3.8

可以画出计算树为

流水为

| pipe\time | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 |
|-----------|---|---|---|---|---|---|---|---|---|----|----|----|----|----|----|----|----|----|
| 5 | | | | 1 | | 2 | | 3 | | 4 | 5 | | | 6 | | | | 7 |
| 4 | | | | | | | | | | 5 | | | 6 | | | | 7 | |
| 3 | | | | | | | | | 5 | | | 6 | | | | 7 | | |
| 2 | | 1 | 1 | 2 | 2 | 3 | 3 | 4 | 4 | | | | | | | | | |
| 1 | 1 | | 2 | | 3 | | 4 | 5 | | | 6 | | | | 7 | | | |

总时间为18△t

 $TP = 7 / (18^{\Delta}t) = 0.39 / ^{\Delta}t$ $S = (4 4^{\Delta}t + 3 4^{\Delta}t) / (18^{\Delta}t) = 14 / 9 = 1.56$ $E = (4 4^{\Delta}t + 3 4^{\Delta}t) / (5 * 18^{\Delta}t) = 14 / 45 = 0.31$

3.9

(1)

先计算F = {1, 3, 4, 8} 对应二进制为

C0:

C0 = 10001101

C2 = 10101111

C5 = 10001101 = C0

C6 = 10001111

C7 = 10001101 = C0

C2:

C2_2 = 10101111 = C2

C2_5 = 10001101 = C0

C2_6 = 10001111 = C6

C2_7 = 10001101 = C0

C6:

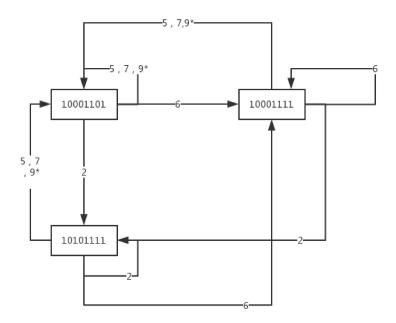
C0 = 10001101

C6_2 = 10101111 = C2

C6_5 = 10001101 = C0

C6_6 = 10001111 = C6

C6_7 = 10001101 = C0



(2)

计算平均启动距离

| Circle | Mean Start Distance |
|---------|---------------------|
| 2, 5 | 3.5 |
| 2, 6, 5 | 4.333 |
| 6, 5 | 5.5 |
| 6, 2, 5 | 4.333 |
| 5 | 5 |

采用(2,5)的间隔

最大吞吐率为 1/(3.5△t) = 0.29/△t

(3)

调度开始处分别为1,3,8,10,15,17 , 最后一个Task结束时刻为25 吞吐率为6/(25 $^{\circ}$ t) = 0.24/ $^{\circ}$ t

3.10

(1)

F = {1, 3, 6}

C0 = 100101

C2 = 101101

C4 = 100111

C5 = 100101 = C0

C2_2 = 101111

C2_4 = 100111 = C4

C2_5 = 100101 = C0

C4_2 = 101101 = C2

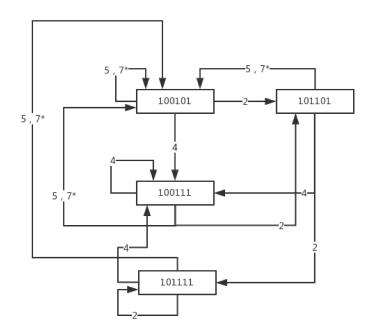
C4_4 = 100111 = C4

C4_5 = 100101 = C0

C2_2_2 = 101111 = C2_2

C2_2_4 = 100111 = C4

C2_2_5 = 100101 = C0



(2)

如果允许不等时间间隔,那么选择2,2,5,最大吞吐率为0.333/at如果不允许不等时间间隔,那么选择4,最大吞吐率为0.25/at

(3)

