

NOTES

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DEPARTMENT OF PHYSICS

Quantum Information

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1 Elements of Quantum Mechanics and Quantum Information

1.1 Tools I from Quantum Mechanics

1.1.1 States and Operators

The harmonic oscillator

$$\hat{H} = \hbar\omega \left(\hat{a}^\dagger \hat{a} + \frac{1}{2} \right) \quad (1)$$

where the unit of $\hbar = 1$ has been applied and the zero-point energy $\omega/2$ has been ignored. With the common relations

$$\hat{a} |n\rangle = \sqrt{n} |n-1\rangle, \quad \hat{a}^\dagger |n\rangle = \sqrt{n+1} |n+1\rangle \quad (2)$$

one can evaluate everything with any knowledge of the wave functions $\psi(x)$ in real space or any other representation.

We can express the operator a as

$$\hat{a} = \sum_{m,n=0}^{\infty} |m\rangle \langle m| \hat{a} |n\rangle \langle n| = \sum_{n=0}^{\infty} \sqrt{n+1} |n\rangle \langle n+1| \quad (3)$$

Similarly, we can express any operator for any quantum mechanical system as

$$\hat{A} = \sum_{i,j} \langle i| \hat{A} |j\rangle |i\rangle \langle j| \quad (4)$$

For a two-dimensional system, *i.e.* a qubit, the corresponding space of operators is in 2×2 matrix, and a common basis is given by the identity $\mathbb{1}$ and the three Pauli matrices.

$$\hat{\sigma}_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad \hat{\sigma}_y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}, \quad \hat{\sigma}_z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}, \quad \mathbb{1} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \quad (5)$$

that implies the **commutator** and **anti-commutator** relation

$$[\hat{\sigma}_\alpha, \hat{\sigma}_\beta] = 2i\varepsilon_{\alpha\beta\gamma} \hat{\sigma}_\gamma, \quad \{\hat{\sigma}_\alpha, \hat{\sigma}_\beta\} = 0 \quad (6)$$

1.1.2 Dynamics

The Schrödinger equation is written as ($\hbar = 1$)

$$i \frac{d}{dt} |\Psi(t)\rangle = \hat{H} |\Psi(t)\rangle \quad (7)$$

with a time-independent Hamiltonian, one obtains that

$$|\Psi(t)\rangle = \hat{U}(t, t_0) |\Psi(t_0)\rangle = e^{-i\hat{H}(t-t_0)} |\Psi(t_0)\rangle \quad (8)$$

One may express this also in terms of the *time-evolution operator* or *propagator*

$$\hat{U}(t, t_0) = e^{-i\hat{H}(t-t_0)} \quad (9)$$

which has several key properties:

- **The propagator \hat{U} and \hat{H} commute, and have the same eigenstates.** With the spectral decomposition $\hat{H} = \sum_j \omega_j |\Psi_j\rangle \langle \Psi_j|$, one obtains

$$\hat{U} = \hat{U} \sum_j |\Psi_j\rangle \langle \Psi_j| = \sum_j e^{-i\omega_j(t-t_0)} |\Psi_j\rangle \langle \Psi_j| \quad (10)$$

- **The propagator is unitary, i.e.**

$$\hat{U}(t, t_0) \hat{U}^\dagger(t, t_0) = \hat{U}(t, t_0)^\dagger \hat{U}(t, t_0) = \mathbb{I} \quad (11)$$

which guarantees conservation of norm of $|\Psi(t)\rangle$, and thus normalisation.

- **The propagator (as operator) satisfies the Schrödinger equation.** The basis states evolve as $|\Psi_j(t)\rangle = \hat{U}(t, t_0) |\Psi_j(t_0)\rangle$. That is, we can write the propagator as

$$\hat{U}(t, t_0) = \hat{U}(t, t_0) \sum_j |\Psi_j(t_0)\rangle \langle \Psi_j(t_0)| = \sum_j |\Psi_j(t)\rangle \langle \Psi_j(t_0)| \quad (12)$$

We thus obtain

$$i \frac{d}{dt} \hat{U}(t, t_0) = \sum_j i \frac{d}{dt} |\Psi_j(t)\rangle \langle \Psi_j(t_0)| = \sum_j \hat{H} |\Psi_j(t)\rangle \langle \Psi_j(t_0)| = \hat{H} \hat{U}(t, t_0) \quad (13)$$

In quantum information, one often uses the term ‘quantum gate’ or simply ‘gate’ instead of propagator.

1.1.3 Measurement

Formally, a measurement is described in terms of projectors. Choose $\{|0\rangle, |1\rangle\}$ as the measurement basis, and the projectors are $\hat{P}_0 = |0\rangle \langle 0|$ and $\hat{P}_1 = |1\rangle \langle 1|$. The state reduction

$$\hat{P}_0 |\psi\rangle = \hat{P}_0 (\alpha |0\rangle + \beta |1\rangle) = \alpha |0\rangle \quad (14)$$

$$\hat{P}_1 |\psi\rangle = \hat{P}_1 (\alpha |0\rangle + \beta |1\rangle) = \beta |1\rangle \quad (15)$$

with probability

$$p_0 = \langle \psi | \hat{P}_0 | \psi \rangle = |\alpha|^2, \quad p_1 = \langle \psi | \hat{P}_1 | \psi \rangle = |\beta|^2 \quad (16)$$

The condition

$$\sum_i p_i = \langle \psi | \sum_i \hat{P}_i | \psi \rangle = 1 \quad (17)$$

must be required.

1.2 Quantum Key Distribution - BB84

Quantum Key Distribution (QKD) permits to share string of random numbers. One of the QKD protocol is *BB84*, which invented by Charles Bennett and Giles Brassard in 1984. Alice uses one qubit and prepares randomly one of the four states

$$|H\rangle = |0\rangle, \quad |V\rangle = |1\rangle, \quad |D\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle), \quad |A\rangle = \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) \quad (18)$$

1. **Sending:** Alice randomly selects a string of bits and a string of basis (+ or \times) of equal length.

$$+ \rightarrow \{|H\rangle, |V\rangle\}, \quad \times \rightarrow \{|D\rangle, |A\rangle\}$$

Then she transmits a photon for each bit with the corresponding polarization to Bob.

2. **Receiving:** Bob randomly chooses a basis for each photon to measure its polarization. If Bob selects the same basis as Alice for a particular photon, he will correctly find the bit Alice wanted to share as he measured the same polarization. If he doesn't guess correctly, he will get a random bit.
3. **Compare:** Bob tells Alice the bases he used to measure each photon. Alice informs Bob of the bases he guessed correctly to measure the encoded bits. After that, Alice and Bob remove the encoded and measured bits on different bases. Now, Alice and Bob have an identical bit-string, the **shifted key**.

Alice's random bits	1	0	0	1	1	0	1	0	0	1
Alice's encoding basis	+	\times	\times	+	\times	+	+	+	+	+
Photons Alice sends	V	D	D	V	A	H	V	H	H	V
Random measurement basis	+	+	\times	+	+	\times	+	+	+	\times
Bits as received by Bob	V	V	D	V	H	A	V	H	H	D
Reveal the sequence of their basis	✓		✓	✓			✓	✓	✓	
Shifted Key	1		0	1			1	0	0	

Table 1: BB84 protocol

1.3 No-Cloning Theorem

Let's assume we have a qubit in a given state $|\Psi\rangle$ (but we don't know the state) and a second qubit in the $|0\rangle$ state. We would like to find a gate such that

$$\hat{U} |\Psi\rangle |0\rangle = |\Psi\rangle |\Psi\rangle \quad (19)$$

for any state $|\Psi\rangle$, so we have

$$\hat{U} |0\rangle |0\rangle = |0\rangle |0\rangle \quad (20)$$

$$\hat{U} |1\rangle |0\rangle = |1\rangle |1\rangle \quad (21)$$

For a general state $|\Psi\rangle = \alpha|0\rangle + \beta|1\rangle$, this implies that

$$\hat{U}(\alpha|0\rangle + \beta|1\rangle)|0\rangle = \alpha|0\rangle|0\rangle + \beta|1\rangle|1\rangle \quad (22)$$

whereas we would have wanted to obtain

$$(\alpha|0\rangle + \beta|1\rangle)(\alpha|0\rangle + \beta|1\rangle) = \alpha^2|0\rangle|0\rangle + \alpha\beta(|0\rangle|1\rangle + |1\rangle|0\rangle) + \beta^2|1\rangle|1\rangle \quad (23)$$

In the above example, the cloning process works if either α or β vanishes; that is, it works for the two orthogonal basis states for which it is defined.

1.4 Tools II from Quantum Mechanics

1.4.1 Bell's Inequality

Assume that Alice has two cards A_0, A_1 having their value either $+1$ or -1 and Bob has his cards B_0 and B_1 also having their values $+1$ or -1 . If all the cards bear the value $+1$, i.e. $A_0 = A_1 = B_0 = B_1 = 1$, then

$$A_0B_0 + A_0B_1 + A_1B_0 - A_1B_1 = 2 \quad (24)$$

The average

$$\langle A_0B_0 \rangle + \langle A_0B_1 \rangle + \langle A_1B_0 \rangle - \langle A_1B_1 \rangle \leq 2 \quad (25)$$

This is called *Bell's inequality* which is obtained for **classical** physics based on local realism.

However, we can prove that a quantum-mechanically correlated state can violate Bell's inequality. Let us assume the singlet state,

$$|\psi\rangle = \frac{1}{\sqrt{2}}(|01\rangle - |10\rangle) \quad (26)$$

and the observables $\hat{A}_0 = \hat{\sigma}_z$, $\hat{A}_1 = \hat{\sigma}_x$, $\hat{B}_0 = -\frac{1}{\sqrt{2}}(\hat{\sigma}_x + \hat{\sigma}_z)$, $\hat{B}_1 = \frac{1}{\sqrt{2}}(\hat{\sigma}_x - \hat{\sigma}_z)$. Then we find that

$$\begin{aligned} \langle \hat{A}_0 \otimes \hat{B}_0 \rangle &= \langle \psi | \hat{A}_0 \otimes \hat{B}_0 | \psi \rangle \\ &= -\frac{1}{2\sqrt{2}}(\langle 01 | - \langle 10 |) \hat{\sigma}_z \otimes (\hat{\sigma}_x + \hat{\sigma}_z) (|01\rangle - |10\rangle) \\ &= -\frac{1}{2\sqrt{2}}(\langle 01 | - \langle 10 |) (|00\rangle - |01\rangle + |11\rangle + |10\rangle) = \frac{1}{\sqrt{2}} \end{aligned} \quad (27)$$

Similarly, we can find

$$\langle \hat{A}_0 \otimes \hat{B}_0 \rangle = \langle \hat{A}_0 \otimes \hat{B}_1 \rangle = \langle \hat{A}_1 \otimes \hat{B}_0 \rangle = -\langle \hat{A}_1 \otimes \hat{B}_1 \rangle = \frac{1}{\sqrt{2}} \quad (28)$$

and the Bell's expectation value is

$$\langle \hat{A}_0 \otimes \hat{B}_0 \rangle + \langle \hat{A}_0 \otimes \hat{B}_1 \rangle + \langle \hat{A}_1 \otimes \hat{B}_0 \rangle - \langle \hat{A}_1 \otimes \hat{B}_1 \rangle = 2\sqrt{2} \quad (29)$$

which clearly **violates** Bell's inequality.

