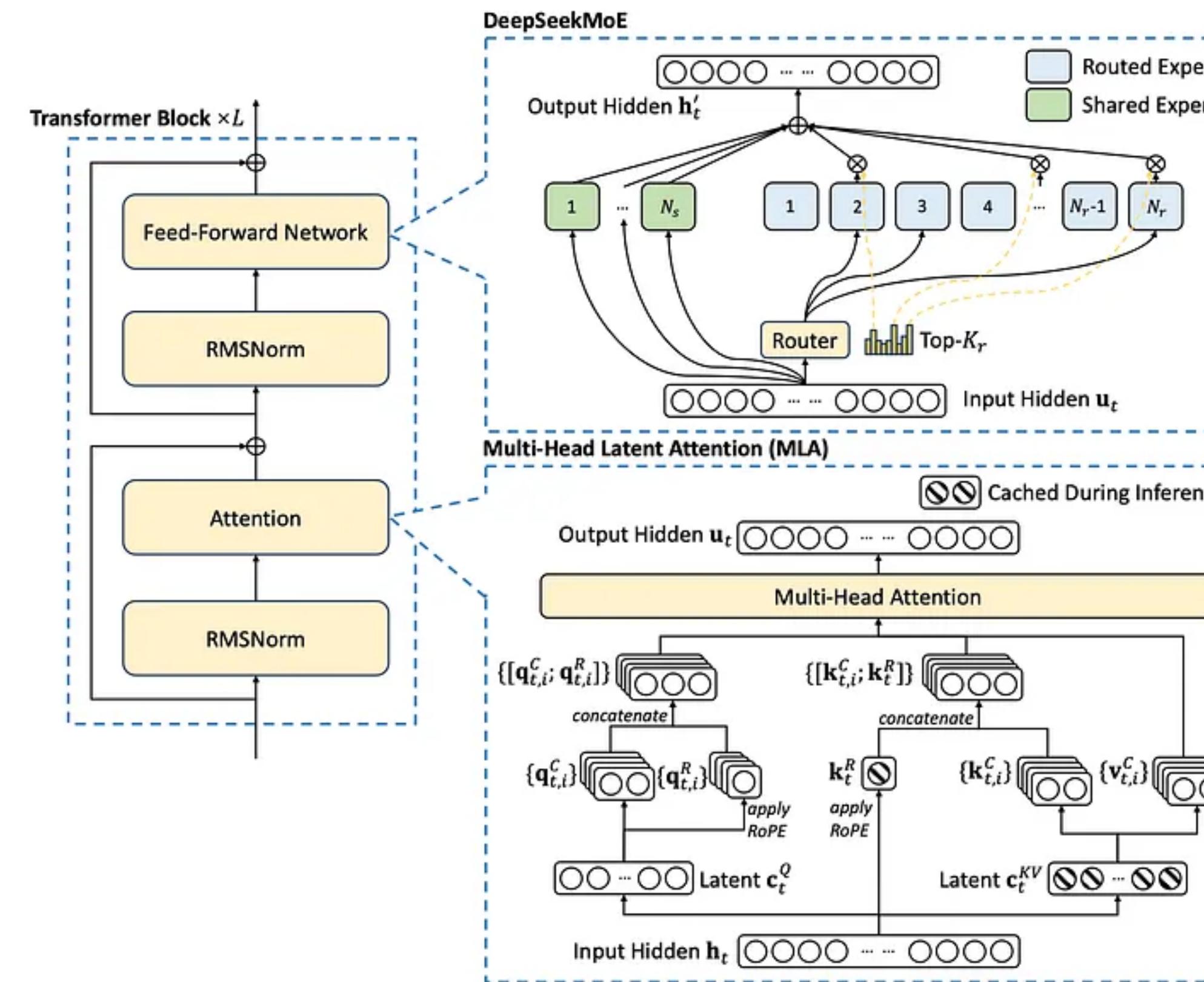


**Hongxun Wu**  
UC Berkeley

# Our Problem

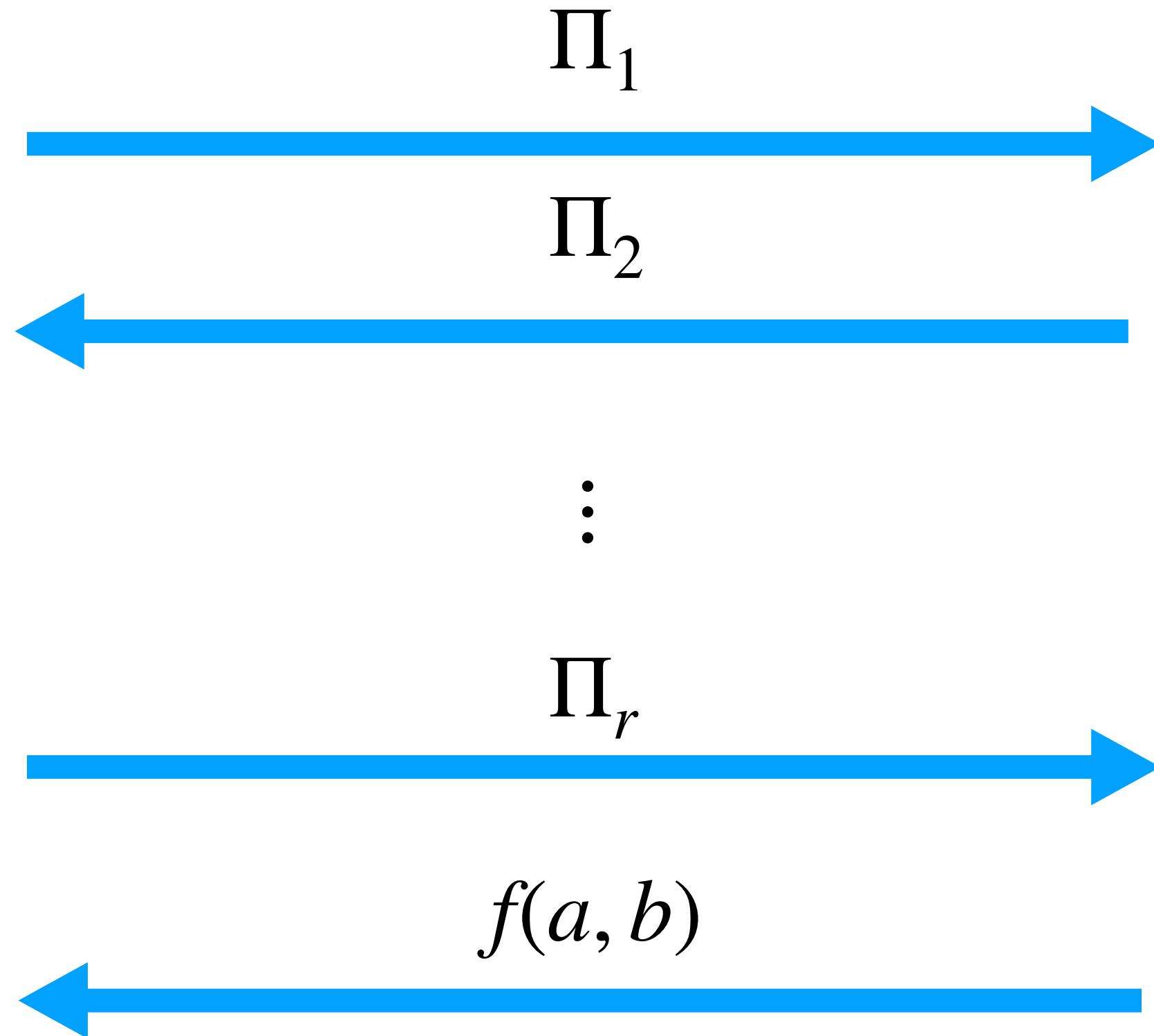
- What tasks can **not** be solved by Transformers without CoT?
- Existing Answer:
  - Transformers are constant-depth threshold circuits  $\text{TC}^0$ . ([Merrill et al.'24](#))
  - It cannot solve inherently sequential tasks.
- In which way is it unsatisfactory (to theoreticians):
  - We do not understand  $\text{TC}^0$ ! ([Complexity Theory's Waterloo](#))
  - Any constant-layer architecture is in  $\text{TC}^0$ . Didn't say much about transformers.

# Transformers



- Many variations!
- MHA / MQA / MLA / ... , tokenization, positional encoding, mixture-of-expert, pre-norm/post-norm/hybrid-norm, activation function, residual connection, gated components...

