免费模板 ~~~

ACM Standard Code Library

Huang Wei

Software Engineering

Computer and Software College

Hangzhou Dianzi University

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ACM 算法模板集

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第一章 常用函数和 STL

一. 常用函数

```
#include <stdio.h>
int getchar( void ); // 读取一个字符 ,一般用来去掉无用字符
char *gets( char *str ); // 读取一行字符串
#include <stdlib.h>
void * malloc( size_t size ); // 动态内存分配 , 开辟大小为 size 的空间
void qsort( void *buf, size_t num, size_t size, int (*compare)( const void *, const
void *) );
                     // 快速排序
Sample:
int compare_ints( const void * a, const void * b)
{
int * arg1 = (int *) a; int * arg2 = (int *) b;
if ( *arg1 < *arg2 ) return -1;
else if ( *arg1 == *arg2 ) return 0;
else return 1;
}
int array[] = \{ -2, 99, 0, -743, 2, 3, 4 \}; int array\_size = 7;
qsort( array, array_size, sizeof (int), compare_ints );
#include <math.h>
// 求反正弦 ,arg [-1, 1], 返回值 [-pi/2, +pi/2]
double asin( double arg );
// 求正弦 , arg 为弧度 , 弧度 = 角度 *Pi/180.0, 返回值 [-1, 1]
double sin( double arg );
// 求 e 的 arg 次方
double exp( double arg );
// 求 num 的对数,基数为 e
double log( double num );
// 求 num 的根
double sqrt( double num );
// 求 base 的 exp 次方
double pow( double base, double exp );
#include <string.h>
// 初始化内存 ,常用来初始化数组
void * memset( void * buffer, int ch, size_t count );
memset( the_array, 0, sizeof (the_array) );
//printf 是它的变形 ,常用来将数据格式化为字符串
int sprintf( char *buffer, const char *format,
         "%d%d", 123, 4567); //s="1234567"
sprintf(s,
```

```
是它的变形 ,常用来从字符串中提取数据
//scanf
int sscanf( const char *buffer, const char *format,
                                                    ... );
Sample:
char result[100]="24
                      hello",
                              str[100];
                                               int num;
               "%d %s", num,str
        result,
                                  ); //num=24;str="hello"
// 字符串比较 , 返回值 <0 代表 str1<str2,
                                       =0 代表 str1=str2,
                                                          >0 代表 str1>str2
int strcmp( const char *str1,
                              const char *str2 );
```

二. 常用 STL

[标准 container 概要]

vector<T> 大小可变的向量 ,类似数组的用法 ,容易实现删除

list<T> 双向链表

priority_queue<T> 优先队列 , empty(), top(), pop(), push()

set<T> 集合

map<key,val> 关联数组 ,常用来作 hash 映射

[标准 algorithm 摘录]

for_each() 对每一个元素都唤起(调用)一个函数

 find()
 查找第一个能与引数匹配的元素

 replace()
 用新的值替换元素 , O(N)

 copy()
 复制 (拷贝) 元素 , O(N)

remove() 移除元素 reverse() 倒置元素

sort() 排序,O(N log(N))

partial_sort() 部分排序 binary_search() 二分查找

merge() 合并有序的序列 ,O(N)

[C++ String 摘录]

 copy()
 从别的字符串拷贝

 empty()
 判断字符串是否为空

 erase()
 从字符串移除元素

find() 查找元素 insert() 插入元素 length() 字符串长度 replace() 替换元素 substr() 取子字符串 swap() 交换字符串

第二章 重要公式与定理

Fibonacci Number

0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, 233, 377, 610 ... Formula:

$$F_0 = 1$$

$$F_1 = 1$$

$$F_{i} = F_{i-1} + F_{i-2}$$

$$F_{n} = \frac{(1+5)^{n} - (1-5)^{n}}{2^{n} 5} = \frac{?}{?} \frac{1}{3} \frac{?}{1+5} \frac{5}{?} \frac{?}{?}$$

Lucas Number

1, 3, 4, 7, 11, 18, 29, 47, 76, 123...

Formula:

$$L_{n} = \stackrel{?}{3} + 5 \stackrel{?}{3} + \stackrel{?}{3} + \stackrel{5}{3} \stackrel{?}{2} \stackrel{?}{3} \stackrel{?}$$

Catalan Number

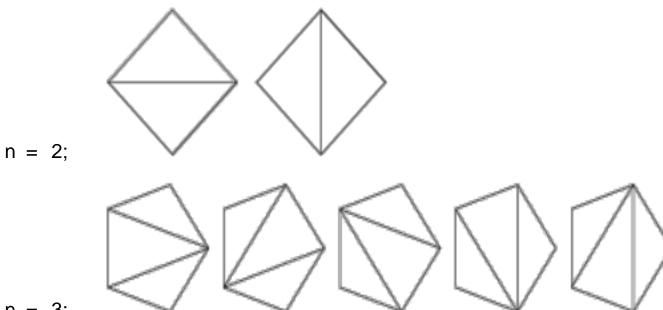
1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012... Formula:

$$Cat_n = \frac{C(2n,n)}{n+1}$$

$$Cat_n = \frac{Cat_n * Cat_{n-1-i}}{cat_n * Cat_{n-1-i}}$$

Application:

1) 将 n + 2 边形沿弦切割成 n 个三角形的不同切割数 Sample:

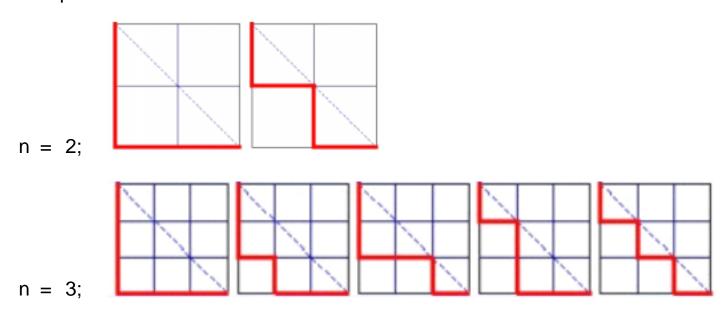


n = 3;

2) n + 1 个数相乘 , 给每两个元素加上括号的不同方法数 Sample:

- 3) n 个节点的不同形状的二叉树数 (严《数据结构》 P.155)
- 4) 从 n * n 方格的左上角移动到右下角不升路径数

Sample:



4. Stirling Number(Second Kind)

S(n, m) 表示含 n 个元素的集合划分为 m 个集合的情况数 或者是 n 个有标号的球放到 m 个无标号的盒子中 ,要求无一为空 ,其不同的方案数

Formula:

$$S_{n,m} ?0 (m = 0 || n < m) ? ?S_{n-1,m-1} + m xS_{n-1,m} (n > m 1)$$

$$S_{n,m} = \frac{1}{m!} \sum_{i=0}^{m} (-1)^{i} \times C(m,i) \times (m-i)^{n}$$

Special Cases:

$$S_{n,0} = 0$$

$$S_{n,1} = 1$$

$$S_{n,2} = 2^{n-1} - 1$$

$$S_{n,3} = \frac{1}{6}(3^n - 3 \times 2^n + 3)$$

$$S_{n,n-1}=C\left(n,2\right)$$

$$S_{n,n} = 1$$

5. Bell Number

n 个元素集合所有的划分数

Formula:

$$B_n = \sum_{i=0}^n \, S_{n,i}$$

6. Stirling's Approximation

$$n! = 2 n ? n ?^{n} ?$$

7. Sum of Reciprocal Approximation

EulerGamma = 0.5 7721566490153286060651209;

$$\int_{i=1}^{n} \frac{1}{i} = \ln(n) + \text{EulerGamma (n)}$$

8. Young Tableau

Young Tableau(杨式图表) 是一个矩阵,它满足条件:如果格子 [i, j] 没有元素,则 [i+1, j] 也一定没有元素如果格子 [i, j] 有元素 a[i, j],则 [i+1, j] 要么没有元素 ,要么 a[i+1, j] > a[i, j]

Y[n] 代表 n 个数所组成的杨式图表的个数

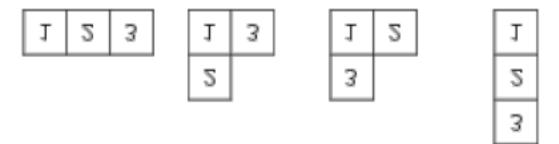
Formula:

 $Y_1 = 1$

 $Y_2 = 2$

$$Y_n = Y_{n-1} + (n-1) x Y_{n-2} (n > 2)$$

Sample:



n = 3;

9. 整数划分

将整数 n 分成 k 份, 且每份不能为空, 任意两种分法不能相同

1) 不考虑顺序

2) 考虑顺序

dp[i][j] = dp[i-k][j-1]; (k=1..i)

3) 若分解出来的每个数均有一个上限 m dp[i][j] = dp[i-k][j-1]; (k=1..m)

10. 错排公式

$$D_1 = 0$$

$$D_2 = 1$$

$$D_n = (n-1) \times (D_{n-1} + D_{n-2})$$

11. 三角形内切圆半径公式

$$p = \frac{a+b+c}{2}$$

$$s = p(p-a)(p-b)(p-c)$$

$$r = \frac{2s}{a+b+c}$$

12. 三角形外接圆半径公式

$$R = \frac{abc}{4s}$$

13. 圆内接四边形面积公式

$$p = a+b+c+d$$

 2
 $s = (p-a)(p-b)(p-c)(p-d)$

14. 基础数论公式

1) 模取幂

$$a^n\%b = ((((a\%b)^* a)\%b)...)\%b$$

2) n 的约数的个数

若 n 满足 n = $p_1^{n_1}$ + $p_2^{n_2}$ + ... + $p_m^{n_m}$, 则 n 的约数的个数为

$$(n_1 + 1)(n_2 + 1)...(n_m + 1)$$

第三章 大数模板

```
typedef int hugeint;
// 应不大于,以防乘法时溢出
const int Base = 1000;
const int Capacity = 1000;
struct xnum
    int Len;
    int Data[Capacity];
    xnum() : Len(0) {}
    xnum( const xnum& V) : Len(V.Len) {
        memcpy(Data, V.Data, Len * sizeof *Data);
    }
    xnum( int V) : Len(0) {
        for (; V > 0; V \neq Base) Data[Len++] = V % Base;
    }
    xnum( char S[]);
    xnum& operator =( const xnum& V) {
        Len = V.Len;
                        V.Data, Len * sizeof *Data);
        memcpy(Data,
        return * this;
    }
    int & operator []( int Index) { return Data[Index];
    int operator []( int Index) const { return Data[Index];
                                                             }
    void print(){
        printf( "%d" ,Len==0?0:Data[Len-1]);
        for (int i=Len-2;i>=0;i--)
            for (int j=Base/10;j>0;j/=10)
                 printf( "%d" ,Data[i]/j%10);
};
xnum::xnum( char S[])
{
    int I, J;
    Data[Len = 0] = 0;
    J = 1;
    for (I = strlen(S)-1; I>=0; I--) {
         Data[Len] += (S[I] - '0') * J;
         J *= 10;
         if (J >= Base) J = 1, Data[++Len] = 0;
```

```
if (Data[Len] > 0) Len++;
}
int compare( const xnum& A, const xnum& B)
{
    int I;
    if (A.Len != B.Len) return A.Len > B.Len ? 1 : -1;
    for (I = A.Len - 1; I >= 0 && A[I] == B[I]; I--);
    if (I < 0) return 0;
    return A[I] > B[I] ? 1 : -1;
}
xnum operator +( const xnum& A, const xnum& B)
{
    xnum R;
    int I;
    int Carry = 0;
    for (I = 0; I < A.Len || I < B.Len || Carry > 0; I++)
        if (I < A.Len) Carry += A[I];
        if (I < B.Len) Carry += B[I];
        R[I] = Carry % Base;
        Carry /= Base;
    R.Len = I;
    return R;
}
xnum operator -( const xnum& A, const xnum& B)
{
    xnum R;
    int Carry = 0;
    R.Len = A.Len;
    int I;
    for (I = 0; I < R.Len; I++)
        R[I] = A[I] - Carry;
        if (I < B.Len) R[I] -= B[I];
        if (R[I] < 0) Carry = 1, R[I] += Base;
        else Carry = 0;
    }
    while (R.Len > 0 \&\& R[R.Len - 1] == 0) R.Len--;
    return R;
```

```
}
xnum operator *( const xnum& A, const int B)
{
    int I;
    if (B == 0) return 0;
    xnum R;
    hugeint Carry = 0;
    for (I = 0; I < A.Len || Carry > 0; I++)
                                            * B;
        if (I < A.Len) Carry += hugeint(A[I])
        R[I] = Carry % Base;
        Carry /= Base;
    R.Len = I;
    return R;
}
xnum operator *( const xnum& A, const xnum& B)
{
    int I;
    if (B.Len == 0) return 0;
    xnum R;
    for (I = 0; I < A.Len; I++)
        hugeint Carry = 0;
        for (int J = 0; J < B.Len || Carry > 0; J++)
        {
            if (J < B.Len) Carry += hugeint(A[I]) * B[J];
            if (I + J < R.Len) Carry += R[I + J];
            if (I + J \ge R.Len) R[R.Len++] = Carry % Base;
            else R[I + J] = Carry \% Base;
            Carry /= Base;
        }
    return R;
}
xnum operator /( const xnum& A, const int B)
{
    xnum R;
    int I;
    hugeint C = 0;
    for (I = A.Len - 1; I >= 0; I--)
```

```
{
        C = C * Base + A[I];
        R[I] = C / B;
        C %= B;
    }
    R.Len = A.Len;
    while (R.Len > 0 \&\& R[R.Len - 1] == 0) R.Len--;
    return R;
}
//div
xnum operator /( const xnum& A, const xnum& B)
{
    int I;
    xnum R, Carry = 0;
    int Left, Right, Mid;
    for (I = A.Len - 1; I >= 0; I--)
    {
        Carry = Carry * Base + A[I];
        Left = 0;
        Right = Base - 1;
        while (Left < Right)
            Mid = (Left + Right + 1) / 2;
            if (compare(B * Mid, Carry) <= 0) Left = Mid;
            else Right = Mid - 1;
        }
        R[I] = Left;
        Carry = Carry - B * Left;
    R.Len = A.Len;
    while (R.Len > 0 \&\& R[R.Len - 1] == 0) R.Len--;
    return R;
}
//mod
xnum operator %( const xnum& A, const xnum& B)
{
    int I;
    xnum R, Carry = 0;
    int Left, Right, Mid;
    for (I = A.Len - 1; I >= 0; I--)
        Carry = Carry * Base + A[I];
```

```
Left = 0;
        Right = Base - 1;
        while (Left < Right)
        {
            Mid = (Left + Right + 1) / 2;
            if (compare(B * Mid, Carry) <= 0) Left = Mid;
            else Right = Mid - 1;
        }
        R[I] = Left;
        Carry = Carry - B * Left;
    R.Len = A.Len;
    while (R.Len > 0 \&\& R[R.Len - 1] == 0) R.Len--;
    return Carry;
}
istream& operator >>(istream& In, xnum& V)
{
    char Ch;
    for (V = 0; In \gg Ch;)
        V = V * 10 + (Ch - '0');
        if (cin.peek() <= ' ') break ;</pre>
    }
    return In;
}
ostream& operator <<(ostream& Out, const xnum& V)
{
    int I;
    Out << (V.Len == 0 ? 0 : V[V.Len - 1]);
    for (I = V.Len - 2; I >= 0; I--)
        for (int J = Base / 10; J > 0; J /= 10)
            Out << V[I] / J % 10;
        return Out;
}
xnum gcd(xnum
                  a,xnum b)
{
    if (compare(b,0)==0)
                           return a;
    else return gcd(b,a%b);
}
int div( char *A, int B)
```

```
{
    int I;
    int C = 0;
    int Alen=strlen(A);
    for (I = 0; I < Alen; I++)
        C = C * Base + A[I]- '0';
        C %= B;
    return C;
}
xnum C(int n, int m)
{
    int i;
    xnum sum = 1;
    for (i = n; i >= n-m+1; i --)
        sum = sum*i;
    for (i = 1; i \le m; i ++)
        sum = sum/i;
    return sum;
}
#define MAXN 9999
#define DLEN 4
class BigNum {
private :
    int a[1000]; // 可以控制大数的位数
    int len; // 大数长度
public :
    BigNum()
               \{len = 1; memset(a,0,
                                       sizeof (a));}
    BigNum(
             const int);
    BigNum(
             const char *);
             const BigNum
    BigNum(
                            &);
    BigNum
             & operator =( const BigNum
                                          &);
    BigNum
                      +( const BigNum
             operator
                                        &) const;
    BigNum
                      -( const BigNum
             operator
                                        &) const;
                               BigNum
    BigNum
             operator
                      *( const
                                        &) const;
    BigNum
             operator /( const int
                                    &) const;
    BigNum
             operator ^( const int
                                     &) const;
          operator %( const int
                                   &) const;
    int
            operator >( const BigNum
    bool
                                       & T) const;
```

```
void print();
};
                     const int b) {
BigNum::BigNum(
    int c,d = b;
    len = 0;
    memset(a,0,
                   sizeof (a));
    while (d > MAXN) {
         c = d - (d / (MAXN + 1)) * (MAXN + 1);
         d = d / (MAXN + 1); a[len++]
                                               = c;
    a[len++]
                = d;
}
BigNum::BigNum(
                     const char *s) {
    int t,k,index,l,i;
    memset(a,0,
                   sizeof (a));
    l=strlen(s);
    len=I/DLEN;
    if (I%DLEN)len++;
    index=0;
    for (i=l-1;i>=0;i-=DLEN)
         t=0;k=i-DLEN+1;
         if (k<0)k=0;
         for (int j=k;j <=i;j++)
             t=t*10+s[j]-
                            '0';
         a[index++]=t;
BigNum::BigNum(
                     const BigNum & T) : len(T.len)
    int i;
    memset(a,0,
                   sizeof (a));
    for (i = 0 ; i < len ; i++)a[i] = T.a[i];
}
BigNum & BigNum:: operator =( const BigNum & n) {
    len = n.len;
    memset(a,0, sizeof (a));
    int i;
    for (i = 0 ; i < len ; i++)
         a[i] = n.a[i];
    return * this;
}
BigNum BigNum:: operator +( const BigNum & T) const {
    BigNum t(* this );
    int i,big; // 位数
```

```
big = T.len > len ? T.len : len;
    for (i = 0 ; i < big ; i++) {
        t.a[i] +=T.a[i];
         if (t.a[i] > MAXN) {
             t.a[i + 1]++;
             t.a[i] -=MAXN+1;
        }
    }
               != 0) t.len = big + 1;
    if (t.a[big]
    else t.len = big;
    return t;
}
BigNum BigNum:: operator -( const BigNum & T) const {
    int i,j,big;
    bool flag;
    BigNum t1,t2;
    if (* this >T) {
        t1=* this;
         t2=T;
         flag=0;
    } else {
         t1=T;
         t2=* this;
         flag=1;
    big=t1.len;
    for (i = 0 ; i < big ; i++) {
         if (t1.a[i] < t2.a[i]) {
             j = i + 1;
             while (t1.a[j] == 0) j++;
             t1.a[j--]--;
             while (j > i) t1.a[j--] += MAXN;
             t1.a[i] += MAXN + 1 - t2.a[i];
         } else t1.a[i] -= t2.a[i];
    }
    t1.len = big;
    while (t1.a[len - 1] == 0 \&\& t1.len > 1) {
         t1.len--;
         big--;
    if (flag)t1.a[big-1]=0-t1.a[big-1];
    return t1;
}
BigNum BigNum:: operator *( const BigNum & T) const {
```

```
BigNum ret;
    int i,j,up;
    int temp,temp1;
    for (i = 0 ; i < len ; i++)
        up = 0;
        for (j = 0 ; j < T.len ; j++) {
             temp = a[i] * T.a[j] + ret.a[i + j] + up;
             if (temp > MAXN) {
                 temp1 = temp - temp / (MAXN + 1) * (MAXN + 1);
                 up = temp / (MAXN + 1);
                 ret.a[i + j] = temp1;
             } else {
                 up = 0;
                 ret.a[i + j] = temp;
             }
        }
        if (up != 0)
            ret.a[i + j] = up;
    }
    ret.len = i + j;
    while (ret.a[ret.len - 1] == 0 \&\& ret.len > 1) ret.len--;
    return ret;
}
BigNum BigNum:: operator /( const int & b) const {
    BigNum ret;
    int i,down = 0;
    for (i = len - 1; i >= 0; i--)
        ret.a[i] = (a[i] + down * (MAXN + 1)) / b;
        down = a[i] + down * (MAXN + 1) - ret.a[i] * b;
    ret.len = len;
    while (ret.a[ret.len - 1] == 0 \&\& ret.len > 1) ret.len--;
    return ret;
}
int BigNum:: operator %( const int & b) const {
    int i,d=0;
    for (i = len-1; i>=0; i--) {
        d = ((d * (MAXN+1))\% b + a[i])\% b;
    return d;
}
BigNum BigNum:: operator ^( const int & n) const {
    BigNum t,ret(1);
    if (n<0)exit(-1);
```

```
if (n==0)
             return 1;
    if (n==1) return * this;
    int m=n;
    while (m>1) {
        t=* this;
         int i;
        for (i=1;i<<1<=m;i<<=1)
             t=t*t;
         }
         m-=i;
         ret=ret*t;
        if (m==1)ret=ret*(* this );
    }
    return ret;
}
bool BigNum:: operator >( const BigNum & T) const {
    int In;
    if (len > T.len) return true;
    else if (len == T.len) {
        ln = len - 1;
        while (a[ln] == T.a[ln] && ln >= 0) ln--;
        if (ln >= 0 \&\& a[ln] > T.a[ln]) return true;
         else return false;
    } else return false;
}
void BigNum::print()
    int i;
    cout << a[len - 1];
    for (i = len - 2 ; i >= 0 ; i--)
         cout.width(DLEN);
         cout.fill( '0');
        cout << a[i];
}
// 读取整数
const int ok = 1;
int get_val( int & ret) {
    ret = 0;
    char ch;
    while ((ch=getchar()) > '9' || ch < '0');
    do {
        ret = ret*10 + ch - '0';
    } while ((ch=getchar()) <= '9' && ch >= '0');
    return ok;
```

```
}
// 带负数
int get_val( int & ret) {
    ret = 0;
    char ch;
    bool neg = false ;
                        > '9' || ch < '0') && ch!= '-');
    while (((ch=getchar())
    if (ch == '-') {
         neg = true ;
        while ((ch=getchar()) > '9' || ch < '0');
    }
    do {
        ret = ret*10 + ch - '0';
    ((ch=getchar()) <= '9' && ch >= '0');
    ret = (neg? -ret : ret);
    return ok;
}
// 读取整数,可判 EOF和 EOL
const int eof = -1;
const int eol = -2;
int get_val( int & ret) {
    ret = 0;
    char ch;
    while (((ch=getchar()) > '9' || ch < '0') \&\& ch!=EOF);
    if (ch == EOF) return eof;
    do {
        ret = ret*10 + ch - '0';
    ((ch=getchar()) <= '9' && ch >= '0');
    if (ch == '\n') return eol;
    return ok;
}
// 读取浮点数
int get_val( double & ret) {
    ret = 0;
    double base = 0.1;
    char ch;
    bool dot = false , neg = false ;
                        > '9' || ch < '0') && ch != '.' && ch != '-');
    while (((ch=getchar())
    if (ch == '-') {
        neg = true ;
                              > '9' || ch < '0') && ch != '.' && ch != '-');
        while (((ch=getchar())
```

```
}
    do {
        if (ch == '.') {
            dot = true ;
            continue ;
        }
        if (dot) {
            ret += (ch- '0') * base;
            base *= 0.1;
        } else ret = ret*10 + (ch- '0' );
    (((ch=getchar()) <= '9' && ch >= '0') || ch == '.');
    ret = (neg? -ret : ret);
    return ok;
}
typedef long long LL;
//LL MultiMod(LL a, LL b, LL c) {
// if (b)
// return (a * (b & 1) % c + (MultiMod(a, b >> 1, c) << 1)) % c;
// return 0;
//}
LL MultiMod(LL a, LL b, LL c) {
    LL ret = 0, d = a;
    for (; b; b >>= 1, d <<= 1, d %= c)
        if ((b & 1))
           ret = (ret + d) \% c;
    return ret;
}
// 128-bits integer's power with mod in O(64*LogN)
LL ModPower(LL
                 base, LL exp, LL mod) {
    LL ret = 1;
    for (; exp; exp >>= 1, base = MultiMod(base,
                                                      base, mod))
        if ((exp & 1))
             ret = MultiMod(ret,
                                 base, mod);
    return ret;
}
```

第四章 数论算法

1. Greatest Common Divisor 最大公约数

```
int GCD( int x, int y)
   int t;
   while (y > 0) {
      t = x \% y;
       x = y;
       y = t;
   }
   return x;
}
     Prime
                素数判断
bool is_prime( int u)
{
   if (u == 0 || u == 1) return false;
   if (u == 2)
                         return true;
   if (u\%2 == 0)
                              return false;
   for (int i=3; i \le sqrt(u)
                             ;i+=2)
       if (u%i==0)
                       return false;
   return true;
}
                          素数筛法
3.
     Sieve Prime
const int M = 1000; // M : size
bool mark[M];
             // true : prime number
void sieve_prime()
{
   memset(mark, true, sizeof (mark));
   mark[0] = mark[1] = false;
   for (int i=2; i \le sqrt(M) ;i++) {
       if (mark[i]) {
           for ( int j=i^*i; j < M ; j+=i)
              mark[j]
                       = false ;
       }
```

}

4. Module Inverse 模逆元

5. Extended Euclid 扩展欧几里德算法

```
// 如果 GCD(a,b)=d, 则存在 x, y, 使 d=ax+by

// extended_euclid(a, b) = ax+by

int extended_euclid( int a, int b, int &x, int &y)

{
    int d;
    if (b == 0) {x = 1;        y = 0;        return a;}
    d = extended_euclid(b, a % b, y, x);
    y -= a / b * x;
    return d;
}
```

6. Modular Linear Equation 模线性方程(同余方程)

```
// 如果 GCD(a, b) 不能整除 c, 则 ax + by = c 没有整数解
      b (mod n) n > 0
// 上式等价于二元一次方程 ax - ny = b
{
   int d, x, y, x0, gcd;
   // 可以减少扩展欧几里德溢出的可能
   gcd = GCD(a, n);
   if (b%gcd != 0) {
       cout << "no solution" << endl;
       return ;
   }
   a /= gcd; b /= gcd; n /= gcd;
   d = extended_euclid(a, n, x, y);
   if ( b\%d == 0) {
      x0 = (x^*(b/d)) \% n; // x0 : basic solution
      int ans = n; // min x = (x0\%(n/d))
                                   + (n/d)) % (n/d)
      for (int i=0; i < d; i++) {
         ans = (x0 + i*(n/d)) % n;
         cout << ans << endl;
```

```
}
       }
              cout << "no solution" << endl;
       else
    }
    7.
         Chinese Remainder Theorem 中国余数定理
    // x b[i] (mod w[i]), i [1, len-1]
    // 前提条件 w[i] > 0,且 w[] 中任意两个数互质
    int chinese_remainder( int b[], int w[], int len)
    {
       int i, d, x, y, m, n, ret;
       ret = 0; n = 1;
       for (i=0; i < len ; i++)   n *= w[i];
       for (i=0; i < len ; i++)
          m = n / w[i] ;
          d = extended_euclid(w[i], m, x, y);
          ret = (ret + y*m*b[i]) % n;
       }
       return (n + ret%n) % n;
    }
// m r[i] (mod a[i])
// a[i] 可以不互素
// -1 表示无解
Pku 2891 Strange Way to Express Integers
      A1 (mod B1), C A2 (mod B2).
假设C
令C = A1 + X1B , 那么 X1B1 A2 - A1 (mod B2)。
用扩展欧几里德算法求出 X1,也就求出 C。
令B = lcm(B1, B2) , 那么上面两条方程就可以被 C ' C (mod B) 代替。
迭代直到只剩下一条方程。
*/
LL chinese_remainder2()
    int i, j;
    if (n == 1)
        return r[0];
    LL m, x, apre;
    x = modular_linear_equation(a[0],
                                     r[1]-r[0],
                                               a[1]);
    if (x == -1)
        return -1;
    m = x*a[0] + r[0];
    apre = LCM(a[0], a[1]);
```

```
for (i=2; i<n; i++)
    {
        x = modular_linear_equation(apre,
                                       r[i]-m,
                                                     a[i]);
        if (x == -1)
            return -1;
        m = x*apre + m;
        apre = LCM(apre, a[i]);
    }
    return m;
}
    8. Euler Function
                                  欧拉函数
// 求 1..n-1 中与 n 互质数的个数
int euler( int n)
{
    int ans = 1;
    int i;
    for (i=2; i*i <= n ; i++) {
        if (n\%i == 0) {
            n /= i;
            ans *= i-1;
            while (n%i == 0) {
                  n /= i;
                  ans *= i;
    if (n > 1) {
        ans *= n-1;
    }
    return ans;
}
          Farey 总数
    9.
// 求MAX 以内所有 Farey 的总数
const int MAX = 1000100;
int n;
bool num[1100];
                 //sqrt(MAX)
int prime[1100],
                 total;
__int64
        f[MAX],
                 inc[MAX];
void cal_prime()
                 {
    int i,j;
```

```
memset(num,
                     false, sizeof (num));
    total = 0;
    for (i=2;i<1100;i++)
         if (!num[i])
              prime[total
                            ++] = i;
              j = i+i;
              while (j < 1100) {
                   num[j] = true ;
                   j += i;
              }
         }
}
void cal_farey()
    int i,j,k;
    inc[1] = 1;
    for (i=2;i<MAX;i++)
         for (j=0;j<total;j++)
                                  {
                              == 0) {
              if (i%prime[j]
                   k = i / prime[j];
                   if(k\%prime[j] == 0) inc[i] = inc[k] * prime[j];
                   else inc[i] = inc[k] * (prime[j]
                                                       -1);
                   break ;
              }
         if (j == total) inc[i] = i -1;
    f[1] = 0;
    for (i=2;i<MAX;i++)
                            f[i] = f[i-1] + inc[i];
}
int main() {
     cal_prime();
    cal_farey();
    while (scanf( "%d" , &n), n) {
         printf( "%I64d\n" , f[n]);
}
     10. Farey 序列构造
// 构造 5000 以内的 Farey 序列
```

```
// 构造 5000 以内的 Farey 序列 const int MAX = 8000000; int total;
```

```
int n,k;
int farey[2][MAX];
void make_farey_seq( int x1, int y1, int x2, int y2)
{
    if (x1+x2 > n || y1+y2 > n) return;
    make\_farey\_seq(x1, y1,x1+x2,
                                      y1+y2);
    total ++;
    farey[0][total]
                    = x1+x2;
    farey[1][total]
                    = y1+y2;
    make\_farey\_seq(x1+x2, y1+y2,x2,y2);
}
int main() {
    int t;
    scanf( "%d %d" , &n, &t);
    total = 1;
    farey[0][1]
                 = 0;
    farey[1][1]
                 = 1;
    make_farey_seq(0,1,1,1);
    farey[0][total+1]
                   = 1;
    farey[1][total+1]
                     = 1;
    total ++;
    while (t --) {
         scanf( "%d" , &k);
         if (k > total) puts( "No Solution" );
         else printf( "%d/%d\n"
                               , farey[0][k],
                                               farey[1][k]);
}
     11. Miller_Rabbin 素数测试, Pollard_rho 因式分解
typedef
       ___int64
                 164;
const char * pformat = "%l64d"
164 big_rand(164 m) {
    I64 x = rand();
    x *= rand();
    if (x < 0) x = -x;
    return x %= m;
}
// x*y % n
I64 mod_mul(I64 x, I64 y, I64 n) {
    if (x == 0 || y == 0) return 0;
    return ( ((x&1)*y)%n + (mod_mul(x>>1,y,n)<<1)%n
                                                       ) % n;
}
// x^y % n
I64 mod_exp(I64 x, I64 y, I64 n) {
```

```
164 \text{ ret} = 1;
    while (y) {
         if (y\&1) ret = mod_mul(ret,x,n);
         x = mod_mul(x,x,n);
         y >>= 1;
    return ret;
}
bool Miller_Rabbin(I64 n) { // O(times * (log N)^3)
    l64 i,j,x,m,k;
    if (n == 2) return true;
    if (n < 2 || !(n&1)) return false;
    m = n - 1; k = 0;
    while (!(m&1)) m >>= 1, k ++; // binary scan
    for (i=0;i<4;i++) { // test times
         x = big_rand(n-2)
                          +2;
         x = mod_exp(x,m,n);
         if (x == 1) continue;
         for (j=0;j< k;j++) {
             if (x == n-1) break;
             x = mod_mul(x,x,n);
         if (j >= k) return false;
    }
    return true;
    /*lrj P.218
    for(i=0;i<20;i++)
    x = big_rand(n-2) +2;
    if(mod\_exp(x,n-1,n) = 1) return false;
    return true;
    */
}
I64 gcd(I64 x, I64 y) {
    if (x > y) std::swap(x,y);
    while (x) {
        164 t = y \% x;
        y = x;
        x = t;
    return y;
}
I64 func(I64 x, I64 m) {
    return (mod_mul(x,x,m)+1)
                                 % m;
```

```
}
164 Pollard(164 n) {
    if (Miller_Rabbin(n))
                             return n;
    if (!(n&1))
                 return 2;
     I64 i,x,y,ret;
    i = 1;
    while (true ) {
         x = i ++;
          y = func(x,n);
          ret = gcd(y-x,n);
          while (ret == 1) {
              x = func(x,n);
              y = func(func(y,n),n);
              ret = gcd((y-x+n)\%n,n)
                                             % n;
          }
          if (0 < ret && ret < n) return
                                            ret;
}
l64 factor[100],
                    nfac, minfac;
void cal_factor(I64
                       n) {
     164 x = Pollard(n);
    if (x == n) {
          //factor[nfac
                         ++] = X;
          minfac = min(minfac,x);
          return ;
     cal_factor(x);
     cal_factor(n/x);
}
void print_factor(I64
                         n) {
     164 i;
     nfac = 0;
     cal_factor(n);
     std::sort(factor,factor
                          + nfac);
    for (i=0;i<nfac;i++)
          if (i > 0) putchar( '');
          printf(pformat,factor[i]);
     puts( "");
}
const I64 lim = 100000;
int main() {
     l64 n,t,i;
    srand(( unsigned )time(NULL));
```

```
scanf(pformat,&t);
     while (t --) {
          scanf(pformat,
                             &n);
                                   puts( "Prime" );
          if (Miller_Rabbin(n))
          else {
               if (!(n&1))
                            puts( "2" );
               else {
                     for (minfac=3; minfac < lim && n % minfac ;minfac+=2)
                     if (minfac >= lim) {
                          164 \text{ rn} = \text{sqrt}(1.0^*\text{n});
                          if (rn * rn == n) {
                               minfac = rn;
                               cal_factor(rn);
                         }
                          else {
                               minfac = n;
                               cal_factor(n);
                          }
                     printf(pformat,minfac);
                     puts( "");
               }
          }
}
      12.
```

