ACM算法模板・一些常用的算法模板-模板合集(打比赛专用)

0.头文件

```
#define _CRT_SBCURE_NO_DEPRECATE
#include <set>
#include <cmath>
#include <queue>
#include <stack>
#include <vector>
#include <string>
#include <cstdio>
#include <cstdlib>
#include <cstring>
#include <iostream>
#include <algorithm>
#include <functional>
using namespace std;
const int maxn = 110;
const int INF = 0x3f3f3f3f;
```

经典

1.埃拉托斯特尼筛法

```
/*
    /埃式筛法/
   /快速筛选素数/
    /16/11/05ztx/
*/
int prime[maxn];
bool is_prime[maxn];
int sieve(int n) {
   int p = 0;
   for (int i = 0; i \le n; ++i)
       is prime[i] = true;
   is prime[0] = is prime[1] = false;
   for (int i = 2; i <= n; ++i) { // 注意数组大小是n
       if(is prime[i]){
           prime[p++] = i;
           for(int j = i + i; j <= n; j += i) // 轻剪枝, j必定是i的倍数
              is prime[j] = false;
   return p; // 返回素数个数
```

2.快速幂

3.大数模拟

大数加法

```
/*
    |大数模拟加法|
    /用string模拟/
    |16/11/05ztx, thanks to caojiji|
string add1(string s1, string s2)
   if (s1 == "" && s2 == "") return "0";
   if (s1 = "") return s2;
   if (s2 = "") return s1;
   string \max = s1, \min = s2;
   if (s1. length() < s2. length()) {
       maxx = s2;
       minn = s1;
   int a = \max x. length() - 1, b = \min n. length() - 1;
   for (int i = b; i >= 0; --i) {
       maxx[a--] += minn[i] - '0'; // a一直在减 , 额外还要减个'0'
   for (int i = \max x. length()-1; i > 0;--i) {
```

```
if (maxx[i] > '9') {
          maxx[i] -= 10;//注意这个是减10
          \max_{i} [i - 1] ++;
       }
   if (\max[0] > '9') {
       \max [0] = 10;
       \max x = 1' + \max x;
   return maxx;
}
大数阶乘
/*
   /大数模拟阶乘/
   |用数组模拟|
   /16/12/02ztx/
#include <iostream>
#include <cstdio>
using namespace std;
typedef long long LL;
const int maxn = 100010;
int num[maxn], len;
/*
   在mult函数中,形参部分: 1en每次调用函数都会发生改变,n表示每次要乘以的
数,最终返回的是结果的长度
   tip: 阶乘都是先求之前的(n-1)!来求n!
   初始化Init函数很重要,不要落下
*/
void Init() {
   1en = 1;
   num[0] = 1;
}
int mult(int num[], int len, int n) {
   LL tmp = 0;
   for (LL i = 0; i < len; ++i) {
        tmp = tmp + num[i] * n;
                               //从最低位开始,等号左边的tmp表示当前
```

位,右边的tmp表示进位(之前进的位)

```
num[i] = tmp % 10; // 保存在对应的数组位置,即去掉进位后的一位数
                                 tmp = tmp / 10; //  NEM = TEM = 
              while(tmp) { // 之后的进位处理
                                 num[1en++] = tmp \% 10;
                                 tmp = tmp / 10;
              return len;
int main() {
              Init();
              int n;
              n = 1977; // 求的阶乘数
              for (int i = 2; i \le n; ++i) {
                             len = mult(num, len, i);
              for(int i = len - 1; i >= 0; --i)
                             printf("%d", num[i]); // 从最高位依次输出,数据比较多采用printf输
Щ
              printf("\n");
              return 0;
4.GCD
/*
               |辗转相除法|
               |欧几里得算法|
               /求最大公约数/
               /16/11/05ztx/
int gcd(int big, int small)
              if (small > big) swap(big, small);
              int temp;
              while (small != 0) { // 辗转相除法
                             if (small > big) swap(big, small);
                             temp = big % small;
                             big = small;
                             small = temp;
             return(big);
```

