

Pre-Announcements

Justin: Alpha Kappa Psi Blockchain Event, next Wednesday in Chevron Auditorium

- Companies like Oracle, Ripple, Coinbase, etc.
- An overview of blockchains and its implications.
- 7 PM next Wednesday in iHouse.
 - Pizza, ice cream, and a raffle

Fliers up front if you want to do this thing.

Announcements

Midterm 2:

- 3/20, 8-10 PM in various rooms.
- Covers material through 3/16 (this Friday).
- Study using study guides.
 - THE KEY IS METACOGNITION: Reflect on your problem solving strategies and those of your fellow students.
 - Understanding a handful of solutions to old midterm problems is less helpful than you might think -- look at answers as late as possible.
- There is an alternate 61C midterm from 6 - 8 in 1 LeConte.

CS61B

Lecture 24: Priority Queues and Heaps

- Priority Queues
- Heaps
- Tree Representations
- Data Structures Summary



The Priority Queue Interface

```
/** (Min) Priority Queue: Allowing tracking and removal of the  
 * smallest item in a priority queue. */  
public interface MinPQ<Item> {  
    /** Adds the item to the priority queue. */  
    public void add(Item x);  
    /** Returns the smallest item in the priority queue. */  
    public Item getSmallest();  
    /** Removes the smallest item from the priority queue. */  
    public Item removeSmallest();  
    /** Returns the size of the priority queue. */  
    public int size();  
}
```

Useful if you want to keep track of the “smallest”, “largest”, “best” etc. seen so far.

Usage Example: Unharmonious Texts

Imagine that you're part of Grand Leader's Information Compliance and Happiness Enhancement (GLICHE) team.

- Your job: Monitor the text messages of the citizens to make sure that they are not having any unharmonious conversations.
- Each day, you prepare a report of the M messages that seem most unharmonious using the HarmoniousnessComparator.

Naive approach: Create a list of all messages sent for the entire day. Sort it using your comparator. Return the M messages that are largest.



Naive Implementation: Store and Sort

```
public List<String> unharmoniousTexts(Sniffer sniffer, int M) {  
    ArrayList<String> allMessages = new ArrayList<String>();  
  
    for (Timer timer = new Timer(); timer.hours() < 24; ) {  
        allMessages.add(sniffer.getNextMessage());  
    }  
  
    Comparator<String> cmptr = new HarmoniousnessComparator();  
    Collections.sort(allMessages, cmptr, Collections.reverseOrder());  
  
    return allMessages.sublist(0, M);  
}
```

Potentially uses a huge amount of memory $\Theta(N)$, where N is number of texts.

- Goal: Do this in $\Theta(M)$ memory using a MinPQ.

```
MinPQ<String> unharmoniousTexts = new HeapMinPQ<Transaction>(cmptr);
```

Better Implementation: Track the M Best

```
public List<String> unharmoniousTexts(Sniffer sniffer, int M) {
    Comparator<String> cmptr = new HarmoniousnessComparator();
    MinPQ<String> unharmoniousTexts = new HeapMinPQ<Transaction>(cmptr);
    for (Timer timer = new Timer(); timer.hours() < 24; ) {
        unharmoniousTexts.add(sniffer.getNextMessage());
        if (unharmoniousTexts.size() > M)
            { unharmoniousTexts.removeSmallest(); }
    }
    ArrayList<String> textlist = new ArrayList<String>();
    while (unharmoniousTexts.size() > 0) {
        textlist.add(unharmoniousTexts.removeSmallest());
    }
    return textlist;
}
```

Can track top M transactions using only M memory. API for MinPQ also makes code very simple (don't need to do explicit comparisons).

How Would We Implement a MinPQ?

Some possibilities:

- Ordered Array
- Bushy BST: Maintaining bushiness is annoying. **Handling duplicate priorities is awkward.**
- HashTable: No good! Items go into random places.

	Ordered Array	Bushy BST	Hash Table	Heap
add	$\Theta(N)$	$\Theta(\log N)$	$\Theta(1)$	
getSmallest	$\Theta(1)$	$\Theta(\log N)$	$\Theta(N)$	
removeSmallest	$\Theta(N)$	$\Theta(\log N)$	$\Theta(N)$	
Caveats		Dups tough		

Worst Case $\Theta(\cdot)$ Runtimes

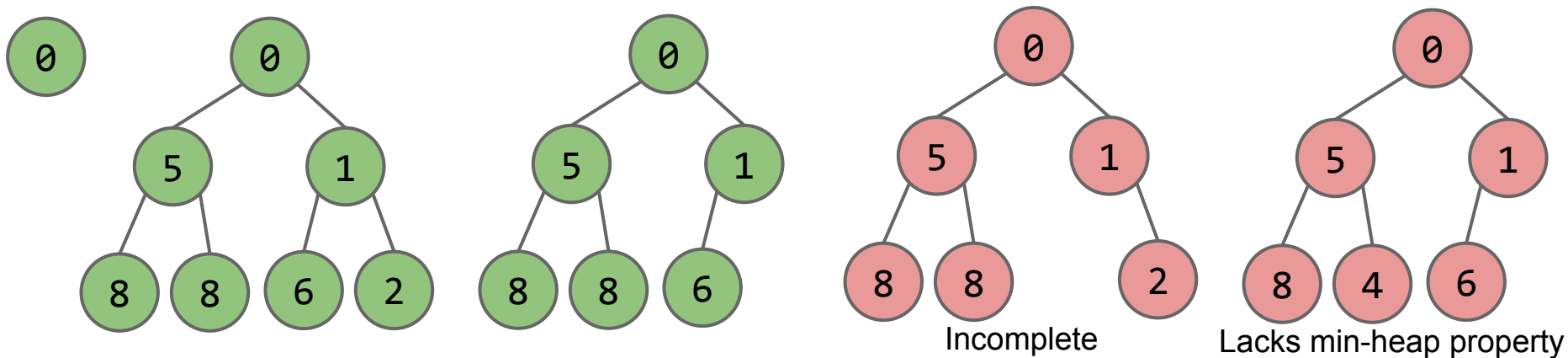
Heaps

Introducing the Heap

BSTs would work, but need to be kept bushy and duplicates are awkward.