第五章 图论算法

1. 最小生成树 (Kruscal 算法)

```
Name :
                           最小生成树 (Kruscal
                                               算法)
    Function
                          ZJU 1203 Swordfish
                                                O(E*LogE)
    Description
#include
          <iostream>
          <algorithm>
#include
#include
          <cstdio>
#include
          <cmath>
using namespace
                   std;
struct struct_edges
   int bv,tv; //bv 起点 tv 终点
   double w; // 权值
};
                               // 边集
struct_edges
              edges[10100];
struct struct_a
    double x;
    double
          у;
};
struct_a arr_xy[101];
int point[101],n,e; //n 顶点数,e 边数(注意是无向网络)
double sum;
int kruscal_f1( int point[],
                           int v)
   int i = v;
   while (point[i] > 0)  i = point[i];
   return i;
}
bool UDlesser(struct_edges
                          a, struct_edges
                                               b)
```

```
{ return a.w < b.w;}
void kruscal() // 只需要准备好 n,e,递增的边集 edges[] 即可使用
   int v1,v2,i,j;
   for (i=0; i<n ;i++) point[i]=0;
   i = j = 0;
   while (j<n-1 && i<e) {
       v1 = kruscal_f1(point,
                             edges[i].bv);
       v2 = kruscal_f1(point,
                             edges[i].tv);
       if (v1 != v2) {
           sum += edges[i].w; // 注意 sum 初始为 0
           point[v1]=v2;
           j++;
       }
       i++;
}
int main()
{
   int k,i,j;
    cin>>n;
    k=0;
    while (n != 0) {
       sum=0;
       k++;
       for (i=0; i< n ; i++)
           cin>>arr_xy[i].x>>arr_xy[i].y;
       e=0;
       for (i=0; i<n ;i++) // 从 0 开始计数
           for (j=i+1; j<n ;j++) // 注意是无向网络
               if (i == j) continue;
               edges[e].bv=i;
               edges[e].tv=j;
               edges[e].w=sqrt((arr_xy[i].x-arr_xy[j].x)*(arr_xy[i].x-arr_xy[j]
.x)+(arr_xy[i].y-arr_xy[j].y)*(arr_xy[i].y-arr_xy[j].y));
               e++;
           }
            sort(edges,edges+e,UDlesser);
                                                // 得到一个递增的边集,注意是从
                                                                                   0 开
始计数
           kruscal();
```

```
#%d:\n",k);
                                        //cout<<"Case
                                                          #"<<k<<":"<<endl;
           printf("Case
                                 distance is: %.2f\n",sum);
                                                                // 输出 sum
           printf("The
                        minimal
           cin>>n;
           if (n != 0) printf("\n");
   }
}
2.
      最小生成树 (Prim
                               算法)
                            最小生成树 (Prim 算法)
    Function
              Name :
                           ZJU 1203 Swordfish
    Description
                                                 O(N^2)
#include
          <iostream>
#include
          <cmath>
          <cstdio>
#include
using namespace
                    std;
double sum, arr_list[101][101],
                                    min;
int i, j, k=0, n;
struct struct_a
   float x;
   float y;
};
struct_a
         arr_xy[101];
struct struct_b
   int point;
   float lowcost;
};
struct_b
         closedge[101];
                            需要准备: n 顶点数 arr_list[][]
                                                          顶点的邻接矩阵也是从
void prim( int n) //prim
                                                                                  0 开
始计数
{
   int i,j,k;
   k=0;
   for (j=0; j< n ; j++) {
       if (j != k) {
```

= k;

closedge[j].lowcost = arr_list[k][j];

closedge[j].point

}

```
}
              closedge[k].lowcost=0;
              for (i=0; i< n ; i++) {
                            min=10000;
                            for (j=0; j< n ; j++) {
                                          if (closedge[j].lowcost
                                                                                                                    != 0 && closedge[j].lowcost < min) {
                                                        k = j;
                                                       min = closedge[j].lowcost;
                            }
                            sum += closedge[k].lowcost; // 不要改成 sum+=min; sum 即为所求值
                            closedge[k].lowcost = 0;
                            for (j=0; j< n; j++)
                                                                                                      {
                                          if (arr_list[k][j]
                                                                                                       < closedge[j].lowcost)
                                                       closedge[j].point
                                                                                                                              = k;
                                                       closedge[j].lowcost = arr_list[k][j];
        arr_list[][]=
                                                           Wij 如果 Vi, Vj 有边
                                               0 如果 i=j
                                               无限大 如果没有边
*/
int main()
              cin>>n;
              while (n != 0) {
                            sum=0;
                            k++;
                            for (i=0; i<n ;i++)
                                          cin>>arr_xy[i].x>>arr_xy[i].y;
                            for (i=0; i< n ; i++)
                                          for (j=0; j<n ;j++) // 得到邻接矩阵 arr_list[][]
                                                           arr\_list[i][j] = arr\_list[j][i] = sqrt((arr\_xy[i].x - arr\_xy[j].x)^*(arr\_x)^*(arr\_x)^*(arr\_x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr\_xy[i].x)^*(arr_xy[i].x)^*(arr_xy[i].x)^*(arr_xy[i].x)^*(arr_xy[i].x)^*(arr_xy[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^*(arr_xx[i].x)^
y[i].x-arr_xy[j].x)+(arr_xy[i].y-arr_xy[j].y)*(arr_xy[i].y-arr_xy[j].y));
                            prim(n);
                                                                                   #"<<k<<":"<<endl;
                            cout<<"Case
                            printf("The
                                                                                                            distance is: %.2f\n",sum);
                                                                          minimal
                            cin>>n;
                                                                          printf("\n");
                            if (n!=0)
              }
```

}

3. 单源最短路径 (Bellman-ford 算法)

```
struct node {
    int e,v;
    node( int a = 0, int b = 0)
         : e(a), v(b) {}
};
vector< vector<node>
                          > path;
int n,p,q;
int dist[1000100];
    SPFA (Shortest Path Faster Algorithm)
    Bellman-Ford 算法的一种队列实现,减少了不必要的冗余计算
    返回值为 false ,说明队列不为空,存在负权环
*/
bool SPFA()
{
    int i,j,k,now,l;
    node next;
     bitset <1000100>
                          vis;
    queue< int > SQ;
    memset(dist,-1,
                      sizeof (dist));
    SQ.push(1);
    vis[1] = true ;
    dist[1] = 0;
    for (i=0;i<=n;i++)
        I = SQ.size();
         if (I == 0) break;
         while (I--) {
             now = SQ.front();
             SQ.pop();
             vis[now]
                       = false ;
             for (j=path[now].size()-1;j>=0;j--)
                  next = path[now][j];
                                         || dist[next.e]
                                                          > dist[now]+next.v)
                  if (dist[next.e]==-1
                      dist[next.e]
                                    = dist[now]+next.v;
                      if (!vis[next.e])
                           SQ.push(next.e);
                           vis[next.e]
                                        = true ;
                  }
             }
```

```
}
    return SQ.empty();
}
          单源最短路径 (Dijkstra
                                         算法)
    4.
                              单源最短路径 (Dijkstra
        Function
                 Name :
                                                    算法)
                             贪心,O(N^2),
                                           不能有负权
        Description
                                 ****/
    int matrix[200][200],n;
                                                30000 表示无限大,即无边.否则为有边,
                                   //matrix[][],
    其值为边的权值
    void Dijkstra( int x, int y) // 起点 Vx 终点 Vy
    {
        int i,j,k,path[40000],mark[40000];
        int min,dist[40000];
        for (i=1;i<=n;i++)
                      {
           mark[i] = 0;
           dist[i] = matrix[x][i];
           path[i] = x;
        }
        mark[x]
                = 1;
        do {
           min=30000;
           k=0;
           for (i=1;i<=n;i++)
               if (mark[i]==0
                             && dist[i]<min)
                  min = dist[i];
                  k = i;
           if (k) {
               mark[k] = 1;
               for (i=1;i<=n;i++)
                  dist[i] = min + matrix[k][i];
                     path[i]
                            = k;
                  }
           }
        } while (k);
                                           的值就是从 Vx 到 Vy 的最短路径值
        cout<<dist[y]<<endl;</pre>
                                 //dist[y]
        // 如果希望得到路径,加入如下代码:
        do {
           cout<<k<<"<--";
           k = path[k];
```

} while (k!=x);

```
cout<<x<endl;
}
      全源最短路径 (Folyd 算法)
5.
                              全源最短路径 (Folyd 算法)
    Function
               Name:
                                 DP, O(N^3)
    Description
                                 ****/
// 初始化
//path[i][j]=j;
void Floyd()
{
   int i,j,k;
   for (k=0;k<vertex_number;k++)
      for (i=0;i<vertex_number;i++)</pre>
         for (j=0;j<vertex_number;j++)
             if ((graph[i][k]==-1)
                                    || (graph[k][j]==-1))
                                                                  continue ;
            if ((graph[i][j]==-1)
                                    || (graph[i][j]
                                                     > graph[i][k]+graph[k][j]))
                            = graph[i][k]+graph[k][j];
                                                               /* 最短路径值 */
                graph[i][j]
                                    /* 最短路径 */
               path[i][j]
                           = k;
      拓扑排序
6.
                              拓扑排序
    Function
               Name
                                 ****/
//degree[]
               每个结点的入度
//f[]
              每个结点所在的层
void Toplogical_sort()
{
    int i,j;
    bool p= true ;
    top=0;
    while (p) {
        p= false ;
        top++;
        for (i=1;i<=n;i++)
```

```
if (degree[i]==0)
                   p= true ;
                   f[i]=top;
           for (i=1;i<=n;i++)
               if (f[i]==top)
                   for (j=1;j<=n;j++)
                      if (map[i][j])
                                       degree[j]--;
                   degree[i]=-1;
        }
       top--;
    7.
          网络预流和最大流
                              最短增广路算法 O(VE^2)
网络中求最大流 Edmonds_Karp
            n 代表网络中节点数 ,第1节点为源点 , 第n 节点为汇点
参数含义:
            net[][] 代表剩余网络 ,0 表示无通路
            path[] 保存增广路径
            neck[] 代表瓶颈,保存增广路径最小容量
返回值:
            最大流量
*/
const int NMAX = 210;
int net[NMAX][NMAX];
int path[NMAX],
                 n;
int bfs()
{
    queue< int > SQ;
    int neck[NMAX],
                     i;
    memset(path,-1,
                     sizeof (path));
    neck[1] = INT_MAX;
    SQ.push(1);
    while (!SQ.empty())
        int now = SQ.front();
        SQ.pop();
        if (now == n) break;
        for (i=1;i<=n;i++)
                           > 0 && path[i]
             if (net[now][i]
                                           == -1) {
                 path[i]
                       = now;
                 neck[i] = min(neck[now],
                                             net[now][i]);
                 SQ.push(i);
```

```
}
    if (path[n] == -1) return -1;
    return neck[n];
}
int Edmonds_Karp()
{
    int now, step;
    int max_flow = 0;
    while ( (step=bfs()) != -1 ) {
        max_flow += step;
        now = n;
        while (now != 1) {
            int pre = path[now];
            net[pre][now]
                          -= step;
            net[now][pre]
                         += step;
            now = pre;
        }
    }
    return max_flow;
}
网络中求最大流 HLPP 高度标号预流推进算法 O(V^2*E^0.5)
            n 代表网络中节点数 ,第0节点为源点 , 第n 节点为汇点
参数含义:
            net[][] 代表剩余网络 ,0 表示无通路
            earn[] 代表各点的盈余
            high[] 代表各点的高度
返回值:
            最大流量
*/
const int NMAX = 110;
int earn[NMAX], net[NMAX][NMAX],
                                   high[NMAX];
int n, m;
queue< int > SQ;
void push( int u, int v) {
    int ex = min(earn[u],
                         net[u][v]);
    earn[u] -= ex;
    net[u][v] -= ex;
            += ex;
    earn[v]
    net[v][u] += ex;
}
void relable( int u) {
```

```
int i, mmin = INT_MAX;
    for (i=0;i<=n;i++)
                      {
                       > 0 \& high[i] >= high[u]) {
          if (net[u][i]
              mmin = min(mmin,
                                      high[i]);
         }
    high[u] = mmin +1;
}
void discharge( int u) {
    int i, vn;
    while (earn[u] > 0) {
          vn = 0;
          for (i=0;i<=n &\& earn[u] > 0;i++)
              if (net[u][i] > 0 \&\& high[u] == high[i]+1)
                   push(u,i);
                    vn ++;
                   if (i != n) SQ.push(i);
          }
          if (vn == 0) relable(u);
    }
}
void init_preflow()
    int i;
     memset(high,0,
                       sizeof (high));
     memset(earn,0,
                       sizeof (earn));
                          SQ.pop();
    while (!SQ.empty())
    high[0] = n+1;
    for (i=1;i<=n;i++)
          if (net[0][i]
                       > 0) {
              earn[i]
                       = net[0][i];
                        -= net[0][i];
               earn[0]
              net[i][0]
                         = net[0][i];
              net[0][i]
                         = 0;
              if (i != n) SQ.push(i);
    }
}
int high_label_preflow_push()
    int i,j;
    init_preflow();
     while (!SQ.empty())
          int overp = SQ.front();
          SQ.pop();
```

```
discharge(overp);
    }
    return earn[n];
}
// 带 gap 优化的高标预流
const int N = 128;
const int INF = 1 << 28;
class Edge {
public:
    int u, v, cuv, cvu, flow;
     Edge() {}
     Edge( int cu, int cv, int ccu, int ccv) : u(cu), v(cv), cuv(ccu), cvu(ccv),
                                                                                   flow(0)
{}
    int other( int p) const { return p == u ? v : u; }
    int cap( int p) const { return p == u ? cuv-flow : cvu+flow; }
    void addFlow( int p, int f) { flow += (p == u ? f : -f); }
};
class NodeList {
private :
    int level, next[N], index[2*N], v;
public:
    void clear( int cv) { v = cv; level = -1; memset(index, -1, sizeof (index)); }
    void insert( int n, int h) { next[n] = index[h]; index[h] = n; level >?= h; }
    int remove();
                  const { return level < 0; }
     bool empty()
};
int NodeList::remove()
     int r = index[level]; index[level]
                                         = next[index[level]];
    while (level >= 0 && index[level]
                                          == -1) level--;
     return r;
}
class Network
private:
     vector<Edge>
                      eg;
    vector<Edge*>
                       net[N];
     int v, s, t;
     NodeList list;
     int h[N],
               hn[2*N],
                          e[N],
                                cur[N];
    void initNet();
     void initFlow();
    void initHeight();
```

```
void push( int );
    void relabel( int );
    void discharge( int );
    void gapHeuristic( int );
public:
    bool build();
    int maxFlow( int , int );
};
if (hn[k] != 0 || k >= v+1) return;
    for (int i = 0; i < v; i++)
         if (h[i] > k \&\& h[i] <= v \&\& i != s)
         \{ hn[h[i]]--; hn[v+1]++; h[i] = v+1; \}
}
void Network::initNet()
    for (int i = 0; i < v; i++) net[i].clear();
    for (int i = eg.size()-1; i >= 0; i--) {
         net[eg[i].u].push_back(&eg[i]);
         net[eg[i].v].push_back(&eg[i]);
    }
}
void Network::initHeight()
    memset(h, 0, sizeof (h)); memset(hn,
                                               0, sizeof (hn));
    memset(e,
                 0, sizeof (e)); e[s] = INF;
    for (int i = 0; i < v; i++)
                                h[i] = v;
    queue< int > Q; Q.push(t);
                                h[t] = 0;
    while (!Q.empty())
         int p = Q.front();
                             Q.pop();
         for (int i = net[p].size()-1; i >= 0; i--) {
             int u = net[p][i] - sother(p), ec = net[p][i] - scap(u);
             if (ec != 0 \&\& h[u] == v \&\& u != s) { h[u] = h[p]+1; Q.push(u);
    for (int i = 0; i < v; i++) hn[h[i]]++;
}
void Network::initFlow()
    initNet();
               initHeight();
    for (int i = 0; i < v; i++) cur[i] = net[i].size()-1;
    list.clear(v);
    for (; cur[s] >= 0; cur[s]--)
                                    push(s);
}
void Network::push(
                      int u) {
    Edge^* te = net[u][cur[u]];
                                 e[u], p = te->other(u);
    int ex = min(te->cap(u),
```

```
if(e[p] == 0 \&\& p != t) list.insert(p,
                                          h[p]);
    te->addFlow(u, ex); e[u] -= ex; e[p] += ex;
}
void Network::relabel(
                     int u) {
    int mh = 2^*v, oh = h[u];
    for (int i = net[u].size()-1;
                              i >= 0; i--) {
         int p = net[u][i]->other(u);
         if (net[u][i]->cap(u) != 0) mh <?= h[p]+1;
    }
    hn[h[u]]--;
                 hn[mh]++; h[u] = mh; cur[u] = net[u].size()-1;
    gapHeuristic(oh);
}
void Network::discharge( int u) {
    while (e[u] > 0)
         if (cur[u] < 0) relabel(u);
         else if (net[u][cur[u]]->cap(u) > 0 && h[u] ==
h[net[u][cur[u]]->other(u)]+1)
                                    push(u);
         else cur[u]--;
}
bool Network::build()
    int m, np, nc;
    int a, b, l, i;
    if (scanf( "%d %d %d %d" , &v, &np, &nc, &m) != 4) return false ;
    v += 2; eg.clear();
    for (i = 0; i < m; i++)
         scanf( "\n(%d,%d)%d"
                                , &a, &b, &l);
         eg.push_back(Edge(a+2, b+2, l, 0));
    for (i = 0; i < np; i++)
         scanf( "\n(\%d)\%d" , \&a, \&l);
         eg.push_back(Edge(0, a+2, l, 0));
    for (i = 0; i < nc; i++)
         scanf( "\n(%d)%d" , &a, &l);
         eg.push_back(Edge(a+2,
                                 1, l, 0));
    }
     return true;
}
int Network::maxFlow(
                         int ss, int tt) {
     s = ss; t = tt; initFlow();
    while (!list.empty())
         int u = list.remove();
         discharge(u);
    }
```

```
return e[t];
}
int main()
{
    Network net;
                  printf( "%d\n" , net.maxFlow(0,
    while (net.build())
                                                   1));
    return 0;
}
网络中求最大流 Dinic 算法 O(V^2E)
适用于稠密图,实际复杂度低于 HLPP 模板
参数含义: n 代表网络中节点数 ,第0节点为源点 , 第n 节点为汇点
            net 代表网络,使用前向星表示法存储边
            dis[] 代表从源点出发的距离标号
            path[] 代表模拟栈中的路径信息
            cur[] 代表模拟栈的现场保存
           最大流量
返回值:
*/
const int NMAX = 21000;
const int MMAX = 250000 << 1;
struct EDGE {
    int u, v, cap, flow;
    int next;
    EDGE( int _u=0, int _v=0, int _c=0, int _f=0)
       : u(_u), v(_v), cap(_c), flow(_f)
};
const int ENDFLAG = -1;
struct EDGELIST {
    int start[NMAX];
    int last[NMAX];
    int tot;
    EDGE arc[MMAX];
    void clear() {
        tot = ENDFLAG +1;
                     ENDFLAG, sizeof (last));
        memset(last,
    }
    void push_back(EDGE
                         edge) {
        edge.next = ENDFLAG;
        arc[tot] = edge;
        if (last[edge.u] != ENDFLAG) arc[ last[edge.u]
                                                    ].next = tot;
        else start[edge.u] = tot;
```

```
last[edge.u]
                         = tot;
          tot ++;
        创建双向弧
     void add_arc(EDGE
                            edge) {
          push_back(edge);
          push_back(EDGE(edge.v,edge.u,edge.cap));
}net;
int que[2][NMAX];
int qf[2],qe[2],qnow;
                         (que[qnow][
                                         qe[qnow]++
#define
          push_que(a)
                                                         ] = (a)
#define
                                           qf[qnow^1]++
                                                              ])
          pop_que2
                        (que[qnow^1][
#define
          switch_que
                        qnow ^= 1;
                        qf[qnow]
                                   = qe[qnow]
#define
                             (qf[qnow^1]
                                             >= qe[qnow^1]
          empty_que2
         size_que2
                        (qe[qnow^1]
                                         qf[qnow^1])
#define
int n, m;
int dis[NMAX];
int path[NMAX],
                    deep;
int cur[NMAX];
bool bfs() {
     int i, j;
     memset(dis,-1,
                       sizeof (dis));
     dis[0] = 0;
     qnow = 0;
     switch_que;
     push_que(0);
     switch_que;
     while (!empty_que2)
         int I = size_que2;
          while (I --) {
              int u = pop_que2;
              for (i=net.start[u];i!=ENDFLAG;i=net.arc[i].next)
                                                                          {
                   int v = net.arc[i].v;
                                     && net.arc[i].cap>net.arc[i].flow)
                   if (dis[v]==-1)
                                                                              {
                        push_que(v);
                        dis[v] = dis[u]+1;
                        if (v == n) return true;
              }
```

```
}
          switch_que;
     return false;
}
int Dinic()
{
     int i, j;
     int u;
     int maxflow
                   = 0;
     while (bfs())
                  {
          memcpy(cur,net.start,
                               sizeof (cur));
          for (deep=u=0;
                             true ;) {
               if (u==n)
                    int neck = INT_MAX,
                                              pos;
                    for (i=0;i<deep;i++)
                         int res = net.arc[path[i]].cap
                                                        net.arc[path[i]].flow;
                         if (res<neck)
                              neck = res;
                              pos = i;
                         }
                    maxflow += neck;
                    for (i=0;i<deep;i++)
                         net.arc[path[i]].flow
                                                  += neck;
                         net.arc[path[i]^1].flow
                                                      -= neck;
                    deep = pos;
                    u = net.arc[path[deep]].u;
               for (i=cur[u];i!=ENDFLAG;i=net.arc[i].next)
                    if (net.arc[i].cap>net.arc[i].flow
                         && dis[u]+1==dis[net.arc[i].v])
                                                                 break ;
               }
               cur[u] = i;
               if (i!=ENDFLAG)
                    path[deep++]
                                      = i;
                    u = net.arc[i].v;
               }
               else {
                    if (deep==0)
                                    break ;
                    dis[u] = -1;
                    u = net.arc[path[--deep]].u;
               }
```

```
}
    return maxflow;
}
         网络最小费用最大流
    8.
网络中最小费用最大流 O(V*E^2)
参数含义:
           n 代表网络中的总节点数 ,第0 节点为源点 , 第n节点为汇点
           net[][] 代表剩余网络
           cost[][] 代表单位费用
           path[] 保存增广路径
           ecost[] 源点到各点的最短路
算法:初始最小费用和最大流均为 0,寻找单位费用最短路
在最短路中求出最大流,即为增广路,再修改剩余网络,直到无可增广路为止
           最小费用,最大流量
返回值:
                           ****/
const int NMAX = 210;
int net[NMAX][NMAX],
                     cost[NMAX][NMAX];
int path[NMAX], ecost[NMAX];
int n;
bool bellman_ford()
    int i,j;
    memset(path,-1,
                   sizeof (path));
    fill(ecost,
             ecost+NMAX,
                        INT_MAX);
    ecost[0] = 0;
    bool flag = true;
    while (flag) {
        flag = false;
        for (i=0;i<=n;i++)
           if (ecost[i] == INT\_MAX) continue;
           for (j=0; j<=n; j++) {
                                                   < ecost[j]) {
               if (net[i][j]
                         > 0 && ecost[i]+cost[i][j]
                   flag = true ;
                           = ecost[i]+cost[i][j];
                   ecost[j]
                   path[j]
                          = i;
           }
    return ecost[n]
                  != INT_MAX;
}
```

```
int min_cost_max_flow()
{
    int i,j;
    int mincost = 0, maxflow = 0;
    while ( bellman_ford()
                      ) {
        int now = n;
        int neck = INT_MAX;
        while (now != 0) {
            int pre = path[now];
            neck = min(neck, net[pre][now]);
            now = pre;
        }
        maxflow += neck;
        now = n;
        while (now != 0) {
            int pre = path[now];
            net[pre][now] -= neck;
            net[now][pre] += neck;
            cost[now][pre]
                         = - cost[pre][now];
            mincost += cost[pre][now] * neck;
            now = pre;
    return mincost;
                              邻接表 SPFA 实现
网络中最小费用最大流 O(V*E^2)
            n 代表网络中的总节点数 ,第s节点为源点 , 第t节点为汇点
参数含义:
            net 代表剩余网络
            path[] 保存增广路径
            ecost[] 源点到各点的最短路
返回值:
            最小费用,最大流量
                            ****/
// POJ 3422
const int NMAX = 5100; // 点数
const int MMAX = 30000; // 边数
const int INF = 0x7f7f7f7f;
int path[NMAX],
                ecost[NMAX];
int n;
int s, t;
struct EDGE {
```

```
int u, v, cap, cost, flow;
    int next;
    EDGE( int _u=0, int _v=0, int _c=0, int _c=0,
                                                        int _f=0
         : u(_u), v(_v), cap(_c), flow(_f),
                                              cost(_ct)
                                                         {}
};
const int ENDFLAG = -1;
struct EDGELIST {
    int start[NMAX];
    int last[NMAX];
    int tot;
    EDGE arc[MMAX];
    void clear() {
         tot = ENDFLAG +1;
                        ENDFLAG, sizeof (last));
         memset(last,
    }
    void push_back(EDGE
                             edge) {
         edge.next = ENDFLAG;
         arc[tot] = edge;
         if (last[edge.u] != ENDFLAG) arc[ last[edge.u]
                                                             ].next = tot;
         else start[edge.u]
                             = tot;
         last[edge.u]
                     = tot;
         tot ++;
       创建双向弧
    void add_arc(EDGE
                          edge) {
         push_back(edge);
         push_back(EDGE(edge.v,edge.u,0,INF));
}net;
int que[2][NMAX];
int qf[2],qe[2],qnow;
#define
                        (que[qnow][
                                      qe[qnow]++ ] = (a))
         push_que(a)
                                        qf[qnow^1]++
                                                          ])
         pop_que2
                       (que[qnow^1][
#define
                       qnow ^= 1;
#define
         switch_que
                       qf[qnow] = qe[qnow]
                           (qf[qnow^1]
         empty_que2
                                          >= qe[qnow^1]
#define
#define
                       (qe[qnow^1]
                                      qf[qnow^1])
         size_que2
bool SPFA()
{
    int i,j;
    bitset <NMAX>
                      vis;
```

```
memset(ecost,
                   0x7f, sizeof (ecost));
                  -1, sizeof (path));
    memset(path,
    bool flag = true ;
    qnow = 1;
    switch_que;
    push_que(s);
    vis[s] = 1;
    ecost[s] = 0;
    for (j=0; j< n \&\& flag; j++)
         flag = false;
         switch_que;
         int I = size_que2;
         while (I --)
             int now = pop_que2;
             vis[now] = 0;
             for (i=net.start[now];
                                  i!=ENDFLAG;
                                                      i=net.arc[i].next)
             {
                   EDGE ed = net.arc[i];
                  if (ed.cap>ed.flow
                                       && ecost[ed.v]>ecost[now]+ed.cost)
                       flag = true ;
                                    = ecost[now]+ed.cost;
                       ecost[ed.v]
                       path[ed.v] = i;
                       if (! vis[ed.v])
                            vis[ed.v] = 1;
                            push_que(ed.v);
                  }
             }
                     != INF;
    return ecost[t]
int min_cost_max_flow()
    int i,j;
    int mincost = 0, maxflow = 0;
```

}

{

```
while (SPFA()) {
       int pre = path[t];
       int neck = INT_MAX;
       while (pre != -1) {
           int res = net.arc[pre].cap
                                 net.arc[pre].flow;
           neck = min(neck,
                            res);
           pre = path[net.arc[pre].u];
       maxflow
               += neck;
       mincost += ecost[t]
                           * neck;
       pre = path[t];
       while (pre != -1) {
           net.arc[pre].flow
                            += neck;
           net.arc[pre^1].flow
                              -= neck;
           net.arc[pre^1].cost
                              = - net.arc[pre].cost;
           pre = path[net.arc[pre].u];
       }
   return mincost;
}
    9.
         网络最大流(高度标号预流推进)
    函数接口: int Relabel_To_Front(int s,int d) O(V^2*sqrt(E))
    参数含义: s 为源点, d 为汇点
    返回值 :
              网络最大流
                      :ver 置为网络中节点的个数 ,
                                                 c[i][j] 代表节点 i 到
    调用函数前的初始化工作
    节点 j 的流量 , vI[i] 存放 i 与相邻的所有节点
    其它全局变量均初始化为零
    */
    const int VEX = 405; // 网络中顶点数
    const int HMAX = 810; // 最大高度的定义 , 只要大于顶点的 2 倍就可以了
    int f[VEX][VEX]; // 流量
    int c[VEX][VEX]; // 边最大容量
    int h[VEX];
                  // 节点高度
    int e[VEX];
              // 节点容量
            // 节点数目
    int ver;
    vector< int > vI[VEX]; // 邻接表 , vI[i] 存放与 i 相邻的节点
    void Push( int u, int v) // 流推进,由节点 u 推向 v
    {
       int cf = c[u][v] - f[u][v]; //u,v 边的容量
```

```
int d = e[u] < cf ? e[u] : cf;
   f[u][v] += d;
   f[v][u] = -f[u][v];
   e[u] -= d;
   e[v] += d;
}
void Relabel( int u) // 对 u 重新标号
{
   int i,t,cf;
   int hmin = HMAX;
   for (i = 0; i < vl[u].size() ; i++){ // 寻找相邻最低点
       t = vl[u][i];
       cf = c[u][t] - f[u][t];
       if (cf > 0 \&\& h[u] <= h[t] \&\& h[t] < hmin)
           hmin = h[t];
    }
   h[u] = hmin + 1;
}
void Init_Preflow( int s) // 初始化网络流 , s 为源点
{
   int i;
    int u;
    h[s] = ver; // 初始化高度
                              ; i++){
   for (i = 0 ; i < vl[s].size()
       u = vl[s][i];
       f[s][u] = c[s][u];
       f[u][s] = -c[s][u];
       e[u] = c[s][u];
       e[s] -= c[s][u];
   }
}
void Discharge( int u)
   int i = 0;
   int cf,v;
   if(vI[u].size() == 0) return ;
   while (e[u] > 0){
```

```
if (i < vl[u].size())
             v = vl[u][i];
             cf = c[u][v]
                            - f[u][v];
        }
         if (i >= vl[u].size()){}
             Relabel(u);
             i = 0;
         }
         else if (cf > 0 && h[u] == h[v] + 1)
             Push(u,v);
         else
             i++;
}
int Relabel_To_Front( int s, int d) //s 为源点, d 为汇点
{
    int u,i,old_h;
    list< int > I;
    list< int >::iterator
                            iter;
    Init_Preflow(s);
    iter = I.begin();
    for (i = 0 ; i < ver ; i++){}
         if (i != s \&\& i != d)
             l.insert(iter,i);
    iter = I.begin();
    while (iter != l.end()){
         u = *iter;
         old_h = h[u];
         Discharge(u);
         if (h[u] > old_h){
             I.erase(iter);
             l.insert(l.begin(),u);
             iter = I.begin();
         }
         iter++;
    return e[ver - 1];
}
```