5.LCM

```
/*
    |辗转相除法|
    |欧几里得算法|
    |求最小公倍数|
    /16/11/05ztx/
int gcd(int big, int small)
   if (small > big) swap(big, small);
   int temp;
   while (small != 0) { // 辗转相除法
       if (small > big) swap(big, small);
       temp = big % small;
       big = small;
       small = temp;
   return(big);
}
6.全排列
/*
    /求1到n的全排列, 有条件/
    |16/11/05ztx, thanks to wangqiqi|
void Pern(int list[], int k, int n) { // k表示前k个数不动仅移动后面n-k位
数
   if (k == n - 1) {
       for (int i = 0; i < n; i++) {
          printf("%d", list[i]);
       printf("\n");
   }else {
       for (int i = k; i < n; i++) { // 输出的是满足移动条件所有全排列
           swap(list[k], list[i]);
          Pern(list, k + 1, n);
          swap(list[k], list[i]);
   }
7.二分搜索
/*
    二分搜索|
    |要求: 先排序|
    16/11/05ztx, thanks to wangxiaocai
*/
```

```
// left为最开始元素, right是末尾元素的下一个数, x是要找的数
int bsearch(int *A, int left, int right, int x) {
   int m;
   while (left < right) {
      m = left + (right - left) / 2;
      if (A[m] >= x) right = m; else left = m + 1;
      // 如果要替换为 upper bound, 改为:if (A[m] <= v) x = m+1; else y =
M;
   return left;
}
/*
   最后left == right
如果没有找到135577找6,返回7
如果找有多少的x,可以用lower bound查找一遍, upper bound查找一遍,下标相减
C++自带的lower bound(a, a+n, x)返回数组中最后一个x的下一个数的地址
upper bound (a, a+n, x) 返回数组中第一个x的地址
如果a+n内没有找到x或x的下一个地址,返回a+n的地址
lower_bound(a, a+n, x) -upper_bound(a, a+n, x)返回数组中x的个数
*/
```

数据结构

并查集

8.并查集

```
int root = x; // 根节点
   while (root != father[root]) { // 寻找根节点
      root = father[root];
   while (x != root) {
      int tmp = father[x];
      father[x] = root; // 根节点赋值
      x = tmp;
   return root;
}
void Union(int x, int y) { // 将x所在的集合和y所在的集合整合起来形成一个
集合。
   int a, b;
   a = findRoot(x);
   b = findRoot(y);
   father[a] = b; // y连在x的根节点上 或father[b] = a为x连在y的根节点
上;
/*
   在findRoot(x)中:
   路径压缩 迭代 最优版
   关键在于在路径上的每个节点都可以直接连接到根上
```

图论

MST

最小生成树

Kruskal

9.克鲁斯卡尔算法

```
void Kruskal() {
    ans = 0;
    for (int i = 0; i < len; i++) {
        if (Find(edge[i].a) != Find(edge[i].b)) {
            Union(edge[i].a, edge[i].b);
            ans += edge[i].len;
        }
    }
}</pre>
```

Prim

10.普里姆算法

```
/*
    |Prim算法|
    |适用于 稠密图 求最小生成树|
    |堆优化版,时间复杂度: 0(e1gn)|
    |16/11/05ztx, thanks to chaixiaojun|
struct node {
   int v, len;
   node(int v = 0, int len = 0) : v(v), len(len) {}
   bool operator < (const node &a) const { // 加入队列的元素自动按距离从小
到大排序
       return len> a.len;
   }
};
vector<node> G[maxn];
int vis[maxn];
int dis[maxn];
void init() {
   for (int i = 0; i < maxn; i++) {
       G[i]. clear();
       dis[i] = INF;
       vis[i] = false;
}
int Prim(int s) {
   priority queue<node>Q; // 定义优先队列
   int ans = 0;
   Q. push (node (s, 0)); // 起点加入队列
   while (!Q. empty()) {
       node now = Q. top(); Q. pop(); // 取出距离最小的点
       int v = now. v;
```

```
if (vis[v]) continue; // 同一个节点,可能会推入2次或2次以上队列,
这样第一个被标记后, 剩下的需要直接跳过。
vis[v] = true; // 标记一下
ans += now.len;
for (int i = 0; i<G[v].size(); i++) { // 开始更新
int v2 = G[v][i].v;
int len = G[v][i].len;
if (!vis[v2] && dis[v2] > len) {
dis[v2] = len;
Q. push(node(v2, dis[v2])); // 更新的点加入队列并排序
}
}
return ans;
}
```