Binary min-heap: Binary tree that is **complete** and obeys **min-heap property**.

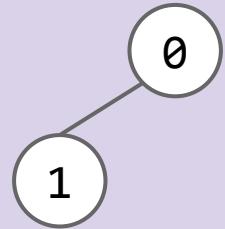
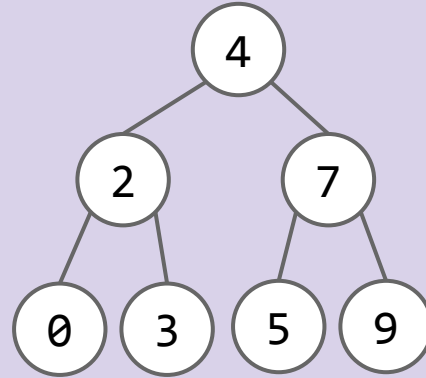
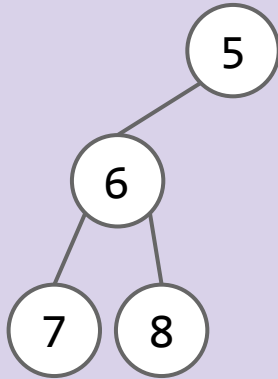
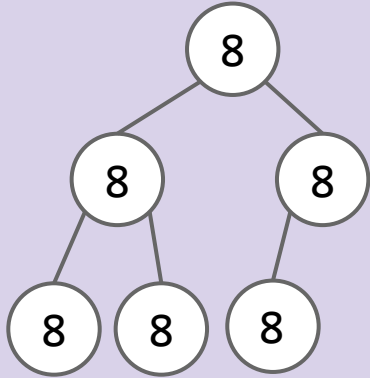
- Min-heap: Every node is less than or equal to both of its children.
- Complete: Missing items only at the bottom level (if any), all nodes are as far left as possible.



Heap Comprehension Test: <http://yellkey.com/present>

How many of these are min heaps?

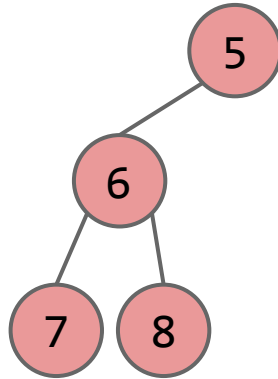
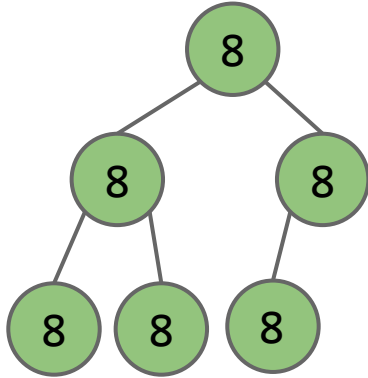
- A. 0
- B. 1
- C. 2
- D. 3
- E. 4



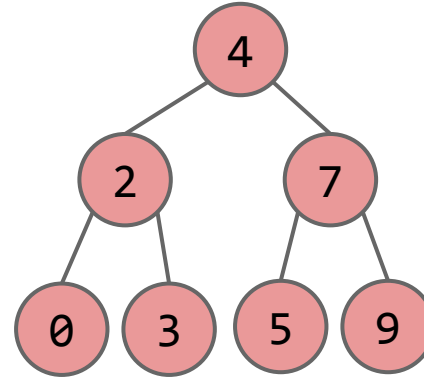
Heap Comprehension Test: <http://yellkey.com/present>

How many of these are min heaps?

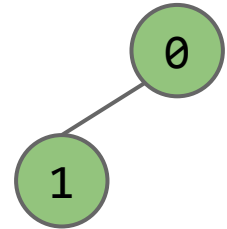
- A. 0
- B. 1
- C. 2
- D. 3
- E. 4



Incomplete



Lacks min-heap property

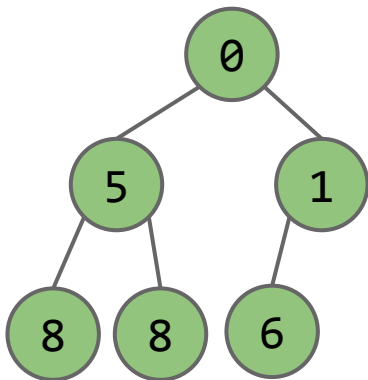


What Good Are Heaps?

Heaps lend themselves very naturally to implementation of a priority queue.

Hopefully easy question:

- How would you support `getSmallest()`?

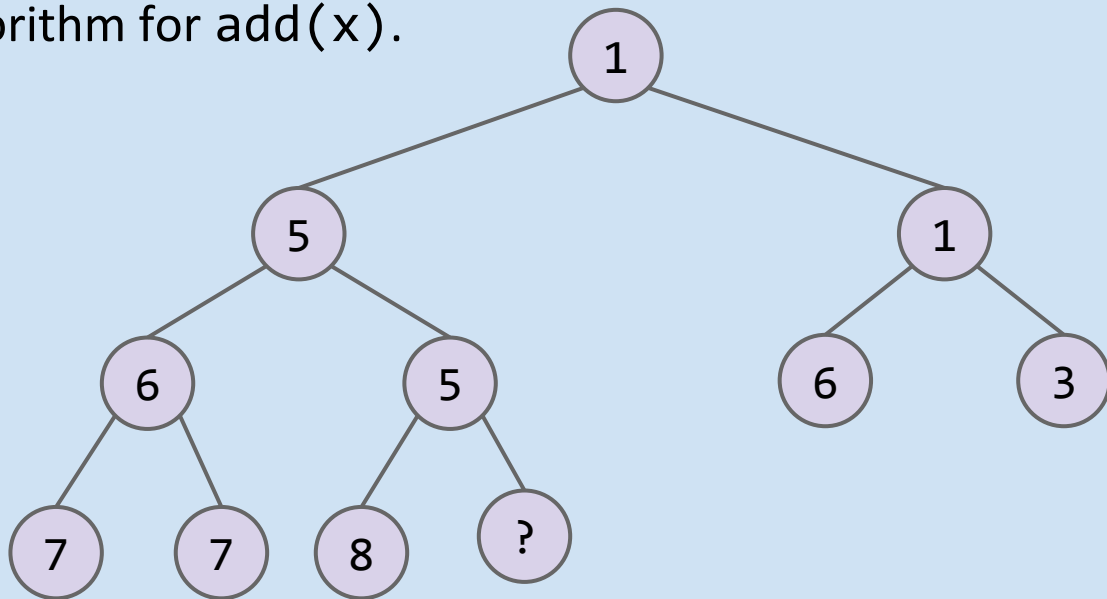


How Do We Add To A Heap?

