

Today's Lecture

* Least Squares (LS)

* Recursive Least Squares (RLS)

Announcements

Homework 6, due Monday 14thFinal, Wednesday 16th 11:30 - 2:30 PMReview session, Monday 14th, 10:00 - 11:30 am

Least Squares

$$\min_x \|Ax - b\|$$

In our setup, we want to determine $\{h[0], \dots, h[N-1]\}$ that minimizes $\sum_{i=0}^{L-1} e[i]^2$ where $L \geq N$. $\Rightarrow d[n] = \sum_{k=0}^{N-1} h[k]x[n-k] + e[n]$

$$\begin{array}{c} \begin{bmatrix} e[0] \\ \vdots \\ e[L-1] \end{bmatrix} \\ L \times 1 \end{array} = \begin{array}{c} \begin{bmatrix} d[0] \\ \vdots \\ d[L-1] \end{bmatrix} \\ L \times 1 (b) \end{array} - \begin{array}{c} \begin{bmatrix} x[0] & x[-1] & \dots & x[-N+1] \\ \vdots & & & \vdots \\ x[L-1] & x[L-2] & \dots & x[-N+L] \end{bmatrix} \\ L \times N (A) \end{array} \begin{array}{c} \begin{bmatrix} h[0] \\ \vdots \\ h[N-1] \end{bmatrix} \\ N \times 1 (x) \end{array}$$

$$e = d - Ah \rightarrow \text{minimize } \|e\|_2^2 = \|d - Ah\|_2^2$$

$$= (d - Ah)^T (d - Ah)$$

$$= d^T d - d^T A h - h^T A^T d + h^T A^T A h$$

$$\text{Gradient} = \underline{0} \Rightarrow \frac{\partial \|e\|_2^2}{\partial h[j]} = 0$$

$$\nabla_h \|e\|_2^2 = -A^T d - A^T d + 2A^T A h = \underline{0}$$

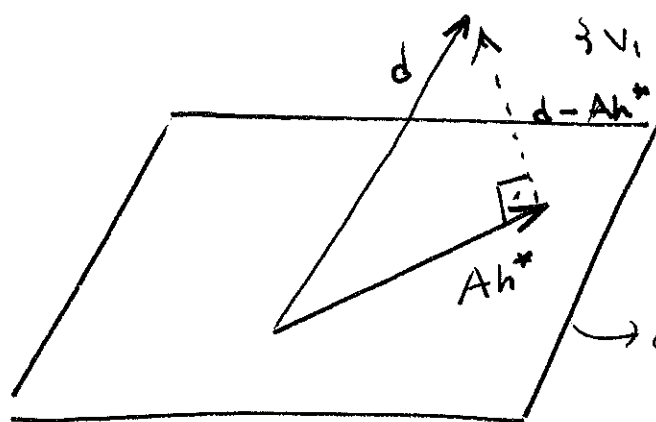
$$A^T A h = A^T d$$

$$A: L \times N \quad (L > N)$$

$$\Rightarrow h^* = \underbrace{(A^T A)^{-1} A^T d}_{\text{pseudo inverse of } A.}$$

Column space of A

$A = [\underline{v}_1 \dots \underline{v}_N]$ → column space is the vector space generated by linear combinations of the column vectors $\{\underline{v}_1, \dots, \underline{v}_N\}$



$$Ah^* = [\underline{v}_1 \dots \underline{v}_N] h^*$$

→ column space of A .

Error is orthogonal to the column space of A .

$$Ah^* \perp (d - Ah^*) \Rightarrow A^T (d - Ah^*) = 0$$

$$d = \underbrace{Ah^*}_{\text{projection to column space of } A.} + \underbrace{(d - Ah^*)}_{\text{error}}$$

$$A^T A = \begin{bmatrix} x[0] & x[1] & \dots & x[L-1] \\ x[-1] & x[0] & \dots & x[L-2] \\ \vdots & \vdots & \ddots & \vdots \\ x[-(N-1)] & x[-(N-2)] & \dots & x[L-N] \end{bmatrix} \begin{bmatrix} x[0] & x[-1] & \dots & x[-(N-1)] \\ x[1] & x[0] & \dots & x[-N] \\ \vdots & \vdots & \ddots & \vdots \\ x[L-1] & x[L-2] & \dots & x[L-N] \end{bmatrix}$$

