6.004 Worksheet Questions L15 – Memory Hierarchy

Keep the most often-used data in a small, fast SRAM (often local to CPU chip). The reason this strategy works: LOCALITY.

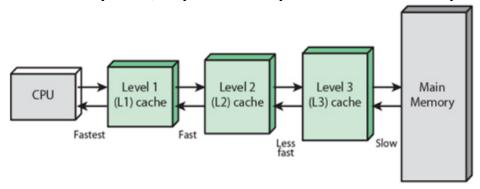
- Temporal locality: If a location has been accessed recently, it is likely to be accessed (reused) soon
- Spatial locality: If a location has been accessed recently, it is likely that nearby locations will be accessed soon

AMAT(Average Memory Access Time) = HitTime + MissRatio * MissPenalty



Problem 1 *

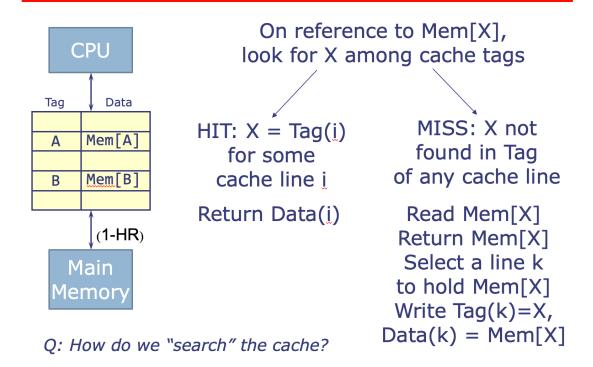
Belly Eyelash is designing a processor and is analyzing the performance of different numbers of cache levels. Without any caches, Belly's main memory has an access time of 140 cycles.



- (A) As a test, Belly adds a 32 KB Level 1 (L1) cache that has single-cycle reads/writes, and runs a long computation, during which she observes a new AMAT of 10. What is the hit ratio for the Level 1 cache during this test?
- (B) Next, Belly adds a 256 KB Level 2 (L2) cache between the L1 cache and main memory. The L2 cache takes 4 cycles to decide if the a memory access is a hit or a miss. After running the same computation with this new memory hierarchy, Belly observes an improved AMAT of 1.45. Assume that the hit rate for the L1 cache is the same as in (A). What is the hit ratio for the Level 2 cache during this test?

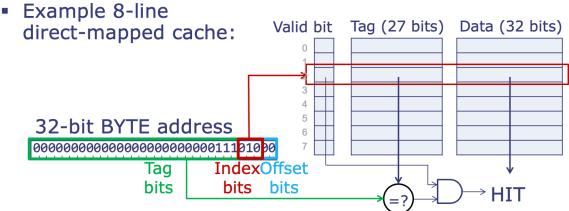
(C) Finally, Belly adds a 10 MB Level 3 (L3) cache between the L2 cache and main memory. After running the same test as before, she observes that the main memory was accessed once for every 140 accesses to the L3 cache. If Belly wants to achieve an AMAT of 1.3 for this computation, what is the maximum number of clock cycles that the Level 3 cache can take to decide if a memory access is a hit or a miss? Assume that the L1 and L2 caches can not be changed and they have the same hit rates observed in (A) and (B).

Basic Cache Algorithm (Reads)



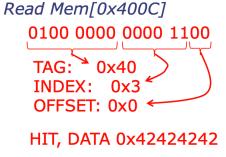
Direct-Mapped Caches

- Each word in memory maps into a single cache line
- Access (for cache with 2^W lines):
 - Index into cache with W address bits (the index bits)
 - Read out valid bit, tag, and data
 - If valid bit == 1 and tag matches upper address bits, HIT



Example: Direct-Mapped Caches

64-line direct-mapped cache \rightarrow 64 indexes \rightarrow 6 index bits



Would 0x4008 hit?

