# **CS570**

# Analysis of Algorithms Spring 2015 Exam I

Name:	
Student ID:	
Email Address:	_
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# \_\_\_\_Check if DEN Student

	Maximum	Received
Problem 1	20	
Problem 2	20	
Problem 3	14	
Problem 4	13	
Problem 5	20	
Problem 6	13	
Total	100	

#### **Instructions:**

- 1. This is a 2-hr exam. Closed book and notes
- 2. If a description to an algorithm or a proof is required please limit your description or proof to within 150 words, preferably not exceeding the space allotted for that question.
- 3. No space other than the pages in the exam booklet will be scanned for grading.
- 4. If you require an additional page for a question, you can use the extra page provided within this booklet. However please indicate clearly that you are continuing the solution on the additional page.

#### 20 pts

Mark the following statements as **TRUE** or **FALSE**. No need to provide any justification.

#### [TRUE/FALSE]

For some graphs BFS and DFS trees can be the same.

#### [TRUE/FALSE]

The number of cycles in a bipartite graph may be odd.

#### [TRUE/FALSE]

Stable matching algorithm presented in class is based on the greedy technique.

#### [TRUE/FALSE]

To delete the i<sup>th</sup> node in a binary min heap, you can exchange the last node with the i<sup>th</sup> node, then check the nodes below the i<sup>th</sup> node to see if the i<sup>th</sup> node should move down the heap to "re-heapify" it.

#### [TRUE/FALSE]

Kruskal's algorithm can fail in the presence of negative cost edges.

#### [TRUE/FALSE]

Consider an instance of the Stable Matching Problem in which there exists a man m and a woman w such that m is ranked first by w, and w is ranked first by m. Then (m, w) must appear in every stable matching.

#### [ TRUE/FALSE ]

If a connected undirected graph G has the same weights for every edge, then a minimum spanning tree can be found in linear time.

#### [TRUE/FALSE]

Given n numbers, one could construct a binary heap using the n numbers, then using the binary heap produce a sorted list of the numbers in O(n) time.

#### [TRUE/FALSE]

In a Fibonacci heap, the insert operation has an amortized cost of O(1) time, but the worst case cost is higher.

#### [TRUE/FALSE]

Function  $10n^{10}2^n + 3^n\log(n)$  is  $O(n^{10}2^n)$ .

## 2) 20 pts

A pharmacist has W pills and n empty bottles. Let  $\{p_1, p_2, \dots, p_n\}$  denote the number of pills that each bottle can hold.

Describe a greedy algorithm, which, given W and  $\{p_1, p_2, ..., p_n\}$ , determines the fewest number of bottles needed to store the pills. Prove that your algorithm is correct.

## 3) 14 pts

We are given a graph G=(V;E) with uniform cost edges and two vertices s and t within G. We are told that the length of the shortest path from s to t is more than n/2 (where n is the number of vertices within G). Give a linear time algorithm to find a vertex v (other than s and t) such that every path from s to t contains v. Justify your solution.

## 4) 13 pts

Arrange the following functions in increasing order of asymptotic complexity. If  $f(n)=\Theta(g(n))$  then put f=g. Else, if f(n)=O(g(n)), put f < g.

$$4n^2$$
,  $\log^2 n$ ,  $5n^2$ ,  $\log^3 n$ ,  $n^n$ ,  $3n$ ,  $nlog(n)$ ,  $2n$ ,  $2n+1$ 

5)	20 pts
	Let $G = (V, E)$ be an (undirected) graph with costs $C_e \ge 0$ on the edges $e \in E$ . Assume
	you are given a minimum-cost spanning tree $T$ in $G$ . Now assume that an edge
	connecting two nodes $v, w \in V$ with cost $c$ is deleted from graph G and let G' be the

new graph. Assume that the graph G' is connected.

a. If the removed edge doesn't appear in the MST T, will T still be the MST of G'? Please justify your answer.

b. If the removed edge appears in the MST T, give an O(|V|+|E|) algorithm to find a MST for the graph G'.

c. Prove that the output produced by your algorithm in part b is an MST

6) 13 pts

Solve the following recurrences by giving tight  $\Theta$ -notation bounds. You do not need to justify your answers, but any justification that you provide will help when assigning partial credit.

i. 
$$T(n)=5T(n/2)+n^2\log n$$

ii. 
$$T(n)=4T(n/2)+n\log n$$

iii. 
$$T(n) = (6006)^{1/2} *T(n/2) + n^{\sqrt{6006}}$$

Additional Space