

VMSP: Efficient Vertical Mining of Maximal Sequential Patterns

Philippe Fournier-Viger¹

Cheng-Wei Wu²

Antonio Gomariz³

Vincent Shin-Mu Tseng²

¹University of Moncton, Canada

²National Cheng Kung University, Taiwan

³University of Murcia



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Introduction

Sequential pattern mining:

- a data mining task with wide applications
- finding frequent subsequences in a **sequence database**.

Example:

minsup = 2

Sequence database

SID	Sequences
1	$\langle \{a, b\}, \{c\}, \{f, g\}, \{g\}, \{e\} \rangle$
2	$\langle \{a, d\}, \{c\}, \{b\}, \{a, b, e, f\} \rangle$
3	$\langle \{a\}, \{b\}, \{f\}, \{e\} \rangle$
4	$\langle \{b\}, \{f, g\} \rangle$



Some sequential patterns

ID	Pattern	Supp.
p1	$\langle \{a\}, \{f\} \rangle$	3
p2	$\langle \{a\}, \{c\} \{f\} \rangle$	2
p3	$\langle \{b\}, \{f, g\} \rangle$	2
p4	$\langle \{g\}, \{e\} \rangle$	2
p5	$\langle \{c\}, \{f\} \rangle$	2
p6...	$\langle \{b\} \rangle$	4

Algorithms

Different approaches to solve this problem

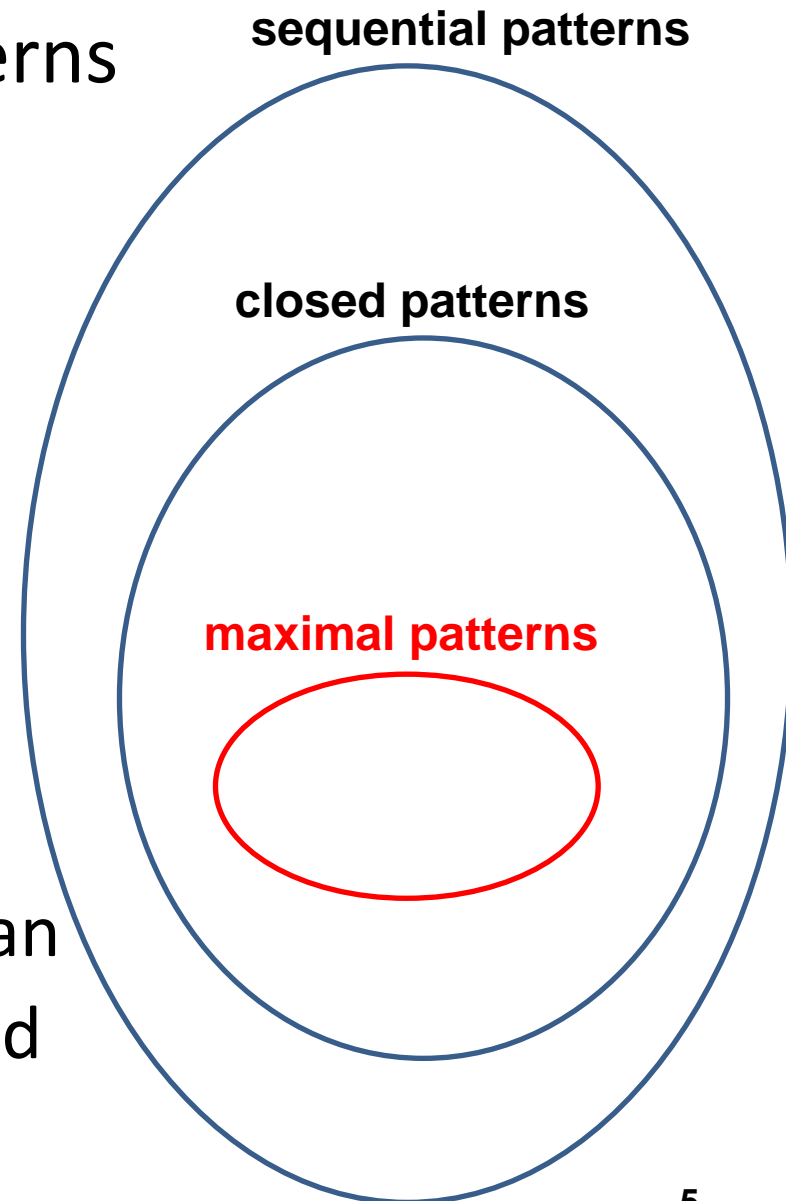
- Apriori-based
(e.g. [GSP](#))
- Pattern-growth
(e.g. [PrefixSpan](#))
- Discovery of sequential patterns using a vertical database representation
(e.g. [SPADE](#), [SPAM](#), [bitSPADE](#)...)

