Page NO 1 unit 2: Design of Combinational Logic (10 hrs) Minimal two level networks - Minimization of Pos and SOP . (2) Design of two level gate networks. (2) 1. Two level NAND-NAND NETWOOK (1) 2. Two level NOR - NOR Network (1) karnaugh maps - Advantages and Limitations . (2) Quine McClusky's Method. (2). Switching function! Expressions are Constructed a Constants and variables with the Boolean Operations. These expressions are also known as Boolean formulas. We use these expressions to describe surtching functions (on Boolean function. Example: (A+B)C F(A,B,c) = (A+B)c (on) F = (A+B)c Let us consider the four variable Boolean function. F(A/B/C/D) = A+Bc+Aco literals. <u>Literals</u>: Fach occurance of a variable in either a complemented or an uncomplemented form is Called a literals.

Product: product of literals

F(A,B,C,D) = (B+D) · (A+B+C) · (A+C).

Sum Sum

There are seven literals, and three sum terms

These literals and terms are arranged in one of the two torms.

- 1. Sum of product form (SOP)
- 2. Product of sum form (POS)

## Sum of Product form:

The word sum and product are derived from the symbolic representations of the DR and AND tunction by + and •

But these are not arithemetic operations.

Here all the product terms are sumed with the operator +.

Frample:

1. 
$$f(A_1B_1C) = ABC + ABC$$

2.  $f(P_1A_1R_1S) = PA + QR + PS$ 

Product of sum;

Sum

\* Sum terms are Anded together.

2. F(PIQIRIS) = (P+Q)P-D(R+S) (P+S)

#### Standard sop and pos form:

we can realise that in the SOP yourn all the individual terms do not involve all the diterals.

Example AB+ ABC

- (i) the first product term do not Contain diteral C.
- (ii) It each term in sop form contains all the literals then the sop form is called standard

Pg. No. 2 Example! 1. Convert the given expression in standard sop form. F(AIBIC) = AC+AB+BC. c is missing Solution: FLA, B, C) = AC + AB + BC B is missing A is missing. AND Product term with (missing literals + its Complement) F(AIBIC) = A( · (B+B) + AB· (C+Z) + BC (A+A) mussing literals and their Complements Expand the terms and reorder literals. FOLDAND F(A,B(C) = ACB + ACB + ABC + ABC + ABC + ABC Reorder F(A1B1C) = ABC + ABC + ABC + ABC + ABC Omit repeated product terms f(A,B,c) = ABC + ABC + ABC + ABC 2. Convert the given expression in standard sop 40aw FCA, B(C) = A + ABC FCAIBIC) = A. CB+B). (C+E) + ABC Expand the term and re-order HAIBIC) = AB+ AB. (C+T)+ABC = ABC + ABC + ABC + ABC + ABC + ABC = ABC + ABC + ABC + ABC F=BZ+Ac in a canonical sop torm. Express solution F=BC+AC = (A+A)BC + ACCB+B) = ABC + ABC + ABC + ABC

Q! Express Fi in Standard Jop torm.  $F_1 = AB + \overline{C}D + A\overline{B}C$  And  $\overline{A}\overline{B}\overline{C}D + \overline{A}B\overline{C}D + \overline{A}B\overline{C}$ 

SOP form.

viii) Fach individual term in sop form is called min term.

## Example you standard sop town:

F(A/BIC) = ABC + ABC + ABC

All the terms consists of all literals in either complemented or uncomplemented form

## Standard Pos torm!

- (i) If each term in Pas form Contains all the literals then the Pos torm is known as Standard or Cononical Pos Jorm.
- (ii) Each individual term in the standard Pos form is called max term Example for standard Pos form.

FLAIBIC) = (A+B+C) · (A+B+C)

All the terms consists of all literals in either Complemented or uncomplemented form.

# Steps to Convert SOP to stoundard SOP torm:

Step1: find the missing literal in each product term if any.

<u>step 2</u>: AND each product term having muxing literals with term torm by or ing the literals and its Complement

Steps: Expand the terms by explying distributive law and recorder the literals in the product term.

Step4: Reduce the expression by omitting repeated Product terms if any. Because A+A = A.

```
Convert the given expressions in standard POS torm. Pg. NO.3
       F(A,B,C) = (A+B). (B+C).
  Solution:
                 , A is missing.
       = (A+B) (B+C)
          c is missing
       = (A+B)+ c.c. (B+C)+A.A
       = (A+B+C) (A+B+C) (A+B+C) (A+B+C)
                            repeated
        = (A+B+C) (A+B+C) (A+B+C)
* convert sop to equivalent Pos.
         ABC + ABC + ABC + ABC + ABC.
    Solution.
         = ABC + ABC + ABC + ABC
               repeated
          = ABC+ ABC+ ABC +ABC,
      Pas Joan
          = (A+B+c) (A+B+c) (A+B+c) (A+B+c).
vconvert (A+B) (A+C) (B+C) untostandard Pos form.
          F(A,BIC) = A+B+C.E A+C+B.B B+C+A.A
    solution:
                   = (A+B+C) (A+B+C) (A+B+C)
                               repeated
                     (A+B+C) (A+B+C), repeated
                   = (A+B+C) (A+B+C) (A+B+C).
      Canonical POS 400 FCAIBIC) = (A+B) (B+C) (A+C)
Obtain
   Solution
          = A+B+C·C B+C+A·A A+C+B·B
          - (A+B+C) (A+B+C) (A+B+C) (A+B+C)
            (A+B+c) (A+B+c) (A+B+c) (A+B+c) (A+B+c)
             the given expression in standard pos torm
  91:
       Convert
         FCPIQIR) = (P+Q) (P+R).
```

Ans: (P+g+R) (P+g+R) (P+g+R)

## Steps to Convert Pos to standard Pos.

Step1: Find the missing literals in each sum term it any.

Step2: DR each sum term having missing literals with

terms form by Anding the literal and its Complement.

Step 8: Expand the terms by applying distributive law and

Yeosder the literals in the sum term.

Step 4: Reduce the expression by omitting repeated sum terms

If any, Because A·A = A.

#### Examples:

1. Convert the given expression in standard Pos form.

(A+B) (B+C) (A+C)

c'is missing B is missing.

 $F(A_1B_1C) = (A+B)(C\cdot C) \cdot B+C+A\cdot \overline{A} \cdot A+C+B\cdot \overline{B}$ Expand and reorder

= (A+B+C) (A+B+C) (B+C+A) (B+C+A) (A+B+C) repeated

F(AIBIC) = (A+B+C) (A+B+C) (A+B+C),

2. Convert the given expression in standard Postorm.  $y = A \cdot (A + B + C)$ 

y = A. (A+B+c) T B,c missing.

= A+B·B+c·c (A+B+c) Expand and reorder.

=(P+B·B+C) (P+B·B+C) (A+B+C)

(A+B) (A+C+B) (A+C+B) (A+C+B) (A+C+B)

(A+B+C) repeated

- (A+B+c) (A+C+B) (A+B+C) (A+C+B)

#### Minterms and Maxterms.

\* Each individual term is standard sop form is valled

\* Each individual term in standard Pos form is valled maxterns.

Three variable logical functions are expressed as  $2^3 = 8$ .

Variables		oles minterms Maxterms		Maxterms	Example (1) F (A; B; C) = ABC + ABC+
A	В	c	me	Me	$\mathbf{A}\mathbf{B}\mathbf{C}$
0	0	0	ABC = Mo	A+B+C = M	= mgm, + mgt m = zm(0,1,3,6) (2) F(A B C)=(A+B+C)(A+B+C) (5+B+C)
0	0	1	ABC = M	A+B+c=Mi	(2) F(AIBIC)=(A+B+C)(A+B+C
0				A+B+C= M2	(MT B (C)
	1	9	4	A+B+&=M3	= M, · M3 · M6. = TI M (1, 3,6)
	D	0	ABC = M4	A+B+C=M4	Sum of products denoted &
1	D		ABC = M5	F+B+C=M5	Sum of products control .
1	1	0	ABC = MG	A+B+C=Mb	Product of Sum denoted by
1	1	1	ABC = MT	A+B+C=M7	T

## Complements of Standard forms;

\* It the given function is

F(ABIC) = mo+m1+m3+m4+m6+m4.

= M2 • M5

: F(AIBIC) = &m(O11, 314, 6,7) = TM (215)

at It the given tunction is

F(AIBICIP) = &m(012141618110112114). then

FLAIBICID) = TM (1,3,5,7,9,11,13,15).

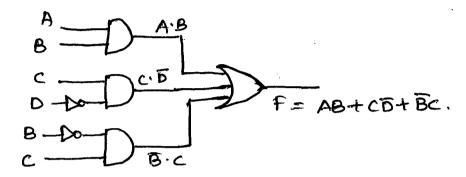
Express the switching function f(BA) = A in terms of minterms f(BA) = A = ACB+B) = AB+AB Express f = A+Bc as sum of minterms. A+BC = A (B+B) (C+C) + B C(A+A) = AB+AB (C+C)+ABC+ABC. = ABC + ABC + ABC + ABC + ABE)+ABC = ABC + ABC + ABC + ABC , f= Em (1,4,5,6,7) Prove that the logical sum of all minterms of a Boolean function of 2 variables is 1. Solution: AB, AB, AB, AB. F = AB+AB+AB+AB. = BCA+A)+ BCA+A)  $=\overline{B}+B$ = 1 Thus proved. the Boolean function F=XY+XZ Express product of max terms.  $F = \chi y + \bar{\chi} Z$ = 24(z+Z) + 2 z(y+y) = xyz+ xyz + xyz+ xyz. = Em (7,6,3,1) = TM (0,2,4,5) = (x+y+z) (x+y+z) (x+y+z) (x+y+z) Express the Boolean function in Pos and soptom (A+B) (B+c) > POS form = AB+B·B+A·C+BC. =AB+AC+BC => SOP torm.