10. 最大团

```
Function Name:
                              最大独立集,最大团
    Description
                         PKU 1419 Graph Coloring
    团: 指G的一个完全子图 , 该子图不包含在任何其他的完全子图当中
    最大独立集 : 补图的最大团
    最大团: 指其中包含顶点最多的团
#include <cstdio>
#include <string>
        NMAX 110
#define
bool path[NMAX][NMAX];
int n, mmax;
int dp[NMAX];
bool v[NMAX];
int seq[NMAX], seq_pos;
//seq 记录最大团集合
bool dfs( int pos, int size)
{
    int i, j, unvis;
    bool tv[NMAX];
    unvis = 0;
    for (i=pos;i< n;i++)
        if (!v[i]) {
            unvis ++;
        }
    if (unvis == 0) \{ //|U| = 0 \}
        if (size > mmax) {
             mmax = size;
            seq_pos = 0;
            seq[ seq_pos ++] = pos+1;
            return true;
        }
        return false;
    }
    for (i=pos; i < n \&\& unvis > 0 ; i++) {
        if (!v[i]) {
             if (unvis + size \leftarrow mmax || dp[i] + size \leftarrow mmax) {
                 return false;
            }
            v[i] = true ; //U = U \setminus \{vi\}
            unvis --;
```

```
memcpy(tv, v, sizeof (v));
             for (j=0;j< n;j++) { //U N(vi);
                  if (!path[i][j]) {
                      v[j] = true ;
                  }
             }
             if ( dfs(i, size+1) ) {
                  seq[ seq_pos ++] = pos+1;
                  return true;
             memcpy(v, tv, sizeof (v));
    } //while U is not empty
    return false;
}
int max_clique()
{
    int i,j;
    mmax = 0;
    for (i=0;i< n;i++) {
         path[i][i] = false;
    }
    for (i=n-1;i>=0;i--) {
         for (j=0;j< n;j++) { //Si
                                     N(vi);
           v[j] = !path[i][j];
         dfs(i, 1);
         dp[i] = mmax;
    return mmax;
}
int main()
{
    int i,j,x,y,e;
    int m,tn;
    scanf( "%d" , &m);
    while (m --) {
         scanf( "%d %d" , &n, &e);
         memset(path,0, sizeof (path));
         for (i=0;i<e;i++) {
             scanf( "%d %d" , &x,&y);
```

```
x--; y--;
           path[x][y] = path[y][x] = true ;
       }
       //max independent
                        set in original
                                       graph
       //max clique in inverse graph
       for (i=0;i< n;i++) {
           for (j=0;j< n;j++) {
               path[i][j] = !path[i][j];
           }
       }
       memset(dp,0, sizeof (dp));
       printf( "%d\n" , max_clique());
       printf( "%d" , seq[0]);
       for (i=1;i < seq_pos;i++) {
           printf( " %d" , seq[i]);
       printf( "\n" );
}
    11. 最大二分图匹配(匈牙利算法)
       Function Name: 最大二分图匹配 (匈牙利算法)
       Description : HDOJ 2063 过山车
       二分图: 指所有顶点分成集合 M 和 N, M 或 N 中任意两个在同一集合中的点互不相连
       匹配:一组边顶点分别在两个集合中,并且任意两条边都没有相同顶点
        最大匹配: 所能得到的最大的边的个数
    #include <cstdio>
    #include <memory>
    #include <vector>
    using namespace
                     std;
    const int Max=1100;
    vector< vector< int > > Bmap;
    int n, m, k, nm;
    int mark[Max];
    bool flag[Max];
    bool dfs( int pos)
      int i, pre, tp;
      for (i=0; i < Bmap[pos].size()
                                   ;i++) {
         tp = Bmap[pos][i];
```

```
if (!flag[tp]
                      ) {
              flag[tp]
                       = true ;
              pre = mark[tp];
              mark[tp]
                         = pos;
              if (pre==-1
                          || dfs(pre))
                                           return true;
              mark[tp]
                         = pre;
        return false;
     }
     inline int Max_Match()
     {
        int mmax = 0, i;
        for (i=1; i \le m ; i++)
           memset(flag,0,
                            sizeof (flag));
           if ( dfs(i) )
                         mmax++;
        }
        return mmax;
     int main()
        int i, j, id, id2;
        while (scanf("%d",
                            scanf("%d%d",&m,
                                   &n);
            nm = n + m;
            Bmap.clear();
                               Bmap.resize(nm+10);
            memset(mark,-1,
                                sizeof (mark));
            for (j=0; j < k; j++)
                            %d", &id, &id2);
                scanf("%d
                id2 += m;
                Bmap[id].push_back(id2);
            }
           printf("%d\n",
                           Max_Match());
     }
// 二分匹配 HopcroftKarp
                          算法 O(sqrt(V)*E)
// 贪心一个初始匹配可以加速
#include
          <iostream>
#include
          <queue>
using namespace
                   std;
const int MAXN = 3002;
```

```
const int INF = 1 << 30;
struct node{
    int x, y;
           U[MAXN];
}G[MAXN],
int n, m, t, nx, ny, dis;
int x[MAXN], y[MAXN],
                          vs[MAXN];
int Isf[MAXN];
bool adj[MAXN][MAXN];
int ds[MAXN], dt[MAXN];
void input(){
     scanf( "%d %d" , &t, &m);
     for (int i = 0; i < m; i++){
         scanf( "%d %d %d" , &G[i].x, &G[i].y, &vs[i]);
         VS[i] *= VS[i];
     }
    scanf( "%d" , &n);
    for (int i = 0; i < n; i++)
         scanf( "%d %d" , &U[i].x,
                                     &U[i].y);
     memset(adj, 0, sizeof (adj));
    if (m < n)
         nx = m, ny = n;
         for (int i = 0; i < m; i++)
              for (int j = 0; j < n; j++){
                   int a = G[i].x - U[j].x,
                                           b = G[i].y - U[j].y;
                   if (a^*a + b^*b) < vs[i]^*t^*t
                       adj[i][j]
                                 = 1;
    } else {
         nx = n, ny = m;
         for (int i = 0; i < n; i++)
              for (int j = 0; j < m; j++){
                   int a = G[j].x - U[i].x, b = G[j].y - U[i].y;
                   if (a^*a + b^*b) \le vs[j]^*t^*t
                       adj[i][j]
                                 = 1;
              }
}
bool Search() {
     memset(ds,
                   -1, sizeof (ds));
     memset(dt, -1, sizeof (dt));
```

```
queue< int > Q; dis = INF;
    for (int i = 0; i < nx; i++)
         if (x[i] == -1){
              Q.push(i);
              ds[i] = 0;
         while (!Q.empty())
              int u = Q.front();
                                  Q.pop();
              if (ds[u] > dis) break ;
              for (int v = 0; v < ny; v++){
                   if (adj[u][v]){
                        if (dt[v] != -1) continue;
                        dt[v] = ds[u]+1;
                        if (y[v] == -1) dis = dt[v];
                        else {
                            ds[y[v]] = dt[v]+1;
                            Q.push(y[v]);
                        }
         }
         return (dis != INF);
}
bool DFS( int u) {
    for (int v = 0; v < ny; v++){
                                 && dt[v] == ds[u] + 1){
         if (!Isf[v]
                    && adj[u][v]
              Isf[v] = true ;
              if (y[v] != -1 \&\& dt[v] == dis) continue;
              if (y[v] == -1 || DFS(y[v])){
                   y[v] = u; x[u] = v;
                   return 1;
         }
    return 0;
}
int HopcroftKarp()
    int cnt = 0;
    for (int i = 0; i < nx; i++) x[i] = -1;
    for (int i = 0; i < ny; i++) y[i] = -1;
    while (Search()){
         memset(Isf,
                        0, sizeof (lsf));
```

```
for (int i = 0; i < nx; i++)
             if(ds[i]) == 0 \&\& DFS(i)
                 cnt++;
    }
    return cnt;
}
int main(){
    int test;
    scanf( "%d" , &test);
    for (int k = 1; k \le test; k++){
        input();
        printf( "Scenario #%d:\n%d\n\n" , k, HopcroftKarp());
    }
    return 0;
}
     12. 带权二分图最优匹配 (KM 算法)
// 二分图带权匹配 O(N^3)
const int MAXN = 509;
const int INF = 0x1fffffff;
int bpCostMatch( int c[][MAXN], int nx, int ny) {
    static int lx[MAXN], ly[MAXN], slack[MAXN];
    static int open[MAXN],
                            prev[MAXN], pnt[MAXN], x[MAXN], y[MAXN];
    int i, j, k, s, head, tail;
    int d, ans = 0;
    if (nx > ny) ny = nx;
    for (i = 0; i < nx; i++) |x[i] = -INF;
    for (i = 0; i < ny; i++)
                             ly[i]
    for (i = 0; i < nx; i++)
        for (j = 0; j < ny; j++)
             if ((|x[i] - c[i][j]) < 0)
                 Ix[i] = c[i][j];
    memset(x, -1, size of (x)); memset(y, -1, size of (y));
    for (i = 0; i < nx; i++)
        memset(prev, -1, sizeof (prev));
        for (j = 0; j < ny; j++) slack[j] = INF;
        open[0] = i; head = 0; tail = 1;
         while (x[i] < 0) {
             for (; head < tail && x[i] < 0; head++)
                 for (s = open[head], j = 0; j < ny && x[i] < 0; j++)
                      if (prev[j] < 0) {
```

```
if ((d = lx[s] + ly[j] - c[s][j]) > 0) {
                               if ((slack[j] - d) > 0) {
                                   slack[j] = d; pnt[j] = s;
                               }
                               continue ;
                           open[tail++] = y[j]; prev[j] = s;
                           if (y[j] >= 0) continue;
                           while (j \ge 0) {
                             s = prev[j]; y[j] = s; k = x[s]; x[s] = j; j = k;
                           }
                      }
                      if (x[i] >= 0) break;
                      for (d = INF, j = 0; j < ny; j++)
                           if (prev[j] < 0 \&\& (d - slack[j]) > 0)
                           d = slack[j];
                      for (j = 0; j < tail; j++) lx[open[j]]
                                                           -= d;
                      for (j = 0; j < ny; j++)
                           if (prev[j] >= 0)
                               ly[j] += d;
                           else if (slack[j] < INF)
                               slack[j] -= d;
                      for (j = 0; j < ny; j++)
                           if (prev[j] < 0 \&\& slack[j] == 0) {
                               open[tail++] = y[j]; prev[j] = pnt[j];
                               if (y[j] >= 0) continue;
                               while (j >= 0) {
                                   s = prev[j]; y[j] = s; k = x[s]; x[s] = j; j = k;
                               }
                               break ;
                           }
    for (i = 0; i < nx; i++)
         if (c[i][x[i]] > -INF) {
             if (c[i][x[i]] < 0)
                  return -1;
             ans += c[i][x[i]];
         } else return -1;
    return ans;
}
int N, M, E;
int c[MAXN][MAXN];
```

```
int cas;
int main() {
    int i, j, a, b, w, ans;
    while (scanf( "%d%d%d" , &N, &M, &E) != EOF) {
        for (i = 0; i < N; i++)
            for (j = 0; j < M; j++)
                c[i][j] = -INF;
        for (i = 0; i < E; i++)
            scanf( "%d%d%d" , &a, &b, &w);
            if (w < 0) continue;
            if (c[a][b] < w)
                c[a][b] = w;
        }
        if (N > M) ans = -1;
        else ans = bpCostMatch(c, N, M);
        printf( "Case %d: %d\n" , ++cas, ans);
    }
    return 0;
}
/*
wywcgs 的KM O(n^3)
只需要把 Graph 里的 n( 顶点数 )和 edge[x][y]( 边权 )赋值 ,第一维为 x点 ,第二维为 y点。
然后调用 KMMatch() 函数即可,返回值为最大权完美匹配。
最小权匹配可将每条边权取反,然后类似求最大权匹配即可。
匹配信息保存在 xmate[] 和 ymate[] 中。其中 xmate[i] 为 x[i] 的匹配点 , ymate[i] 为 y[i] 的匹
配点。
*/
#include
         <cstdio>
#include
         <queue>
#include <algorithm>
using namespace std;
const int N = 310;
const int INF = 1 \ll 28;
class Graph {
private :
   bool xckd[N], yckd[N];
   int n, edge[N][N],
                     xmate[N],
                                ymate[N];
   int lx[N], ly[N], slack[N], prev[N];
   queue< int > Q;
   bool bfs();
```

```
void agument( int );
public:
   bool make();
   int KMMatch();
};
bool Graph::bfs() {
   while (!Q.empty()) {
      int p = Q.front(), u = p>>1; Q.pop();
      if (p&1)
         if (ymate[u] == -1) { agument(u); return true ; }
         else { xckd[ymate[u]] = true ; Q.push(ymate[u]<<1);
      } else {
         for (int i = 0; i < n; i++)
             if (yckd[i]) continue ;
             else if (lx[u]+ly[i] != edge[u][i])
                int ex = lx[u]+ly[i]-edge[u][i];
                if (slack[i] > ex) { slack[i] = ex; prev[i] = u; }
             } else {
                yckd[i] = true ; prev[i] = u;
                Q.push((i << 1)|1);
             }
      }
   return false;
void Graph::agument(      int u) {
   while (u != -1) {
      int pv = xmate[prev[u]];
      ymate[u] = prev[u]; xmate[prev[u]] = u;
      u = pv;
   }
}
int Graph::KMMatch()
   memset(ly, 0, sizeof (ly));
   for (int i = 0; i < n; i++) {
      Ix[i] = -INF;
      for (int j = 0; j < n; j++) lx[i] >?= edge[i][j];
   }
   memset(xmate,
                     -1, sizeof (xmate));
                                            memset(ymate,
                                                            -1, sizeof (ymate));
   bool agu = true;
   for (int mn = 0; mn < n; mn++) {
      if (agu) {
                          false, sizeof (xckd));
          memset(xckd,
          memset(yckd,
                          false, sizeof (yckd));
```

```
for (int i = 0; i < n; i++) slack[i] = INF;
          while (!Q.empty()) Q.pop();
          xckd[mn] = true ; Q.push(mn << 1);
       }
       if (bfs()) { agu = true ; continue ; }
       int ex = INF; mn--; agu = false;
       for (int i = 0; i < n; i++)
          if (!yckd[i]) ex <?= slack[i];</pre>
       for (int i = 0; i < n; i++) {
          if (xckd[i])
                      lx[i] -= ex;
          if (yckd[i])
                      ly[i] += ex;
          slack[i] -= ex;
       for (int i = 0; i < n; i++)
          if (!yckd[i] \&\& slack[i] == 0) \{ yckd[i] = true ; Q.push((i<<1)|1); \}
   }
   int cost = 0;
   for (int i = 0; i < n; i++) cost += edge[i][xmate[i]];
   return cost;
}
bool Graph::make()
    int i, j;
    while (scanf( "%d" , &n) == 1)
         for (i = 0; i < n; i++)
              for (j = 0; j < n; j++)
                   scanf( "%d" , &edge[i][j]);
         return true;
    return false;
}
int main()
   Graph g;
   while (g.make())
                      printf( "%d\n" , g.KMMatch());
   return 0;
}
```

13. 强连通分量 (Kosaraju 算法)

```
有向图的强连通分量 Kosaraju 算法 O(E)
参数含义:
            使用邻接表来保存图
             path 原图 , npath 逆图
             scc 强连通个数
             id[x]=y 表示第 x个顶点属于 y强连通
*/
        NMAX 11000
#define
vector<
        vector< int > > path;
        vector<
                 int > > npath;
vector<
int n,m, scc;
int order[NMAX], order_pos,
                              id[NMAX];
bool vis[NMAX];
void dfs( int pos)
{
    int i,j,l;
    vis[pos] = true;
    I = path[pos].size();
    for (i=0;i<1;i++) {
        j = path[pos][i];
        if (!vis[j]) {
             dfs(j);
        }
    order[ order_pos ++ ] = pos; //make
                                            order
}
void ndfs( int pos)
    int i,j,l;
    vis[pos] = true;
    id[pos] = scc;
    I = npath[pos].size();
    for (i=0;i<1;i++) {
        j = npath[pos][i];
        if (!vis[j]) {
             ndfs(j);
}
void Kosaraju()
{
```

```
int i,j;
    //dfs in original graph
    memset(vis,
                 0, sizeof (vis));
    order_pos = 0;
    for (i=1; i <= n ; i++)
        if (!vis[i])
             dfs(i);
        }
    }
    //dfs in inverse graph
    memset(vis, 0, sizeof (vis));
    memset(id,
                0, sizeof (id));
    scc = 1;
    for (i=order_pos-1; i>=0 ;i--) {
        if (!vis[ order[i]
                         ]) {
             ndfs(order[i]);
             SCC ++;
        }
    }
    SCC --;
}
     14. 强连通分量 (Gabow 算法)
有向图的强连通分量 Gabow 算法 O(E)
            使用邻接表来保存图
参数含义:
             path 原图
             scc 强连通个数
             id[x]=y 表示第 x个顶点属于 y强连通
*/
#define
        NMAX 11000
        vector< int > > path;
vector<
int n,m, scc, step;
int order[NMAX],
                  order_pos,
                              id[NMAX];
int order2[NMAX],
                   order2_pos;
int vis[NMAX];
void dfs( int pos)
{
    int i,j,next,l,pre;
    vis[pos] = step ++;
    order[ order_pos ++ ] = pos;
    order2[ order2_pos
                         ++ ] = pos;
    I = path[pos].size();
```

```
for (i=0;i<1;i++) {
        next = path[pos][i];
         if (vis[next]
                      == 0) {
             dfs(next);
        }
         else if (id[next] == 0) { //have a circle and belong to nothing
             while (vis[ order2[order2_pos -1] ] > vis[next])
                 order2_pos --;
         }
    } //for i
    if (order2[order2_pos
                         -1] == pos) { //if pos back to begin of scc
        order2_pos --;
    }
    else {
         return ;
    do { //record scc
        pre = order[order_pos
                                -1];
        id[pre] = scc;
        order_pos --;
    } while (pre != pos);
    SCC ++;
}
void Gabow()
    int i,j,l;
    //dfs in original graph
                 0, sizeof (id));
    memset(id,
    memset(vis,
                   0, sizeof (vis));
    scc = step = 1;
    order_pos = order2_pos = 0;
    for (i=1; i <= n ; i++) {
        if (vis[i] == 0) {
             dfs(i);
        }
    SCC --;
}
```

15. 无向图割边割点和双连通分量

```
#define mclear(x) memset((x), 0, size of ((x))) const int MAX = 5100;
```

```
int n,m,deep;
vector< int > path[MAX];
int vis[MAX],
               low[MAX];
vector< int > cutpoint; // 割点
         pair< int , int > > bridge; // 割边 ,桥
vector<
int nbcc; // 双连通分量数
stack< pair< int, int > > order;
vector< int > bcc[MAX]; // 双连通分量
void dfs( int pos, int father) {
    int i,j, total = 0;
    bool cut = false;
    int reback = 0; // 处理平行边
    vis[pos] = low[pos] = deep ++;
    int Is = path[pos].size();
    for (j=0;j<1s;j++)
         i = path[pos][j];
         if (i == father) reback ++;
         if (vis[i] == 0) {
              pair< int ,int > e(pos, i);
              order.push(e);
              dfs(i, pos);
              if (low[i] >= vis[pos]) {
                  nbcc ++;
                  bcc[nbcc].clear();
                   pair < int , int > r;
                   do {
                       r = order.top();
                       order.pop();
                       bcc[nbcc].push_back(r.second);
                  } while (e != r);
                  bcc[nbcc].push_back(r.first);
              }
              total ++;
              low[pos] = min(low[i], low[pos]);
              if ((vis[pos] == 1 && total > 1) ||
                   (vis[pos] != 1 \&\& low[i] >= vis[pos])) cut = true;
                       > vis[pos])
                                     bridge.push_back(e);
              if (low[i]
         else if (i != father) {
              low[pos] = min(vis[i],
                                        low[pos]);
         }
    }
    if (reback > 1) low[pos] = min(low[pos], vis[father]);
```

```
if (cut) cutpoint.push_back(pos);
}
void find_cut() {
   int i;
   mclear(vis);
              mclear(low);
                  bridge.clear();
   cutpoint.clear();
   nbcc = 0;
                       order.pop();
   while (!order.empty())
   for (i=1;i<=n;i++) {
       if (vis[i] == 0) {
          deep = 1;
          dfs(i, i);
       }
   }
}
/*********
图的 DFS 信息构建 by oyjpArt
g矩阵: g[i][j] -> 0: 无边
          1: 可重复访问边
          -1: 非可重复访问边
说明:以为在无向图中 u->v 访问之后就不能再从 v->u 访问了
故{u, v} 访问了之后 {v, u} 要置 -1
如果是有向图则没有这个规则
gc 矩阵 :gc[i][j]-> 0 : 无边
1: 树枝边
2: 反向边
3: 正向边
4: 交叉边
d数组: 顶点的开始访问时间表
f数组: 顶点的结束访问时间表
c数组: 顶点颜色表 0白色-1 灰色 1 黑色
p数组: 顶点的前驱表
I数组: 顶点的 L值(最顶层的祖先层数 )
b数组: 顶点的度数表
关于标号函数 LOW()
LOW(U) 代表的是与 U以及 U的子孙直接相连的结点的最高辈分(深度)
                  U首次被访问时
d[U]
LOW[U] =
                        d[W])
                                 访问边 {U,W}
            min(LOW[U],
                     U的儿子 S的关联边全部被访问时
min(LOW[U],
            LOW[S])
/*************************/
const int maxn = 100;
int n, g[maxn][maxn],
                    gc[maxn][maxn];
```

```
p[maxn], c[maxn], b[maxn];
int d[maxn],
           f[maxn],
                    l[maxn],
int time;
void dfs_visit( int u) { // 递归搜索以 U为根的深度优先树
   int v;
   c[u] = -1; // 置顶点为灰色 // 去掉这句之后适用于有向图 (后面设置不可访问亦
同)
   time++; d[u] = time, l[u] = time;
   for (v = 1; v \le n; v + +)
       if (g[u][v] > 0)
          if (c[v] == 0) { // 如果 v是白色节点
                              // 不可再访问
              g[v][u] = -1;
                              // 树枝边
              gc[u][v] = 1;
                                // 度数
              b[u]++;
                             // 记录父亲节点
              p[v] = u;
              dist_visit(v); // 递归搜索以 v 为根的深度优先树
              if (I[v] < I[u]) //v 是u的后代
                 |[u] = |[v]; //u 的儿子 v的关联边搜索完后计算父亲的 | low 值
              g[v][u] = 1; // 恢复可访问标志
          } else {
              if (c[v] < 0) { // 若顶点为灰色
                 if (l[v] < l[u]) //u 与v相连
                     I[u] = I[v];
                 gc[u][v] = 2; // 反向边
                               // 黑色
              } else {
                 if (d[v] > d[u])
                     gc[u][v] = 3; // 正向边
                 else
                     gc[u][v] = 4; // 交叉边
              }
                       //DFS 完毕置黑色吧
   c[u] = 1;
   time++; f[u] = time;
}
void dfs() {
   int u;
              0, sizeof (gc));
   memset(gc,
             0, sizeof (c));
   memset(c,
             0, sizeof (b));
   memset(b,
   time = 0;
   for (u = 1; u \le n; u++)
       if (c[u] == 0) {
          p[u] = 0;
```

```
dfs_visit(u);
        }
}
     16. 最小树形图 O(N^3)
最小树形图 O(N^3)
             使用邻接矩阵来保存图,邻接表
参数含义:
                                        O ( VE )
             path 原图
             pre 保存最小入弧的权
             del 表示被缩去的点
             fpre 保存最小树形图的逆路径
例题:
             TJU 2248 Channel Design
*/
const int NMAX = 110;
const int INF = 0x7f7f7f7f;
int n;
int path[NMAX][NMAX];
int pre[NMAX];
                 del[NMAX];
bool vis[NMAX],
int min_cost;
int fold[NMAX],
                fpre[NMAX];
void dfs( int pos) {
    int i;
    vis[pos] = true ;
    for (i=1;i<=n;i++)
        if (path[pos][i]
                     != INF && !vis[i])
                                            dfs(i);
bool is_connect( int root) {
    int i;
    memset(vis,
                 0, sizeof (vis));
    dfs(root);
    for (i=1;i<=n;i++)
        if (!vis[i])
                  return false;
    }
    return true;
}
//O(N^3)
bool min_tree_graph(
                      int root) {
    int i,j,k;
    //make sure every node(except
                                      root) have in-arc
```

if (!is_connect(root))

return false;

```
memset(del, 0, sizeof (del));
    min\_cost = 0;
    for (i=0;i<=n;i++)
                          fold[i] = fpre[i] = i;
    while (true ) {
         for (i=1;i<=n;i++)
              if (del[i]
                     || i == root) continue ;
             pre[i] = i;
              path[i][i] = INF; //delete
                                          self-cycle
             for (j=1;j<=n;j++) {
                  if (del[j])
                           continue ;
                             < path[ pre[i] ][i])
                  if (path[j][i]
                                                     pre[i] = fpre[fold[i]]
             }
         } //find min in-arc
         for (i=1;i<=n;i++)
             if (del[i]
                      || i == root) continue ;
             j = i;
              memset(vis, 0, sizeof (vis));
              while (!vis[j] && j != root) {
                  vis[j] = true ;
                  j = pre[j];
             }
              if (j == root) continue ; //no cycle
             i = j; //cycle begin node
             min_cost += path[ pre[i] ][i];
                          j != i ;j=pre[j])
              for (j=pre[i];
                  del[j] = true ; //fold
                                        cycle
                  min_cost += path[ pre[j]
                                                ][j]; //add cycle cost
             for (j=1;j<=n;j++)
                  if (del[j])
                             continue ;
                  if (path[j][i]
                             != INF) path[j][i]
                                                    -= path[ pre[i] ][i];
             } //i is new fold node
             for (k=1;k<=n;k++) {
                                  continue ;
                       if (del[k])
                                   = min(path[i][k], path[j][k]);
                       path[i][k]
                       if (path[k][j] != INF && path[k][i]
                                                             > path[k][j]
path[ pre[j]
             ][j]) {
                            path[k][i]
                                       = path[k][j]
                                                      path[ pre[j] ][j];
                            fold[i] = j; //record fold node
                            fpre[i] = j;
                       }
                  }
             } //make new graph
```

```
break ;
         }
         if (i > n) {
             for (i=1;i<=n;i++) {
                  if(del[i] || i == root) continue;
                  min_cost += path[ pre[i] ][i];
             break ;
         } //graph no cycle
    } //while have cycle
    return true;
}
//print path in min tree graph
void print_mtg( int root) {
    int i, total = n;
    memset(vis, 0, sizeof (vis));
    for (i=1;i \le n;i++) vis[fpre[i]]
                                   = true ;
    for (i=1;i<=n;i++) {
         if (!vis[i]) {
             int pos = i;
             while (pos != root) {
                  printf( "%d <- ", pos);</pre>
                  pos = fpre[pos];
             printf( "%d\n" , root);
int main() {
    int i,m;
    while (scanf( "%d %d", &n,&m),
                                       !(n==0 && m==0))
         memset(path,
                         0x7f, sizeof (path));
         while (m --) {
             int x,y,z;
             scanf( "%d
                        %d %d" , &x,&y,&z);
                          = min(path[x][y],
             path[x][y]
                                              z);
         }
         if ( !min_tree_graph(1)
                              ) puts( "impossible"
         else printf( "%d\n" , min_cost);
}
     17. 最小树形图 O(VE)
const int NMAX = 1500;
```

```
const int INF = 0x7f7f7f7f;
struct LINKT {
     int Is;
     int adj[NMAX];
     void clear() \{ls = 0;\}
                   [] (const int pos) { return
     int operator
                                                  adj[pos];}
     int size() { return ls;}
     void push_back(
                        const int pos) {adj[ls
                                                  ++] = pos;}
};
int n;
int path[NMAX][NMAX];
LINKT epath[NMAX],
                         nepath[NMAX];
int pre[NMAX];
bool vis[NMAX],
                   del[NMAX];
int min_cost;
int fold[NMAX],
                  fpre[NMAX];
void dfs( int pos) {
     int i;
     vis[pos] = true ;
     for (i=0;i<epath[pos].ls;i++)
          if (!vis[ epath[pos].adj[i]
                                       ]) dfs(epath[pos].adj[i]);
     }
}
bool is_connect(
                   int root) {
     int i;
                    0, sizeof (vis));
     memset(vis,
     dfs(root);
     for (i=1;i<=n;i++)
          if (!vis[i])
                     return false;
     return true;
}
//O(VE)
bool min_tree_graph( int root) {
     int i,j,k;
     //make sure every node(except
                                         root) have in-arc
                             return false;
     if (!is_connect(root))
     memset(del, 0, sizeof (del));
     min\_cost = 0;
                           fold[i] = fpre[i] = i;
     for (i=0;i<=n;i++)
     while (true ) {
          for (i=1;i<=n;i++)
                           {
              if (del[i] || i == root) continue;
              pre[i] = i;
```

```
path[i][i]
                 = INF; //delete self-cycle
     for (j=0;j<nepath[i].ls;j++)
          int t = nepath[i].adj[j];
          if (del[t])
                      continue ;
          if (path[t][i]
                                    pre[i] ][i])
                          < path[
                                                   pre[i] = fpre[fold[i]]
                                                                           = t;
} //find min in-arc
for (i=1;i<=n;i++)
     if (del[i]
               || i == root) continue ;
    j = i;
     memset(vis, 0, sizeof (vis));
     while (!vis[j] && j != root) {
          vis[j] = true ;
          j = pre[j];
     }
     if (j == root) continue ; //no cycle
     i = j; //cycle begin node
     min_cost += path[ pre[i] ][i];
     for (j=pre[i];
                     j != i ;j=pre[j])
          del[j] = true ; //fold cycle
          min_cost += path[ pre[j]
                                           ][j]; //add cycle cost
     }
     for (j=0;j<nepath[i].ls;j++)
          int t = nepath[i].adj[j];
          if (del[t])
                      continue ;
                       -= path[ pre[i]
          path[t][i]
     } //i is new fold node
                    j != i ;j=pre[j])
     for (j=pre[i];
          for (k=0;k<epath[j].ls;k++)
                                               {
               int t = epath[j].adj[k];
               if (del[t])
                           continue ;
                if (path[i][t]
                               == INF) {
                     epath[i].push_back(t);
                     nepath[t].push_back(i);
               }
                path[i][t]
                            = min(path[i][t],
                                                   path[j][t]);
          }
          for (k=0;k<nepath[j].ls;k++)
                int t = nepath[j].adj[k];
               if (del[t])
                           continue ;
                if (path[t][i]
                               == INF) {
                     epath[t].push_back(i);
                     nepath[i].push_back(t);
                }
```

```
if (path[t][i] > path[t][j]
                                                 - path[ pre[j] ][j]) {
                           path[t][i] = path[t][j]
                                                 path[ pre[j] ][j];
                           fold[i] = j; //record
                                                 fold node
                           fpre[i] = j;
                      }
             } //make
                       new graph
             break ;
         if (i > n) {
             for (i=1;i<=n;i++) {
                  if(del[i] || i == root) continue;
                  min_cost += path[ pre[i] ][i];
             }
             break ;
         } //graph no cycle
    } //while have cycle
    return true;
}
     18.
```

第六章 几何算法

```
* COMPUTATIONAL GEOMETRY ROUTINES
* WRITTEN BY: LIU Yu (C) 2003
*************
   叉乘
//
   两个点的距离
//
   点到直线距离
   返回直线 Ax + By + C =0 的系数
//
   线段
//
//
   员
   两个圆的公共面积
//
//
   矩形
   根据下标返回多边形的边
//
   两个矩形的公共面积
//
   多边形 ,逆时针或顺时针给出 x,y
//
//
   多边形顶点
   多边形的边
//
   多边形的周长
//
```

/************

```
判断点是否在线段上
//
    判断两条线断是否相交,端点重合算相交
//
//
     判断两条线断是否平行
    判断两条直线断是否相交
//
    直线相交的交点
//
     判断是否简单多边形
//
    求多边形面积
//
//
    判断是否在多边形上
    判断是否在多边形内部
//
    点阵的凸包,返回一个多边形
//
    最近点对的距离
//
#include
         <cmath>
#include
         <cstdio>
         <memory>
#include
         <algorithm>
#include
         <iostream>
#include
using namespace
                  std;
        double TYPE;
typedef
        Abs(x)
                (((x)>0)?(x):(-(x)))
#define
                (((x)<0)?(-1):(1))
#define
        Sgn(x)
                  (((a)>(b))?(a):(b))
        Max(a,b)
#define
                  (((a)<(b))?(a):(b))
        Min(a,b)
#define
        Epsilon 1e-10
#define
        Infinity
                 1e+10
#define
        Pi 3.14159265358979323846
#define
TYPE Deg2Rad(TYPE
                     deg)
{ return (deg * Pi / 180.0);}
TYPE Rad2Deg(TYPE
                      rad)
{ return (rad * 180.0 / Pi);}
TYPE Sin(TYPE deg)
{ return sin(Deg2Rad(deg));}
TYPE Cos(TYPE deg)
{ return cos(Deg2Rad(deg));}
TYPE ArcSin(TYPE
{ return Rad2Deg(asin(val));}
```

```
TYPE ArcCos(TYPE val)
{ return Rad2Deg(acos(val));}
TYPE Sqrt(TYPE val)
{ return sqrt(val);}
struct POINT
{
   TYPE x;
   TYPE y;
   TYPE z;
   POINT() : x(0), y(0), z(0) {};
   POINT(TYPE _{x_{-}}, TYPE _{y_{-}}, TYPE _{z_{-}} = 0) : x(_{x_{-}}), y(_{y_{-}}), z(_{z_{-}}) {};
};
// cross product of (o->a) and (o->b)
// 叉乘
TYPE Cross( const POINT & a, const POINT & b, const POINT & o)
{ return (a.x - o.x) * (b.y - o.y) - (b.x - o.x) * (a.y - o.y);}
// planar points'
                   distance
// 两个点的距离
TYPE Distance( const POINT & a, const POINT & b)
{ return Sqrt((a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y - b.y) + (a.z - b.z) * (a.z
 - b.z));}
struct LINE
   POINT a;
   POINT b;
   LINE() {};
   \label{eq:line_point} \mbox{LINE(POINT } \mbox{$\_a$}, \mbox{ POINT } \mbox{$\_b$}) : \mbox{$a(\_a$}), \mbox{$b(\_b$}) \mbox{$\{\}$};
};
// 点到直线距离
double PointToLine(POINT p0 ,POINT p1 ,POINT p2 ,POINT &cp)
{
    double d = Distance(p1, p2);
    double s = Cross(p1 ,p2 ,p0) / d;
    cp.x = p0.x + s*( p2.y-p1.y) / d;
    cp.y = p0.y - s*( p2.x-p1.x) / d;
    return Abs(s);
}
```

```
// 返回直线 Ax + By + C =0 的系数
void Coefficient( const LINE & L, TYPE & A, TYPE & B, TYPE & C)
{
   A = L.b.y - L.a.y;
   B = L.a.x - L.b.x;
   C = L.b.x * L.a.y - L.a.x * L.b.y;
}
void Coefficient( const POINT & p, const TYPE a, TYPE & A, TYPE & B, TYPE & C)
{
   A = Cos(a);
   B = Sin(a);
   C = - (p.y * B + p.x * A);
}
// 线段
struct SEG
{
   POINT a;
   POINT b;
   SEG() {};
   SEG(POINT _a_, POINT _b_):a(_a_),b(_b_)
                                                   {};
};
// 圆
struct CIRCLE
   TYPE x;
   TYPE y;
   TYPE r;
   CIRCLE() {}
   \label{eq:circle} \mathsf{CIRCLE}(\mathsf{TYPE} \quad \_x\_, \quad \mathsf{TYPE} \quad \_y\_, \quad \mathsf{TYPE} \quad \_r\_) \quad : \ x(\_x\_), \qquad y(\_y\_), \qquad r(\_r\_) \quad \{\}
};
POINT Center( const CIRCLE & circle)
{ return POINT(circle.x,
                             circle.y);}
TYPE Area( const CIRCLE & circle)
{ return Pi * circle.r * circle.r;}
// 两个圆的公共面积
TYPE CommonArea(
                       const CIRCLE & A, const CIRCLE & B)
   TYPE area = 0.0;
```

```
const CIRCLE & M = (A.r > B.r) ? A : B;
  const CIRCLE & N = (A.r > B.r) ? B : A;
  TYPE D = Distance(Center(M), Center(N));
  if ((D < M.r + N.r) && (D > M.r - N.r))
     TYPE cosM = (M.r * M.r + D * D - N.r * N.r) / (2.0 * M.r * D);
     TYPE cosN = (N.r * N.r + D * D - M.r * M.r) / (2.0 * N.r * D);
     TYPE alpha = 2.0 * ArcCos(cosM);
     TYPE beta = 2.0 * ArcCos(cosN);
     TYPE TM = 0.5 * M.r * M.r * Sin(alpha);
     TYPE TN = 0.5 * N.r * N.r * Sin(beta);
     TYPE FM = (alpha / 360.0) * Area(M);
     TYPE FN = (beta / 360.0) * Area(N);
     area = FM + FN - TM - TN;
  else if (D \le M.r - N.r)
     area = Area(N);
  return area;
    矩形
    矩形的线段
    3 |
                 | 1
    a -----
    0
struct RECT
                                    // 左下点
  POINT a;
  POINT b;
                                    // 右上点
  RECT() {};
  RECT( const POINT & _a_, const POINT & _b_)
```