The problem of redundancy

- **Observation:** if $\{a\}, \{c\}, \{f\}$ is frequent, then the pattern $\{c\}, \{f\}$, the pattern $\{a\}$, the pattern $\{c\}$... are frequent.
- Consider a frequent pattern of 20 distinct items.
- Its $2^{20}-1$ subsequences are also frequent!
- Because of redundancy,
 - very time-consuming to analyze patterns,
 - require much more storage space.

A solution

- **Closed sequential patterns:** patterns that are not included in another pattern *having the same support*.
 - lossless
 - this set is still quite large for some applications
- **Maximal sequential patterns:** patterns that are not included in another pattern.
 - lossless with an extra database scan
 - generally much smaller than closed patterns



Multiple applications

- discovering frequent longest common subsequences in texts,
- analyzing DNA sequences,
- data compression,
- web log mining.

Example

A sequence database

SID	Sequences
1	$\langle \{a, b\}, \{c\}, \{f, g\}, \{g\}, \{e\} \rangle$
2	$\langle \{a, d\}, \{c\}, \{b\}, \{a, b, e, f\} \rangle$
3	$\langle \{a\}, \{b\}, \{f\}, \{e\} \rangle$
4	$\langle \{b\}, \{f, g\} \rangle$

Patterns found for *minsup* = 2

Pattern	Sup.	Pattern	Sup.
$\langle \{a\} \rangle$	3 C	$\langle \{b\}, \{g\}, \{e\} \rangle$	2 CM
$\langle \{a\}, \{g\} \rangle$	2	$\langle \{b\}, \{f\} \rangle$	4 C
$\langle \{a\}, \{g\}, \{e\} \rangle$	2 CM	$\langle \{b\}, \{f, g\} \rangle$	2 CM
$\langle \{a\}, \{f\} \rangle$	3 C	$\langle \{b\}, \{f\}, \{e\} \rangle$	2 CM
$\langle \{a\}, \{f\}, \{e\} \rangle$	2 CM	$\langle \{b\}, \{e\} \rangle$	3 C
$\langle \{a\}, \{c\} \rangle$	2	$\langle \{c\} \rangle$	2
$\langle \{a\}, \{c\}, \{f\} \rangle$	2 CM	$\langle \{c\}, \{f\} \rangle$	2
$\langle \{a\}, \{c\}, \{e\} \rangle$	2 CM	$\langle \{c\}, \{e\} \rangle$	2
$\langle \{a\}, \{b\} \rangle$	2	$\langle \{e\} \rangle$	3
$\langle \{a\}, \{b\}, \{f\} \rangle$	2 CM	$\langle \{f\} \rangle$	4
$\langle \{a\}, \{b\}, \{e\} \rangle$	2 CM	$\langle \{f, g\} \rangle$	2
$\langle \{a\}, \{e\} \rangle$	3 C	$\langle \{f\}, \{e\} \rangle$	2
$\langle \{a, b\} \rangle$	2 CM	$\langle \{g\} \rangle$	3
$\langle \{b\} \rangle$	4	$\langle \{g\}, \{e\} \rangle$	2
$\langle \{b\}, \{g\} \rangle$	3 C		

C = Closed M = Maximal

Algorithms

- **MSPX**: approximate solution
- **DISMAPS**: for strings with no repeating items
- for the general problem:
 - AprioriAdjust, MSPX, MFSPAN**
 - AprioriAdjust is based on Apriori,
 - they all need to maintain a large set of intermediate candidates in memory during the mining process
- **MaxSP**:
 - most recent algorithm
 - does not maintain intermediate candidates in memory
 - only explore patterns occurring in the DB

Our proposal

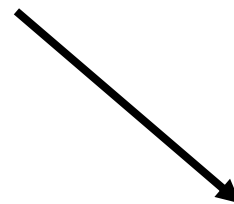
VMSP:

- discovers maximal sequential patterns,
- based on the **SPAM** search procedure
- integrates three novel strategies:
 - **EFN**: Efficient Filtering of Non-Maximal Patterns
 - **FME**: Forward Maximal Extension Checking
 - **CPC**: Candidate Pruning by Co-Occurrence Map

The SPAM search procedure

Step 1: creates a vertical representation of the database (SID lists):

SID	Sequences
1	$\langle \{a, b\}, \{c\}, \{f, g\}, \{g\}, \{e\} \rangle$
2	$\langle \{a, d\}, \{c\}, \{b\}, \{a, b, e, f\} \rangle$
3	$\langle \{a\}, \{b\}, \{f\}, \{e\} \rangle$
4	$\langle \{b\}, \{f, g\} \rangle$



a	
SID	Itemsets
1	1
2	1,4
3	1
4	

b	
SID	Itemsets
1	1
2	3,4
3	2
4	1

c	
SID	Itemsets
1	2
2	2
3	
4	

d	
SID	Itemsets
1	
2	1
3	
4	

e	
SID	Itemsets
1	5
2	4
3	4
4	

f	
SID	Itemsets
1	3
2	4
3	3
4	2

g	
SID	Itemsets
1	3,4
2	
3	
4	2

The SPAM search procedure (2)

Step 2:

- identify frequent patterns containing a single item.
- recursively append items to each frequent pattern to generate larger patterns.
 - s-extension: $\langle l_1, l_2, l_3 \dots l_n \rangle$ with $\{a\}$ is $\langle l_1, l_2, l_3 \dots l_n, \{a\} \rangle$
 - i-extension: $\langle l_1, l_2, l_3 \dots l_n \rangle$ with $\{a\}$ is $\langle l_1, l_2, l_3 \dots l_n \cup \{a\} \rangle$
- The support of a larger pattern is calculated by intersecting SID lists:

a	
SID	Itemsets
1	1
2	1,4
3	1

support = 3

b	
SID	Itemsets
1	1
2	3,4
3	2
4	1

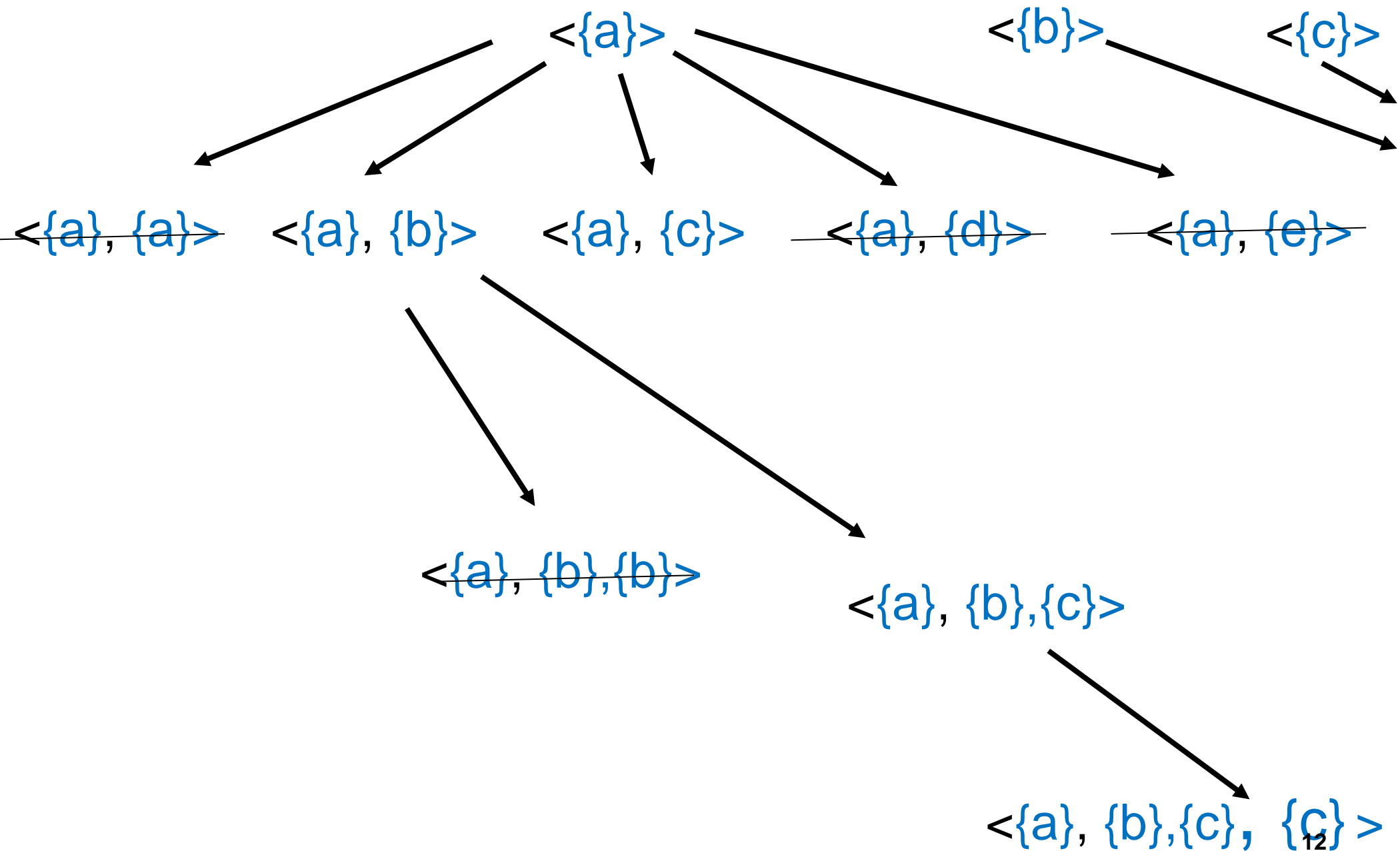
support = 4



$\langle \{a\}, \{b\} \rangle$	
SID	Itemsets
1	1
2	1, 3
3	2

support = 3

The SPAM search procedure (3)