# Communication Game



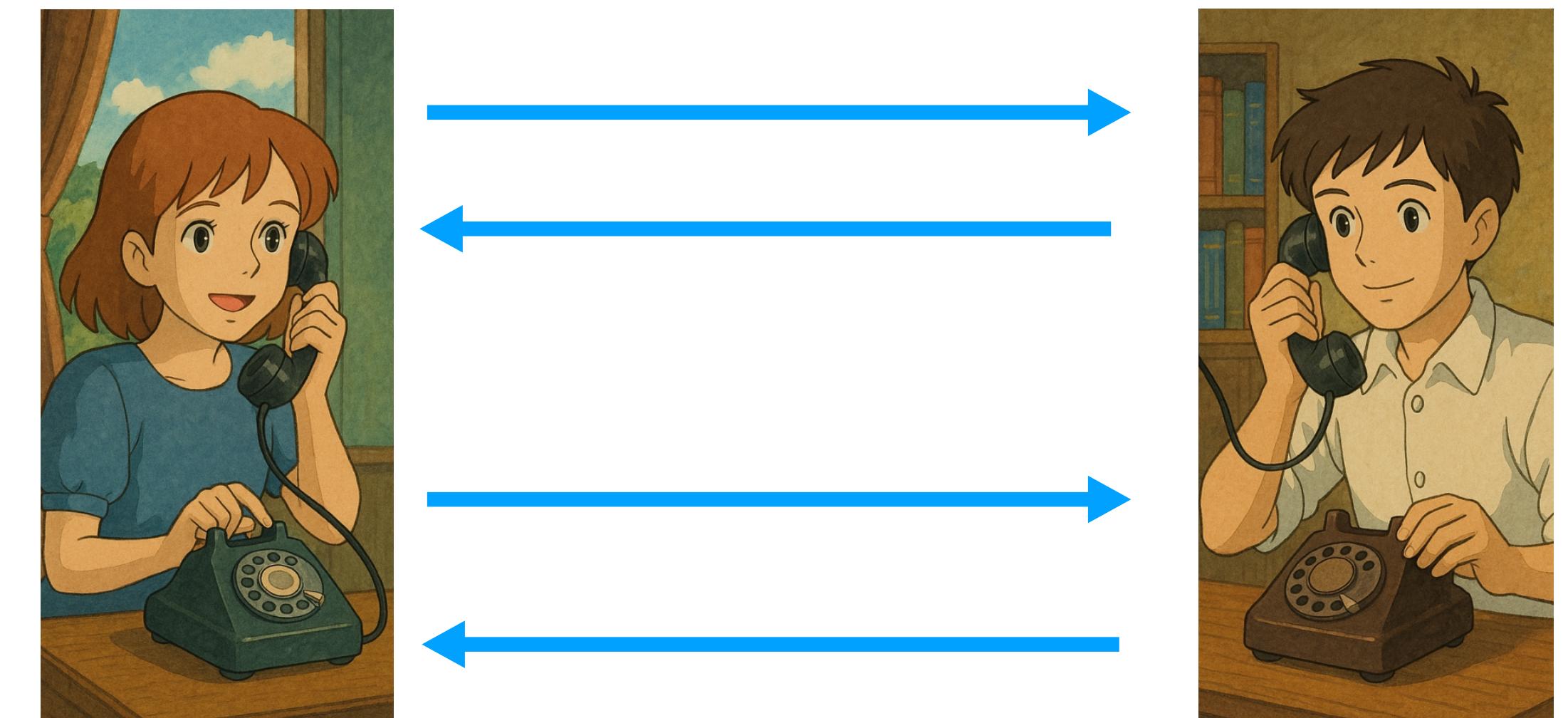
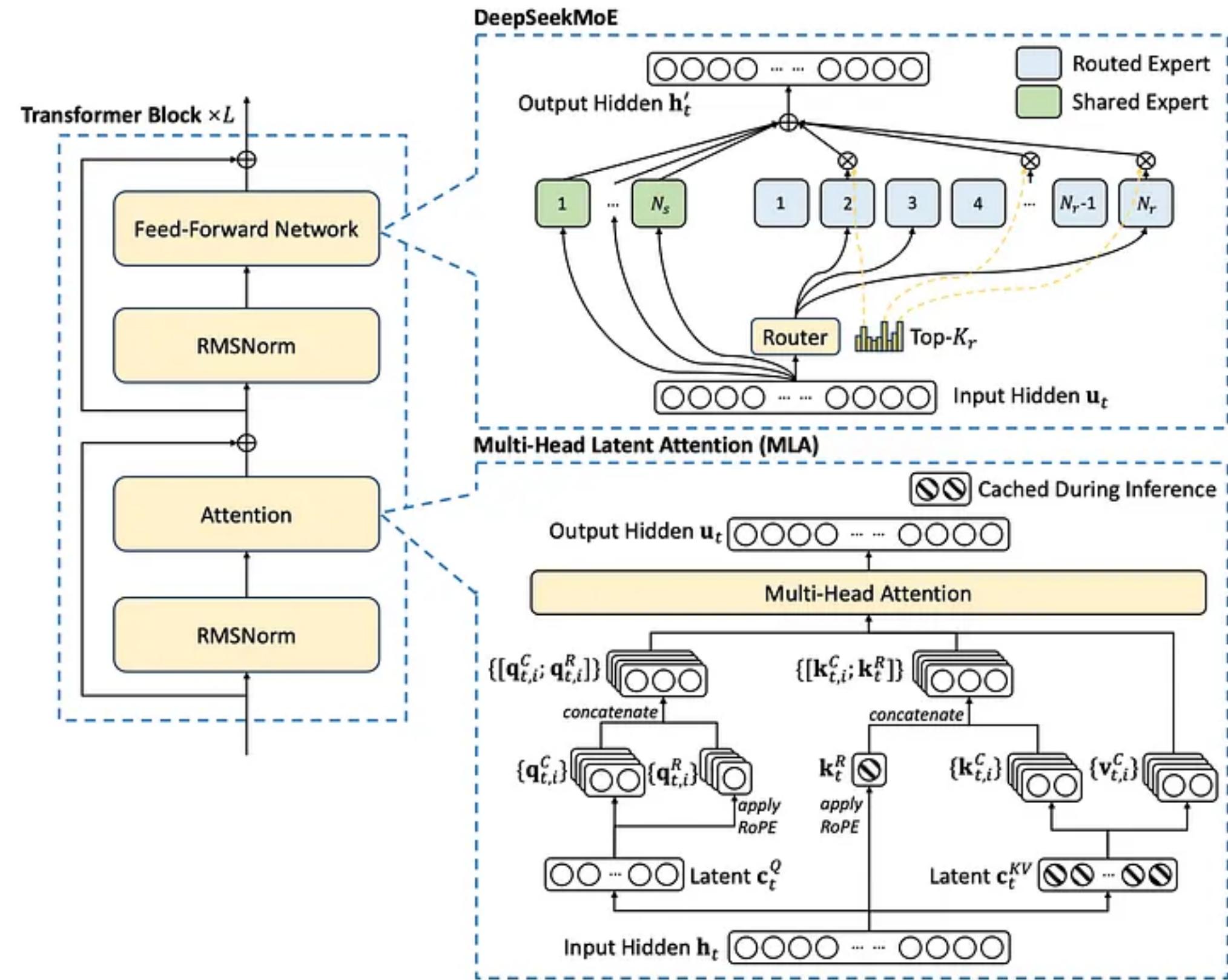
$$a \in \{0,1\}^n$$

$$\text{minimize } |\Pi_1| + |\Pi_2| + \dots + |\Pi_r|$$

$$b \in \{0,1\}^n$$

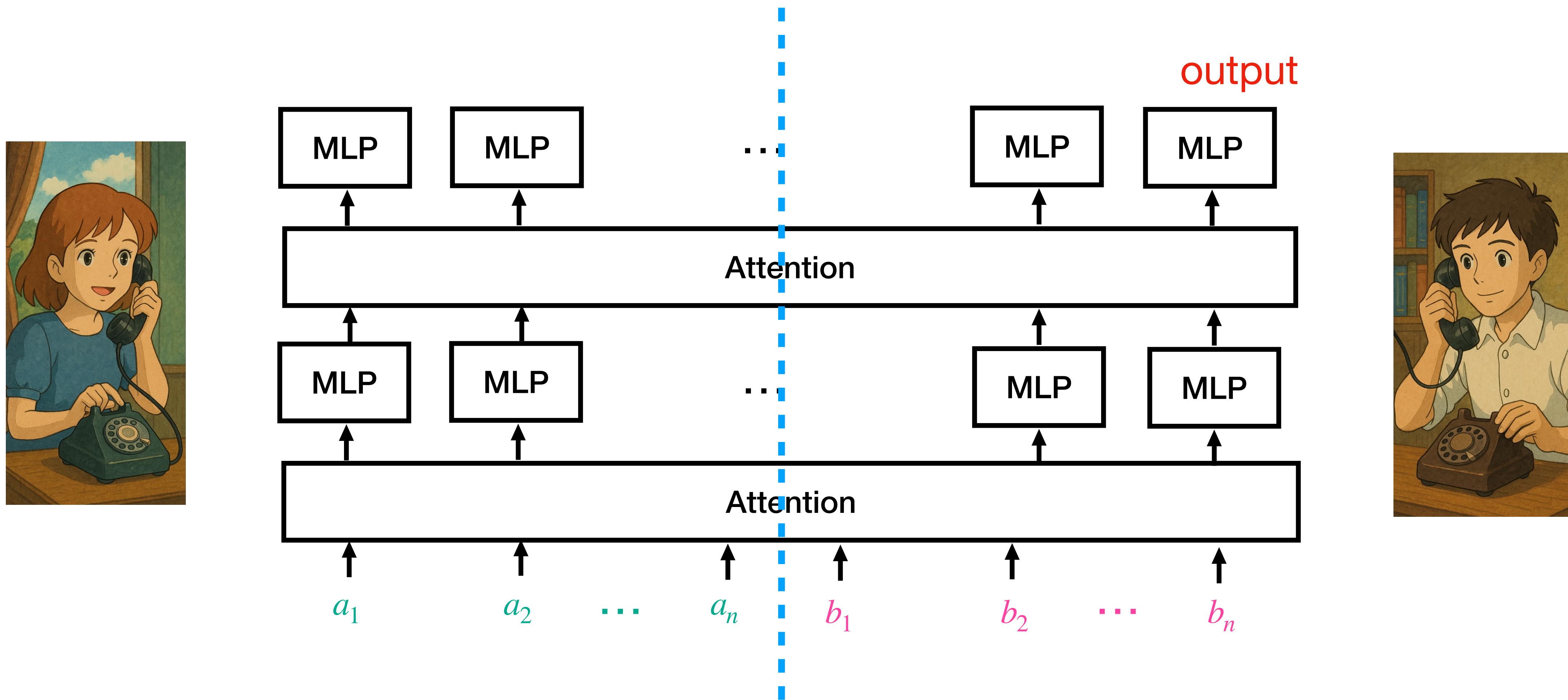
# Theme

- Understand the limitations of a transformer by understand its **inside communication!**



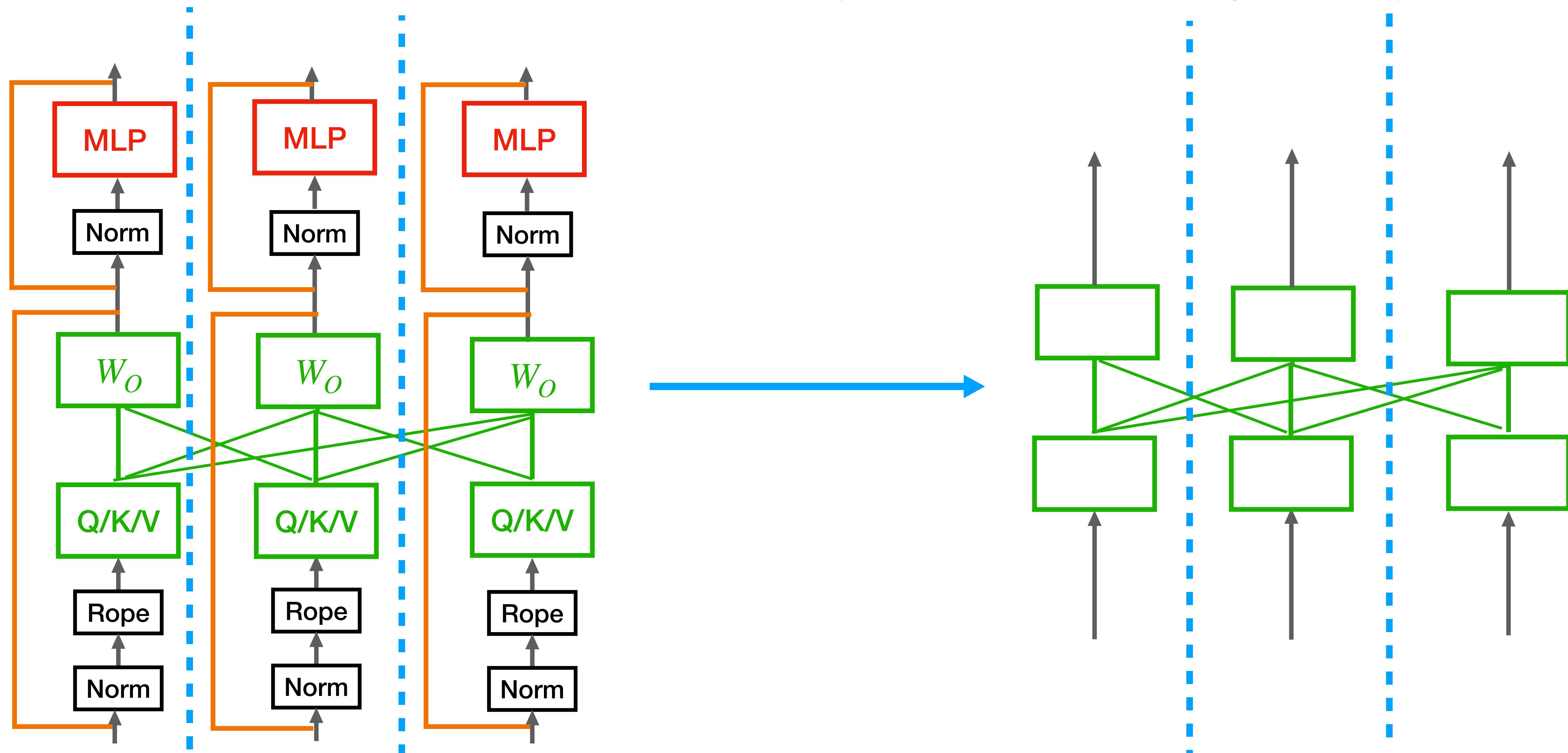
# Theme

- Understand the limitations of a transformer by understand its **inside communication!**



# Theme

- When we focus on **communicaiton**, many **architectural choice** goes away.



