

for $f \in I$ (really $f = \sum C_{\alpha, f} y^\alpha \Rightarrow \pi_*^{-1}(f, m) = [\sum C_{\alpha, f} y^\alpha]_S^m$)

where

$$C_{\alpha, f} y^\alpha = y_q^{|\alpha|-m} C_{\alpha, f} (y_1 y_2, \dots, y_q, \dots, y_m)$$

while

$$f|_S = C_{\alpha, f}|_S \quad \text{and} \quad \pi_*^{-1}(f, m)|_S = C_{\alpha, f} y^\alpha|_S,$$

$$\begin{aligned} \pi_*^{-1}(f, m)|_S &= (C_{\alpha, f})|_S = y_q^{|\alpha|-m} (\underbrace{\pi^* C_{\alpha, f}}_{\text{composition}})|_S = y_q^{|\alpha|-m} \pi|_S^*(C_{\alpha, f}|_S) \\ &= y_q^{|\alpha|-m} \pi|_S^*(f|_S) = \pi|_{S^*}^{-1}(f|_S, m) \end{aligned}$$

$$\Rightarrow \pi_*^{-1}(I, m)|_S = \pi|_{S^*}^{-1}(I, m)|_S.$$

(4). Notations as in (3).

As before, $\text{Supp } (\pi_*^{-1} I, m) \cap S' = \bigcap_{\substack{f \in I \\ |\alpha| \leq m}} \text{Supp } (C_{\alpha, f}|_S), m = |\alpha|$

$$\text{Now } C_{\alpha, f} = y_m^{-m+|\alpha|} \pi^*(C_{\alpha, f}) \subseteq \pi_*^{-1} D^{\otimes m}(I, m).$$

$$\begin{aligned} \text{back to LHS} &\supseteq \bigcap_{0 \leq |\alpha| \leq m} \text{Supp } (\pi_*^{-1} D^{\otimes m}(I, m)|_S) \\ &= \bigcap_{i=0}^m \text{Supp } (\pi|_{S^*}^{-1}(D^{\otimes i}(I, m)|_S)). \\ &= \text{Supp } (\pi|_{S^*}^{-1} C(I, m)|_S). \end{aligned}$$

$$\text{so } S' \cap \text{Supp } (\pi_*^{-1} C(I, m)) \supseteq S' \cap \text{Supp } (\pi_*^{-1}(I, m)) \supseteq \text{Supp } (\pi|_{S^*}^{-1} C(I, m)|_S).$$

$$\begin{array}{c} \text{Supp } (\pi_*^{-1} C(I, m)|_S) \\ \parallel \\ \text{Supp } (\pi|_{S^*}^{-1} (C(I, m)|_S)). \end{array}$$

□

Remark: The above proposition says that, $S \subset X$ not contained in $\text{supp}(I.m)$, an order reduction for $C(I.m)|_S$ on S lifts naturally to an "order reduction" on X .

To be more specific

$$Z \subset S \subset X$$

$$\begin{array}{ll} B : & \pi : X_r \rightarrow \dots \rightarrow X_0 \\ & \downarrow \text{lift} \quad \downarrow \quad \downarrow \\ B_S : & \pi_S : S_r \rightarrow \dots \rightarrow S_0 \supseteq C(I.m)|_S \end{array}$$

$$(1) Z_i \subseteq \text{Supp}([C(I.m)|_S]_i) \Rightarrow Z_i \subseteq \text{Supp}[C(I.m)_i] \cap S_i$$

$$(2) \text{Supp}([C(I.m)|_S]_r) = \emptyset \Rightarrow \text{Supp}[C(I.m)_r] \cap S_r = \emptyset.$$

i.e. $\text{Supp}(I.m)_r = \text{supp}(C(I.m))_r$ is disjoint with S_r .

(3) If B_S is functorial resp to smooth morphisms, then the natural lifting is also functorial resp to smooth morphism.