Challenge: Come up with an algorithm for `add(x)`.

- How would we insert 3?

Runtime must be logarithmic.



Bonus: Come up with an algorithm for `removeSmallest()`.

Solution: See <https://goo.gl/wBKdFQ> for an animated demo.

Heap Operations Summary

Given a heap, how do we implement PQ operations?

- `getSmallest()` - return the item in the root node.
- `add(x)` - place the new employee in the last position, and promote as high as possible.
- `removeSmallest()` - assassinate the president (of the company), promote the rightmost person in the company to president. Then demote repeatedly, always taking the 'better' successor.

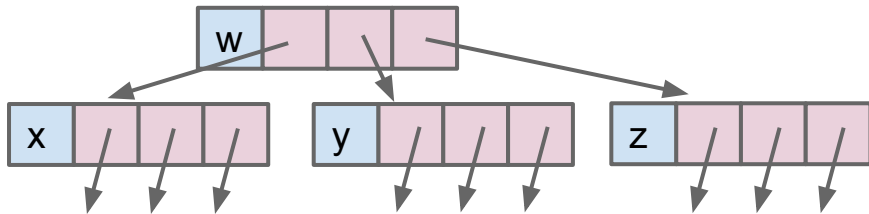
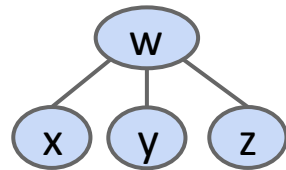
See <https://goo.gl/wBKdFQ> for an animated demo.

Remaining question: How would we do all this in Java?

Tree Representations

How do we Represent a Tree in Java?

Approach 1a, 1b and 1c: Create mapping from node to children.



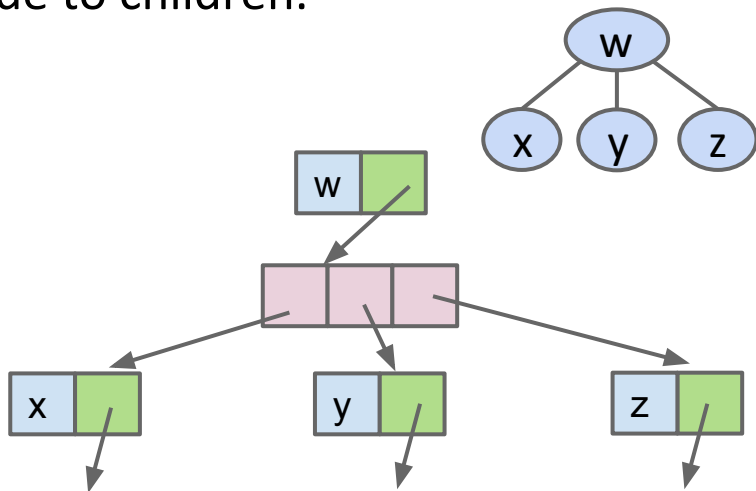
1a: Fixed-Width Nodes (BSTMap used this approach)

```
public class Tree1A<Key> {  
    Key k; // e.g. 0  
    Tree1A left;  
    Tree1A middle;  
    Tree1A right;  
    ...  
}
```

How do we Represent a Tree in Java?

Approach 1a, 1b and 1c: Create mapping from node to children.

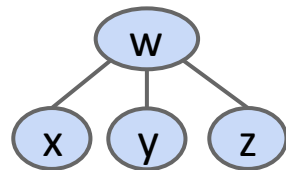
```
public class Tree1B<Key> {  
    Key k; // e.g.  $\emptyset$   
    Tree1B[] children;  
    ...  
}
```



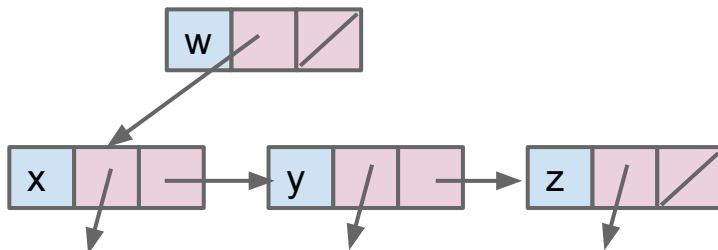
1b: Variable-Width Nodes

How do we Represent a Tree in Java?

Approach 1a, 1b and 1c: Create mapping from node to children.



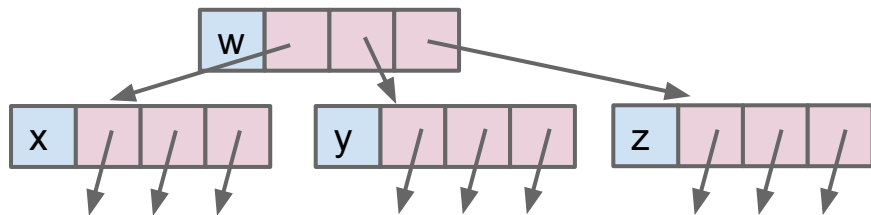
```
public class Tree1C<Key> {  
    Key k; // e.g. 0  
    Tree1C favoredChild;  
    Tree1C sibling;  
    ...  
}
```



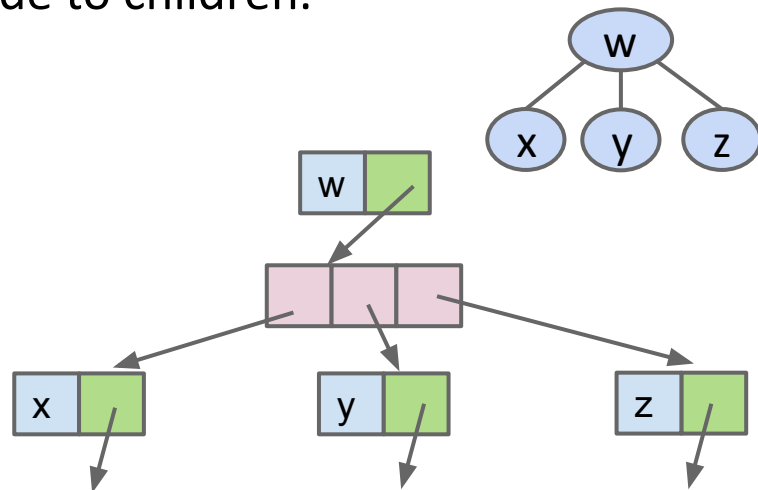
1c: Sibling Tree

How do we Represent a Tree in Java?

