

CX1107 Data Structures and Algorithms

2020/21 Semester 2

Tutorial 4: Algorithm Analysis
School of Computer Science and Engineering
Nanyang Technological University
Q1 Given the precedence of some operators, $bc * \rightarrow abc * d \% + e \Rightarrow$

| Operators | Precedence |
|------------|------------|
| $*, /, \%$ | highest |
| $+, -$ | |
| $<<, >>$ | |
| $\&\&$ | lowest |
| $=$ | |

(a) $x = a+b*c \% d \rightarrow$

$x = a+b*c \% d \rightarrow$

$x = a+b*c \% d \rightarrow$

(b) $y = a \& b \& c \rightarrow$

$y = a \& b \& c \rightarrow$

(c). $xabc * d \% + e \Rightarrow$

\downarrow infix

$x = a+b*c \% d \Rightarrow$

\downarrow prefix

$= x \Rightarrow +a \% * b c d e$

该表达式先有起来类似于数组在 postfix 中

(a) convert an infix expression, $x = a + b * c \% d \rightarrow e$, to a postfix expression

(b) convert a prefix expression, $= y \&\& << ab >> c + de$, to an infix expression

(c) convert a postfix expression, $xabc * d \% + e \Rightarrow =$, to a prefix expression

Q2 The function subset() below takes two linked lists of integers and determines whether the first is a subset of the second. Give the worst-case running time of subset as a function of the lengths of the two lists. When will this worst case happen?

```

1  typedef struct _listnode{
2      int item;
3      struct _listnode *next;
4  } ListNode;
5
6  //Check whether integer X is an element of linked list Q
7  int element (int X, ListNode* Q)
8  {
9      int found; //Flag whether X has been found
10     found = 0;
11     while ( Q != NULL && !found) {
12         found = Q->item == X;
13         Q = Q->next;
14     }
15     return found;
16 }
17
18 // Check whether L is a subset of M
19 int subset (ListNode* L, ListNode* M)
20 {
21     int success; // Flag whether L is a subset so far
22     success = 1;
23     while ( L != NULL && success) {
24         success = element(L->item, M);
25         L = L->next;
26     }
27     return success;
28 }
```

$M: N \text{ items} \quad | \quad 2 \quad 3 \quad 4 \quad \dots \quad n$

$L: N \text{ items} \quad | \quad n \quad \dots \quad 4 \quad 3 \quad 2 \quad 1$

$$n+(n-1)+\dots+1 = \frac{n(n+1)}{2} = \frac{n^2+n}{2}$$

Worst case scenario:

- the first $|L|-1$ items maps to the last $|M|-1$ in reverse
- last element of L not in M .

i.e. $L (|L| \text{ items}) \quad | \quad 1 \quad 2 \quad 3 \quad \dots \quad |L|$

$M (|M| \text{ items}) \quad | \quad 3 \quad 2 \quad 1$

$$|M| + |M-1| + |M-2| + \dots + |M| - (|L|-2) + |M|$$

$\underbrace{\quad \quad \quad}_{\text{共 } |L|-1 \text{ 次}}$

$$\begin{aligned} &= \frac{(2|M|-|L|+2) \times (|L|-1)}{2} + |M| \\ &= |L||M| - \frac{(|L|-2)(|L|-1)}{2} = \Theta(|L||M|) \end{aligned}$$

Q3 Find the number of printf used in the following functions. Write down its time complexity in Θ notation in terms of N .

```

1 void Q3a (int N)
2 {
3     int j, k;
4     for (j=1; j<=N; j*=3)
5         for(k=1; k<=N; k*=2)
6             printf("CZ1107\n");
7 }
```

$$\begin{aligned}
 & \text{Time complexity: } \Theta(f(N)) = \log_3 N \times \log_2 N \\
 & = (\log_2 N)^2
 \end{aligned}$$

```

1 void Q3b (int N)
2 {
3     int i;
4     if(N>0)
5     {
6         for(i=0; i<N; i++)
7             printf("CZ1107\n");
8         Q3b(N-1);
9         Q3b(N-1);
10    }
11 }
```

$$\begin{aligned}
 W(N) &= 2W(N-1) + N \\
 &= 2(2W(N-2) + N-1) + N \\
 &= 2^2 W(N-2) + 2N - 2 + N \\
 &= 2^2 (2W(N-3) + N-2) + 3N - 2 \\
 &= 2^3 W(N-3) + 2^2 N + 2N + N - 2 - 2 - 1. \\
 &= 2^n W(N-n) + \underbrace{(2^{n-1} + 2^{n-2} + \dots + 1)}_{\text{geometric series}} \times (N-1) \\
 &= 0 + (2^n - 1) \times (N-1)
 \end{aligned}$$

$$\Theta(W(N)) = (2^n - 1)(N-1)$$



CX1107 Data Structures and Algorithms Solution 4: Algorithm Analysis

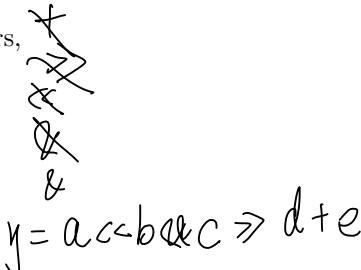
2020/21 Semester 2

School of Computer Science and Engineering

Nanyang Technological University

Q1 Given the precedence of some operators,

| Operators | Precedence |
|------------|------------|
| $*, /, \%$ | highest |
| $+, -$ | |
| $<<, >>$ | |
| $\&\&$ | |
| $=$ | lowest |