## Implementation of Logic function using gates.

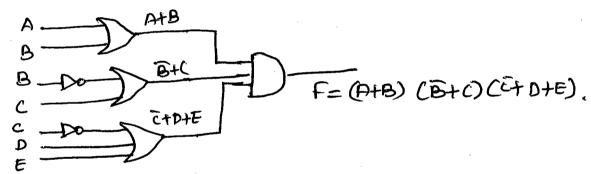
Boolean algebra is used to express the olp of any Combinational network. Such network is implemented using logic gates.

#### Example

Implementation of SOP expression.

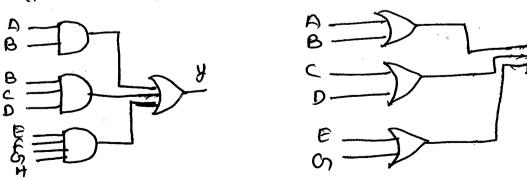


Implementation of Pos expression,

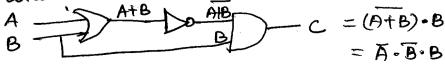


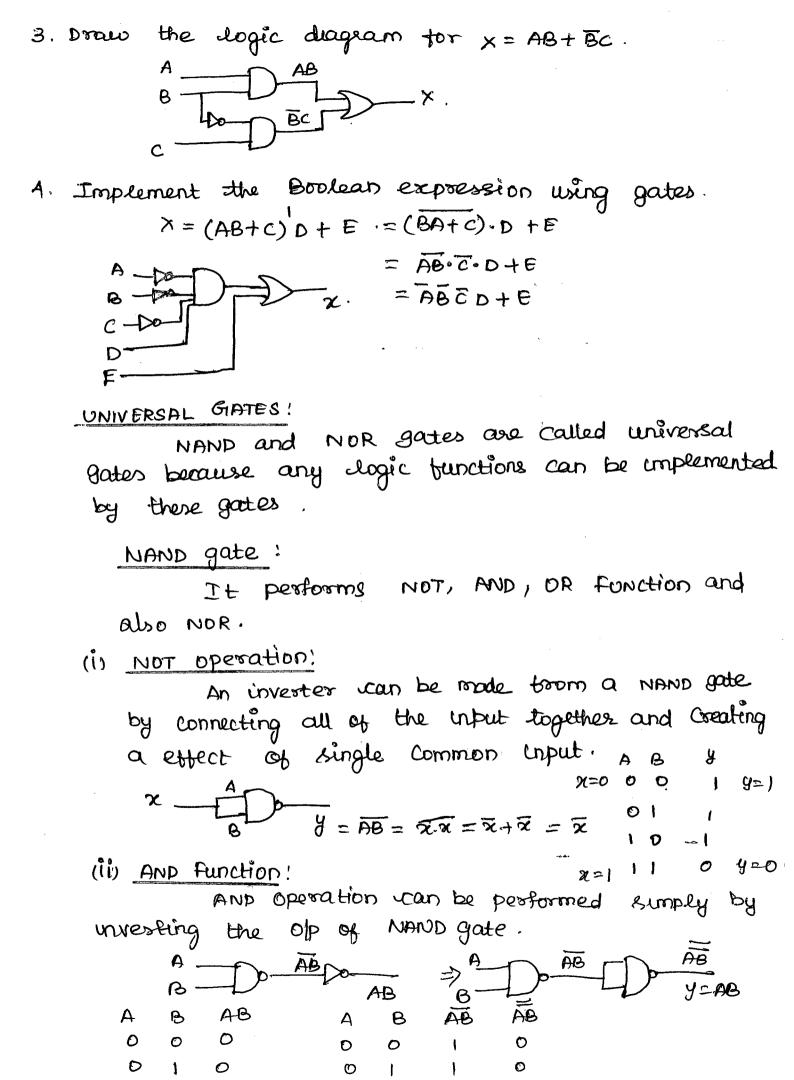
## (1) Implement the expression.

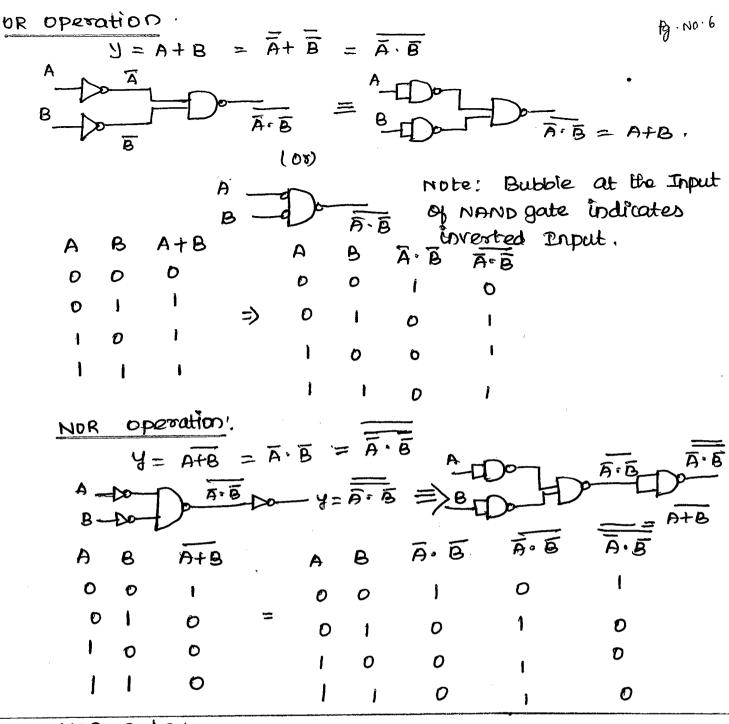
a) AB+BCD+FFGH .



(2) write the Boolean expression for the elp of the system





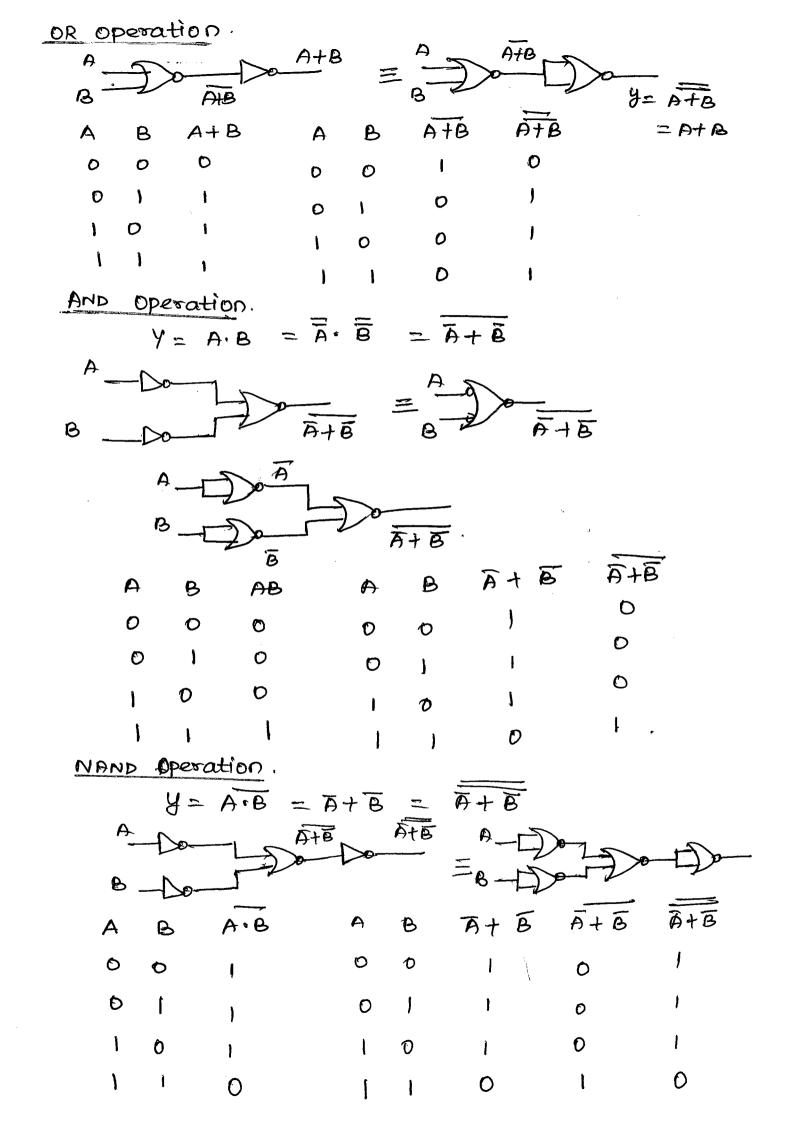


NOR Gate:

NOR gate performs NOT, AND, OR, NAMO operations so called universal gate.

