The Conjugate Gradient Method

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Lecture Objectives

- describe when CG can be used to solve Ax = b
- relate CG to the method of conjugate directions
- describe what CG does geometrically
- explain each line in the CG algorithm

We are interested in solving the linear system

$$Ax = b$$

where $x, b \in \mathbb{R}^n$ and $A \in \mathbb{R}^{n \times n}$

Matrix is symmetric positive-definite (SPD)

$$A^{T} = A$$
 (symmetric)
 $x^{T}Ax > 0$, $\forall x \neq 0$ (positive-definite)

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- discretization of elliptic PDEs
- optimization of quadratic functionals
- nonlinear optimization problems

When A is SPD, solving the linear system is the same as minimizing the quadratic form

$$f(x) = \frac{1}{2}x^T A x - b^T x.$$

Why?

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Why? If x^* is the minimizing point, then

$$\nabla f(x^{\star}) = Ax^{\star} - b = 0$$

and, for $x \neq x^*$

$$f(x) - f(x^*) > 0$$
. (homework)

Definitions

Let x_i be the approximate solution to Ax = b at iteration i.

error:
$$e_i \equiv x_i - x$$

residual:
$$r_i \equiv b - Ax_i$$

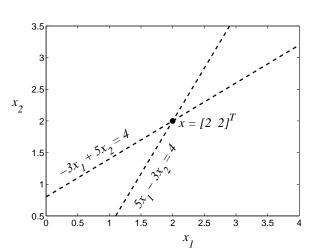
The following identities for the residual will be useful later.

$$r_i = -Ae_i$$

 $r_i = -\nabla f(x_i)$

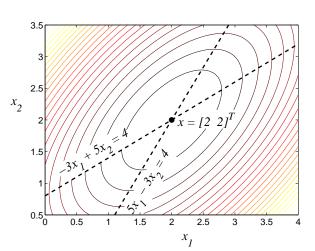
Model problem

$$\begin{bmatrix} 5 & -3 \\ -3 & 5 \end{bmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} 4 \\ 4 \end{pmatrix}$$



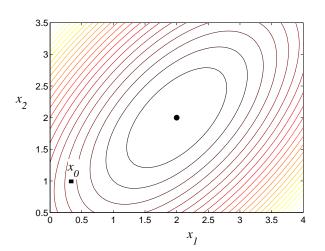
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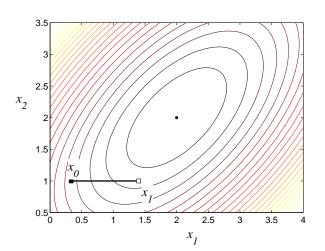
Review: Steepest Descent Method

Qualitatively, how will steepest descent proceed on our model problem, starting at $x_0 = (\frac{1}{3}, 1)^T$?



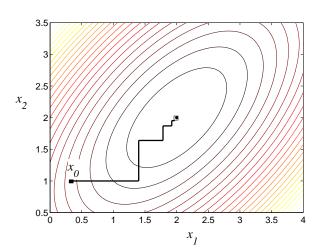
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How can we eliminate this zig-zag behaviour?

To find the answer, we begin by considering the easier problem

$$\begin{bmatrix} 2 & 0 \\ 0 & 8 \end{bmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} 4\sqrt{2} \\ 0 \end{pmatrix},$$

$$f(x) = x_1^2 + 4x_2^2 - 4\sqrt{2}x_1.$$

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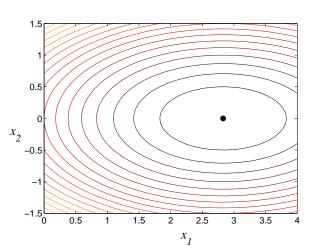
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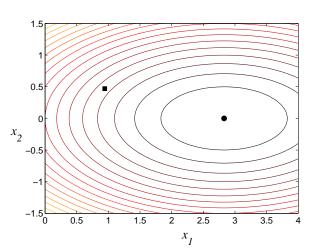
$$f(x) = x_1^2 + 4x_2^2 - 4\sqrt{2}x_1.$$

Here, the equations are decoupled, so we can minimize in each direction independently. What do the contours of the corresponding quadratic form look like?