1.4.2 Partial Trace

The trace of the operator \hat{A} reads

$$\text{tr } \hat{A} = \sum_i \langle i | \hat{A} | i \rangle \quad (30)$$

for any orthonormal basis $\{|i\rangle\}$. For an operator $\hat{A} \otimes \hat{B}$ (on $\mathcal{H}_a \otimes \mathcal{H}_b$) we can define the partial traces

$$\hat{A} \text{ tr}_b \hat{B} = \text{tr}_b (\hat{A} \otimes \hat{B}) \quad (31)$$

$$\hat{B} \text{ tr}_a \hat{A} = \text{tr}_a (\hat{A} \otimes \hat{B}) \quad (32)$$

The partial trace naturally appears in expressions of the form $\text{tr} \left((\hat{A} \otimes \mathbb{I}) \hat{C} \right)$, where \hat{A} acts on \mathcal{H}_a , \mathbb{I} acts on \mathcal{H}_b and \hat{C} acts on $\mathcal{H}_a \otimes \mathcal{H}_b$.

$$\begin{aligned} \text{tr} \left((\hat{A} \otimes \mathbb{I}) \hat{C} \right) &= \sum_{ij} \langle i |_a \otimes \langle j |_b \left((\hat{A} \otimes \mathbb{I}) \hat{C} \right) |i\rangle_a \otimes |j\rangle_b \\ &= \sum_{ij} \left(\langle i |_a \hat{A} \right) \otimes \langle j |_b \hat{C} (|i\rangle_a \otimes |j\rangle_b) \\ &= \sum_i \langle i |_a \hat{A} \left(\sum_j \langle j |_b \hat{C} |j\rangle_b \right) |i\rangle_a \\ &= \text{tr}_a \left(\hat{A} \text{tr}_b \hat{C} \right) \end{aligned} \quad (33)$$

2 Quantum Algorithms

2.1 Quantum Interferometers

Let's take the Mach-Zehnder interferometer as an example.

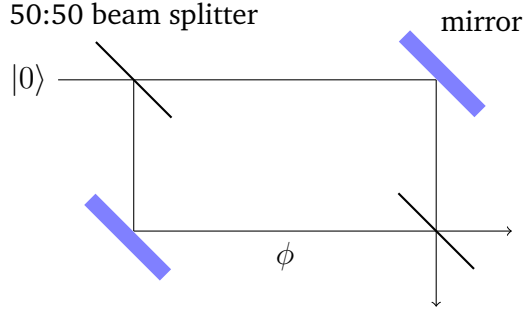


Figure 1: The Mach-Zehnder interferometer.

The Mach-Zehnder interferometer is described by

$$\begin{aligned}
 \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & -1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & e^{i\phi} \end{pmatrix} \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ -1 & 1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} &= \frac{1}{2} \begin{pmatrix} 1 & -1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & e^{i\phi} \end{pmatrix} \begin{pmatrix} 1 \\ -1 \end{pmatrix} \\
 &= \frac{1}{2} \begin{pmatrix} 1 & -1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} 1 \\ -e^{i\phi} \end{pmatrix} \\
 &= \frac{1}{2} e^{i\phi/2} \begin{pmatrix} e^{-i\phi/2} + e^{i\phi/2} \\ e^{-i\phi/2} - e^{i\phi/2} \end{pmatrix} \\
 &= e^{i\phi/2} \begin{pmatrix} \cos \phi/2 \\ -i \sin \phi/2 \end{pmatrix}
 \end{aligned} \tag{34}$$

where phase shift is described by the rotation

$$\hat{R}_\phi = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\phi} \end{pmatrix} = e^{i\phi/2} \begin{pmatrix} e^{-i\phi/2} & 0 \\ 0 & e^{i\phi/2} \end{pmatrix} \tag{35}$$

and the beam splitter is described by

$$\hat{B} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ -1 & 1 \end{pmatrix} = \frac{1}{\sqrt{2}} (\mathbb{1} + \hat{\sigma}_y) = e^{i\pi\hat{\sigma}_y/4} \tag{36}$$

2.2 Basic Gate Operators

- **Pauli X, Y, Z gates** - single qubit gates

$$\hat{X} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad \hat{Y} = \begin{pmatrix} 0 & i \\ -i & 0 \end{pmatrix}, \quad \hat{Z} = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \tag{37}$$

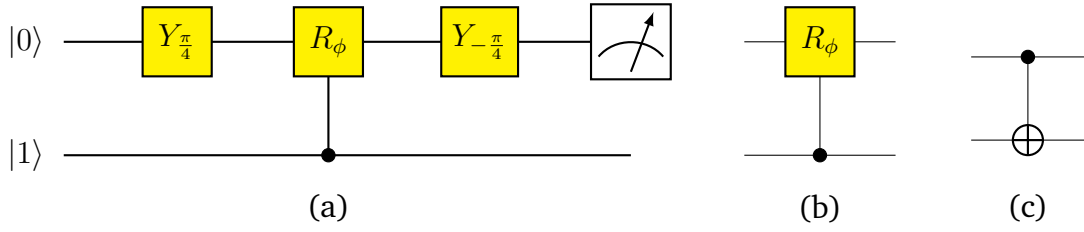


Figure 2: (a) For the input state $|0\rangle$, the quantum circuit for the Macj-Zehnder interferometer. (b) The controlled-phase gate. (c) The Controlled-NOT gate.

• **Controlled-phase gate**

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & e^{i\phi} \end{pmatrix} \quad (38)$$

• **Controlled-NOT gate**

$$\begin{aligned} \hat{U}_c &= |00\rangle \langle 00| + |01\rangle \langle 01| + |10\rangle \langle 11| + |11\rangle \langle 10| \\ &= |0\rangle \langle 0| \otimes \mathbb{1} + |1\rangle \langle 1| \otimes \hat{\sigma}_x \end{aligned} \quad (39)$$

which generates

input		output	
control	target	control	target
0	0	0	0
0	1	0	1
1	0	1	1
1	1	1	0

Table 2: Controlled-NOT gate.

Compare with the XOR gate (for conventional computer)

input		output
0	0	0
0	1	1
1	0	1
1	1	0

Table 3: The XOR gate in conventional computer.

NOTE: Any unitary on a system with several qubits as a sequence of *single qubit gates* and *CNOT* gates. A general single qubit unitary gate and two-qubit controlled-gate are called a *universal* quantum gate.

2.3 Deutsch Jozsa Algorithm

The task of this algorithm is to probe whether a function $f : \{0, 1\} \rightarrow \{0, 1\}$ is constant. *i.e.* if $f(0) = f(1)$, or if it is balanced, *i.e.* if $f(0) \neq f(1)$. The function f can be implemented in terms of a two qubit gate

$$|i\rangle \otimes |j\rangle \rightarrow \hat{U}_o |ij\rangle = |i\rangle \otimes |j \oplus f(i)\rangle \quad (40)$$

where ' \oplus ' denotes the addition modulo 2, *i.e.* $0 \oplus A = A$ and $1 \oplus A = \bar{A}$, where \bar{A} denotes 'not A '.

$$\hat{U}_o |00\rangle = |0f(0)\rangle \quad (41)$$

$$\hat{U}_o |01\rangle = |0\overline{f(0)}\rangle \quad (42)$$

$$\hat{U}_o |10\rangle = |1f(1)\rangle \quad (43)$$

$$\hat{U}_o |11\rangle = |1\overline{f(1)}\rangle \quad (44)$$

With the initial state

$$\frac{|0\rangle + |1\rangle}{\sqrt{2}} \otimes \frac{|0\rangle - |1\rangle}{\sqrt{2}} = \frac{1}{2} (|00\rangle - |01\rangle + |10\rangle - |11\rangle) \quad (45)$$

one obtains

$$\hat{U}_o \frac{|0\rangle + |1\rangle}{\sqrt{2}} \otimes \frac{|0\rangle - |1\rangle}{\sqrt{2}} = \frac{1}{2} (|0f(0)\rangle - |0\overline{f(0)}\rangle + |1f(1)\rangle - |1\overline{f(1)}\rangle) = |\Psi_0\rangle \quad (46)$$

If $f(0) = f(1)$, this reduces to

$$\begin{aligned} |\Psi_0\rangle &= \frac{1}{2} [(|0\rangle + |1\rangle) \otimes |f(0)\rangle - (|0\rangle + |1\rangle) \otimes |\overline{f(0)}\rangle] \\ &= \frac{|0\rangle + |1\rangle}{\sqrt{2}} \otimes \frac{|0\rangle - |\overline{f(0)}\rangle}{\sqrt{2}} = |+\rangle \otimes \frac{|0\rangle - |\overline{f(0)}\rangle}{\sqrt{2}} \end{aligned} \quad (47)$$

If $f(0) = \overline{f(1)}$, this reduces to

$$\begin{aligned} |\Psi_0\rangle &= \frac{1}{2} [(|0\rangle - |1\rangle) \otimes |f(0)\rangle - (|0\rangle - |1\rangle) \otimes |\overline{f(0)}\rangle] \\ &= \frac{|0\rangle - |1\rangle}{\sqrt{2}} \otimes \frac{|0\rangle - |f(1)\rangle}{\sqrt{2}} = |-\rangle \otimes \frac{|0\rangle - |f(1)\rangle}{\sqrt{2}} \end{aligned} \quad (48)$$

A measurement on the first qubit in the $\hat{\sigma}_x$ -basis permits to distinguish between these two cases.