$N \times L$ $L \times N$

$$= \begin{bmatrix} x[0]^2 + x[1]^2 + \dots + x[L-1]^2 & \dots & \dots \\ x[-1]x[0] + x[0]x[1] + x[1]x[2] + \dots + x[L-2]x[L-1] & \dots & \dots \\ \vdots & \ddots & \vdots \end{bmatrix}$$

$\overset{\text{WSS}}{\approx} LR$ \rightarrow autocorrelation

$$(A^T A) h^* = A^T d \Rightarrow LR h^* = A^T d \approx Lp \rightarrow \text{cross correlation}$$

$$R h^* \approx p \Rightarrow h^* \approx R^{-1} p$$

Wiener-Hopf.

* This is a deterministic approximation to the Wiener filter (if the process is WSS)

For LS, $h^* = (A^T A)^{-1} A^T d$

We want to get $h_{n+1}^* = (A_{n+1}^T A_{n+1})^{-1} A_{n+1}^T d_{n+1}$

from $h_n^* = (A_n^T A_n)^{-1} A_n^T d_n$

We can obtain h_{n+1}^* from h_n^* using Recursive Least Squares (RLS)

$$\min \sum_{i=0}^{n-1} \lambda^{(n-1)-i} e^2[i] = C(h_n)$$

$$0 < \lambda < 1 \quad (\text{forgetting factors})$$

$$= 1 \cdot e^2[n-1] + \lambda e^2[n-2] + \lambda^2 e^2[n-3] + \dots$$

$$\frac{\partial C(h_{n-1})}{\partial h_{n-1}(k)} = \sum_{i=0}^{n-1} 2 \lambda^{(n-1)-i} e(i) \frac{\partial e(i)}{\partial h_{n-1}(k)}$$

$$e(i) = d(i) - \sum_{l=0}^{N-1} h_{n-1}(l) x[i-l], \quad i=0, \dots, n-1$$

$$\frac{\partial e[i]}{\partial h_{n-1}(k)} = -x[i-k]$$

$$\begin{aligned} \frac{\partial C}{\partial h_{n-1}(k)} &= - \sum_{i=0}^{n-1} 2 \lambda^{(n-1)-i} \left[d(i) - \sum_{l=0}^{N-1} h_{n-1}(l) x[i-l] \right] x[i-k] \\ &= 0 \end{aligned}$$

$$\Rightarrow \sum_{i=0}^{n-1} \lambda^{(n-1-i)} d[i] x[i-k] = \sum_{l=0}^{n-1} h_{n-1}(l) \left[\sum_{i=0}^{n-1} \lambda^{(n-1-i)} x[i-l] x[i-k] \right]$$

$$\begin{aligned} \text{Let } \phi_{n-1} &= \sum_{i=0}^{n-1} \lambda^{(n-1-i)} u_i u_i^T && \text{autocorrelation} \\ z_{n-1} &= \sum_{i=0}^{n-1} \lambda^{(n-1-i)} u_i d[i] && \text{cross correlation} \end{aligned} \quad \left. \vphantom{\begin{aligned} \phi_{n-1} \\ z_{n-1} \end{aligned}} \right\} \begin{array}{l} \text{same as before} \\ \text{but with } \lambda \\ \text{weightings.} \end{array}$$

$$\Rightarrow \begin{array}{ccc} \phi_{n-1} & h_{n-1} & = z_{n-1} \quad \text{at time } n-1 \\ \downarrow & \downarrow & \downarrow \\ \begin{array}{c} N \times N \\ \text{matrix} \end{array} & \begin{array}{c} N \times 1 \\ \text{vector} \end{array} & \begin{array}{c} N \times 1 \\ \text{vector} \end{array} \end{array} \quad (R h = p)$$

h_{n-1}^* is the optimal solution at time $n-1$.