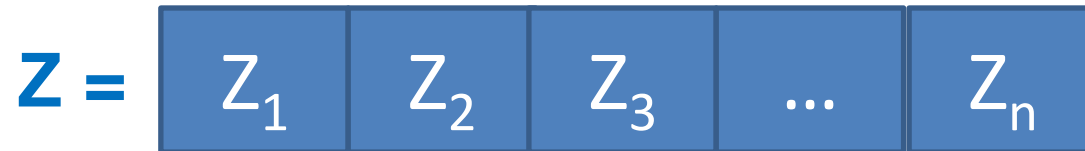


EFN: Efficient Filtering of Non-Maximal Patterns

- A structure **Z**
 - for storing maximal patterns
 - is initialized as empty.
- For each pattern $S = \{a_1, a_2, \dots a_n\}$ found
 - **super-pattern checking:** if **S** is a subsequence of a pattern **X** in **Z**, then **S** is not maximal and is not added to **Z**.
Otherwise, it is added to **Z**.
 - **sub-pattern checking:** if **S** is a super-sequence of a pattern **X** in **Z**, then **X** is not maximal and is removed from **Z**.

EFN: Efficient Filtering of Non-Maximal Patterns (cont'd)

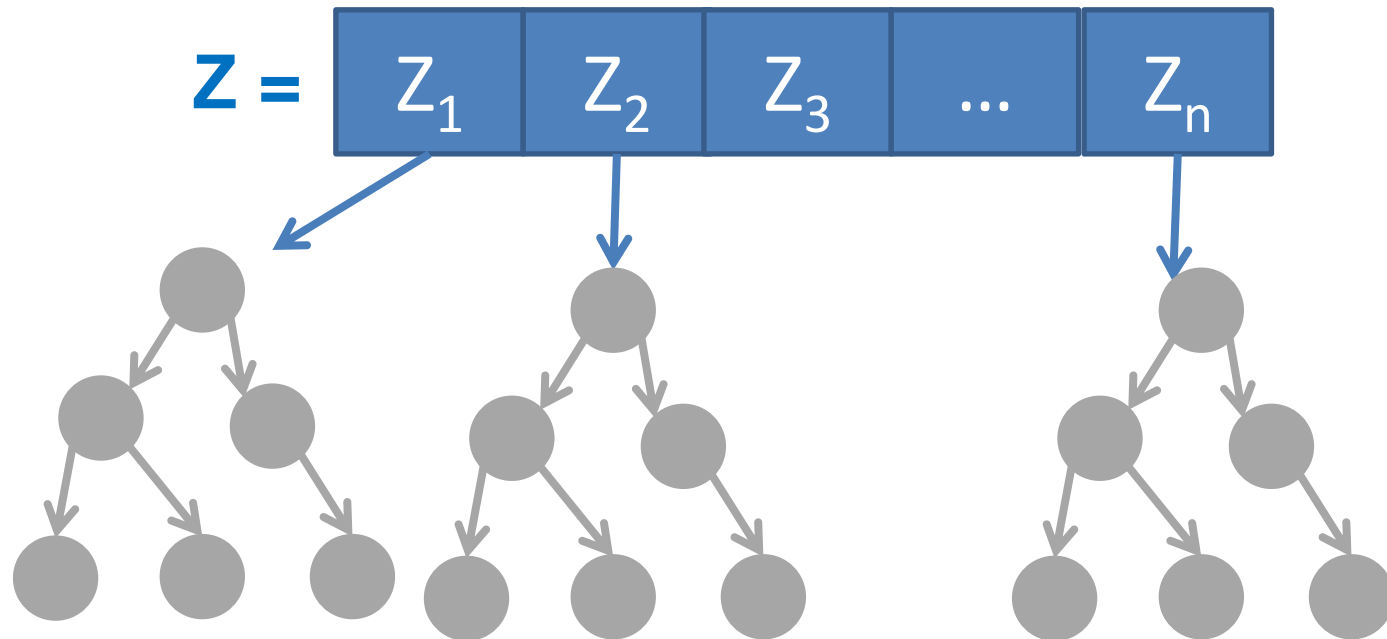
We implement **Z** as a List of heaps



The **k-th** list entry contains patterns of size **k**

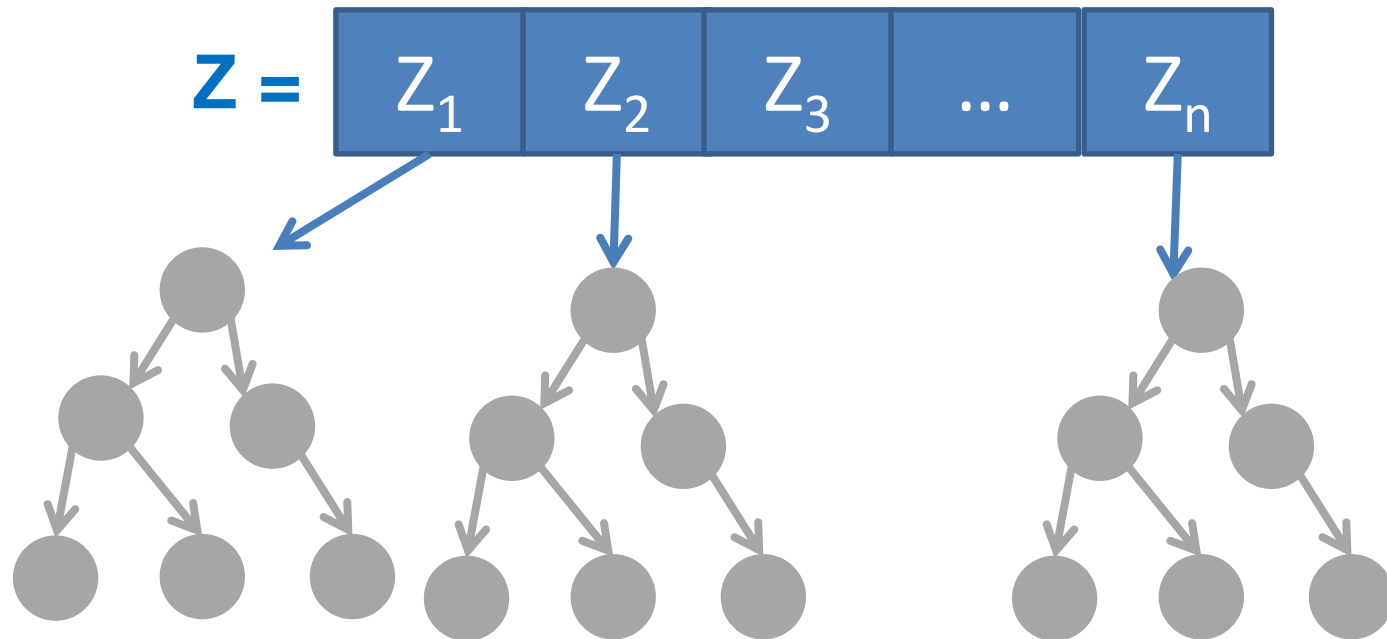
This allows to perform super-pattern checking and sub-pattern checking only with smaller and larger patterns

EFN: Efficient Filtering of Non-Maximal Patterns (cont'd)



- The **sum of items** in each pattern is calculated.
- Each **heap** orders patterns by decreasing sum of items.
- For each pattern S_a found and pattern S_b in Z_k , if $\text{sum}(S_a) < \text{sum}(S_b)$ we don't need to perform super-pattern checking with S_b and any following patterns in Z_k .
- Similar for sub-pattern-checking

EFN: Efficient Filtering of Non-Maximal Patterns (cont'd)



- **Support check optimization:**
 - A pattern cannot be contained in another pattern if its support is smaller.
 - A pattern cannot contain another pattern if its support is larger.

FME: Forward Maximal Extension Checking