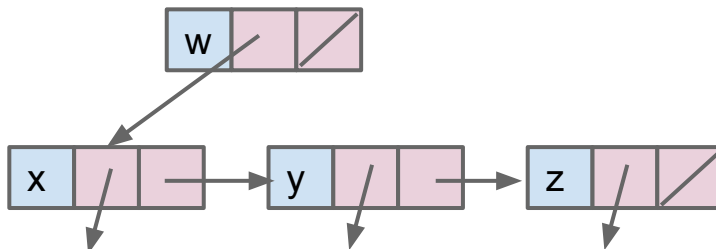
Approach 1a, 1b and 1c: Create mapping from node to children.



1a: Fixed-Width Nodes



1b: Variable-Width Nodes



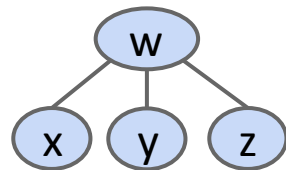
1c: Sibling Tree

How do we Represent a Tree in Java?

Approach 2: Store keys in an array. Store parentIDs in an array.

- Similar to what we did with disjointSets.

```
public class Tree2<Key> {  
    Key[] keys;  
    int[] parents;  
    ...  
}
```

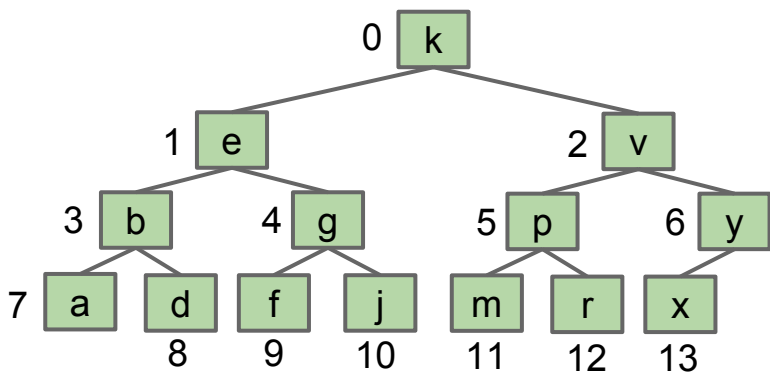


Key[] keys

w	x	y	z
---	---	---	---

int[] parents

0	0	0	0
0	1	2	3



Key[] keys

k	e	v	b	g	p	y	a	d	f	j	m	r	x
0	1	2	3	4	5	6	7	8	9	10	11	12	13

int[] parents

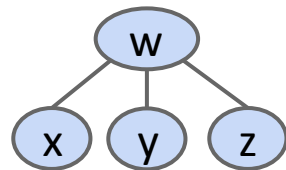
0	0	0	1	1	2	2	3	3	4	4	5	5	6
0	1	2	3	4	5	6	7	8	9	10	11	12	13

How do we Represent a Tree in Java?

Approach 3: Store keys in an array. Don't store structure anywhere.

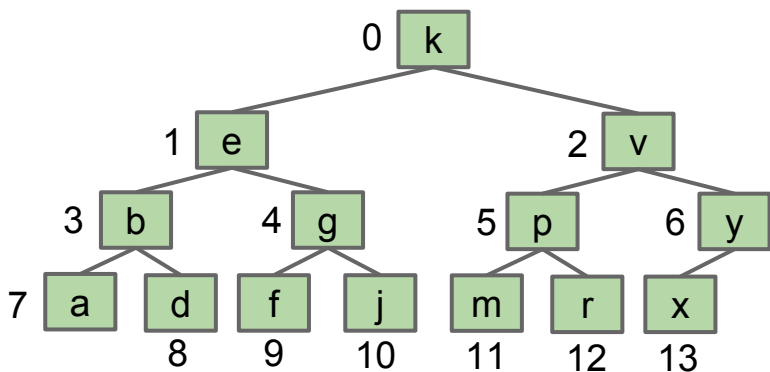
- To interpret array: Simply assume tree is complete.
- Obviously only works for “complete” trees.

```
public class Tree3<Key> {  
    Key[] keys;  
    ...  
}
```



Key[] keys

w	x	y	z
0	1	2	3



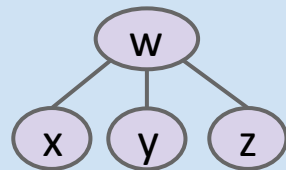
Key[] keys

k	e	v	b	g	p	y	a	d	f	j	m	r	x
0	1	2	3	4	5	6	7	8	9	10	11	12	13

A Deep Look at Approach 3

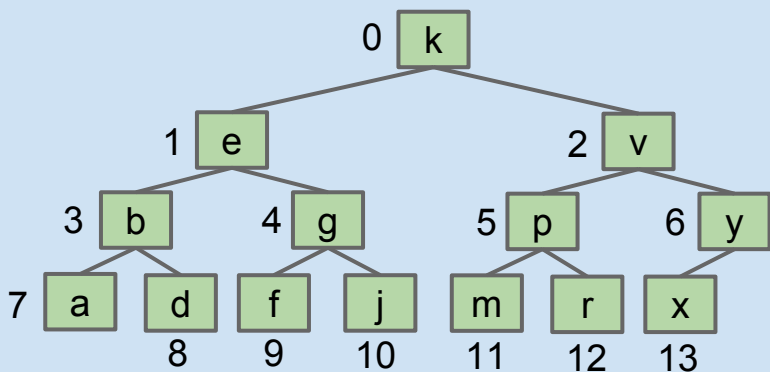
Challenge: Write the `parent(k)` method for approach 3.

```
public void swim(int k) {  
    if (keys[parent(k)] > keys[k]) {  
        swap(k, parent(k));  
        swim(parent(k));  
    }  
}
```