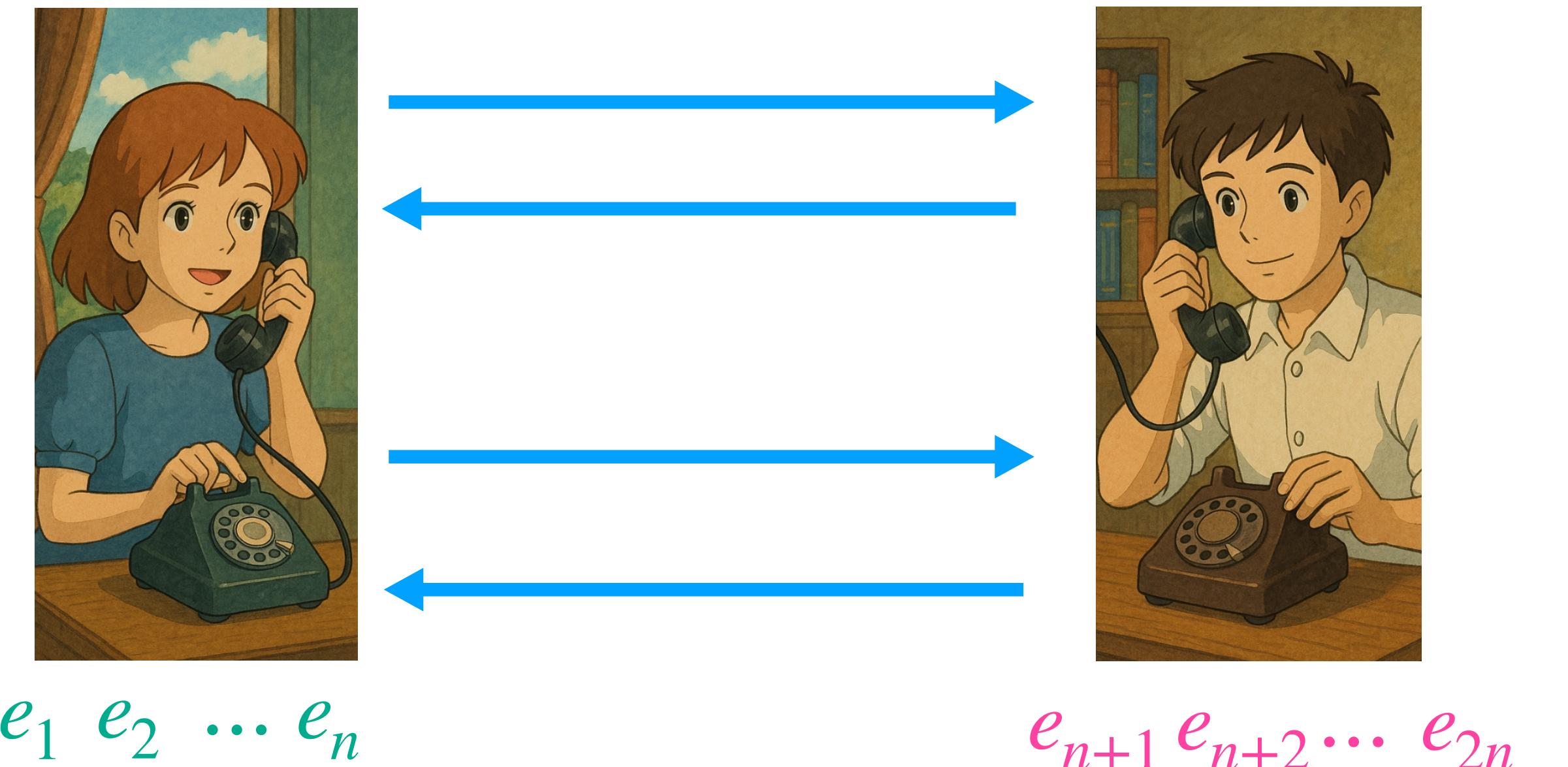
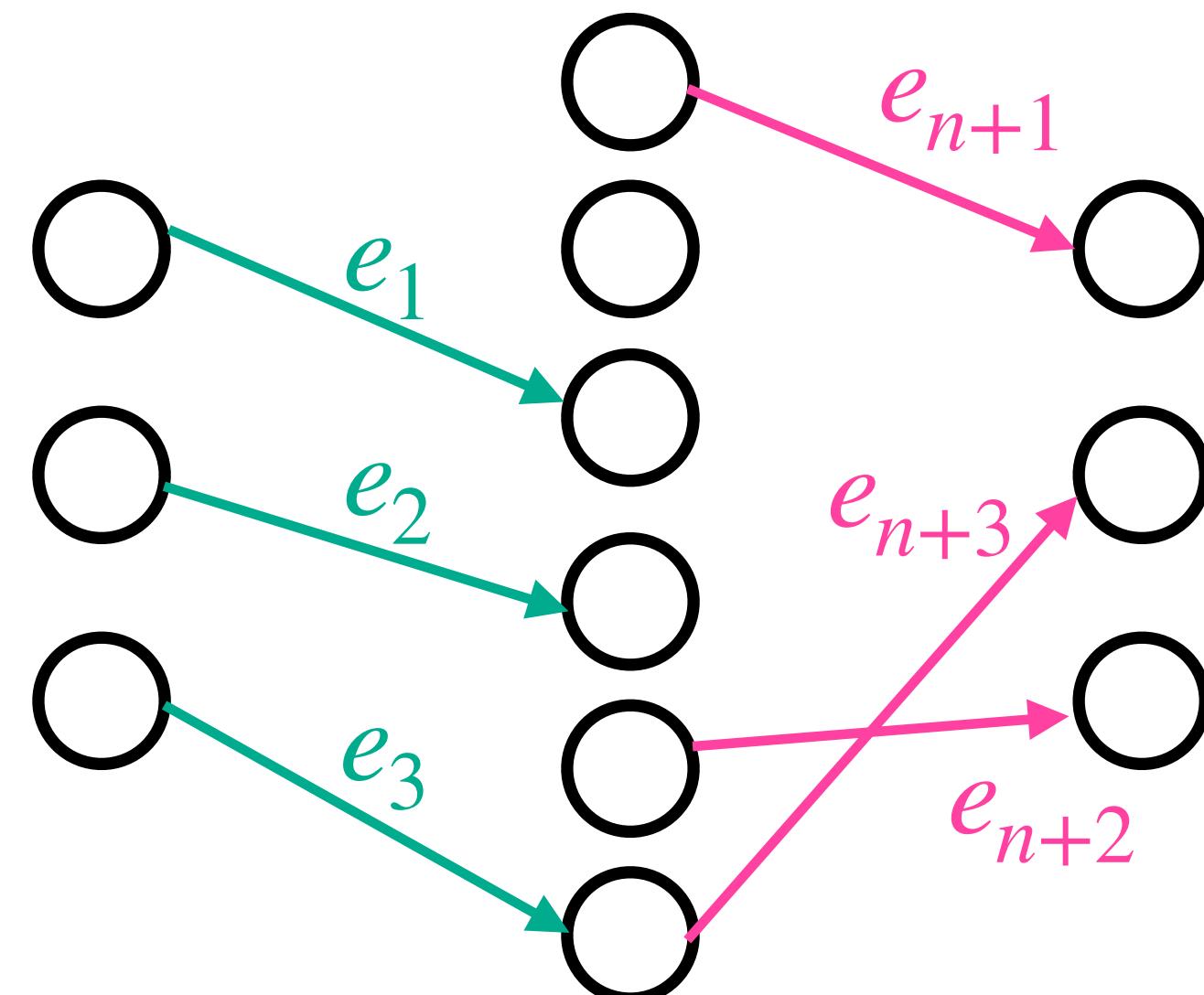
# Warm up: One Layer - Example 1

$n$ : prompt length  
 $d$ : hidden dimension  
 $H=1$ : attention head  
 $p$ : precision (bits)

Theorem (Mixture of Parrot by Jelassi et al.'25)

A single-layer transformer cannot find a length-2 path in a graph unless  $dp = \Omega(n)$ .

Problem (Length-2 Path)



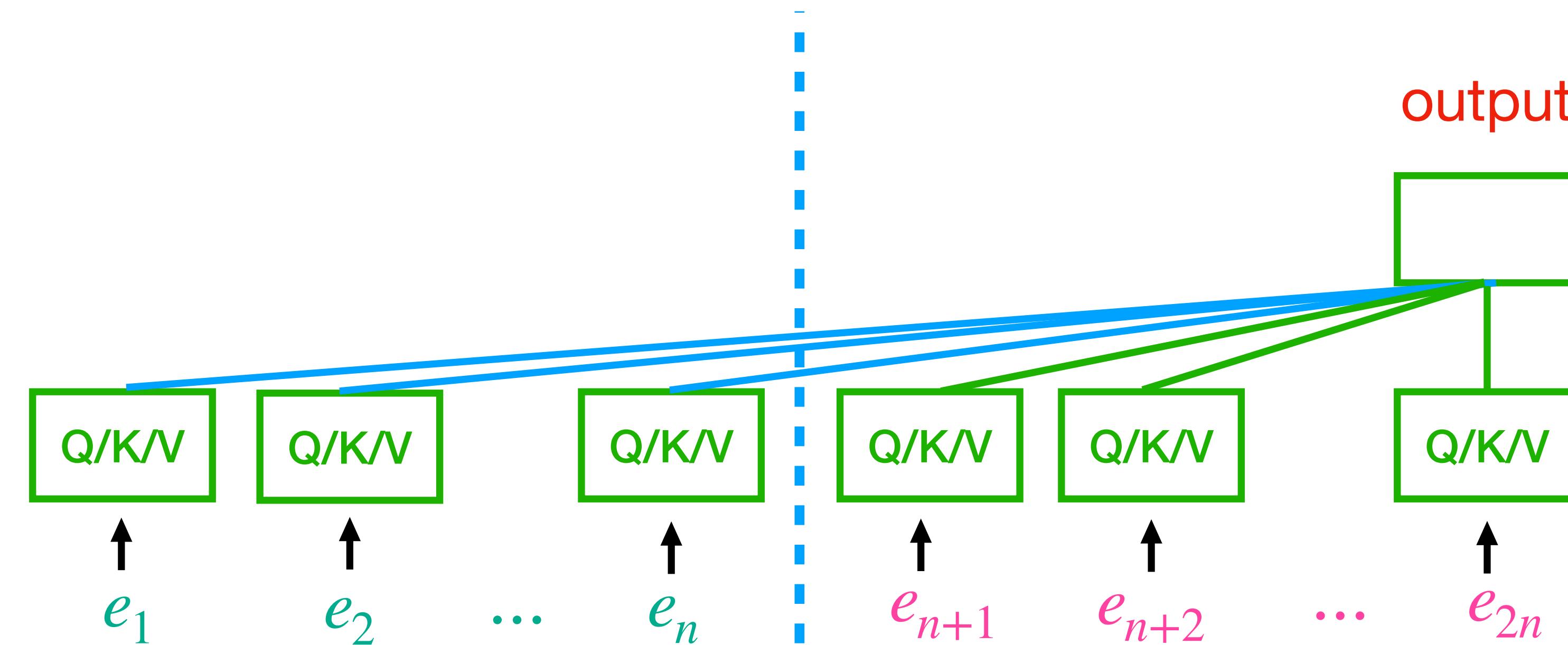
Theorem: To decide whether there is a Length-2 Path, they must exchange at least  $\Omega(n)$  bits of information!