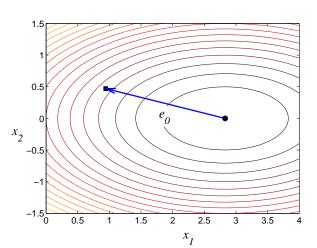
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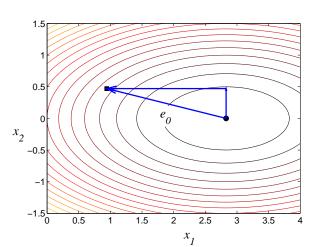
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Method of Orthogonal Directions

Idea: Express error as a sum of *n* orthogonal search directions

$$e \equiv x_0 - x = \sum_{i=0}^{n-1} \alpha_i d_i.$$

At iteration i + 1, eliminate component $\alpha_i d_i$.

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How would we apply the method of orthogonal directions to a non-diagonal matrix?

Review of Inner Products

The search directions in the method of orthogonal directions are orthogonal with respect to the dot product.

The dot product is an example of an inner product.

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Inner Product

For $x, y, z \in \mathbb{R}^n$ and $\alpha \in \mathbb{R}$, an inner product $(,): \mathbb{R}^n \times \mathbb{R}^n \to \mathbb{R}$ satisfies

- symmetry: (x, y) = (y, x)
- linearity: $(\alpha x + y, z) = \alpha(x, z) + (y, z)$
- positive-definiteness: $(x, x) > 0 \Leftrightarrow x \neq 0$

Fact: $(x, y)_A \equiv x^T A y$ is an inner product

A-orthogonality (conjugacy)

We say two vectors $x, y \in \mathbb{R}^n$ are A-orthogonal, or conjugate, if

$$(x,y)_A = x^T A y = 0.$$

What happens if we use *A*-orthogonality rather than standard orthogonality in the method of orthogonal directions?

Let $\{p_0, p_1, \dots, p_{n-1}\}$ be a set of n linearly independent vectors that are A-orthogonal. If p_i is the ith column of P, then

$$P^TAP = \Sigma$$

where Σ is a diagonal matrix.

Substitute x = Py into the quadratic form:

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Substitute x = Py into the quadratic form:

$$f(Py) = y^T \Sigma y - (P^T b)^T y.$$

We can apply the method of orthogonal directions in *y*-space.

New Problem: how do we get the set $\{p_i\}$ of conjugate vectors?

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Gram-Schmidt Conjugation

Let $\{d_0, d_1, \dots, d_{n-1}\}$ be a set of linearly independent vectors, e.g., coordinate axes.

- set $p_0 = d_0$
- for i > 0

$$p_i = d_i - \sum_{j=0}^{i-1} \beta_{ij} p_j$$

where $\beta_{ij} = (d_i, p_i)_A/(p_i, p_i)_A$.

Force the error at iteration i + 1 to be conjugate to the search direction p_i .

$$p_i^T A e_{i+1} = p_i^T A (e_i + \alpha_i p_i) = 0$$

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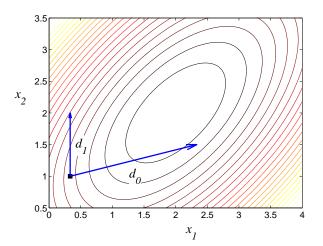
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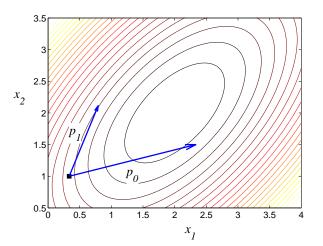
$$= \frac{p_i^T r_i}{p_i^T A p_i}$$

- never need to search along p_i again
- converge in n iterations!

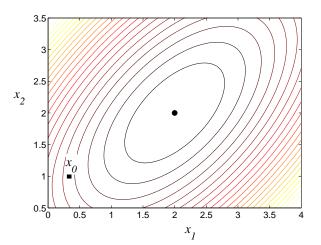
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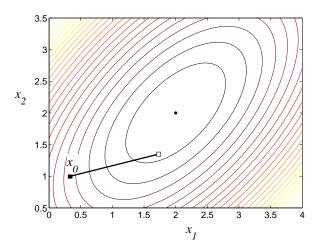
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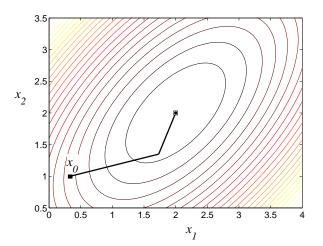
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The Method of Conjugate Directions is well defined, and avoids the "zig-zagging" of Steepest Descent.

What about computational expense?

