

Finding Interpolant as Broad Race-free Condition

- $raceCons(\text{race constraint}) =$
 $parallelCond(\text{parallel condition}) \wedge$
 $pathCond(\text{path condition}) \wedge raceCond(\text{race condition})$
- $raceFreeCond$ (race free variable condition):
 - e.g. $size=5 \wedge chunk=2 \wedge tM=2$
 - Initially $raceFreeCond = false$
- **Race-free** property: $raceFreeCond \wedge raceCons \rightarrow UNSAT$
- $itpl$: Craig interpolant
 - $raceFreeCond \wedge raceCons \rightarrow UNSAT$ s.t.
 - $raceFreeCond \rightarrow itpl$
 - $raceCons \wedge itpl \rightarrow UNSAT$
 - $var(itpl) = var(raceFreeCond) \cap var(raceCons)$



The Algorithm, in Readable Text

1. *Initializing: $\text{raceFreeCond} := \text{False}$ (**RaceFree** := {False})*
2. *Negative learning with case assignment: $\text{raceFreeConsUnexpl}(\text{raceFreeCond}) := \text{raceFreeCons} \setminus \sim \text{raceFreeCond}$*
 - a. *if SAT \Rightarrow goto step 3 with $\text{vars}_{RF} := \text{SAT model}$*
 - b. *if UNSAT \Rightarrow return raceFreeCond as maximal race free condition*
 - c. *if unknown/timed-out*
 - i. *doing regression: goto step 2 with $\text{raceFreeConsUnexpl}(\text{raceFreeCond}_{\text{regr}})$, where $\text{raceFreeCond}_{\text{regr}}$ temporarily combinatorially hides some known race-free cases, including False, from raceFreeCond*
 - ii. *doing more existential quantification: goto step 2 with $\mathbf{var}(\text{raceFreeCond}) := \mathbf{var}(\text{raceFreeCond}) \cup \{v \mid v \in \mathbf{var}(\text{raceCons}) \setminus v \notin \mathbf{var}(\text{raceFreeCond})\}$*
 - iii. *return raceFreeCond as best-known race free condition if existential quantification upper bound is reached*
3. *Case expansion: $\text{raceFreeCond} := \text{raceFreeCond} \setminus / \text{vars}_{RF}$*