# Warm up: One Layer - Example 1

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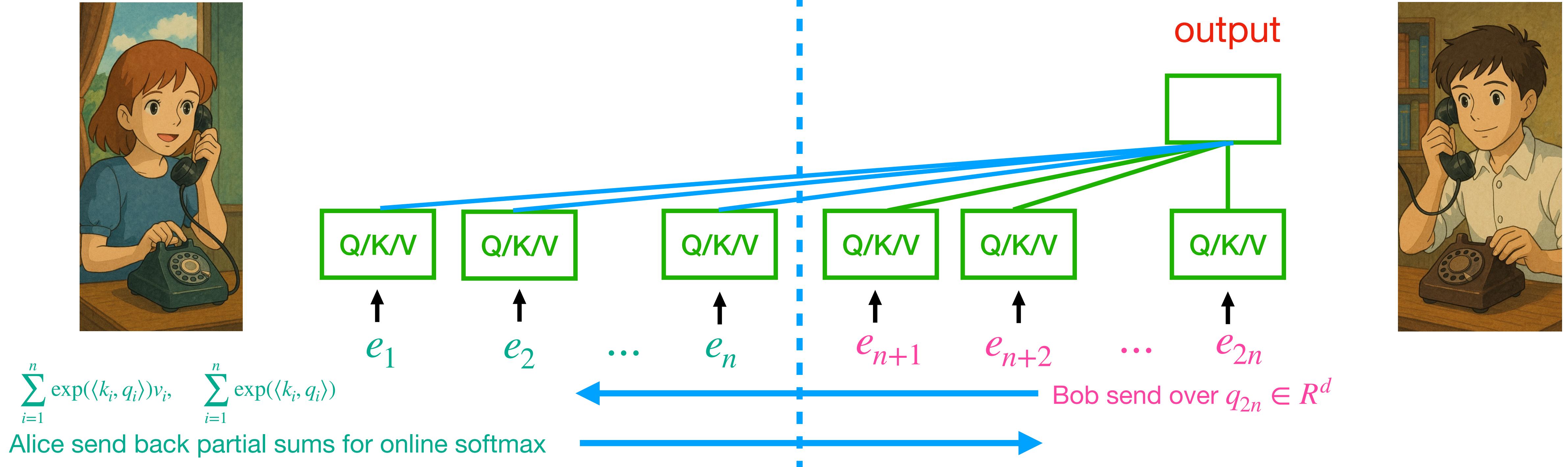
Key: For 1 layer, only the last queries matter!

# Warm up: One Layer - Example 1

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A single-layer transformer cannot find a length-2 path in a graph unless  $dp = \Omega(n)$ .



# Warm up: One Layer - Example 2

$n$ : prompt length  
 $d$ : hidden dimension  
 $H=1$ : attention head  
 $p$ : precision (bits)

Theorem (Sanford et al.'24)

A single-layer transformer cannot solve the **induction head task** unless  $dp = \Omega(n)$ .

Problem (**Induction-head**)

a b c d c a \* b d a b \*

Find the character before the last occurrence of the same token.

Problem (**Indexing**)

$s = [a, b, c, d, c, a, b, d, a, b]$   
 $s[i] = ?$

Alice has  $s$  while Bob has  $i$

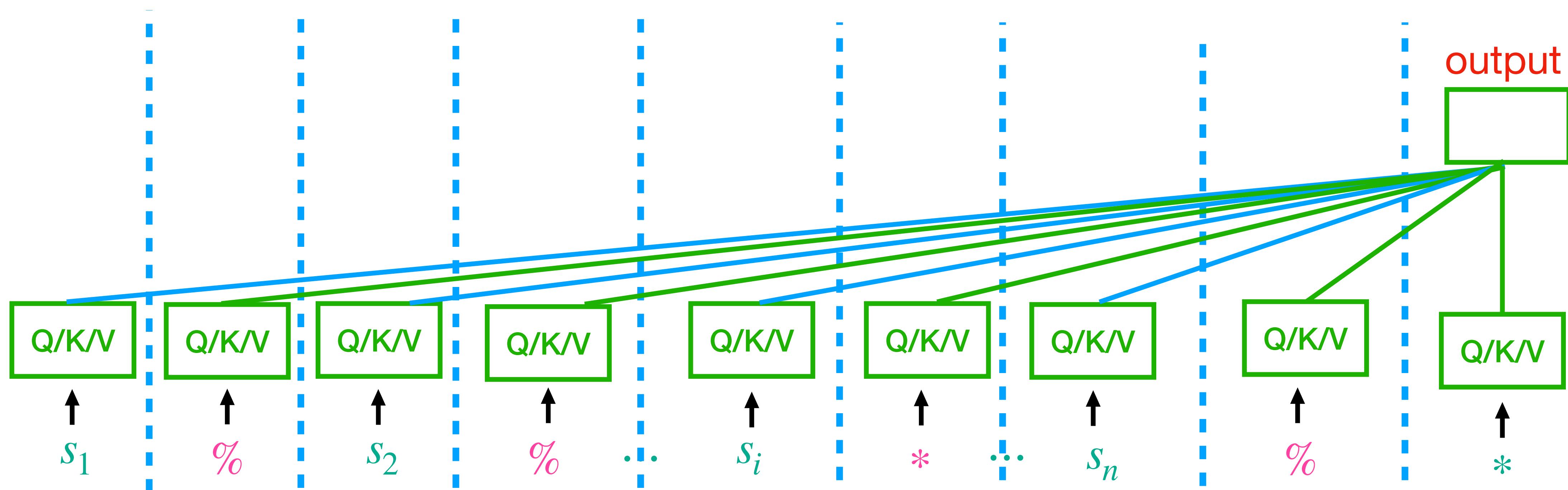
Theorem: One-way (Alice  $\rightarrow$  Bob)  
communication complexity of indexing =  $\Omega(n)$

# Warm up: One Layer - Example 2

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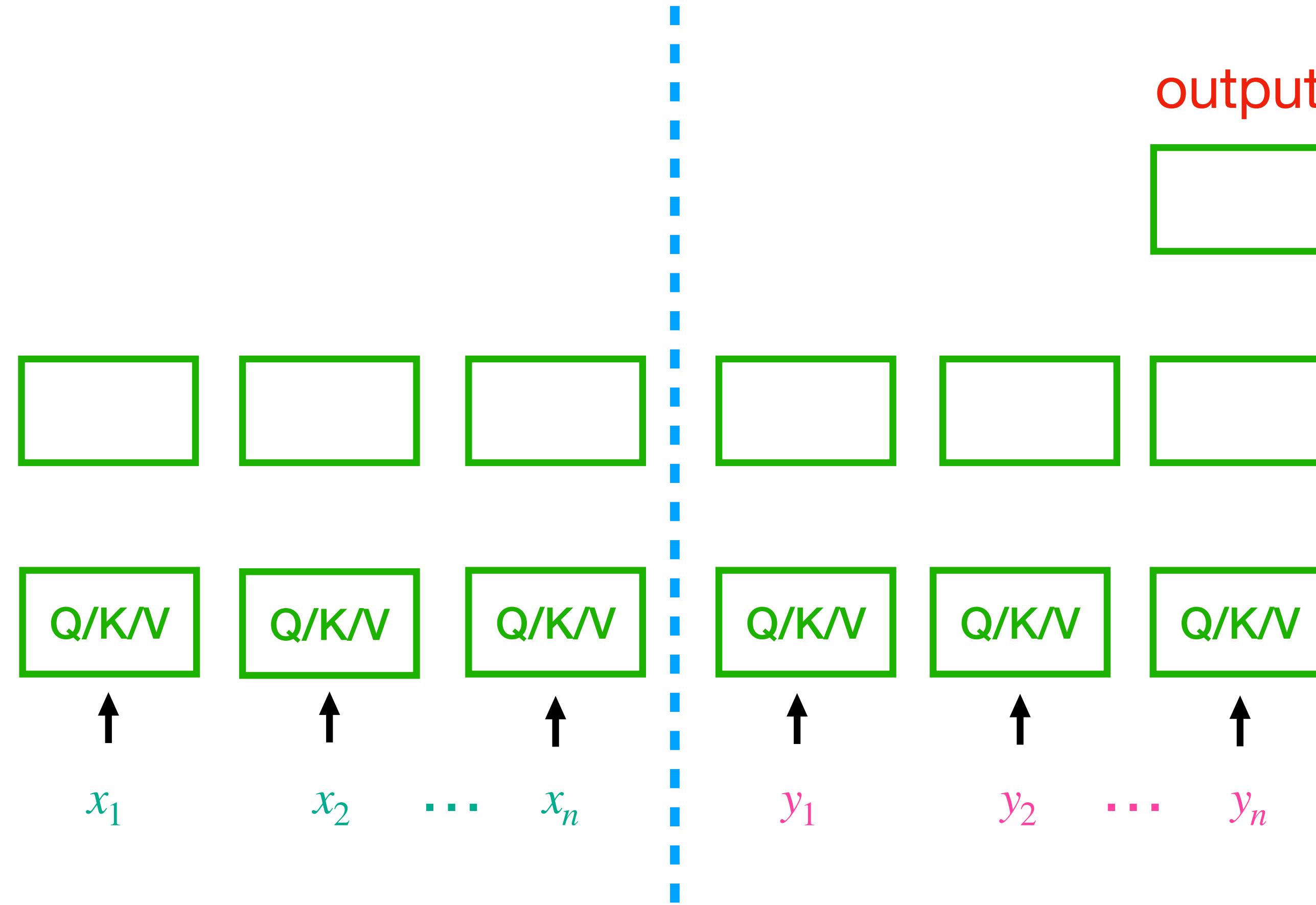


Alice has all the **odd positions**, while Bob has all the **even positions**.

# Why did they stop at 1-layer?

- Constant-layer encoder can simulate constant-depth threshold circuits  $\text{TC}^0$ .
- We do not understand  $\text{TC}^0$  or even  $\text{THR} \circ \text{THR}$ ! ([Complexity Theory's Waterloo](#))
- But **why** not even 2-layer decoders?