}

//

//

//

//

//

//

```
{a = \_a_; b = \_b_;}
};
// 根据下标返回多边形的边
SEG Edge( const RECT & rect, int idx)
   SEG edge;
   while (idx < 0) idx += 4;
   switch (idx % 4)
   {
   case 0:
      edge.a = rect.a;
      edge.b = POINT(rect.b.x,
                               rect.a.y);
      break ;
   case 1:
      edge.a = POINT(rect.b.x,
                               rect.a.y);
      edge.b = rect.b;
      break ;
   case 2:
      edge.a = rect.b;
      edge.b = POINT(rect.a.x,
                                rect.b.y);
      break ;
   case 3:
      edge.a = POINT(rect.a.x,
                                rect.b.y);
      edge.b = rect.a;
      break ;
   default :
      break ;
   return edge;
}
TYPE Area( const RECT & rect)
{ return (rect.b.x
                  - rect.a.x) * (rect.b.y - rect.a.y);}
// 两个矩形的公共面积
TYPE CommonArea(
                     const RECT & A, const RECT & B)
{
   TYPE area = 0.0;
   POINT LL(Max(A.a.x,
                          B.a.x),
                                   Max(A.a.y,
                                               B.a.y));
   POINT UR(Min(A.b.x,
                          B.b.x), Min(A.b.y,
                                              B.b.y));
   if ((LL.x \leftarrow UR.x) \&\& (LL.y \leftarrow UR.y))
```

```
{
     area = Area(RECT(LL, UR));
   return area;
}
// 多边形 ,逆时针或顺时针给出 x,y
struct POLY
  int n; //n 个点
  TYPE * x; //x,y 为点的指针,首尾必须重合
  TYPE * y;
  POLY(): n(0), x(NULL), y(NULL) {};
   POLY( int _n_, const TYPE * _x_, const TYPE * _y_)
   {
     n = _n_;
     x = new TYPE[n + 1];
     memcpy(x, _x_n, _x, _x);
     x[n] = _x_[0];
     y = new TYPE[n + 1];
     memcpy(y, _y_-, _x sizeof (TYPE));
     y[n] = y_[0];
  }
};
// 多边形顶点
POINT Vertex( const POLY & poly, int idx)
  idx %= poly.n;
  return POINT(poly.x[idx], poly.y[idx]);
}
// 多边形的边
SEG Edge( const POLY & poly, int idx)
  idx %= poly.n;
  return SEG(POINT(poly.x[idx], poly.y[idx]),
     POINT(poly.x[idx + 1], poly.y[idx + 1]));
}
// 多边形的周长
TYPE Perimeter( const POLY & poly)
{
```

```
TYPE p = 0.0;
   for (int i = 0; i < poly.n; i++)
      p = p + Distance(Vertex(poly, i), Vertex(poly, i + 1));
   return p;
}
bool IsEqual(TYPE a, TYPE b)
{ return (Abs(a - b) < Epsilon);}
bool IsEqual( const POINT & a, const POINT & b)
{ return (IsEqual(a.x, b.x) && IsEqual(a.y, b.y));}
bool IsEqual( const LINE & A, const LINE & B)
{
   TYPE A1, B1, C1;
   TYPE A2, B2, C2;
   Coefficient(A,
                 A1, B1, C1);
                 A2, B2, C2);
   Coefficient(B,
   return IsEqual(A1 * B2, A2 * B1) &&
      IsEqual(A1 * C2, A2 * C1) &&
     IsEqual(B1 * C2, B2 * C1);
}
// 判断点是否在线段上
bool IsOnSeg( const SEG & seg, const POINT & p)
{
   return (IsEqual(p,
                    seg.a) || IsEqual(p, seg.b)) ||
      (((p.x - seg.a.x) * (p.x - seg.b.x) < 0 ||
      (p.y - seg.a.y) * (p.y - seg.b.y) < 0) &&
      (IsEqual(Cross(seg.b,
                              p, seg.a),
                                        0)));
}
// 判断两条线断是否相交,端点重合算相交
bool IsIntersect(
                 const SEG & u, const SEG & v)
{
                                              v.b, u.a) >= 0) &&
   return (Cross(v.a,
                      u.b, u.a) * Cross(u.b,
      (Cross(u.a,
                                         u.b, v.a) >= 0) &&
                  v.b, v.a) * Cross(v.b,
                                         v.b.x))
      (Max(u.a.x,
                          >= Min(v.a.x,
                   u.b.x)
                                                 &&
      (Max(v.a.x,
                         >= Min(u.a.x,
                  v.b.x)
                                         u.b.x))
                                                 &&
      (Max(u.a.y,
                   u.b.y)
                         >= Min(v.a.y,
                                         v.b.y))
                                                 &&
      (Max(v.a.y,
                  v.b.y)
                             Min(u.a.y,
                                         u.b.y));
                         >=
}
```

```
// 判断两条线断是否平行
bool IsParallel( const LINE & A, const LINE & B)
{
   TYPE A1, B1, C1;
   TYPE A2, B2, C2;
                 A1, B1, C1);
   Coefficient(A,
   Coefficient(B,
                 A2, B2, C2);
   return (A1 * B2 == A2 * B1) &&
     ((A1 * C2 != A2 * C1) || (B1 * C2 != B2 * C1));
}
// 判断两条直线断是否相交
bool IsIntersect( const LINE & A, const LINE & B)
{ return !IsParallel(A,
                      B);}
// 直线相交的交点
                    const LINE & A, const LINE & B)
POINT Intersection(
   TYPE A1, B1, C1;
   TYPE A2, B2, C2;
   Coefficient(A,
                 A1, B1, C1);
   Coefficient(B,
               A2, B2, C2);
   POINT I(0, 0);
   I.x = - (B2 * C1 - B1 * C2) / (A1 * B2 - A2 * B1);
   I.y = (A2 * C1 - A1 * C2) / (A1 * B2 - A2 * B1);
   return I;
}
bool IsInCircle( const CIRCLE & circle, const RECT & rect)
{
   return (circle.x - circle.r >= rect.a.x) &&
             + circle.r <= rect.b.x)
     (circle.x
                                      &&
     (circle.y - circle.r >= rect.a.y) &&
     (circle.y + circle.r <= rect.b.y);
}
// 判断是否简单多边形
bool IsSimple( const POLY & poly)
{
```

```
if (poly.n < 3)
      return false;
   SEG L1, L2;
   for (int i = 0; i < poly.n - 1; i++)
      L1 = Edge(poly, i);
      for (int j = i + 1; j < poly.n; j++)
         L2 = Edge(poly, j);
         if (j == i + 1)
         {
                                  if (IsOnSeg(L1, L2.b) || IsOnSeg(L2,
                                                                                 L1.a))
return
       false;
        }
         else if (j == poly.n - i - 1)
         {
                                  if (IsOnSeg(L1, L2.a) || IsOnSeg(L2,
                                                                                 L1.b))
       false;
return
        }
         else
         {
            if (IsIntersect(L1, L2))
                                     return false;
        }
     } // for j
   } // for i
   return true;
}
// 求多边形面积
TYPE Area( const POLY & poly)
{
   if (poly.n < 3) return TYPE(0);
   double s = poly.y[0] * (poly.x[poly.n - 1] - poly.x[1]);
   for (int i = 1; i < poly.n; i++)
   {
      s += poly.y[i] * (poly.x[i - 1] - poly.x[(i + 1) % poly.n]);
   return s/2;
}
// 判断是否在多边形上
bool IsOnPoly( const POLY & poly, const POINT & p)
   for (int i = 0; i < poly.n; i++)
```

```
{
     if (IsOnSeg(Edge(poly, i), p)) return true;
   return false;
}
// 判断是否在多边形内部
bool IsInPoly( const POLY & poly, const POINT & p)
   SEG L(p, POINT(Infinity,
                           p.y));
  int count = 0;
  for (int i = 0; i < poly.n; i++)
     SEG S = Edge(poly,
                         i);
     if (IsOnSeg(S,
                    p))
                       // 如果想让 在 poly 上则返回 true, 则改为 true
        return false;
     }
     if (!IsEqual(S.a.y, S.b.y))
        POINT & q = (S.a.y > S.b.y)?(S.a):(S.b);
        if (IsOnSeg(L,
                      q))
           ++count;
        else if (!IsOnSeg(L, S.a) && !IsOnSeg(L, S.b) && IsIntersect(S,
                                                                     L))
           ++count;
  return (count % 2 != 0);
}
// 点阵的凸包,返回一个多边形
POLY ConvexHull( const POINT * set, int n) // 不适用于点少于三个的情况
{
   POINT * points = new POINT[n];
   memcpy(points, set, n * sizeof (POINT));
  TYPE * X = new TYPE[n];
  TYPE * Y = new TYPE[n];
  int i, j, k = 0, top = 2;
```

```
for (i = 1; i < n; i++)
   if ((points[i].y
                     < points[k].y)
                                          &&
      ((points[i].y
                      == points[k].y)
      (points[i].x
                     < points[k].x)))
      k = i;
std::swap(points[0],
                          points[k]);
for (i = 1; i < n - 1; i++)
{
   k = i;
   for (j = i + 1; j < n; j++)
      if ((Cross(points[j],
                                                         > 0) |
                               points[k],
                                            points[0])
          ((Cross(points[j],
                                points[k],
                                             points[0])
                                                          == 0) &&
          (Distance(points[0],
                                   points[j])
                                                < Distance(points[0],
                                                                           points[k]))))
          k = j;
      }
   std::swap(points[i],
                            points[k]);
                        Y[0] = points[0].y;
X[0] = points[0].x;
                        Y[1] = points[1].y;
X[1]
     = points[1].x;
X[2]
                        Y[2] = points[2].y;
    = points[2].x;
for (i = 3; i < n; i++)
   while (Cross(points[i], POINT(X[top], Y[top]),
      POINT(X[top - 1], Y[top - 1])) >= 0 && top>0)
      top--;
   ++top;
   X[top] = points[i].x;
   Y[top] = points[i].y;
}
delete [] points;
```

```
POLY poly(++top,
                   X, Y);
   delete [] X;
   delete [] Y;
   return poly;
}
// 最近点对的距离 , Written By PrincessSnow
#define MAXN 100000
POINT pt[MAXN];
bool cmp(POINT n1, POINT n2)
{ return (n1.x < n2.x || n1.x == n2.x && n1.y < n2.y);}
double Get( double dis, int mid, int start, int end)
{
   int s=mid, e=mid, i, j;
   double t;
   while (s > start && pt[mid].x - pt[s].x <= dis)
                                                       S--;
   while (e < end && pt[e].x - pt[mid].x <= dis)
                                                      e++;
   for (i=s; i \le e; i++)
     for (j=i+1; j \le e \&\& j \le i+7; j++) {
        t = Distance(pt[i],
                            pt[j]);
        if (t < dis)
                      dis=t;
     }
    return dis;
}
double ClosestPairDistance(
                          int start, int end)
   int m = end-start+1, mid, i;
   double t1, t2, dis=-1,
                          t;
   if (m <= 3)
     for (i=start; i < end; i++) {
       t = Distance(pt[i], pt[i+1]);
       if (t < dis || dis == -1) dis = t;
     }
                         , pt[end]);
     t = Distance(pt[start]
      if (t < dis) dis=t;
      return dis;
   }
   if (m\%2 == 0) mid = start + m/2 - 1;
        mid = start + m/2;
   if (m\%2 == 0) {
```

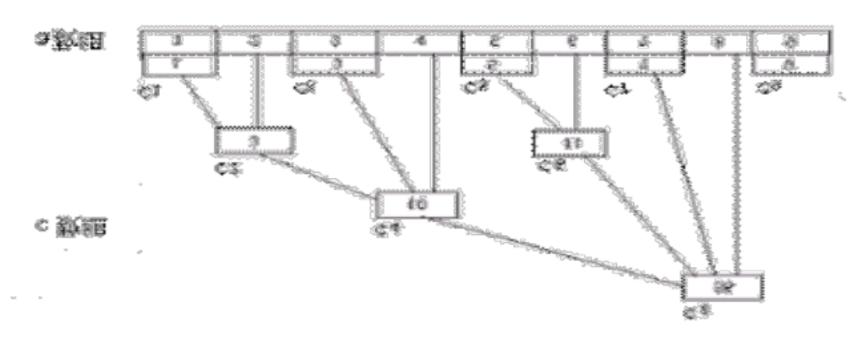
```
t1 = ClosestPairDistance(start,
                                      mid);
     t2 = ClosestPairDistance(mid+1,
                                        end);
  }
   else {
     t1 = ClosestPairDistance(start,
                                      mid);
     t2 = ClosestPairDistance(mid+1,
                                        end);
  }
   if (t1 < t2) dis = t1;
              dis = t2;
   else
   dis = Get(dis, mid, start, end);
   return dis;
}
          球面上两点最短距离
    1.
    // 计算圆心角 lat 表示纬度 , -90 <= w <= 90, lng 表示经度
    // 返回两点所在大圆劣弧对应圆心角
                                   , 0 <= angle <= pi
    double angle( double lng1, double lat1, double lng2, double lat2)
    {
       double dlng = fabs(lng1 - lng2) * pi / 180;
       while (dlng >= pi+pi)
                                dlng -= pi+pi;
       if (dlng > pi)
                       dlng = pi + pi - dlng;
       lat1 *= pi / 180, lat2 *= pi / 180;
       return acos( cos(lat1)*cos(lat2)*cos(dlng)
                                                  + sin(lat1)*sin(lat2)
                                                                          );
    }
    // 计算距离 , r 为球半径
    double line_dist( double r, double lng1, double lat1, double lng2, double lat2)
       double dlng = fabs(lng1 - lng2) * pi / 180;
        while (dlng >= pi+pi)
                                 dlng -= pi+pi;
       if (dlng > pi)
                       dlng = pi + pi - dlng;
       lat1 *= pi / 180, lat2 *= pi / 180;
                                                               + sin(lat1)*sin(lat2)
       return r * sqrt( 2 - 2*( cos(lat1)*cos(lat2)*cos(dlng)
     );
    }
    // 计算球面距离 , r 为球半径
    double sphere_dist( double r, double lng1, double lat1, double lng2, double lat
    2)
    {
       return r * angle(lng1, lat1, lng2, lat2);
    }
```

2. 三点求圆心坐标

```
double GetRadiusBy3Points(
                        double x1, double y1,
                   double x2, double y2,
                   double x3, double y3,
                   double &x, double &y)
{
   // 由 (x - x1)^2 + (y - y1)^2 = (x - x2)^2 + (y - y2)^2 得
   // 2*( x2 - x1 )*x + 2*( y2 - y1 )*y = x2^2 - x1^2 + y2^2 - y1^2
   // 同理得
   // 2*( x3 - x2 )*x + 2*( y3 - y2 )*y = x3^2 - x2^2 + y3^2 - y2^2
   // 由行列式解方程得 x , y
   double a11, a12, a21, a22, b1, b2;
   double d, d1, d2;
   a11 = 2 * (x3 - x2);
   a12 = 2 * (y3 - y2);
   a21 = 2 * (x2 - x1);
   a22 = 2 * (y2 - y1);
   b1 = x3*x3 - x2*x2 + y3*y3 - y2*y2;
   b2 = x2*x2 - x1*x1 + y2*y2 - y1*y1;
   d = a11*a22 - a12*a21;
   d1 = b1*a22 - a12*b2;
   d2 = a11*b2 - b1*a21;
  // x , y 是圆心坐标
   x = d1 / d;
   y = d2 / d;
   return (x1 - x)^*(x1 - x) + (y1 - y)^*(y1 - y);
}
    三角形几个重要的点
3.
设三角形的三条边为 a, b, c , 且不妨假设 a <= b <= c
三角形的面积可以根据海伦公式算得,如下:
s = sqrt(p * (p - a) * (p - b) * (p - c)), p = (a + b + c) / 2
1. 费马点(该点到三角形三个顶点的距离之和最小 )
有个有趣的结论:若三角形的三个内角均小于 120 度,
那么该点连接三个顶点形成的三个角均为 120 度;若三角形存在一个内角
大于 120 度,则该顶点就是费马点)
计算公式如下 :
若有一个内角大于 120 度 (这里假设为角 C) ,则距离为 a + b
若三个内角均小于 120 度,则距离为
sqrt((a * a + b * b + c * c + 4 * sqrt(3.0) * s) / 2), 其中
```

2. 内心 ---- 角平分线的交点

第七章 专题讨论



1. 树状数组

* Function Name : 树状数组
* Description : HDOJ 1166 敌兵布阵
* 减少冗余统计 , 是线段树的一种变化
**** **** **** **** ****/
#include <cstdio>
int data[50001], s[50001], T[50001];

inline int lowbit(int t)
{ return t & (-t);}

inline int sum(int end)

```
{
   int sum = 0;
   while (end > 0) {
       sum += T[end];
       end -= lowbit(end);
   return sum;
}
inline void plus( int pos, int num, int count)
   while (pos <= count) {
       T[pos] += num;
       pos += lowbit(pos);
   }
}
int main()
{
   char buffer[10];
   int i, j, t, n, a, b;
   scanf("%d", &t);
   for (i=1; i \le t ; i++) {
       scanf("%d", &n);
       T[0] = s[0] = data[0] = 0;
       for (j=1; j \le n ; j++)
           scanf("%d", &data[j]);
           s[j] = s[j - 1] + data[j];
           T[j] = s[j] - s[j - lowbit(j)];
       printf("Case
                     %d:\n",
                              i);
       while (scanf("%s",
                         buffer) == 1 && buffer[0] != 'E') {
           scanf("%d%d",
                            &a, &b);
           switch (buffer[0])
                              {
           case 'Q':
               printf("%d\n",
                            sum(b) - sum(a) + data[a]);
                                                                break ;
           case 'A':
               plus(a,
                       b, n); break;
           case 'S':
                       -b, n); data[a] -= b; break ;
               plus(a,
       }
   }
}
```

2. 字典树

```
Name : 字典树(多路查找树)
    Function
    Description
                               HDOJ 1075 What Are You Talking About
                          易于字符保存 ,插入和查找 ,时间复杂度都是线性
          <cstdio>
#include
#include
          <string>
using namespace
                   std;
struct trie
   trie * next[26];
   int index;
};
trie *thead;
char dic[1000000][20];
inline trie * newnode()
{
   int i;
   trie *t;
   t=(trie*)malloc(
                    sizeof (trie));
                sizeof (trie));
   memset(t,0,
   return t;
void insert(trie * s, char x[], int pos)
   int i;
   trie *t;
   for (i=0; x[i] ; i++) {
      if ( s->next[ x[i]-'a' ] ) s=s->next[ x[i]-'a' ];
      else {
       t=newnode();
        s->next[ x[i]-'a' ]=t;
         s=t;
   } //for
   s->index=pos;
}
void deltrie(trie * s)
```

```
{
   int i;
   for (i=0; i < 26; i++) {
      if ( s->next[i]
          deltrie(s->next[i]);
   free(s);
   s=NULL;
int find(trie *s, char x[])
   int i;
   if (x[0] == 0) return -1;
   for (i=0; x[i] ; i++) {
      if ( s->next[
                     x[i]-'a' ] ) s=s->next[
                                                     x[i]-'a' ];
      else
                                     break ;
   if (x[i]==0)
                     return s->index;
                return -1;
   else
}
int main()
   int t,n,i,j,all;
   char word[20],mars[20],ch;
   thead=newnode();
   while (scanf("%s",word)==1)
      if (word[0]=='S')
                               break ;
    i=1;
                                        && dic[i][0]!='E')
    while (scanf("%s",dic[i])==1
        scanf("%s",mars);
        insert(thead,mars,i);
        i++;
    }
    all=i;
    while (scanf("%s",word)==1)
        if (word[0]=='S')
                                 break ;
    getchar();
                    j=0;
                                    && ch!='E') {
    while (scanf("%c",&ch)==1
                     && ch<='z')
        if (ch>='a'
             mars[j]=ch;
                                j++;
        }
```

```
else {
    mars[j]=0;
    t=find( thead , mars );
    j=0;
    if (t>0)    printf("%s",dic[t]);
    else if (mars[0]!=0)    printf("%s",mars);
    printf("%c",ch);
    }
} //while
    deltrie(thead);
}
```

3. 后缀树

```
Name:
                               后缀树
    Function
    Description
                          PKU 2774 Long Long Message
                           有效的支持字符串匹配和查询
#include
         <cstdio>
#include
         <string>
#define
         NUM
                             27
         STARTCHAR
                             'a'
#define
         SPECIALCHAR
#define
         ERROR
#define
                            -1
         TYPE1
#define
                           1
                           2
#define
         TYPE2
        LEAF
                          1
#define
#define
         NOTLEAF
                              2
struct SuffixTrie
   int Start, End;
   SuffixTrie\\
              * Next[NUM];
   SuffixTrie
              * Link;
   SuffixTrie
              * Father;
   int Flag;
```

```
int Length;
};
char str[100010], buf[100010];
SuffixTrie head;
SuffixTrie^*P \quad , \ ^*G, \quad ^*U, \quad ^*V, \quad ^*q;
int W[3], len, len2;
void CreateNode(SuffixTrie * & Node) {
   int i;
   Node = (SuffixTrie * ) malloc( sizeof (SuffixTrie));
   Node -> Start = Node -> End = Node -> Length = ERROR;
   for (i = 0; i < NUM; i++) Node -> Next[i] = NULL;
   Node -> Link = Node -> Father = NULL;
   Node -> Flag = LEAF;
}
void Init(SuffixTrie & h, char s[]) {
   int i;
   h.Start = h.End = ERROR;
   for (i = 0; i < NUM; i++) h.Next[i]
                                        = NULL;
   h.Link = & h;
   h.Father = NULL;
   h.Flag = LEAF;
   h.Length = 0;
   len = strlen(s);
   s[len] = SPECIALCHAR;
   s[len + 1] = '\0';
   len++;
int FindV( char s[]) {
   int old;
   SuffixTrie * t, * newt;
   t = U \rightarrow Next[s[W[0]] - STARTCHAR];
   old = 0;
   while (W[2] > (t -> End) - (t -> Start) + 1 + old) {
     old += (t -> End - t -> Start + 1);
    t = t \rightarrow Next[s[W[0] + old] - STARTCHAR];
  }
   if (W[2] == (t -> End) - (t -> Start) + 1 + old) {
     V = t;
```

```
P \rightarrow Link = V;
     return TYPE1;
  } else {
     CreateNode(newt);
     newt -> Start = t -> Start;
     newt -> End = t -> Start + W[2] - old - 1;
     newt -> Father = t -> Father;
     newt ->
           Length = newt -> Father -> Length + newt -> End - newt ->
           Start + 1;
     t -> Father -> Next[s[t -> Start] - STARTCHAR] = newt;
     t \rightarrow Start = newt \rightarrow End + 1;
     newt -> Next[s[t -> Start] - STARTCHAR] = t;
     t -> Father = newt;
     V = newt;
      P \rightarrow Link = V;
      return TYPE2;
}
int Insert(SuffixTrie * Node, int start, char s[]) {
  int i, posbegin,
                   posend;
   SuffixTrie * t;
  if (Node -> Next[s[start]
                             - STARTCHAR] == NULL) {
      CreateNode(Node
                      -> Next[s[start]
                                        STARTCHAR]);
     Node -> Next[s[start]
                             - STARTCHAR] -> Start = start;
      Node -> Next[s[start]
                             - STARTCHAR] -> End = len - 1;
      Node -> Next[s[start]
                             STARTCHAR] -> Father = Node;
      Node -> Next[s[start]
                             - STARTCHAR]
           Length = Node -> Length + len - start;
      Node -> Flag = NOTLEAF;
      P = Node;
      return TYPE1;
  } else {
     posbegin = Node -> Next[s[start] - STARTCHAR] -> Start;
     posend = Node -> Next[s[start] - STARTCHAR] -> End;
     for (i = posbegin; i \le posend; i++) {
        if (s[i] != s[start + i - posbegin])
                                            break ;
     if (i == posend + 1) {
         return Insert(Node -> Next[s[start] - STARTCHAR], start + i - posbe
gin, s);
     } else {
        CreateNode(t);
```

```
t -> Start = posbegin;
         t -> End = i - 1;
         t -> Flag = NOTLEAF;
         Node -> Next[s[start] - STARTCHAR] -> Start = i;
                        - STARTCHAR] = Node -> Next[s[start] - STARTCHAR];
         t -> Next[s[i]
         t -> Next[s[i] - STARTCHAR] -> Father = t;
         Node -> Next[s[start] - STARTCHAR] = t;
         t -> Father = Node;
         t -> Length = Node -> Length + t -> End - t -> Start + 1;
         Insert(t, start + i - posbegin, s);
         G = Node;
         P = t;
         return TYPE2;
      }
   }
}
int Select( int start, char s[], int type) {
   int result1, result2, result;
   if (type == TYPE1) {
      U = P \rightarrow Link;
      result = Insert(U, start + U -> Length, s);
   } else {
      U = G \rightarrow Link;
      if (G \rightarrow Link == G) {
         W[0] = P \rightarrow Start + 1;
         W[1] = P \rightarrow End;
         W[2] = P \rightarrow End - P \rightarrow Start;
      } else {
         W[0] = P \rightarrow Start;
         W[1] = P \rightarrow End;
         W[2] = P -> End - P -> Start + 1;
      }
      if (W[2] == 0) {
        V = G;
         P \rightarrow Link = V;
         result = Insert(V, start, s);
      } else {
         result1 = FindV(s);
         result2 = Insert(V, start + V -> Length, s);
         if (result1 == TYPE2) {
           G = P \rightarrow Father;
           result = result1;
         } else result = result2;
```

```
}
  return result;
}
void BuildSuffixTrie(SuffixTrie & h, char s[]) {
  int i;
  int type;
  len = strlen(s);
  CreateNode(h.Next[s[0] - STARTCHAR]);
  h.Next[s[0] - STARTCHAR] -> Start = 0;
  h.Next[s[0] - STARTCHAR] -> End = len - 1;
  h.Next[s[0] - STARTCHAR] -> Father = & h;
   h.Next[s[0] - STARTCHAR] -> Length = h.Length + h.Next[s[0] - STARTC
HAR] -> End - h.Next[s[0] - STARTCHAR] -> Start + 1;
  h.Flag = NOTLEAF;
  type = TYPE1;
  P = \& h;
  for (i = 1; i < len; i++) type = Select(i, s, type);
}
int i;
  for (i = 0; i < NUM; i++) {
     if (Node -> Next[i] != NULL) {
        DeleteSuffixTrie(Node -> Next[i]);
        Node -> Next[i] = NULL;
  free(Node);
int FindString( int start, char s[]) {
  int result;
  int i;
  int temp;
  SuffixTrie * x;
  x = P \rightarrow Next[s[start]]
                        - STARTCHAR];
  result = P -> Length;
  if (x == NULL) {
     P = P \rightarrow Link;
     return result;
```

```
}
   temp = 0;
   for (i = start; i < len2; i++) {
      if (x -> Start + i - start - temp > x -> End) {
         temp = i - start;
         P = x;
         x = x \rightarrow Next[s[start + temp] - STARTCHAR];
         if (x == NULL) break;
      }
      if (s[i] != str[x -> Start + i - start - temp]) break;
      result++;
   P = P \rightarrow Link;
   return result;
}
int Search(SuffixTrie & h, char s[]) {
   int result;
   int i;
   int temp;
   len2 = strlen(s);
   result = 0;
   P = \& head;
   for (i = 0; i < len2; i++) {
      temp = FindString(i + P -> Length,
                                               s);
      if (result < temp) result = temp;
      if (result >= len2 - i) break;
   return result;
}
int main() {
   int result;
   while (scanf("%s",
                        str) != EOF) {
      Init(head,
                  str);
      BuildSuffixTrie(head,
                              str);
      scanf("%s",
                    buf);
      result = Search(head,
                                buf);
      printf("%d\n",
                       result);
}
```

4. 线段树

```
[1,10]
                        [5,10]
      [1,5]
                            [7,10]
          [3,5]
                    [5,7]
   [1,3]
                 [5,6] [6,7]
                           [7,8] [8,10]
 [1,2] [2,3] [3,4] [4,5]
                           [8,9] [9,10]
    Function
                               线段树
              Name:
    Description
                          HDOJ 1542 Atlantis
                          用于表示区间线段
#include <cstdio>
#include <algorithm>
using namespace
                   std;
typedef struct ITREE_NODE {
   ITREE_NODE * pLChild, * pRChild;
                             // 左端点,右端点
   double left, right;
                              // 测度
   double measure;
   int count;
                          // 覆盖计数器
                               // 独立线段数
   int lines;
   int Ibound, rbound; // 覆盖左、右顶点的线段数目
}*PITREE_NODE;
inline void safe_add( int & v, int value) {
   v += value;
   if (v < 0) v = 0;
}
void itree_splite( const double * pList, PITREE_NODE pParent, const int iLeft,
const int iRight) {
  if (iRight - iLeft <= 1) return ;
  int iMid = (iLeft + iRight) >> 1;
   pParent -> pLChild = new ITREE_NODE;
   pParent -> pRChild = new ITREE_NODE;
                     -> pLChild, 0, sizeof (ITREE_NODE));
   memset(pParent
   memset(pParent
                     -> pRChild, 0, sizeof (ITREE_NODE));
   pParent -> pLChild -> left = pList[iLeft];
   pParent -> pLChild -> right = pList[iMid];
   pParent -> pRChild -> left = pList[iMid];
   pParent -> pRChild -> right = pList[iRight];
  itree_splite(pList,
                      pParent -> pLChild, iLeft, iMid);
```

```
itree_splite(pList,
                     pParent -> pRChild, iMid, iRight);
}
PITREE_NODE itree_generate( const double * pList, const int iListCount) {
  PITREE_NODE pRoot = new ITREE_NODE;
  memset(pRoot, 0, sizeof (ITREE_NODE));
  pRoot -> left = pList[0];
  pRoot -> right = pList[iListCount - 1];
                   pRoot, 0, iListCount - 1);
  itree_splite(pList,
  return pRoot;
}
void itree_destroy(PITREE_NODE
                                  pParent) {
  if (pParent == NULL) return ;
  if (pParent -> pLChild) itree_destroy(pParent -> pLChild);
  if (pParent -> pRChild) itree_destroy(pParent
                                              -> pRChild);
  delete pParent;
}
inline void itree_measure(PITREE_NODE
                                         pNode) {
  if (pNode -> count > 0)
     pNode -> measure = pNode -> right - pNode -> left;
  else if (pNode -> pLChild && pNode -> pRChild)
      pNode -> measure = pNode -> pLChild -> measure + pNode -> pRChil
d -> measure;
  else
       pNode -> measure = 0;
}
inline void itree_lines(PITREE_NODE
                                     pNode) {
  if (pNode -> count > 0) {
     pNode \rightarrow lines = 1;
  } else if (pNode -> pLChild && pNode -> pRChild) {
     if (pNode -> pLChild -> rbound && pNode -> pRChild -> lbound) {
        pNode -> lines = pNode -> pLChild -> lines + pNode -> lines - 1;
     } else {
        pNode -> lines = pNode -> pLChild -> lines + pNode -> lines;
     }
  } else {
     pNode -> lines = 0;
}
// 插入的时候        value = 1, 删除的时候        value = -1
```

```
const double left, const double
void itree_update(PITREE_NODE
                                      pParent,
t,
              int value) {
   if (pParent -> left == left && pParent -> right == right) {
      safe_add(pParent
                                      value);
                           -> count,
      safe_add(pParent
                                        value);
                             lbound,
      safe_add(pParent
                           -> rbound,
                                        value);
      itree_measure(pParent);
      itree_lines(pParent);
   } else {
      if (pParent -> pLChild -> right > left) {
         if (pParent -> pLChild -> right >= right)
            itree_update(pParent
                                     -> pLChild,
                                                  left, right, value);
         } else {
            itree_update(pParent
                                     -> pLChild,
                                                   left,
                   pParent -> pLChild -> right, value);
            itree_update(pParent
                                     -> pRChild,
                                                   pParent -> pRChild -> left,
                   right, value);
         }
      } else {
         itree_update(pParent
                                  -> pRChild,
                                                left, right,
                                                            value);
      itree_measure(pParent);
      itree_lines(pParent);
      if (left == pParent -> left) safe_add(pParent
                                                         -> lbound,
                                                                      value);
       if (right == pParent -> right) {
         safe_add(pParent
                              -> rbound,
                                           value);
}
                                    pParent, const double left, const double right)
void itree_insert(PITREE_NODE
                           left, right, 1); }
{itree_update(pParent,
void itree_delete(PITREE_NODE
                                    pParent, const double left, const double right)
                           left, right, -1); }
{itree_update(pParent,
struct EVENT {
    double x, y1, y2;
    int type;
};
bool cmp( const EVENT & a, const EVENT & b)
\{ return a.x < b.x; \}
```

```
PITREE_NODE pRoot;
EVENT env[200];
double Y[200];
double tsize = 0.0;
int main() {
  double x1, x2, y1, y2;
  int i, n, n2, cas = 0;
   while (scanf("%d", & n) == 1 && n) {
     cas++;
     n2 = n << 1;
     for (i = 0; i < n2; i += 2) {
        scanf("%lf%lf%lf%lf", & x1, & y1, & x2, & y2);
        env[i].x = x1;
        env[i].y1 = y1;
        env[i].y2 = y2;
        env[i].type = 1;
        env[i + 1].x = x2;
        env[i + 1].y1 = y1;
        env[i + 1].y2 = y2;
        env[i + 1].type = -1;
        Y[i] = y1;
        Y[i + 1] = y2;
     sort(env, env + n2, cmp);
     sort(Y, Y + n2);
     pRoot = itree_generate(Y , n2);
     for (i = 0; i < n2; ++i) {
        if (i > 0) tsize += pRoot -> measure * (env[i].x - env[i - 1].x);
        else tsize = 0.0;
        itree_update(pRoot,
                          env[i].y1, env[i].y2,
                                                    env[i].type);
     itree_destroy(pRoot);
      printf("Test
                  case #%d\nTotal
                                    explored area: %.2lf\n\n",
                                                                 cas, tsize);
   return 0;
}
      并查集
5.
                              并查集
    Function
             Name:
    Description
                              集合操作,并,除,判断
                              ****/
```

```
const int Max=1000;
typedef int ElemType;
int Parent[Max],Rank[Max];
int Find( int x)
   int temp = x, root, w;
   // 搜寻根节点
   while (Parent[x]!=0)
                              x=Parent[x];
   root=x;
   x=temp;
   // 压缩路径
   while (Parent[x]!=0)
                          {
      w=Parent[x];
      Parent[x]=root;
      x=w;
   return root;
int Union( int x, int y)
   int u, v, root;
   u=Find(x);
   v=Find(y);
               <= Rank[v])
   if (Rank[u]
      root = Parent[u] = v;
      if (Rank[u]
                   == Rank[v])
                                      Rank[v]++;
           root=Parent[v]=u;
   else
   return root;
}
```