In practice, the algorithm would be broken down in the elementary steps:

1. Prepare of the initial state $|0\rangle \otimes |0\rangle$.

2. Application of the gate $e^{-i\frac{\pi}{4}\hat{\sigma}_y} \otimes e^{i\frac{\pi}{4}\hat{\sigma}_y}$ (if we want to use the Hadamard gate $\hat{H} \otimes \hat{H} \hat{\sigma}_x$)

$$\left[\exp\left(-i\frac{\pi}{4}\hat{\sigma}_y\right) \otimes \exp\left(i\frac{\pi}{4}\hat{\sigma}_y\right) \right] (|0\rangle \otimes |0\rangle) = |+\rangle \otimes |-\rangle \quad (49)$$

3. Query to the oracle

- If $f(0) = f(1)$, then the first qubit is in $|+\rangle$.
- If $f(0) \neq f(1)$, then the first qubit is in $|-\rangle$.

4. Application of the gate $\exp(-i\frac{\pi}{4}\hat{\sigma}_y) \otimes \mathbb{1}$, which will bring $|+\rangle$ to $|1\rangle$ and $|-\rangle$ to $|0\rangle$.

5. Measurement on the first qubit in the $\hat{\sigma}_z$ -basis.

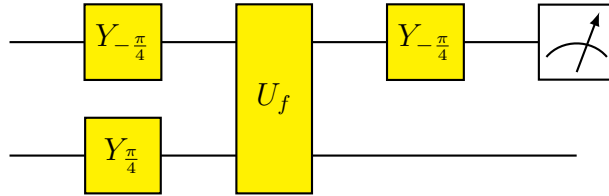


Figure 3: The quantum circuit of the Deutsch-Jozsa algorithm.

2.4 Basic Gate Operation II

2.4.1 Controlled-Z Gate and Controlled-Unitary Gate

$$\hat{U}_{cz} = |0\rangle\langle 0| \otimes \mathbb{1} + |1\rangle\langle 1| \otimes \hat{\sigma}_z \quad (50)$$

similar to a CNOT gate which reads

$$\hat{U}_c = \hat{U}_{cx} = |0\rangle\langle 0| \otimes \mathbb{1} + |1\rangle\langle 1| \otimes \hat{\sigma}_x \quad (51)$$

With $\hat{\sigma}_z = \exp(i\frac{\pi}{4}\hat{\sigma}_y)\hat{\sigma}_x\exp(-i\frac{\pi}{4}\hat{\sigma}_y)$ one can see that \hat{U}_{cz} can be realised as the gate sequence

$$\begin{aligned} \hat{U}_{cz} &= |0\rangle\langle 0| \otimes \mathbb{1} + |1\rangle\langle 1| \otimes \exp\left(i\frac{\pi}{4}\hat{\sigma}_y\right)\hat{\sigma}_x\exp\left(-i\frac{\pi}{4}\hat{\sigma}_y\right) \\ &= \left[\mathbb{1} \otimes \exp\left(i\frac{\pi}{4}\hat{\sigma}_y\right)\right] (|0\rangle\langle 0| \otimes \mathbb{1}) \left[\mathbb{1} \otimes \exp\left(-i\frac{\pi}{4}\hat{\sigma}_y\right)\right] \\ &\quad + \left[\mathbb{1} \otimes \exp\left(i\frac{\pi}{4}\hat{\sigma}_y\right)\right] (|1\rangle\langle 1| \otimes \hat{\sigma}_x) \left[\mathbb{1} \otimes \exp\left(-i\frac{\pi}{4}\hat{\sigma}_y\right)\right] \\ &= \left[\mathbb{1} \otimes \exp\left(i\frac{\pi}{4}\hat{\sigma}_y\right)\right] \hat{U}_c \left[\mathbb{1} \otimes \exp\left(-i\frac{\pi}{4}\hat{\sigma}_y\right)\right] \end{aligned} \quad (52)$$

so we have

$$\hat{U}_{cz} |00\rangle = |00\rangle, \quad \hat{U}_{cz} |01\rangle = |01\rangle, \quad \hat{U}_{cz} |10\rangle = |10\rangle, \quad \hat{U}_{cz} |11\rangle = -|11\rangle \quad (53)$$

We can generalize two qubit operations to the controlled-unitary operation:

$$\hat{U}_{cu} = |0\rangle\langle 0| \otimes \mathbb{1} + |1\rangle\langle 1| \otimes \hat{U} \quad (54)$$

NOTE: In a controlled gate operation, the action on one qubit is dependent on the state of another qubit. So they cannot be written in

$$\hat{U}_1 \otimes \hat{U}_2 \quad (55)$$

2.4.2 Three Qubits - Controlled-Controlled-Unitary Gate

A generalisation of the controlled-unitary gate to a system of three qubits is a controlled-controlled-unitary gate

$$(\mathbb{1} \otimes \mathbb{1} - |1\rangle\langle 1| \otimes |1\rangle\langle 1|) \otimes \mathbb{1} + |1\rangle\langle 1| \otimes |1\rangle\langle 1| \otimes \hat{U} \quad (56)$$

The controlled-not gate operation can be written in

$$|i\rangle|j\rangle \rightarrow |i\rangle|i \oplus j\rangle \quad (57)$$

The controlled-unitary gate operation can be written in

$$|i\rangle|j\rangle \rightarrow |i\rangle\hat{U}^i|j\rangle \quad (58)$$

The three qubits gate can be decomposed into single and two-qubit gates as depicted in the quantum circuit Fig.(4)

1. $|i\rangle|j\rangle\hat{V}^j|k\rangle$
2. $|i\rangle|i \oplus j\rangle\hat{V}^j|k\rangle$
3. $|i\rangle|i \oplus j\rangle(\hat{V}^\dagger)^{i \oplus j}\hat{V}^j|k\rangle$
4. $|i\rangle|i \oplus (i \oplus j)\rangle(\hat{V}^\dagger)^{i \oplus j}\hat{V}^j|k\rangle = |i\rangle|j\rangle(\hat{V}^\dagger)^{i \oplus j}\hat{V}^j|k\rangle$
5. $|i\rangle|j\rangle\hat{V}^i(\hat{V}^\dagger)^{i \oplus j}\hat{V}^j|k\rangle$

So we have

$$\begin{aligned} \hat{U}|00\rangle \otimes |\phi\rangle &= |00\rangle \otimes |\phi\rangle \\ \hat{U}|01\rangle \otimes |\phi\rangle &= |01\rangle \otimes |\phi\rangle \\ \hat{U}|10\rangle \otimes |\phi\rangle &= |10\rangle \otimes |\phi\rangle \\ \hat{U}|11\rangle \otimes |\phi\rangle &= |11\rangle \otimes \hat{U}|\phi\rangle \end{aligned} \quad (59)$$

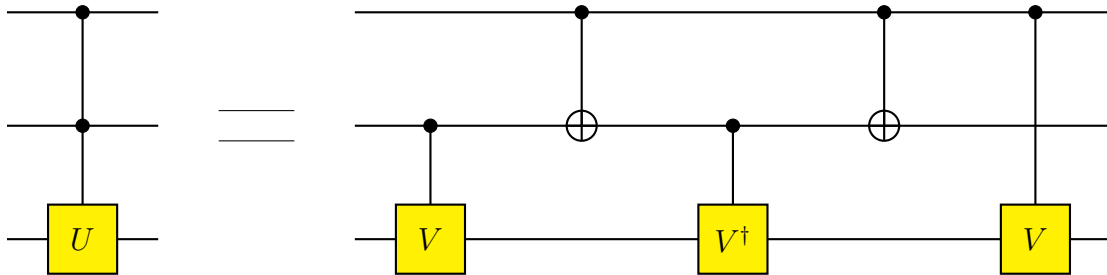


Figure 4: A controlled-controlled- U operation can be realised in terms of CNOT and controlled- V operations for $\hat{U} = \hat{V}^2$

2.4.3 NOT Gate

The Pauli $\hat{\sigma}_x$ operator is also called as the NOT gate as it switches

$$\hat{\sigma}_x |0\rangle = |1\rangle, \quad \hat{\sigma}_x |1\rangle = |0\rangle \quad (60)$$

In quantum mechanics, do we have a universal-NOT gate? Let's consider an arbitrary state

$$|\Psi\rangle = \alpha |0\rangle + \beta |1\rangle \quad (61)$$

then

$$\text{NOT } |\Psi\rangle = \alpha |1\rangle + \beta |0\rangle = |\Psi'\rangle \quad (62)$$

as

$$\langle \Psi' | \Psi \rangle = (\alpha^* \langle 1| + \beta^* \langle 0|) (\alpha |0\rangle + \beta |1\rangle) = \alpha^* \beta + \beta^* \alpha \neq 0 \quad (63)$$

NOT operation cannot to bring the initial state $|\psi\rangle$ to its orthogonal state $|\Psi^\perp\rangle$ ($\langle \Psi^\perp | \Psi \rangle = 0$), *i.e.* we don't have a universal-NOT gate in quantum mechanics.

2.4.4 SWAP Gate

The SWAP gate is to swap two qubits

$$\text{SWAP } |\Psi\rangle |\Phi\rangle = |\Phi\rangle |\Psi\rangle \quad (64)$$

with

$$\text{SWAP} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \quad (65)$$

2.5 Grover's Algorithm (Quantum Search Algorithm)

The goal of the algorithm is to find a solution to the search problem. Let's consider a system with n qubits; that is, we are working in an $N = 2^n$ dimensional Hilbert space. The Grover's algorithm follows the following steps:

- 1) The system begins with the initial state $|0\rangle |0\rangle \dots |0\rangle = |0\rangle^{\otimes n}$
- 2) Apply Hadamard transform $\hat{H}^{\otimes n}$

$$|s\rangle = \hat{H} |0\rangle \dots \hat{H} |0\rangle \hat{H} |0\rangle = |+\rangle^{\otimes n} = \frac{1}{\sqrt{N}} \sum_{i=0}^{N-1} |i\rangle \quad (66)$$

- 3) Apply $\hat{U} = \mathbb{1} - 2 |q\rangle \langle q|$

$$\begin{aligned} |\Psi\rangle &= \hat{U} |s\rangle = (\mathbb{1} - 2 |q\rangle \langle q|) |s\rangle \\ &= |s\rangle - 2 |q\rangle \langle q| \frac{1}{\sqrt{N}} \sum_{i=0}^{N-1} |i\rangle \\ &= |s\rangle - \frac{2}{\sqrt{N}} |q\rangle \end{aligned} \quad (67)$$

4) Apply $\hat{V} = 2|s\rangle\langle s| - \mathbb{1}$

$$\begin{aligned}
\hat{V}|\Psi\rangle &= (2|s\rangle\langle s| - \mathbb{1})(|s\rangle - \frac{2}{\sqrt{N}}|q\rangle) \\
&= 2|s\rangle - |s\rangle - \frac{4}{\sqrt{N}}|s\rangle\langle s|q\rangle + \frac{2}{\sqrt{N}}|q\rangle \\
&= |s\rangle - \frac{4}{N}|s\rangle \sum_{i=0}^{N-1} \langle i|q\rangle + \frac{2}{\sqrt{N}}|q\rangle \\
&= \frac{N-4}{N}|s\rangle + \frac{2}{\sqrt{N}}|q\rangle \\
&= \frac{N-4}{N\sqrt{N}} \sum_{i \neq q}^{N-1} |i\rangle + \left(\frac{N-4}{N\sqrt{N}} + \frac{2}{\sqrt{N}} \right) |q\rangle
\end{aligned} \tag{68}$$

We find that the probability of finding $|q\rangle$

$$\left(\frac{N-4}{N\sqrt{N}} + \frac{2}{\sqrt{N}} \right)^2 > \frac{1}{N} \tag{69}$$

when $N > 2$.