# Why did they stop at 1-layer?

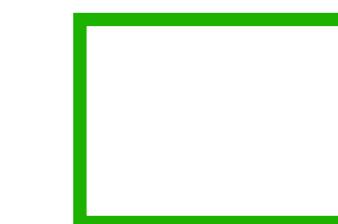
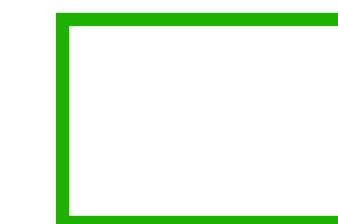
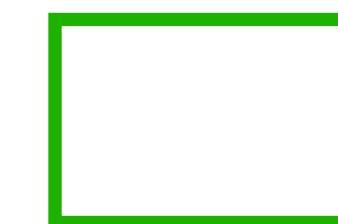


Now we need to first compute **every hidden state** on the second level.  
=> **Every query matters now!**

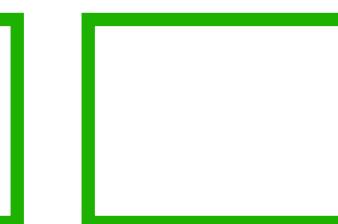
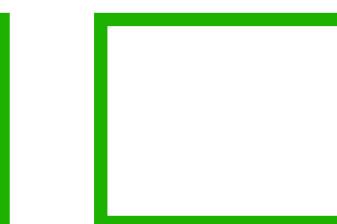
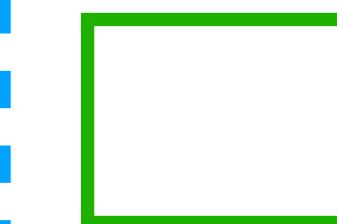
# Why did they stop at 1-layer?



Either Alice sends over  
 $k_1, v_1, k_2, v_2, \dots, k_m, v_m \in \mathbb{R}^d$



output



$x_1$

$x_2$

$\dots$

$x_m$

$y_1$

$y_2$

$\dots$

$y_n$

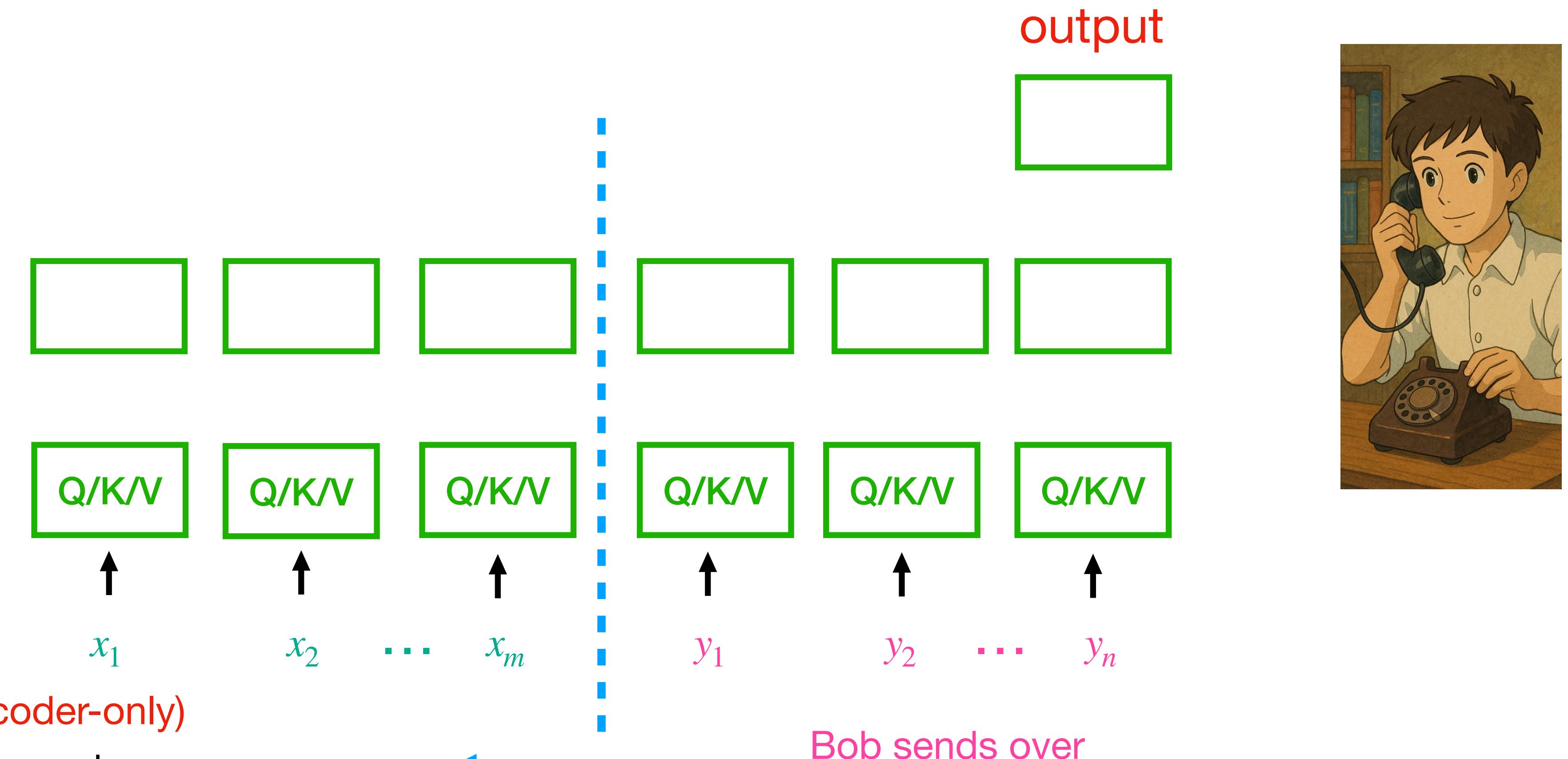


They have to send  $\Omega(dm) / \Omega(dn)$  information,  
which is more than enough to reveal their input.

Or Bob sends over  
 $q_{m+1}, q_{m+2}, \dots, q_{m+n} \in \mathbb{R}^d$



# Autoregressive (“forgetful”) communication



Our Simple Observation (decoder-only)

Communication game fails to capture:

After Bob sends **his queries**, and Alice replies with her **partial sums**, Alice “forgets” these **queries**!

Bob sends over  
 $q_{m+1}, q_{m+2}, \dots, q_{m+n} \in \mathbb{R}^d$

# Autoregressive (“forgetful”) communication

- $N$  players, where the player  $i \in [N]$  receives  $z_i$  as input.
- $L$  epochs of communication
- For  $\ell = 1, 2, \dots, L$ , the  $\ell$ -th epoch proceeds as follows
  - For each player  $i \in [N]$ , player  $i$  sends a message  $\Gamma_{i,j}^{(\ell)}$  to previous player  $j \in [1 : i - 1]$ 
    - The player  $j$ , based on its own information state  $X_j^{(\ell-1)}$ , and the message  $\Gamma_{i,j}^{(\ell)}$ , replies with a message  $\Pi_{j,i}^{(\ell)}$  to player  $i$
  - The player  $i$  updates its information state as

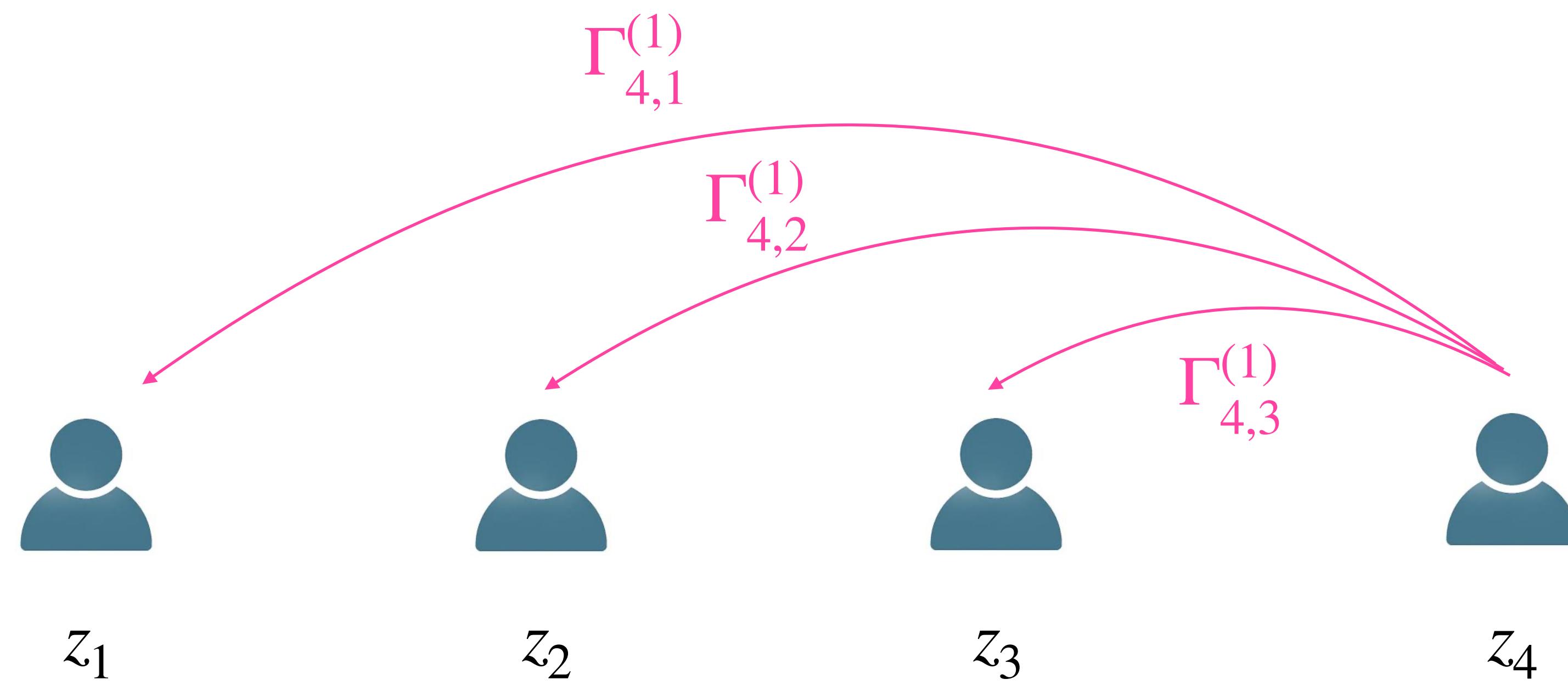
$$X_i^{(\ell)} := X_i^{(\ell-1)} \cup \bigcup_{j < i} \Pi_{j,i}^{(\ell)}$$