Bellman-Ford

单源最短路

Dijkstra

11.迪杰斯特拉算法

```
/*
   |Dijkstra算法|
   |适用于边权为正的有向图或者无向图|
   |求从单个源点出发,到所有节点的最短路|
   |优化版: 时间复杂度 0(e1bn)|
   |16/11/05ztx, thanks to chaixiaojun|
*/
struct node {
   int v, len;
   node(int v = 0, int len = 0) : v(v), len(len) {}
   bool operator < (const node &a)const { // 距离从小到大排序
       return len > a.len;
};
vector<node>G[maxn];
bool vis[maxn];
int dis[maxn];
void init() {
   for (int i = 0; i < maxn; i++) {
       G[i].clear();
       vis[i] = false;
```

```
dis[i] = INF;
}
int dijkstra(int s, int e) {
   priority_queue<node>Q;
   Q. push (node (s, 0)); // 加入队列并排序
   dis[s] = 0;
   while (!Q.empty()) {
       node now = Q. top();
                          // 取出当前最小的
       Q. pop();
       int v = now. v;
       if (vis[v]) continue; // 如果标记过了,直接continue
       vis[v] = true:
       int v2 = G[v][i].v;
          int len = G[v][i]. len;
          if (!vis[v2] \&\& dis[v2] > dis[v] + 1en) {
              dis[v2] = dis[v] + 1en;
              Q. push (node(v2, dis[v2]));
          }
       }
   return dis[e];
```

SPFA

12.最短路径快速算法 (Shortest Path Faster Algorithm)

```
int v1, v2, weight;
   queue <int> Q;
   memset (inqueue, false, sizeof (inqueue)); // 标记是否在队列中
   memset(cnt, 0, sizeof(cnt)); // 加入队列的次数
   dist[s] = 0;
   Q. push(s); // 起点加入队列
   inqueue[s] = true; // 标记
   while (!Q. empty()) {
       v1 = Q. front();
       Q. pop();
       inqueue[v1] = false; // 取消标记
       for(int i = 0; i < G[v1]. size(); ++i){ // 搜索v1的链表
          v2 = G[v1][i].vex:
          weight = G[v1][i]. weight;
          if(dist[v2] > dist[v1] + weight){ // 松弛操作
              dist[v2] = dist[v1] + weight;
              if (inqueue[v2] == false) { // 再次加入队列
                  inqueue[v2] = true;
                  //cnt[v2]++; // 判负环
                  //if(cnt[v2] > n) return -1;
                  Q. push (v2);
              } } }
   return dist[e];
}
/*
   不断的将s的邻接点加入队列,取出不断的进行松弛操作,直到队列为空
如果一个结点被加入队列超过n-1次,那么显然图中有负环
*/
```

Floyd-Warshall

13.弗洛伊德算法

```
for (int j = 0; j < n; j++) {
         dis[i][j] = min(dis[i][j], dis[i][k] + dis[k][j]);
}
}</pre>
```

二分图

14.染色法

```
/*
    |交叉染色法判断二分图|
    /16/11/05ztx/
int bipartite(int s) {
    int u, v;
    queue<int>Q;
    color[s] = 1;
    Q. push(s);
    while (!Q.empty()) {
        u = Q. front();
        Q. pop();
        for (int i = 0; i < G[u].size(); i++) {
            v = G[u][i];
            if (color[v] == 0) {
                color[v] = -color[u];
                Q. push (v);
            else if (color[v] == color[u])
                return 0;
    return 1;
}
```