5) Repeat the process 2) and 3) to increase the probability of finding the system in $|q\rangle$.

2.6 Quantum Fourier Transform

The quantum Fourier transform \mathcal{QF} for an N -dimensional system is defined as

$$\boxed{\mathcal{QF}|\Phi_p\rangle = \frac{1}{\sqrt{N}} \sum_{q=0}^{N-1} e^{\frac{2\pi i}{N}pq} |\Phi_q\rangle} \tag{70}$$

For $N = 2^n$ one can realise it with n qubits.

- **A single qubit**

$$\mathcal{QF}|p\rangle = \frac{1}{\sqrt{2}} \sum_{q=0}^1 e^{\frac{2\pi i}{2}pq} |q\rangle = \frac{1}{\sqrt{2}} (|0\rangle + e^{i\pi p} |1\rangle) \tag{71}$$

This is also know as Hadamard gate, with

$$\mathcal{QF}|0\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) = |+\rangle \tag{72}$$

$$\mathcal{QF}|1\rangle = \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) = |-\rangle \tag{73}$$

- **Many qubits**

$$\mathcal{QF}_N = \frac{1}{\sqrt{N}} \begin{pmatrix} 1 & 1 & 1 & \dots & 1 \\ 1 & \omega & \omega^2 & \dots & \omega^{N-1} \\ 1 & \omega^2 & \omega^4 & \dots & \omega^{2(N-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \omega^{N-1} & \omega^{2(N-1)} & \dots & \omega^{(N-1)^2} \end{pmatrix} \quad (74)$$

where $\omega = \exp(2\pi i/N)$ is the N^{th} root of unity. The state for n qubits can be written in

$$\begin{aligned} |\Psi\rangle &= a_0 |0\rangle + a_1 |1\rangle + \dots + a_{N-1} |N-1\rangle \\ &= a_0 \begin{pmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix} + a_1 \begin{pmatrix} 0 \\ 1 \\ \vdots \\ 0 \end{pmatrix} + \dots + a_{N-1} \begin{pmatrix} 0 \\ 0 \\ \vdots \\ 1 \end{pmatrix} = \begin{pmatrix} a_0 \\ a_1 \\ \vdots \\ a_{N-1} \end{pmatrix} \end{aligned} \quad (75)$$

- **Two qubits**

$$\mathcal{QF}_4 = \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ 1 & -1 & 1 & -1 \\ 1 & -i & -1 & i \end{pmatrix} \quad (76)$$

We have some examples

- i) $|f\rangle = \frac{1}{2}(|0\rangle + |1\rangle + |2\rangle + |3\rangle) \Rightarrow |\tilde{f}\rangle = \mathcal{QF}_4 |f\rangle = |0\rangle$
- ii) $|g\rangle = |0\rangle \Rightarrow |\tilde{g}\rangle = \mathcal{QF}_4 |g\rangle = \frac{1}{2}(|0\rangle + |1\rangle + |2\rangle + |3\rangle)$
- iii) $|h\rangle = |1\rangle \Rightarrow |\tilde{h}\rangle = \mathcal{QF}_4 |h\rangle = \frac{1}{2}(|0\rangle + i|1\rangle - |2\rangle - i|3\rangle)$

2.6.1 Properties of Quantum Fourier Transform

1. QFT is unitary

Proof: an operator is unitary if its columns are orthonormal.

$$\frac{1}{N} \sum_{n=0}^{N-1} \omega^{ni} \omega^{nj*} = \frac{1}{N} \sum_{n=0}^{N-1} (\omega^{i-j})^n = \begin{cases} 1 & i = j \\ 0 & i \neq j \end{cases} \quad (77)$$

2. **Linear shift** as shown in the example above $|\tilde{g}\rangle$ and $|\tilde{h}\rangle$.

3. **Period/wave length relationship**

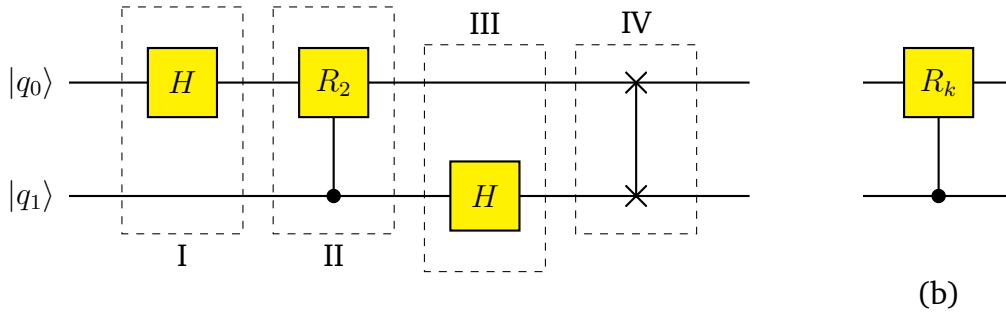


Figure 5: (a) The quantum circuit for quantum Fourier transform. (b) The part II in the circuit (a).

2.6.2 Quantum Circuit of Quantum Fourier Transform

The quantum circuit for the two-qubit QFT see in Fig.5(a).

I. The part I of the circuit reads

$$\hat{H} \otimes \mathbb{1} = \frac{1}{\sqrt{2}} \begin{pmatrix} \mathbb{1} & \mathbb{1} \\ \mathbb{1} & -\mathbb{1} \end{pmatrix} = \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & -1 & 0 \\ 0 & 1 & 0 & -1 \end{pmatrix} \quad (78)$$

II. The part II (see Fig.5(b)) reads

$$\mathbb{1} \otimes |0\rangle\langle 0| + (|0\rangle\langle 0| + e^{2\pi i/2^k} |1\rangle\langle 1|) \otimes |1\rangle\langle 1| = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & e^{2\pi i/2^k} \end{pmatrix} \quad (79)$$

III. The part III reads

$$\mathbb{1} \otimes \hat{H} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & -1 \end{pmatrix} \quad (80)$$

IV. The part IV reads

$$\text{SWAP} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \quad (81)$$

In general, the QFT acts on a quantum state $|q^{(2)}\rangle$

$$\mathcal{QF} |p\rangle = \frac{1}{\sqrt{N}} \sum_{q=0}^{N-1} e^{\frac{2\pi i}{N} pq} |q^{(2)}\rangle \quad (82)$$

$q^{(2)}$ is the binary basis of q , which satisfies the following relationship

$$q = \sum_{i=0}^{n-1} q_i^{(2)} 2^i \quad (83)$$

For example, $(1, 0, 1) = 1 \times 2^0 + 0 \times 2^1 + 1 \times 2^2 = 5$. So we can expand

$$\begin{aligned} \mathcal{QF} |p\rangle &= \frac{1}{\sqrt{N}} \sum_{q=0}^{N-1} \exp \left[\frac{2\pi i}{N} p \left(\sum_{i=0}^{n-1} q_i^{(2)} 2^i \right) \right] |q^{(2)}\rangle \\ &= \frac{1}{\sqrt{N}} \sum_{q=0}^{N-1} \exp \left[\frac{2\pi i}{N} p \left(q_0^{(2)} + \sum_{i=1}^{n-1} q_i^{(2)} 2^i \right) \right] |q^{(2)}\rangle \\ &= \frac{1}{\sqrt{N}} \sum_{q=0}^{N-1} \exp \left[\frac{2\pi i}{N} p \left(0 \times 1 + 2 \sum_{i=0}^{n-2} q_i^{(2)} 2^i \right) \right] (|0\rangle \otimes |q^{(2)}\rangle) \\ &\quad + \frac{1}{\sqrt{N}} \sum_{q=0}^{N-1} \exp \left[\frac{2\pi i}{N} p \left(1 \times 1 + 2 \sum_{i=0}^{n-2} q_i^{(2)} 2^{i-1} \right) \right] (|1\rangle \otimes |q^{(2)}\rangle) \\ &= \frac{1}{\sqrt{N}} \left(|0\rangle + e^{\frac{2\pi i}{N} p \times 1} |1\rangle \right) \otimes \sum_{q=0}^{N'-1} e^{\frac{2\pi i}{N'} p q'} |q^{(2)}\rangle \end{aligned} \quad (84)$$

where $N' = 2^{n-1}$. Thus, by performing the same procedure, we obtain the following separable state

$$\frac{1}{\sqrt{N}} \left(|0\rangle + e^{\frac{2\pi i}{2^n} p} |1\rangle \right) \otimes \left(|0\rangle + e^{\frac{2\pi i}{2^{n-1}} p} |1\rangle \right) \otimes \cdots \otimes \left(|0\rangle + e^{\pi i p} |1\rangle \right) \quad (85)$$

and $(|0\rangle + e^{\pi i p} |1\rangle)$ can be achieved by Hadamard gate

$$H |p\rangle = \frac{1}{\sqrt{2}} (|0\rangle + e^{\pi i p} |1\rangle) \quad (86)$$

The idea of QFT is thus to perform the following actions

1. Initial states: $|p_{n-1}\rangle \otimes |p_{n-2}\rangle \otimes \cdots \otimes |p_0\rangle$
2. Apply Hadamard gate H on p_0 : $|p_{n-1}\rangle \otimes |p_{n-2}\rangle \otimes \cdots \otimes (|0\rangle + e^{\pi i p} |1\rangle)$

3 Physical Realisation - Trapped Ions

3.1 Di-Vincenzo Criteria

- i) Well-defined qubits
- ii) Initialisation $|0\rangle |0\rangle \cdots |0\rangle$
- iii) Universal set of quantum gates
- iv) Measurements of each qubit
- v) Sufficiently long coherence times