 $y = a << b \&\& c >> d + e$

- (a) convert an infix expression, $x = a + b * c \% d >> e$, to a postfix expression
- (b) convert a prefix expression, $= y \&\& << ab >> c + de$, to an infix expression
- (c) convert a postfix expression, $xabc * d \% + e >> =$, to a prefix expression

S1 (a) $xabc * d \% + e >> =$

- (b) $y = a << b \&\& c >> d + e$
- (c) $= x >> +a \% * bcde$

Q2 The function subset() below takes two linked lists of integers and determines whether the first is a subset of the second. Give the worst-case running time of subset as a function of the lengths of the two lists. When will this worst case happen?

```

1  typedef struct _listnode{
2      int item;
3      struct _listnode *next;
4  } ListNode;
5
6  //Check whether integer X is an element of linked list Q
7  int element (int X, ListNode* Q)
8  {
9      int found; //Flag whether X has been found
10     found = 0;
11     while ( Q != NULL && !found) {
12         found = Q->item == X;
13         Q = Q->next;
14     }
15     return found;
16 }
17
18 // Check whether L is a subset of M
19 int subset (ListNode* L, ListNode* M)
20 {
21     int success; // Flag whether L is a subset so far
22     success = 1;
23     while ( L != NULL && success) {
24         success = element(L->item, M);
25         L = L->next;
26     }
27     return success;
28 }
```

S2 Let $|L|$ and $|M|$ indicate the length of the linked lists, L and M , respectively. The worst-case running time of subset:

- the first $|L| - 1$ elements of L from the last $|L| - 1$ elements of M in reverse order.
- the last element of L not in M

\therefore Total number of comparisons between elements of L and M

$$\begin{aligned}
 &= |M| + (|M| - 1) + (|M| - 2) + \dots + (|M| - (|L| - 2)) + |M| \\
 &= |L||M| - (1 + 2 + 3 + \dots + (|L| - 2)) \\
 &= |L||M| - \frac{(|L| - 2)}{2}(1 + |L| - 2) \\
 &= |L||M| - \frac{(|L| - 2)(|L| - 1)}{2} \\
 &= \Theta(|L||M|)
 \end{aligned}$$

Here we assume that $|L| < |M|$

Q3 Find the number of printf used in the following functions. Write down its time complexity in Θ notation in terms of N .

```

1 void Q3a (int N)
2 {
3     int j, k;
4     for (j=1; j<=N; j*=3)
5         for(k=1; k<=N; k*=2)
6             printf("CZ1107\n");
7 }
```

```

1 void Q3b (int N)
2 {
3     int i;
4     if(N>0)
5     {
6         for(i=0; i<N; i++)
7             printf("CZ1107\n");
8         Q3b(N-1);
9         Q3b(N-1);
10    }
11 }
```

S3a let K denote the number of iterations for the inner loop and J denote the number of iteration for the outer loop.

For the inner loop, we have

$$\begin{aligned}
 2^{K-1} &\leq N < 2^K \\
 (K-1) &\leq \log_2 N < K \\
 K &\leq \log_2 N + 1 < K + 1 \\
 K &= \lfloor \log_2 N + 1 \rfloor = \lfloor \log_2 N \rfloor + 1
 \end{aligned}$$

For the outer loop, we have

$$\begin{aligned}
 3^{J-1} &\leq N < 3^J \\
 (J-1) &\leq \log_3 N < J \\
 J &\leq \log_3 N + 1 < J + 1 \\
 J &= \lfloor \log_3 N + 1 \rfloor = \lfloor \log_3 N \rfloor + 1
 \end{aligned}$$

\therefore The number of printf is $JK = (\lfloor \log_3 N \rfloor + 1)(\lfloor \log_2 N \rfloor + 1) = \Theta((\log N)^2)$

S3b let $W(N)$ denote the number of printf used in the function with problem size of N

$$\begin{aligned}
 W(N) &= 2W(N-1) + N \\
 &= 2(2W(N-2) + (N-1)) + N \\
 &= 2^2W(N-2) + 2(N-1) + N \\
 &= 2^2(2W(N-3) + (N-2)) + 2(N-1) + N \\
 &= \dots \\
 &= 2^{N-1}(1) + 2^{N-2}(2) + \dots + 2^3(N-3) + 2^2(N-2) + 2(N-1) + N \\
 &= \sum_{t=0}^{n-1} 2^t(n-t) \\
 &= 2^{N+1} - 2 - N
 \end{aligned}$$

\therefore The number of printf is $\Theta(2^N)$