Zscaler

1. Why Zscaler

Zscaler (/ˈziːˌskeɪlər/) is an **information security company** that provides **network security services** all based on the cloud. Its flagship services, **Zscaler Internet Access**, and **Zscaler Private Access** create fast, secure connections between users and applications, regardless of device, location, or network. Glassdoor rating of 4.7/5.0 + 98% CEO Approval is incredible.

As a new graduate student, I'm always fascinated by new technologies. Working at Zscaler would have opportunities to solve the most challenging and interesting problems on Internet Security. I could learn a lot and making things with my team because I know the culture in Zscaler is amazing. Moreover, the growth of this company is pretty impressive. I believe there will be a lot of growth opportunities for engineers working at Zscaler.

2. Interview Question

1. Integer palindrome

Determine whether an integer is a palindrome.

```
if 1 != r:
    return False

x %= div

x //= 10
    div /= 100

return True
```

2. Reverse Linked List

3. Reverse Bits

```
1. class Solution:
2.  # @param n, an integer
3.  # @return an integer
4.  def reverseBits(self, n):
5.     res = 0
6.     for i in range(32):
7.         res = res << 1
8.         res = res | (n >> i & 1)
9.     return res
```

4. First Missing Positive

```
class Solution(object):
   def firstMissingPositive(self, nums):
       :type nums: List[int]
       :rtype: int
       11 11 11
       n, i = len(nums), 0
       if n == 0: return 1
       while i < n:
           # 获取当前位置的数据,减去1是为了得到 在list要插入的位置
           w = nums[i] - 1
           # 0 < nums[i] <= n 判断是否出界
           # nums[i] != nums[w] 如果相等或者是本身就没必要替换了,避免死循环
           if 0 < nums[i] <= n and nums[i] != nums[w]:</pre>
               nums[i], nums[w] = nums[w], nums[i]
           else:
           # 当前位置上的数,没有找到合适的位置,进行下一个位置
               i += 1
       for i in range(n): # 遍历返回
           if i + 1 != nums[i]:
              return i + 1
       return (n + 1)
```

5. Minimum spanning tree

5.1 Prim Algorithm

Prim算法:设图G = (V, E), 其生成树的顶点集合为U。

- ①、把v0放入U。
- ②、在所有 $u \in U$, $v \in V-U$ 的边 $(u, v) \in E$ 中找一条最小权值的边,加入生成树。

③、把②找到的边的v加入U集合。如果U集合已有n个元素,则结束,否则继续执行②。

```
# Prim's Minimum Spanning Tree (MST) algorithm.
# The program is for adjacency matrix representation of the graph
import sys
class Graph():
    def init (self, vertices):
        self.V = vertices
        self.graph = [[0 for column in range(vertices)]
                    for row in range(vertices)]
 # A utility function to print the constructed MST stored in parent[]
    def printMST(self, parent):
        print "Edge \tWeight"
        for i in range(1, self.V):
            print parent[i], "-",i,"\t", self.graph[i][ parent[i] ]
    # A utility function to find the vertex with
    # minimum distance value, from the set of vertices
    # not yet included in shortest path tree
    def minKey(self, key, mstSet):
        # Initilaize min value
        min = sys.maxint
        for v in range(self.V):
            if key[v] < min and mstSet[v] == False:</pre>
                min = key[v]
                min index = v
        return min index
    # Function to construct and print MST for a graph
    # represented using adjacency matrix representation
    def primMST(self):
        #Key values used to pick minimum weight edge in cut
        key = [sys.maxint] * self.V
        parent = [None] * self.V # Array to store constructed MST
        # Make key 0 so that this vertex is picked as first vertex
        key[0] = 0
        mstSet = [False] * self.V
```

```
parent[0] = -1 # First node is always the root of
         for cout in range(self.V):
              # Pick the minimum distance vertex from
              # the set of vertices not yet processed.
              # u is always equal to src in first iteration
             u = self.minKey(key, mstSet)
              # Put the minimum distance vertex in
              # the shortest path tree
             mstSet[u] = True
              # Update dist value of the adjacent vertices
              # of the picked vertex only if the current
              # distance is greater than new distance and
              # the vertex in not in the shotest path tree
             for v in range(self.V):
                  # graph[u][v] is non zero only for adjacent vertices of
 m
                  \# mstSet[v] is false for vertices not yet included in M
 ST
                  # Update the key only if graph[u][v] is smaller than ke
 y[v]
                 if self.graph[u][v] > 0 and mstSet[v] == False and key[v]
 ] > self.graph[u][v]:
                          key[v] = self.graph[u][v]
                          parent[v] = u
         self.printMST(parent)
q = Graph(5)
 g.graph = [0, 2, 0, 6, 0],
             [2, 0, 3, 8, 5],
              [0, 3, 0, 0, 7],
              [6, 8, 0, 0, 9],
              [0, 5, 7, 9, 0]]
 g.primMST();
```

5.2 Kruskal's Algorithm

- 1. 把图中的所有边按代价从小到大排序;
- 2. 把图中的n个顶点看成独立的n棵树组成的森林;

- 3. 按权值从小到大选择边,所选的边连接的两个顶点ui,viui,vi,应属于两颗不同的树,则成为最小生成树的一条边,并将这两颗树合并作为一颗树。
- 4. 重复(3),直到所有顶点都在一颗树内或者有n-1条边为止

```
from collections import defaultdict
#Class to represent a graph
class Graph:
    def init (self, vertices):
        self.V= vertices #No. of vertices
        self.graph = [] # default dictionary
                                # to store graph
   # function to add an edge to graph
   def addEdge(self,u,v,w):
        self.graph.append([u,v,w])
    # A utility function to find set of an element i
    # (uses path compression technique)
   def find(self, parent, i):
        if parent[i] == i:
            return i
        return self.find(parent, parent[i])
    # A function that does union of two sets of x and y
    # (uses union by rank)
    def union(self, parent, rank, x, y):
        xroot = self.find(parent, x)
        yroot = self.find(parent, y)
        # Attach smaller rank tree under root of
        # high rank tree (Union by Rank)
        if rank[xroot] < rank[yroot]:</pre>
            parent[xroot] = yroot
        elif rank[xroot] > rank[yroot]:
           parent[yroot] = xroot
        # If ranks are same, then make one as root
        # and increment its rank by one
        else :
            parent[yroot] = xroot
           rank[xroot] += 1
```

```
# The main function to construct MST using Kruskal's
    # algorithm
def KruskalMST(self):
    result =[] #This will store the resultant MST
   i = 0 # An index variable, used for sorted edges
    e = 0 # An index variable, used for result[]
        # Step 1: Sort all the edges in non-decreasing
            # order of their
            # weight. If we are not allowed to change the
            # given graph, we can create a copy of graph
   self.graph = sorted(self.graph, key=lambda item: item[2])
   parent = [] ; rank = []
    # Create V subsets with single elements
    for node in range(self.V):
       parent.append(node)
       rank.append(0)
    # Number of edges to be taken is equal to V-1
   while e < self.V -1:
        # Step 2: Pick the smallest edge and increment
                # the index for next iteration
       u,v,w = self.graph[i]
       i = i + 1
        x = self.find(parent, u)
       y = self.find(parent , v)
        # If including this edge does't cause cycle,
                    # include it in result and increment the index
                    # of result for next edge
       if x != y:
            e = e + 1
            result.append([u,v,w])
            self.union(parent, rank, x, y)
        # Else discard the edge
    # print the contents of result[] to display the built MST
   print "Following are the edges in the constructed MST"
    for u, v, weight in result:
        \#print str(u) + " -- " + str(v) + " == " + str(weight)
```