

GPU Teaching Kit

Accelerated Computing



Module 23 – Dynamic Parallelism

Lecture 23 - In depth study of Dynamic Parallelism

Dynamic Parallelism

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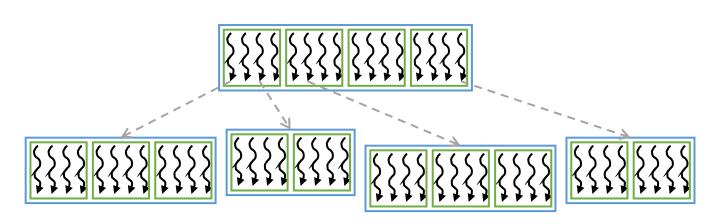


Objective

- To learn about CUDA Dynamic Parallelism
- Applications that benefit from Dynamic parallelism
- Dynamic parallelism in action with BFS algorithm
- Dynamic parallelism Optimization

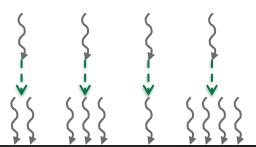
Dynamic Parallelism

 CUDA dynamic parallelism refers to the ability of threads executing on the GPU to launch new grids



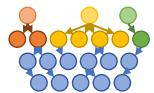
Nested Parallelism

- Dynamic parallelism is useful for programming applications with nested parallelism where each thread discovers more work that can be parallelized
- Dynamic parallelism is particularly useful when the amount of nested work is dynamically determined at execution time, so enough threads cannot be launched up front



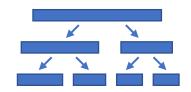
Applications of Dynamic Parallelism

- Applications whose amount of nested work may be unknown before execution time:
 - Nested parallel work is irregular (varies across threads)
 - e.g., graph algorithms (each vertex has a different #neighbors)
 - e.g., Bézier curves (each curve needs different #points to draw)
 - Nested parallel work is recursive with datadependent depth
 - e.g., tree traversal algorithms (e.g., quadtrees and octrees)
 - e.g., divide and conquer algorithms (e.g., quicksort)