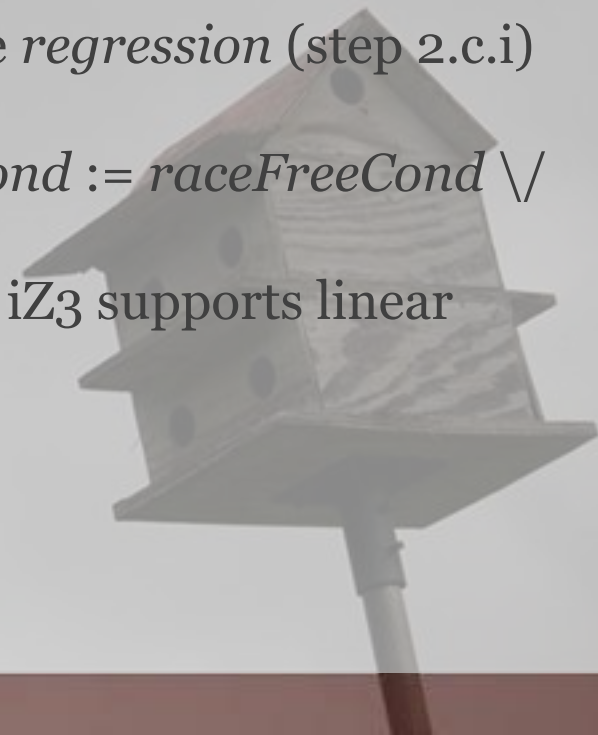


4. *Interpolation*: $itpl := interpolant(raceFreeCond, raceCons)$
 - a. if SAT/unknown/timed-out/ $itpl = raceFreeCond \Rightarrow$
 $itpl_{turbo_heu} := interpolant(raceFreeCond_{turbo_heu}, raceCons)$, where
 $raceFreeCond_{turbo_heu} :=$
 - i. *one without totally conflicting inv*: if $raceFreeCond_{turbo_heu}$ doesn't totally conflict inv , goto step 4.c with $raceFreeCond := raceFreeCond_{turbo_heu} \setminus itpl_{turbo_heu}$ or $raceFreeCond_{turbo_heu}$ (if $itpl_{turbo_heu} = raceFreeCond_{turbo_heu}$)
 - ii. *subsumption learning*: $raceFreeCond \setminus subsume_heu(raceFreeCond)$, where $subsume_heu(raceFreeCond) = /\ name(\mathbf{vars}_{RFi}[j]) </= /> value(\mathbf{vars}_{RFi}[j])$ (subsumption):
While there's any single, double, triple, quad, ..., n -tuple variable combination \mathbf{vars}_{RFi} ,
 1. bypassing (and caching) combinatorially redundant computation
 2. *analogical variable bounded subsumption*
 3. *simple bounded subsumption*
 4. including (unhiding) the cases hidden in the *regression* (step 2.c.i) and computing again

iii. *simple bounded linear learning: $\text{raceFreeCond} \setminus \text{linear_heu}(\text{raceFreeCond})$, where $\text{linear_heu}(\text{raceFreeCond}) = (\text{boundary}) \setminus \text{linear relation}$:*

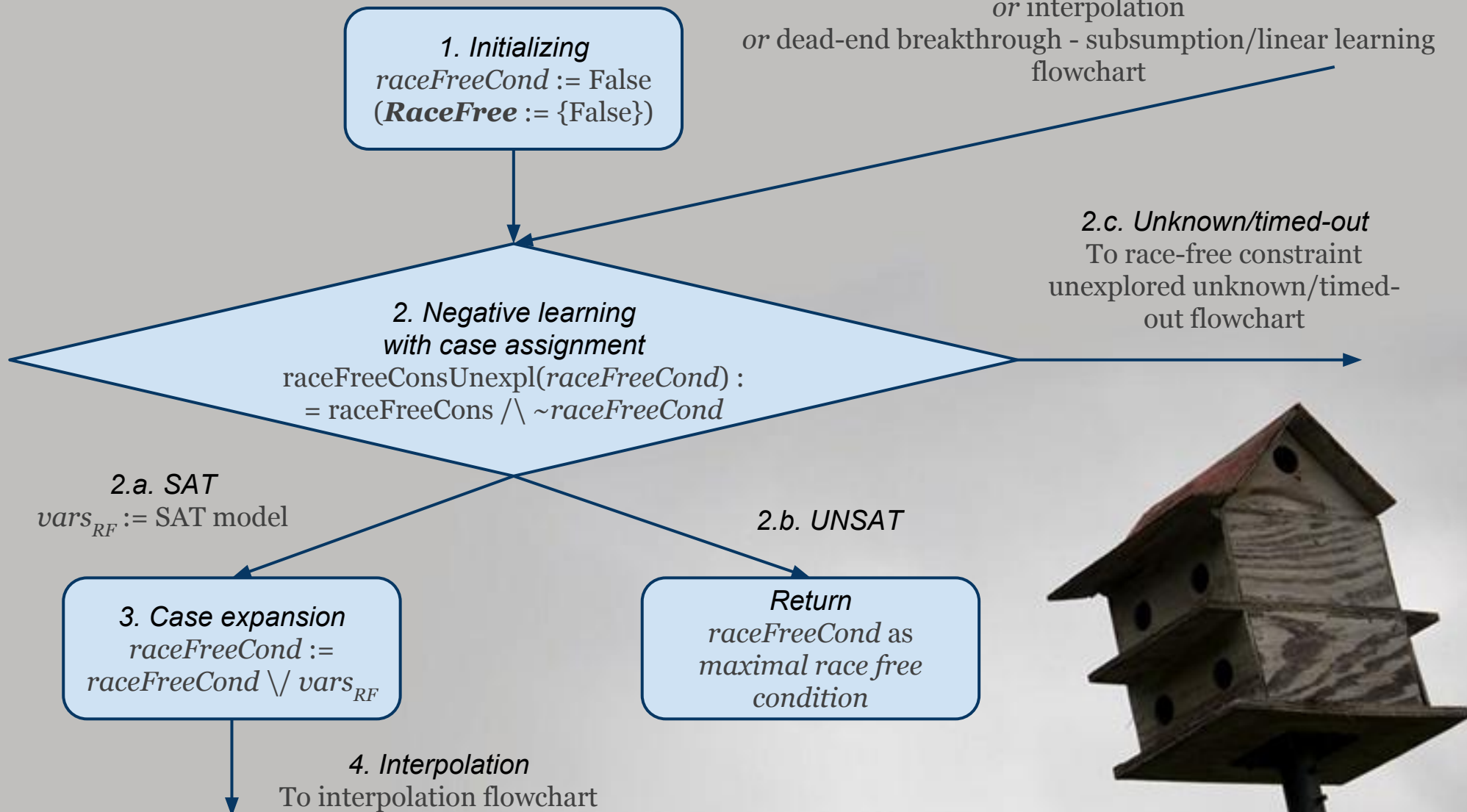
While there's any pair of single, double, triple, quad, ..., n -tuple variable vector ***vars_pair***_{*RFi*},