6. 二叉堆

```
/**** **** **** **** **** ****

* Function Name : 二叉堆

* Description : 父结点的键值总是大於或等於任何一个子节点的键值

* 便於寻找父节点和子节点

**** **** **** **** ****/
```

```
const int Max=1000;
typedef int ElemType;
ElemType Heap[Max];
int Sift_Up( int i)
                   // 上移
   ElemType
             temp;
   bool flag;
   flag = true ;
   if (i == 1)
                 return 0;
   do {
      if (Heap[i] > Heap[i/2])
      {temp=Heap[i];
                            Heap[i]=Heap[i/2];
                                             Heap[i/2]=temp;}
              flag = false;
      else
      i /= 2;
   } while (i>1 || flag);
   return 1;
}
int Sift_Down( int i, int n) // 下移
{
   bool flag;
   ElemType temp;
   flag = false ;
   if (2*i > n)
                  return 0;
   do {
      i*=2;
      if (i+1 <= n && Heap[i+1]
                                   > Heap[i])
                                                   i++;
                    < Heap[i])
      if (Heap[i/2]
      {temp=Heap[i];
                                                     Heap[i/2]=temp;}
                           Heap[i]=Heap[i/2];
              flag = false;
      else
   } while (2*i \le n \quad || flag);
   return 1;
}
int Insert( int &n,ElemType x) // 插入元素
   Heap[++n] = x;
   if ( Sift_Up(n) ) return n;
}
int Delete( int &n, int i) // 输出元素
{
```

```
ElemType x,y;
  x = Heap[i]; y = Heap[n];
   n--;
  if (i == n+1) return x;
  Heap[i] = y;
                Sift_Up(i);
  if (y >= x)
             Sift_Down(i,n);
   else
   return x;
}
int Delete_Max( int &n) // 输出最大值
  ElemType x;
  x = Heap[1];
  Delete(n,1);
  return x;
int Make_Heap( int n) // 转换为大顶堆
{
  int i;
  for (i=n/2; i \ge 1; i--) Sift_Down(i,n);
  return n;
}
int HeapSort( int n) // 非降序排序
  int i;
  ElemType temp;
  Make_Heap(n);
  for (i=n; i >= 2; i--) {
      temp=Heap[i];
                         Heap[i]=Heap[1];
                                               Heap[1]=temp;
      Sift_Down(1,i-1);
  }
   return 1;
}
      逆序数(归并排序)
7.
                           逆序数 (归并排序 )
    Function
             Name:
                              N*Log(N)
    Description
                              ****/
// 逆序数值存放在
                anti 中
int p[MAX], t[MAX],
                     anti = 0;
```

```
void merge( int first, int last)
{
   int mid = (first+last)/2;
    int i1 = 0, i2 = first, i3 = mid+1;
    while (i2 <= mid && i3 <= last) {
       if (p[i2] > p[i3]) {
           t[i1++]
                    = p[i3++];
           anti += mid-i2+1;
       }
       else t[i1++]=p[i2++];
    }
    while (i2 <= mid) t[i1++] = p[i2++];
    while (i3 <= last)
                       t[i1++]
                                  = p[i3++];
   i1 = first; i2 = 0;
    while (i2 < last-first+1)
                           p[i1++] = t[i2++];
}
void merge_sort( int first, int last)
{
   int mid;
    if (first<last)
               {
       mid = (first+last)/2;
       merge_sort(first,
                           mid);
       merge_sort(mid+1,
                              last);
       merge(first,
                     last);
   }
}
      树状 DP
8.
                             树状 DP
             Name :
    Function
    Description
                           HDOJ 1561 The more, The Better
                                ****/
#include <cstdio>
         <memory>
#include
#include <queue>
using namespace
                    std;
#define
       Max 210
int n,m,a[Max][Max];
struct inf
{
   int I,r,p;
   int v;
}tree[Max];
```

```
int tp,now;
queue< int > SQ;
char v[Max];
int main()
{
   int i,j;
   int root,pt,tv;
   while (scanf("%d%d",&n,&m))
      if (n==0 &\& m==0)
                                  break ;
                     sizeof (tree));
      memset(tree,0,
      memset(a,0,
                     sizeof (a));
                     sizeof (v));
      memset(v,3,
      while (!SQ.empty())
                                 SQ.pop();
      for (i=1; i \le n; i++)
         scanf("%d%d",
                            &root, &tree[i].v);
          if (tree[root].l
                          == 0) {
             tree[root].l
                           = i;
             v[root]--;
             tree[i].p
                       = root;
         }
          else {
             pt = tree[root].l;
                             != 0) pt = tree[pt].r;
             while (tree[pt].r
             tree[pt].r
                       = i;
             v[pt] -= 2;
                     = pt;
             tree[i].p
      for (i=1;i<=n;i++)
          if (v[i]==3)
                           SQ.push(i);
        while (!SQ.empty())
            now = SQ.front();
            SQ.pop();
            a[now][1]
                         = tree[now].v;
            for (i=1; i \le m ; i++)
                            = a[now][i]
                                           < a[ tree[now].r
                                                               ][i] ? a[ tree[now].r
                 a[now][i]
                                                                                         ][i]
: a[now][i];
            for (i=2; i \le m ; i++)
                for (j=1; j \le i ; j++) {
                     tv = a[tree[now].l ][j-1] + tree[now].v + a[tree[now].r]
                                                                                          ][i
-j];
                    a[now][i] = a[now][i] < tv ? tv : a[now][i];
```

```
}
                 tree[now].p ].l == now) v[ tree[now].p
          if (tree[
                                                           ]++;
          else
                                        v[ tree[now].p
                                                      ] += 2;
          if (v[tree[now].p] == 3) SQ.push(tree[now].p);
      }
      printf("%d\n",a[
                   tree[0].l
                                ][m]);
}
9.
     欧拉路
                         欧拉路
    Function
            Name :
   Description
                        ZJU 2730 Necklace
   欧拉路的构造方法:
   若图连同且度为奇数的节点不超过
                                 2 个,则该图可以构造出欧拉路
   先选一个度为奇数的节点
                        (若没有就任选一个度为偶数的节点
   再以该节点为起点 ,用 dfs 遍历所有的弧 (每条弧只遍历一次 ),遇到死胡同就回溯
    在每次回溯时将所在弧按顺序记录下来,这组弧的排列就组成了一条欧拉路
#include <stdio.h>
#define MAXN 50
void find_path_euler(
                    int n, int mat[][MAXN], int now, int & step, int * path)
{
   int i;
   for (i=n-1; i >= 0; i--)
      while (mat[now][i])
          mat[now][i]--,
                         mat[i][now]--;
          find_path_u(n,
                         mat, i, step, path);
         step++ ]=now;
   path[
int main()
{
   int n;
   int a[MAXN][MAXN];
   int i, j, cnt, mmin;
   int b[10000],c[10000];
   while (scanf("%d",&n)!=EOF)
      for (i=0; i< n ; i++)
          for (j=0; j< n; j++)
             if (j == i)
                              = 0;
                         a[i][j]
                        a[i][j]
                              = a[j][i]
                                       = 1;
             else
```

```
cnt = 0;
       mmin = 2000000000;
       for (i=0; i< n ; i++)
           find_path_u(n,
                           a, i, cnt, b);
           if (cnt < mmin) {
              mmin = cnt;
              for (j=0; j<mmin
                                 ;j++)
                  c[j] = b[j];
              break ;
           }
       printf("%d\n",
                       mmin-1);
       for (i=0; i<mmin-2; i++)
           printf("%d
                       ",c[i]);
       printf("%d",c[i]);
       printf("\n");
}
10. 八数码
                           八数码 Eight(Special
    Function
            Name :
                                                 Judge)
    Description
                       搜索 + 状态 hash
    PKU(1077)
                     HDOJ(1043)
                                        ZOJ(1217)
                广搜
                                          HDOJ(TLE)
                           PKU(312ms)
    BFS
                                                            ZOJ(TLE)
    BFS2
               双向广搜
                           PKU(31ms)
                                          HDOJ(1325ms)
                                                            ZOJ(TLE)
    以上均是每次计算的运行耗时,
                                 ZOJ 的可以全部计算后保存状态
#include
        <cstdio>
         <string>
#include
         <memory>
#include
#include <queue>
using namespace std;
char input[100];
int state[10],
                       e10[10],
                                 fac_n[10];
            s_num,
                                      hash_T2[400000];
char hash_T[400000],
                        step[10000],
struct inf
{
  int pos;
  char mode;
};
queue< int > SQ;
queue<inf>
             SQ2;
```

```
int num_pos( int num, int x, int y)
{
   int temp=(x-1)*3+y;
   if (temp == num\%10) return 9;
   if (temp > num%10)
                       return (num / e10[9-temp+1])
                                                             %10;
   else
          return (num / e10[9-temp]
                                       )%10;
}
int state_pos( int num, int x, int y)
   int temp=(x-1)*3+y;
   if (temp == state[9])
                        return 9;
   if (temp > state[9])
                          return state[temp-1];
   else
          return state[temp];
}
inline int move( int num, char op)
{
   int t0,t1,t2;
   switch (op)
   {
   case 'r':
     if (num%10%3 == 0)
                             return 0;
      return num+1;
   case 'l':
      if ((num-1)%10%3
                          ==0)
                                   return 0;
      return num-1;
   case 'u':
      if (num%10 -3 <= 0)
                               return 0;
      t0 = 9-num\%10
                        + 1;
      t1 = num / e10[t0];
      t2 = t1\%1000;
      t1= t1- t2 + (t2 % 100) * 10 + t2 / 100;
      t1*= e10[t0];
      return (t1 + ((num \% e10[t0]) - 3));
   case 'd':
      if (num\%10 +3 > 9)
                              return 0;
      t0 = 9-num\%10 + 1 -3;
      t1 = num / e10[t0];
      t2 = t1\%1000;
      t1= t1- t2 + (t2 % 10) * 100 + t2 / 10;
      t1*= e10[t0];
      return (t1 + ((num \% e10[t0]) + 3));
   }
```

```
}
bool be_solved()
   int i,j,anti=0;
   for (i=1;i<=8;i++)
      for (j=1;j< i;j++)
                       < state[j] )
          if ( state[i]
             anti++;
    if (anti%2)
                    return false;
                 return true ;
    else
inline int hash( int num)
{
   int dig[10],i=9,j,sum,anti;
   if (num==0)
                    return -1;
                    dig[i]=num%10
                                     , num/=10 , i-- ;
   while (num)
   sum=(9-dig[9])*fac_n[8];
   for (i=1;i<9;i++)
      for (anti=0,j=1;j<i;j++)
          if (dig[i]
                   < dig[j])
             anti++;
        sum += anti*fac_n[i-1];
   return sum;
void BFS()
   int k,to_num,to_hash,i;
   memset(hash_T
                      ,0, sizeof (hash_T));
   while (!SQ.empty())
                              SQ.pop();
   SQ.push(123456789);
   hash_T[ hash(123456789)
                                    ]='e';
   while (!SQ.empty())
      k=SQ.front();
       SQ.pop();
      to_num=move(k,'r');
                                   to_hash=hash(to_num);
                         && hash_T[ to_hash ]==0)
      if (to_hash>=0
                              ]='r' , SQ.push(to_num);
          hash_T[ to_hash
      to_num=move(k,'l');
                                   to_hash=hash(to_num);
```

```
if (to_hash>=0
                        && hash_T[ to_hash
                            ]='I' , SQ.push(to_num);
         hash_T[ to_hash
      to_num=move(k,'u');
                                  to_hash=hash(to_num);
      if (to_hash>=0
                        && hash_T[ to_hash ]==0)
         hash_T[ to_hash
                             ]='u' , SQ.push(to_num);
      to_num=move(k,'d');
                                   to_hash=hash(to_num);
      if (to_hash>=0
                        && hash_T[ to_hash ]==0)
         hash_T[ to_hash
                             ]='d' , SQ.push(to_num);
}
void BFS2()
   int to_num,to_hash,i;
   char *phash,*phash2;
   char op;
   inf k,t;
                     ,0, sizeof (hash_T));
   memset(hash_T
                          sizeof (hash_T2));
   memset(hash_T2,0,
                               SQ2.pop();
   while (!SQ2.empty())
   k.pos=s_num;
                        k.mode=1;
   SQ2.push(k);
   k.pos=123456789;
                             k.mode=2;
   SQ2.push(k);
             hash(s_num)
   hash_T[
                             ]='s';
   hash_T2[
             hash(123456789)
                                    ]='e';
   while (!SQ2.empty())
      k=SQ2.front();
      SQ2.pop();
      to_hash=hash(k.pos);
      if (k.mode==1)
         if (hash_T2[
                       to_hash
                                 ]!=0)
                                           break ;
                                  ,phash2=hash_T2;
         else
                 phash=hash_T
            if (k.mode==2)
                            to_hash
                if (hash_T[
                                     ]!=0)
                                                break;
                        phash=hash_T2,phash2=hash_T;
                else
            t=k;
            t.pos=move(k.pos,'r');
                                          to_hash=hash(t.pos);
            if (to_hash>=0
                             && phash[ to_hash ]==0)
                                 ]='r' , SQ2.push(t);
                phash[ to_hash
            t.pos=move(k.pos,'l');
                                          to_hash=hash(t.pos);
                             && phash[ to_hash ]==0)
            if (to_hash>=0
                                  ]='I' , SQ2.push(t);
                phash[ to_hash
            t.pos=move(k.pos,'u');
                                          to_hash=hash(t.pos);
```

```
if (to_hash>=0
                               && phash[ to_hash
                                   ]='u' , SQ2.push(t);
                 phash[ to_hash
            t.pos=move(k.pos,'d');
                                             to_hash=hash(t.pos);
            if (to_hash>=0
                               && phash[
                                            to_hash ]==0)
                 phash[ to_hash
                                   ]='d'
                                          , SQ2.push(t);
   }
   i=0;
             = hash(k.pos);
   to_hash
   to_num
             = k.pos;
                               ] != 's' ) {
   while ( hash_T[ to_hash
      switch ( step[i++]=hash_T[
                                        to_hash
                                                  ] ) {
             'r':
                    op='l'; break;
       case
             Ή:
                   op='r'; break ;
       case
                    op='d';
             'u':
                             break;
       case
                    op='u';
             'd':
                             break;
       case
      to_num=move(to_num,op);
      to_hash=hash(to_num);
   }
   while (i>0)
                    printf("%c",step[--i]);
   to_hash=hash(k.pos);
   to_num=k.pos;
   while ( hash_T2[
                       to_hash ]!='e'
      switch ( hash_T2[
                            to_hash
                                      ] ) {
                    op='l';
             'r':
                           break ;
       case
             Ή:
                   op='r';
                           break ;
       case
       case
             'u':
                    op='d';
                             break ;
                    op='u';
             'd':
       case
                             break;
      printf("%c",op);
       to_num=move(to_num,
                                    op );
      to_hash=hash(to_num);
int main()
{
   int i,j;
   for (e10[0]=1,i=1;i<=9;i++)
      e10[i] =e10[i-1]*10;
   for (fac_n[0]=0,fac_n[1]=1,i=2;i<=9;i++)
      fac_n[i]
                 =fac_n[i-1]*i;
   while ( gets(input)
                         ) {
      for (i=strlen(input)-1,j=8;i>=0;i--)
```

}

```
if (input[i]!='
                      ') {
            if (input[i]=='x')
              state[9]=j+1;
            else state[j--]=input[i]-'0';
        }
     }
     for (s_num=0,i=9,j=1;i>0;i--,j*=10)
        s_num += state[i]*j;
     if ( !be_solved()
        printf("unsolvable\n");
      else {
        BFS2();
        printf("\n");
}
     高斯消元法
11.
/****
                             高斯消元法
     Function
             Name :
    Description
                          求解线性方程组
    void exchange_col(int
                            p1,int p2,int n)
    交换 p1 行和 p2 行的所有数据
    bool gauss(int
                    n)
    求解系数矩阵为
                   n 的线性方程组 , 方程组无解返回
                                                 false , 否则 true
                            牛顿迭代法
   x1 = x0 - f(x0)/f'(x0)
                              ****/
const int num = 100;
                      + 1]; // 系数矩阵 , 从 0 开始
double matrix[num][num
double ans[num];
                             // 结果数组
void exchange_col( int p1, int p2, int n) // 交换 p1 行和 p2 行的所有数据
{
   double t;
   int i;
   for (i = 0 ; i \le n ; i++)
       t = matrix[p1][i], matrix[p1][i]
                                = matrix[p2][i],matrix[p2][i]
                                                                        = t;
}
bool gauss( int n) // 求解系数矩阵为 n 的线性方程组
```

```
{
   int i,j,k;
   int p;
   double r;
   for (i = 0 ; i < n - 1 ; i++)
       p = i;
       for (j = i + 1; j < n; j++) { // 寻找 i 列绝对值最大值位置
           if (abs(matrix[j][i]) > abs(matrix[p][i]))
              p = j;
       }
      if (p != i) exchange_col(i,p,n);
    if (matrix[i][i] == 0) return false;
       for (j = i + 1; j < n; j++) { // 剩余列进行消元
           r = matrix[j][i] / matrix[i][i];
           for (k = i ; k \le n ; k++)
              matrix[j][k] -= r * matrix[i][k];
       }
   }
   for (i = n - 1; i >= 0; i--) { // 获得结果
       ans[i] = matrix[i][n];
       for (j = n - 1; j > i; j--)
           ans[i] -= matrix[i][j] * ans[j];
    if (matrix[i][i] == 0) return false;
       ans[i] /= matrix[i][i];
   return true;
12. 字符串匹配 (KMP 算法)
    Function Name :
                         字符串匹配 (KMP 算法)
    Description :
                       O(N+M)
void get_nextval( const string & s, int * p)
{
   int i = 0, j = -1;
   p[0] = -1;
   while (i < s.size()) {
       if (j == -1 || s[i] == s[j]) {
           ++i,++j;
           if (s[i] != s[j]) p[i] = j;
           else p[i] = p[j];
       }
```

```
else j = p[j];
    }
}
int Index_KMP( const string & s, const string & s1, int pos)
    int i = pos - 1, j = 0;
    int * next = new int [s1.size()];
    get_nextval(s1,next);
    while (i \leq s.size() && j \leq s1.size()) {
        if (j == -1 || s[i] == s1[j]) ++i,++j;
        else j = next[j];
    }
    if(j > s1.size()) return i - s1.size();
             return -1;
    else
}
13. 全排列,全组合
                               全排列,全组合
     Function
              Name :
void createper( int n) // 全排列
    int total,i,j,k,t,*a=
                          new int [n],top;
    total=1;
    for (i=1;i<=n;i++)
                      {
        a[i]=i;
        total*=i;
    for (i=1;i< n;i++)
                           printf("%d
                                          ",a[i]);
    printf("%d\n",a[n]);
    for (i=1;i<total;i++)
                             {
        j=n;
        while (a[j]< a[j-1])
                              j--;
        k=n;
        while (a[j-1]>a[k])
                               k--;
        t=a[j-1];
        a[j-1]=a[k];
        a[k]=t;
        top=(j+n-1)/2;
        for (k=j;k \le top;k++)
                                   {
            t=a[k];
            a[k]=a[n-k+j];
            a[n-k+j]=t;
```

```
}
              for (j=1;j< n;j++)
                                     printf("%d
                                                   ",a[j]);
              printf("%d\n",a[n]);
         }
     }
                        int m, int n) // 全组合
     void createfab(
     {
                               new int [n+2];
         int i,j,lcount,*a=
         for (i=1;i<=n;i++)
                                      a[i]=i;
         a[n+1]=m+1;
         for (j=1;j< n;j++)
                                    printf("%d
                                                  ",a[j]);
          printf("%d\n",a[n]);
          lcount=1;
          while (a[1]<m-n+1)
                                    {
              for (i=n;i>0;i--)
                   if (a[i] < a[i+1]-1)
                       a[i]++;
                       for (j=i;j< n;j++)
                                                a[j+1]=a[j]+1;
                                                                ",a[j]);
                       for (j=1;j< n;j++)
                                                 printf("%d
                       printf("%d\n",a[n]);
                       lcount++;
                       break ;
              }
      Function
                                二维线段树 RMQ
                  Name
                              HDOJ 1823 Luck and Love
      Description
                                    ****/
           <cstdio>
#include
#include
           <string>
           <algorithm>
#include
using namespace
                      std;
          NMAX 500000
#define
          MQ(x,y)
#define
                     ((x)>(y)?(x):(y))
struct node {
     node * pleft, * pright;
     node * ytree;
     int left, right;
     int M;
```

```
}mem[NMAX];
int mem_pos;
node * new_node()
    node * pt = &mem[mem\_pos]
                                    ++];
    memset(pt,0,sizeof(node));
    pt ->M = -1; //maximum
                                or minimum
    return pt;
}
node * create_tree( int x1, int y1, int x2, int y2, bool flag)
{
    node * root = new_node();
    if (flag) { // first dimension
         root \rightarrowleft = x1;
         root -> right = y1;
         root ->ytree = create_tree(x1, y1, x2, y2, false);
         if (x1 != y1) {
             int mid = (x1+y1)/2;
             root ->pleft = create_tree(x1, mid, x2, y2, true );
             root ->pright = create_tree(mid+1, y1, x2, y2, true );
         }
     else { // second dimension
         root -> left = x2;
         root -> right = y2;
         if (x2 != y2) {
             int mid = (x2+y2)/2;
             root ->pleft = create_tree(x1, y1, x2, mid, false);
             root ->pright = create_tree(x1,
                                              y1, mid+1, y2, false );
    return root;
}
void update(node * root, int d1, int d2, int v, bool flag)
{
    int mid = (root -> left + root -> right)/2;
    if(flag) { // first dimension
         update(root ->ytree, d1, d2, v, false);
         if (root ->left < root ->right) {
             if (d1 <= mid) {
                  update(root ->pleft, d1, d2, v, true );
```

```
}
             else {
                  update(root ->pright, d1, d2, v, true );
             }
         }
    else { // second dimension
         if (root ->left == root ->right) {
             root ->M = MQ(root ->M, v);
         }
         else {
             if (d2 \le mid)
                  update(root ->pleft, d1, d2, v, false);
             }
             else {
                  update(root ->pright, d1, d2, v, false );
             root ->M = MQ(root ->pleft ->M, root ->pright ->M);
         }
    }
}
int query(node * root, int x1, int y1, int x2, int y2, bool flag)
{
    int lmq, rmq;
    int mid = (root -> left + root -> right)/2;
    if (flag) { // first dimension
         if (root ->left == x1 \&\& root ->right == y1) {
             return query(root ->ytree, x1, y1, x2, y2, false);
         else {
             if (y1 <= mid) {
                  return query(root ->pleft, x1, y1, x2, y2, true);
             }
             if (x1 > mid) {
                  return query(root ->pright, x1, y1, x2, y2, true );
             }
             Imq = query(root
                                ->pleft, x1, mid, x2, y2, true );
             rmq = query(root ->pright, mid+1, y1, x2, y2, true );
         }
    }
    else { // second dimension
         if (root ->left == x2 \&\& root ->right == y2) {
             return root ->M;
```

```
}
         else {
             if (y2 <= mid) {
                  return query(root ->pleft, x1, y1, x2, y2, false);
             }
             if (x2 > mid) {
                  return query(root ->pright, x1, y1, x2, y2, false);
             Imq = query(root
                                 ->pleft,
                                          x1, y1, x2, mid, false);
             rmq = query(root
                               ->pright, x1, y1, mid+1, y2, false );
         }
    return MQ(Imq, rmq);
}
int main()
{
    int m;
    char cmd;
    while (scanf( "%d" , &m), m) {
         mem_pos = 0;
         node * root = create_tree(100,200,0,1000,
                                                        true );
         while (m --) {
             getchar();
             cmd = getchar();
             if (cmd == 'I') {
                  int h, ia, il;
                  double a,l;
                  scanf( "%d %lf %lf" , &h, &a, &l);
                  ia = 10*(a+0.05);
                  il = 10*(l+0.05);
                  update(root, h, ia, il, true);
             }
             else {
                  int h1, h2, ia1, ia2;
                  double a1, a2;
                  scanf( "%d %d %lf %lf" , &h1, &h2, &a1, &a2);
                  ia1 = 10*(a1+0.05);
                  ia2 = 10*(a2+0.05);
                  if (h1 > h2) {
                      swap(h1, h2);
                  }
                  if (ia1 > ia2) {
                      swap(ia1, ia2);
```

```
int t = query(root,
                                      h1, h2, ia1, ia2, true );
                  if (t == -1)
                      puts( "-1" );
                  }
                  else {
                      printf( "%.1lf\n"
                                       , t / 10.0);
         }
     15.
          稳定婚姻匹配
     Function Name:
                           稳定婚姻匹配 gale_shapley
                                                      算法
                          HDOJ 1522
                                      Marriage is Stable
     Description
                               ****/
//rmw[i][j]
           代表 i男对女生的喜欢排名
//lwm[i][j]
           代表 i 女对 j 男的喜欢程度
const int MAX = 510;
int w,m,n;
int rmw[MAX][MAX];
int Imw[MAX][MAX],
                       lwm[MAX][MAX];
int couple[MAX];
char sman[MAX][110],
                         swoman[MAX][110];
queue< int > SQ;
void gale_shapley()
    int i,man,woman;
    while (!SQ.empty())
         SQ.pop();
                         sizeof (couple));
    memset(couple,-1,
    for (i=1;i<=n;i++)
                         {
         SQ.push(i);
    }
    while (!SQ.empty())
         man = SQ.front();
         for (i=1;i<=n;i++)
             if (rmw[man][i]
                              != -1) {
                  // 选择为被拒绝且最喜欢的女生
                  woman = rmw[man][i];
                  rmw[man][i]
                                 = -1;
```

```
int pre = couple[woman];
                  if (pre == -1) {
                      couple[woman]
                                        = man;
                      SQ.pop();
                      break ;
                  else {
                      if (lwm[woman][man]
                                              > lwm[woman][pre])
                           SQ.pop();
                           SQ.push(pre);
                           couple[woman]
                                            = man;
                           break ;
                      }
         }
    } //while
}
     16. 后缀数组
                           后缀数组 O(NLogN)
     Function
               Name :
     Description
                          PKU 2774 Long Long Message
          <cstdio>
#include
#include
          <string>
     namespace
using
                   std;
const int MAX = 250000;
char txt[MAX];
int mem[3][MAX],
                    c[MAX],
                             height[MAX];
int * SA, * nSA, * Rank, * nRank;
int len, I1, I2;
//O(NlogN)
//SA[ rank ] = who;
//Suffix(SA[i]) < Suffix(SA[i+1]) , 1 i<n
//Rank[ who ] = rank;
//k-Rank[i] 代表加上满足 Suffix(j) <k Suffix(i) 的j的个数
void init()
{
    11 = strlen(txt);
    txt[l1] = 1; // 特殊结尾
    gets(txt + I1+1);
    12 = strlen(txt + 11+1);
    len = 11 + 12+1;
```

```
txt[len ++] = 1; // 特殊结尾
}
// 性质 .1 对 k n , Suffix(i) <k Suffix(j) 等价于 Suffix(i) < Suffix(j)
// 性质 .2 Suffix(i) =2k Suffix(j) 等价于
//Suffix(i)
           =k Suffix(j)
                         且Suffix(i+k)
                                       =k Suffix(j+k)
// 性质 .3 Suffix(i) <2k Suffix(j)
                               等价于
                        或 (Suffix(i) = k Suffix(j) 且Suffix(i+k) <k Suffix(j+k))
           <k Suffix(j)
//Suffix(i)
void suffix_array()
{
    int i, j, k;
    SA = mem[0]; nSA = mem[1]; Rank = mem[2];
    memset(c, 0, sizeof (c));
    for (i=0;i<len;i++) {
        c[ txt[i] ] ++;
    }
    for (i=0;i<128;i++) {
        c[i+1] += c[i];
    }
    for (i=0;i<len;i++) {
        SA[ -- c[txt[i]] ] = i;
    }
    Rank[ SA[0] ] = 0;
    for (i=1;i<len;i++) {
         Rank[SA[i]] = Rank[SA[i-1]];
         if (txt[ SA[i] ] != txt[ SA[i-1] ]) {
             Rank[ SA[i] ] ++;
         }
    for (k=1;k<len && Rank[SA[len-1]]<len-1;k*=2)
                     0, sizeof (c));
         memset(c,
         for (i=0;i<len;i++) {
             c[ Rank[SA[i]] ] ++;
         }
         for (i=1;i<len;i++) {
             c[i] += c[i-1];
         }
         for (i=len-1;i>=0;i--) {
             if (SA[i] >= k) {
                  nSA[ -- c[ Rank[SA[i]-k] ] ] = SA[i] - k;
         for (i=len-k;i<len;i++) {
             nSA[ -- c[ Rank[i] ] ] = i;
         }
```

```
nRank = SA;
        nRank[ nSA[0] ] = 0;
        for (i=1;i<len;i++) {
            nRank[ nSA[i] ] = nRank[ nSA[i-1] ];
                         != Rank[nSA[i-1]] || Rank[nSA[i]+k] !=
            if (Rank[nSA[i]]
Rank[nSA[i-1]+k])  {
                nRank[nSA[i]]
                            ++;
            }
        SA = nSA;
        nSA = Rank;
        Rank = nRank;
    }
}
//LCP(i,j)=lcp(Suffix(SA[i]),Suffix(SA[j])
//height[i]=LCP(i,i+1) , i<n
int getlcp()
{
    int i, j, k, rs;
    for (i = 0, k = 0; i < len; i++)
        if (Rank[i] == len - 1) {
            height[Rank[i]] = k = 0;
        }
        else {
            if (k > 0) {
                k --;
            j = SA[Rank[i] +1];
            while (txt[i + k] == txt[j + k]) {
                k ++;
            }
            height[Rank[i]] = k;
        }
    }
    for (i = 0, rs = 0; i < len - 1; i++)
        rs = height[i];
        }
    int t = min(11,12);
    return min(t, rs);
}
int main()
```

```
{
    gets(txt);
    init();
    suffix_array();
    printf( "%d\n"
                 , getlcp());
    return 0;
}
     17. 左偏树
                         左偏树
     Function
             Name :
     Description
                        HDOJ 1512 Monkey King
                          二叉堆的变形,方便堆的合并
                             ****/
         <cstdio>
#include
#include
         <string>
         <queue>
#include
         <algorithm>
#include
using namespace
                  std;
const int MAX = 101000;
struct node {
    int v, dis; // 键值,距离
    node * pl, * pr; // 左右子树
    node * pf; // 父节点
}mem[MAX];
int mem_pos;
int value[MAX],
              n;
node * new_node()
    node * pt = mem + (mem_pos
                                     ++);
    memset(pt,
                0, sizeof (node));
    return pt;
}
// 清除节点休息
inline void clear(node * pos) {
    if (pos == NULL) return ;
    pos->pl = pos->pr = pos->pf = NULL;
    pos->dis = 0;
}
// 合并堆 O(log N)
node * merge(node * pa, node * pb) {
    if (pa == NULL) return pb;
    if (pb == NULL) return pa;
    //maximum vertex heap
```

```
if (pb->v > pa->v) std::swap(pa,
                                       pb);
    pa->pr = merge(pa->pr,
                               pb);
    if (pa->pr) {
        if (pa->pl == NULL || pa->pr->dis > pa->pl->dis) {
             std::swap(pa->pl, pa->pr);
        }
    }
    if (pa-pr == NULL) pa-pdis = 0;
    else pa->dis = pa->pr->dis
                                  +1;
    if(pa->pl) pa->pl->pf = pa;
    if(pa->pr) pa->pr->pf = pa;
    return pa;
}
// 插入节点
inline node * insert(node * root, node * val) {
    return merge(root, val);
}
// 删除最大顶
inline node * delete_max(node * root) {
    node * pt = root;
    root = merge(root ->pl, root ->pr);
    if(root) root->pf = NULL;
    clear(pt);
    return root;
}
// 取得最大值
inline int get_max(node
                       * root) {
    return root ->v;
// 构建左偏树 O(N)
inline node * make_leftist_tree()
    queue<node
                  *> SQ;
    node * ptemp;
    int i;
    ptemp = new_node();
    for (i=0;i<n;i++)
         ptemp->v = value[i];
         SQ.push(ptemp);
    while (!SQ.empty())
         int I = SQ.size();
         if (I == 1) return SQ.front();
         while (I --) {
             node * pa = SQ.front();
```

```
SQ.pop();
             node * pb = SQ.front();
             SQ.pop();
             SQ.push(merge(pa,
                                   pb));
         }
}
// 删除已知任意点 O(log N)
inline void delete_any(node
                            * pos) {
    node * ppre = pos->pf;
    node * pnew = delete_max(pos);
    if(pnew) pnew->pf = ppre;
    if (ppre) {
         if (ppre->pl == pos) ppre->pl = pnew;
         else ppre->pr = pnew;
    }
    while (ppre) {
         int vl = -1, vr = -1;
         if (ppre->pl) vl = ppre->pl->dis;
         if (ppre->pr) vr = ppre->pr->dis;
         if (vl < vr) std::swap(ppre->pl,
                                           ppre->pr);
         if (vr +1 == ppre->dis)
                                return ;
         ppre->dis = vr +1;
         pnew = ppre;
         ppre = ppre->pf;
}
node ltree[MAX];
int main() {
    int i,j;
    int m,t;
    while (scanf( "%d" , &n)==1)
         for (i=0;i< n;i++) {
             scanf( "%d" , &t);
             ltree[i].v
                       = t;
             Itree[i].dis = 0;
             ltree[i].pl = ltree[i].pr = ltree[i].pf
                                                    = NULL;
         scanf( "%d" , &m);
         int a,b;
         while (m --) {
             scanf( "%d %d" , &a,&b);
             a --; b --;
```

```
node * pa, * pb;
              pa = ltree +a;
              pb = ltree +b;
              while (pa->pf)
                              pa = pa->pf;
              while (pb->pf)
                              pb = pb - pf;
                                puts( "-1" );
              if (pa == pb)
              else {
                   node * p1 = delete_max(pa);
                  node * p2 = delete_max(pb);
                  pa->v /= 2;
                   pb->v /= 2;
                   p1 = insert(p1,
                                     pa);
                   p1 = insert(p1,
                                     pb);
                   p1 = merge(p1,
                                     p2);
                  printf( "%d\n" , get_max(p1));
              }
         }
}
     18. 标准 RMQ-ST
      Function Name:
                            标准 RMQ-ST
                           PKU 3264 Balanced Lineup
      Description
                                ****/
#include
          <cstdio>
          <string>
#include
          <algorithm>
#include
using namespace
                    std;
const int MAX = 51000;
const int LOGMAX = 16;
int n,q;
int st_max[LOGMAX][MAX],
                                st_min[LOGMAX][MAX];
void make_st()
{
    int i,j,k;
    for (j=1; (1<< j) <= n ; j++) {
         k = 1 << (j-1);
         for (i=0; i+k < n; i++)
              st_max[j][i]
                            = \max(st_max[j-1][i],
                                                      st_max[j-1][i+k]);
              st_min[j][i]
                            = min(st_min[j-1][i],
                                                     st_min[j-1][i+k]);
         }
```

```
}
int rmq( int a, int b, int flag)
{
    int dis = abs(b-a) +1;
    int k;
                       <= dis ;k++) ;
    for (k=0; (1<< k))
    k --;
    if (flag > 0) {
         return \max(st_{\max}[k][a], st_{\max}[k][b-(1<< k)+1]);
    }
    else {
         return min(st_min[k][a],
                               st_min[k][b-(1<<k)+1]);
    }
}
int main()
{
    while (scanf( "%d %d", &n,&q)==2)
         int i;
         for (i=0;i< n;i++)
             scanf( "%d" , &st_max[0][i]);
             st_min[0][i]
                         = st_max[0][i];
         make_st();
         for (i=0;i<q;i++)
             int a,b;
             scanf( "%d %d" , &a,&b);
             printf( "%d\n", rmq(a-1,b-1,1) - rmq(a-1,b-1,-1));
         }
}
     19. 度限制最小生成树
                           度限制最小生成树
     Function Name:
     Description
                          PKU 1639 Picnic Planning
                           有一个顶点有度限制,如果所有点都有限制,当限制
                                                                             >4 时是 NP
#include
          <cstdio>
#include
          <string>
#include
          <queue>
#include
          <vector>
```

```
#include <map>
#include <algorithm>
using namespace std;
const int MAX = 50;
int t,n,m;
          , int > names;
map<string
int path[MAX][MAX];
//dmax[i]:
                    ,不与 park 相连的边的最大权值
          vi->park
int dmax[MAX];
struct node {
    int s,t;
    int dis;
    bool operator < (const node & tt) const {
        return dis > tt.dis;
    }
};
bool vis[MAX];
//block[i]: vi 所属连通分量
//bs: 连通分量数目
//v0min[i][2]: park 与第 i个连通分量的最小权值 [0] ,连接顶点 [1]
int block[MAX], v0min[MAX][2],
                                  bs;
       度限制生成树
//mst:
vector< int > mst[MAX];
queue< int > sq;
// 最小花费 , park 下标 , 限制度数
int cost, park, deg;
//O(NlogN)
           prime 求所有连通分量 mst
void prime_all_mst() {
    int i,j;
    priority_queue<node>
                            pq;
    node now, next;
               0, sizeof (vis));
    memset(vis,
    for (i=0;i<=n;i++)
                        mst[i].clear();
    vis[park] = true;
    block[park] = 1; bs = 1; //park 为第个连通分量
    cost = 0;
    for (i=1;i<=n;i++)
        if (!vis[i]) {
             bs ++;
             while (!pq.empty())
                                 pq.pop();
             now.s = i; now.t = i; now.dis = 0;
             pq.push(now);
             while (!pq.empty())
                 now = pq.top();
```

```
pq.pop();
                  if (vis[now.t])
                                  continue ;
                  vis[now.t]
                              = true ;
                  mst[now.s].push_back(now.t);
                  mst[now.t].push_back(now.s);
                   block[now.t]
                                 = bs;
                   cost += now.dis;
                  next.s = now.t;
                  for (j=1;j<=n;j++)
                       if (!vis[j]
                                 && path[next.s][j]
                                                      != -1) {
                            next.t = j;
                            next.dis = path[next.s][j];
                            pq.push(next);
                       }
             }
         }
}
//O(N)
        park 连接各连通分量
bool connect_block()
    int i,j,k;
    // 选取连接相邻连通分量的最小边
    for (i=2;i<=bs;i++)
                          v0min[i][0]
                                         = INT_MAX;
    while (!sq.empty())
                          sq.pop();
    for (i=1;i<=n;i++)
                           != -1 && v0min[block[i]][0]
         if (path[park][i]
                                                           > path[park][i])
             v0min[block[i]][0]
                                   = path[park][i];
             v0min[block[i]][1]
                                   = i;
    k = 0;
    for (i=2;i<=bs;i++) {
         if (v0min[i][0] != INT_MAX) {
             cost += v0min[i][0];
             path[park][ v0min[i][1] ] = -1;
              dmax[ v0min[i][1]
                                ] = INT_MIN;
             sq.push(v0min[i][1]);
                                 // 用来初始化 dmax
             k ++; // 能连通的分量数
         }
    }
    // 图连通,且限制度数大于等于连通分量数
    deg -= bs-1;
    return k \ge bs-1 \&\& deg \ge 0;
```

```
}
//O(N) 计算 dmax
void cal_dmax()
     int i;
     memset(vis,
                    0, sizeof (vis));
     while (!sq.empty())
          int now = sq.front();
          sq.pop();
          vis[now]
                    = true ;
          for (i=0;i<mst[now].size();i++)
               int next = mst[now][i];
               if (!vis[next])
                    dmax[next]
                                   = max(dmax[now],
                                                           path[now][next]);
                    sq.push(next);
                    vis[next] = true ;
               }
          }
}
//O(N)
         差额最小删除操作
void del_path( int pos, int val) {
     int i;
     queue< int > sq2;
     memset(vis,
                    0, sizeof (vis));
     sq2.push(pos);
     vis[pos] = true ;
     while (!sq2.empty())
          int now = sq2.front();
          sq2.pop();
          for (i=0;i<mst[now].size();i++)</pre>
                                                 {
               int next = mst[now][i];
               if (!vis[next])
                    if (val == path[now][next])
                         mst[now].erase(mst[now].begin()
                                                                    +i);
                         return ;
                    }
                    sq2.push(next);
                    vis[next]
                               = true ;
          }
}
//O(deg*N)
bool deg_limit_mst()
```

```
int i,j,v;
    int minv, minp;
    cal_dmax();
    for (i=0;i<deg;i++) {
         minv = INT_MAX;
                             minp = -1;
         for (j=1;j<=n;j++)
              if (path[park][j] != -1) { // 差额最小选择操作
                   if (minv > path[park][j]

    dmax[j])

                       minv = path[park][j]

    dmax[j];

                       minp = j;
             }
         v = cost + minv;
         if (minp == -1 \mid \mid v >= cost) return false;
         cost = v;
                         = -1; // 差额最小添加删除操作
         path[park][minp]
         del_path(minp,
                           dmax[minp]);
         mst[park].push_back(minp);
         while (!sq.empty())
                              sq.pop();
         sq.push(minp);
         dmax[minp]
                       = INT_MIN;
         cal_dmax();
    for (i=0;i<mst[park].size();i++)
                                         mst[ mst[park][i]
                                                             ].push_back(park);
    return true;
}
int main() {
    int i,j;
    char n1[20],
                   n2[20];
    names.clear();
    scanf( "%d" , &m);
    memset(path, -1, sizeof (path));
    n = 1;
    for (i=0;i< m;i++)
                     {
         int x,y,z;
         scanf( "%s %s %d" , n1,n2,&z);
         x = names[string(n1)];
         y = names[string(n2)];
         if (x == 0) names[string(n1)] = x = n ++;
         if (y == 0) names[string(n2)] = y = n ++;
         if (strcmp(n1, "Park") == 0) park = x;
         else if (strcmp(n2, "Park") == 0) park = y;
```

```
= path[y][x]
         path[x][y]
                                    = z;
    }
    n --;
    scanf( "%d" , &deg);
    prime_all_mst();
    connect_block();
    deg_limit_mst();
    printf( "Total miles driven: %d\n", cost);
}
     20. 最优比率生成树
                           最优比率生成树 (迭代法)
     Function Name:
     Description
                          PKU 2728 Desert King
                                ****/
#include
          <cstdio>
          <string>
#include
          <cmath>
#include
          <algorithm>
#include
using namespace
                   std;
const int MAX = 1100;
int n;
struct point {
    int x,y,z;
}vi[MAX];
struct node {
    int s, t;
    double dis;
    bool operator
                    < (const node & tt) const {
         return dis > tt.dis;
};
double dist[MAX][MAX];
bool vis[MAX];
double rate;
double prime() {
    double cost = 0;
    double len = 0;
    double \quad d[MAX], v;
    int pre[MAX];
    int i,j;
    memset(vis,
                  0, sizeof (vis));
```

```
vis[0] = true ;
    for (i=1;i<n;i++) {
         d[i] = abs(vi[0].z-vi[i].z) - rate*dist[0][i];
         pre[i] = 0;
     }
    for (i=1;i< n;i++) {
         double minv = INT_MAX;
         int minp = -1;
         for (j=1;j< n;j++) {
              if (!vis[j] &\& minv > d[j]) {
                   minv = d[j];
                   minp = j;
              }
         }
         vis[minp] = true ;
         cost += abs(vi[pre[minp]].z - vi[minp].z);
         len += dist[pre[minp]][minp];
         for (j=1;j< n;j++) {
              if (!vis[j] \&\& d[j] > (v=abs(vi[minp].z-vi[j].z) - rate*dist[minp][j]))
                                                                                       {
                   d[j] = v;
                   pre[j] = minp;
              }
         }
    return cost / len;
}
int main() {
    int i,j;
    while (scanf( "%d" , &n), n) {
         for (i=0;i<n;i++) {
              scanf( "%d %d %d" , &vi[i].x, &vi[i].y, &vi[i].z);
         }
         for (i=0;i<n;i++) {
              dist[i][i] = 0;
              for (j=i+1;j<n;j++)
                   dist[i][j]
                             = dist[j][i]
                                           = sqrt(1.0*(vi[i].x-vi[j].x)*(vi[i].x-vi[j].x)
(vi[i].y-vi[j].y)*(vi[i].y-vi[j].y));
         }
         rate = 0;
         while (true ) {
              double pre = rate;
              rate = prime();
```

```
if (fabs(rate - pre) < 0.001) break ;
}
printf( "%.3lf\n" , rate);
}</pre>
```