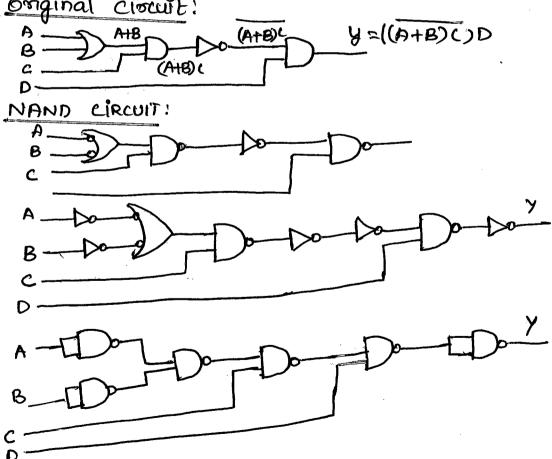
#### NOT function!

An invester can be made from a NOR gate by Connecting all the Inputs together and Greating the effect of single common input. Y= A+B



# Conversion of ANDIDRINOT Logic to NANDINOR Logic using graphical procedure.

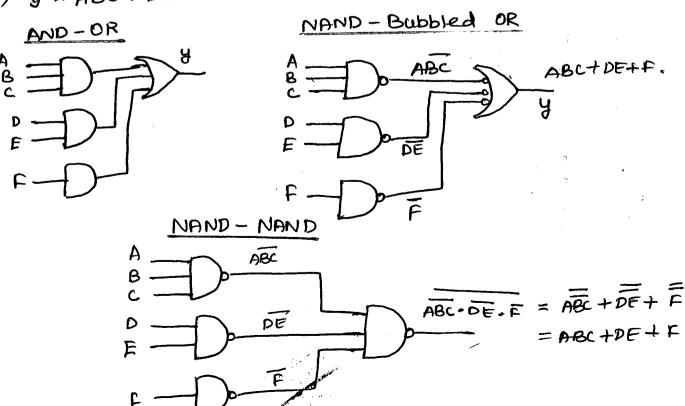
- (1) Draw ANDLOR logic.
- (2) If NAND hardware has been Chorsen, add bubbles on the output of each AND gate; and bubbles on input side to all DR gates.
- (3) If NOR hardware has been choosen, add bubbles on output of each OR gate and bubble on input of each AND gate
- 4. Add or subtract an investes on each line that received a bubble in step2 or 3
- 5. Replace bubbled or by NAND and bubbled AND by NOR
  - 6. Eliminate double inversions.
- \* Boolean expression: ((A+B)C)D original circuit:



## Design of two level Network:

NAND-NAND Implementation.

- \* Requires the expression sumplified to SOP term
- between Ano-or and NAND-NAND Logic is \* Relation explained
  - 1) y = ABC+ DE+F.

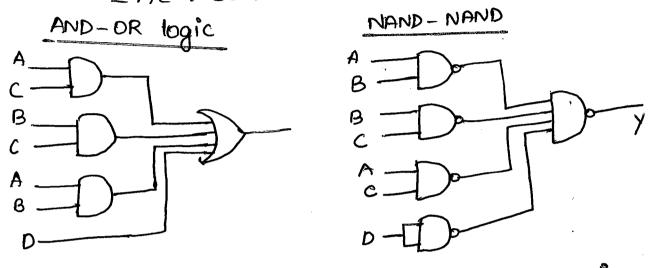


Implement the following Boolean function with NAND-NAND dogic  $Y = AC + ABC + \overline{ABC} + \overline{ABC$ 

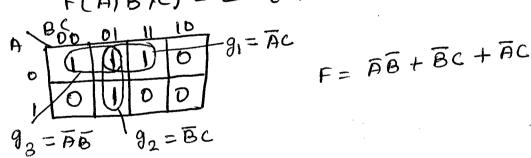
= AC+ ABC + ABC+AB+D.

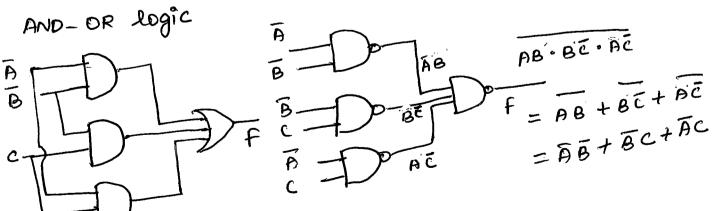
= AC + BCCA+A) +AB+D.

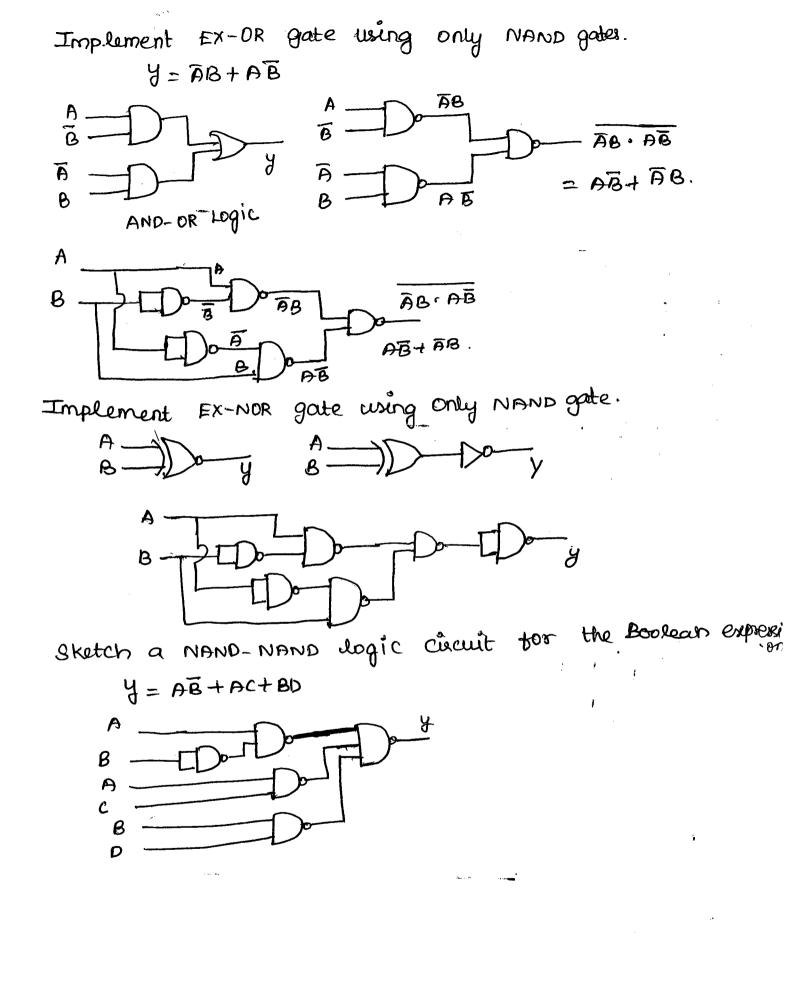
= AC + BC + AB+D.



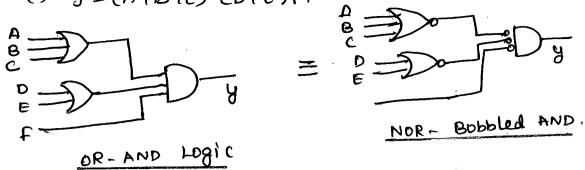
Implement the following Boolean tunction with NAND-NAND logic.

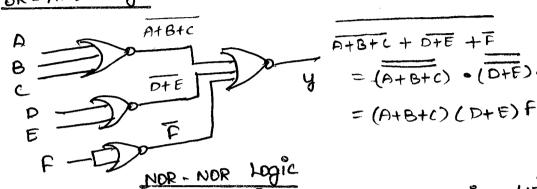












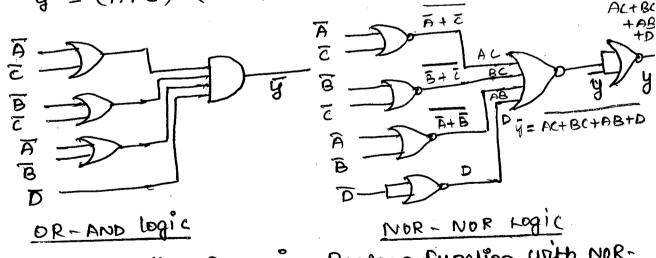
(2) Implement the tollowing Boolean tunction with NOR-NOR logic Y = Ac + Bc + AB + D.

Solution: Convert to Pos torm

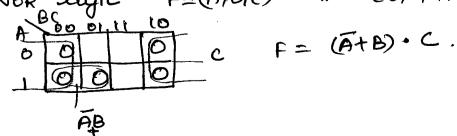
By Dualility \* Change . to +, + to .