1. eliminating shared variables
 2. *computing linear relations ($\text{vars_pair}_{RFi}[1] = a * \text{vars_pair}_{RFi}[2] + b$) for each pair of the rest variable vectors*
 3. *combinatorially bounding all non-eliminated variables if there's no interpolant for the heuristics with just linear relation*
 4. including (unhiding) the cases hidden in the *regression* (step 2.c.i) and computing again
- b. else (if UNSAT and $\text{itpl} \neq \text{raceFreeCond}$) $\Rightarrow \text{raceFreeCond} := \text{raceFreeCond} \setminus \text{itpl}$,
- c. double-check $\text{raceFreeCond} \setminus \text{raceCons}$ by Z3 (for that iZ3 supports linear integer algorithm only)
- i. should be UNSAT \Rightarrow goto step 2,
 - ii. if SAT/unknown/timed-out \Rightarrow warning buggy iZ3

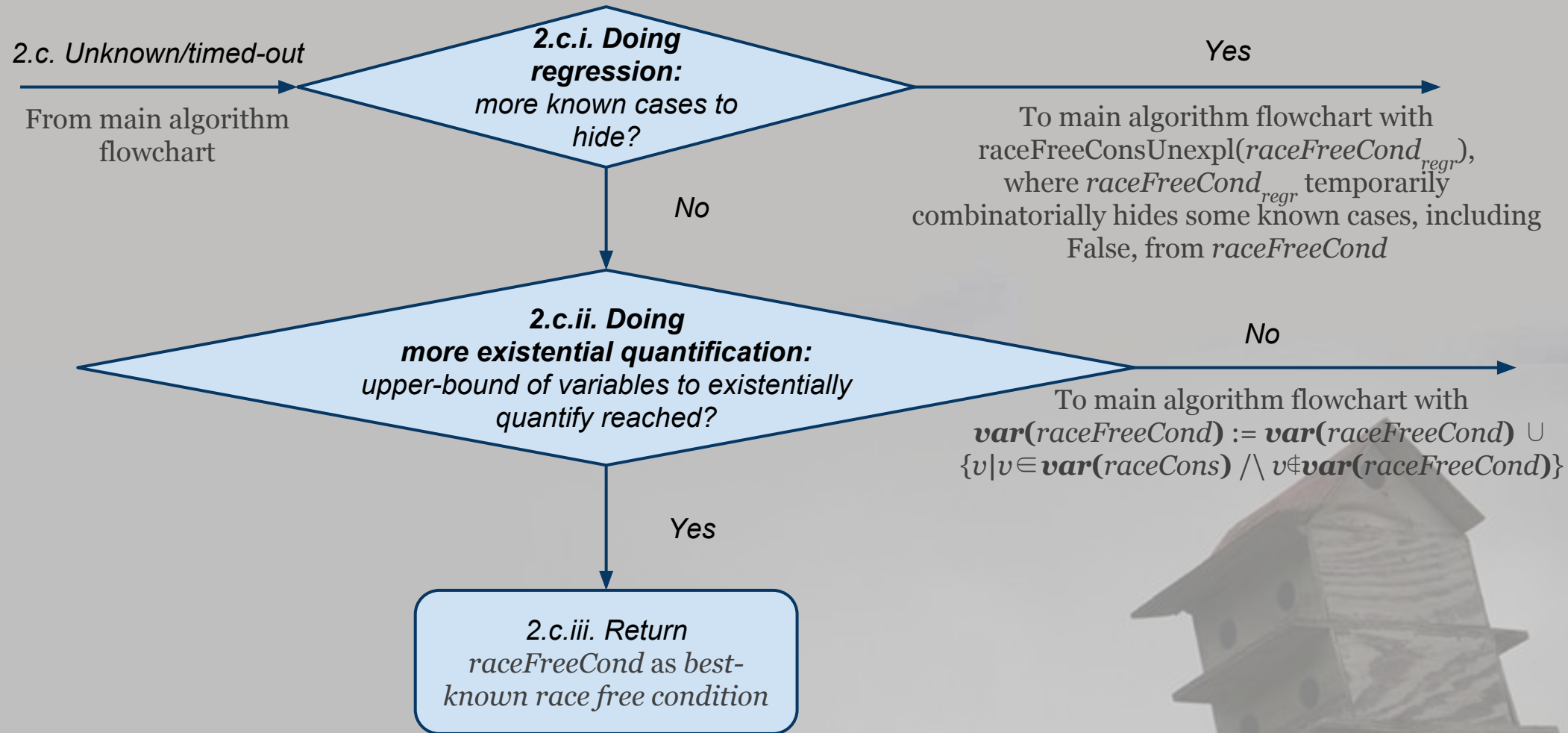


Main Algorithm Flowchart

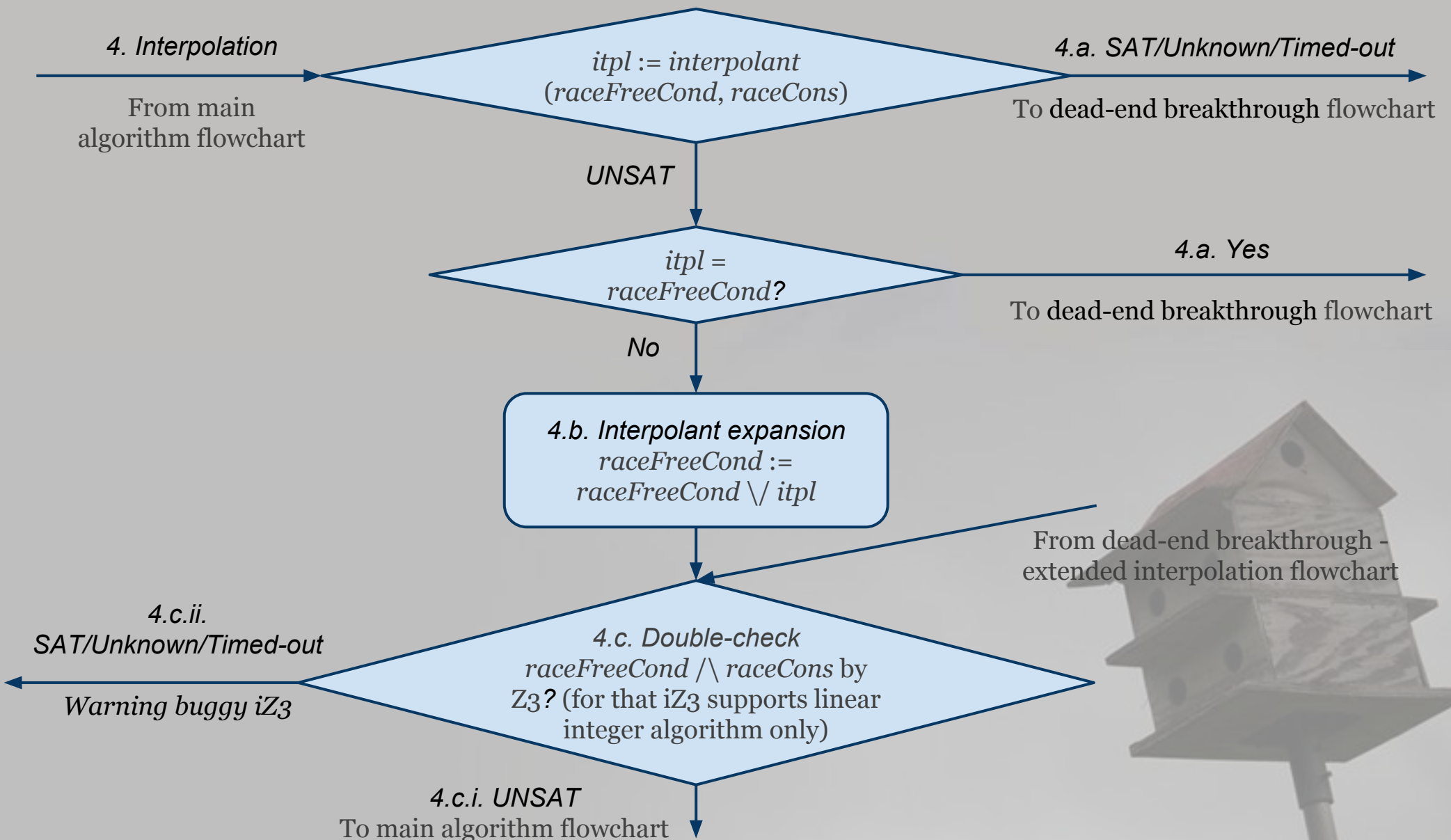
From race-free constraint unexplored unknown/timed-out
or interpolation
or dead-end breakthrough - subsumption/linear learning
flowchart