21. 最小花费置换

```
//Cow Sorting
// 对一个轮换进行处理的时候,
                             应该考虑在轮换内进行交换, 或与轮换外的元素交换之后,
                                                                                   使代
价值更小
#include
         <cstdio>
         <string>
#include
#include
         <functional>
         <algorithm>
#include
using namespace
                   std;
int g[10100],n;
bool vis[10100];
int pos[101000];
struct node {
    int v,p;
    bool operator < (const node & t) const {
         return v < t.v;
    }
}g2[10100];
int main()
    int i,j;
    int sum, mmin;
    while (scanf( "%d" , &n)==1)
         sum = 0;
         mmin = INT_MAX;
         for (i=1;i<=n;i++)
             scanf( "%d" , g+i);
             sum += g[i];
             g2[i].v = g[i];
             g2[i].p = i;
             mmin = min(mmin,
                                  g[i]);
             vis[i] = false;
         }
         sort(g2+1,g2+n+1);
         for (i=1;i<=n;i++) {
             pos[ g2[i].v ] = g2[i].p;
         }
         for (i=1;i<=n;i++) {
             if (!vis[i]) {
```

```
int tpos = i;
                  int len = 0;
                  int tmin = INT\_MAX;
                  do {
                      tmin = min(tmin,
                                          g[tpos]);
                      vis[tpos]
                                 = true ;
                      tpos = pos[g2[tpos].v
                                                ];
                      len ++;
                  } while (tpos != i);
                  // 选择两种方案中的最优方案
                  sum += min( (len-2)*tmin,
                                                 (len+1)*mmin
                                                                  +tmin);
         }
         printf( "%d\n" , sum);
    }
}
    22. 区间 K 大数
//POJ 2104
#include
          <cstdio>
#include
          <string>
#include
          <vector>
         <algorithm>
#include
using namespace
                   std;
const int NMAX = 100000;
const int LOGNMAX = 17 +1;
int sortseq[LOGNMAX][NMAX];
int num[NMAX];
struct node {
    int I,r,d;
    node * pl,*
                 pr;
}mem[(NMAX<<1)+100];</pre>
int mempos,n,m;
node * root;
node * make_tree( int I, int r, int d) {
    node * rt = mem+(mempos
    rt->I = I; rt->r = r; rt->d = d;
    if (I == r) {
         sortseq[d][l] = num[l];
         return rt;
    }
    int mid = (l+r) \gg 1;
    rt->pl = make_tree(l,mid,d+1);
    rt->pr = make_tree(mid+1,r,d+1);
```

```
int i=l,j=mid+1,k=l;
     while (i \le mid \&\& j \le r) {
           \text{if } (\mathsf{sortseq}[\mathsf{d+1}][\mathsf{i}] \qquad < \mathsf{sortseq}[\mathsf{d+1}][\mathsf{j}]) \qquad \mathsf{sortseq}[\mathsf{d}][\mathsf{k++}] \qquad = \\ \\ \\
sortseq[d+1][i++];
          else sortseq[d][k++] = sortseq[d+1][j++];
     while (i \le mid) sortseq[d][k++] = sortseq[d+1][i++];
    while (j \le r) sortseq[d][k++] = sortseq[d+1][j++];
     return rt;
}
int s,t,rank;
int query(node * rt, int val) {
    int i,mid,ret;
     if (s \le rt > l & rt > r \le t) {
          if (val <= sortseq[rt->d][rt->l]) return 0;
          else if (sortseq[rt->d][rt->r]
                                        < val) return rt->r - rt->l +1;
          else if (sortseq[rt->d][rt->r] == val) return rt->r - rt->l;
          int I = rt->I, r = rt->r, mid;
          while (l \ll r) {
              mid = (I+r) >> 1;
              if (val \le sortseq[rt->d][mid])  r = mid-1;
              else I = mid+1;
          return | - rt->l;
     else {
          ret = 0;
          mid = (rt->l+rt->r) >> 1;
          if (s <= mid) ret += query(rt->pl,val);
          if (mid+1 <= t) ret += query(rt->pr,val);
          return ret;
}
// 二分查找时遇到相同值的处理非常重要
int main() {
    int i,j,l,r;
     scanf( "%d %d" ,&n,&m);
     for (i=0;i<n;i++) scanf( "%d" ,num+i);
     mempos = 0;
     root = make_tree(0,n-1,0);
     while (m --) {
         s = get_val()-1; t = get_val()-1; rank = get_val()-1;
         l = 0, r = n-1;
          while (I \ll r) {
```

```
int mid = (l+r) \gg 1;
             // 二分查找 sortseq[0][mid]
                                           在区间 [s,t] 中的排名
             int pos = query(root,sortseq[0][mid]);
             if (rank < pos) r = mid-1;
             else I = mid+1;
         printf( "%d\n" ,sortseq[0][r]);
}
     23. LCA - RMQ-ST
//POJ 3417
         O(nlogn)-O(1)
//online
#include
          <cstdio>
          <string>
#include
#include
          <queue>
#include
          <algorithm>
using namespace
                    std;
typedef __int64
                  bigint;
const int MAX = 100010;
const int STMAX = 200010;
const int LOGMAX = 18;
int n,m;
const int ENDFLAG = 0;
struct EDGELIST {
    int start[MAX];
    int last[MAX];
                              //pos,listnext
    int edge[MAX<<1][2];</pre>
    int tot;
    void clear() {
         tot = ENDFLAG +1;
         memset(last,ENDFLAG,
                                   sizeof (last));
         memset(start,ENDFLAG,
                                    sizeof (start));
    }
    void push_back(
                       int s, int t) {
         edge[tot][0]
                        = t;
         edge[tot][1]
                        = ENDFLAG;
         if (last[s]
                    != ENDFLAG) {
             edge[ last[s] ][1] = tot;
         }
         else {
             start[s]
                       = tot;
```

```
}
          last[s] = tot;
          tot ++;
          //swap
          if (s == t) return ;
          edge[tot][0]
                         = s;
          edge[tot][1]
                         = ENDFLAG;
          if (last[t] != ENDFLAG) {
              edge[ last[t] ][1] = tot;
          }
          else {
              start[t] = tot;
          last[t] = tot;
          tot ++;
}tree;
int cov[MAX];
bool vis[MAX];
bigint cnt[2];
int root[MAX],son[MAX];
int stn;
int st_min[LOGMAX][STMAX];
int order[STMAX],first[MAX],deep[STMAX];
void make_st()
     int i,j,k;
     for (i=0;i<stn;i++)
                       st_min[0][i]
                                            = i;
     for (j=1; 1 << j <= stn ; j++) {
          k = 1 << (j-1);
          for (i=0; i+k < stn ; i++) {
                                           < deep[st_min[j-1][i+k]])
              if (deep[st_min[j-1][i]]
                    st_min[j][i]
                                  = st_min[j-1][i];
              else
                   st_min[j][i]
                                  = st_min[j-1][i+k];
          }
     }
}
int rmq( int a, int b) {
     int dis = abs(b-a)
                          +1;
     int k;
    for (k=0; (1<< k) <= dis ;k++);
```

```
k --;
     int ret = st_min[k][a];
     if (deep[ret]
                   > deep[st_min[k][b-(1<<k)+1]])
         ret = st_min[k][b-(1<< k)+1];
     return ret;
}
int lca( int a, int b) {
    int x = first[a],y = first[b];
    if (x > y) swap(x,y);
    return order[rmq(x,y)];
}
int ordcnt;
int sq[MAX];
int qf,qe;
void dfs( int pos, int d) {
    int i,j;
    vis[pos] = true;
    first[pos]
                = ordcnt;
     deep[ordcnt]
                    = d;
     order[ordcnt
                    ++] = pos;
    for (i=tree.start[pos];
                              i != ENDFLAG ;i=tree.edge[i][1])
         int next = tree.edge[i][0];
         if (vis[next])
                         continue ;
         son[pos]
                    ++;
         root[next]
                      = pos;
         dfs(next,d+1);
         deep[ordcnt]
                         = d;
         order[ordcnt
                         ++] = pos;
    if (son[pos] == 0) sq[qe ++] = pos;
}
void treedp() {
    son[0] = -1;
    while (qf < qe) {
         int now = sq[qf ++];
         if (cov[now] \ll 1) cnt[cov[now] ++;
         son[ root[now] ] --;
         cov[ root[now] ] += cov[now];
         if (son[ root[now] ] == 0)  sq[qe ++] = root[now];
    }
}
```

```
int get_val() {
    int ret = 0;
    char ch;
    while ((ch=getchar()) > '9' || ch < '0');
    do {
         ret = ret*10 + ch - '0';
    } while ((ch=getchar()) <= '9' && ch >= '0');
    return ret;
}
int main() {
    int i,j,a,b,rt;
    n = get_val();
    m = get_val();
    if (n == 1) {
         puts( "0" );
         return 0;
    }
    tree.tot = ENDFLAG +1;
    qf = qe = 0;
    for (i=0;i< n-1;i++) {
         a = get_val();
         b = get_val();
         tree.push_back(a,b);
         rt = a;
    }
    ordcnt = 0;
     dfs(rt,0);
     stn = ordcnt;
    make_st();
    for (i=0;i< m;i++)
         a = get_val();
         b = get_val();
         cov[a] ++;
         cov[b] ++;
         cov[lca(a,b)]
                      -= 2;
    }
    treedp();
    cnt[0] --;
```

```
printf( "%I64d\n" ,cnt[0]*m + cnt[1]);
}
```

24. LCA - Tarjan

```
//POJ 3417
//offline
          O(na(n))
#include
          <cstdio>
          <string>
#include
#include
          <vector>
          <algorithm>
#include
using namespace
                    std;
        __int64
typedef
                  bigint;
const int MAX = 100010;
int n,m;
const int ENDFLAG = 0;
struct EDGELIST {
    int start[MAX];
    int last[MAX];
    int edge[MAX<<1][2];</pre>
                              //pos,listnext
    int tot;
    void clear()
         tot = ENDFLAG +1;
         memset(last,ENDFLAG,
                                    sizeof (last));
         memset(start,ENDFLAG,
                                     sizeof (start));
    void push_back(
                       int s, int t) {
         edge[tot][0]
                        = t;
         edge[tot][1] = ENDFLAG;
         if (last[s] != ENDFLAG) {
             edge[ last[s] ][1] = tot;
         }
         else {
            start[s] = tot;
         last[s] = tot;
         tot ++;
         //swap
         if (s == t) return ;
         edge[tot][0] = s;
         edge[tot][1] = ENDFLAG;
```

```
if (last[t] != ENDFLAG) {
              edge[ last[t] ][1] = tot;
         }
         else {
              start[t] = tot;
         last[t] = tot;
         tot ++;
    }
}tree,newed;
int cov[MAX];
bool vis[MAX];
bigint cnt[2];
int father[MAX];
int ancestor[MAX];
int find_set( int x) {
    if (father[x] == x) return x;
    return father[x] = find_set(father[x]);
}
void union_set( int x, int y) {
    father[ find_set(y) ] = x;
}
void tarjan( int pos, int pre) {
    int i,j;
    father[pos]
                  = pos;
     ancestor[pos]
                     = pos;
    for (i=tree.start[pos]; i != ENDFLAG ;i=tree.edge[i][1])
         int next = tree.edge[i][0];
         if (next == pre) continue ;
         tarjan(next,pos);
         union_set(pos,next);
         cov[pos] += cov[next];
    }
    vis[pos] = true;
    for (j=newed.start[pos]; j != ENDFLAG ;j=newed.edge[j][1])
                                                                           {
         int next = newed.edge[j][0];
                                          find_set(next) ] ] -= 2;
         if (vis[next]) cov[ ancestor[
```

```
if (cov[pos] \ll 1) cnt[cov[pos] ++;
}
int get_val() {
    int ret = 0;
    char ch;
    while ((ch=getchar()) > '9' || ch < '0');
    do {
      ret = ret*10 + ch - '0';
    return ret;
}
int main() {
    int i,j,a,b,rt;
    n = get_val();
    m = get_val();
    if (n == 1) {
        puts( "0");
        return 0;
    }
    tree.tot = newed.tot = ENDFLAG +1;
    for (i=0;i< n-1;i++)
        a = get_val();
        b = get_val();
        tree.push_back(a,b);
        rt = a;
    for (i=0;i< m;i++)
        a = get_val();
         b = get_val();
         newed.push_back(a,b);
         cov[a] ++;
        cov[b] ++;
    tarjan(rt,rt);
    cnt[0] --;
    printf( "%l64d\n"
                      ,cnt[0]*m
                                 + cnt[1]);
}
```

25. 指数型母函数

```
/*
HDOJ 1521 排列组合
有n种物品,并且知道每种物品的数量。要求从中选出 m 件物品的排列数。 <=m,n<=10
数量较少时,直接用除法
*/
#include <cstdio>
#include <string>
#define MAX 100
double cal[2][MAX];
double *pre,*now,*pt;
int n,m;
int a[11];
double fac[100];
int main()
{
    int i,j,k,sum;
    fac[0] = fac[1] = 1;
    for (i=2;i<=20;i++) {
        fac[i] = fac[i-1] * i;
    while (scanf( "%d %d" ,&n,&m)==2)
                                        {
        memset(cal,0, sizeof (cal));
        for (i=0;i< n;i++) {
            scanf( "%d" ,&a[i]);
        }
        pre = cal[0];
        now = cal[1];
        pre[0] = 1;
        for (i=1;i<=a[0];i++) {
            pre[i] = 1.0 / fac[i];
        }
        for (i=1;i< n;i++) {
            for (j=0;j<MAX;j++)
                if (pre[j] > 0) {
                    for (k=0;k<=a[i];k++)
                        now[k+j] += pre[j] / fac[k];
                }
            }
```

```
pt = now;
             now = pre;
             pre = pt;
             memset(now,0,
                              sizeof (cal[0]));
             pre[0] = 1;
         }
         printf( "%.0lf\n"
                         ,fac[m] * pre[m]);
}
     26. 指数型母函数(大数据)
#include <iostream>
using namespace
                   std;
int mm[1000][100];
__int64 a[1000],
                   b[1000];
inline __int64 gcd( __int64 x, __int64 y) // 求公约数
{
    __int64 temp;
    while (x % y) {
        temp = x \% y;
        x = y;
        y = temp;
    return y;
}
int main() {
    int n, m, i, j, k;
    __int64 tmp, tmp1;
    while (scanf( "%d %d" , &n, &m) != EOF) {
         for (i = 0; i < n; i ++) {
             scanf( "%d" , &mm[i][0]);
                                       j ++)
             for (j = 1; j \le mm[i][0];
                 scanf( "%d" , &mm[i][j]);
         }
         memset(a, 0, sizeof (\underline{\phantom{a}}int64 )*(m + 1));
         for (i = 1; i \le mm[0][0];
             a[mm[0][i]]
                         = 1;
         for (i = 1; i < n; i ++) {
             memset(b, 0, sizeof (\underline{\phantom{a}}int64 )*(m + 1));
             for (j = 0; j \le m; j ++)
                 for (k = 1; j + mm[i][k] <= m && k <= mm[i][0]; k ++) {
                      if (mm[i][k] != 0) {
```

```
tmp = 1; tmp1 = 1;
                        int w = j + mm[i][k];
                        int x = mm[i][k] < j ? mm[i][k] : j; //x 是较小的数
                        int y = w - x;
                        __int64 z;
                        while (w > y) {
                            tmp *= w;
                           tmp1 *= x;
                            z = gcd(tmp, tmp1);
                            if (z > 1) {
                               tmp /= z;
                               tmp1 \neq z;
                            }
                            W--;
                            X--;
                        b[j + mm[i][k]] += tmp * a[j];
                }
            for (j = 0; j \le m; j ++)
                a[j] += b[j];
       }
        printf( "%I64d\n" , a[m]);
    return 0;
}
    27. AC 自动机 (字典树 +KMP )
const int NMAX = 10000;
const int LMAX = 1000001;
const int MMAX = 51;
const int MEMMAX = 500000;
char s[LMAX];
char p[MMAX];
int n, m;
struct NODE
    int nsuffix;
    char chword;
    NODE * next, * father;
```

```
NODE * son[26];
}mem[MEMMAX];
int total;
NODE * root;
NODE * new_node()
{
    NODE * ret = &mem[total
    memset(ret, 0, sizeof (NODE));
    return ret;
}
// O(n MMAX)
void insert(NODE * rt, char * p)
{
    //puts(p);
    if (*p == 0)
    {
         rt->nsuffix
                      ++;
         return ;
    }
                  - 'a' ] == NULL)
    if (rt->son[*p
                      - 'a' ] = new_node();
         rt->son[*p
         rt->son[*p
                     - 'a' ]->father = rt;
                      - 'a' ]->chword = *p;
         rt->son[*p
                    - 'a' ], p+1);
    insert(rt->son[*p
}
// O(n MMAX)
void bfs()
    int i, j;
    queue <NODE *> sq;
    sq.push(root);
    while (!sq.empty())
         NODE * now = sq.front();
         sq.pop();
         if (now->father
                           == root)
              now->next
                          = root;
```

```
else
         {
             NODE * shift = (now->father)->next;
                                                                    - 'a' ] == NULL)
              while (shift != root && shift->son[now->chword
                  shift = shift->next;
                           = shift->son[now->chword
                                                          - 'a' ];
              now->next
             if (now->next
                              == NULL)
                  now->next
                               = root;
         }
         for (i=0; i<26; i++)
         {
             if (now->son[i]
                            != NULL)
                  sq.push(now->son[i]);
         }
    }
}
// O(LMAX)
int solve()
{
    int i,j;
    int ret = 0;
    NODE * now = root;
    NODE * psuffix;
    root->father
                   = root;
    bfs();
    for (i=0;
               s[i]; i++)
    {
         while (now != root && now->son[ s[i] - 'a' ] == NULL)
             now = now->next;
         now = now->son[ s[i] - 'a'];
         if (now == NULL)
             now = root;
         // add same suffix
         psuffix = now;
         while (psuffix != root && psuffix->nsuffix != -1)
             ret += psuffix->nsuffix;
             psuffix->nsuffix
                            = -1;
             psuffix = psuffix->next;
         }
```

```
return ret;
}
    28. FFT (大数乘法)
const int BASE = 100000;
const int N_DIGIT = 5;
const int N = 32768;
const double PI = acos(-1.0);
struct Complex
{
    double real, imag;
};
Complex omega[N >> 1];
Complex
         a[N];
Complex b[N];
char s[1000003];
int d[N], len;
void bitReverse(Complex
                           a[])
{
    int i, j = 1, k;
    Complex t;
    for (i = 1; i < len; ++ i)
        if (i < j)
            t.real = a[j - 1].real;
            t.imag = a[j - 1].imag;
             a[j - 1].real = a[i - 1].real;
             a[j - 1].imag = a[i - 1].imag;
             a[i - 1].real = t.real;
             a[i - 1].imag = t.imag;
        }
        k = len \gg 1;
        while (k < j)
        {
            j -= k;
            k >>= 1;
        }
        j += k;
```

```
}
void calOmega()
{
    double unit = 2 * PI / len;
    int n = len \gg 1;
   for (int i = 0; i < n; ++ i)
       double t = unit * i;
       omega[i].real = cos(t);
       omega[i].imag = sin(t);
}
void fft(Complex a[], bool inverse = false)
{
    bitReverse( a);
   int s = len \gg 1;
   int m, k, j;
   int up, t, step;
    int i1, i2;
    Complex tmp;
   if (inverse)
       for (j = 0; j < s; ++ j)
           omega[j].imag = - omega[j].imag;
    s = 1;
    for (m = 2; m \le len; m \le 1)
       !!!!!!!!
       for (k = 0; k < len; k += m)
       {
           step = 0;
           for (j = 0; j < up; ++ j)
               i1 = k + j;
               i2 = i1 + up;
               tmp.real = omega[step].real * a[i2].real - omega[step].imag *
a[i2].imag;
                        = omega[step].real * a[i2].imag + omega[step].imag *
               tmp.imag
```

```
a[i2].real;
                 a[i2].real
                            = a[i1].real - tmp.real;
                 a[i2].imag
                            = a[i1].imag
                                          tmp.imag;
                 a[i1].real
                            += tmp.real;
                 a[i1].imag
                             += tmp.imag;
                 step += t;
           }
        }
         ++ s;
    }
    if (inverse)
        double t = 1.0 / len;
        for (j = 0; j < len; ++ j)
             a[j].real *= t;
    }
}
int convert( int d[], char s[])
{
    int sLen = strlen( s );
    int dLen = ((sLen - 1) / N_DIGIT) + 1, i = 0, n;
    char *pRight = s + sLen - 1, *pLeft = pRight - (N_DIGIT - 1);
    memset(d, 0, sizeof (int) * dLen);
    while (i < dLen && pRight >= s)
        if (pLeft < s) pLeft = s;
        n = 0;
        while (pLeft <= pRight)
         {
             n = n * 10 + (*pLeft & & 15);
             ++ pLeft;
        }
        d[i ++] = n;
        pRight -= N_DIGIT;
        pLeft = pRight - (N_DIGIT - 1);
    return dLen;
}
bool init()
{
    int i, j;
```

```
//read a
    if (scanf( "%s" , s) != 1)
        return false;
    int aLen = convert(d, s); //length of a
    for (i = 0; i < aLen; ++ i)
        a[i].real
                 = d[i];
        a[i].imag = 0;
    }
    //read b
    scanf( "%s" , s);
    int bLen = convert(d, s); //length of b
    for (j = 0; j < bLen; ++ j)
        b[j].real = d[j];
        b[j].imag = 0;
    }
    len = 1;
                                    //length of product who uses int
    while (len < aLen + bLen)
                                  len <<= 1;
             + i, 0, sizeof (Complex) * (len - i));
    memset(a
             + j, 0, sizeof (Complex) * (len - j));
    memset(b
    calOmega();
    return true;
}
void mul()
{
    for (int i = 0; i < len; ++ i)
    {
        double real = a[i].real * b[i].real - a[i].imag * b[i].imag;
        double imag = a[i].real * b[i].imag + a[i].imag * b[i].real;
        a[i].real
                 = real;
        a[i].imag
                  = imag;
}
void print()
{
    double carry = 0, t;
```

```
static char format[10];
    int i;
    for (i = 0; i < len; ++ i)
       t = carry + a[i].real;
        carry = floor((t + 0.5) / BASE);
        d[i] = int (t - carry * BASE + 0.5);
    for (i = len - 1; i > 0 \&\& d[i] == 0; -- i);
    sprintf(format, "%%.%dd" , N_DIGIT);
    printf( "%d" , d[i]);
    for (-- i; i >= 0; -- i)
        printf(format,
                      d[i]);
    printf( "\n" );
}
int main()
{
    while (init())
        fft( a );
        fft( b);
        mul();
        fft(a, true);
        print();
    return 0;
}
    29. 二分图网络最大流最小割
// PKU 2125 Destroying The Graph
// 二分图最小点权覆盖集,求割集
// 1. 设置一个集合 A, 最开始 A={s}, 按如下方法不断扩张 A:
// 2. 若存在一条边 (u,v), 其流量小于容量,且 u属于 A,则 v加入 A
// 3. 若存在 (v,u), 其流量大于 0 , 且 u 属于 A,则 v加入 A
// 4.A 计算完毕,设 B=V-A ,最小割集 E={(u,v) | u A,v
                                                      B}
// Character '+' means that Bob removes all arcs incoming into the specified vertex
and '-' that Bob removes all arcs outgoing from the specified vertex.
bool S[MAX];
void dfs( int pos) {
    int i;
    S[pos] = 1;
```

```
for (i=1;i<=m;i++)
          if (!S[i] && net[pos][i])
                                      dfs(i);
     }
}
struct node {
     int num;
     char sign;
}cs[MAX];
void find_cut() {
     int i,ps = 0;
     memset(S, 0, sizeof (S));
     dfs(1);
     for (i=2;i<=n+1;i++)
          if (!S[i] \&\& net[1][i] == 0) {
               //printf("%d
                            -\n", i-1);
               cs[ps].num
                             = i-1;
                             = '-' ;
               cs[ps].sign
               ps ++;
          }
     }
     for (i=n+2;i<=2*n+1;i++)
          if (S[i] \&\& net[i][m] == 0) {
               //printf("%d
                             +\n", i-n-1);
               cs[ps].num
                             = i-n-1;
               cs[ps].sign
                             = '+' ;
               ps ++;
          }
     printf( "%d\n" , ps);
                           printf( "%d %c\n" , cs[i].num,
     for (i=0;i<ps;i++)
                                                                cs[i].sign);
     //puts("----");
     //for(i=1;i<=m;i++)
                               if(S[i])
                                                      ",i);
                                        printf("%d
     //puts("\n----");
}
                 wout[MAX];
int win[MAX],
int main() {
     int i,j;
     while (scanf( "%d %d" , &n,&m)==2)
          memset(net,
                          0, sizeof (net));
                                     scanf( "%d" , win+i);
          for (i=2;i<=n+1;i++)
          for (i=2;i<=n+1;i++)
                                   scanf( "%d" , wout+i);
          while (m --) {
```

```
int x,y;
             scanf( "%d %d" , &x,&y);
             x ++; y += n+1;
             net[x][y] = INT\_MAX;
        for (i=2;i<=n+1;i++) {
            net[1][i] = wout[i];
        for (i=n+2;i<=2*n+1;i++) {
            net[i][2*n+2] = win[i-n];
         }
         m = 2*n + 2;
         printf( "%d\n" , Edmonds_Karp());
        find_cut();
}
    30. 混合图欧拉回路
// 1637 PKU
bool solve()
    int i, j;
    for (i=2; i<n; i++)
        x[i] = indeg[i] - outdeg[i];
        if (x[i] % 2)
             return false;
        if (x[i] > 0)
             net[i][m] += x[i] >> 1;
         else if (x[i] < 0)
             net[1][i]
                    += (-x[i]) >> 1;
    int cap = 0;
```

for (i=2; i<n; i++)

for (i=2; i < m; i++)

*/

cap += net[1][i];

int flow = Edmonds_Karp();

for (j=2; j< m; j++)

if (net[i][j] != 0)