In fact $S_Y \hookrightarrow Y$ $h : Y \rightarrow X$ smooth $\Rightarrow h_Y : S_Y \rightarrow S_X$ smooth.

$$\begin{array}{ccc} h_Y \downarrow & \downarrow h \text{ smooth} & \text{lift } B_S \text{ to } B, \text{ blow-up center is } Z. \\ Z \subset S \hookrightarrow X & & \text{functoriality for } B_S \text{ imply blow-up center for } S_Y \text{ is } h_Y^{-1}(Z). \\ & & \text{lift } \not\rightarrow Y, \text{ blow-up center is again } h_Y^{-1}(Z) \subset Y. \end{array}$$

□.

Rem: In previous case, we only consider the restriction \rightarrow ord reduction \rightarrow lifting that end up with $S_r \cap \text{Supp}(I.m)_r = \emptyset$.

Key: If we can find $S \supseteq \text{Supp}(I.m)$ such that each time.

maxi cont. $\longrightarrow S_i \supseteq \text{Supp}[C(I.m)_i]$ then we end up with $\emptyset = S_r \cap \text{Supp}(I.m)_r = \text{Supp}(I.m)_r$!

Def-Prop 2.7 (Hypersurface of Maximal contact).

The maxi contact ideal sheaf of $(I.m)$ is $(T(I))_i = D^m(I.m)$ $m = \text{maxord } I$.

For any $x \in \text{Supp}(I.m) = \text{Supp}(T(I))$, \exists open neighbor $x \in U_x$, and

a smooth element $h \in T(I)(U)$ ($V(h) \cong H$ is sm hypersurface on U_x)

with $I|_H \neq 0$, we call H a hypersurface of maximal contact.

Exam: x^2+y^3 , maxord=2, $D((x^2+y^3)) = (x,y^2)$ $x+cy^2$ is a hysurf of m.c.

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Now, $\pi: Bl_Z U \rightarrow U$ a sm blow up with $Z \subseteq \text{Supp}(I.m) \cap H$, we have

$$\text{Supp}(\pi_*^{-1}(I.m)) \subset \pi^{-1}H.$$

Proof: $\text{Supp}(\pi(I), 1) \subseteq V(h) = H$.

$$\text{Supp}(I.m)$$

$$\text{since } \pi_*^{-1}(h, 1) \subseteq \pi_*^{-1}(\pi(I), 1) \subseteq (\pi(\pi_*^{-1}(I)), 1)$$

$$\Rightarrow \text{Supp}(\pi_*^{-1}I, m) = \text{Supp}(\pi(\pi_*^{-1}I), 1) \subseteq \text{Supp}(\pi_*^{-1}h) = h^*H.$$

Rem: the maximal contact hypersurface is local and depends on choice of h .

(That is where $H(I)$ plays a role).

Lem 2.8 Let (X, I, m, E) be a marked triple, $m = \max \text{ord } I$.

for any $u, v \in T(X, m)_x$ at $x \in \text{Supp}(I.m)$ that are smooth and snc with E . Then we have automorphism

$$\overset{\wedge}{\phi}_{uv} \text{ of } \hat{X}_x = \text{Spec } \overset{\wedge}{\mathcal{O}}_{x,x}$$

s.t. (1) $\overset{\wedge}{\phi}_{uv}^*(H(I))_x = (H(I)_x)$

(2) $\overset{\wedge}{\phi}_{uv}^* E = E$

(3) $\overset{\wedge}{\phi}_{uv}^*(u) = v$

(4) $\text{Supp}(\hat{I}.m) = V(T(\hat{I}, m))$ is in the fixed point set of $\overset{\wedge}{\phi}_{uv}$.

Proof: Step 1 construction.

Take $u = u_1, u_2, \dots, u_n$ s.t. both u or v , u_2, u_3, \dots, u_n form local coordinates and is compatible with E .

Set $\overset{\wedge}{\phi}_{uv}(u) = v$ $\overset{\wedge}{\phi}_{uv}(u_i) = u_i$ for $i > 0$.

Step 2: Variation.