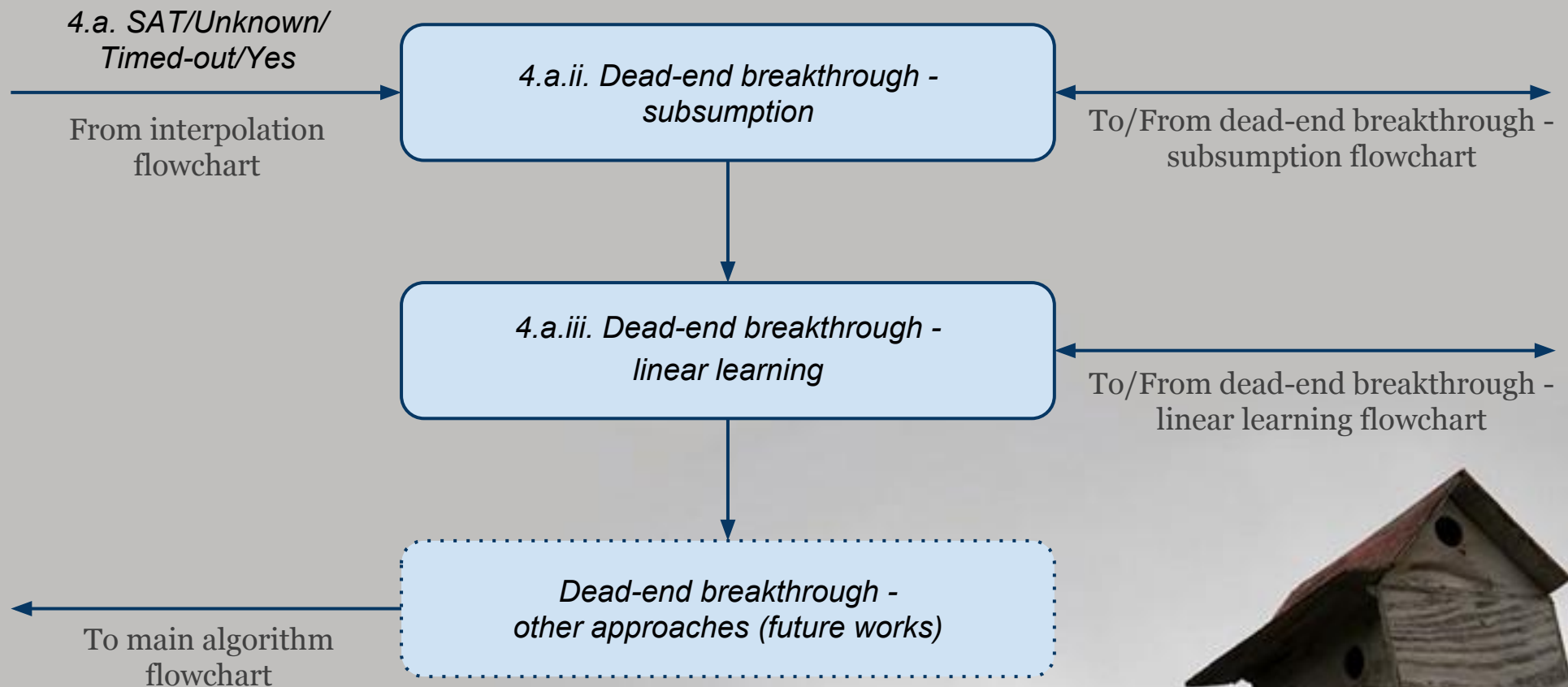
Race-free Constraint Unexplored Unknown/Timed-out Flowchart