15..匈牙利算法

```
/*
    /求解最大匹配问题/
    /递归实现/
    /16/11/05ztx/

*/
vector<int>G[maxn];
bool inpath[maxn]; // 标记
int match[maxn]; // 记录匹配对象
void init()
{
```

```
memset (match, -1, sizeof (match));
    for (int i = 0; i < maxn; ++i) {
       G[i].clear();
bool findpath(int k) {
    for (int i = 0; i < G[k].size(); ++i) {
       int v = G[k][i];
       if (!inpath[v]) {
            inpath[v] = true;
           if (match[v] == -1 | findpath(match[v])) { // 递归
               match[v] = k; // 即匹配对象是 "k妹子" 的
               return true:
   return false;
}
void hungary() {
    int cnt = 0;
    for (int i = 1; i <= m; i++) { // m为需要匹配的"妹子"数
       memset(inpath, false, sizeof(inpath)); // 每次都要初始化
       if (findpath(i)) cnt++;
   cout << cnt << end1;
}
/*
    |求解最大匹配问题|
    |dfs实现|
    /16/11/05ztx/
*/
int v1, v2;
bool Map[501][501];
bool visit[501];
int link[501];
int result;
bool dfs(int x) {
    for (int y = 1; y \le v2; ++y) {
       if (Map[x][y] && !visit[y]) {
           visit[y] = true;
           if (link[y] = 0 \mid | dfs(link[y])) {
               link[y] = x;
               return true;
           } } }
```

```
return false;
void Search() {
    for (int x = 1; x \le v1; x++) {
        memset(visit, false, sizeof(visit));
        if (dfs(x))
           result++;
}
动态规划
背包
16.17.18背包问题
    /01背包/
    /完全背包/
    |多重背包|
    /16/11/05ztx/
*/
// 01背包:
void bag01(int cost, int weight) {
    for (i = v; i \ge cost; --i)
    dp[i] = max(dp[i], dp[i-cost]+weight);
}
// 完全背包:
void complete(int cost, int weight) {
    for (i = cost ; i \le v; ++i)
    dp[i] = max(dp[i], dp[i - cost] + weight);
   多重背包:
void multiply(int cost, int weight, int amount) {
    if (cost * amount \geq v)
        complete(cost, weight);
    else{
       k = 1;
        while (k < amount) {
           bag01(k * cost, k * weight);
           amount -= k;
```

```
k += k;
        bag01(cost * amount, weight * amount);
}
// other
int dp[1000000];
int c[55], m[110];
int sum;
void CompletePack(int c) {
    for (int v = c; v \le sum / 2; ++v) {
        dp[v] = \max(dp[v], dp[v - c] + c);
}
void ZeroOnePack(int c) {
    for (int v = sum / 2; v >= c; --v) {
        dp[v] = max(dp[v], dp[v - c] + c);
}
void multiplePack(int c, int m) {
    if (m * c > sum / 2)
        CompletePack(c);
    else{
        int k = 1;
        while (k < m) {
            ZeroOnePack(k * c);
            m = k;
            k \ll 1;
        if (m != 0) {
            ZeroOnePack(m * c);
}
```

LIS

19.最长上升子序列

```
/*
|最长上升子序列|
|状态转移|
```

```
/16/11/05ztx/
*/
/*
   状态转移dp[i] = max\{1.dp[j] + 1\}; j < i; a[j] < a[i];
   d[i]是以i结尾的最长上升子序列
   与i之前的 每个a[j]<a[i]的 j的位置的最长上升子序列+1后的值比较
*/
void solve(){ // 参考挑战程序设计入门经典;
   for (int i = 0; i < n; ++i) {
      dp[i] = 1;
      for (int j = 0; j < i; ++j) {
          if(a[j] < a[i])
             dp[i] = max(dp[i], dp[j] + 1);
          } } }
}
/*
   优化方法:
   dp[i]表示长度为i+1的上升子序列的最末尾元素
   找到第一个比dp末尾大的来代替
void solve() {
      for (int i = 0; i < n; ++i) {
          dp[i] = INF;
      for (int i = 0; i < n; ++i) {
          *lower bound(dp, dp + n, a[i]) = a[i]; // 返回一个指针
      printf("%d\n", *lower_bound(dp, dp + n, INF) - dp;
/*
   函数lower_bound()返回一个 iterator 它指向在[first, last)标记的有序序列中
可以插入value,而不会破坏容器顺序的第一个位置,而这个位置标记了一个不小于
value的值。
*/
```

LCS

20.最长公共子序列

/* |求最长公共子序列|

```
/遊推形式/

/16/11/05ztx/

*/

void solve() {
    for (int i = 0; i < n; ++i) {
        for (int j = 0; j < m; ++j) {
            if (s1[i] == s2[j]) {
                dp[i + 1][j + 1] = dp[i][j] + 1;
            } else {
                dp[i + 1][j + 1] = max(dp[i][j + 1], dp[i + 1][j]);
            } } }
}
```