Let $h = v - u \in T(I)$. $\forall f \in \hat{I}$

$$\overset{\wedge}{\phi}_{uv}^* f = f(u+h, u_2, \dots, u_n)$$

$$= f(u_1, \dots, u_n) + \frac{\partial f}{\partial u_1} h + \frac{1}{2!} \frac{\partial^2 f}{\partial u_1^2} h^2 + \dots$$

$$\subseteq \hat{I} + \hat{D}\hat{I} \cdot \hat{T}\hat{I} + \dots + \hat{D}^i\hat{I} \cdot \hat{T}\hat{I}:$$

$$\frac{\partial^i f}{\partial u_1^i} h^i$$

$$\Rightarrow \overset{\wedge}{\phi}_{uv}^* \hat{I} \subset H \hat{I}. \quad \text{Similarly } \overset{\wedge}{\phi}_{uv}^* (D^i \hat{I}) \subset H D^i \hat{I} \quad \overset{\wedge}{\phi}_{uv}^* T(\hat{I}) \subset T(H \hat{I})$$

$$T(\hat{I})$$

to sum up, $\hat{\phi}_{uv}^*(D^i \hat{I} \cdot \hat{T}(I)^i) \subset D^i \hat{I} \cdot \hat{T}(I)^i + \cdots + D^{m-1} \hat{I} \cdot \hat{T}(I)^{m-1} \cdot \hat{T}(I)^m \subset H(\hat{I})$.
 $\Rightarrow \hat{\phi}_{uv}^* H(\hat{I}) \subset H(\hat{I})$ Noetherian properties guarantees that
 $\hat{\phi}_{uv}^{*n}(H(\hat{I})) = \hat{\phi}_{uv}^{*(n)}(H(\hat{I})) \Rightarrow (1) \checkmark$.

(2) (3) \checkmark by construction

(4) $h=0$ is fixed by $\hat{\phi}_{uv}^*$ $\Rightarrow \text{Supp}(T(I).1)$ is fixed $\Rightarrow \text{Supp}(I, m)$ fixed.

Formal local uniqueness imply étale equivalence.

Lem 2.9 Settings as in Lem 2.8.

Then there exists étale neighborhoods

$$\phi_u, \phi_v : U \xrightarrow{\psi} X \text{ of } x = \phi_u(\tilde{x}) = \phi_v(\tilde{x})$$

$$\tilde{x} \in U \xrightarrow{\phi_u} X \xrightarrow{\phi_v}$$

s.t. (1) $\phi_u^*(X, H(I), m, E) = \phi_v^*(X, H(I), m, E) := (\tilde{X}, \tilde{H}(\tilde{I}), m, \tilde{E})$

(2) $\phi_u^*(u) = \phi_v^*(v)$

(3) $IB : X_r \rightarrow \dots \rightarrow X_0$ be a seq of sm blow-up with Z_i in $\text{Supp}(I, m)$

then $\phi_u^* IB(X, H(I), m, E) = \phi_v^* IB(X, H(I), m, E) : \tilde{X}_r \rightarrow \dots \rightarrow \tilde{X}_0$

$\phi_{ui} \circ \phi_{vi} : \tilde{X}_i \rightarrow X_i$ satisfies

$$\phi_{ui}^*(V_{(W,i)}) = \phi_{vi}^{-1}(V_{(W,i)}) \text{ and}$$

$$\phi_{ui}(\tilde{y}_i) = \phi_{vi}(\tilde{y}_i) \quad \forall \tilde{y} \in \text{Supp}(\tilde{I}_i, m).$$

Remark: Lemma 2.9 allow us to glue restricted resolution! $\forall x \in X$.

that is, $\forall U_{(u,x)}$ and $U_{(v,x)}$ two open set that restricted to $V_{(u)}$, $V_{(v)}$
and def blow up seq and lift to Blow up seq $B_u(U_{(u,x)})$, $B_v(U_{(v,x)})$

$$\exists \quad U_{(uv,x)} \xrightarrow{\phi_u} U_{(u,x)} \cap U_{(v,x)}.$$

$$\text{s.t. } \phi_u^* B_u(U_{(u \cap v)}) = \phi_v^* B_v(U_{(u \cap v)}).$$

\Rightarrow restricted to $U_{(u \cap v)}$, the blow up center for $B_u(U)$, $B_v(U)$ coincide!

