## Running Time of Programs

Exercise 1. Suppose you have the choice between three algorithms:

- (a) Algorithm A solves your problem by dividing it into five subproblems of half the size, recursively solving each subproblem, and then combining the solutions in linear time.
- (b) Algorithm B solves problems of size n by recursively solving two subproblems of size n-1 and then combining the solutions in constant time.
- (c) Algorithm C solves problems of size n by dividing them into nine subproblems of size  $\frac{n}{3}$ , recursively solving each subproblem, and then combining the solutions in  $\mathcal{O}(n^2)$  time.

Estimate the running times of each of these algorithms. Which one would you choose?

**Exercise 2.** Recall the recurrence for Mergesort: T(1) = 0;  $T(n) = 2T(\frac{n}{2}) + (n-1)$ , for n > 1. Prove by induction that

$$T(n) = n \cdot (\log_2 n - 1) + 1$$
 for  $n = 2^k$ ,  $k > 1$ 

**Exercise 3.** Analyse the complexity of the following recursive algorithm to test whether a number x occurs in an *unordered* list  $L = [x_1, x_2, \ldots, x_n]$  of size n. Take the cost to be the number of list element comparison operations.

**Search** $(x, L = [x_1, x_2, ..., x_n])$ :

- if  $x_1 = x$  then return yes
- else if n > 1 then return **Search** $(x, [x_2, ..., x_n])$
- else return no

\*Exercise 4. Analyse the complexity of the following recursive algorithm to test whether a number x occurs in an *ordered* list  $L = [x_1, x_2, \ldots, x_n]$  of size n. Take the cost to be the number of list element comparison operations.

BinarySearch $(x, L = [x_1, x_2, \dots, x_n])$ :

- if n = 0 then return no
- else
  - if  $x_{\lceil \frac{n}{2} \rceil} > x$  then return  $\mathbf{BinarySearch}(x, [x_1, \dots, x_{\lceil \frac{n}{2} \rceil 1}])$
  - else if  $x_{\lceil \frac{n}{2} \rceil} < x$  return **BinarySearch** $(x, [x_{\lceil \frac{n}{2} \rceil+1}, \dots, x_n])$
  - else return yes