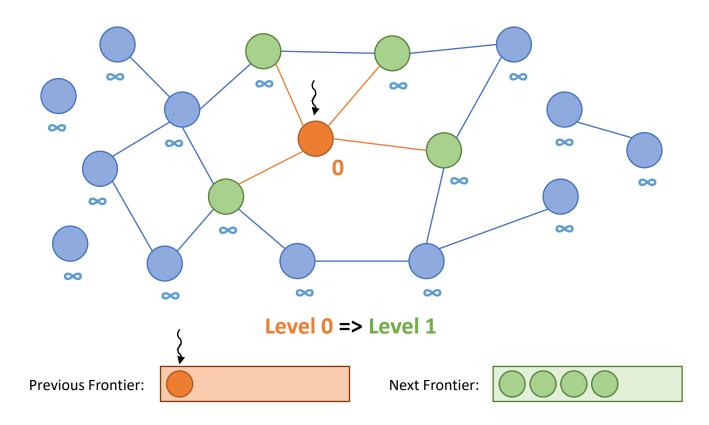




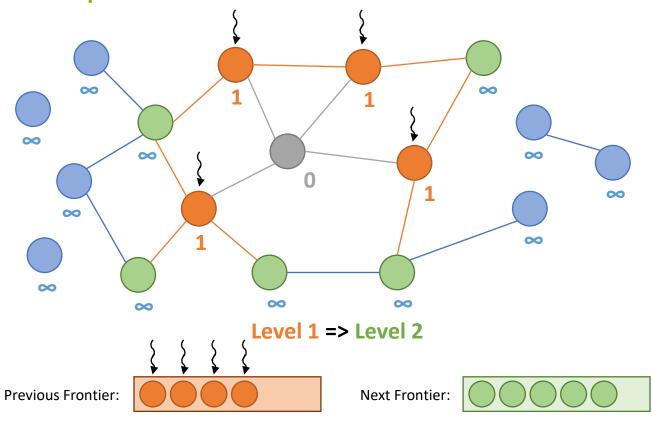




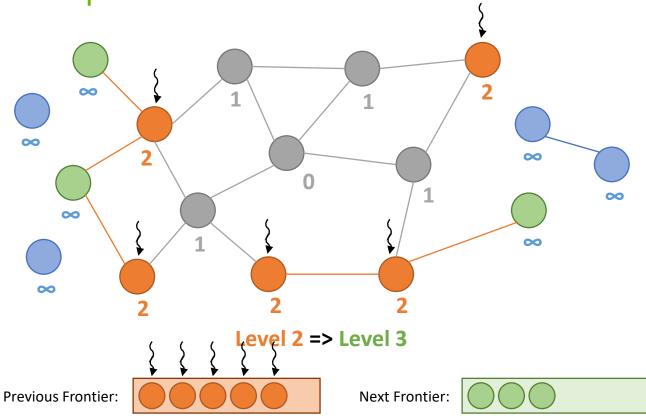
Example: BFS

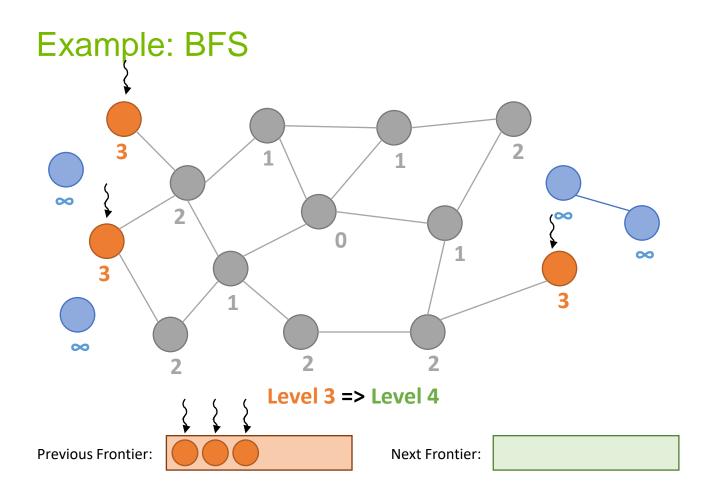


Example: BFS



Example: BFS





BFS Code

```
_global___ void bfs_kernel(CSRGraph csrGraph, unsigned int* nodeLevel, unsigned int* prevFrontier,
                  unsigned int* currFrontier, unsigned int numPrevFrontier, unsigned int*
numCurrFrontier.
                  unsigned int level) {
   unsigned int i = blockIdx.x*blockDim.x + threadIdx.x;
   if(i < numPrevFrontier) {</pre>
        unsigned int node = prevFrontier[i];
        unsigned int start = csrGraph.nodePtrs[node];
        unsigned int numNeighbors = csrGraph.nodePtrs[node + 1] - start;
        for(unsigned int i = 0; i < numNeighbors; ++i) {</pre>
            unsigned int edge = start + i:
            unsigned int neighbor = csrGraph.neighbors[edge];
            if(atomicCAS(&nodeLevel[neighbor], UINT_MAX, level) == UINT_MAX) { // Not previously visited
                unsigned int currFrontierIdx = atomicAdd(numCurrFrontier, 1);
                currFrontier[currFrontierIdx] = neighbor;
            }
                                                     Loop over
                                                neighbors can be
```

parallelized

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Dynamic Parallelism API

- The device code for calling a kernel to launch a grid is the same as the host code
- Memory is needed for buffering grid launches that have not started executing
 - The limit on the number of dynamic launches is referred to as the pending launch count
 - By default, the runtime supports 2048 launches, and exceeding this limit will cause an error
 - The limit can be increased as follows:

cudaDeviceSetLimit(cudaLimitDevRuntimePendingLaunchCount, < new limit >);

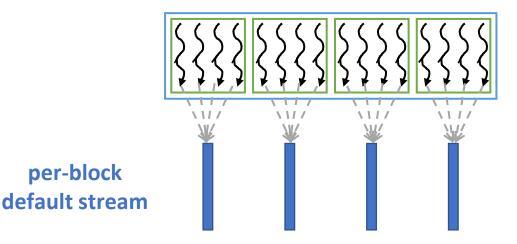


BFS Code with Dynamic Parallelism

```
__qlobal__ void bfs_neighbors_kernel(CSRGraph csrGraph, unsigned int* nodeLevel, unsigned int*
currFrontier.
        unsigned int* numCurrFrontier, unsigned int level, unsigned int start, unsigned int numNeighbors)
                                                                           Loop index becomes a
   unsigned int i = blockIdx.x*blockDim.x + threadIdx.x;
                                                                                 thread index
   if(i < numNeighbors) {</pre>
        unsigned int edge = start + i;
        unsigned int neighbor = csrGraph.neighbors[edge]:
        if(atomicCAS(&nodeLevel[neighbor], UINT_MAX, level) == UINT_MAX) { // Not previously visited
            unsigned int currFrontierIdx = atomicAdd(numCurrFrontier, 1);
            currFrontier[currFrontierIdx] = neighbor;
__global__ void bfs_kernel(CSRGraph csrGraph, unsigned int* nodeLevel, unsigned int* prevFrontier,
                  unsigned int* currFrontier, unsigned int numPrevFrontier, unsigned int* numCurrFrontier,
                  unsigned int level) {
   unsigned int i = blockIdx.x*blockDim.x + threadIdx.x;
   if(i < numPrevFrontier) {</pre>
                                                                                     Loop becomes a
        unsigned int node = prevFrontier[i]:
        unsigned int start = csrGraph.nodePtrs[node]:
                                                                                         kernel call
        unsigned int numNeighbors = csrGraph.nodePtrs[node + 1] - start;
        unsigned int numThreadsPerBlock = 1024;
        unsigned int numBlocks = (numNeighbors + numThreadsPerBlock 1)/numThreadsPerBlock;
        bfs_neighbors_kernel <<< numBlocks, numThreadsPerBlock
                         (csrGraph, nodeLevel, currFrontier, numCurrFrontier, level, start, numNeighbors);
```