We can glue blow up center and globalize it as in L1.

And sm func preserved.

§ 3. Algorithm of Resolution

Now we are ready to prove

$$\text{Ord II in dim } \leq n-1 \xrightarrow{T_1} \text{Ord I in dim } n \xrightarrow{T_2} \text{Ord II in dim } n.$$

(div smooth)

T_1 : We start from $(X, I, E) \quad E = (E^1, \dots, E^s)$, $\max \text{ord } I \leq m$
 If $\max \text{ord } I < m$, the process is trivial. We assume $\max \text{ord } I = m$

Step 1: Construct $T_{I_1}: (X_{r_1}, I_{r_1}, E_{r_1}) \rightarrow \dots \rightarrow (X, I, E)$ s.t.

$$\text{Supp } T_{I_1}^{-1} E \cap \text{Supp } (I_{r_1}, m) = \emptyset.$$

Consider the equivalent ideal $C(I)$ now

1.1. Let Z_0 be the union of all irr comp of E' contained in $\text{Supp}(I, m)$.
 Blow up Z_0 , we have (for irr comp E'^k in E' for example) $T_{I_0}: X \rightarrow X$
 $\max \text{ord}_{E'^k} I = m \Rightarrow \max_{T_{I_0}^{-1} E'^k} \text{ord } I \leq m - m = 0$.
 $\Rightarrow T_{I_0}^{-1} E'^k$ and $\text{Supp}(I_{r_1}, m)$ are disjoint.

1.2. Now, set $S = E'$, $E_S = (E - E')|_S = (0, E^2|_S, \dots, E^s|_S)$

consider $(S, I|_S, m, E|_S)$ apply Ord II in dim $\leq n-1$, we get
 $B\Omega_1: (S_{E_1}, I_{E_1}|_{S_{E_1}}, m, E|_S)$ lift $(X_{r_{E_1}}, I_{r_{E_1}}, m, E_S) \rightarrow \dots \rightarrow (\dots)$
 s.t. $\text{Supp } (I_{r_{E_1}}, m) \cap T_{E_1}^{-1} S = \emptyset$

Inductively, we have $T_{I_1}: (\dots) \rightarrow \dots (\dots)$ s.t. $\text{Supp } T_{I_1}^{-1} E \cap \text{Supp } (I_{r_1}, m) = \emptyset$
 Functionality follows from previous remark.

Step 2: Start from $J = H(C(I_{r_1}))$, $Y = X_{r_1}$, $F = (T_{E_1}^{-1} E, E^{s+1}, \dots, E^{s+r_1})$
 (Y, J, m, F).

$\forall y \in \text{Supp } (J, m)$, $\exists U_y$ h.t. $y \in U_y$. Since in step 1, all blow-up is snc,
 we can take h s.t. H_h snc with F ($T_{E_1}^{-1} E$ away from $\text{Supp}(I, m)$).
 Locally consider $(H_h, J|_{H_h}, m, F|_{H_h})$

Apply Ord II in dim $< n$ and lift it and globalized it to

$$T_{I_2}: (X_{r_2}, I_{r_2}, E_{r_2}) \rightarrow (X_{r_1}, I_{r_1}, E_{r_1})$$

$$\text{s.t. } \text{Supp } (I_{r_2}, m) = \emptyset !$$

(Note, all blow up seq is also for $C(I)$ and I , $T_H(C(I)) = C(I)$)

Functionality follows from previous rem.

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□.

T2: We start with a marked triple (X, I, m, E) $E = (E^1, \dots, E^s)$. 2.

Step 0:

We may write $I = N(I) M(I)$, $M(I) = \mathcal{O}_X(-\sum_{i=1}^s E_i)$ and $\text{Supp } N(I)$ does not contain any of E^i !

Rem: if $E = \emptyset$, $M(I) = \emptyset$ and $I = N(I)$.