$$\phi_n h_n^* = z_n \Rightarrow h_n^* = \phi_n^{-1} z_n$$

$$\text{where } \phi_n = \sum_{i=0}^n \lambda^{n-i} u_i u_i^T$$

$$z_n = \sum_{i=0}^n \lambda^{n-i} u_i d[i]$$

$$\phi_n = \underbrace{\sum_{i=0}^{n-1} \lambda^{n-i} u_i u_i^T}_{\lambda \phi_{n-1}} + u_n u_n^T = \lambda \phi_{n-1} + u_n u_n^T$$

$$z_n = \lambda z_{n-1} + u_n d[n]$$

Matrix Inversion Lemma (MIL)

A, B positive definite $M \times M$

C $M \times N$

D positive definite $N \times N$

If $A = B^{-1} + C D^{-1} C^T$

Then $A^{-1} = B - B C (D + C^T B C)^{-1} C^T B$

Let $A = \Phi_n$
 $B^{-1} = \lambda \Phi_{n-1}$
 $C = u_n$
 $D^{-1} = 1$

$\underbrace{A}_{\Phi_n} = \underbrace{B^{-1}}_{\lambda \Phi_{n-1}} + \underbrace{C}_{u_n} \underbrace{C^T}_{u_n^T}$
 $\Phi_n = \lambda \Phi_{n-1} + u_n u_n^T$

Using MIL,

$$\Phi_n^{-1} = \lambda^{-1} \Phi_{n-1}^{-1} - \lambda^{-1} \Phi_{n-1}^{-1} u_n (1 + u_n^T \lambda^{-1} \Phi_{n-1}^{-1} u_n)^{-1} u_n^T \lambda^{-1} \Phi_{n-1}^{-1}$$

$$P_n = \Phi_n^{-1}$$

$$P_n = \frac{1}{\lambda} P_{n-1} - \frac{1}{\lambda^2} \left(\frac{P_{n-1} u_n u_n^T P_{n-1}}{1 + \frac{1}{\lambda} u_n^T P_{n-1} u_n} \right)$$

$$= \frac{1}{\lambda} P_{n-1} - \frac{1}{\lambda} k_n u_n^T P_{n-1}$$

where $k_n = \frac{1}{\lambda} \frac{P_{n-1} u_n}{1 + \frac{1}{\lambda} u_n^T P_{n-1} u_n}$ is the gain vector.

$$k_n = \left(\frac{1}{\lambda} P_{n-1} - \frac{1}{\lambda} k_n u_n^T P_{n-1} \right) u_n = P_n u_n \quad (\text{try to show this})$$

$$\Rightarrow h_n^* = \phi_n^{-1} z_n = P_n z_n$$

$$= P_n (\lambda z_{n-1} + u_n d[n])$$

$$= \lambda P_n z_{n-1} + P_n u_n d[n]$$

$$= \lambda \left(\frac{1}{\lambda} P_{n-1} - \frac{1}{\lambda} k_n u_n^T P_{n-1} \right) z_{n-1} + P_n u_n d[n]$$

$$= P_{n-1} z_{n-1} - k_n u_n^T P_{n-1} z_{n-1} + P_n u_n d[n]$$

$$= h_{n-1}^* - k_n u_n^T h_{n-1}^* + k_n d[n]$$

$$= h_{n-1}^* - k_n (u_n^T h_{n-1}^* - d[n])$$

$$= h_{n-1}^* + k_n \xi_n \quad \text{where} \quad \xi_n = d[n] - u_n^T h_{n-1}^*$$

Recursive Least Squares (RLS)

$$1. \text{ Gain vector} \quad k_n = \frac{\frac{1}{\lambda} P_{n-1} u_n}{1 + \frac{1}{\lambda} u_n^T P_{n-1} u_n}$$

2. Error from previous estimate

$$\xi_n = d[n] - h_{n-1}^{*T} u_n = d[n] - u_n^T h_{n-1}^*$$

3. Update filter

$$h_n^* = h_{n-1}^* + k_n \xi_n$$

$$4. \text{ Update } \phi_n^{-1} = P_n \quad \text{where} \quad P_n = \frac{1}{\lambda} P_{n-1} - \frac{1}{\lambda} k_n u_n^T P_{n-1}$$

In practice, RLS converges much faster than LMS.

It does not depend on eigenvalues of Φ_n .