Inductively defined, initially  $X_j^{(0)} := z_j$

# Example

Number of players  $N = 4$   
Epochs of communication  $L = 1$

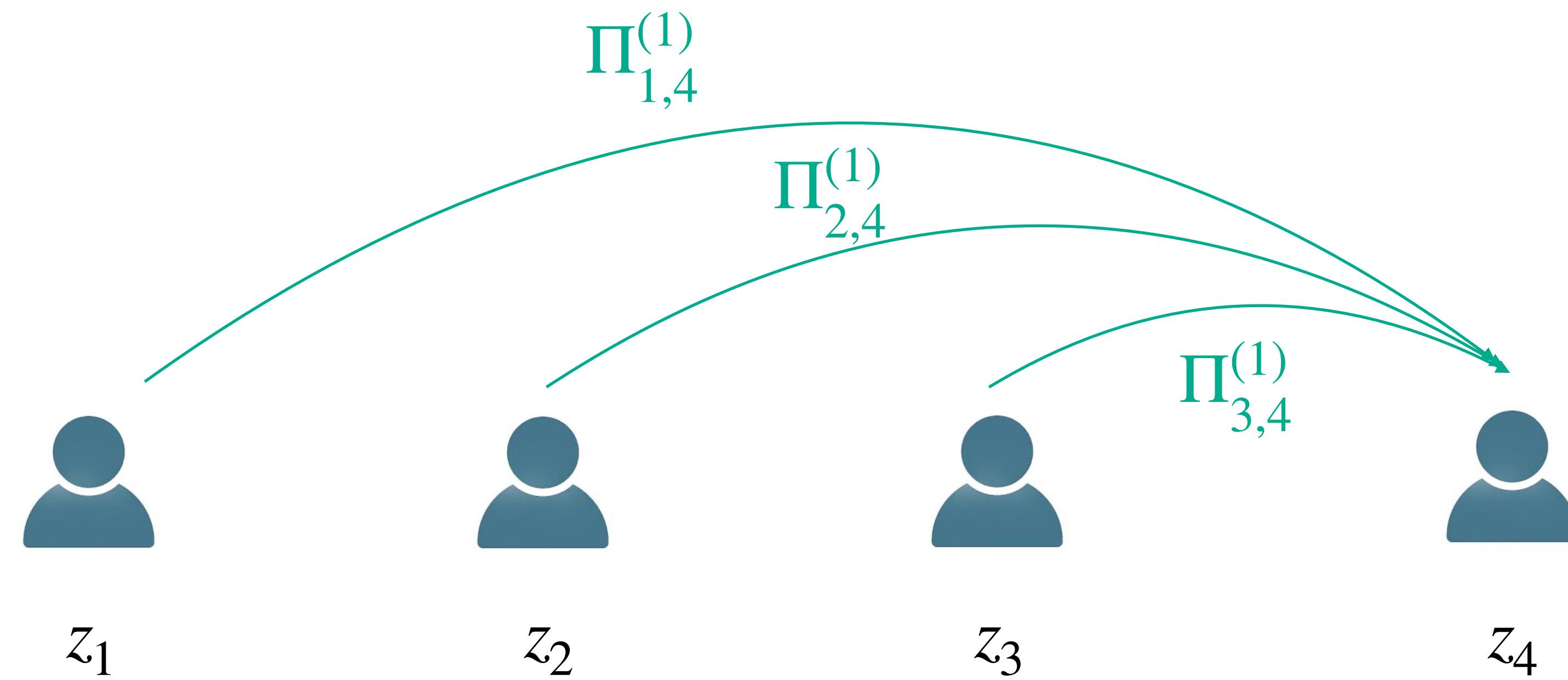
$\Gamma_{4,1}^{(1)}, \Gamma_{4,2}^{(1)}, \Gamma_{4,3}^{(1)}$  depends only on  $X_1^{(0)} := z_1$



# Example

Number of players  $N = 4$   
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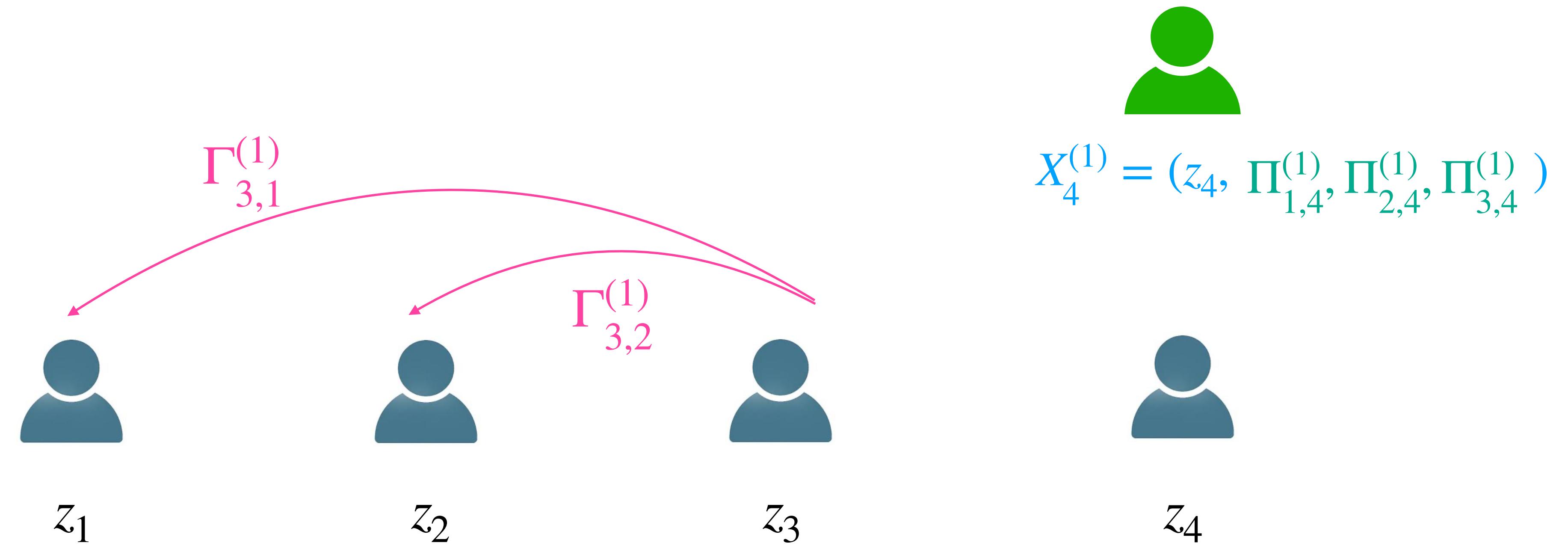
$\Pi_{1,4}^{(1)}$  depends on  $X_1^{(0)} := z_1$  and  $\Gamma_{4,1}^{(1)}$



# Example

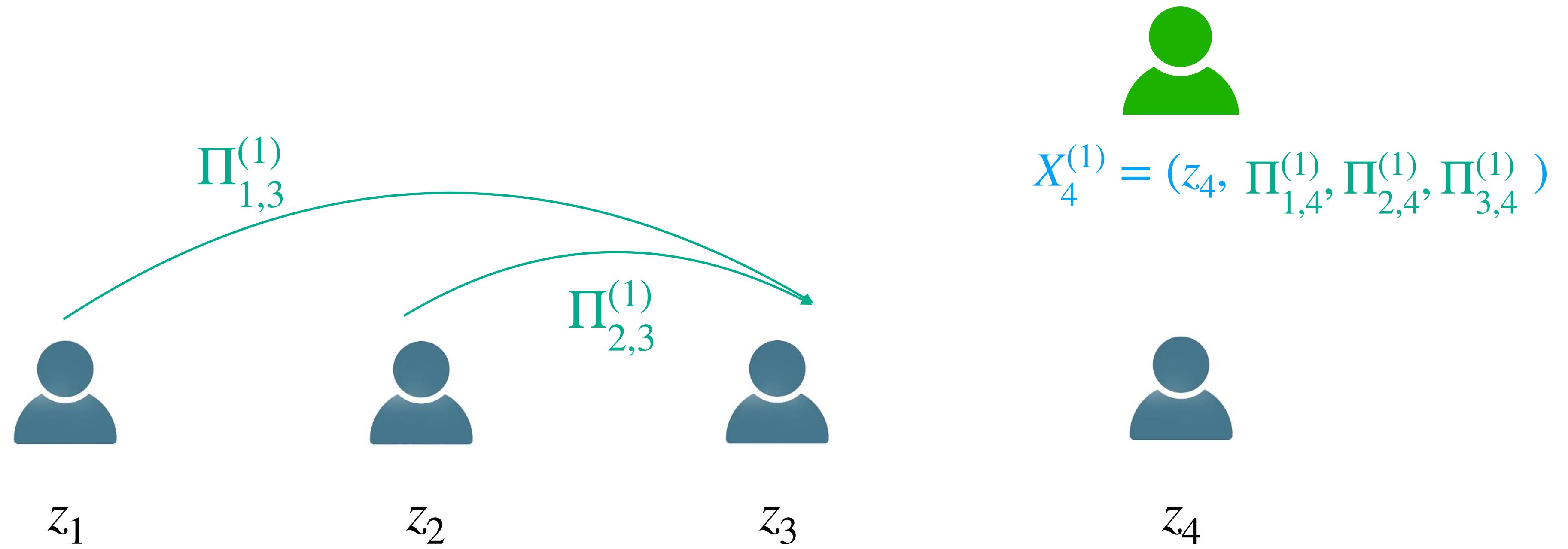
Number of players  $N = 4$   
Epochs of communication  $L = 1$