- The algorithm performs a depth-first search (it grows patterns by appending items to smaller patterns one item at a time).
- We can avoid super-pattern checking for a pattern **S** if the recursive call to the search procedure with **S** produces a frequent pattern.

CPC: Candidate Pruning by Co-occurrence Map

- A structure **CMAPI** stores every items that succeeds each item by **i-extension** at least *minsup* times.
- A similar structure **CMAS** stores every items that succeeds each item by **s-extension** at least *minsup* times.

<i>CMAPI</i>		<i>CMAS</i>	
item	is succeeded by (i-extension)	item	is succeeded by (s-extension)
<i>a</i>	{ <i>b</i> }	<i>a</i>	{ <i>b, c, e, f</i> }
<i>b</i>	\emptyset	<i>b</i>	{ <i>e, f, g</i> }
<i>c</i>	\emptyset	<i>c</i>	{ <i>e, f</i> }
<i>e</i>	\emptyset	<i>e</i>	\emptyset
<i>f</i>	{ <i>g</i> }	<i>f</i>	{ <i>e, g</i> }
<i>g</i>	\emptyset	<i>g</i>	\emptyset

CPC: Candidate Pruning by Co-occurrence Map

- **Pruning:** for a pattern **S**, an i-extension (s-extension) with an item **x** will result in an infrequent patterns if there exists a pair of items in the resulting pattern that is not in $CMAP_i$ ($CMAP_s$).
- This avoid performinig costly SID lists intersections.

$CMAP_i$		$CMAP_s$	
item	is succeeded by (i-extension)	item	is succeeded by (s-extension)
<i>a</i>	$\{b\}$	<i>a</i>	$\{b, c, e, f\}$
<i>b</i>	\emptyset	<i>b</i>	$\{e, f, g\}$
<i>c</i>	\emptyset	<i>c</i>	$\{e, f\}$
<i>e</i>	\emptyset	<i>e</i>	\emptyset
<i>f</i>	$\{g\}$	<i>f</i>	$\{e, g\}$
<i>g</i>	\emptyset	<i>g</i>	\emptyset

Other optimizations

- SID lists are implemented as bitsets as in the **SPAM** and **BitSpade** algorithms.