INDEX: $0x2 \rightarrow tag \ mismatch$

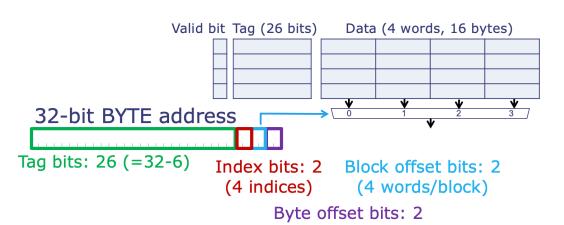
→ MISS

Valid bit		Tag (24 bits)	Data (32 bits)	
0	1	0x000058	0xDEADBEEF	
1	1	0x000058	0x00000000	
2	1	0x000058	0x00000007	
3	1	0x000040	0x42424242	
4	0	0x000007	0x6FBA2381	
	:	:	:	
63	1	0x000058	0xF7324A32	

Part of the address (index bits) is encoded in the location Tag + Index bits unambiguously identify the data's address

Block Size

- Take advantage of spatial locality: Store multiple words per data line
 - Always fetch entire block (multiple words) from memory
 - Another advantage: Reduces size of tag memory!
 - Potential disadvantage: Fewer indices in the cache
- Example: 4-block, 16-word direct-mapped cache



Problem 2

The RISC-V Engineering Team is working on the design of a cache. They've decided that the cache will have a total of $2^{10} = 1024$ data words, but are still thinking about the other aspects of the cache architecture.

First assume the team chooses to build a direct-mapped cache with a block size of 4 words.

(A) Please answer the fe	following questions:
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Number of lines in the cache:	
Number of bits in the tag field for each cache entry:	

(B) This cache takes 2 clock cycles to determine if a memory access is a hit or a miss and, if it's a hit, return data to the processor. If the access is a miss, the cache takes 20 additional clock cycles to fill the cache line and return the requested word to the processor. If the hit rate is 90%, what is the processor's average memory access time in clock cycles?

Average memory access time assuming 90% hit rate (clock cycles): _____

Problem 3 *

(A) The timing for a particular cache is as follows: checking the cache takes 1 cycle. If there's a hit the data is returned to the CPU at the end of the first cycle. If there's a miss, it takes 10 *additional* cycles to retrieve the word from main memory, store it in the cache, and return it to the CPU. If we want an average memory access time of 1.4 cycles, what is the minimum possible value for the cache's hit ratio?

	Minimum	possible	value of hit ratio:	
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(B) If the cache block size, i.e., words/cache line, is doubled but the total number of data words in the cache is unchanged, how will the following cache parameters change? Please circle the best answer.

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# of offset bits: UNCHANGED ... +1 ... -1 ... 2x ... 0.5x ... CAN'T TELL

# of tag bits: UNCHANGED ... +1 ... -1 ... 2x ... 0.5x ... CAN'T TELL

# of cache lines: UNCHANGED ... +1 ... -1 ... 2x ... 0.5x ... CAN'T TELL
```

Consider a direct-mapped cache with 64 total data words with 1 word/cache line.

(C) If cache line number 5 is valid and its tag field has the value 0x1234, what is the address in main memory of the data word currently residing in cache line 5?

Main memory address of data word in cache line 5: 0x_____

Problem 4 ★

	Co	nsider a	direct-mapped	d cache w	ith 64	total data	a words	with 1	word/cache	line.
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The program shown on the right repeatedly executes an inner loop that sums the 16 elements of an array that is stored starting in location 0x310. The program is executed for many iterations, then a measurement of the cache statistics is made during one iteration through all the code, i.e., starting with the execution of the instruction labeled outer_loop: until just before the next time that instruction is executed.	mv x1, x0 loop: subi x4, x4, 1 slli x2, x4, 2 lw x3, 0x310(x2) add x1, x1, x3 bne x4, x0, loop	<pre>// tell assembler to start at // address 0 // initialize loop index J // x1 holds sum, initially 0 // add up elements in array // decrement index // convert to byte offset // load value from A[J] // add to sum // loop until all words are summed // perform test again!</pre>			
(A) In total, how many instruction fetches occur during one complete iteration of the outer loop? How many data reads?					
(B) How many instruction fetch misses of How many data read misses? Hint:	Number of instructions				
(C) What is the hit ratio measured after of	one complete iteration of	f the outer loop? Hit ratio:			