Key[] keys

w	x	y	z
0	1	2	3



Key[] keys

k	e	v	b	g	p	y	a	d	f	j	m	r	x
0	1	2	3	4	5	6	7	8	9	10	11	12	13

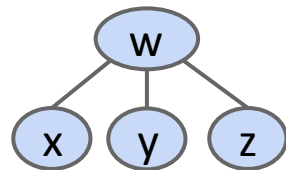
```
public class Tree3<Key> {  
    Key[] keys;  
    ...  
}
```

A Deep Look at Approach 3

Challenge: Write the `parent(k)` method for approach 3.

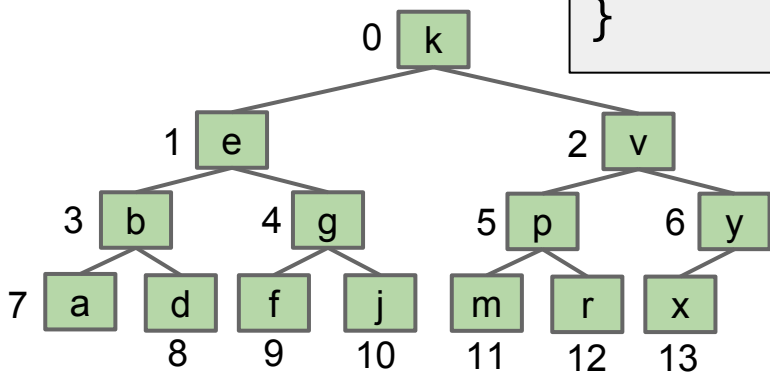
```
public void swim(int k) {  
    if (keys[parent(k)] > keys[k]) {  
        swap(k, parent(k));  
        swim(parent(k));  
    }  
}
```

```
public int parent(int k) {  
    return (k - 1) / 2;  
}
```



Key[] keys

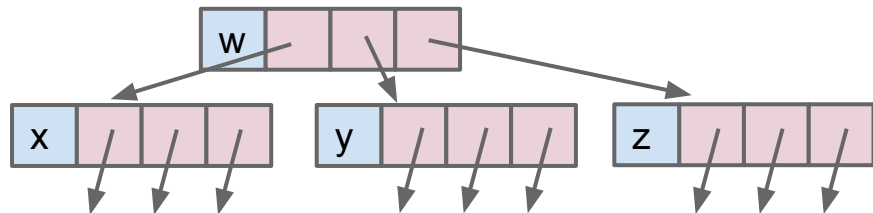
w	x	y	z
0	1	2	3



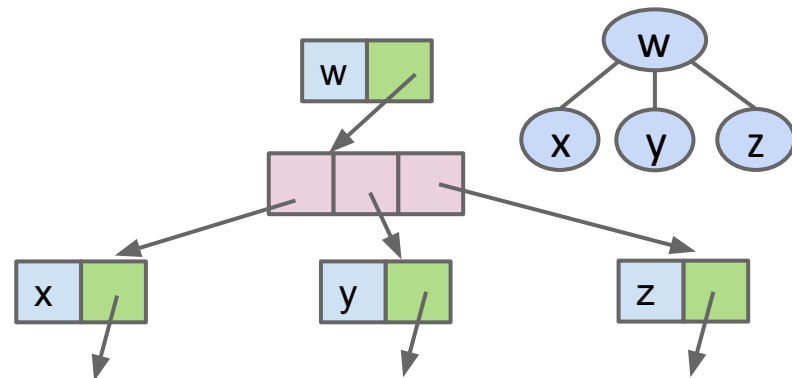
k	e	v	b	g	p	y	a	d	f	j	m	r	x
0	1	2	3	4	5	6	7	8	9	10	11	12	13

```
public class Tree3<Key> {  
    Key[] keys;  
    ...  
}
```

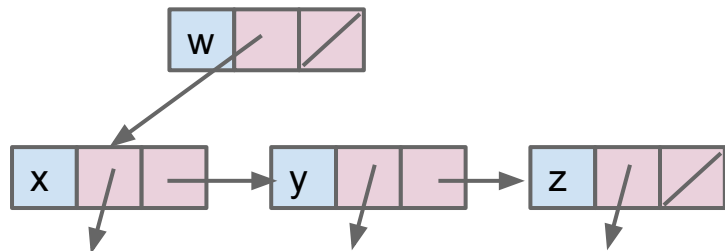

Tree Representations (Summary)



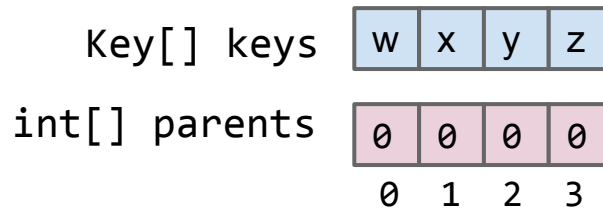
1a: Fixed Number of Links (One Per Child)



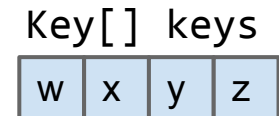
1b: Array of Child Links



1c: FirstBorn/Sibling Links



2: Array of Keys, Array of Structure

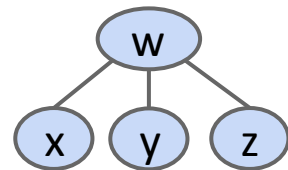


3: Array of Keys

Approach 3B (book implementation): Leaving One Empty Spot

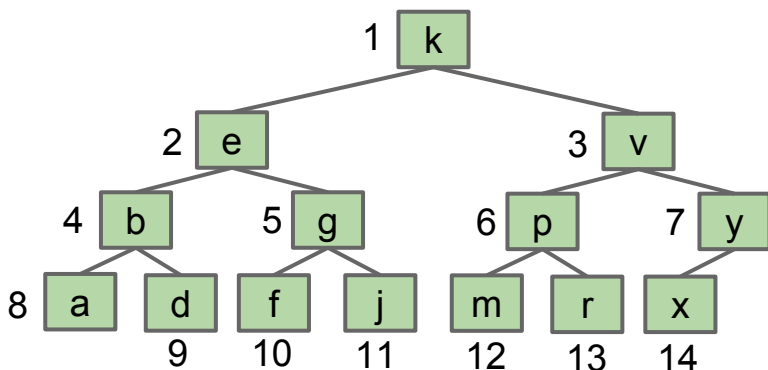
Approach 3b: Store keys in an array. Offset everything by 1 spot.