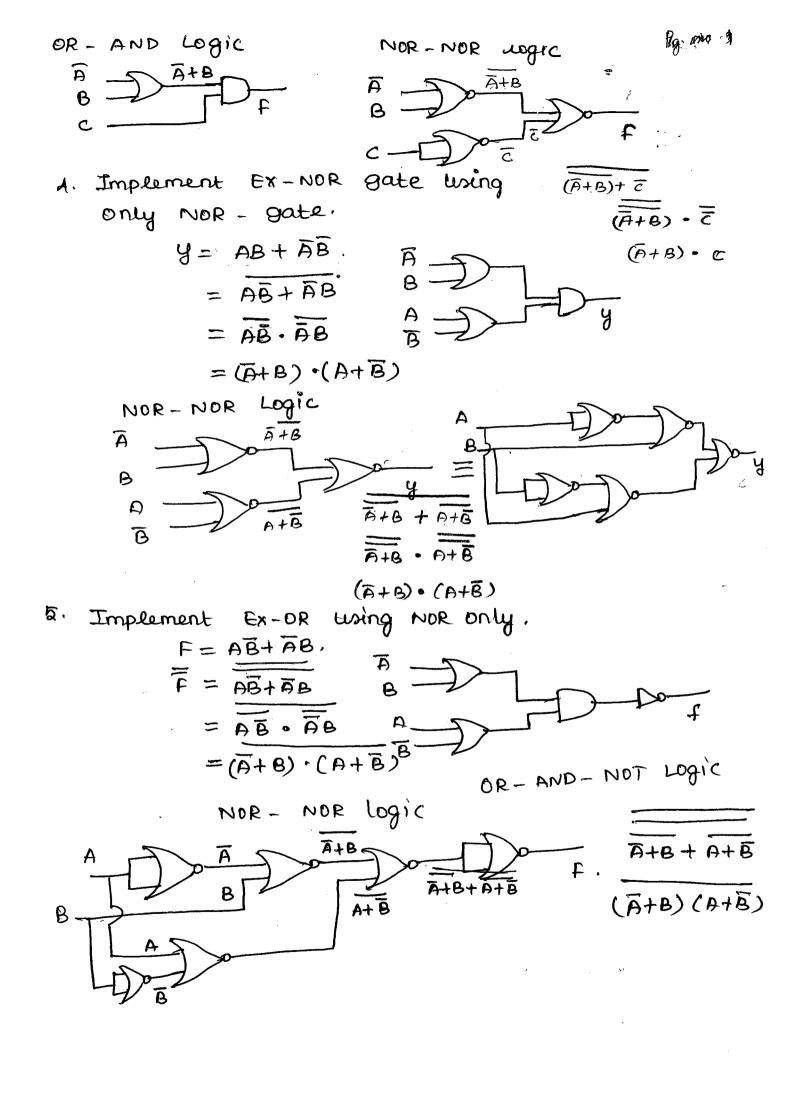
\* Take Complement of the function.

$$\overline{y} = (\overline{A} + \overline{c}) \cdot (\overline{B} + \overline{c}) \cdot (\overline{A} + \overline{B}) + \overline{D}$$

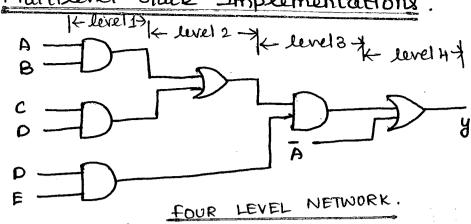


3. Implement the following Boolean function with NOR-NOR Logic f=(A,B,C)=TM(0,2/4/5/6).







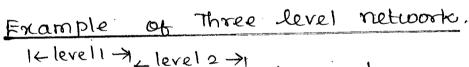


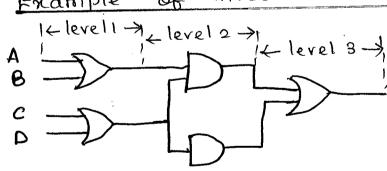
we know that logic gates can be carcaded to get the desired output. The maximum number of gates cascaded in between a network ilp and the output is referred to number of levels of gates.

\* Thus the Boolean function written in sop form and Pos form are the two level gate networks.

\* usually it is assumed that all variables and their complements are available as network inputs.

\* Thus we will not normally count investers Connected directly to input variables as a seperate levels in the network.

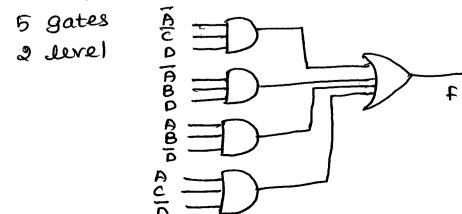




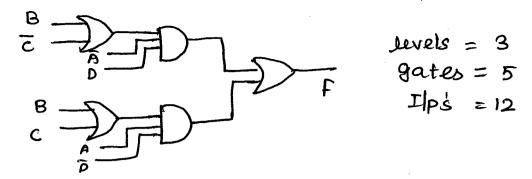
#### Example:

F = ACD + ABD + ABD + ACD

\* 5 gates

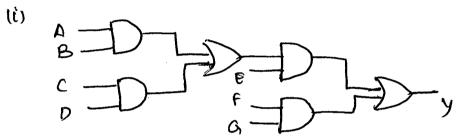


F = ACD + ABD + ACD = AD(B+C) + AD(B+C)



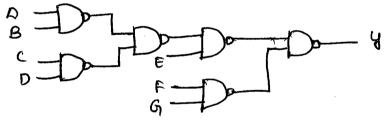
## Multilevel NAND - NOR Implementations

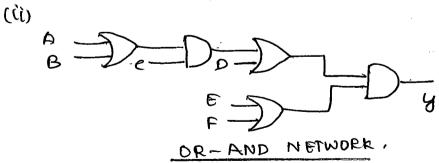
AND- OR Network can be converted to NAND-NAND and also to NOR- NOR Network. From OR- AND



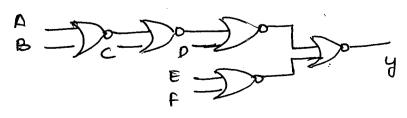
AND-OR- Network.

NAND-NAND Implementations.

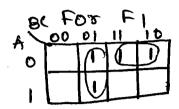




NOR- NOR Implementation.

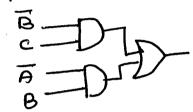


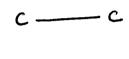
1. Simplify the tollowing functions and draw the logic diagram for the same.

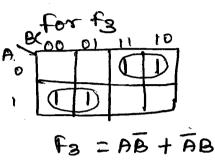


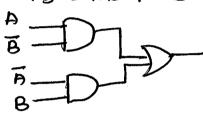




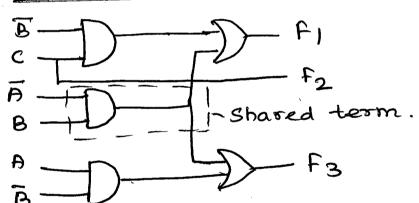






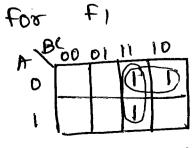


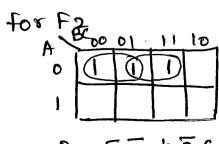
Combined ciscuit



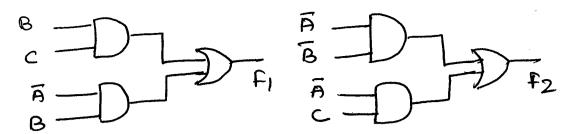
2. Simplify the following functions and draw the logic diagram for the same.

$$F_1 = F(A_1B_1C) = £(2_13_17)$$
  
 $F_2 = F(A_1B_1C) = £(0_11_13).$ 

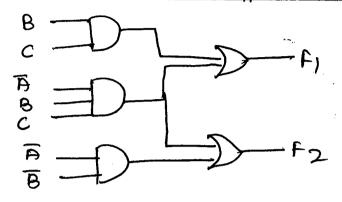




F2= FB+FC



multiple of Considered.

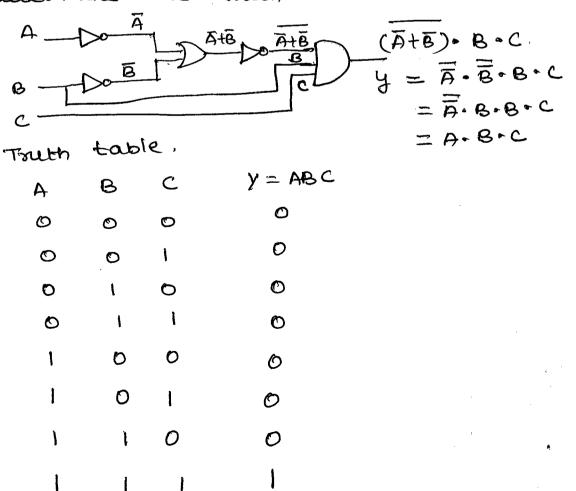


common term has to be found for maltiple output Implementation.

Note: It multiple outputs are Considered together does not gurantee a munimum No. of Ic's.

#### Extra!

1. write a Boolean expression for the output y and determine the truth table.

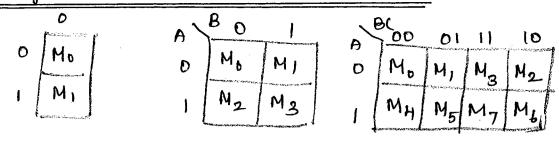


Pg. No. 12 Karnaugh Map Boolean expressions can be simplified by Boolean algebra needs better understanding of Boolean laws, rules and theorems. during simplification we have to predict each sucerive steps. can also simplify the expressions by We mapping also by systamatic approach. 1. one variable 2. Two variable 3. three variable. H. four variable. 2 variable map contains - 2 = 4 cells. 3 variable map Contains - 23 = 8 cells A variable map. 4 variable map contains = 2 = 16 cells. 2 variable mapsion of 11 00 10 01 1 van able map 3 variable map 1 terms are assigned to the cells. \* Product BIBE BC BC BE Q2 Q3 ABG ABG ABG ABG ABED ABED ABED ABOABL PB A ABE ABE ABE ABE A ABO ABO FECD ABLD ABLD AB 1 variable map DEGIDERED PECD ABOD AB ABOO ABOO ABOU ABOU for sop expression. represent ways to BC 00 D 01 11 . 10 00 A WP 0 Mo mi m2 00/ Mo m, m3 m2 Ma Đ W<sup>D</sup> m, Ó 18)<sub>7</sub> m mę 16)<sub>4</sub> m5 M ന്തള m5 m7 B) H m2 01 W12 mis mis 1) terms. map the to Mg Sum

8 8 B	B+C B+C B+C B+C
A A A A B A B A	AIBK AIBK AFBR AIBH
A A A ATB ATB A	AHBY AHBY AFBY
A service of the serv	,
A second	AHB! AHB! AHB! AFB!