3.2 Trapped Ion Hamiltonian

The free Hamiltonian is composed of

$$\hat{H}_0 = \underbrace{\frac{\omega_0}{2} \hat{\sigma}_z}_{\text{atomic internal energy}} + \underbrace{\omega_t \hat{a}^\dagger \hat{a}}_{\text{trapped motion energy}} \quad (87)$$

where ω_t is the trap frequency. The interaction between the ion and a laser can be expressed as

$$\hat{H}_I = \Omega_R \hat{\sigma}_x \cos(\omega t - k\hat{x} + \phi) \quad (88)$$

where Ω_R is the coupling strength between the ion and laser, ω is the laser frequency, and $\hat{x} = \sqrt{\hbar/2m\omega_t}(\hat{a} + \hat{a}^\dagger)$ is the position of ion. In the interaction picture¹

$$\begin{aligned} \hat{H} &= e^{i\hat{H}_0 t} \hat{H}_I e^{-i\hat{H}_0 t} \\ &= \frac{1}{2} (\hat{\sigma}_- e^{-i\omega_0 t} + \hat{\sigma}_+ e^{i\omega_0 t}) \\ &\quad \times \left[\Omega_R e^{i\phi} \exp \left(i \left(\omega t - k \sqrt{\frac{\hbar}{2m\omega_t}} (\hat{a} e^{-i\omega_t t} + \hat{a}^\dagger e^{i\omega_t t}) \right) \right) + \text{h.c.} \right] \end{aligned} \quad (89)$$

Here we introduce the *Lamb-Dicke parameter*

$$\eta = k \sqrt{\frac{\hbar}{2m\omega_t}} = \frac{\hbar k}{\sqrt{2m\hbar\omega_t}} = \frac{p_{\text{photon}}}{p_{\text{phonon}}} \quad (90)$$

Assume $\eta \ll 1$ (which is typically of the order of 1/10), then

$$\hat{H} \approx \frac{\Omega_R}{2} (\hat{\sigma}_- e^{-i\omega_0 t} + \hat{\sigma}_+ e^{i\omega_0 t}) [e^{-i\omega t} (\mathbb{1} - i\eta (\hat{a} e^{-i\omega_t t} + \hat{a}^\dagger e^{i\omega_t t})) + \text{h.c.}] \quad (91)$$

for $\phi = 0$. For sufficiently weak driving, after rotating wave approximation (RWA), one obtains the following approximations

$$1 e^{i\hat{A}\theta} \hat{B} e^{-i\hat{A}\theta} = \hat{B} + \theta [\hat{A}, \hat{B}] + \frac{\theta^2}{2} [\hat{A}, [\hat{A}, \hat{B}]] + \cdots$$

- Carrier transition $\omega = \omega_0$

$$\hat{H}_c = \frac{\Omega_R}{2} (\hat{\sigma}_+ + \hat{\sigma}_-) = \frac{\Omega_R}{2} \hat{\sigma}_x \quad (92)$$

- Red sideband $\omega = \omega_0 - \omega_t$

$$\hat{H}_r = -i\eta \frac{\Omega_R}{2} (\hat{\sigma}_- \hat{a}^\dagger - \hat{\sigma}_+ \hat{a}) \quad (93)$$

- Blue sideband $\omega = \omega_0 + \omega_t$

$$\hat{H}_b = -i\eta \frac{\Omega_R}{2} (\hat{\sigma}_- \hat{a} - \hat{\sigma}_+ \hat{a}^\dagger) \quad (94)$$

\hat{H} induces single qubit gates; \hat{H}_r and \hat{H}_b affect the ion and oscillation jointly.

3.3 Cirac-Zoller Gate

Taking $\lambda = -i\eta\Omega e^{i\phi/2}/2$ and $\phi = \pi/2$, then

$$\hat{H}_r = \lambda (\hat{\sigma}_- \hat{a}^\dagger + \hat{\sigma}_+ \hat{a}) \quad (95)$$

Solving the dynamic equation

$$|g, n\rangle \rightarrow \cos \lambda \sqrt{n} t |g, n\rangle - i \sin \lambda \sqrt{n} t |e, n-1\rangle \quad (96)$$

$$|e, n-1\rangle \rightarrow \cos \lambda \sqrt{n} t |e, n-1\rangle - i \sin \lambda \sqrt{n} t |g, n\rangle \quad (97)$$

where $|g\rangle, |e\rangle$ are ground and excited state and $|n\rangle$ is the n phonon number (trap state).

Lets consider two ionic states and phononic state

- A laser shining on ion 1 for $\lambda t = \pi/2$.
- Another laser shining on ion 2 for $\lambda t = \pi$.
- A laser shining on ion 1 for $\lambda t = \pi/2$.

	\hat{U}_1		\hat{U}_2		\hat{U}_1	
$ g, g, 0\rangle$	\rightarrow	$ g, g, 0\rangle$	\rightarrow	$ g, g, 0\rangle$	\rightarrow	$ g, g, 0\rangle$
$ g, e, 0\rangle$	\rightarrow	$ g, e, 0\rangle$	\rightarrow	$ g, e, 0\rangle$	\rightarrow	$ g, e, 0\rangle$
$ e, g, 0\rangle$	\rightarrow	$-i g, g, 1\rangle$	\rightarrow	$i g, g, 1\rangle$	\rightarrow	$ e, g, 0\rangle$
$ e, e, 0\rangle$	\rightarrow	$-i g, e, 1\rangle$	\rightarrow	$-i g, e, 1\rangle$	\rightarrow	$- e, e, 0\rangle$

Table 4: Cirac-Zoller Gate

Now we discuss the ground state motion. The average excitation within thermal equilibrium is that

$$\bar{n} = \frac{1}{e^{\hbar\omega/k_B T} - 1} \quad (98)$$

The frequency $\omega \sim \text{MHz}$, and the room temperature $T \sim 300\text{K}$, so the average excitation in room temperature $\bar{n} \sim 10^7$. So we need to cool the ions down to $T \sim 10^{-5}\text{K}$.

4 Decoherence and Quantum Error Correction

4.1 Density Matrices

If a pure state $|\Psi_j\rangle$ is prepared with probability p_j , an expectation value of an observable \hat{A} can be calculated

$$\begin{aligned}
 \langle \hat{A} \rangle &= \sum_j p_j \langle \Psi_j | \hat{A} | \Psi_j \rangle \\
 &= \sum_{j,k} p_j \langle \Psi_j | \hat{A} | k \rangle \langle k | \Psi_j \rangle \quad \left(\text{when } \sum_k |k\rangle \langle k| = \mathbb{1} \right) \\
 &= \sum_{j,k} p_j \langle k | \Psi_j \rangle \langle \Psi_j | \hat{A} | k \rangle \\
 &= \sum_k \langle k | \sum_j p_j |\Psi_j\rangle \langle \Psi_j| \hat{A} | k \rangle \\
 &= \text{Tr}(\hat{\rho} \hat{A})
 \end{aligned} \tag{99}$$

with the density matrix

$$\hat{\rho} = \sum_j p_j |\Psi_j\rangle \langle \Psi_j| \tag{100}$$

The properties of density matrix

1. $\hat{\rho}$ is hermitian: $\hat{\rho} = \hat{\rho}^\dagger$
2. $\hat{\rho}$ has unit trace: $\text{Tr } \hat{\rho} = 1$ ($\sum_j p_j = 1$)
3. $\hat{\rho}$ is positive semi-definite

$$\langle \phi | \hat{\rho} | \phi \rangle = \sum_j p_j \langle \phi | \Psi_j \rangle \langle \Psi_j | \phi \rangle = \sum_j p_j |\langle \phi | \Psi_j \rangle|^2 \geq 0 \tag{101}$$

4.1.1 Reduced States

Let's consider a two-body system

$$|\Psi\rangle = \sum_{ij} \Psi_{ij} |\varphi_i\rangle \otimes |\phi_j\rangle \tag{102}$$

For the observable $\hat{A} = \hat{O} \otimes \mathbb{1}$, the expectation value can be calculated

$$\begin{aligned}
 \langle \hat{A} \rangle &= \langle \Psi | \hat{O} \otimes \mathbb{1} | \Psi \rangle \\
 &= \text{Tr}(\hat{O} \otimes \mathbb{1} |\Psi\rangle \langle \Psi|) \\
 &= \text{Tr}(\hat{O} \text{Tr}_2 |\Psi\rangle \langle \Psi|) \\
 &= \text{Tr}(\hat{O} \hat{\rho}_1)
 \end{aligned} \tag{103}$$

where the *reduced density matrix* for system 1

$$\hat{\rho}_1 = \text{Tr}_2 |\Psi\rangle \langle\Psi| \quad (104)$$

The reduced density matrix is

- i) Hermitian
- ii) Positive semidefinite
- iii) $\text{Tr } \hat{\rho} = 1$

4.1.2 Population and Coherence

Consider the initial state $|+\rangle = (|0\rangle + |1\rangle)/\sqrt{2}$, the density matrix

$$\hat{\rho}(0) = |+\rangle \langle+| = \frac{1}{2}(|0\rangle \langle 0| + |1\rangle \langle 1| + |0\rangle \langle 1| + |1\rangle \langle 0|) \quad (105)$$

measurements on $\{|0\rangle, |1\rangle\}$ basis show

$$\langle 0| \hat{\rho}(0) |0\rangle = \frac{1}{2}, \quad \langle 1| \hat{\rho}(0) |1\rangle = \frac{1}{2} \quad (106)$$

Recall the dynamics $\hat{U} |+\rangle$ with $\hat{U} = e^{-i\omega\hat{\sigma}_z t}$ from the Ramsey scheme

$$\begin{aligned} \rho(t) &= \hat{U} \hat{\rho}(0) \hat{U}^\dagger = \frac{1}{2} \left(\hat{U} |0\rangle \langle 0| \hat{U}^\dagger + \hat{U} |1\rangle \langle 1| \hat{U}^\dagger + \hat{U} |0\rangle \langle 1| \hat{U}^\dagger + \hat{U} |1\rangle \langle 0| \hat{U}^\dagger \right) \\ &= \frac{1}{2} (|0\rangle \langle 0| + |1\rangle \langle 1| + e^{-i\omega t} |0\rangle \langle 1| + e^{i\omega t} |1\rangle \langle 0|) \end{aligned} \quad (107)$$

The probability to project on the state $|+\rangle$

$$\langle+| \hat{\rho}(t) |+\rangle = \frac{1}{2}(1 + \cos \omega t) \quad (108)$$

Now we consider the initial state

$$\hat{\rho}_\eta(0) = \frac{1}{2}(|0\rangle \langle 0| + |1\rangle \langle 1| + \eta |0\rangle \langle 1| + \eta |1\rangle \langle 0|) \quad (109)$$

with a real parameter $\eta \in [0, 1]$

$$\hat{\rho}_\eta(t) = \hat{U} \hat{\rho}_\eta(0) \hat{U}^\dagger = \frac{1}{2} (|0\rangle \langle 0| + |1\rangle \langle 1| + \eta e^{-i\omega t} |0\rangle \langle 1| + \eta e^{i\omega t} |1\rangle \langle 0|) \quad (110)$$

The probability to project on the state $|+\rangle$

$$\langle+| \hat{\rho}_\eta(t) |+\rangle = \frac{1}{2}(1 + \eta \cos \omega t) \quad (111)$$

We can write $\hat{\rho}_\eta(0)$ also

$$\begin{aligned} \hat{\rho}_\eta(0) &= \frac{1-\eta}{2} |0\rangle \langle 0| + \frac{1-\eta}{2} |1\rangle \langle 1| + \frac{\eta}{2} (|0\rangle \langle 0| + |1\rangle \langle 1| + |0\rangle \langle 1| + |1\rangle \langle 0|) \\ &= \frac{1-\eta}{2} |0\rangle \langle 0| + \frac{1-\eta}{2} |1\rangle \langle 1| + \frac{\eta}{2} |+\rangle \langle+| \end{aligned} \quad (112)$$

4.1.3 Pure States and Mixed States

A density matrix with a single non-vanishing eigenvalue describes a pure state $|\Psi\rangle\langle\Psi|$; a density matrix with at least two non-vanishing eigenvalues describe a mixed state.