- Same as 3, but leave spot 0 empty.
- Makes computation of children/parents “nicer”.
 - $\text{leftChild}(k) = k * 2$
 - $\text{rightChild}(k) = k * 2 + 1$
 - $\text{parent}(k) = k / 2$



Key[] keys

-	w	x	y	z
0	1	2	3	4



Key[] keys

-	k	e	v	b	g	p	y	a	d	f	j	m	r	x
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

In next week's lab, you'll implement a MinPQ.

Heap Implementation of a Priority Queue

	Ordered Array	Bushy BST	Hash Table	Heap
add	$\Theta(N)$	$\Theta(\log N)$	$\Theta(1)$	$\Theta(\log N)$
getSmallest	$\Theta(1)$	$\Theta(\log N)$	$\Theta(N)$	$\Theta(1)$
removeSmallest	$\Theta(N)$	$\Theta(\log N)$	$\Theta(N)$	$\Theta(\log N)$

Notes:

Items with same priority hard to handle.

- Why “priority queue”? Can think of position in tree as its “priority.”
- Heap is $\log N$ time AMORTIZED (some resizes, but no big deal).
- BST can have constant getSmallest if you keep a pointer to smallest.
- Heaps handle duplicate priorities much more naturally than BSTs.
- Array based heaps take less memory (very roughly about 1/3rd the memory of representing a tree with approach 1a).

Some Implementation Questions

1. How does a PQ know how to determine which item in a PQ is larger?
 - a. What could we change so that there is a default comparison?
2. What constructors are needed to allow for different orderings?

```
/** (Min) Priority Queue: Allowing tracking and removal of the  
 * smallest item in a priority queue. */  
public interface MinPQ<Item> {  
    /** Adds the item to the priority queue. */  
    public void add(Item x);  
    /** Returns the smallest item in the priority queue. */  
    public Item getSmallest();  
    /** Removes the smallest item from the priority queue. */  
    public Item removeSmallest();  
    /** Returns the size of the priority queue. */  
    public int size();  
}
```

Data Structures Summary

The Search Problem

Given a stream of data, retrieve information of interest.

- Examples:
 - Website users post to personal page. Serve content only to friends.
 - Given logs for thousands of weather stations, display weather map for specified date and time.

My Friends: (98) [Edit Friends]

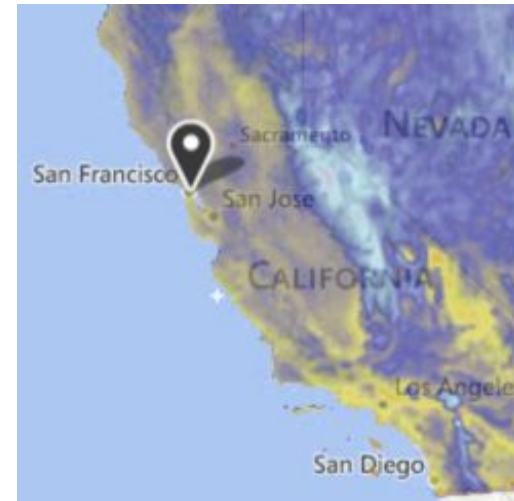
<u>Stephanie</u> 	<u>Hank</u> 	<u>Nikki</u> 	<u>Di</u> 	<u>Claire</u> 	<u>Johnny</u> 
<u>Kate</u> 	<u>Sujit</u> 	<u>Alain</u> 	<u>Dan</u> 	<u>Katharine</u> 	<u>Larry</u> 

[See My 98 Friends]

Waiting for confirmation from 3 friends
[Review/Cancel]

Recent Bulletin Board Posts from your Friends:

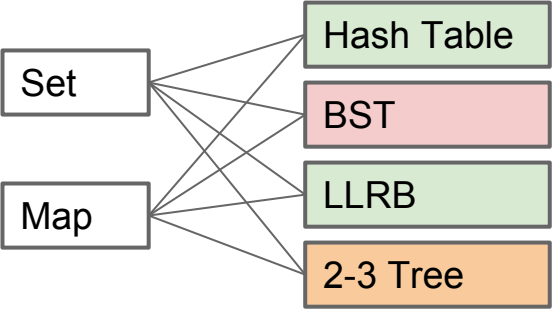
From	Date	Subject
<u>Jojo</u>	10/12/2004	<u>Adam Arcuragi wants you to make out with somebody!</u>
<u>Maria</u>	10/11/2004	<u>anyone looking for loft in greepoint?</u>
<u>Larry</u>	10/09/2004	<u>My new band's debut - 10/16</u>
<u>Jojo</u>	10/09/2004	<u>me and bitter, bitter weeks</u>



Search Data Structures (The particularly abstract ones)

Name	Storage Operation(s)	Primary Retrieval Operation	Retrieve By:
List	<code>add(key)</code> <code>insert(key, index)</code>	<code>get(index)</code>	index
Map	<code>put(key, value)</code>	<code>get(key)</code>	key identity
Set	<code>add(key)</code>	<code>containsKey(key)</code>	key identity
PQ	<code>add(key)</code>	<code>getSmallest()</code>	key order (a.k.a. key size)
Disjoint Sets	<code>connect(int1, int2)</code>	<code>isConnected(int1, int2)</code>	two int values

Searching Data Structures:



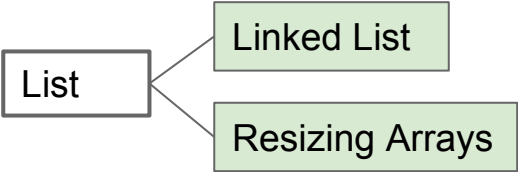
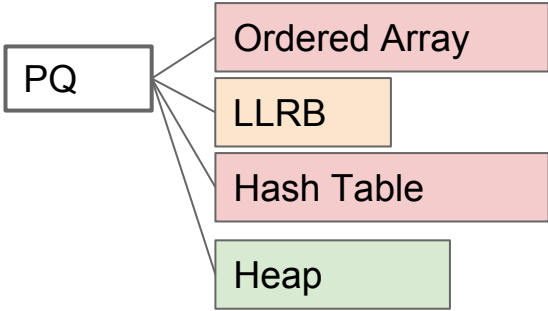
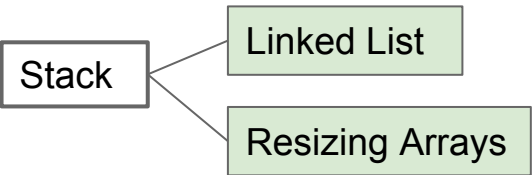
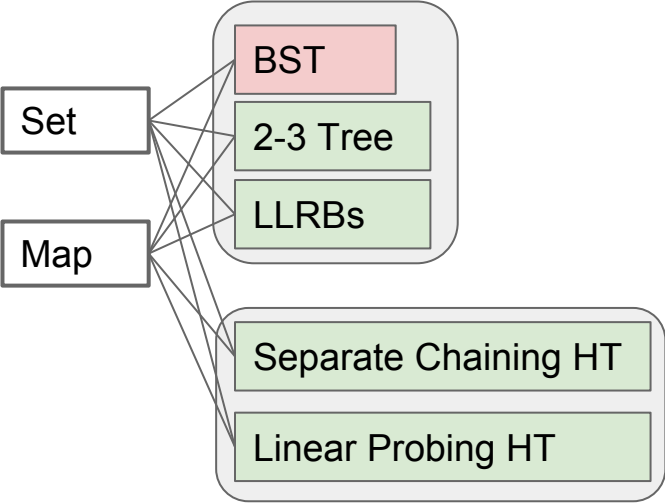
Stack