	CHD			<u>c</u> +p
AtB	A+B+ C+D	A+B+ C+D	A+B +C+D	A+B+ C+D
			A+B+ C+ D	AT B TC+D
PHB	45 0+0	A+B +c+B	A+B+ C+D	A+B+
A+B	At Bt C+D	A+B++(+	At B tct	A+B+ C+D

ways to represent Pos expression.



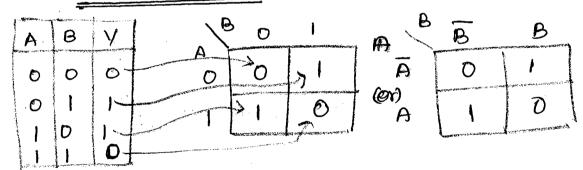
200	Pob	01	11	lo
ו פוניו	l L	M,	Mg	M <sub>2</sub>
0)	M <sub>4</sub>	M <sub>5</sub>	Mg	Mb
()	M <sub>12</sub>	Mis	M <sub>15</sub>	M <sub>14</sub>
lo	Mg	Mq	Mij	Mio

Plotting of K-Map.

- (1) for muth table
- @ Sop expression.
- (8) Pos expression.

#### Touth Table!

(1) & - Variable.



## (a) 3 Variable.

A	В	1 c	У	] A	00	ol	11	vo
0	0	0	0-	n	0	0	9	1
0	0	1	0	And the second s	migration of the second of the			0
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A	В	С	D	У	AB \	CD	0)	11	w
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b	Ø	1	Ь	O	0)	7	١,	}	0
0	٥	ſ	1	1-	.11	71		to profingular or	
D	1	0	O	1		7 1	0	7	2
Ø	ļ	D	1	1.	yo .	0	Spark A		
0		1	O	0			/ O	0	71
O	1	}	1	10					and the second section of the second
١	-0	O	0	6	/ /05	/	, and the second second		
	0	0		0	AB X	73	CP	CD	CD
1	Ø		O	1/	PB	1	0	1	0
	O		The state of the s	0	FB	1	1	I	0
. <b>1</b>	140	0.	0	<b>†</b>	84	Î	0	1	1
1	1	O	1	0		•			
ľ	1	1	O	1	A B	0	0	0	
1	1	1	1	1	<u>}</u>			- design of the second	A

Representation of sop on k-map.

H Boolean expressions in the sop form can be plotted on k-map by placing 1' in the map for the terms in the expression and remaining yoro.

Example; Plot Boolean expression  $Y = ABC + ABC + \overline{ABC}$  on the K-map.

ABC= 001. Solution: ABC = 110 ) ABC = 111 01 11 10 Bc BC  $\sigma_{\mathcal{G}}$ BC BT A BBC F (or) D ABC A

Plot Boolean expression

Y= ABCD + ABCD + ABCD + ABCD + ABCD on k-map.

<b>*</b>	OD SCD	01	11	w
00	O	0	D	0
01	AB CD	0	0	Fe vo
11	0	A870	0	O
10	0	0	. 1	AB CD

ABCD = 0100 ABCD = 1010 ABCD = 1011 ABCD = 1101

Representation of pos on k-map

\* Sum terms in the expression one represented by b' in kmap.

\* remaining by '1'

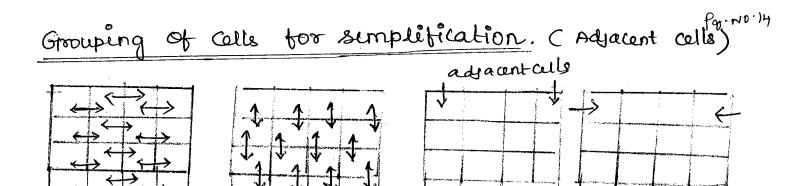
Frample: plot Boolean expression y=(A+B+c)(A+B+c)
(A+B+c) (A+B+c) on the K-map.

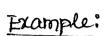
golution	<u> </u>			) a	
A+ B+c	=>	010	$\Rightarrow$ M <sub>2</sub>	A 00 01 11	LO Manageral
A+B+C	=>	011	⇒ M3	0 1 0 0	0
A+B+c	=>	110	=> M6		0
A+B+c	=>	001	=> M ,	emperation community and a second community of the commun	rije sikaj rati krisad

Plot Boolean expression y = (A+B+C+D)(A+B+C+D)(A+B+C+D)(A+B+C+D)(A+B+C+D), on K-map.

A+B+C+D = 0001  $\Rightarrow$  M<sub>1</sub> A+B+C+D  $\Rightarrow$  1101  $\Rightarrow$  M<sub>18</sub> A+B+C+D  $\Rightarrow$  0110  $\Rightarrow$  M<sub>6</sub> A+B+C+D  $\Rightarrow$  1110  $\Rightarrow$  M<sub>14</sub> A+B+C+D  $\Rightarrow$ 

A+B+c+D = 0011 => Mg
AB
00 1 0 0 1
11 1 0
11 1 0



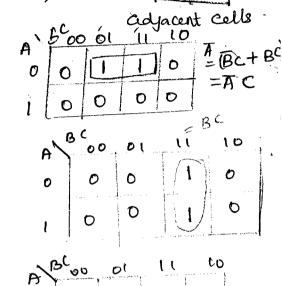


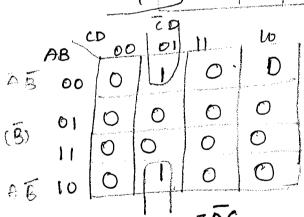
(a) 
$$Y = \overline{ABC} + \overline{ABC}$$
  
=  $\overline{AC}(B + \overline{B})$   
=  $\overline{AC}$ 

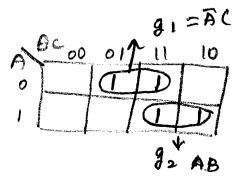
d) 
$$Y = \overline{A}\overline{B}\overline{C}D + A\overline{B}\overline{C}D$$
.  
 $= \overline{B}\overline{C}D + A\overline{B}\overline{C}D$ .

= AC+BC.

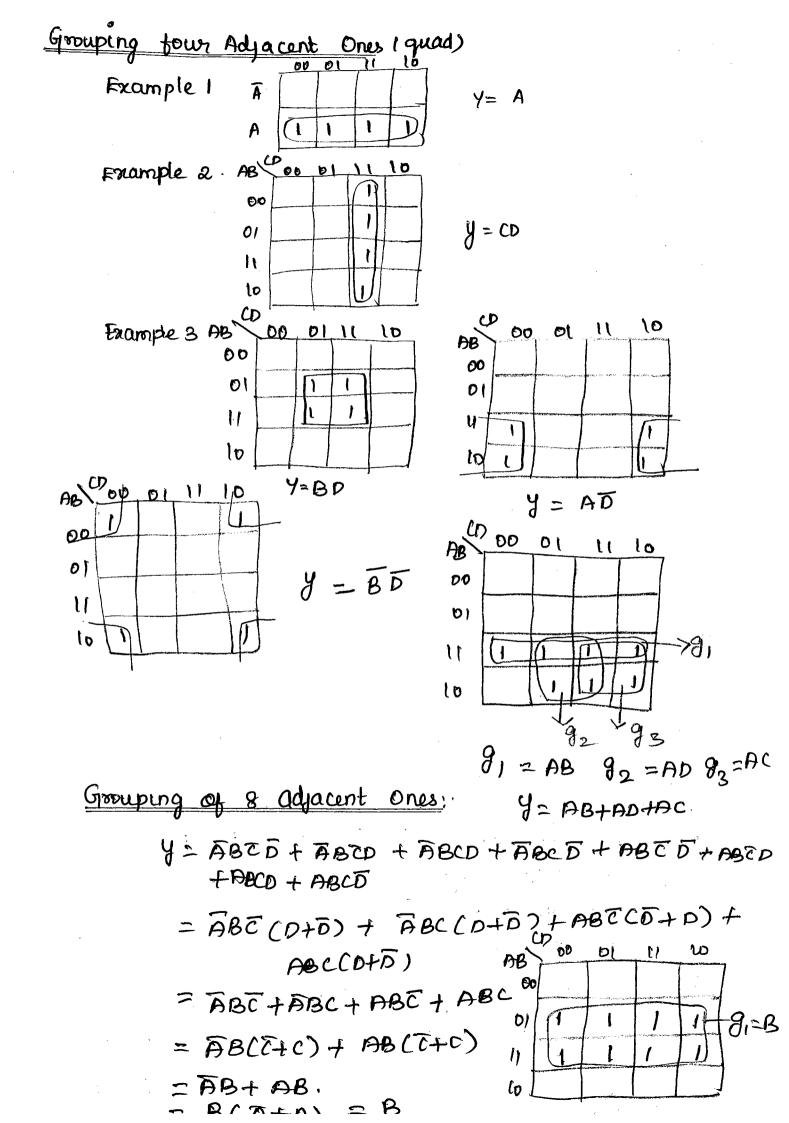
= 001 +011 + 111 = ABC+ ABC A Beo 01 8211 W + ABCHARC AO A+A=A = AC(B+B) + ABC AI = AC + ABC + ABC = Ac + BCCA+A)