计算几何

21.向量基本用法

```
/*
    /16/11/06ztx/
struct node {
    double x; // 横坐标
    double y; // 纵坐标
};
typedef node Vector;
Vector operator + (Vector A, Vector B) { return Vector(A.x + B.x, A.y +
B. y); }
Vector operator - (Point A, Point B) { return Vector(A.x - B.y, A.y -
B. y); }
Vector operator * (Vector A, double p) { return Vector(A. x*p, A. y*p); }
Vector operator / (Vector A, double p) { return Vector (A. x / p, A. y*p); }
double Dot(Vector A, Vector B) { return A. x*B. x + A. y*B. y; } // 向量点乘
double Length(Vector A) { return sqrt(Dot(A, A)); } // 向量模长
double Angle (Vector A, Vector B) { return acos (Dot (A, B) / Length (A) /
double Cross(Vector A, Vector B) { // 叉积计算 公式
   return A. x*B. y - A. y*B. x;
Vector Rotate(Vector A, double rad) // 向量旋转 公式 {
    return Vector (A. x*cos (rad) - A. y*sin (rad), A. x*sin (rad) +
A. y*cos (rad));
```

```
Point getLineIntersection(Point P, Vector v, Point Q, Vector w) { // 两直线
交点t1 t2计算公式
    Vector u = P - Q;
    double t = Cross(w, u) / Cross(v, w); // 求得是横坐标
   return P + v*t; // 返回一个点
}
22.求多边形面积
/*
    /16/11/06ztx/
node G[maxn];
int n;
double Cross (node a, node b) { // 叉积计算
   return a. x*b. y - a. y*b. x;
int main()
    while (scanf("%d", &n) != EOF && n) {
       for (int i = 0; i < n; i++)
           scanf("%lf %lf", &G[i].x, &G[i].y);
       double sum = 0;
       G[n].x = G[0].x;
       G[n].y = G[0].y;
       for (int i = 0; i < n; i++) {
               sum += Cross(G[i], G[i + 1]);
       // 或者
           //for (int i = 0; i < n; i++) {
               //sum += fun(G[i], G[(i+1)\% n]);
           //}
       sum = sum / 2.0;
       printf("\%.1f\n", sum);
    system("pause");
   return 0;
23..判断线段相交
```

/16/11/06ztx/

```
node P[35][105];
double Cross_Prouct (node A, node B, node C) { // 计算BA叉乘CA
    return (B. x-A. x)*(C. y-A. y)-(B. y-A. y)*(C. x-A. x);
bool Intersect (node A, node B, node C, node D) { // 通过叉乘判断线段是否相
交;
    if (\min(A. x, B. x) \leq \max(C. x, D. x) \&\&
                                            // 快速排斥实验;
       min(C. x, D. x) \le max(A. x, B. x) &&
       min(A. y, B. y) \le max(C. y, D. y) &&
       min(C. y, D. y) \le max(A. y, B. y) &&
       Cross_Prouct(A, B, C)*Cross_Prouct(A, B, D)<0&& // 跨立实验;
       Cross_Prouct(C, D, A)*Cross_Prouct(C, D, B)<0) // 叉乘异号表示在
两侧;
      return true;
    else return false;
}
24.求三角形外心
/*
    /16/11/06ztx/
Point circumcenter(const Point &a, const Point &b, const Point &c) { //返回
三角形的外心
   Point ret;
    double a1 = b.x - a.x, b1 = b.y - a.y, c1 = (a1*a1 + b1*b1) / 2;
    double a2 = c. x - a. x, b2 = c. y - a. y, c2 = (a2*a2 + b2*b2) / 2;
    double d = a1*b2 - a2*b1;
    ret. x = a. x + (c1*b2 - c2*b1) / d;
    ret. y = a. y + (a1*c2 - a2*c1) / d;
   return ret;
}
24.极角排序
/*
    /16/11/06ztx/
*/
double cross(point pl, point p2, point q1, point q2) { // 叉积计算
    return (q2. y - q1. y)*(p2. x - p1. x) - (q2. x - q1. x)*(p2. y - p1. y);
bool cmp(point a, point b) {
    point o;
   0. x = 0. y = 0;
   return cross(o, b, o, a) < 0; // 叉积判断
}
```

字符串

kmp

25.克努特-莫里斯-普拉特操作

```
/*
    /kmp算法/
    /字符串匹配/
    /17/1/21ztx/
void getnext(char str[maxn], int nextt[maxn]) {
   int j = 0, k = -1;
   nextt[0] = -1;
   while (j < m) {
       if (k == -1 || str[j] == str[k]) {
           j++;
           k++;
           nextt[j] = k;
       else
           k = nextt[k];
}
void kmp(int a[maxn], int b[maxn]) {
    int nextt[maxm];
   int i = 0, j = 0;
   getnext(b, nextt);
   while (i < n) {
       if (j == -1 || a[i] == b[j]) { // 母串不动,子串移动
           j++;
           i++;
       else {
           // i不需要回溯了
           // i = i - j + 1;
           j = nextt[j];
       if (j == m) {
           printf("%d\n", i - m + 1); // 母串的位置减去子串的长度+1
           return:
   printf("-1 \ ");
```