$\Gamma_{3,1}^{(1)}, \Gamma_{3,2}^{(1)}$ , depends only on  $X_3^{(0)} := z_3$



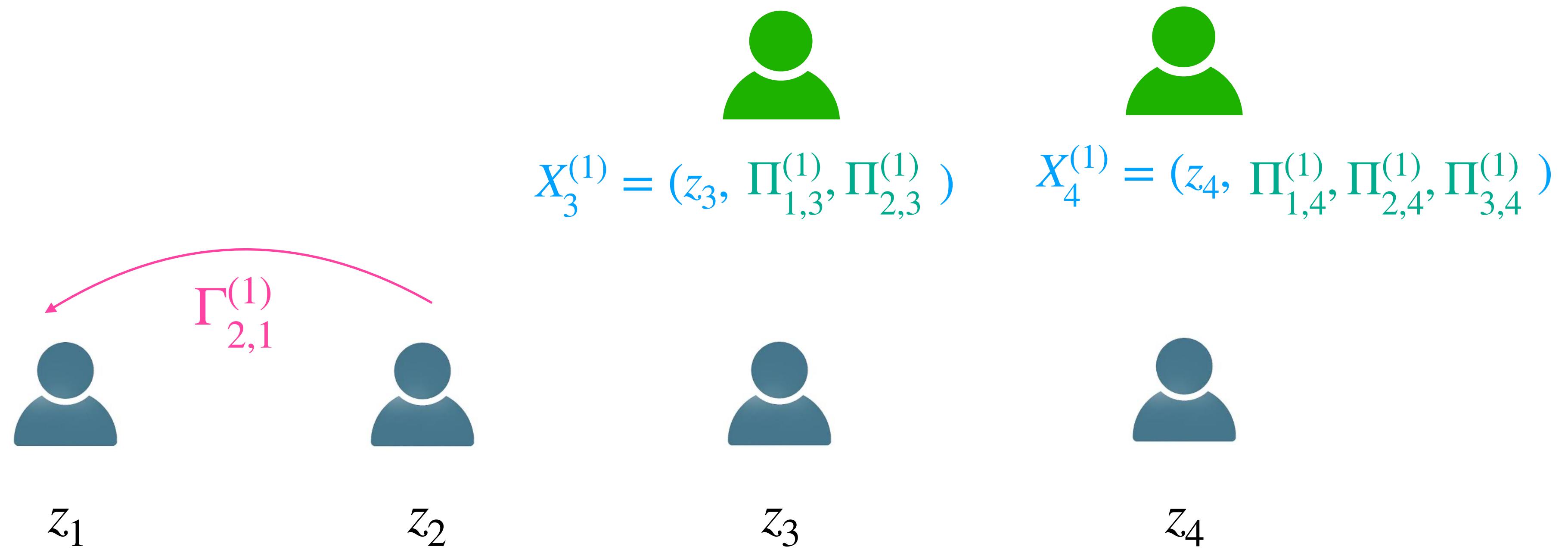
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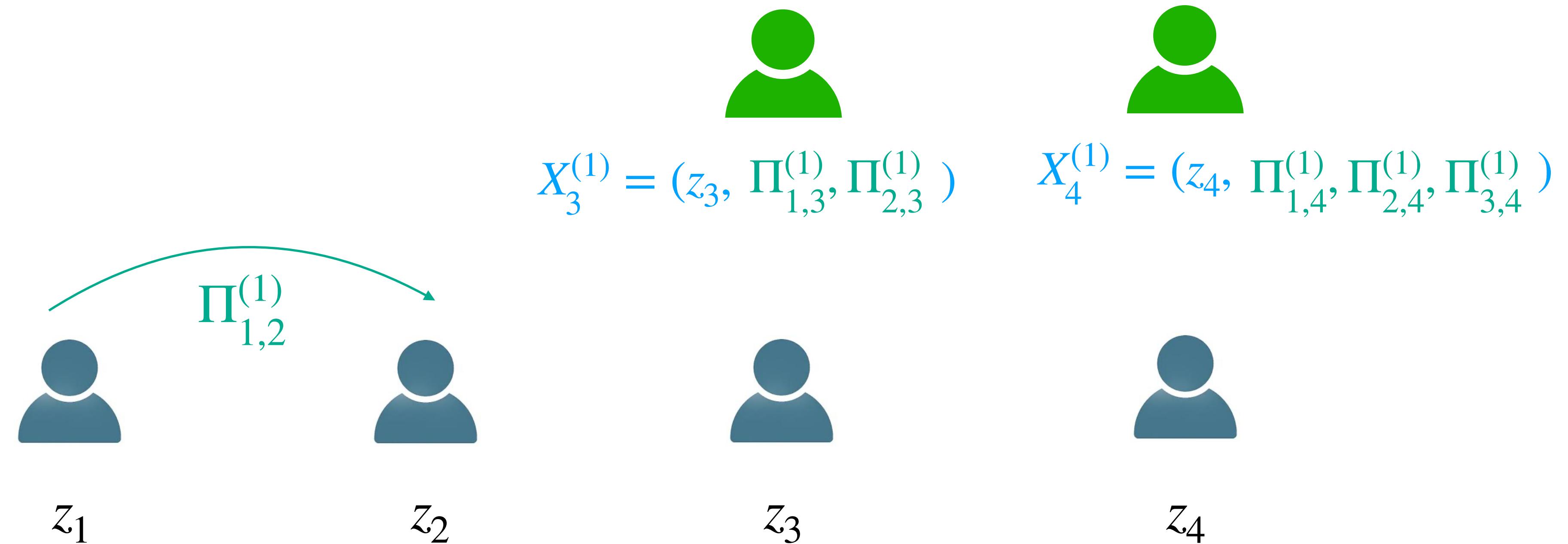
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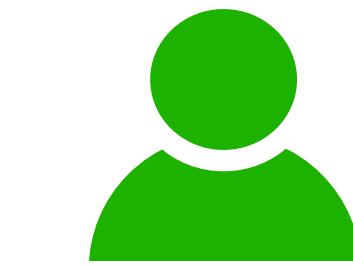
Number of players  $N = 4$   
Epochs of communication  $L = 1$



$$X_1^{(1)} = z_1$$



$$X_2^{(1)} = (z_2, \Pi_{1,2}^{(1)}) \quad X_3^{(1)} = (z_3, \Pi_{1,3}^{(1)}, \Pi_{2,3}^{(1)}) \quad X_4^{(1)} = (z_4, \Pi_{1,4}^{(1)}, \Pi_{2,4}^{(1)}, \Pi_{3,4}^{(1)})$$



$$z_1$$



$$z_2$$



$$z_3$$



$$z_4$$

# Our Results

$n$ : prompt length  
 $d$ : hidden dimension  
 $H$ : attention head  
 $p$ : precision

- **Theorem.** An  $L$ -layer decoder-only Transformer could not solve  $L$ -sequential function composition unless  $Hdp \geq n^{2^{-4L}}$



$L$ -step composition requires  $\Omega(L)$  layers of Transformer (asymptotically)

# Composition

- **$L$ -Sequential function composition:** Given  $L$  functions  $f_1, \dots, f_L$  and a query  $w = (w_1, \dots, w_L)$ , compute

$$i_1 = f_1(w_1), i_2 = f_2(w_2, i_1), \dots, i_L = f_L(w_L, i_{L-1})$$

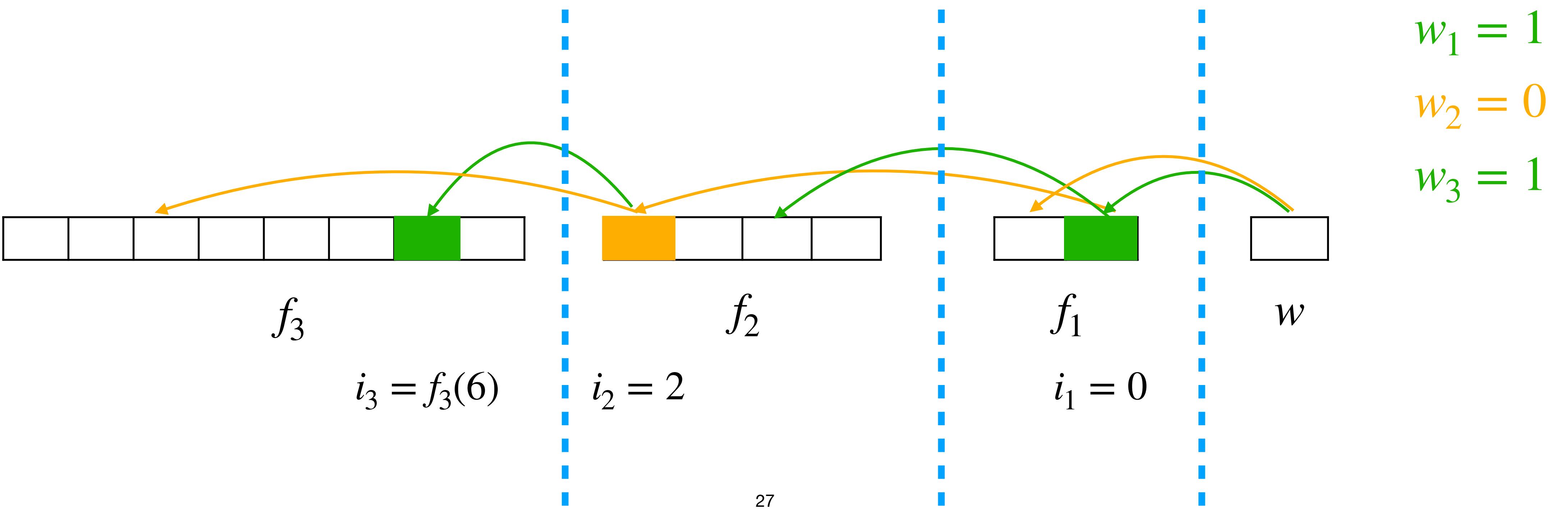
and output  $i_L$

- The input prompt explicitly describes the functions in the order of  $f_L, f_{L-1}, \dots, f_1$ , followed by the query  $w$ .

# Example

- $L = 3, w_1, w_2, w_3 \in \{0,1\}$
- $i_1 = f_1(w_1), i_2 = f_2(w_2, i_1), i_3 = f_3(w_3, i_2)$
- $f_1 : [2] \rightarrow [2], f_2 : [4] \rightarrow [4], f_3 : [8] \rightarrow [8]$

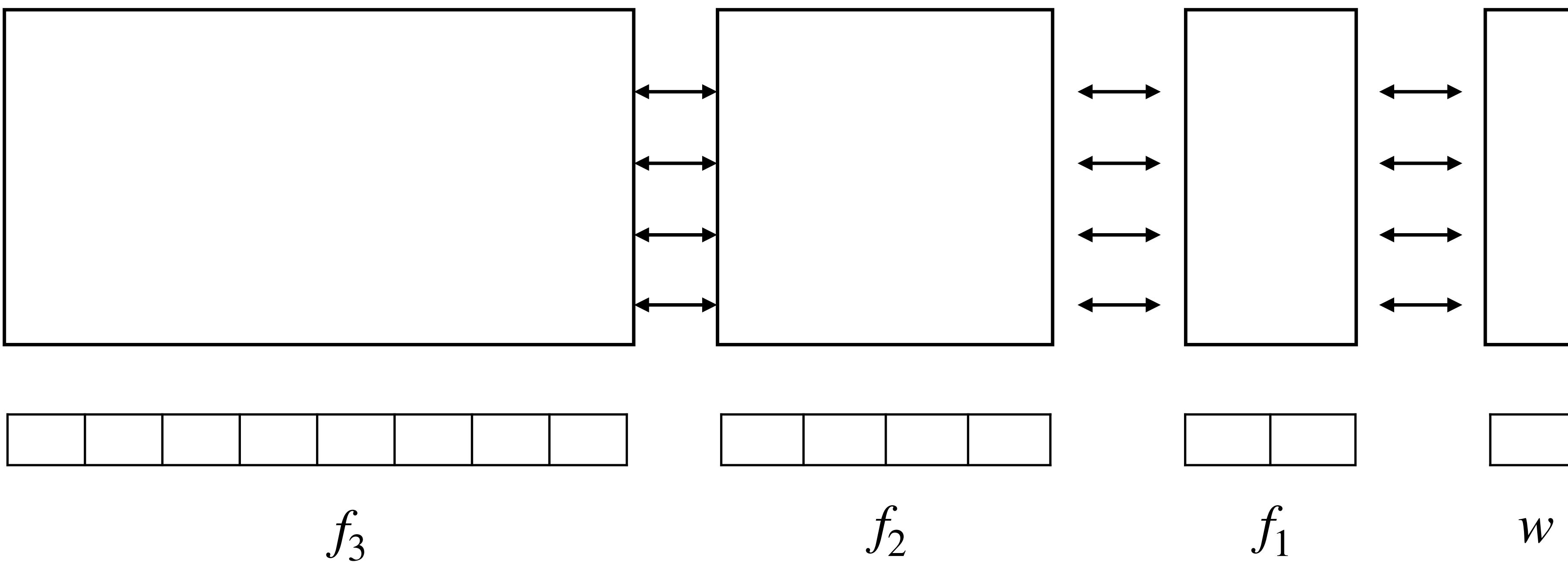
$$w = w_1 w_2 w_3 = \textcolor{green}{1} \textcolor{orange}{0} \textcolor{blue}{1}$$



# Intuition

- When important information is presented at the end of sequence, **autoregressive Transformer** fails to retrieve information efficiently