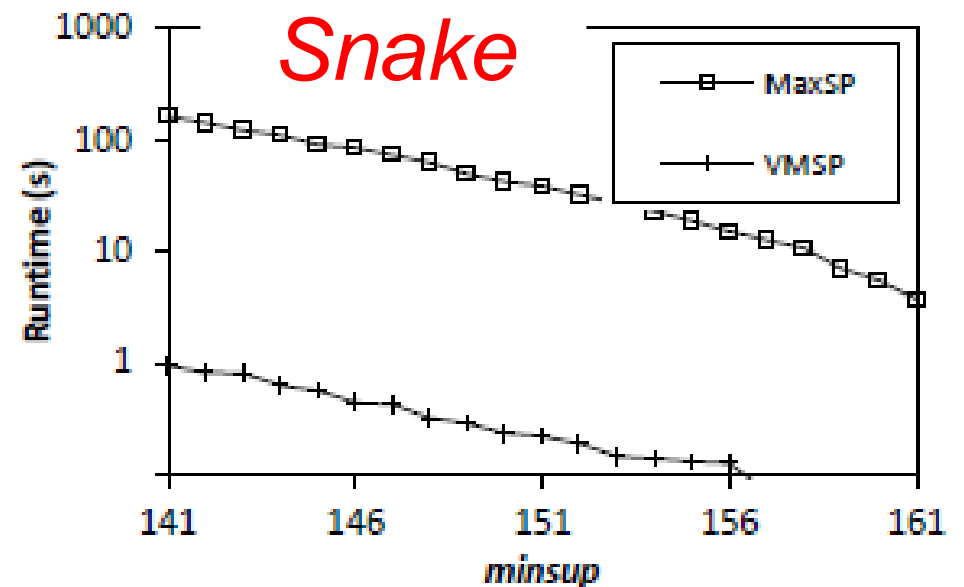
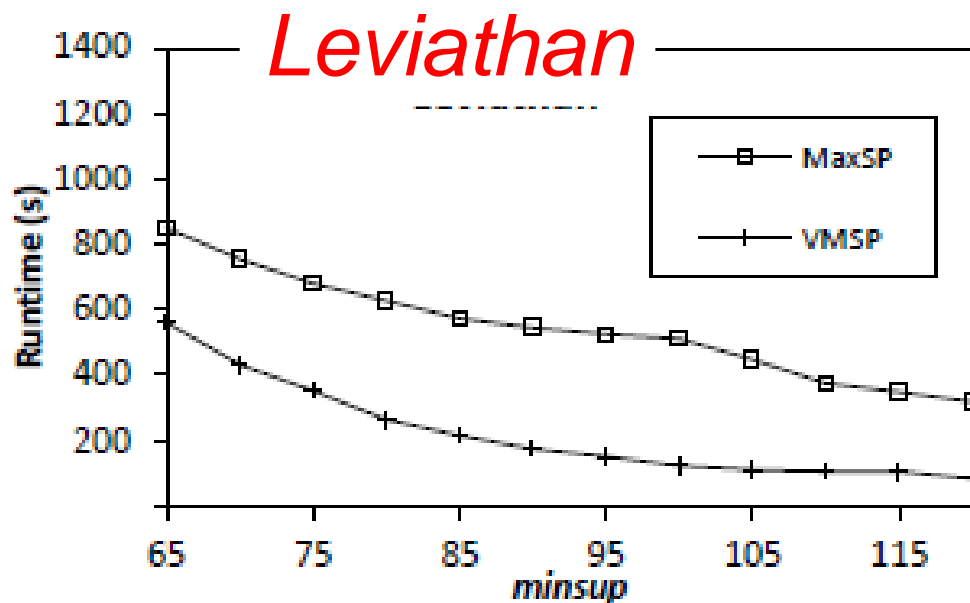
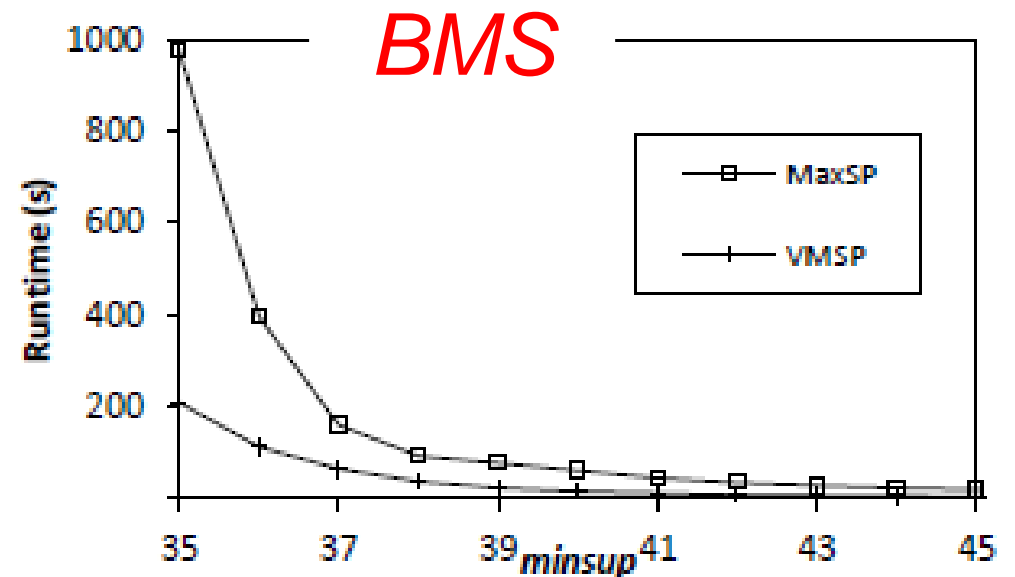
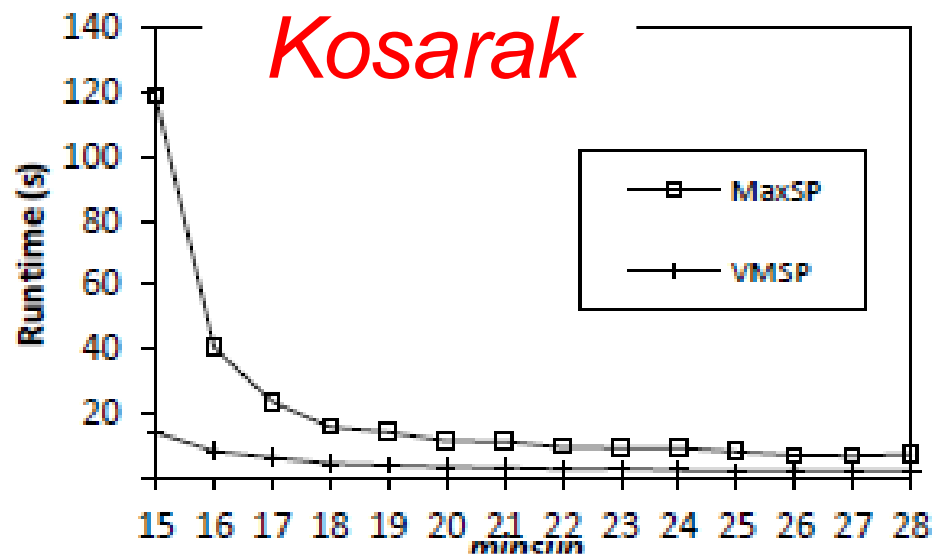
Experimental Evaluation