26.kmp扩展

```
/16/11/06ztx/
*/
#include iostream
#include<cstring>
using namespace std;
const int MM=100005;
int next[MM], extand[MM];
char S[MM], T[MM];
void GetNext(const char *T) {
    int len = strlen(T), a = 0;
    next[0] = 1en;
    while (a < len - 1 \&\& T[a] == T[a + 1]) a++:
    next[1] = a;
    a = 1;
    for (int k = 2; k < 1en; k ++) {
        int p = a + next[a] - 1, L = next[k - a];
        if ((k-1) + L) = p) {
            int j = (p - k + 1) > 0? (p - k + 1) : 0;
            while (k + j < len \&\& T[k + j] == T[j]) j++;
            next[k] = j;
            a = k;
        else next[k] = L;
void GetExtand(const char *S, const char *T) {
    GetNext(T):
    int slen = strlen(S), tlen = strlen(T), a = 0;
    int MinLen = slen < tlen ? slen : tlen;</pre>
    while (a \leq MinLen && S[a] == T[a]) a++;
    extand[0] = a;
    a = 0;
    for (int k = 1; k < slen; k ++) {
        int p = a + extand[a] - 1, L = next[k - a];
        if ((k-1) + L) = p) {
            int j = (p - k + 1) > 0? (p - k + 1) : 0;
            while (k + j < slen && j < tlen && S[k + j] == T[j]) j ++;
            extand[k] = j;
            a = k;
```

字典树

27.字典树

```
/*
    16/11/06ztx
*/
struct Trie{
    int cnt;
    Trie *next[maxn];
    Trie() {
        cnt = 0;
        memset(next, 0, sizeof(next));
};
Trie *root;
void Insert(char *word) {
    Trie *tem = root;
    while(*word != '\0') {
        int x = *word - 'a';
        if(tem->next[x] == NULL)
             tem->next[x] = new Trie;
        tem = tem \rightarrow next[x];
        tem->cnt++;
        word++;
```

```
}
int Search(char *word) {
    Trie *tem = root;
    for (int i=0; word[i]!='\0'; i++) {
        int x = word[i]-'a';
        if(tem->next[x] == NULL)
             return 0;
        tem = tem \rightarrow next[x];
    }
    return tem->cnt;
void Delete(char *word, int t) {
    Trie *tem = root;
    for (int i=0; word[i]!='\0'; i++)
        int x = word[i] - a';
        tem = tem \rightarrow next[x];
         (tem->cnt)-=t;
    for (int i=0; i < \max; i++)
        tem->next[i] = NULL;
int main() {
    int n;
    char str1[50];
    char str2[50];
    while (scanf ("%d", &n) !=EOF) {
        root = new Trie;
        while (n--)
             scanf ("%s %s", str1, str2);
             if(str1[0]=='i') {
                 Insert(str2);
             else if(str1[0] == 's')
                 if (Search(str2))
                      printf("Yes\n");
                 else
                      printf("No\n");
             }else {
                 int t = Search(str2);
                 if(t)
                     Delete(str2, t);
             } } }
    return 0;
```