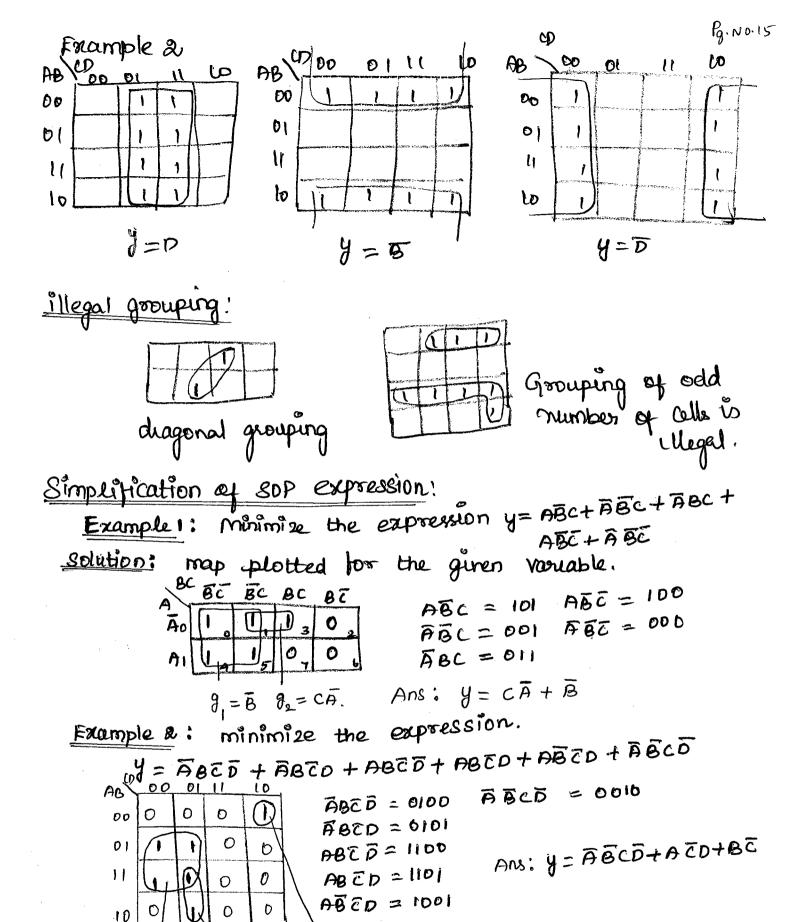






 $\theta_1 = BC$ 





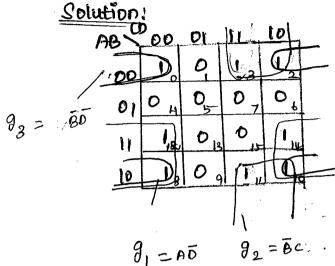
g=ACD 93= ABCD

9,=BC

#### Example 3:

Reduce the tollowing four variables function to its

Y = ABCO+ ABCO + ABCO+ ABCO+ ABCO+ ABCO+ ABCO

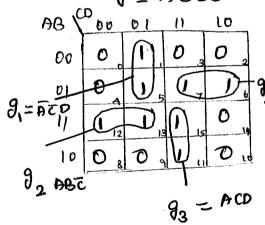


Solution: Y= BO+ AD+ BC.

#### Example 4!

Reduce the tollowing function to its minimum sum of products form.

Y= ABCO+ ABCO+ ABCO+ ABCO+ ABCO+



ABCD + ABCD

PBCD = 000 |= | ABED = 110 |= 13

D 9= 000 |= | ABED = 1111 = 15

D 1 ABCD = 011 |= | ABCD = 1011 = 11

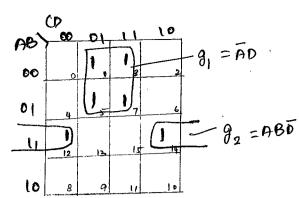
ABCD = 0110 = 12

D 10 ABCD = 1100 = 12

Answer: FBC + ACD + ABC + ACD

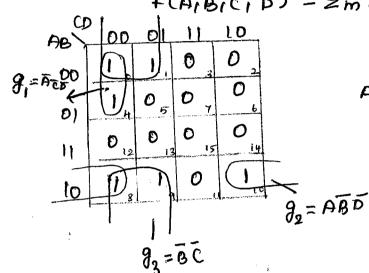
Example to: Reduce the tollowing function using k-map technique and implement using gates:

FLABCD) = ABD + ABDD + ABCD.



Solution: y= ABD+ AD.

6. Reduce the following function using k-map techniques,  $F(A_1B_1C_1D) = \sum_{m} (0,11,4,8,9,10).$ 



Ans: F(A,B,C,D) = ACD+BC+ABD

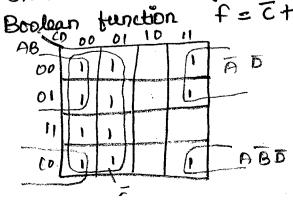
7. Plot the tollowing Bodean function on a koomangh map and simplify it.

F(w/x/4/z)= 2 (0/)/2/4/5/6/8/9/12/13/14).

C	, 1	- (m)	X 9/	7)	_	
CI	00	01	11	10	-	
60	II.	1	0 3	12	<del>&gt;</del> 92=	ĀĎ
01	可	15	0 7	1 6		<b>.</b> =
<u>I</u>	1/8	1	D 15	14		3=85
10	18	والم	0	0 10		
	1	9=	$\overline{c}$			<b>6</b>

Answer = 
$$8\overline{0} + \overline{A}\overline{0} + \overline{C}$$
  
=  $2\overline{2} + \overline{L}\overline{0}\overline{C} + \overline{y}$ 

8. Show the Karnaugh map with the encurcled groups for the Boolean function  $f = \overline{C} + \overline{A} \, \overline{D} + \overline{A} \, \overline{B} \, \overline{D}$ .



### Incompletely specified functions (Don't fare terms)

\* In some logic circuits, certain input vonditions never

occur, therefore the corresponding output never appears.

\* The output level is not defined, either high or Low.

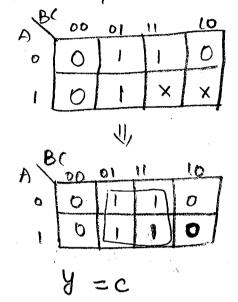
d'indicated by 'x' or d' called don't care Conditions or incompletely specified function.

Example:

FCAIBICID) = &m(1/2/4/7/8)+ d(10,11,12,13/14,15). Minimization of incompletely specified functions.

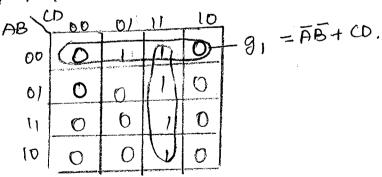
at don't care condition, it is either "O' or 1".

		1	
A	В	C	У
D	0	O	0
0	0	1	1
0	ì	0	0
0	)	ĵ	)
)	O	0	0
)	0	١	ı
1	)	0	×
	1	1	×

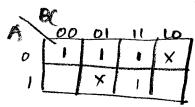


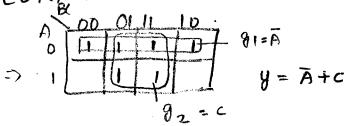
Example: 1. Find the reduced SOP form of the following function. FCAIBICID) = &m(1,3,7,11,15)+2d(0,2,4).

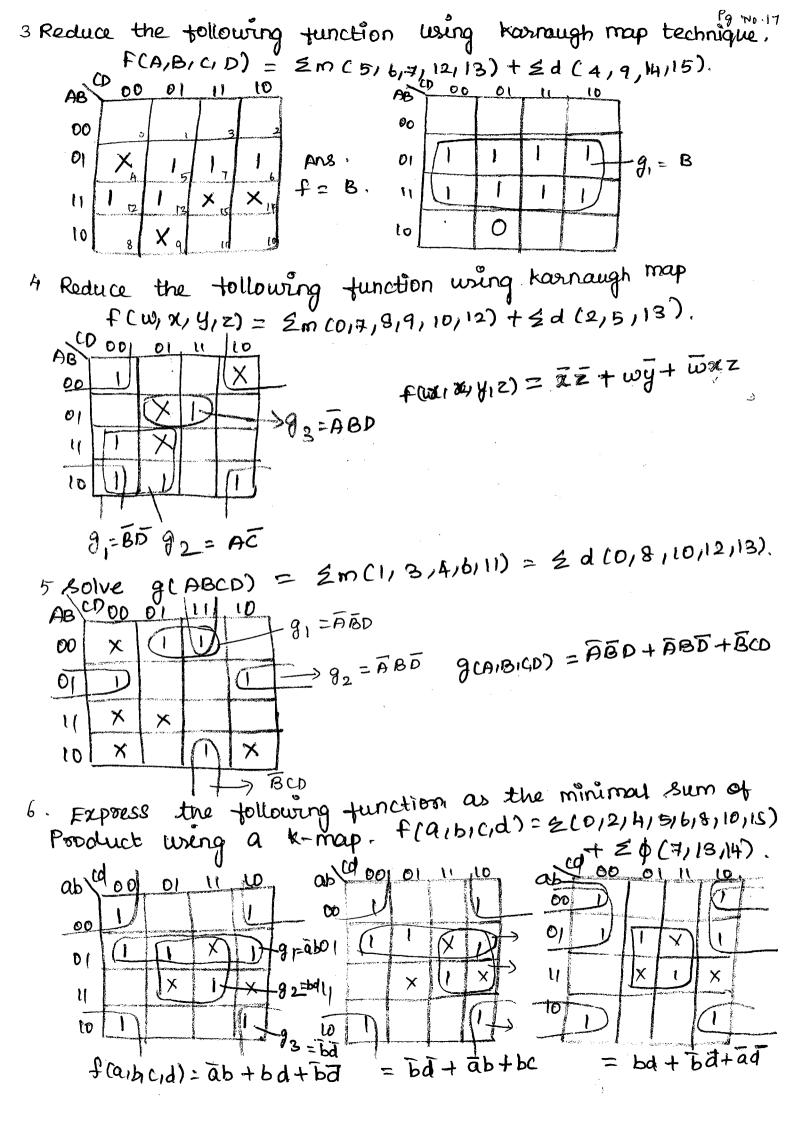
- 0			•	
AR	D 00	ot	11	lo
	X	1		×
00	C			2
٥Į	X	5		, d
Ŋ	l2	13	115	14
to	8	9	1 4	lo
lj to,	12	9	1 11	10