Massive computation are “wasted” at the beginning of the sequence



# Remarks on the proof

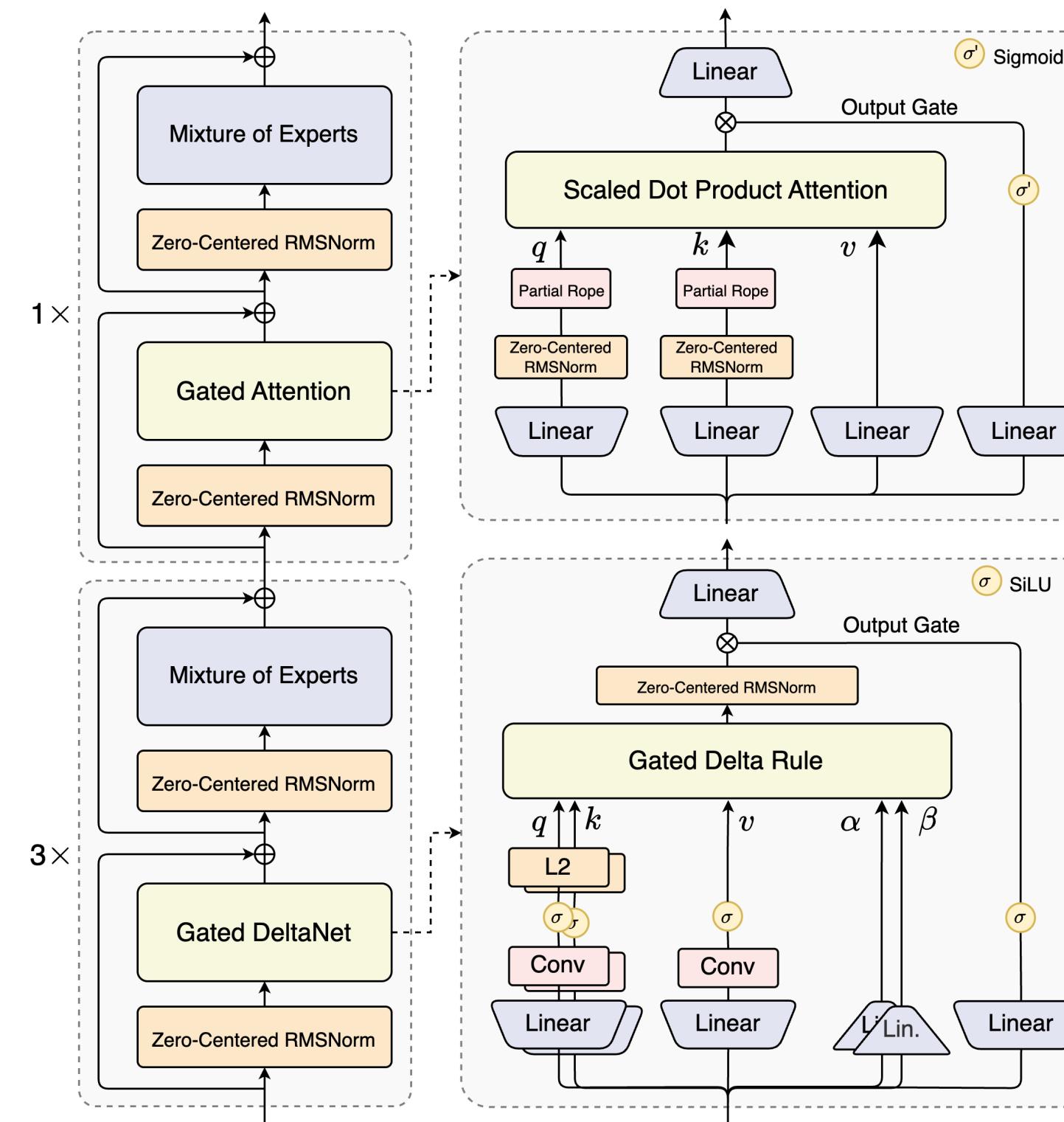
- How we actually prove the lower bound
  - Indistinguishable decomposition: Find  $R_{\leq \ell}$  and  $Z_{> \ell}$ , such that for every **inputs**  $z_{> \ell} \in Z_{> \ell}$  of **players**  $> \ell$ , all **assignments** in  $R_{\leq \ell}$  of **players**  $[1 : \ell]$  are **indistinguishable** after  $\ell$  epochs (i.e., they lead to the same transcripts)
  - We additional need the indistinguishable decomposition we find to have **large coverage** on  $i_\ell$



# Take away

- The first **unconditional lower bound** of multi-layer (decoder-only) Transformer
- Key concept idea: **Autoregressive communication**
  - To study computation model of Transformer, it is crucial to study new communication model (this leads to stronger and unconditional lower bound)
- The **communication perspective** offers yet another way of looking at architectures.

# Discussions - Hybrid models



Qwen3-Next

# Discussions - KV-cache sharing

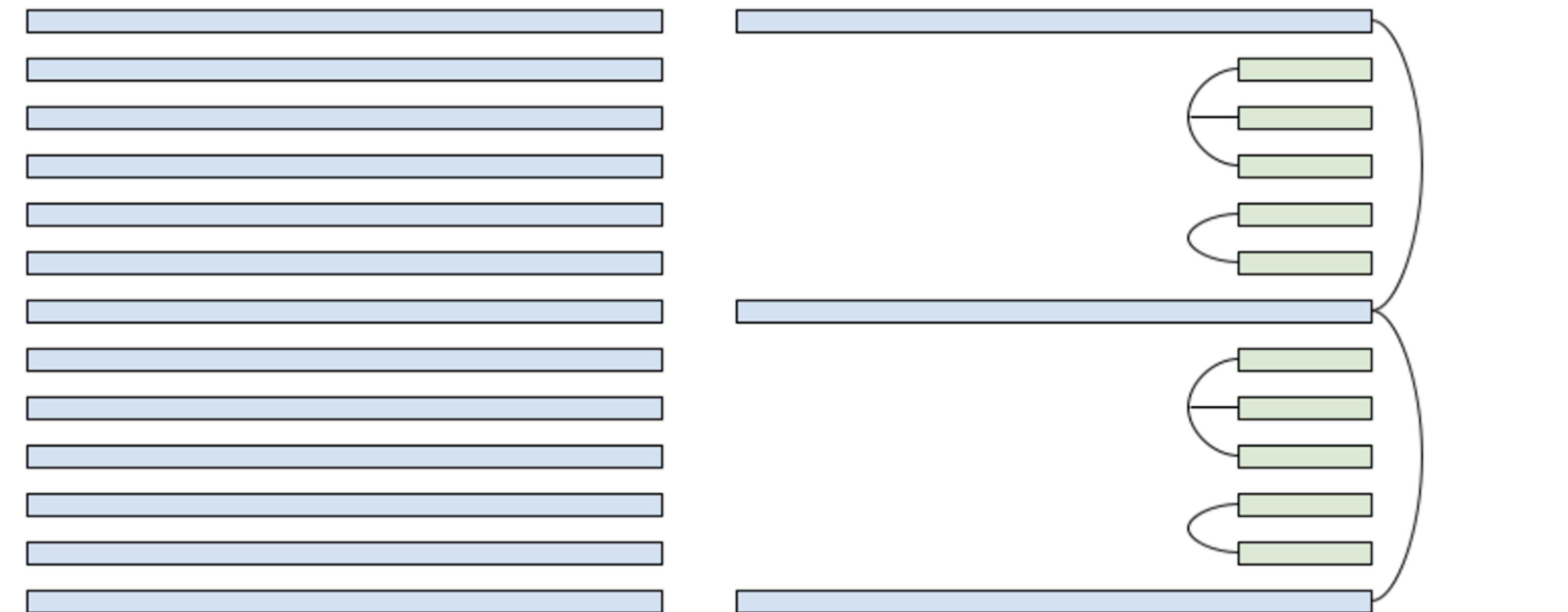
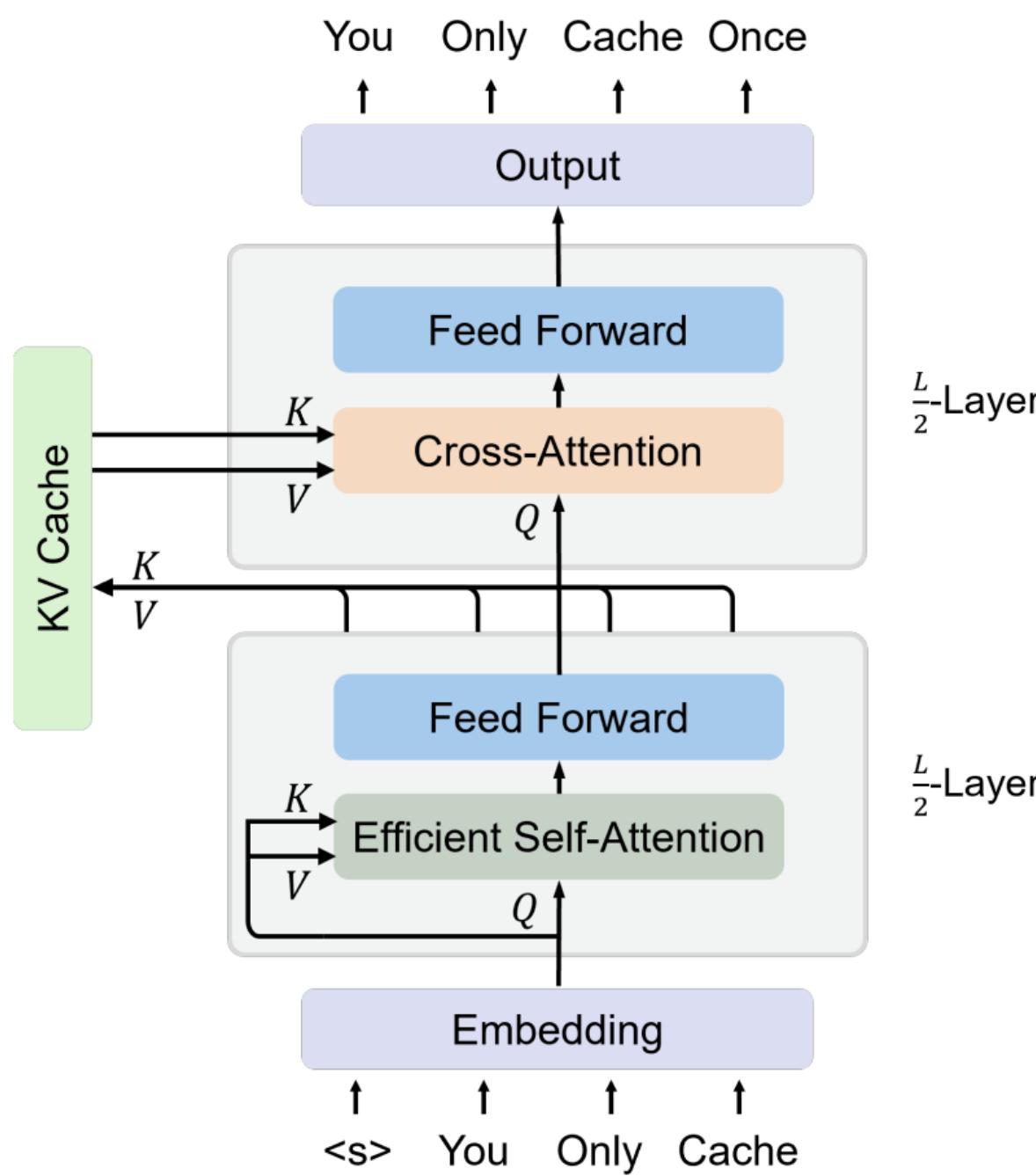


Figure 1. Left: Standard transformer design where every attention is global attention. Right: The attention design in our production model. Blue boxes indicate global attention, green boxes indicate local attention, and curves indicate KV-sharing. For global attention layers, we share KV across multiple non-adjacent layers. This illustration depicts only a subset of the layers in the full model.

YOCO [Sun et al.]

Character.ai blog

# Thanks!