28.AC自动机

```
/*
    16/11/06ztx
*/
#include iostream
#include<cstdio>
#include < cstring >
#include<string>
using namespace std;
#define N 1000010
char str[N], keyword[N];
int head, tail;
struct node {
    node *fail;
    node *next[26];
    int count;
    node() { //init
        fail = NULL;// 默认为空
        count = 0;
        for (int i = 0; i < 26; ++i)
            next[i] = NULL;
*_{q}[N];
node *root;
void insert(char *str) { // 建立Trie
    int temp, len;
    node *p = root;
    len = strlen(str);
    for (int i = 0; i < 1en; ++i) {
        temp = str[i] - 'a';
        if(p->next[temp] == NULL)
            p->next[temp] = new node();
        p = p-next[temp];
    p->count++;
}
void build ac() { // 初始化fail指针, BFS 数组模拟队列:
    q[tail++] = root;
    while(head != tail) {
```

```
node *p = q[head++]; // 弹出队头
        node *temp = NULL;
        for (int i = 0; i < 26; ++i) {
             if(p->next[i] != NULL)
                 if (p == root) { // 第一个元素fail必指向根
                     p\rightarrow next[i]\rightarrow fail = root;
                 }else {
                     temp = p->fail; // 失败指针
                     while(temp != NULL) { // 2种情况结束: 匹配为空or找到匹
配
                          if(temp->next[i]!= NULL) { // 找到匹配
                              p\rightarrow next[i]\rightarrow fail = temp\rightarrow next[i];
                              break:
                          temp = temp->fail;
                     if(temp == NULL) // 为空则从头匹配
                         p->next[i]->fail = root;
                 q[tail++] = p->next[i]; // 入队
            } } }
}
int query() // 扫描
    int index, len, result;
    node *p = root; // Tire\lambda□
    result = 0;
    len = strlen(str);
    for (int i = 0; i < len; ++i)
        index = str[i] - 'a';
        while(p->next[index] == NULL && p != root) // 跳转失败指针
            p = p \rightarrow fai1;
        p = p \rightarrow next[index];
        if(p == NULL)
            p = root:
        node *temp = p; // p不动, temp计算后缀串
        while (\text{temp} != \text{root \&\& temp-}) count != -1)
            result += temp->count;
            temp \rightarrow count = -1;
            temp = temp->fail;
    return result;
}
```

```
int main() {
    int num;
    head = tail = 0;
    root = new node();
    scanf ("%d", &num);
    getchar();
    for(int i = 0; i < num; ++i) {
       scanf ("%s", keyword);
        insert(keyword);
    build ac();
    scanf("%s", str);
    if (query())
        printf("YES\n");
    else
        printf("NO\n");
    return 0;
}
```

/* 假设有N个模式串,平均长度为L;文章长度为M。 建立Trie树: 0(N*L) 建立fail指针: 0(N*L) 模式匹配: 0(M*L) 所以,总时间复杂度为:0((N+M)*L)。*/