- For pure states $\text{Tr } \hat{\rho}^2 = 1$
- For mixed states $\text{Tr } \hat{\rho}^2 < 1$

4.1.4 Expansion in Operator Basis

The expansion of an operator \hat{A}

$$\begin{aligned}\hat{A} &= \begin{pmatrix} A_{00} & A_{01} \\ A_{10} & A_{11} \end{pmatrix} \\ &= \begin{pmatrix} \frac{1}{2}(A_{00} + A_{11}) + \frac{1}{2}(A_{00} - A_{11}) & \frac{1}{2}(A_{01} + A_{10}) + \frac{1}{2}(A_{01} - A_{10}) \\ \frac{1}{2}(A_{01} + A_{10}) - \frac{1}{2}(A_{01} - A_{10}) & \frac{1}{2}(A_{00} + A_{11}) - \frac{1}{2}(A_{00} - A_{11}) \end{pmatrix} \\ &= \frac{1}{2}(A_{00} + A_{11})\mathbb{1} + \frac{1}{2}(A_{00} - A_{11})\hat{\sigma}_z + \frac{1}{2}(A_{01} + A_{10})\hat{\sigma}_x + i\frac{1}{2}(A_{01} - A_{10})\hat{\sigma}_y \\ &= \frac{1}{2}\text{Tr } \hat{A}\mathbb{1} + \frac{1}{2}\text{Tr}(\hat{A}\hat{\sigma}_z)\hat{\sigma}_z + \frac{1}{2}\text{Tr}(\hat{A}\hat{\sigma}_x)\hat{\sigma}_x + \frac{1}{2}\text{Tr}(\hat{A}\hat{\sigma}_y)\hat{\sigma}_y\end{aligned}\quad (113)$$

where

$$\text{Tr}(\hat{A}\hat{B}^\dagger) = \sum_{ij} \langle i|\hat{A}|j\rangle \langle j|\hat{B}^\dagger|i\rangle = \sum_{ij} A_{ij}B_{ji}^\dagger \quad (114)$$

Using $\hat{\sigma}_i^2 = \mathbb{1}$ and $\hat{\sigma}_i^\dagger = \hat{\sigma}_i$, and consider $\text{Tr } \hat{A} = 1$ (normalised), then

$$\text{Tr}(\hat{A}^2) = \frac{1}{4}\text{Tr } \mathbb{1} + \frac{1}{4}\sum_{i=1}^3 \mu_i^2 \text{Tr } \hat{\sigma}_i^2 = \frac{1}{2} + \frac{1}{2}\sum_{i=1}^3 \mu_i^2 = \frac{1}{2} + \frac{1}{2}\mathbf{u} \cdot \mathbf{u} \quad (115)$$

where $\mathbf{u} = \text{Tr } \hat{A}\hat{\sigma}_i^\dagger$ is the Bloch vector. The operator basis

$$\hat{\sigma}_{1,2,3} = \hat{\sigma}_{x,y,z}, \quad \hat{\sigma}_4 = \mathbb{1} \quad (116)$$

4.2 Open Quantum Systems and Decoherence

4.2.1 Quantum Channels and Open Quantum Dynamics

Let's assume that an atom is the system $a|e\rangle + b|g\rangle$ surrounded by the vacuum fields $|0_\lambda\rangle$. The total system

$$|\Psi(0)\rangle = (a|e\rangle + b|g\rangle)|0_\lambda\rangle \rightarrow |\Psi(t)\rangle = a(\alpha|e\rangle|0_\lambda\rangle + \beta|g\rangle|1\rangle_\lambda) + b|g\rangle|0_\lambda\rangle \quad (117)$$

The density operator of the system

$$\begin{aligned}\hat{\rho}_{\text{system}} &= \text{Tr}_{\text{env}} |\Psi\rangle\langle\Psi| \\ &= (a\alpha|e\rangle + b\beta|g\rangle)(a^*\alpha^*\langle e| + b^*\beta^*\langle g|) + \beta\beta^*|g\rangle\langle g|\end{aligned}\quad (118)$$

In general, the total system $|\Psi\rangle$ and the evolution operator $\hat{U}(t) = \exp(-i\hat{H}t)$. Then the evolution of the total system is given by

$$|\Psi(t)\rangle = \exp(-i\hat{H}t) |\Psi(0)\rangle = \hat{U}(t) |\Psi(0)\rangle \quad (119)$$

Let's assume the initial total system is composed of the system and its environment $\hat{\rho}_{0,\text{total}} = \hat{\rho}_0 \otimes |\phi_{\text{env}}\rangle \langle \phi_{\text{env}}| = |\phi_{\text{env}}\rangle \hat{\rho}_0 \langle \phi_{\text{env}}|$. To find the density operator for the system

$$\begin{aligned} \hat{\rho}(t) &= \text{Tr}_{\text{env}} |\Psi(t)\rangle \langle \Psi(t)| \\ &= \sum_i \langle i| \hat{U}(t) |\Psi(0)\rangle \langle \Psi(0)| \hat{U}^\dagger |i\rangle \\ &= \sum_i \langle i| \hat{U}(t) |\phi_{\text{env}}\rangle \hat{\rho}_0 \langle \phi_{\text{env}}| \hat{U}^\dagger |i\rangle \\ &= \sum_i \hat{F}_i(t) \hat{\rho}_0 \hat{F}_i^\dagger(t) \end{aligned} \quad (120)$$

is called a *quantum channel*, with

$$\hat{F}_i(t) = \langle i| \hat{U}(t) |\phi_{\text{env}}\rangle \quad (121)$$

are called *Kraus operators*, and

$$\sum_i \hat{F}_i^\dagger \hat{F}_i = \mathbb{1} \quad (122)$$

4.2.2 Properties of Quantum Channels

Quantum channels

$$\hat{\Lambda}(\rho) = \sum_i \hat{F}_i \hat{\rho} \hat{F}_i^\dagger \quad (123)$$

have the properties

1) Trace preserving

$$\text{Tr}(\hat{\Lambda}(\rho)) = \text{Tr} \hat{\rho} \quad (124)$$

due to $\sum_i \hat{F}_i^\dagger \hat{F}_i = \mathbb{1}$.

2) Positive maps

$$\langle \Phi | \sum_i \hat{F}_i \hat{\rho} \hat{F}_i^\dagger | \Phi \rangle = \sum_i \langle \Phi | \hat{F}_i \hat{\rho} \hat{F}_i^\dagger | \Phi \rangle = \sum_i \langle \Phi_i | \hat{\rho} | \Phi_i \rangle \geq 0 \quad (125)$$

where $|\Phi_i\rangle = \hat{F}_i^\dagger |\Phi\rangle$.

3) Completely positive maps

$$(\hat{\Lambda} \otimes \mathbb{1})\hat{\rho} = \sum_i (\hat{F}_i \otimes \mathbb{1})\hat{\rho}(\hat{F}_i^\dagger \otimes \mathbb{1}) \quad (126)$$

When this is positive, we call it a completely positive map.

4.2.3 Exemplary Quantum Channels

1) Dissipation channel

Consider the Kraus operators

$$\hat{F}_1 = |0\rangle\langle 0|, \quad \hat{F}_2 = |0\rangle\langle 1| \quad (127)$$

which satisfies $\sum_i \hat{F}_i^\dagger \hat{F}_i = \mathbb{1}$. Let us apply \hat{F}_i to $\hat{\rho}$

$$\hat{F}_1 \hat{\rho} \hat{F}_1^\dagger = \rho_{00} |0\rangle\langle 0| \quad (128)$$

$$\hat{F}_2 \hat{\rho} \hat{F}_2^\dagger = \rho_{11} |0\rangle\langle 0| \quad (129)$$

$$\sum_i \hat{F}_i \hat{\rho} \hat{F}_i^\dagger = (\rho_{00} + \rho_{11}) |0\rangle\langle 0| = |0\rangle\langle 0| \quad (130)$$

We can turn this into a time-dependent process

$$\hat{F}_0(t) = \sqrt{1-p(t)} \mathbb{1}, \quad \hat{F}_1(t) = \sqrt{p(t)} |0\rangle\langle 0|, \quad \hat{F}_2(t) = \sqrt{p(t)} |0\rangle\langle 1| \quad (131)$$

$$\sum_i \hat{F}_i^\dagger(t) \hat{F}_i(t) = (1-p(t)) \mathbb{1} + p(t) |0\rangle\langle 0| + p(t) |1\rangle\langle 1| = \mathbb{1} \quad (132)$$

$$\sum_i \hat{F}_i(t) \hat{\rho}(0) \hat{F}_i^\dagger(t) = (1-p(t)) \hat{\rho}(0) + p(t) |0\rangle\langle 0| \quad (133)$$