Step 1: Write $(N(I), 1) + (I, m) = (J, s)$ here s a number,

Write $m_J = \max \text{ord } J$, Run $\text{Ord } I$ to (X, J, E) with m_J , we get (X_1, J_1, E_1) s.t. $\max \text{ord } J_1 = m_{J_1} < m_J$.

(Note, $\text{Supp } (J, m_J)_K \subseteq \text{Supp } (J, s)_K = \text{Supp } (N(I), 1)_K \cap \text{Supp } (I, m)_K$)

Inductively we get

(X_r, J_r, E_r) s.t. $\max \text{ord } J_r < s$.

this implies $\emptyset = \text{Supp } (J, s)_r = \text{Supp } (N(I), 1)_r \cap \text{Supp } (I, m)_r$
 $\text{Supp } (N(I)_r)$.

To sum up, we have $(X_r, I_r, m, E_r) \rightarrow \dots \rightarrow (X, I, m, E)$

s.t. $N(I)_r \cap \text{Supp } (I, m)_r = \emptyset$.

Rem: $N(I)_r$ $N(I_r)$ differs by some exceptional comps $E^{k (k > s)}$
 $M(I)_r$ $M(I_r)$ and is contained in $M(I_r)$.

Step 2: $I = M(I) = \mathcal{O}_X(-\sum a_j E^j)$ $E = (E^1, E^2, \dots, E^s, \dots)$.

2.1 Sub { $E^1, E^2, \dots, E^s, \dots$ } has a lexicographic order.

($x \dots$).

2.2. $\forall x \in X$, set $p(x) = (\{E^{j_1}, \dots, E^{j_k}\})$ the maximal subset (in above order) satisfying

(1) E^{j_i} pass x $\forall 1 \leq i \leq k$.

(2) $\sum a_{j_i} \geq m$ (3) $a_{j_1} + \dots + \hat{a}_{j_i} + \dots + a_{j_k} < m$.

$D_{p(x)} = \bigcap_{i=1}^k E^{j_i}$, and it is the focus that is a maximal component of

$\text{Supp } (I, m)$.

$v = (\max \text{ord } I, \text{ member of maximal comp of } \text{Supp } (I, m) \text{ attain maxord})$
 $= (m, n)$

Each time, we blow up $D_{p(x)}$ $\Rightarrow D_{p(x)} = \{x_{ji} = \dots = x_{jk} = 0\}$.

$$\begin{aligned} & \forall x \in D_{p(x)}, \quad \text{since } \sum_{i=1}^k q_{ji} - m < a_{ji} \quad \forall |s| \leq k, \\ & Ix = \prod_{t \in s} x_t^{a_t} \\ & \text{and } \underbrace{Ix}_{\text{in cod } \prod_{t \neq j_i} x_t^{a_t}} = \prod_{t \neq j_i} x_t^{a_t} \cdot x_{j_i}^{\sum a_{ji} - m} < q_{ji}^{a_{ji} + \hat{a}_{ji} + \dots + a_{jk} - m + a_{ji}} \end{aligned}$$

$$\Rightarrow \text{ord}_y T_x^{-1}(I) < \text{ord}_x I \quad \forall \pi(y)=x.$$

\Rightarrow v decrease strictly in the lexicographic order \Rightarrow the procedure terminates with $v = (< m, \#)$. $\Rightarrow \max \text{ord } I_r < m$ eventually.

Functionality follows from the sm invariant property of $p(x)$ and v . \square .

OCE: Commute resp to closed embedding or $\text{Ord } II \underset{Is}{\rightarrow} E \neq \emptyset$

$$T: S \rightarrow X$$

Now, $T^* Is$ in X we have locally

$T^* Is$ is adding some smooth element $\{u\}$.

$$\text{Ord}(T^* Is, 1) = 1 \quad T(Is) = T^* Is \leftarrow \text{sm element.}$$

Recall T_1 : We restrict to sm by sf in T to do induction, so for this case, we just restrict $T^* Is$ to $\{u\}$ and exactly get I , the procedure commute with closed embedding. \square