2 Reduce the following function using kornaughs map FCAIBIC) = 2m Coil, 3,7) + 2d(2,5).







#### Note:

\* After grouping the cells, the sum terms which appear in the k-map are called prime implicants group.

\* A circuit designer is tree to make the Olp of any don't care Conditions either a o' ora 'i' in order to produce the simplest olp expression.

Simplification of POS Expressions.

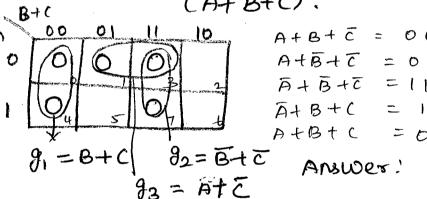
marked as zero in K-map.

a grouping of or as to be done.

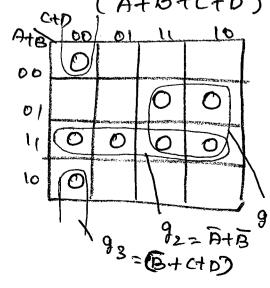
## Example:

1. minimize the expression

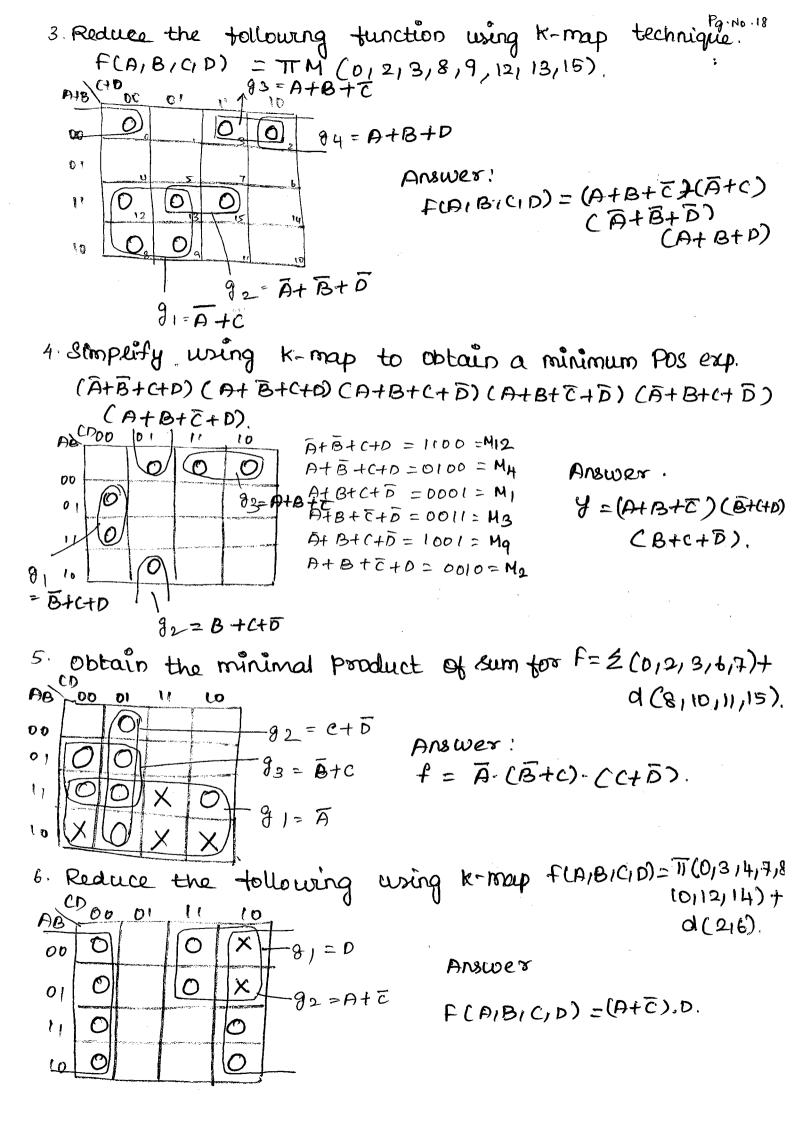
Y=(A+B+C)(A+B+C)(A+B+C)(A+B+C)
(A+B+C).



2. Minimize the following expressions in POS torm.  $U = (\overline{A} + \overline{B} + C + D) (\overline{A} + \overline{B} + \overline{C} + \overline{D}) (\overline{A} + \overline{B} + C + \overline{D}) (\overline{A} + \overline{D} + \overline{D}) (\overline{A} + \overline{D$ 



Ans: y= (B+c+D) (B+c)
(B+B)



Rules for simplifying clogic function using k-map.

- 1. group should not enclude any cell containing a xow.
- 2. The number of cells in a group must be a power of 2, such as 1,211,8 00 lb.
- 3. Group may be horizontal, Vertical but not diagonal.
- 4. Cell Containing 1 must be included in atteast one group
- 5. Groups may overlap.
- 6. Each group should be as large as possible to get maximum simplification.
- I. Groups may be wrapped around the map. ]
- 8. Cell may be grouped more than once.
- 9. We need not group all don't care cells,
- 10. the above rules are for SOP and for Pas except "o" she sest are same.

## Limitations of k-Map:

- of Convenient as long as the no of Variables is is orb \* for 7,8,9 variables it is a impossible
- tasks.
  - \* Another unportant point is that the K-map technique simplification is manual technique and simplification process is heavily depends on human abilities.
    - \* TO meet this need we go for Tabulation method to simplify called
      - "Quine Mc cluskey or tabular Method".

## Quine-Mccluskey Method of Minimization.

1. Simplify the following Boolean function by using a Quine Mccluskey method.

F(A/B/C/D) = 2m (012/3/6/7/8,10/12/13).

	solution	•									
	minterms	Binan	y representa	tion.							
	$m_o$	00	•	mo		0,2	- 00-0				
	$m_2$	00	10	Wo	0010	018	-000				
	m <sub>g</sub>	00	[ ]	mg	1000	213	001-				
	w <sup>e</sup>	011	O	mg	0011	216	0-10				
	ws w <sup>1</sup>	0111		wf	01100	2110	_ 010				
	W <sup>IP</sup>	100		$m_{lb}$	1010	8,10	10-0				
		$\frac{m_{12}}{m_{13}}$ 1100		m12	1100	8112	1-00				
	ml3			m4	0111	317	m 1 1				
	Ü	110	1	w 13	1101	617					
	•					1211	3 1.0-				
	0,2,8,10 -0-0				0/2/8/10 -0-0						
	21316,7 0-1-				2,3,6,7 0-1-						
				0			1 -00				
	21316	0 - 1			112	110_					
	8/12		1 - 00	)	1 2	1 13					
	12/1	3	110_	_							
			Ö	2 3	6 7	8 10	12 13				
	8,12	ACD				*	$\times$ $\times$				
١	2113 /	ABZ					<b>x (r)</b>				
0,	218110	BD	X	Х	,	×, (X)					
) i	31617	ÃС		ХХ	$\otimes$						
21	2101			_	_						

The final expression is

F(A,B,C,D) = ABC + BD + AC

```
the expression using McCluskey Method.
2. Minimize
     Y = ABOD + ABOD + ABOD + ABOD + ABOD + ABOD + ABOD +
        = 0100 + 0101 + 1100 + 1101 + 1001 + 0010.
        = m4 + m5 + m12+ m13+ m9 + m2.
                                                 4,5,12,13 -10-
                                        010-
                                (4,5)
                        0100
  m4
        0100
                  m4
                                                (411215113)-10-
                                        110 -
                                 (12113)
                         0010
  ME
        0101
                  W 2
                                        _ 100
                                                   9113 1-01
                                 (H112)
                         1100
                  m12
                                        _ 101
  m12
        1100
                                 (5113)
                         0101
                   m 5
  MIZ
                                 (9,13) 1-01
        1101
                        1001
                   ma
   ma
         1001
                         1101
                   m13
   m2
         0010
                             BE
                 _10 -
     415/12/13
                             ACD
                  1 - 01
        9,13
                             ABCD
                  0010
          2
                                              13
                                         12
                                     9
                                5
                            4
                       2
                                               X
                                X
          415,12,13
    BE
                                              X
            9,13
   ACD
   ABCD
              expression
                             is
        Final
              ABCD + ACD + BC
            the following equation using the tabular method.
3. Reduce
          F = mo+m2+m3+m5+m8+m10+m11+m13.
                                         00-0
     30
                                  (012)
                           0000
           0000
                    200
                                         -000
                                  (018)
                           0010
     m_2
                    m 2
           0010
                                          001-
                                  (213)
                           1000
                     m &
     m3
           0011
                                          -010
                                  (2110)
                           0011
                     m 3
     m 5
           0101
                                          10-0
                           0101
                                  (8,10)
                     m 5
     W &
           1000
                     m 10
                           1010
                                  (10,11)
                                           101-
                      11 0
     10
           1010
                           1011
                                   (311)
                                           -011
                      M13
```