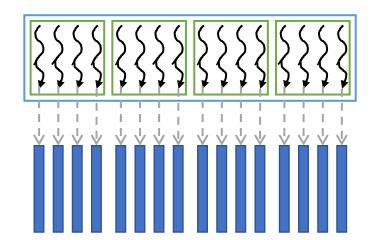
Streams

- Recall: without specifying a stream when calling a kernel, grids get launched into a default stream
- For device launches, threads in the same block share the same default stream
 - Launches by threads in the same block are serialized



Per-Thread Streams

- Parallelism can be improved by creating a different stream per thread
 - Approach #1: Use stream API just like on host
 - Approach #2: Use compiler flag --default-stream per-thread



per-thread stream

Optimizations

– Pitfalls:

- Launching very small grids may not be worth the overhead (more efficient to serialize)
- Launching too many grids causes queueing delays
- Optimization: apply a threshold to the launch
 - Only launch the large grids that are worth the overhead and serialize the rest
 - Threshold value is data dependent and can be tuned
- Optimization: aggregate launches
 - Have one thread collect the work of multiple threads and launch a single grid on their behalf

BFS Code with Threshold

```
global void bfs kernel(CSRGraph csrGraph, unsigned int* nodeLevel, unsigned int* prevFrontier.
                  unsigned int* currFrontier, unsigned int numPrevFrontier, unsigned int*
numCurrFrontier,
                  unsigned int level) {
    unsigned int i = blockIdx.x*blockDim.x + threadIdx.x;
    if(i < numPrevFrontier) {</pre>
                                                                                   Check if the
        unsigned int node = prevFrontier[i];
                                                                                 threshold is met
        unsigned int start = csrGraph.nodePtrs[node]:
        unsigned int numNeighbors = csrGraph.nodePtrs[node + 1] - start;
        if(numNeighbors > 1200) {__
            unsigned int numThreadsPerBlock = 1024;
            unsigned int numBlocks = (numNeighbors + numThreadsPerBlock - 1)/numThreadsPerBlock;
            bfs neighbors kernel <<< numBlocks. numThreadsPerBlock >>>
                         (csrGraph, nodeLevel, currFrontier, numCurrFrontier, level, start,
numNeighbors);
        } else {
            for(unsigned int i = 0; i < numNeighbors; ++i) {</pre>
                unsigned int edge = start + i:
                unsigned int neighbor = csrGraph.neighbors[edge];
                if(atomicCAS(&nodeLevel[neighbor], UINT_MAX, level) == UINT_MAX) {
                    unsigned int currFrontierIdx = atomicAdd(numCurrFrontier, 1);
                    currFrontier[currFrontierIdx] = neighbor;
                }
            }
        }
```

Offloading Host Driver (Control) Code

- In some applications, the host code that drives the computation launches multiple consecutive grids to synchronize across all threads between launches
 - e.g., BFS launches a new grid for each level
- Another application of dynamic parallelism is to offload the driver code to the device
 - Main advantage is to free up the host to do other things

BFS Driver Kernel with Dynamic Parallelism

```
__global__ void bfs_driver_kernel(CSRGraph csrGraph, unsigned int* nodeLevel, unsigned int* prevFrontier,
    unsigned int* currFrontier, unsigned int* numCurrFrontier) {
   unsigned int numPrevFrontier = 1;
   unsigned int numThreadsPerBlock = 256;
   for(unsigned int level = 1: numPrevFrontier > 0: ++level) {
        // Visit nodes in previous frontier
        *numCurrFrontier = 0;
        unsigned int numBlocks = (numPrevFrontier + numThreadsPerBlock - 1)/numThreadsPerBlock:
        bfs_kernel <<< numBlocks, numThreadsPerBlock >>>
               (csrGraph, nodeLevel, prevFrontier, currFrontier, numPrevFrontier, numCurrFrontier,
level);
        cudaDeviceSynchronize();
        // Swap buffers
        unsigned int* tmp = prevFrontier;
        prevFrontier = currFrontier;
        currFrontier = tmp;
        numPrevFrontier = *numCurrFrontier:
    }
```

Launch a single thread to drive the computation

Host code:

bfs_driver_kernel <<< 1, 1 >>> (csrGraph, nodeLevel, prevFrontier, currFrontier, numCurrFrontier);

Memory Visibility

Operations on **global memory** made by a parent thread before the launch are visible to the child

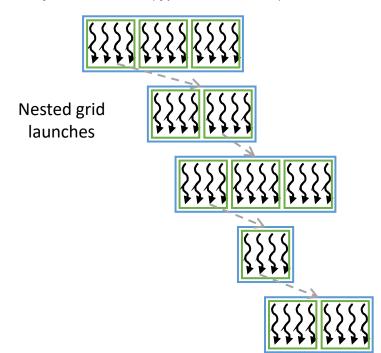
Operations made by the child are visible to the parent after the child returns and the parent has synchronized

A thread's **local memory** and a block's **shared memory** cannot be accessed by child threads

Child threads launched by a parent thread may run on a different SM

Nesting Depth

- The nesting depth refers to how deeply dynamically launched grids may launch other grids
 - Determined by the hardware (typical value is 24)





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