// when flow==cap, say it exist euler circuit

 $printf("%d -> %d\n", i-1,j-1);$

// print the undirected edge's direction

```
return (flow == cap);
}
int main()
{
    int i, j, cas;
    scanf( "%d" , &cas);
    while (cas --)
    {
        memset(indeg, 0, sizeof (indeg));
        memset(outdeg, 0, sizeof (outdeg));
        memset(net, 0, sizeof (net));
        scanf( "%d %d" , &m, &s);
        for (i=0; i< s; i++)
            int x, y, d;
            scanf( "%d %d %d" , &x, &y, &d);
            x ++; y ++;
            // if d=0, make x->y
            indeg[y] ++;
            outdeg[x]
                     ++;
            if (d == 0)
                net[x][y]
                          ++;
        n = m + 2;
        puts(solve()
                    ? "possible"
                                : "impossible" );
}
    31. 无源汇上下界网络流
    2314 ZJU Reactor Cooling
    无源汇上下界网络流
    (1) 新建 S , T
    (2) D(u) = B(i,u) - B(u,i)
        D(u) > 0, 建弧 (S,u) ,权值为 D(u)
        D(u) < 0, 建弧 (u,T) ,权值为 -D(u)
    (3) 求最大流,判定是否满流
*/
struct NODE
{
    int x, y;
    int b, c;
    NODE (int _x = 0, int _y = 0, int _b = 0, int _c = 0)
```

```
: x(_x), y(_y), b(_b), c(_c) {}
};
vector <NODE>
                   nodes;
int make_net()
{
    int i, j;
     int D[NMAX] = \{0\};
     memset(net, 0, sizeof (net));
     n += 2;
     vector <NODE>::iterator
                                   iter;
     foreach (iter,nodes)
         i = iter->x;
         j = iter->y;
         net[i][j] = iter->c - iter->b;
         D[j] += iter->b;
         D[i] = iter->b;
     int ret = 0;
     for (i=2; i<n; i++)
         if (D[i] > 0)
              net[1][i] = D[i];
         else if (D[i] < 0)
            net[i][n]
                      = - D[i];
         ret += net[1][i];
     return ret;
}
void solve()
{
     int i, j;
     int cap = make_net();
     int flow = Edmonds_Karp();
     if (flow != cap)
          puts( "NO\n" );
     else
         puts( "YES" );
         vector <NODE>::iterator
                                       iter;
         foreach (iter,nodes)
              printf( "%d\n" , iter->c - net[iter->x][iter->y]);
     }
```

```
}
int main()
{
    int i, j, cas;
    scanf( "%d" , &cas);
    while (cas --)
    {
        scanf( "%d %d" , &n, &m);
         nodes.clear();
        for (i=0; i< m; i++)
             int x, y, I, cap;
             scanf( "%d %d %d %d" , &x, &y, &l, &cap);
             x ++; y ++;
             nodes.push_back(NODE(x,y,l,cap));
         }
        solve();
}
     32. 二分图最小点权覆盖
// 3308 PKU Paratroopers
// 2874 ZJU
double R[MAX], C[MAX];
// 二分图最小点权覆盖 -> 网络最大流
void make_net()
{
    int i, j;
    memset(net, 0, sizeof (net));
    // C(S,x) = W[x]
    for (i=0; i< n; i++)
       net[0][i+1] = log(R[i]);
    // C(y,T) = W[y]
    for (i=0; i< m; i++)
        net[n+i+1][n+m+1]
                               = log(C[i]);
    // C(x,y) = \inf
    for (i=0; i<1; i++)
        int x, y;
        scanf( "%d %d" , &x, &y);
        net[x][y+n] = 1e99;
    }
    n = n + m + 1;
```

```
}
double solve()
{
    int i, j;
    double ret;
    make_net();
    ret = Edmonds_Karp();
    return exp(ret);
}
     33. 带约束的轨道计数 (Burnside
                                                     引理)
// PKU 2888
#include
          <stdio.h>
#include
          <math.h>
const int MOD = 9973;
bool A[32000];
int prim[3500],T[10][10];
int total,n,m;
void init()
    int i,j;
    total = 0;
    for (i = 2; i < 32000; i++)
         if (!A[i])
             prim[total++] = i;
             for (j = 2*i; j < 32000; j += i)
                  A[j] = true ;
}
int phi( int x)
{
    int temp,i,num;
    if (x == 1) return 1;
    temp = 1;
    for (i = 0; i < total; i++)
         num = prim[i];
         if (x \% num == 0)
         {
```

```
temp *= num-1;
             temp %= MOD;
             x /= num;
             while (x \% num == 0)
                  x /= num;
                  temp *= num;
                  temp %= MOD;
             if (x == 1) break;
         }
    if (x != 1)
         temp *= x - 1;
         temp %= MOD;
    return temp;
}
void GT( int TT[][10], int p)
    int i,j,sum,k;
    if (p == 1)
         for (i = 0; i < m; i++)
             for (j = 0; j < m; j++)
                  TT[i][j]
                          = T[i][j];
         return ;
    int t2[10][10];
    GT(t2,p/2);
    for (i = 0; i < m; i++)
         for (j = 0; j < m; j++)
         {
             sum = 0;
             for (k = 0; k < m; k++)
                  sum += t2[i][k]*t2[k][j];
             TT[i][j] = sum \% MOD;
         }
    if (p \% 2 == 0) return;
    int t[10][10];
    for (i = 0; i < m; i++)
         for (j = 0; j < m; j++)
```

```
{
             sum = 0;
             for (k = 0; k < m; k++)
                  sum += TT[i][k]*T[k][j];
             t[i][j] = sum \% MOD;
         }
    for (i = 0; i < m; i++)
         for (j = 0; j < m; j++)
             TT[i][j] = t[i][j];
}
int Tr( int p)
{
    int sum = 0,i;
    int TT[10][10];
    GT(TT,p);
    for (i = 0; i < m; i++)
        sum += TT[i][i];
    return sum % MOD;
}
int gn()
{
    int temp, sum = 0, i, all;
    temp = (int)(sqrt(1.0*n)+0.4) + 1;
    sum = (phi(1)*Tr(n)%MOD
                                   + phi(n)*Tr(1)%MOD)
                                                        % MOD;
    for (i = 2; i < temp; i++)
         if (n \% i == 0)
             if (i^*i == n)
             {
                  sum += phi(i)*Tr(i)%MOD;
                  sum %= MOD;
             }
             else
             {
                  sum += phi(i)*Tr(n/i)%MOD
                                                   + phi(n/i)*Tr(i)%MOD;
                  sum %= MOD;
    all = sum / n;
    sum %= n;
```

```
all++;
         sum = n - sum;
         sum %= MOD;
         if (sum == 0) break;
         sum = MOD - sum;
    return all;
}
int main()
{
    int cas,k,x,y,i,j;
    scanf( "%d" ,&cas);
    init();
    while (cas--)
    {
         scanf( "%d %d %d" ,&n,&m,&k);
         for (i = 0; i < m; i++)
             for (j = 0; j < m; j++)
                 T[i][j] = 1;
         while (k--)
             scanf( "%d %d" ,&x,&y);
             X--;
             y--;
             T[x][y] = 0;
             T[y][x] = 0;
         printf( "%d\n" ,gn());
    return 0;
}
    34. 三分法求函数波峰
// linle 专场考研路茫茫 —— 早起看书
const int MMAX = 11000;
const double EPS = 1e-4;
int x[MMAX], y[MMAX];
int n, m;
#define f(dt) k^*(dt-x[p-1]) + y[p-1] + 1.0*n/dt/dt;
```

while (sum)

```
double triple_search(
                       int p)
    double mmin = 1e99;
           k = 1.0 * (y[p]-y[p-1]) / (x[p]-x[p-1]);
     double
           xI = x[p-1], \quad xr = x[p];
     double
    double Im, rm, flm, frm;
    while (fabs(xr-xl)
                      > EPS)
         lm = (2.0*xl + xr) / 3.0;
         rm = (2.0*xr)
                       + xl) / 3.0;
         flm = f(lm);
         frm = f(rm);
         if (frm > flm)
              xr = rm, mmin = min(mmin,
                                               flm);
         else
              xI = Im, mmin = min(mmin,
                                              frm);
    }
    return mmin;
}
double solve()
    int i, j, k;
    double mmin = 1e99;
    for (i=1; i < m; i++)
         double ret = triple_search(i);
         mmin = min(mmin,
                             ret);
    return mmin;
}
int main()
{
    int i, j;
    while (scanf( "%d %d" , &m, &n) == 2)
         for (i=0; i< m; i++)
              scanf( "%d %d" , &x[i],
                                       &y[i]);
         printf( "%.3lf\n"
                         , solve());
    }
}
```

35. 单词计数 , DFA 自动机 , Trie 图

```
// linle 专场考研路茫茫 —— 单词情结
// 由正则表达式构造 NFA , NFA 转 DFA , 最小化 DFA
// 构造状态转移矩阵,矩阵乘法
        unsigned
                  long long ULL;
typedef
        foreach (it,c) for (it=(c).begin();
                                        it!=(c).end();
#define
                                                           it++)
                     for (it=0; it<(c).size();
#define
        forsize(it,c)
                                               it++)
const int NMAX = 6;
int n, l;
char rt[NMAX][6];
const int SMAX = 80;
       ADD(a,x)
#define
                   ((a)=((a)+(x)))
struct MATRIX
{
    ULL mat[SMAX][SMAX];
    int n;
    MATRIX (int _n = SMAX)
         n = _n;
         memset(mat,
                    0, sizeof (mat));
    void to_E( int nn)
         int i;
         n = nn;
                        0, sizeof (mat));
         memset(mat,
         for (i=0; i< n; i++)
             mat[i][i] = 1;
    void fill( const MATRIX & mt, int x, int y)
    {
         int i, j;
         for (i=0; i< mt.n; i++)
             for (j=0; j<mt.n; j++)
                  mat[i+x][j+y] = mt.mat[i][j];
    MATRIX operator * (const MATRIX & mt)
         MATRIX ret;
         int i, j, k;
```

```
for (i=0; i< n; i++)
             for (j=0; j< n; j++)
             {
                  ret.mat[i][j] = 0;
                  for (k=0; k<n; k++)
                                     mat[i][k] * mt.mat[k][j]);
                      ADD(ret.mat[i][j],
             }
         ret.n = n;
         return ret;
    MATRIX operator ^ (int ex)
         int i;
         MATRIX ret, tmp;
         ret = *this;
         tmp.to_E( this ->n);
         while (ex > 1)
             if (ex & 1)
              tmp = tmp * ret;
             ret = ret * ret;
             ex >>= 1;
         ret = ret * tmp;
         return ret;
};
const int NFAMAX = 60;
struct EDGE
    char ch;
    int next;
    EDGE (char _c = 0, int _n = 0)
         : ch(_c), next(_n) {}
};
                 nfa[NFAMAX];
vector <EDGE>
vector <EDGE>
                 dfa[NFAMAX];
vector <EDGE>
                 mindfa[NFAMAX];
int nfact;
int dfasn;
int mindfasn;
vector < int > dfact;
vector < int > mindfact;
```

```
#define
          BADD(x,p)
                        ((x) |= ((ULL)1 << (p)))
          BSUB(x,p)
                        ((x) \&= \sim ((ULL)1 << (p)))
#define
                        ((x) & ((ULL)1<<(p)))
          BGET(x,p)
#define
void make_nfa()
{
     int i, j, k;
     for (i=0; i<NFAMAX;
                              i++)
          nfa[i].clear();
     for (i= 'a'; i \leftarrow 'z'; i++)
          nfa[0].push_back(EDGE(i,0));
     nfact = 1;
     int lend[NMAX];
     for (i=0; i< n; i++)
     {
          nfa[0].push_back(EDGE(
                                         '$' ,nfact++));
          for (j=0; rt[i][j];
                               j++)
               nfa[nfact-1].push_back(EDGE(rt[i][j],nfact));
               nfact ++;
          lend[i] = nfact - 1;
     }
                i<n; i++)
     for (i=0;
          nfa[ lend[i] ].push_back(EDGE(
                                                 '$' ,nfact));
     for (i= 'a'; i <= 'z'; i++)
          nfa[nfact].push_back(EDGE(i,nfact));
     nfact ++;
}
bitset <NFAMAX>
                       vis;
ULL e_closure( int now)
{
     int i, j;
     ULL ret = 0;
     vector <EDGE>::iterator
                                     iter;
     BADD(ret,
                  now);
     if (vis[now])
                  ret;
          return
```

```
= true ;
     vis[now]
     foreach (iter, nfa[now])
          if (iter->ch == '$')
              ret |= e_closure(iter->next);
     return ret;
}
ULL e_closure2(ULL
                        now)
{
     int i, j;
     ULL ret = now;
     vis.reset();
     for (i=0; i<nfact;
                         i++)
          if (BGET(now,
                           i))
              ret |= e_closure(i);
     return ret;
}
map < ULL, int > hash;
void dfs(ULL now)
{
     int i, j;
     vector <EDGE>::iterator
                                   iter;
            < int >::iterator
     vector
                                iter2;
     vector < int > nxt[30];
                i<nfact; i++)
     for (i=0;
     {
          if (BGET(now,i))
               foreach (iter, nfa[i])
               {
                    if (iter->ch
                                  == '$')
                         continue ;
                                   - 'a' ].push_back(iter->next);
                    nxt[iter->ch
          }
     int stag = hash[now];
     for (i= 'a'; i<= 'z'; i++)
```

```
if (nxt[i- 'a' ].empty())
               continue ;
          ULL next = 0;
          foreach (iter2,
                            nxt[i- 'a'])
               BADD(next,
                              *iter2);
          next = e_closure2(next);
          bool flag = false;
          int ntag = hash[next];
          if (ntag == 0)
                                      = dfasn ++, flag = true ;
               ntag = hash[next]
          dfa[stag-1].push_back(EDGE(i,ntag-1));
          if (flag)
               if (BGET(next,
                                 nfact-1))
                    dfact.push_back(ntag-1);
               dfs(next);
          }
}
void nfa_dfa()
     int i, j, k;
     dfasn = 1;
     vis.reset();
     hash.clear();
     dfact.clear();
     for (i=0; i< NFAMAX;
                               i++)
          dfa[i].clear();
     ULL bs = e_{closure}(0);
     hash[bs]
                = dfasn ++;
     dfs(bs);
}
void min_dfa()
{
     int i, j, k;
     vector < vector < int > > split;
     vector <EDGE>::iterator
                                    iter;
     int belg[NFAMAX];
```

```
for (i=0; i<dfasn;
                       i++)
     vector < int > newi;
     newi.push_back(i);
     split.push_back(newi);
     belg[i] = i;
}
bool flag = true ;
while (flag)
     flag = false;
     for (i=0; i<split.size();
                                 i++)
     {
                        j<split.size();
          for (j=i+1;
                                           j++)
          {
                vector < pair < char , int > > ibel, jbel;
                for (k=0; k<split[i].size();
                                                  k++)
                     foreach (iter, dfa[ split[i][k]
                                                         ])
                          ibel.push_back(make_pair(iter->ch,
                                                                         belg[iter->next]));
                for (k=0; k<split[j].size();
                                                  k++)
                     foreach (iter, dfa[ split[j][k]
                                                         ])
                          jbel.push_back(make_pair(iter->ch,
                                                                         belg[iter->next]));
                sort(ibel.begin(),
                                      ibel.end());
                sort(jbel.begin(),
                                      jbel.end());
                if (ibel == jbel)
                     flag = true ;
                     break ;
          if (flag)
                break ;
     }
     if (flag)
     {
          int s1 = belg[ split[i][0]
                                        ], s2 = belg[ split[j][0]
                                                                       ];
          for (k=0; k<dfasn;
                                   k++)
                if (belg[k]
                             == s2)
                     belg[k] = s1;
          split[i].insert(split[i].end(),
                                             split[j].begin(),
                                                                  split[j].end());
          split.erase(split.begin()
                                         + j);
     }
}
```

```
for (i=0; i<split.size(); i++)
         for (j=0; j<split[i].size(); j++)
              belg[ split[i][j]
                               ] = i;
     bitset <NFAMAX>
                          acts;
     for (i=0; i<dfact.size();
                              i++)
         acts[ dfact[i] ] = true ;
     mindfact.clear();
     mindfasn = split.size();
    for (i=0; i<mindfasn;
                              i++)
         int go[30];
                       -1, sizeof (go));
          memset(go,
          mindfa[i].clear();
          flag = false;
         for (j=0; j<split[i].size(); j++)
          {
              if (acts[ split[i][j] ])
                   flag = true ;
              foreach (iter, dfa[ split[i][j]
                   go[iter->ch - 'a'] = belg[iter->next];
         for (j= 'a'; j <= 'z'; j++)
              if (go[j- 'a'] != -1)
                   mindfa[i].push_back(EDGE(j,go[j-
                                                    'a' ]));
          if (flag)
              mindfact.push_back(i);
}
MATRIX T;
MATRIX TT;
MATRIX BT;
MATRIX E;
void make_matrix()
{
    int i, j;
    vector <EDGE>::iterator
                                   iter;
     E.to_E(mindfasn);
```

```
T.n = mindfasn;
     memset(T.mat,
                       0, sizeof (T.mat));
     for (i=0; i<mindfasn;
                               i++)
          foreach (iter, mindfa[i])
              T.mat[i][iter->next]
                                       ++;
// 构造等比矩阵
     BT.n = mindfasn << 1;
     memset(BT.mat,
                         0, sizeof (BT.mat));
     BT.fill(T, 0, 0);
     BT.fill(E,
               0, mindfasn);
               mindfasn,
     BT.fill(E,
                            mindfasn);
}
ULL solve()
{
     int i, j;
     ULL ret = 0;
     vector <EDGE>::iterator
                                   iter;
     make_nfa();
     nfa_dfa();
     dfasn = hash.size();
     min_dfa();
     make_matrix();
     BT = BT \wedge I;
     TT.n = mindfasn;
               i<mindfasn;
     for (i=0;
                               i++)
          for (j=0; j<mindfasn;
                                    j++)
                         = BT.mat[i][j+mindfasn];
              TT.mat[i][j]
     T = T * TT;
     for (i=0; i<mindfact.size(); i++)
         ret += T.mat[0][
                              mindfact[i] ];
     return ret;
}
int main()
{
     int i, j;
     while (scanf( "%d %d" , &n, &l) == 2)
          for (i=0; i<n; i++)
```

```
scanf( "%s" , rt[i]);
         printf( "%I64u\n" , solve());
    }
}
// HDU 2471 History of Languages
// DFA 同构判断
#include
         <iostream>
         <string>
#include
#include
         <algorithm>
using namespace
                 std;
const int NMAX = 2010;
const int TMAX = 26;
const int FAIL = -1;
int sigma; // 字符集大小
bool dis[NMAX][NMAX];
int lx[NMAX*NMAX], ly[NMAX*NMAX];
struct DFA {
    int X[NMAX][TMAX];
    bool F[NMAX];
    int n;
    void read() {
         int i, j, k;
         scanf( "%d" , &n);
         for (i = 0; i < n; i++) {
             scanf( "%d" , &k);
             F[i] = (k == 1);
             for (j = 0; j < sigma; j++) {
                  scanf( "%d" , &X[i][j]);
                 if (X[i][j] == -1) X[i][j] = n;
             }
         }
         // 虚拟一个非接受态节点,再补边
         F[n] = false;
         for (i = 0; i < sigma; i++)
             X[n][i] = n;
         //n++;
    void Minization(DFA
                        &dst) {
         memset(dst.X,
                        -1, sizeof (dst.X));
```

```
int i, j, ch;
int p, q;
bool vis[NMAX] = { false };
vis[0] = true;
fill(vis, vis + n, true );
fill(dis[0], dis[n], false);
int Q[NMAX], Qf, Qr;
Qf = Qr = 0;
Q[Qr++] = 0;
while (Qf < Qr) {
    p = Q[Qf++];
    for (ch = 0; ch < sigma; ch++) {
         q = X[p][ch];
         if (q != -1 \&\& !vis[q]) {
             vis[q] = true;
             Q[Qr++] = q;
}
for (i = 0; i < n; i++) {
    if (!vis[i]) continue ;
    for (j = i + 1; j < n; j++) {
         if (!vis[j])
                   continue ;
        if (F[i] != F[j]) dis[i][j] = dis[j][i] = true;
}
while (1) {
    bool update = false;
    for (i = 0; i < n; i++)
         if (!vis[i]) continue ;
         for (j = i + 1; j < n; j++) {
             if (!vis[j]) continue ;
                       == true ) continue ;
             if (dis[i][j]
             for (ch = 0; ch < sigma; ch++) {
                  p = X[i][ch];
                  q = X[j][ch];
                  if (p == -1 \&\& q == -1) continue;
                  if (p == -1 || q == -1 || dis[p][q]) {
```

= true ;

dis[i][j] = dis[j][i]

```
update = true ;
                           break ;
                       }
                  } //update
             } //for __j
        } //for i
         if (!update)
                      break ;
    } //while
    int id[NMAX], cnt = 0;
    fill(id, id + n, -1);
    for (i = 0; i < n; i++)
         if (!vis[i]) continue;
         if (id[i] != -1) continue;
         for (j = 0; j < n; j++) {
             if (!vis[j])
                        continue ;
             if (dis[i][j]
                        == false ) id[j] = cnt;
         }
         cnt++;
    dst.n = cnt;
    for (i = 0; i < n; i++)
         if (!vis[i]) continue ;
         if (id[i] != -1) continue ;
         p = id[i];
         dst.F[p] = F[i];
         for (ch = 0; ch < sigma; ch++) {
              q = X[i][ch];
              if (q != FAIL) q = id[q];
              dst.X[p][ch]
                          = q;
         }
    }
}
void show() {
    int i, j;
    for (i = 0; i < n; i++) {
         cout << F[i] << "";
         for (j = 0; j < sigma; j++)
              cout << X[i][j]
                             << " ";
```

```
cout << endl;
         }
         cout << endl;
    }
    bool equals(DFA
                       &dfa2) {
                      0 , sizeof (dis));
         memset(dis,
         int p, q;
         p = q = 0;
         lx[q] = 0;
         ly[q] = 0;
         dis[0][0] = true ;
         q ++;
         while (p < q) {
             if (F[ lx[p] ] != dfa2.F[ ly[p] ]) break;
             for (int k = 0; k < sigma;k++)
                  if (! dis[ X[lx[p]][k] ][ dfa2.X[ly[p]][k]
                                                          ]) {
                                                         ] = true ;
                       dis[ X[lx[p]][k] ][ dfa2.X[ly[p]][k]
                       Ix[q] = X[Ix[p]][k];
                       ly[q] = dfa2.X[ ly[p] ][k];
                       q ++;
              p ++;
         return p == q;
};
DFA a, b;
int main() {
    int cas = 1;
    while (scanf( "%d" , &sigma), sigma) {
         a.read();
         b.read();
         printf( "Case #%d: ", cas++);
         if (a.equals(b)) puts( "Yes" );
         else puts( "No" );
    return 0;
}
/*
```

```
Trie 图
在Tire 树的基础上补边(类似 AC 自动机)
可用于多字符串匹配和自动机的构造
图可转换为矩阵,或拓扑排序,用作统计或动态规划之用
*/
typedef pair < int , int > PII;
typedef vector <PII> TRANS;
const int VMAX = 200;
const int SIGMAX = 50;
char d2c[110] = "ACGT"; // 字符集
struct TrieGraph
{
   struct NODE
       int suffix; // 后缀节点指针
        int father; // 父节点指针
        int next[SIGMAX]; // 儿子节点指针
        bool mark; // 标记是否出现过
        char ch; // 入边信息
   };
   NODE mem[VMAX];
   int vn, root;
   char c2d[300]; // 字符集 hash 表
   int siglen; // 字符集大小
   int new_node( char ch, int fat) {
        memset(mem+vn,
                          0, sizeof (NODE));
        mem[vn].ch
                 = ch;
        mem[vn].father = fat;
        return vn ++;
   // 初始化(字符集)
   void init( char * pstr) {
        vn = 0;
        siglen = strlen(pstr);
        for (int i=0; pstr[i];
                           i++)
           c2d[pstr[i]] = i;
       // 初始化节点,用作安全转移
        root = new_node( '$', 0);
        for (int i=0; i < siglen; i++)
           mem[root].next[i] = new_node(pstr[i],
                                                  root);
   }
   // 构造 Trie
   void insert( char * pstr) {
```

```
int i;
          for (i=root;
                              pstr++)
                        *pstr;
                                         {
              int x = c2d[*pstr];
              if (mem[i].next[x]
                                         0)
                   mem[i].next[x]
                                       = new_node(*pstr,
                                                               i);
              i = mem[i].next[x];
         mem[i].mark
                          = true ;
    int get_suffix(
                     int idx) {
         int fat = mem[idx].father;
         if (fat == root) return
                                   root;
         int ich = c2d[mem[idx].ch];
         for (fat=mem[fat].suffix;
                                                    && mem[fat].next[ich]==0;
                                        fat!=root
fat=mem[fat].suffix)
         if (mem[fat].next[ich]
                                     == 0) return
                                                     root;
         return mem[fat].next[ich];
    // 构造 Trie Graph
    void construct()
         queue < int > sq;
         for (int i=0; i<siglen;
                                    i++) {
              if (mem[root].next[i]
                                         == 0) continue ;
                                                               ];
              NODE & son = mem[ mem[root].next[i]
                          = get_suffix(mem[root].next[i]);
              son.suffix
              sq.push(mem[root].next[i]);
         while (! sq.empty())
              int idx = sq.front();
                                       sq.pop();
              NODE & now = mem[idx];
              for (int i=0; i<siglen;
                                         i++)
                   int sonidx = now.next[i];
                    NODE & son = mem[sonidx];
                   if (sonidx == 0) continue;
                   sq.push(sonidx);
                                = get_suffix(sonidx);
                    son.suffix
                                            || mem[son.suffix].mark;
                               = son.mark
                    son.mark
              }
              for (int i=0; i<siglen;
                                         i++) {
                   if (now.next[i]
                                     != 0) continue ;
                   now.next[i]
                                  = mem[now.suffix].next[i];
              }
         }
```

```
// 构造安全图
    TRANS make_safe_graph()
         queue < int > sq;
         bitset <VMAX>
                          vis;
         TRANS ret;
         sq.push(root);
         vis[root] = true ;
         while (! sq.empty())
             int idx = sq.front();
                                   sq.pop();
             NODE & now = mem[idx];
             for (int i=0; i<siglen; i++) {
                 int sonidx = now.next[i];
                               || mem[sonidx].mark)
                  if (sonidx==0
                      now.next[i] = 0; // 更新 Trie Graph
                      continue ;
                  }
                  ret.push_back(PII(idx,
                                          sonidx));
                  if (vis[sonidx])
                                  continue ;
                  sq.push(sonidx);
                 vis[sonidx]
                            = true ;
             }
         }
         return
                ret;
    void print()
         printf( "%6s%6s%6s...\n"
                                         , "Node" , "Suff" , "Mark" , "Son" );
         for (int i=0; i<vn; i++) {
             printf( "%6d%6d%6d"
                                     , i, mem[i].suffix,
                                                        mem[i].mark);
             for (int j=0; j<siglen;
                                     j++)
                  printf( "%6d" , mem[i].next[j]);
             puts( "");
        }
    }
};
/*
PKU 1625 Censored!
求长度为 m , 字符集为 n 且不含 p 个不良单词的字符串的数目 ,
就是求在安全图中从根结点出发走
                                m 步有多少种走法。
                表示从根结点出发走 step 步到结点 x 的走法数。
用count[step,x]
fillchar(count,sizeof(count),0);
count[0, 根]:=1;
for step:=1
            to m do
    for 安全图中每条边 (i,j) do
         inc(count[step,j],count[step-1,i]);
```

```
ans:=0;
for 安全图中每个结点 x do
    inc(ans,count[m,x]);
*/
/*
PKU 2778 DNA Sequence
用矩阵做状态转移,矩阵二分求答案
int n, m, p;
xnum dp[2][VMAX];
TrieGraph tg;
void solve()
{
    int i, j;
    xnum ret = 0;
    TRANS tr = tg.make_safe_graph();
    for (i=0; i< tg.vn; i++) dp[0][i] = dp[1][i] = 0;
    dp[0][0]
              = 1;
    for (i=1; i \le m; i++)
         for (j=0; j<tr.size(); j++)
             dp[i\&1][tr[j].second] = dp[i\&1][tr[j].second]
dp[(i&1)^1][tr[j].first];
         for (j=0; j<tg.vn; j++) dp[(i&1)^1][j] = 0;
    }
    for (i=0; i< tg.vn; i++)
         ret = ret + dp[m&1][i];
                     "" );
    ret.print();puts(
}
int main()
{
    int i, j;
    char str[110];
    while (scanf( "%d %d %d" , &n, &m, &p) == 3)
         scanf( "%s" , d2c);
         tg.init(d2c);
         for (i=0; i< p; i++)
             scanf( "%s" , str);
```

```
tg.insert(str);
         }
         tg.construct();
         //tg.print();
         solve();
}
     36. 字符串和数值 hash
// 整数 hash
// 104729,
            224737,
                      350377,
                                479909,
                                          611953,
                                                    882377
// 1020379,
              1299709,
                         1583539,
                                     1870667,
                                                2015177
// 4256233,5800079,7368787,
                                   10570841,
                                                15485863
const int MOD = 20023;
bool bhash[MOD];
int vhash[MOD];
int cnt[MOD];
bool find_hash( int & pos) {
    int val = pos;
    pos %= MOD;
    for (; bhash[pos];
                         pos=(pos+1)%MOD)
         if (vhash[pos]
                         == val)
             return true;
    return false;
}
int make_hash(
                 int val) {
    int pos = val;
    if (! find_hash(pos))
                           {
         bhash[pos]
                     = true ;
         vhash[pos]
                     = val;
         cnt[pos] = 0;
    cnt[pos]
              ++;
    return pos;
}
// 字符串 hash
const int MOD = 20023;
bool bhash[MOD];
char vhash[MOD][45];
char str[45];
```

```
int cal_str() {
    int i, j, pos;
    for (i=pos=0,j=1; str[i]; i++,j=(j*27)&INT_MAX,pos&=INT_MAX)
        int num = str[i] - 'a';
        if (str[i] == ' ')
            num = 26;
        pos += j*num;
    return pos % MOD;
}
bool find_hash( int & pos) {
    pos = cal_str();
    for (; bhash[pos]; pos=(pos+1)%MOD)
        if (strcmp(vhash[pos], str) == 0)
            return true;
    return false;
}
int make_hash()
    int pos;
    if (! find_hash(pos))
                        {
        bhash[pos] = true ;
        strcpy(vhash[pos],
                             str);
    return pos;
}
    37. 滚动队列,前向星表示法
int que[2][2000];
int qf[2],qe[2],qnow;
#define
                       (que[qnow][
                                     qe[qnow]++
         push_que(a)
                                                  ] = (a)
                      (que[qnow^1][
                                       qf[qnow^1]++
         pop_que2
                                                       ])
#define
                      qnow ^= 1;
#define
         switch_que
                      qf[qnow]
                                = qe[qnow]
                                             = 0;
#define
         empty_que2
                          (qf[qnow^1]
                                        >= qe[qnow^1]
#define
                      (qe[qnow^1]
                                    - qf[qnow^1])
         size_que2
前向星表示法
空间 O(E+N)
存储所有边,并用链表来实现读取
                               s为起点的有向边
```

```
方便插入和遍历所有边 , 删除是 O(E)
*/
const int ENDFLAG = -1;
struct EDGELIST {
    int start[NMAX];
    int last[NMAX];
    int edge[MMAX][2];
                       //pos,listnext
    int tot;
    void clear() {
         tot = ENDFLAG +1;
         memset(last,ENDFLAG, sizeof (last));
    }
    void push_back(
                     int s, int t) {
         edge[tot][0]
                     = t;
         edge[tot][1]
                     = ENDFLAG;
         if (last[s] != ENDFLAG) edge[ last[s] ][1] = tot;
         else start[s] = tot;
         last[s] = tot;
         tot ++;
    }
    int get_start( int s) {
         return start[s];
    int get_next( int & p) {
         p = edge[p][1];
         return edge[p][0];
    void erase( int s, int t) {
         int i,pre = ENDFLAG;
         int p,v;
         for (p=start[s]; p!=ENDFLAG; p=edge[p][1]) {
             v = edge[p][0];
             if (v == t) {
                  if (pre == ENDFLAG) start[s] = edge[p][1];
                  else edge[pre][1]
                                   = edge[p][1];
             }
             else pre = p;
         last[s] = pre;
};
```