2) Dephasing channel

Consider two Kraus operators

$$\hat{F}_1 = \sqrt{\frac{1+p}{2}} \mathbb{1}, \quad \hat{F}_2 = \sqrt{\frac{1-p}{2}} \hat{\sigma}_z \quad (134)$$

with $\hat{\rho} = \rho_{00} |0\rangle\langle 0| + \rho_{11} |1\rangle\langle 1| + \rho_{01} |0\rangle\langle 1| + \rho_{10} |1\rangle\langle 0|$ one obtains

$$\hat{F}_1 \hat{\rho} \hat{F}_1^\dagger = \frac{1+p}{2} (\rho_{00} |0\rangle\langle 0| + \rho_{11} |1\rangle\langle 1| + \rho_{01} |0\rangle\langle 1| + \rho_{10} |1\rangle\langle 0|) \quad (135)$$

$$\hat{F}_2 \hat{\rho} \hat{F}_2^\dagger = \frac{1-p}{2} (\rho_{00} |0\rangle\langle 0| + \rho_{11} |1\rangle\langle 1| - \rho_{01} |0\rangle\langle 1| - \rho_{10} |1\rangle\langle 0|) \quad (136)$$

$$\hat{\Lambda}(\hat{\rho}) = \sum_i \hat{F}_i \hat{\rho} \hat{F}_i^\dagger = \rho_{00} |1\rangle\langle 1| + \rho_{11} |1\rangle\langle 1| + p(\rho_{01} |0\rangle\langle 0| + \rho_{10} |1\rangle\langle 0|) \quad (137)$$

Consider another set of Kraus operators

$$\hat{F}_1 = \sqrt{p} \mathbb{1}, \quad \hat{F}_2 = \sqrt{1-p} |0\rangle\langle 0|, \quad \hat{F}_3 = \sqrt{1-p} |1\rangle\langle 1| \quad (138)$$

Then

$$\hat{\Lambda}(\hat{\rho}) = \sum_i \hat{F}_i \hat{\rho} \hat{F}_i^\dagger = \rho_{00} |0\rangle\langle 0| + \rho_{11} |1\rangle\langle 1| + p(\rho_{01} |0\rangle\langle 1| + \rho_{10} |1\rangle\langle 0|) \quad (139)$$

Consider the Kraus operators

$$\hat{F}_1 = \frac{1}{\sqrt{2}} \exp\left(i\frac{\varphi}{2} \hat{\sigma}_z\right), \quad \hat{F}_2 = \frac{1}{\sqrt{2}} \exp\left(-i\frac{\varphi}{2} \hat{\sigma}_z\right) \quad (140)$$

Then

$$\hat{\Lambda}(\hat{\rho}) = \sum_i \hat{F}_i \hat{\rho} \hat{F}_i^\dagger = \rho_{00} |0\rangle \langle 0| + \rho_{11} |1\rangle \langle 1| + \cos \varphi (\rho_{01} |0\rangle \langle 1| + \rho_{10} |1\rangle \langle 0|) \quad (141)$$

We see that there is non-uniqueness of Kraus operators.

3) Depolarising channel

Consider the Kraus operators

$$\hat{F}_0 = \sqrt{1-p} \mathbb{I}, \quad \hat{F}_1 = \sqrt{\frac{p}{3}} \hat{\sigma}_x, \quad \hat{F}_2 = \sqrt{\frac{p}{3}} \hat{\sigma}_y, \quad \hat{F}_3 = \sqrt{\frac{p}{3}} \hat{\sigma}_z \quad (142)$$

Then

$$\hat{\Lambda}(\hat{\rho}) = \sum_i \hat{F}_i \hat{\rho} \hat{F}_i^\dagger = (1-p) \hat{\rho} + \frac{p}{3} (\hat{\sigma}_x \hat{\rho} \hat{\sigma}_x + \hat{\sigma}_y \hat{\rho} \hat{\sigma}_y + \hat{\sigma}_z \hat{\rho} \hat{\sigma}_z) \quad (143)$$

The depolarising channel brings the original quantum state to the identity \mathbb{I} .

4.2.4 Generalised Measurements

1) Von-Neumann measurements

Consider the operator $\hat{\sigma}_z$

$$\hat{\sigma}_z |0\rangle = +1 |0\rangle, \quad \hat{\sigma}_z |1\rangle = -1 |1\rangle \quad (144)$$

Then

$$\hat{\sigma}_z = |0\rangle \langle 0| - |1\rangle \langle 1| \quad (145)$$

The probability that $\hat{\sigma}_z$ give the result +1

$$P(+1) = \langle 0 | \hat{\rho} | 0 \rangle = \text{Tr } \hat{\rho} |0\rangle \langle 0| \quad (146)$$

In general, we define a projector

$$\hat{P} = |\lambda_i\rangle \langle \lambda_i| \quad (147)$$

The properties of projectors are

- (a) They are Hermitian
- (b) They are positive operators
- (c) They are complete
- (d) They are orthonormal

These also the conditions for the von-Neumann measurements.

2) Generalised measurements

Let's consider the four measurement bases for two qubits

$$|\Phi_1\rangle = \frac{1}{\sqrt{2}} \left(|1\rangle_1 |0\rangle_2 + \tan \theta |0\rangle_1 |0\rangle_2 + \sqrt{1 - \tan^2 \theta} |1\rangle_1 |1\rangle_2 \right)$$

$$\begin{aligned} |\Phi_2\rangle &= \frac{1}{\sqrt{2}} \left(|1\rangle_1 |0\rangle_2 - \tan \theta |0\rangle_1 |0\rangle_2 - \sqrt{1 - \tan^2 \theta} |1\rangle_1 |1\rangle_2 \right) \\ |\Phi_3\rangle &= \sqrt{1 - \tan^2 \theta} |0\rangle_1 |0\rangle_2 - \tan \theta |1\rangle_1 |1\rangle_2 \\ |\Phi_4\rangle &= |0\rangle_1 |1\rangle_2 \end{aligned}$$

where $0 \leq \theta \leq \pi/4$, and the projectors are $|\Phi_i\rangle \langle \Phi_i|$. We assume that we prepare the second qubit in $|0\rangle$. Then the four projectors will result in the projection the first qubit on

$$\begin{aligned} |\Psi_1\rangle &= \frac{1}{\sqrt{2}} (|1\rangle + \tan \theta |0\rangle) \\ |\Psi_2\rangle &= \frac{1}{\sqrt{2}} (|1\rangle - \tan \theta |0\rangle) \\ |\Psi_3\rangle &= \sqrt{1 - \tan^2 \theta} |0\rangle \end{aligned}$$

while $|\Psi_i\rangle \langle \Psi_i|$ satisfy the properties 1,2 and 3 above while they are not orthogonal to each other. We can also consider

$$\hat{\Pi}_0 = (1 - p) |0\rangle \langle 0| + p |1\rangle \langle 1| \quad (148)$$

$$\hat{\Pi}_1 = (1 - p) |1\rangle \langle 1| + p |0\rangle \langle 0| \quad (149)$$

when $p = 0$, then $\hat{\Pi}_0 = |0\rangle \langle 0|$ and $\hat{\Pi}_1 = |1\rangle \langle 1|$.

4.2.5 System-Environment Interaction

The Hamiltonian for system-environment interaction

$$H_{SE} = \sum_i |i\rangle \langle i| \otimes \hat{h}_i \quad (150)$$

so the propagator

$$\hat{U}_i = \exp \left(-it \sum_i |i\rangle \langle i| \otimes \hat{h}_i \right) = \sum_i |i\rangle \langle i| \otimes \hat{U}_i(t) \quad (151)$$

where $\hat{U}_i = \exp(-i\hat{h}_i t)$. Applying the evolution operators to the system-environment

$$|\psi(t)\rangle = \sum_i \psi_i |i\rangle \otimes \hat{U}_i(t) |\Phi\rangle \quad (152)$$

where

$$|\psi(0)\rangle = \sum_i \psi_i |i\rangle_{\text{sys}} \otimes |\Phi\rangle_{\text{env}} \quad (153)$$

The density operator for the system

$$\hat{\rho} = \sum_{ij} \psi_i \psi_j^* |i\rangle \langle j| \langle \Phi| \hat{U}_j^\dagger(t) \hat{U}_i(t) |\Phi\rangle \quad (154)$$

4.3 Error Correction

4.3.1 Classical Error Correction

For instance, three qubits realise one logical bit

$$\begin{matrix} \bar{0} \\ \text{logical qubit} \end{matrix} = \begin{matrix} 000 \\ \text{physical qubits} \end{matrix}, \quad \bar{1} = 111 \quad (155)$$

The probability of two or three errors (error probability $p \ll 1$)

$$p^3 + 3p^2(1 - p) \ll 1 \quad (156)$$

4.3.2 Quantum Error Correction

1. **Bit flip error** $|0\rangle \rightarrow |1\rangle, |1\rangle \rightarrow |0\rangle$

Let us measure on $\{|0\rangle, |1\rangle\}$ basis, the qubit can be expressed as

$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle \quad (157)$$

then the logical qubits

$$|\bar{\psi}\rangle = \alpha|000\rangle + \beta|111\rangle \quad (158)$$

Let us find out how to achieve the logical qubits. Consider the quantum circuits fig.6(a)

- 1st CNOT

$$\text{CNOT}(\alpha|0\rangle + \beta|1\rangle)|0\rangle = \alpha|00\rangle + \beta|11\rangle \quad (159)$$

- 2nd CNOT

$$\text{CNOT}(\alpha|00\rangle + \beta|11\rangle)|0\rangle = \alpha|000\rangle + \beta|111\rangle \quad (160)$$

Then we consider the bit flip errors (fig.6(b)). If a bit-flip error happens on Q_1 , then the state before first CNOT gate is $\alpha|100\rangle + \beta|011\rangle$

- 1st CNOT

$$\alpha|100\rangle|1\rangle + \beta|011\rangle|0\rangle \quad (161)$$

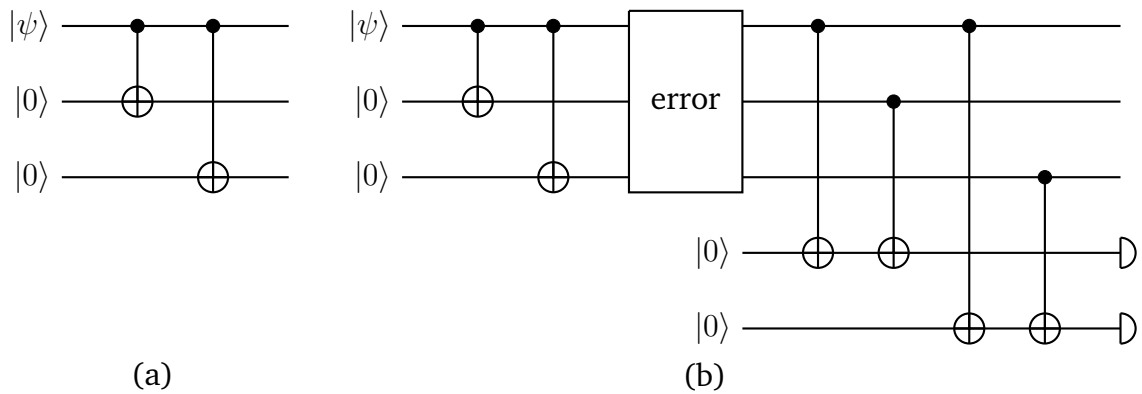


Figure 6: (a) Prepare the logical qubits. (b) Bit flip error detection.