Datasets' characteristics

dataset	sequence count	item count	avg. seq. length (items)	type of data
Leviathan	5834	9025	33.81 (std= 18.6)	book
Snake	163	20	60 (std = 0.59)	protein sequences
FIFA	20450	2990	34.74 (std = 24.08)	web click stream
BMS	59601	497	2.51 (std = 4.85)	web click stream
Kosarak10k	10000	10094	8.14 (std = 22)	web click stream

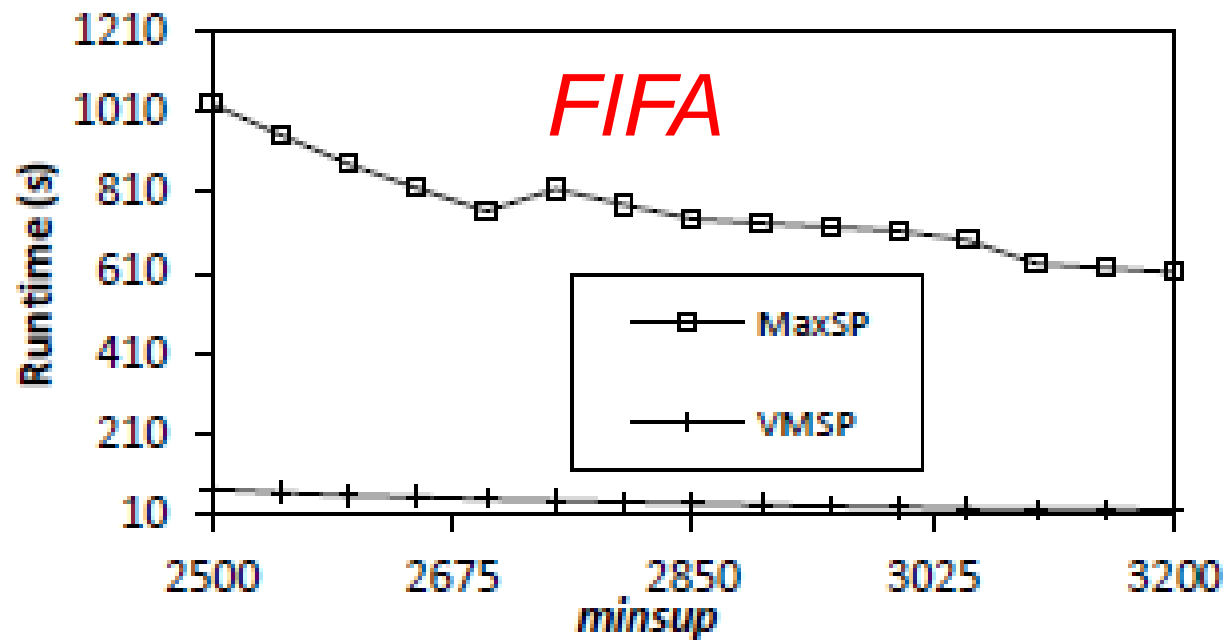
- VMSP vs MaxSP
- All algorithms implemented in Java
- Windows 7, 5 GB of RAM