线段树

29.线段树

1) 点更新

```
tree[m].right = r;
    if (1 == r) \{
        tree[m]. \max = a[1];
        tree[m]. sum = a[1];
        return;
    }
    int mid = (1 + r) >> 1;
    build (m << 1, 1, mid);
    build(m << 1 | 1, mid + 1, r);
    tree[m]. max = max(tree[m << 1]. max, tree[m << 1 | 1]. max);
    tree[m].sum = tree[m << 1].sum + tree[m << 1 | 1].sum;
}
void update(int m, int a, int val) //a 是 节点位置, val 是 更新的值(加减的
值)
    if (tree[m].left == a \&\& tree[m].right == a) {
        tree[m].max += val;
        tree[m].sum += val;
        return;
    int mid = (tree[m].left + tree[m].right) >> 1;
    if (a \leq mid)
        update (m << 1, a, val);
    }
    else{
        update (m << 1 | 1, a, val);
    tree[m]. max = max(tree[m << 1]. max, tree[m << 1 | 1]. max);
    tree[m].sum = tree[m << 1].sum + tree[m << 1 | 1].sum;
}
int querySum(int m, int 1, int r)
    if (1 == tree[m].left \&\& r == tree[m].right)
        return tree[m].sum;
    int mid = (tree[m].left + tree[m].right) >> 1;
    if (r \leq mid) {
        return querySum(m << 1, 1, r);
    else if (1 > mid) {
        return querySum(m \langle\langle 1 \mid 1, 1, r \rangle\rangle;
    return querySum(m \ll 1, 1, mid) + querySum(m \ll 1 | 1, mid + 1, r);
}
```

```
int queryMax(int m, int 1, int r)
    if (1 == tree[m].left \&\& r == tree[m].right) {
        return tree[m].max;
    int mid = (tree[m].left + tree[m].right) >> 1;
    if (r \leq mid) {
        return queryMax(m << 1, 1, r);
    else if (1 > mid) {
        return queryMax (m << 1 | 1, 1, r);
    return max(queryMax(m << 1, 1, mid), queryMax(m << 1 | 1, mid + 1, r));
}
build(1, 1, n);
update(1, a, b);
query (1, a, b);
2)区间更新
/*
    /16/11/06ztx/
typedef long long 11;
const int maxn = 100010;
int t, n, q;
11 anssum;
struct node{
    11 1, r;
    11 addv, sum;
} tree [maxn<<2];
void maintain(int id) {
    if(tree[id].1 >= tree[id].r)
    tree[id].sum = tree[id << 1].sum + tree[id << 1|1].sum;
}
void pushdown(int id) {
    if(tree[id].1 >= tree[id].r)
        return ;
    if(tree[id].addv) {
        int tmp = tree[id].addv;
        tree[id<<1].addv += tmp;
```

```
tree[id << 1 | 1]. addv += tmp;
         tree[id << 1].sum += (tree[id << 1].r - tree[id << 1].1 + 1)*tmp;
         tree[id << 1|1].sum += (tree[id << 1|1].r - tree[id << 1|1].1 + 1)*tmp;
         tree[id].addv = 0;
}
void build(int id, 11 1, 11 r) {
    tree[id].1 = 1;
    tree[id].r = r;
    tree[id].addv = 0;
    tree[id].sum = 0;
    if(1==r) {
         tree[id].sum = 0;
        return :
    11 \text{ mid} = (1+r) >> 1;
    build(id<<1, 1, mid);
    build(id<<1|1, mid+1, r);
    maintain(id);
}
void updateAdd(int id, 11 1, 11 r, 11 val) {
    if(tree[id].1 \ge 1 \&\& tree[id].r \le r)
    {
         tree[id].addv += val;
         tree[id].sum += (tree[id].r - tree[id].1+1)*val;
        return ;
    }
    pushdown(id);
    11 mid = (tree[id].1+tree[id].r)>>1;
    if(1 \le mid)
         updateAdd(id<<1, 1, r, val);
    if(mid < r)
         updateAdd(id<<1|1,1,r,va1);
    maintain(id);
void query(int id, 11 1, 11 r) {
    if(tree[id].1 >= 1 \&\& tree[id].r <= r) {
         anssum += tree[id].sum;
        return ;
    }
    pushdown(id);
    11 \text{ mid} = (\text{tree[id]}.1 + \text{tree[id]}.r) >> 1;
    if(1 \le mid)
         query (id<<1, 1, r);
```

```
if(mid < r)
        query (id<<1 | 1, 1, r);
    maintain(id);
}
int main() {
    scanf ("%d", &t);
    int kase = 0;
    while (t--) {
        scanf ("%d %d", &n, &q);
        build(1, 1, n);
        int id;
        11 x, y;
        11 val;
        printf("Case %d:\n", ++kase);
        while (q--) {
             scanf ("%d", &id);
             if(id==0) {
                 scanf ("%11d %11d %11d", &x, &y, &val);
                 updateAdd(1, x+1, y+1, val);
             }
             else{
                 scanf ("%11d %11d", &x, &y);
                 anssum = 0;
                 query (1, x+1, y+1);
                 printf("%11d\n", anssum);
    return 0;
}
30.树状数组
/*
    /16/11/06ztx/
*/
#include<iostream>
#include<cstdio>
#include<cstring>
#include<string>
#include<cmath>
using namespace std;
typedef long long 11;
const int maxn = 50005;
```

```
int a[maxn];
int n;
int lowbit(const int t) {
   return t & (-t);
}
void insert(int t, int d) {
    while (t \le n) {
        a[t] += d;
       t = t + lowbit(t);
   }
}
11 getSum(int t) {
    11 \text{ sum} = 0;
    while (t > 0) {
        sum += a[t];
       t = t - lowbit(t);
    return sum;
}
int main() {
    int t, k, d;
    scanf("%d", &t);
    k=1;
    while (t--) {
        memset(a, 0, sizeof(a));
        scanf("%d", &n);
        for (int i = 1; i \le n; ++i) {
            scanf("%d", &d);
            insert(i, d);
        string str;
        printf("Case %d:\n", k++);
        while (cin >> str) {
            if (str == "End") break;
            int x, y;
            scanf("%d %d", &x, &y);
            if (str == "Query")
                printf("%11d\n", getSum(y) - getSum(x - 1));
            else if (str == "Add")
                insert(x, y);
            else if (str == "Sub")
                insert(x, -y);
        }
```

```
return 0;
}
```

其他

31.中国剩余定理(孙子定理)

```
/*
    /16/11/06ztx/
*/
int CRT(int a[], int m[], int n) {
    int M = 1;
    int ans = 0;
    for (int i=1; i \le n; i++)
        M = m[i];
    for(int i=1; i<=n; i++) {
        int x, y;
        int Mi = M / m[i];
        extend_Euclid(Mi, m[i], x, y);
        ans = (ans + Mi * x * a[i]) % M;
    if (ans < 0) ans += M;
    return ans;
}
void extend_Euclid(int a, int b, int &x, int &y) {
    if(b == 0) {
        X = 1;
        y = 0;
        return;
    extend_Euclid(b, a % b, x, y);
    int tmp = x;
    X = y;
    y = tmp - (a / b) * y;
}
```