- 2nd CNOT

$$\alpha |100\rangle |1\rangle + \beta |011\rangle |1\rangle = (\alpha |100\rangle + \beta |011\rangle) |1\rangle \quad (162)$$

Similarly, we have

Error	data⟩ ancilla⟩
no	$(\alpha 000\rangle + \beta 111\rangle) 00\rangle$
Q_1	$(\alpha 100\rangle + \beta 011\rangle) 11\rangle$
Q_2	$(\alpha 010\rangle + \beta 101\rangle) 10\rangle$
Q_3	$(\alpha 001\rangle + \beta 110\rangle) 01\rangle$

Table 5: The results of the bit flip detection quantum circuits.

To correct the bit-flip errors, we can perform $\hat{\sigma}_x$ on the qubits, which can realise $|0\rangle \rightarrow |1\rangle$ and $|1\rangle \rightarrow |0\rangle$.

2. Phase flip error $|+\rangle \rightarrow |-\rangle, |-\rangle \rightarrow |+\rangle$

We encode the information on basis $\{|+\rangle, |-\rangle\}$, then the physical qubit

$$|\psi\rangle = \alpha |+\rangle + \beta |-\rangle \quad (163)$$

and the logical qubits

$$|\bar{0}\rangle = |+++ \rangle, \quad |\bar{1}\rangle = |-- - \rangle \quad (164)$$

and

$$|\bar{\psi}\rangle = \alpha |+++ \rangle + \beta |-- - \rangle \quad (165)$$

which are generated by Hadamard gate (see in fig.7)

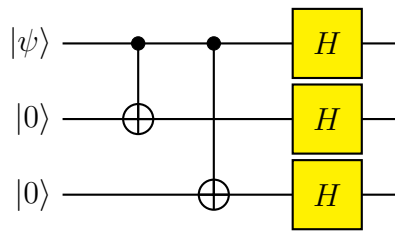


Figure 7: Prepare the logical qubits for phase error detection.

4.4 Stabiliser Formalism

4.4.1 Stabilisers

A state $|\psi\rangle$ is defined to be stabilised by operators \hat{K} if it is a +1 eigenstate of \hat{K}

$$\hat{K} |\psi\rangle = |\psi\rangle \quad (166)$$

For qubit systems, let's consider the Pauli group

$$P = \left\{ \pm \mathbb{1}, \pm i \mathbb{1}, \pm \hat{X}, \pm i \hat{X}, \pm \hat{Y}, \pm i \hat{Y}, \pm \hat{Z}, \pm i \hat{Z} \right\} \quad (167)$$

For N qubits

$$P_N = P^{\otimes N} \quad (168)$$

An N -qubit stabiliser state $|\psi\rangle$ is then defined by the N generators of an Abelian (all elements commute) sub-group, G , of the N -qubit Pauli group

$$G = \left\{ \hat{K}_i | \hat{K}_i |\psi\rangle = |\psi\rangle, [\hat{K}_i, \hat{K}_j] = 0, \forall (i, j) \right\} \quad (169)$$

where \hat{K}_i is Hermitian and $\hat{K}_i \hat{K}_i = \mathbb{1}$. If we encode k logical qubits into n physical qubits, the number of stabilisers is $l = n - k$.

Example.

1)

$$\frac{1}{\sqrt{2}}(|000\rangle + |111\rangle) \quad (170)$$

is stabilised by

$$\hat{K}_1 = \hat{X} \hat{X} \hat{X}, \quad \hat{K}_2 = \hat{Z} \hat{Z} \mathbb{1}, \quad \hat{K}_3 = \mathbb{1} \hat{Z} \hat{Z} \quad (171)$$

Note. Where is $\hat{K}_4 = \hat{Z} \mathbb{1} \hat{Z}$? The four Paulis $\hat{K}_1, \hat{K}_2, \hat{K}_3, \hat{K}_4$ are not mutual exclusive generators, because $\hat{K}_4 = \hat{K}_2 \hat{K}_3$ (or any other permutation). Hence, to generate this specific stabilizer we only need \hat{K}_1 and any two of the set $\hat{K}_2, \hat{K}_3, \hat{K}_4$.

2) Shor code

$$|\bar{0}\rangle = \frac{1}{2\sqrt{2}}(|000\rangle + |111\rangle) \otimes (|000\rangle + |111\rangle) \otimes (|000\rangle + |111\rangle) \quad (172)$$

$$|\bar{1}\rangle = \frac{1}{2\sqrt{2}}(|000\rangle - |111\rangle) \otimes (|000\rangle - |111\rangle) \otimes (|000\rangle - |111\rangle) \quad (173)$$

3) Consider the following 7-qubit code

$$|\bar{0}\rangle = \frac{1}{\sqrt{8}}(|0000000\rangle + |1010101\rangle + |0110011\rangle + |0001111\rangle \\ + |1011010\rangle + |0111100\rangle + |1101001\rangle + |1100110\rangle) \quad (174)$$

$$|\bar{1}\rangle = \frac{1}{\sqrt{8}}(|1111111\rangle + |0101010\rangle + |1001100\rangle + |1110000\rangle \\ + |0100101\rangle + |1000011\rangle + |0010110\rangle + |0011001\rangle) \quad (175)$$

The stabiliser set for the 7-qubit code is specified by

$$\hat{K}_1 = \mathbb{1} \mathbb{1} \mathbb{1} \hat{X} \hat{X} \hat{X} \hat{X}, \quad \hat{K}_2 = \hat{X} \mathbb{1} \hat{X} \mathbb{1} \hat{X} \mathbb{1} \hat{X} \quad (176)$$

$$\hat{K}_3 = \mathbb{1} \hat{X} \hat{X} \mathbb{1} \mathbb{1} \hat{X} \hat{X}, \quad \hat{K}_4 = \mathbb{1} \mathbb{1} \mathbb{1} \hat{Z} \hat{Z} \hat{Z} \hat{Z} \quad (177)$$

$$\hat{K}_5 = \hat{Z} \mathbb{1} \hat{Z} \mathbb{1} \hat{Z} \mathbb{1} \hat{Z}, \quad \hat{K}_6 = \mathbb{1} \hat{Z} \hat{Z} \mathbb{1} \mathbb{1} \hat{Z} \hat{Z} \quad (178)$$

4.4.2 Error Measurements

By the definition of stabiliser codes

$$\hat{K}_i |\psi\rangle = |\psi\rangle \quad (179)$$

The errors are assumed to happen in the form $\hat{X}, \hat{Y}, \hat{Z}$. Denoting the error by \hat{E} , we know that \hat{E} and \hat{K} are Pauli operators, so

$$[\hat{K}_i, \hat{E}] = 0, \quad \{\hat{K}_i, \hat{E}\} = 0 \quad (180)$$

we can write

$$\hat{K} \hat{E} |\psi\rangle = -\hat{E} \hat{K} |\psi\rangle = -\hat{E} |\psi\rangle \quad (181)$$

Since an error-free state is already a +1 eigenstate of all the stabilisers, errors which anticommute with any of the stabilisers will flip the relevant eigenstate giving the eigenvalue -1.

4.4.3 Preparation of Eigenstates for a Pauli group operator

Consider fig.8. Before the measurement

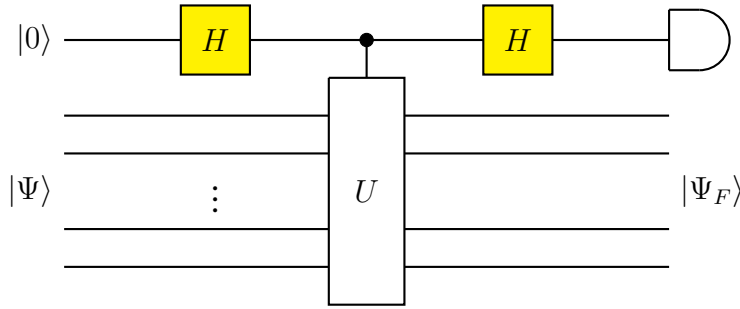


Figure 8: Preparation of eigenstates for \hat{U} , where \hat{U} is a unitary gate, and $\hat{U}^2 = \mathbb{1}$.

$$\begin{aligned} |0\rangle |\Psi\rangle &\xrightarrow{H} \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) |\Psi\rangle \\ &\xrightarrow{C-U} \frac{1}{\sqrt{2}}(|0\rangle |\Psi\rangle + |1\rangle U |\Psi\rangle) \\ &\xrightarrow{H} \frac{1}{2} [(|0\rangle + |1\rangle) |\Psi\rangle + (|0\rangle - |1\rangle) U |\Psi\rangle] \\ &= \frac{1}{2} [|0\rangle (|\Psi\rangle + U |\Psi\rangle) + |1\rangle (|\Psi\rangle - U |\Psi\rangle)] \end{aligned} \quad (182)$$

If the measurement outcome is 0

$$|\Psi_{F1}\rangle = \frac{1}{\sqrt{2}}(|\Psi\rangle + U |\Psi\rangle) \quad (183)$$

If the measurement outcome is 1

$$|\Psi_{F2}\rangle = \frac{1}{\sqrt{2}}(|\Psi\rangle - U |\Psi\rangle) \quad (184)$$

The two states are eigenstates of \hat{U}

$$U |\Psi_{F1}\rangle = \frac{1}{\sqrt{2}}(U |\Psi\rangle + U^2 |\Psi\rangle) = \frac{1}{\sqrt{2}}(|\Psi\rangle + U |\Psi\rangle) = |\Psi_{F1}\rangle \quad (185)$$

$$U |\Psi_{F2}\rangle = \frac{1}{\sqrt{2}}(U |\Psi\rangle - U^2 |\Psi\rangle) = -\frac{1}{\sqrt{2}}(|\Psi\rangle - U |\Psi\rangle) = -|\Psi_{F2}\rangle \quad (186)$$

For instance, if $\hat{U} = \hat{K}_1$, we generate an eigenstate of \hat{K}_1 .