38. 最小点基,最小权点基

```
// HDOJ 1827 Summer Holiday
// 点基:通过点基的点,能够到达有向图全部点
// 最小权点基:有向图顶点有权值
void Gabow()
{
    int i,j,l;
    //dfs in original graph
    memset(id, 0, sizeof (id));
    memset(vis, 0, sizeof (vis));
    scc = step = 1;
    order_pos = order2_pos = 0;
    for (i=1; i <= n ; i++) {
        if (vis[i] == 0) {
             dfs(i);
        }
    }
    SCC --;
}
void top_sort()
    int i,j,k,l,m = 0;
    memset(out_degre,0, sizeof (out_degre));
                  0x7f, sizeof (sel));
    memset(sel,
    I = SQ.size();
    while (I --) {
         SQ.pop();
    for (i=1;i<=n;i++)
        int id1 = id[i];
        I = path[i].size();
         for (j=0;j<1;j++)
             int id2 = id[ path[i][j]
             if (id1 != id2) {
                 out_degre[id2]
                                   ++;
                 dag[id1].push_back(id2);
         }
    for (i=1;i<=scc;i++)
         if (out_degre[i]
                           == 0) {
             SQ.push(i);
        }
    }
```

```
while (!SQ.empty()) {
         int now = SQ.front();
         SQ.pop();
        I = dag[now].size();
         for (i=0;i<1;i++) {
             int next = dag[now][i];
             out_degre[next]
             if (out\_degre[next] == 0) {
                  SQ.push(next);
                  m --; //find non-hightest
                                              SCC
                  sel[next] = -1;
         }
    for (i=1;i<=n;i++)
        if (sel[ id[i] ] != -1) { //selection mininum cost in the highest scc
             sel[id[i]] = sel[id[i]] > cost[i] ? cost[i] : sel[id[i]];
         }
    min\_cost = 0;
    for (i=1;i<=scc;i++) {
        if (sel[i] != -1) {
             min_cost += sel[i];
        }
    min_num
               = scc+m;
}
int main()
    int i,x,y;
    path.resize(NMAX);
    dag.resize(NMAX);
    while (scanf( "%d %d" ,&n, &m)==2)
         for (i=0;i<=n;i++)
             path[i].clear();
             dag[i].clear();
         }
         for (i=1;i<=n;i++)
             scanf( "%d" , cost+i);
         }
         for (i=1;i<=m;i++)
             scanf( "%d %d" , &x,&y);
             path[x].push_back(y);
```

```
}
        Gabow();
        top_sort();
       //min_num
                   : mininum
                              vertex number
        //min_cost
                   : mininum
                              cost
        printf( "%d %d\n" , min_num,
                                      min_cost);
}
    39. LCSubsequence O(N^2/logN)
// 1210 WHU
    LCSubsequence O(N^2/logN)
    这个解法是在字符集不大的情况下,先预处理,再用位运算做状态转移。
*/
typedef UL data_type; // 变量存储类型
const int NMAX = 31000; // 字符串长度
const int BITLEN = sizeof (data_type)*8; // 变量存储位长度
const int BINLEN = 5; // 2^BINLEN = BITLEN
const int MMAX = (NMAX/BITLEN) +1; // 申请空间长度
// ((x)/BITLEN)
#define GETBLOCK(x) ((x)>>BINLEN)
char str1[NMAX];
char str2[NMAX];
struct BITSET
    data_type
              dat[MMAX];
    int len;
    int bs;
    BITSET (int I = 0) {
        len = I;
        bs = GETBLOCK(I+BITLEN-1);
                     0, sizeof (dat));
        memset(dat,
    }
                  [] (int p) {
    bool operator
        return (dat[GETBLOCK(p)]
                                   & ((data_type)1<<(p%BITLEN)));
    void set( int p, bool flag) {
        if (! flag)
                                = \sim ((data_type)1 << (p\%BITLEN));
            dat[GETBLOCK(p)]
```

```
else
         dat[GETBLOCK(p)] |= ((data_type)1<<(p%BITLEN));
}
void reset( int I) {
    len = I;
    bs = GETBLOCK(I+BITLEN-1);
    memset(dat, 0, sizeof (dat));
BITSET operator ~ () {
    BITSET ret = * this ;
    int i;
    for (i=0; i< bs; i++)
         ret.dat[i] = \sim ret.dat[i];
     return ret;
BITSET operator & (const BITSET & a) {
    BITSET ret = * this;
    int i;
    for (i=0; i< bs; i++)
         ret.dat[i] &= a.dat[i];
     return ret;
BITSET operator | (const BITSET & a) {
    BITSET ret = * this;
    int i;
    for (i=0; i< bs; i++)
                 |= a.dat[i];
         ret.dat[i]
     return ret;
BITSET operator ^ (const BITSET & a) {
    BITSET ret = * this;
     int i;
    for (i=0; i< bs; i++)
         ret.dat[i]
                    ^= a.dat[i];
    return ret;
}
BITSET & operator <<= (int I) {
    int i, j;
    int II = I \% BITLEN;
    I /= BITLEN;
    for (i=bs-l;
                I && i>=0; i--)
         dat[i] = dat[i-l];
    for (i=bs-l; II && i>0; i--)
         dat[i] = (dat[i] << II) | (dat[i-1] >> (BITLEN-II));
```

```
dat[0]  <<= II;
         return * this ;
    }
     BITSET operator - (const BITSET & a) {
         BITSET ret = * this ;
         int i, borw = 0, tborw;
         for (i=0; i< bs; i++) {
              if (ret.dat[i] < a.dat[i] + borw)</pre>
                  tborw = 1;
              else
                  tborw = 0;
              ret.dat[i] -= a.dat[i] + borw;
              borw = tborw;
         }
         return ret;
    }
    int count( int I = 0) {
         int i, j, ret = 0;
         I = (I==0) ? len : I;
         for (i=0; i<bs && l>0; i++,l-=BITLEN)
              data_type tmp = dat[i];
              int tl = l;
              for (; tmp && tl; tmp>>=1,tl--)
                 ret += (tmp & 1);
         return ret;
};
BITSET ext[300];
BITSET row, X;
int BitLCS( char * s1, char * s2)
{
    int i, j;
     int len1 = strlen(s1);
     int len2 = strlen(s2);
    for (i=0; i<300;
                      i++)
         ext[i].reset(len1);
     row.reset(len1);
    X.reset(len1);
    for (i=0; i<len1; i++)
```

```
ext[ s1[i] ].set(i, 1);
    for (i=0; i<len2; i++)
        X = row \mid ext[ s2[i] ];
        row <<= 1;
        row.set(0, 1);
        row = X & ((X-row) & ^{\wedge} X);
    }
    return row.count(len1);
}
int main()
{
    while (scanf( "%s %s" , str1, str2) == 2)
        printf( "%d\n" , BitLCS(str1, str2));
}
    40. 伸展树
/*
伸展树
二叉查找树的改进,平摊复杂度都是 O(log n)
维护序列,适用于统计对象次序发生大规模变化
有翻转和移动时,线段树不适用,且效率高于块状链表
HDOJ 1890 Robotic Sort
*/
const int MMAX = 101000;
struct NODE
    int key, cnt; // 键值,重复次数
    NODE * pl, * pr;
    NODE * pf;
};
NODE mem[MMAX];
int mempos;
NODE * root;
inline NODE * new_node()
{
    NODE * pt = mem[mempos]
                                 ++];
               0, sizeof (NODE));
    memset(pt,
    return pt;
}
// x = L[y]
```

```
inline void Zig(NODE * y)
{
    NODE * x = y - pl;
    NODE * z = y - pf;
    y->pl = x->pr;
    if (y->pI) y->pI->pf = y;
    x->pr = y; y->pf = x;
    if (! z)
    {
       x - pf = NULL;
        return ;
    }
    if (z->pl == y) z->pl = x;
    else z->pr = x;
    x \rightarrow pf = z;
}
// x = R[y]
inline void Zag(NODE * y)
{
    NODE * x = y - pr;
    NODE * z = y-pf;
    y->pr = x->pl;
    if (y->pr) y->pr->pf = y;
    x->pl = y; y->pf = x;
    if (! z)
         x \rightarrow pf = NULL;
         return ;
    if (z->pl == y) z->pl = x;
    else z \rightarrow pr = x;
    x \rightarrow pf = z;
}
// y = L[z], x = L[y]
inline void ZigZig(NODE * z)
{
    NODE * y = z \rightarrow pl;
    NODE * x = y - pI;
    NODE * gz = z - pf;
    y->pl = x->pr;
```

```
if (y->pl) y->pl->pf = y;
     z \rightarrow pl = y \rightarrow pr;
     if (z->pl) z->pl->pf = z;
     x->pr = y; y->pf = x;
     y-pr = z; z-pf = y;
     if (! gz)
     {
      x->pf = NULL;
        return ;
     }
     if (gz \rightarrow pl = z) gz \rightarrow pl = x;
     else gz \rightarrow pr = x;
     x - pf = gz;
}
// y = R[z], x = L[y]
inline void ZigZag(NODE * z)
{
     NODE * y = z \rightarrow pr;
     NODE * x = y-pI;
     NODE * gz = z - pf;
     y->pl = x->pr;
     if (y->pl) y->pl->pf = y;
     z \rightarrow pr = x \rightarrow pl;
     if (z->pr) z->pr->pf = z;
     x->pl = z; z->pf = x;
     x->pr = y; y->pf = x;
     if (! gz)
     {
          x->pf = NULL;
          return ;
     }
     if (gz \rightarrow pl = z) gz \rightarrow pl = x;
     else gz \rightarrow pr = x;
     x \rightarrow pf = gz;
}
// y = R[z], x = R[y]
inline void ZagZag(NODE * z)
{
     NODE * y = z - pr;
     NODE * x = y-pr;
```

```
NODE * gz = z - pf;
     y->pr = x->pl;
     if (y->pr) y->pr->pf = y;
     z \rightarrow pr = y \rightarrow pl;
     if (z->pr) z->pr->pf = z;
     x->pl = y; y->pf = x;
     y->pl = z; z->pf = y;
     if (! gz)
     {
        x-pf = NULL;
         return ;
     }
     if (gz \rightarrow pl = z) gz \rightarrow pl = x;
     else gz \rightarrow pr = x;
     x - pf = gz;
}
// y = L[z], x = R[y]
inline void ZagZig(NODE * z)
{
     NODE * y = z \rightarrow pl;
     NODE * x = y - pr;
     NODE * gz = z - pf;
     y->pr = x->pl;
     if (y-pr) y-pr-pf = y;
     z \rightarrow pl = x \rightarrow pr;
     if (z->pl) z->pl->pf = z;
     x->pl = y; y->pf = x;
     x \rightarrow pr = z; z \rightarrow pf = x;
     if (! gz)
     {
          x->pf = NULL;
          return ;
     }
     if (gz \rightarrow pl = z) gz \rightarrow pl = x;
     else gz \rightarrow pr = x;
     x \rightarrow pf = gz;
}
NODE * splay_slow(NODE * x)
{
```

```
if (! x) return NULL;
    while (x->pf)
    {
         NODE * y = x-pf;
         if (y-p) = x) Zig(y);
         else Zag(y);
    }
    return x;
}
NODE * splay(NODE * x)
{
    if (! x) return NULL;
    while (x->pf)
    {
         NODE * y = x-pf;
         NODE * z = y-pf;
         if (z)
             if (z->pl == y)
             {
                 if (y->pl == x) ZigZig(z);
                 else ZagZig(z);
             }
             else
                 if (y->pr == x) ZagZag(z);
                  else ZigZag(z);
             }
         else
             if (y-p) == x) Zig(y);
             else Zag(y);
        }
    return x;
}
NODE * find( int val, NODE * rt)
    NODE * x = rt;
    NODE * pre = rt;
    while (x)
```

```
{
         if (x->key == val) return x;
         pre = x;
         if (val < x->key) x = x->pl;
         else x = x - pr;
    return pre;
}
// make sure all_elem(rt1) <= all_elem(rt2)
NODE * join(NODE * rt1, NODE * rt2)
{
    if (rt1) rt1->pf = NULL;
    if (rt2) rt2->pf = NULL;
    if (! rt1) return rt2;
    if (! rt2) return rt1;
    NODE * x = find(INT\_MAX, rt1);
    rt1 = splay(x);
    rt1->pr = rt2;
    rt2->pf = rt1;
    return rt1;
}
NODE * split( int val)
{
    NODE * x = find(val, root);
    if (x == NULL || x->key != val) return NULL;
    root = splay(x);
    NODE * newroot = root->pr;
    newroot->pf = NULL;
    root = root->pl;
    root->pf = NULL;
    return newroot;
}
void insert( int val)
    if (root == NULL)
        root = new_node();
        root->key = val;
        root->cnt = 1;
        return ;
```

```
NODE * x = find(val, root);
    if (x->key == val) x->cnt ++;
    else
    {
         NODE * pnew = new_node();
         pnew->key = val;
         pnew->cnt = 1;
         pnew->pf = x;
         if (val < x->key) x->pl = pnew;
         else x \rightarrow pr = pnew;
}
void remove( int val)
{
    NODE * x = find(val, root);
    root = splay(x);
    if (root && root->key == val)
         root->cnt --;
         if (root->cnt == 0)
            root = join(root->pl, root->pr);
}
void print_tree(NODE * rt)
{
    if (rt == NULL) return ;
    printf( " (" );
    print_tree(rt->pl);
    if (rt->pf == NULL)
         printf( " [%d] ", rt->key);
    else
         printf( " %d ", rt->key);
    print_tree(rt->pr);
    printf( ") ");
}
void test_tree()
{
    int v;
    char cmd[2];
    root = NULL;
```

```
mempos = 0;
    while (scanf( "%s %d" , cmd, &v) == 2)
        if (cmd[0] == 'i')
            insert(v);
        else if (cmd[0]
                       == 'r' )
            remove(v);
        else if (cmd[0]
                       == 's')
            root = splay(find(v,
                                root));
        print_tree(root);
        puts( "");
}
int main()
{
    test_tree();
    return 0;
}
    41. Treap
/*
Treap
是有随机数满足堆的性质的二叉搜索树
其结构相当于以随机顺序插入的二叉搜索树
其基本操作的期望复杂度为 O(log n)
其特点是实现简单,效率高于伸展树并且支持大部分基本功能,性价比很高
*/
#define
        MAX 100
typedef
        struct
    int I,r,key,fix;
}node;
class treap
public :
    node p[MAX];
    int size,root;
    treap()
    {
        srand(time(0));
        size=-1;
        root=-1;
```

```
void rot_I( int &x)
     int y=p[x].r;
     p[x].r=p[y].l;
     p[y].l=x;
     x=y;
void rot_r( int &x)
     int y=p[x].I;
     p[x].l=p[y].r;
     p[y].r=x;
     x=y;
void insert( int &k, int tkey)
{
     if (k==-1)
          k=++size;
          p[k].l=p[k].r=-1;
          p[k].key=tkey;
           p[k].fix=rand();
     }
     else
          if (tkey<p[k].key)
                insert(p[k].l,tkey);
                if (p[ p[k].l ].fix>p[k].fix)
                      rot_r(k);
          }
           else
          {
                insert(p[k].r,tkey);
                if (p[p[k].r].fix>p[k].fix)
                      rot_l(k);
          }
}
void remove( int &k, int tkey)
```

```
{
     if (k==-1)
                    return ;
     if (tkey<p[k].key)
           remove(p[k].l,tkey);
     else if (tkey>p[k].key)
           remove(p[k].r,tkey);
     else
                              && p[k].r==-1)
           if (p[k].l==-1
                k=-1;
           else if (p[k].l==-1)
                k=p[k].r;
           else if (p[k].r==-1)
                k=p[k].I;
           else
                if (p[ p[k].I ].fix < p[ p[k].r
                                                     ].fix)
                      rot_l(k);
                      remove(p[k].I,tkey);
                }
                else
                      rot_r(k);
                      remove(p[k].r,tkey);
     }
int find( int k, int r)
     if (r \le p[p[k].l].size)
           return
                  find(p[k].l,r);
     else if (r> p[p[k].l].size+p[k].cnt)
           return \qquad find(p[k].r,r-(p[p[k].l].size+p[k].cnt)); \\
     return p[k].key;
}
void print( int k)
     if (p[k].l!=-1)
           print(p[k].l);
     cout << p[k].key
                             << " : " << p[k].fix << endl;
     if (p[k].r!=-1)
           print(p[k].r);
```

```
};
treap T;
int main()
{
    int i;
    for (i=3;i>=1;i--)
         T.insert(T.root,i);
    T.print(T.root);
    for (i=3;i>=1;i--)
    {
         cout << endl;
         T.remove(T.root,i);
         T.print(T.root);
    }
    return 0;
}
                  分数规划 K 约束
     42. 0/1
// 2976 PKU Dropping
                        tests
// 0-1 Fractional Programing
                               with K drop limit
(a_i) / (b_i) >= x
               a_i >= (x^*b_i)
       <=>
               (a_i - x^*b_i) >= 0.
       <=>
let w_i = a_i - x^*b_i.
\max Q(x) = \max (w_i).
so drop K smallest w_i is ok.
*/
const int NMAX = 1100;
int n, k;
int A[NMAX],
              B[NMAX];
struct NODE
{
    double w;
    int a, b;
    bool operator < (const NODE & nt) const
    {
         return w < nt.w;
    }
```

```
}W[NMAX];
#define
         EQ(a,b) \quad (fabs((a)-(b))<1e-4)
// Dinkelbach
               iterative
                         algorithm
int solve()
{
    int i, j;
    double x = 1.0;
    for (i=0; i<100; i++)
    {
         for (j=0; j< n; j++)
         {
             W[j].w = 1.0*A[j]
                                 - x*B[j];
             W[j].a = A[j];
             W[j].b = B[j];
         }
         sort(W, W+n);
         double sum = 0;
         double sa, sb;
         sa = sb = 0;
         for (j=k; j< n; j++)
             sum += W[j].w;
             sa += W[j].a;
             sb += W[j].b;
         if (EQ(sum, 0)) break;
         x = 1.0 * sa / sb;
    return (x*100+0.5);
}
// binary enum
int solve2()
{
    int i, j;
    double lb = 0, ub = 1;
    double x = 1, prex;
    for (i=0; i<100; i++)
         prex = x;
         x = (lb+ub) / 2;
```

if (EQ(x, prex)) break;

for (j=0; j< n; j++)

```
sort(W, W+n);
         double sum = 0;
         for (j=k; j< n; j++)
             sum += W[j].w;
         if (sum >= 0) lb = x;
         else ub = x;
    return (x*100+0.5);
}
int main()
{
    int i, j;
    while (scanf( "%d %d" , &n, &k), n|k)
    {
         for (i=0; i< n; i++)
             scanf( "%d" , A+i);
         for (i=0; i< n; i++)
             scanf( "%d" , B+i);
         printf( "%d\n" , solve());
}
     43. 表达式求值
// HDU 2127 Polish notation
#include
          <cstdio>
          <string>
#include
#include
          <algorithm>
using namespace
                   std;
typedef __int64
                  int64;
const int MAX = 1100;
char exp[MAX];
int priority[MAX];
int len;
bool output;
int64 dfs( int spos, int epos) {
    int i;
    int64 op1,op2,ans;
    char opr;
    int minv, minp = -1;
    if (spos > epos || spos >= len || epos < 0) {
```

W[j].w = 1.0*A[j] - x*B[j];

```
return 0;
}
for (i=epos;i>=spos;i--) {
    if (priority[i] != 0) {
         if (minp == -1) {
             minp = i;
             minv = priority[i];
         else if (minv > priority[i])
             minv = priority[i];
             minp = i;
    }
}
ans = 0;
if (minp == -1) {
    for (i=spos;i<=epos;i++) {
         ans = ans * 10 + \exp[i] - '0';
    }
    if (output) putchar( '');
    printf( "%d" ,ans);
    output = true ;
else {
    opr = exp[minp];
    if (opr != '(' && opr != ')' ) {
         if (output) putchar( ' ');
         putchar(opr);
         output = true ;
         op1 = dfs(spos,minp-1);
         op2 = dfs(minp+1,epos);
         switch (opr) {
         case '+' :
             ans = op1 + op2;
             break ;
         case '-' :
             ans = op1 - op2;
             break ;
         case '*' :
             ans = op1 * op2;
             break ;
         }
    }
```

else {

```
ans = dfs(spos,minp-1) + dfs(minp+1,epos);
        }
    }
    return ans;
}
int main() {
    int i,pre,t = 1;
    while (gets(exp)) {
        len = strlen(exp);
        pre = 0;
        // +,- 1
        // *,/ 2
        // - 3
        // () 4
        for (i=0;i<len;i++) {
             if (exp[i] == '*' || exp[i] == '/') priority[i] = pre + 2;
             else if (exp[i]
                           == '+' ) priority[i] = pre + 1;
                           == '(' ) priority[i] = pre = pre + 4;
             else if (exp[i]
             else if (exp[i]
                           == ')' ) priority[i] = pre, pre -= 4;
             else if (exp[i] == '-') {
                 if (exp[i-1]>= 'a' \&\& exp[i-1]<= 'z' || exp[i-1]== ')') priority[i]
pre + 1;
                 else priority[i] = pre + 3;
             else priority[i] = 0;
        printf( "Case %d:\n" , t ++);
         output = false;
        printf( \n\n%164d\n" , dfs(0,len-1));
}
    44. 乘除法博弈 ,Wythoff
                                        博弈
/*
PKU 2633 Funny Games
给定 f[1..n] 和 x
两人轮流选择一个 f[i], 使得 x=x*f[i]
当x<=1.0 时胜利
普通的博弈搜索难于 x过大,又是浮点数
*/
typedef pair < double , double > pdd;
typedef pair < double , double > * ppdd;
```

```
EQ(a,b)
                   (fabs((a)-(b))
                                    <= 1e-8)
#define
         LES(a,b)
#define
                    ((a) < (b))
#define
         LEQ(a,b)
                    ((a)+1e-8)
                                 <= (b))
pdd win[10000];
                              begin,
ppdd interval_union(ppdd
                                      ppdd end, ppdd dest)
    sort(begin,
                  end);
    for (; begin != end; dest++)
         *dest = *begin;
                          begin != end; begin++)
         for (begin++;
              if (LEQ(dest->second,
                                        begin->first))
                   break ;
                              = max(dest->second,
              dest->second
                                                         begin->second);
         }
    return dest;
}
char * solve2()
    int i, j;
    pdd lose;
    double maxf;
    lose.second
                 = \max f = 0;
    for (i=0; i< k; i++)
         win[i] = make_pair(1,
                                    1.0/f[i]);
         maxf = max(maxf,
    }
    int nwin;
    for (i=0,nwin=k;
                       LES(lose.second,
                                             x); i++)
    {
         // union the interval
         nwin = interval_union(win+i,
                                            win+nwin,
                                                         win+i) - win;
         // lose <- win
         lose.first = win[i].second;
                       = win[i].second
                                          / maxf; // it's min, and must
         lose.second
         // if the win have many interval
         if (i < nwin-1)
                          = min(lose.second,
                                                   win[i+1].first);
              lose.second
```

// win <- lose

```
for (j=0; j< k; j++)
            win[nwin ++] = make_pair(lose.first/f[j], lose.second/f[j]);
    }
    if (LES(x, lose.first)) return "Nils";
    return "Mikael" ;
}
// 有 2 堆石子,一次可以取任意个在一堆中或者任意个在两堆中取相同数目,取完者胜利
int swap( int &x, int &y)
{
    int t;
    t=x;
    x=y;
    y=t;
}
int main()
{
    double alpha = (1.0 + sqrt(5.0)) / 2.0;
    double beta = (3.0 + sqrt(5.0)) / 2.0;
    int big, small, n, temp1, temp2;
    while (cin>>big>>small)
        if (big < small)
            swap(big, small);
        n = ceil(big / beta);
        temp1 = alpha * n;
        temp2 = beta * n;
        if (small == temp1 && big == temp2)
            cout<<0<<endl;
        else cout<<1<<endl;
    return 0;
}
    45. 状态压缩的积木型 DP
// 1038 PKU Bugs Integrated,
                               Inc.
// DP with state compression
       MAX(a,b) ((a)>(b)?(a):(b))
#define
int bad[160];
int n, m, k;
int e3[] = {
    1, 3, 9, 27, 81, 243, 729, 2187, 6561, 19683, 59049,
                                                           177147
};
short dp[2][60000];
```

```
-- X(N)
Y(M)
0: [ ][ ]
1: [#][ ]
2: [#][#]
当维数扩展时 ,可以相应的扩展进制数
此题积木可分解为 1*3 影响列,则记为3进制数.
当M 小时, 也可用进制直接保存列 ...
基于状态压缩的积木填充型 DP , 有种实现方法 :
1) 直接构造 fm 和 gn 状态,如此题.
2)for(fm) 再构造 gn 状态.
3) 预处理,保存 <fm,gn> 为边.
优劣:
1)2) 易实现,但状态量大时会比较耗时
3) 保存边信息后 ,状态转移的时间消耗少
但边分布不均 ,需要动态数组或链表 ,且要处理去重操作
*/
void dfs( int x, int y, int fm, int gn, short v)
{
   if (y > m) return;
   if (y == m)
      dp[(x&1)^1][gn] = MAX(v + dp[(x&1)][fm], dp[(x&1)^1][gn]);
      return ;
   }
   int mask23 = 7 \ll y;
   int mask32 = 3 \ll y;
   if (x+1 < n)
      dfs(x, y+3, fm, gn+13*e3[y], v+1); // 2*3
      if (x+2 < n)
      {
         && !(bad[x+2]&mask32))
             dfs(x, y+2, fm, gn+8*e3[y], v+1); // 3*2
      }
   dfs(x, y+1, fm, gn, v); // 0->0
   if (!(bad[x]&(1<< y)))
```

```
{
         dfs(x, y+1, fm+1*e3[y], gn, v); // 1->0
         if (!(bad[x+1]&(1<< y)))
             dfs(x, y+1, fm+2*e3[y], gn+1*e3[y],
                                                  v); // 2->1
    }
}
int solve()
{
    int i, j;
    memset(dp, 0, sizeof (dp));
    for (i=0; i< n; i++)
         dfs(i, 0, 0, 0, 0);
         memset(dp[i&1], 0, sizeof (dp[0]));
    }
    return dp[i&1][0];
}
int main()
{
    int i, j;
    int cas;
    for (scanf( "%d" , &cas); cas; cas--)
         scanf( "%d %d %d" , &n, &m, &k);
         memset(bad, 0, sizeof (bad));
         for (i=0; i< k; i++)
             int x, y;
             scanf( "%d %d" , &x, &y);
             x --; y --;
             bad[x] = 1 \ll y;
         }
         printf( "%d\n" , solve());
}
     46. 解一般线性方程组 (消元法)
typedef int INT;
INT gcd(INT a, INT b) {
    return (b == 0)?a:gcd(b, a % b);
}
struct Fraction {
```

```
INT up, down;
    Fraction():up(0), down(1) {};
    Fraction(INT a, INT b = 1):up(a),
                                        down(b) {};
    Fraction( const Fraction& a) {
        up = a.up;
        down = a.down;
    }
    Fraction operator - () const {
         return Fraction(-up, down);
    }
    Fraction& operator = (const Fraction& a) {
         up = a.up;
         down = a.down;
        return * this ;
    }
    void reduce() {
        INT g = gcd(abs(up), abs(down));
         up /= g; down /= g;
};
Fraction abs( const Fraction& a) {
    return Fraction(abs(a.up), abs(a.down));
}
Fraction operator + (Fraction a, Fraction b) {
    INT u1 = a.up * b.down + a.down * b.up;
    INT u2 = a.down * b.down;
    INT g = gcd(abs(u1), abs(u2));
    return Fraction(u1 / g, u2 / g);
Fraction operator - (Fraction a, Fraction b) {
    return a + (-b);
}
Fraction operator * (Fraction a, Fraction b) {
    INT u1 = a.up * b.up;
    INT u2 = a.down * b.down;
    INT g = gcd(abs(u1), abs(u2));
    return Fraction(u1 / g, u2 / g);
}
Fraction operator / (Fraction a, Fraction b) {
    INT u1 = a.up * b.down;
    INT u2 = a.down * b.up;
    if (u2 < 0) u1 = -u1, u2 = -u2;
    int g = gcd(abs(u1), abs(u2));
```

```
return Fraction(u1 / g, u2 / g);
}
bool operator > (const Fraction& a, const Fraction& b) {
    return (a - b).up > 0;
}
bool operator == (const Fraction& a, const Fraction& b) {
    return (a - b).up == 0;
}
bool operator != (const Fraction& a, const Fraction& b) {
    return !(a == b);
}
bool operator < (const Fraction& a, const Fraction& b) {
    return (b - a).up > 0;
}
ostream& operator << (ostream& out, const Fraction& a) {
    if (a.down == 1) out << a.up;
    else out << a.up << '/' << a.down;
    return out;
}
const int nSize = 101;
const int mSize = 26;
typedef Fraction fEquation[nSize];
typedef fEquation fMatrix[mSize];
解一般形式的线性方程组 (只输出其中一组解 )
negn 个方程, nvar 个变量
矩阵表示如下
B0 = A0,0 * X0 + A0,1 * X1 + ... + A0,n-1 * Xn-1
B1 = A1,0 * X0 + A1,1 * X1 + ... + A1,n-1 * Xn-1
Bm = Am,0 * X0 + Am,1 * X1 + ... + Am,n-1 * Xn-1
*/
struct EqnGauss {
    int neqn, nvar;
    fMatrix f;
    Fraction avail[nSize];
    EqnGauss():
                 neqn(0), nvar(0) {}
    EqnGauss( const EqnGauss& a): neqn(a.neqn),
                                                    nvar(a.nvar)
        memcpy(f, a.f, sizeof (fMatrix));
    }
    void build() {
        for (int i = 0; i < neqn; ++ i)
```

```
for (int j = 0; j <= nvar; ++ j) {
              int x;
              cin >> x;
              f[i][j] = Fraction(x);
         }
}
bool rebuild() {
     int cur = 0;
    for (int i = 1; i \le nvar; ++ i) {
         bool found = false ;
         for (int j = cur; j < neqn; ++ j)
              if (f[j][i] != 0) {
                  found = true;
                  fEquation tmp;
                  memcpy(tmp,
                                   f[cur], sizeof (fEquation));
                  memcpy(f[cur],
                                     f[j], sizeof (fEquation));
                  memcpy(f[j], tmp, sizeof (fEquation));
              }
              if (!found) continue ;
              f[cur][0] = f[cur][0] / f[cur][i];
              for (int j = i + 1; j \le nvar; ++ j)
                  f[cur][j] = f[cur][j] / f[cur][i];
              f[cur][i] = 1;
              for (int j = 0; j < neqn; ++ j)
                  if (j != cur && f[j][i] != 0) {
                       f[j][0] = f[j][0] - f[j][i] * f[cur][0];
                       for (int k = i + 1; k \le nvar; ++ k)
                            f[j][k] = f[j][k] - f[j][i] * f[cur][k];
                       f[j][i] = 0;
              ++ cur;
     return (cur != nvar);
}
void solve() {
    int cur = 0;
    for (int i = 1; i \le nvar; ++ i)
         if (f[cur][i] == 0) {
              INT ulcm = 1;
              for (int j = 0; j < cur; ++ j)
                                             f[j][i].down) * f[j][i].down;
                  ulcm = ulcm / gcd(ulcm,
              avail[i] = ulcm;
         } else ++ cur;
    cur = 0;
```

```
for (int i = 1; i \le nvar; ++ i)
             if (f[cur][i] == 1) {
                  avail[i] = f[cur][0];
                  for (int j = i + 1; j \le nvar; ++ j)
                      avail[i] = avail[i] - f[cur][j] * avail[j];
                  ++ cur;
             }
};
ostream& operator << (ostream& out, const EqnGauss& a) {
    for (int i = 1; i \le a.nvar; ++ i)
         out << a.avail[i] << endl;
    return out;
}
EqnGauss eqns;
int main() {
    while (true ) {
         cin >> eqns.nvar >> eqns.neqn;
         eqns.build();
         if (eqns.rebuild())
             eqns.solve();
             cout << eqns;
         }
         else
             cout << "no solution" << endl;
    return 0;
}
4 2
10 1 1 1 1
3 1 1 0 0
4 2
10 1 1 0 0
3 1 1 0 0
*/
    47. 块状链表
// 块状链表
#include <iostream>
```

```
#include <iostream>
#include <cmath>
#define MAX 2900
struct block { //Type of block
```

```
int nos; //Number
                          of elements
                                       in this block
     char em[MAX];
                      //elements
     block *be,*su; //previous&successor
}
*first;
struct cursor { //Type of cursor
    int n; //The postion of cursor(in
                                          one block)
     block *II; //The block the cursor
}
cur;
int tot,n; //tot:the
                     total number
                                     of elements,n:the
                                                           number of operations
inline void clean(block
                           *op) { //Clean a new block
     op->nos=0;
     memset(op->em,0,
                            sizeof op->em);
    op->be=NULL;
    op->su=NULL;
}
inline void Spilt(block
                          *a, int newsize) { //Break a block into two blocks,one's
                                                                                           size
equal to 'newsize', and
                          another's
                                      equal to 'a->nos-newsize'
     if (!newsize) return ;
    int tmp=a->nos;
     a->nos=newsize;
     block *tps=a->su;
     a->su= new (block);
     clean(a->su);
     a->su->be=a;
    if (tps) {
          a->su->su=tps;
          a->su->su->be=a->su;
     a->su->nos=tmp-newsize;
     block *tt=a->su;
    for (int i=newsize+1;i<=tmp;i++)tt->em[i-newsize]=a->em[i];
}
inline void Merge(block
                                            a & a->su
                            *a) { //Merge
     if (a->su==NULL)
                         return
     int tmp=a->nos;
    for (int i=1;i<=a->su->nos;i++)a->em[a->nos+i]=a->su->em[i];
     a->nos+=a->su->nos;
     block *oo=a->su;
     if (cur.ll==oo)
         cur.ll=cur.ll->be;
          cur.n+=tmp;
    }
```

```
a->su=a->su->su;
     delete (oo);
     if (a->su)a->su->be=a;
}
inline void Balance()
                         { //Make
                                     these blocks' size balance
     block *k=first;
     int kk=( int )sqrt(tot);
     for (;k!=NULL;)
          for (;k->nos
                         < kk/2 || k->nos
                                               > 2*kk;)
                                    { //the
               if (k->nos< kk/2)
                                             block is too small?
                    if (k->su)Merge(k);
                    else break;
               } else if (k->nos>kk*2)
                                            { //the block is too big?
                     Spilt(k,(k->nos)>>1);
                    if((cur.ll==k)&&(cur.n>k->nos))
                          cur.ll=k->su;
                          cur.n-=k->nos;
                    k=k->su;
          }
          k=k->su;
}
inline void Insert(block
                            *lk, int x, int k) { //Insert
                                                           text behind the cursor
     block *oo= new (block);
     clean(oo);
     block *gg=oo;
     int rr=k;
     tot+=rr;
     int bt=( int )sqrt(tot);
     for (int i=0;i< rr;i++)
          char gg;
          scanf( "%c" ,&gg);
                                                 "%c" ,&gg));
          for (;(gg>126)||(gg<32);scanf(
          oo->em[++oo->nos]=gg;
          if ((oo->nos>=bt)&&(i<rr-1))
                                               {
               oo->su= new (block);
               clean(oo->su);
               oo->su->be=oo;
               00=00->su;
          }
     }
     if (x) {
```

```
Spilt(lk,x);
     } else {
           if (!cur.ll->be)
                block *jj;
                jj=first;
                first= new (block);
                clean(first);
                first->su=jj;
                jj->be=first;
           }
           cur.ll=cur.ll->be;
           cur.n=cur.ll->nos;
           lk=lk->be;
     block *tmp=lk->su;
     lk->su=gg;
     oo->su=tmp;
     if (oo->su)oo->su->be=oo;
     lk->su->be=lk;
     cur.ll=lk->su;
     cur.n=0;
     Balance();
}
inline void Remove(block
                                 *lk, int x, int num) { //Delete
                                                                                                     the
                                                                        'num'
                                                                                elements
                                                                                             behind
cursor
     if (x) {
           Spilt(lk,x);
           lk=lk->su;
     } else {
           if (!cur.ll->be)
                block *jj;
                jj=first;
                first= new (block);
                clean(first);
                first->su=jj;
                jj->be=first;
           }
           cur.ll=cur.ll->be;
           cur.n=cur.ll->nos;
     }
     tot-=num;
     int ttt=num;
     block *tmp;
     block *ii;
```

```
for (tmp=lk;tmp&&((num-tmp->nos)>=0);tmp=ii)
                                                                        {
           ii=tmp->su;
           num-=tmp->nos;
           delete (tmp);
     if (num&&tmp)
           Spilt(tmp,num);
           cur.ll->su=tmp->su;
           cur.ll->su->be=cur.ll;
           if (cur.ll->be)cur.ll->be->su=cur.ll;
     } else {
           cur.ll->su=tmp;
           if (cur.ll->be)cur.ll->be->su=cur.ll;
           if (cur.ll->su)cur.ll->su->be=cur.ll;
     Balance();
}
inline void Print(block
                             *lk, int x, int num) { //print
                                                                                            the cursor
                                                               'num'
                                                                        elements
                                                                                    behind
     block *cp=lk;
     for (;num-(cp->nos-x)>0;cp=cp->su)
                                                       {
                                                            "%c" ,cp->em[i]);
           for (int i=x+1;i<=cp->nos;i++)printf(
           num-=(cp->nos-x);
           x=0;
     for ( int i=1;i \le num;i++)printf(
                                                    ,cp->em[i+x]);
     printf( "\n" );
}
inline void prev()
                       { //cursor
                                    move forward
     cur.n--;
     if (cur.n<0)
           cur.ll=cur.ll->be;
           cur.n=cur.ll->nos-1;
}
inline void next()
                      { //cursor
                                    move backward
     cur.n++;
     if ((cur.n>=cur.ll->nos)&&(cur.ll->su))
           cur.ll=cur.ll->su;
           cur.n-=cur.ll->be->nos;
     }
}
inline void move( int k) { //move
                                          the cursor to the postion
                                                                          'k'
     cur.ll=first;
     for (;(cur.ll)&&(k-cur.ll->nos>0);cur.ll=cur.ll->su)
```

```
k-=cur.ll->nos;
    cur.n=k;
}
         editor O(sqrt(n))
NOI2003
                         将光标移动到第 k个字符之后,如果 k=0 ,将光标移到文本开头
MOVE(k)
            Move k
INSERT(n,
                             在光标处插入长度为 n 的字符串 s , 光标位置不变 ,
                 Insert n S
                         删除光标后的 n个字符,光标位置不变,
DELETE(n)
            Delete n
                                                             n >= 1
                         输出光标后的 n个字符,光标位置不变, n >= 1
GET(n)
            Get n
                         光标前移一个字符
PREV()
            Prev
                         光标后移一个字符
NEXT()
            Next
*/
int main() {
    freopen( "editor.in"
                       ,"r" ,stdin);
                        ,"w" ,stdout);
    freopen( "editor.out"
    first= new (block);
    clean(first);
    char str[100];
    cur.ll=first;
    cur.n=0;
    scanf( "%d\n" , &n);
    int k;
    for (int i=0;i< n;i++)
        scanf( "%s" , str);
        switch (str[0])
                      { //deal with operations
        case 'M' : scanf( "%d" , &k); move(k);
                                               break ;
                                    Insert(cur.ll,cur.n,k);
        case 'I' : scanf( "%d" , &k);
                                                          break ;
        case 'D' : scanf( "%d" , &k);
                                    Remove(cur.II,cur.n,k);
                                                             break ;
        case 'G' : scanf( "%d" , &k);
                                    Print(cur.II,cur.n,k);;
                                                           break ;
        case 'P' : prev();
                          break ;
        case 'N' : next();
                          break ;
    fclose(stdin);
    fclose(stdout);
}
    48. Factor
                    Oracle
Factor Oracle
后缀自动机构造过于复杂
可利用 Factor Oracle 实现基于子串的搜索
实现两个串公共最长子串 / 单串最长重复子串的 O(n) 算法
(1) 能在 O(|u|) 识别 p 的子串 u
```

```
(2) 可以识别 p的所有子串 ,可能会误识别长度小于 |p| 的子串!!
(3) 在O(|p|) 时间内构造
*/
/*
abaaabbabaa
abab
Find, it 'vsrong. (2)
*/
const int INIT = 1;
const int FAIL = 0;
const int MMAX = 201000;
const int SIGMAX = 30;
2774 PKU Long Long Message
求出 Irs 和S以后,问题就好解决了
对于公共子串,扫描后半部分的 Irs,加上S的限制,防止重复串在同一个串中
若是单串,直接取 Irs 中的最大值
*/
char str[MMAX];
struct ORACLE
    int T[MMAX][SIGMAX];
    int S[MMAX];
    int LRS[MMAX];
    int SN;
    int c2d[256];
    void init() {
        SN = INIT + 1;
        S[INIT] = FAIL;
        LRS[INIT]
                  = 0;
                   FAIL, sizeof(T)); // 节省清空的时间开销
//
        memset(T,
//
        memset(c2d, FAIL, sizeof(c2d));
    }
    void add( char ch) {
        int m = SN - 1;
        SN ++;
        ch = c2d[ch];
        T[m][ch] = m + 1;
        int k = S[m];
        int pre = m;
        while (k \vdash FAIL \&\& T[k][ch] == FAIL) {
           T[k][ch] = m + 1;
           pre = k;
```

```
k = S[k];
         }
         int shift;
         if (k == FAIL) shift = INIT;
         else shift = T[k][ch];
         S[m+1] = shift;
         LRS[m+1] = len_repeat_suffix(pre,
                                            S[m+1]);
    void construct( char * p) {
         for (int i=0,j=FAIL+1; p[i]; i++) {
             if (c2d[p[i]] = FAIL)
                  c2d[p[i]] = j ++;
             add(p[i]);
         }
    }
    int len_common_suffix( int p1, int p2) {
         if (S[p1] == p2) return LRS[p1];
         while (S[p1] != S[p2]) p2 = S[p2];
         return min(LRS[p1], LRS[p2]);
    }
    int len_repeat_suffix( int p1, int m) {
         if (m == INIT) return 0;
         return len_common_suffix(p1, m-1) + 1;
    }
};
ORACLE fo;
int main()
{
    gets(str);
    int len = strlen(str);
    str[len] = '\$';
    gets(str+len+1);
    fo.init();
    fo.construct(str);
    int ans = 0;
    for (int i=len+INIT+2;
                              i<fo.SN;
                                         i++)
         if (fo.S[i] \rightarrow = len+INIT+2)
                                        continue ;
         ans = max(ans,
                           fo.LRS[i]);
    printf( "%d\n" , ans);
}
```

49.