CSE 326: Data Structures

Priority Queues: Leftist Heaps

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Leftist Heaps

Idea:

Focus all heap maintenance work in one small part of the heap

Leftist heaps:

- 1. Binary trees
- 2. Most nodes are on the left
- 3. All the merging work is done on the right

New Heap Operation: Merge

Given two heaps, merge them into one heap

 first attempt: insert each element of the smaller heap into the larger.

runtime:

 second attempt: concatenate binary heaps' arrays and run buildHeap.

runtime:

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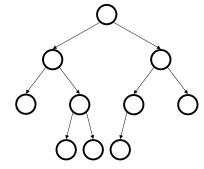
Definition: Null Path Length

null path length (npl) of a node x = the number of nodes between x and a null in its subtree

OR

npl(x) = min distance to a descendant with 0 or 1 children

- npl(null) = -1
- npl(leaf) = 0
- npl(single-child node) = 0



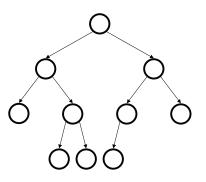
Equivalent definition:

$$npl(x) = 1 + \min\{npl(left(x)), npl(right(x))\}\$$

Definition: Null Path Length

Another useful definition:

npl(x) is the height of the largest perfect binary tree that is both itself rooted at x and contained within the subtree rooted at x.

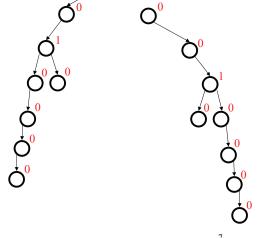


Leftist Heap Properties

- Order property
 - parent's priority value is ≤ to childrens' priority values
 - result: minimum element is at the root
 - (Same as binary heap)
- Structure property
 - For every node x, $npl(left(x)) \ge npl(right(x))$
 - result: tree is at least as "heavy" on the left as the right

(Terminology: we will say a leftist heap's tree is a leftist tree.)

Are These Leftist?



Observations

Are leftist trees always...

- complete?
- balanced?

Consider a subtree of a leftist tree...

- is it leftist?

Right Path in a Leftist Tree is Short (#1)

<u>Claim</u>: The right path (path from root to rightmost leaf) is as short as *any* in the tree.

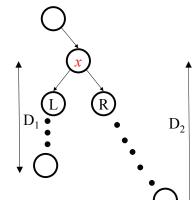
<u>Proof</u>: (By contradiction)

Pick a shorter path: $D_1 < D_2$ Say it diverges from right path at x

 $npl(L) \le D_1-1$ because of the path of length D_1-1 to null

 $npl(R) \ge D_2-1$ because every node on right path is leftist

Leftist property at *x* violated!



Right Path in a Leftist Tree is Short (#2)

Claim: If the right path has \mathbf{r} nodes, then the tree has at least $2^{r}-1$ nodes.

Proof: (By induction)

Base case : r=1. Tree has at least $2^1-1 = 1$ node

Inductive step: assume true for r-1. Prove for tree with right

path at least r.

1. Right subtree: right path of **r-1** nodes

 \Rightarrow 2^{r-1}-1 right subtree nodes (by induction)

2. Left subtree: also right path of length at least **r-1** (prev. slide)

 \Rightarrow 2^{r-1}-1 left subtree nodes (by induction)

 \Rightarrow Total tree size: (2^{r-1}-1) + (2^{r-1}-1) + 1 = 2^r-1

Why do we have the leftist property?

Because it guarantees that:

- the *right path is really short* compared to the number of nodes in the tree
- A leftist tree of N nodes, has a right path of at most log₂(N+1) nodes

Idea – perform all work on the right path

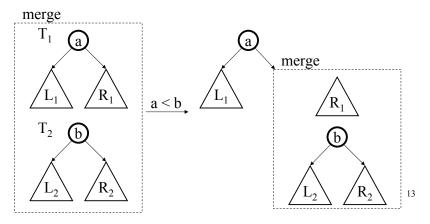
Merge two heaps (basic idea)

- Put the root with smaller value as the new root.
- Hang its left subtree on the left.
- <u>Recursively</u> merge its right subtree and the other tree.
- Before returning from recursion:
 - Update npl of merged root.
 - Swap left and right subtrees just below root, if needed, to keep leftist property of merged result.

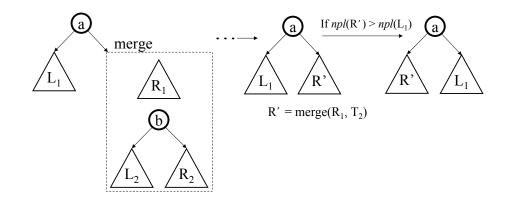
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Merging Two Leftist Heaps

Recursive calls to $merge(T_1,T_2)$: returns one leftist heap containing all elements of the two (distinct) leftist heaps T_1 and T_2

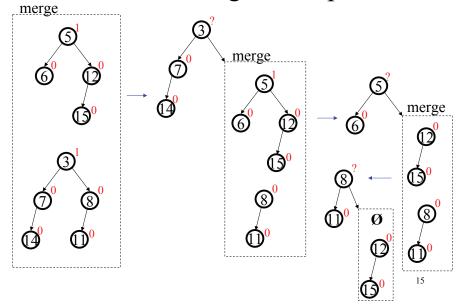


Merge Continued

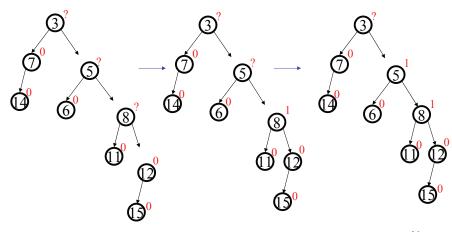


Note special case: merge(null, T) = merge(T, null) = T runtime:

Leftest Merge Example



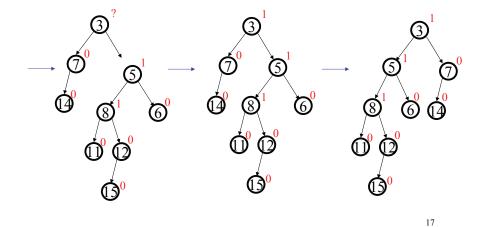
Sewing Up the Example



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Sewing Up the Example



Other Heap Operations

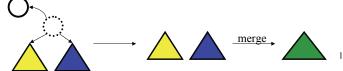
- insert?
- deleteMin?

Operations on Leftist Heaps

- merge with two trees of total size n: O(log n)
- insert with heap size n: O(log n)
 - pretend node is a size 1 leftist heap
 - insert by merging original heap with one node heap



- <u>deleteMin</u> with heap size n: O(log n)
 - remove and return root
 - merge left and right subtrees



Leftist Heaps: Summary

Good

- •
- •

Bad

- •
- •

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