Interpolation Flowchart



Dead-end Breakthrough Flowchart



Dead-end Breakthrough - Subsumption Flowchart

From dead-end breakthrough/
dead-end breakthrough - extended interpolation flowchart

- (1) SAT (no interpolant) \wedge Unknown \rightarrow Unknown;
 - (2) UNSAT (got interpolant) \wedge Unknown \rightarrow UNSAT \rightarrow **TDD (tertiary decision diagram) cacheable**;
 - (3) SAT \vee UNSAT \rightarrow SAT (which means no interpolants, even given a widening range of subsumption) \rightarrow **TDD cacheable**;
 - (4) SAT \vee SAT \rightarrow SAT \rightarrow **TDD cacheable**;
 - (5) UNSAT \vee UNSAT \rightarrow UNSAT \rightarrow **TDD cacheable**
- ex., $(\text{raceFreeCond} \vee (\text{size} > 6 \wedge \text{chunk} > 3 \wedge \text{tM} = 2)) \wedge \text{raceCons}$ is SAT since $(\text{raceFreeCond} \vee (\text{size} > 5 \wedge \text{chunk} > 2 \wedge \text{tM} = 2)) \wedge \text{raceCons}$ is SAT and $(\text{raceFreeCond} \vee (\text{size} = 6 \wedge \text{chunk} = 3 \wedge \text{tM} = 2)) \wedge \text{raceCons}$ is UNSAT

4.a.ii.1. Bypassing (and caching) combinatorially redundant computation:
deducing cases surely with/without interpolants

4.a.ii.2. Analogical variable bounded subsumption
ex., subsumption singleton $\text{tM} = 2$ or
triple $(\text{size} < 5) \wedge (\text{chunk} < 1) \wedge (\text{tM} = 2)$ from cases
 $(\text{size}, \text{chunk}, \text{tM}) = (5, 2, 2)$ and $(6, 3, 2)$, which sharing $\text{tM} = 2$

4.a.ii.3. Simple bounded subsumption
ex., subsumption $\text{size} < 5$
from case $(\text{size}, \text{chunk}, \text{tM}) = (5, 2, 2)$

4.a.ii.4.
More cases hidden in
Doing regression
(step 2.c.i) to include?

$\text{raceFreeCond}_{\text{Turbo_heu}} =$
 $\text{raceFreeCond} \vee \text{subsume_heu}$
 $(\text{raceFreeCond}),$
where
 $\text{subsume_heu}(\text{raceFreeCond}) =$
 $\wedge \text{name}(\text{vars}_{\text{RFi}}[j]) < / = > \text{value}$
 $(\text{vars}_{\text{RFi}}[j])$ (subsumption)

4.a.i.

To dead-end breakthrough -
extended interpolation flowchart

Yes
(single, dual,
triple, quad,
..., n-tuple)

More
variable
combination
 vars_{RFi} to
increase?

Yes

To main algorithm
flowchart

No

Dead-end Breakthrough - Linear Learning Flowchart

From dead-end breakthrough flowchart

Simple bounded linear learning:

4.a.iii.1. eliminating shared variables

-> 4.a.iii.2. computing linear relations

$$(\text{vars_pair}_{RFi}[1] = a * \text{vars_pair}_{RFi}[2] + b)$$

for each pair of the rest variable vectors

ex., size = chunk + 3

from cases (size, chunk, tM) = (5,2,2) and (6,3,2)

4.a.iii.3. Combinatorially bounding all non-eliminated variables if returned without any interpolant

$$\text{raceFreeCond}_{\text{turbo_heu}} = \text{raceFreeCond} \setminus \text{linear_heu}(\text{raceFreeCond}), \text{ where}$$

$$\text{linear_heu}(\text{raceFreeCond}) = (\text{boundary}) \setminus \text{linear relation}$$

4.a.i.

To/From dead-end breakthrough - extended interpolation flowchart

Yes

Any more pair of variable vector vars_pair_{RFi} ?

No

Yes

(single, dual, triple, quad, ..., n-tuple)

No

4.a.iii.4.

More cases hidden in **Doing regression** (step 2.c.i) to include ?

To main algorithm flowchart



Dead-end Breakthrough - Extended Interpolation Flowchart