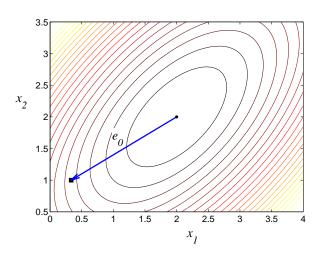
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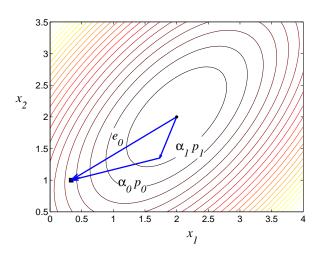
- If we choose the d_i in Gram-Schmidt conjugation to be the coordinate axes, the Method of Conjugate Directions is equivalent to Gaussian elimination.
- Keeping all the p_i is the same as storing a dense matrix!

Can we find a smarter choice for d_i ?

Error Decomposition Using p_i



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$$e_i = \sum_{k=i}^{n-1} \alpha_k p_k,$$

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so the error must be conjugate to p_j for j < i:

$$\label{eq:posterior} \begin{split} \rho_j^T A e_i &= 0, \quad \Rightarrow \rho_j^T r_i = 0, \end{split}$$

but from Gram-Schmidt conjugation we have

$$p_j = d_j - \sum_{k=0}^{j-1} \beta_{jk} p_k.$$

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but from Gram-Schmidt conjugation we have

$$p_j^T r_i = d_j^T r_i - \sum_{k=0}^{j-1} \beta_{jk} p_k^T r_i$$
$$0 = d_j^T r_i, \quad j < i.$$

Thus, the residual at iteration i is orthogonal to the vectors d_i used in the previous iterations:

$$d_i^T r_i = 0, \quad j < i$$

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$$d_i^T r_i = 0, \quad j < i$$

Idea: what happens if we choose $d_i = r_i$?

- residuals become mutually orthogonal
- r_i is orthogonal to p_i , for $j < i^*$
- r_{i+1} becomes conjugate to p_i , for j < i

This last point is not immediately obvious, so we will prove it. This result has significant implications for Gram-Schmidt conjugation.

^{*}we showed this is true for any choice of d_i

$$x_{j+1} = x_j + \alpha_j p_j$$

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Next, take the dot product of both sides with an arbitrary residual r_i :

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We can show that the first case (i = j) contains no new information (homework). Divide the remaining cases by $p_i^T A p_j$ and insert the definition of α_{i-1} :

$$\frac{r_i^T A p_j}{p_j^T A p_j} = \begin{cases} -\frac{r_i^T r_i}{r_{i-1}^T r_{i-1}}, & i = j+1\\ 0, & \text{otherwise.} \end{cases}$$

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We recognize the L.H.S. as the coefficients in Gram-Schmidt conjugation

only one coefficient is nonzero!

Set
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 and $i = 0$

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$$r_{i+1} = r_i - \alpha_i A p_i$$
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$$\beta_{i+1,i} = -(r_{i+1}^T r_{i+1})/(r_i^T r_i)$$
 (G.S. coeff.)

$$\alpha_{i} = (p_{i}^{T}r_{i})/(p_{i}^{T}Ap_{i}) \qquad \text{(step length)}$$

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$$p_{i+1} = r_{i+1} - \beta_{i+1,i} p_{i} \qquad \text{(Gram Schmidt)}$$

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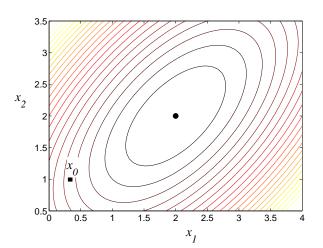
i := i + 1

$$Ap_i$$

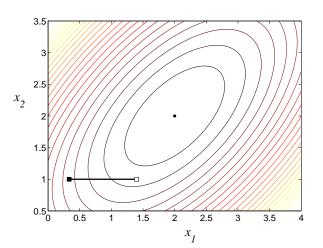
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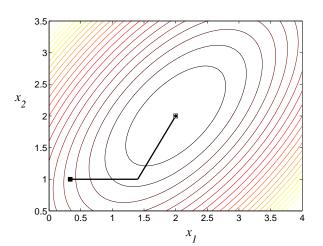
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- explain each line in the CG algorithm

References

- Saad, Y., "Iterative Methods for Sparse Linear Systems", second edition
- Shewchuk, J. R., "An introduction to the Conjugate Gradient method without the agonizing pain"