PQ

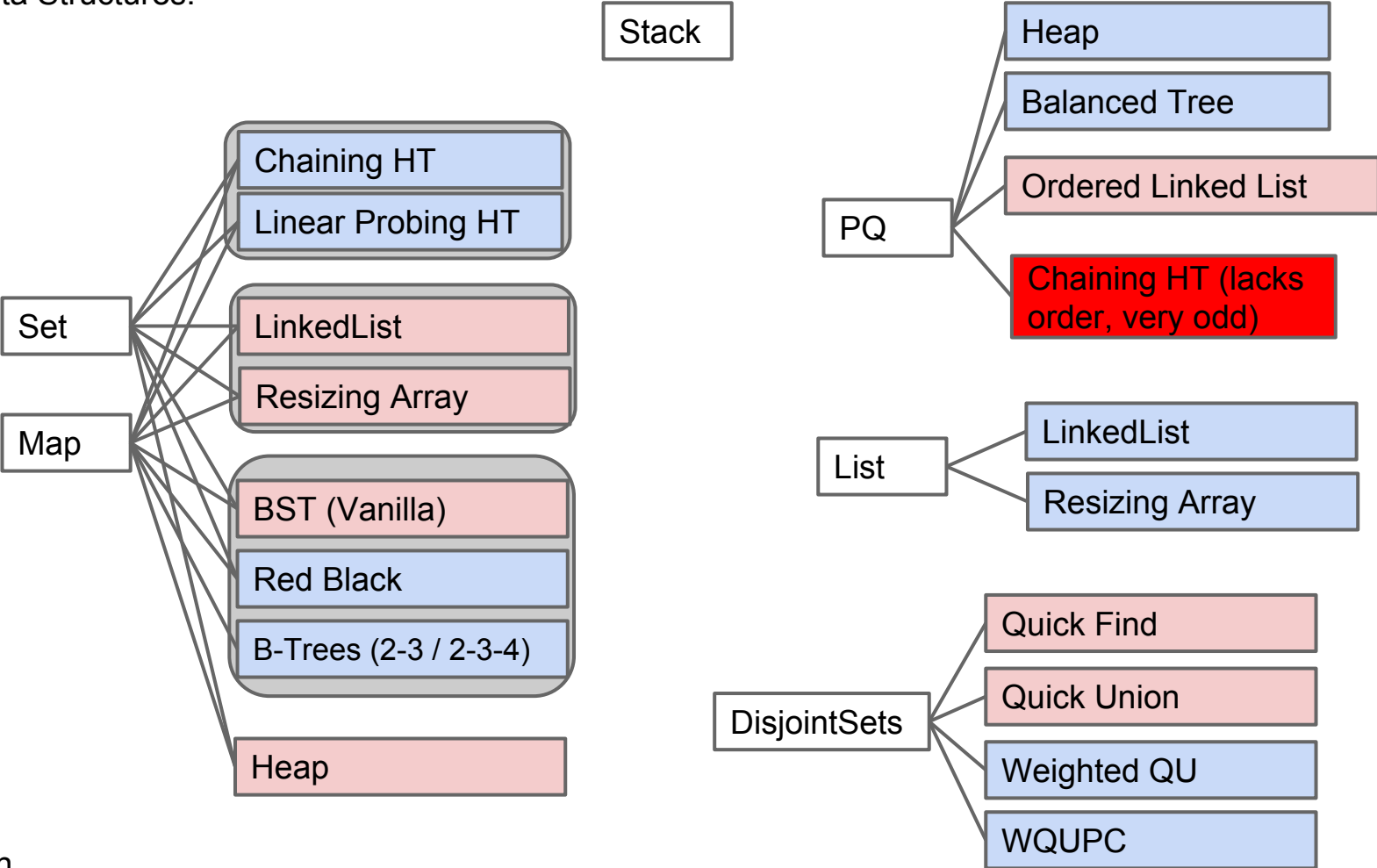
List

DisjointSets

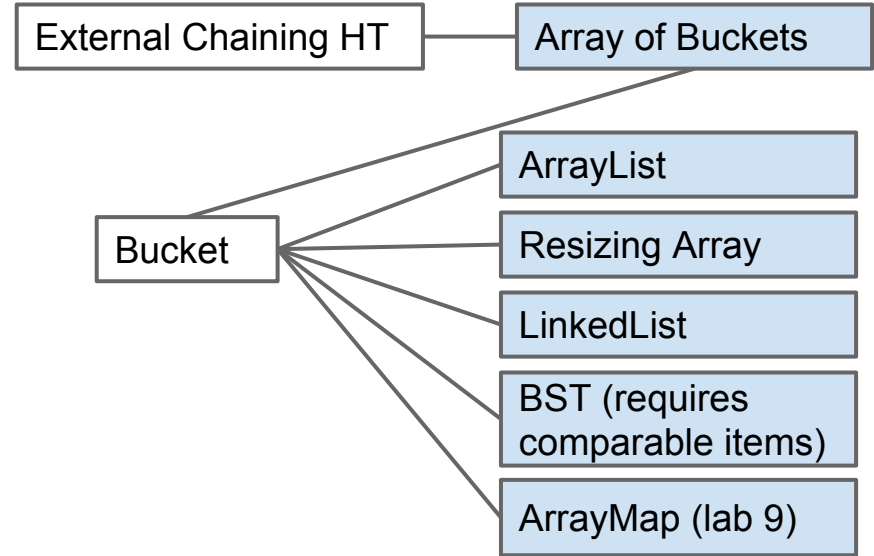
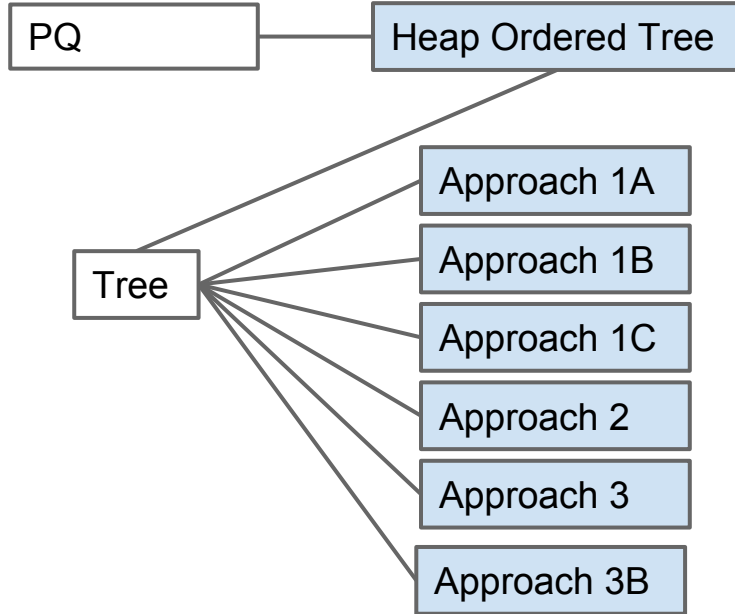
Searching Data Structures:



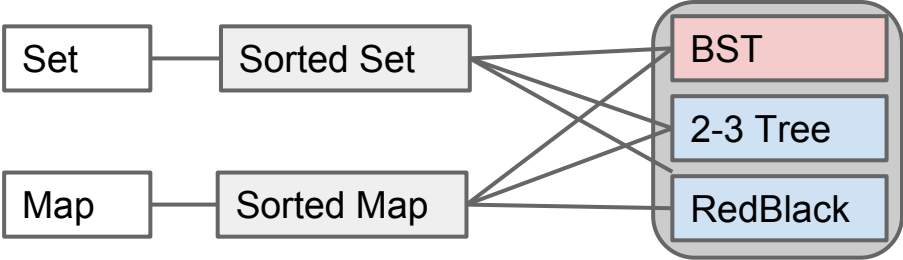
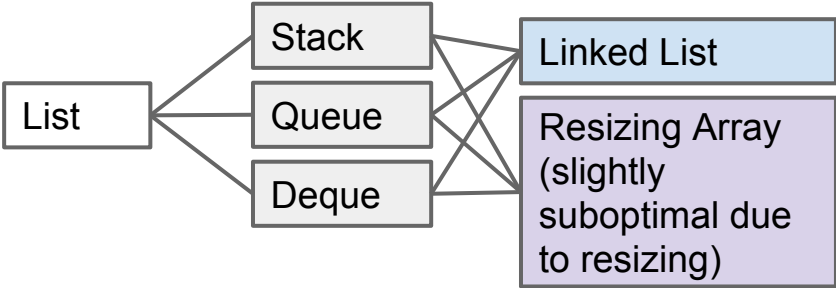
Searching Data Structures:



Abstraction often happens in layers!



Specialized Searching Data Structures:



In Java:
`java.util.SortedSet`

`java.util.SortedMap`

PQ

Don't usually consider MinPQ and MaxPQ to be different data structures, since we can just provide the opposite comparator.

Data Structures

Data Structure: A particular way of organizing data.

- We've covered many of the most fundamental abstract data types, their common implementations, and the tradeoffs thereof.
- We'll do two more in this class:
 - The quadtree (Friday).
 - The graph (starting after exam).

V•T•E	Data structures	[hide]
Types	Collection • Container	
Abstract	Associative array • Double-ended priority queue • Double-ended queue • List • Map • Multimap • Priority queue • Queue • Set (multiset) • Disjoint Sets • Stack	
Arrays	Bit array • Circular buffer • Dynamic array • Hash table • Hashed array tree • Sparse array	
Linked	Association list • Linked list • Skip list • Unrolled linked list • XOR linked list	
Trees	B-tree • Binary search tree (AA • AVL • red-black • self-balancing • splay) • Heap (binary • binomial • Fibonacci) • R-tree (R* • R+ • Hilbert) • Trie (Hash tree)	
Graphs	Binary decision diagram • Directed acyclic graph • Directed acyclic word graph	

Citations

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