如果允许不等时间间隔,开始位置为1,3,5,10,12,14,19,21,23,28.结束位置34,吞吐率为0.29/ $^{\perp}$ t,加速比为70/34=2.06 如果不允许不等时间间隔,开始位置为1,5,9,13,17,21,25,29,33,37,结束位置43,吞吐率为0.23/ $^{\perp}$ t,加速比为70/43=1.63

3.11

(1)

| instruction\time | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 |
|------------------|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|
| LW | lF | ID | EX | ME | WB | | | | | | | | | | | | | | | |
| DADDUI | | IF | S | S | ID | EX | ME | WB | | | | | | | | | | | | |
| SW | | | | | | IF | S | S | ID | EX | ME | WB | | | | | | | | |
| DADDUI | | | | | | | | IF | ID | EX | ME | WB | | | | | | | | |
| DSUB | | | | | | | | | lF | S | S | ID | EX | ME | WB | | | | | |
| BNEZ | | | | | | | | | | | | IF | S | S | ID | EX | ME | WB | | |
| LW(2nd loop) | | | | | | | | | | | | | | | S | IF | ID | EX | ME | WB |

需要一共99个loop,每个loop需要15个时钟周期,总计15*99+3 = 1488

(2)

| instruction\time | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 |
|------------------|----|----|----|----|----|----|----|----|----|----|----|----|----|----|
| LW | IF | ID | EX | ME | WB | | | | | | | | | |
| DADDUI | | IF | ID | S | EX | ME | WB | | | | | | | |
| SW | | | IF | S | ID | EX | ME | WB | | | | | | |
| DADDUI | | | | | IF | ID | EX | ME | WB | | | | | |
| DSUB | | | | | | IF | ID | EX | ME | WB | | | | |

| BNEZ | | | | IF | ç | ID | EX | ME | WB | | |
|--------------|--|--|--|----|---|----|----|------|----|----|----|
| DIVEZ | | | | 11 | 3 | ID | LA | IVIL | WD | | |
| LW(2nd loop) | | | | | | S | IF | ID | EX | ME | WB |

一共需要9 * 99 + 3 = 894个时钟周期

(3)

更改指令为

```
LW R1,0(R2)

DADDIU R2,R2,#4

DSUB R4,R3,R2

DADDIU R1,R1,#1

BNEZ R4,LOOP

SW R1,-4(R2)
```

| instruction\time | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 |
|------------------|----|----|----|----|----|----|----|----|----|----|----|
| LW | lF | ID | EX | ME | WB | | | | | | |
| DADDUI | | IF | ID | EX | ME | WB | | | | | |
| DSUB | | | lF | ID | EX | ME | WB | | | | |
| DADDUI | | | | IF | ID | EX | ME | WB | | | |
| BNEZ | | | | | IF | ID | EX | ME | WB | | |
| SW | | | | | | lF | ID | EX | ME | WB | |
| LW(2nd loop) | | | | | | | IF | ID | EX | ME | WB |

需要6 * 99 + 4 = 598个时钟周期

Q1

Q1: Write down the bypass condition for the path

between M (Memory) -> D (Decode) stages into register B. (The path is shown with a dotted line in the figure.)

```
Bypass MEM->ID(B) = Case opcode M

LW => rdata
... => addr
```

BSrc need to choose bypass by the rd2 and wd, if rd2 == rd E or rd M or rd W then choose bypass ALU->ID(B) or bypass Memory->ID(B) or bypass Writeback->ID(B)

Q2

Q2 : Write down the stall condition in which stalls are only caused by data hazards. Stall = (rsD=wsE). (opcodeE=LWE).(wsE \neq 0).re2D + (rtD=wsE). (opcodeE=LWE).(wsE \neq 0).re2D

Q3

```
LW r2 r1 (r2 <- MEM[r1])
NOP
ADD r3 r1 r2 (r3 <- r1 + r2)
NOP
NOP
```

In the 3th instruction, the register B is r2 and r2 should use the data from instruct 'LW r2 r1'.