1101

-101

(5)13)

110

1011

```
Pg. No. 20
                - 0 - 0
  01218,10
                             0,218,10
                                            _ 0_0
                                            -01-
  0,8,2,10
                -0-0
                              213,10,11
                                            -011
  213/10/11
                - 01 -
                                   3,11
                                             _ 101
                - 01L
   ` · · (3, 11)
                                    5113
     (5/13)
                -101
                                                  13
                                         10
                                              11
                                     8
                                5
                           3
                       2
                  0
 BD (01218110) ~ (X)
                                    \bigotimes
                                          X
                       χ
 BC (213,10,11)
                                             X
                            X
                                          Х
                        X
 BCD (3111)
                                              X
                             X
 BCD (5/13) ~
                                                   (X)
 minterns not-covered
                                             > selected
                                 2 variables
                        01-
                                                 (minèmem Variable)
    (2, 3, 10, 11)
                                 3 variables.
                      - 011
       (3,11)
         expression is
              = BC+BD+ BCD.
4. Simplify F(w,x,y,z) = &m(1,2,3,5,9,12,14,15)+&d(4,8,11),
                                         00-1- (1,3,9,11)-0-1
                                  (113)
                         0001
                                                  (11913111) - 0-1
                   mi
         0001
  mI
                                         0-01
                                  (115)
                          0010
                   m 2
         0010
                                          -001~
  m2
                                   (119)
                          0100
                   4mb
         100
   m3
                                         001-
                           1000
                   d ma
                                   (213)
   m 5
          0101
                                   (415) 010-
                           0011
                    നദ
          1001
   ma
                                    (4,12) -100
                           0101
                     m 5
   W15
          1100
                           1001
                     ma
                                           100 -
                                    (819)
   W14
                            1100
                     m12
                                    (8112) $ 100-
          1110
                            1110
                     MIH
                                           -011
    m 15
          111)
                                     (3,11)
                            1011
                     dmil
                                     (11,15) 1-11
   dm4
          0100
                            1111
                      m15
                                     (14, 15) 111 -
   dma
          1000
                                      (12114) 11-0
   dm11
           1011
                                            10-1-
```

(9,11)

```
(1, 3, 9, 11)
               -0-1/
               0-01
  (1,5)
                0011 -
  (213)
                010 -
  (415)
               - 100
  (H112)
                100 -
  (9,9)
                1-00
  (8112)
                1112
  (H15)
                f(I) = t_{\mathbb{R}^n}
  (11/15)
                11-0/
 (121 14)
```

,		allege and the control of the contro			رد پرون پرون در	- Section of the Policy Co. Section 100 Se	والمعادف المارات المارات والمسارات	ani anangkani	nanda er erendende	d	cl	d
		1	2	3	5	9	12	14	15	4	8	1)
	1,3,9,11	X	the section of the fourth that	X	age control of the	×					-	*
V	115	Х		}	*	<u>.</u>						
~	213		$\otimes$	×		ž.				-	and the state of t	
	4,5	; ;			Х				,	X		
	4,12						$\bigotimes$			X	· · · · · · · · · · · · · · · · · · ·	
	8,9		- -			X	A CONTRACTOR OF THE CONTRACTOR	To the sales			X	
	8112	1 2 2	i i i i i i i i i i i i i i i i i i i			v.						
V	14,15					3	(8)	X	X		×	
	11/15		As controlled to			1 3 &			×			X
~	121.14	\ \			)		$\otimes$	×				

```
(2,3) for final expression.
                        select
2 -> checked once
                   30
```

1) " twice so select 
$$(1,3,9,11)$$
,  $(1,5)$  for final exp.

already selected

already selected 5 ->

8 elected already 9 -)

12 -> Checked thiree So select (14,12), not (8,112) or (4,112) 4, 8, -> don't car

Select (14115) only not (11,15) 15 -> Udon't care.

The final expression is

F = BD + ACD + ABC + ABC

1. Simplify tollowing logical expression using karnaugh maps. Y=ABC+ABC+ABC+ABC Ans: AB+C

2. Simplify the tollowing function.

f, (AIBICID) = 2m (DI 3/5/6/9/10/12/15) Ans (ABB) O(COD)

3. Simplify the following function.

f3(A/BCD) = Zm(0,1,2,3,1,12,14,15) Ans: AB+ABD+ ACD

4. Simplify the following using k-map.

X = AB+ ABC+ ABC+ ABC. ANS: AC+B= X

5. Simplify the following function using k-map technique.

F(A1B1C1D) = & (012/3/6,7/8,10,12/13) Ans ABC+BB+AC

simplify the following Boolean function using 4- variable map f(w,x,y,2)=至(213,10,11,12,13,14,15) Ans: wx+ xy.

7. Simplify the following surtching function using k-map FCAIBICID) = & CO15,71819/10/11/14/15) + \$C14/13)

8. Determine the minimal sum of product form of +(w, x, y,z) = Em(4,5,7,12,14,15)+d(3,8,10).

Ans: Wxy+xyz+wz

9. simplify the following Boolean tunction for minimum Pos form.

FCW121912) = TM (415/6/7/8/12) +d(1/2/3/9/11/14)

Ans: (w+x) (w+y+z)

10. Reduce using k-map.

F(A1B1C) = TM(0/172/3/4/7)

Ans = A. (B+C). (B+c)

11. simplify the following using Boolean algebra. f(x1412) = TM(315,7). AN: (y+2)(x+2)

12 Reduce using Tabulation method

- (1) f(A/B/C/D) = \( \int (0/1,2,3,4/6,8/10,12) \),

  ANS! \( \overline{A} = \overline{D} \) ANS! AB+D
- (2) F(A/B/C) = & (0/1/4/6/7)

(3) FCAIBICID) = & C21416, 7,9/11,132 + ARC+ ABD+ ACD+

- H. FCA,B, (ID) = &(D)1,2,3,5,6,7,8,11,13).
  AND: AB+AD+AC+BCD+BCD+BCD+BCD
- 5. FCA, B1(1D) = 2m (21315/7,19,11,12/13,14,15)
  AND ABC # 2000+ BD+ AD+AB.
- b. F(A/B/CID) = &m(0/1/2/3/6/7/13/15)
- AM: ABD+ AB+ AC FCA1B1C1D) = Sm (0/1/2/3/4/7/9/10)
- 8. CIA.A.C.D) AM: ACD+BCD+BCD+AB
- 8. FLAIBICIP) = Em (011,3,6,7,8,9,13,15)
- 9. FCA/BICID) = &m (3,415,7,9,13)14,15)
- 10. F= mo+m2+m4+m8+m9+m9+m10+m11+m12+m13
  Past-A,
  Ans: BD+ABC+ACD+ACD+ABC
  Ans: BD+ABC+ACD+ABC
  Ans: BD+ABC+ACD+ABC
  Ans: BD+ABC+ACD+ABC
- 1. Give the steps for simplification of sop expression.
- 2. what do you mean by essential prime implicants
- 3. what is don't care conditions?
- A. Give the steps for sumplification of POS exp.
- 5. State the rules for K-map sumplification.
- 6. Explain the Limitations of Ravnaugh map.
- 7. State the advantage and disadvantage of Qunine McClurky method of simplification.

## Possible Questions:

- 1. Define switching function
- 2. Define literal, product term and sum term.
- 3. Explain sum of product form.
- 4. What do you mean by standard sop and pos form.
- 5. Explain how to Convert sop or Pos expressions in their standard forms.
- 6. what do you mean by minterms and maxterms.
- 1. Express  $F_1$  in standard sop form  $F_1 = AB + CD + ABC$
- 2. Determine the canonical sop form of  $F(x_1y_1z) = (xy+\overline{z}) (y+x\overline{z})$
- 3. Convert the given expression in standard Pos form  $F(P, Q, R) = (P + \overline{Q}) (P + R)$



Questions:

- 1. Implement the following function with NAND gates  $f(x_1y_1z) = \mathcal{E}(o_1b)$
- 2. Design a logic circuit to simulate the function F(A/B/C) = A(B+C) by using only NAND gates.
- 3. Realize is AND gotte (ii) NOR gate wing only NAND gates.
- 4. Realize i or gate ii Ex-OR gate using NAND gates.
- 5. Realize (i) OR gate (ii) AND gate using only NOR gates.
- 6. write a short note on multilexel NAND and NOR Implementations.
- 7. write a short note on mulioutput gate Implementation.
- 8. show that y=ABC can be implemented with One two input NOR and one two Input NAND gate