Execution time



VMSP is up to 100 times faster than MaxSP

Execution time (cont'd)



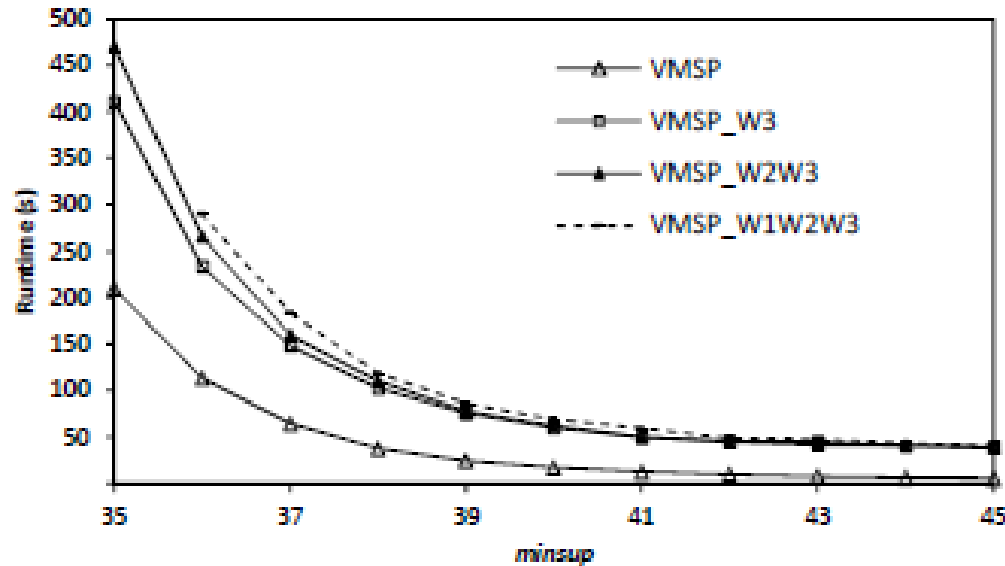
Maximum Memory Usage (MB)

Dataset	VMSP	MaxSP
BMS	840	403
Snake	45	380
Kosarak	1600	393
Leviathan	911	1150
FIFA	611	970

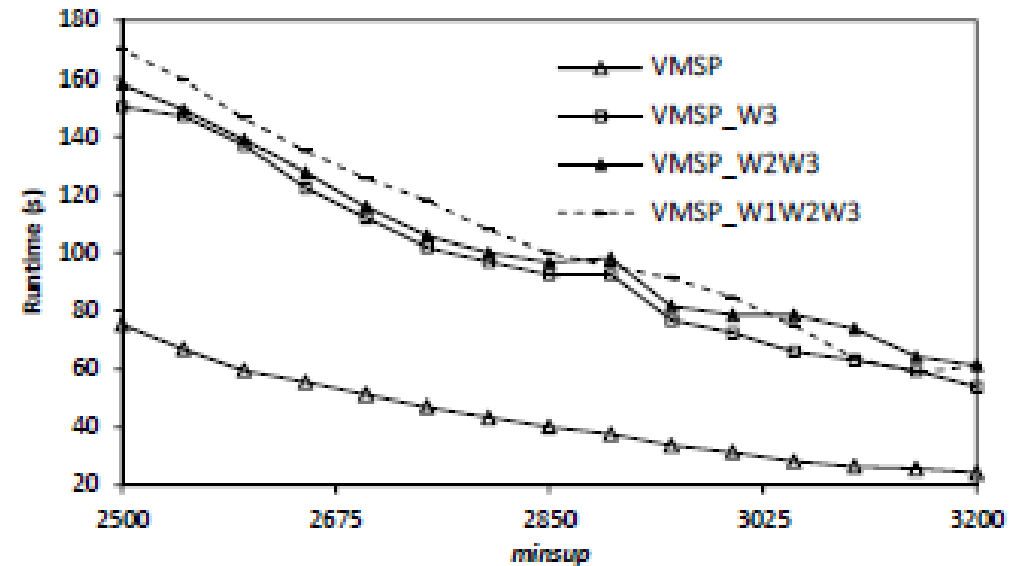
VMSP has the lowest memory consumption for 3 out of 5 datasets

Influence of the strategies

BMS



FIFA



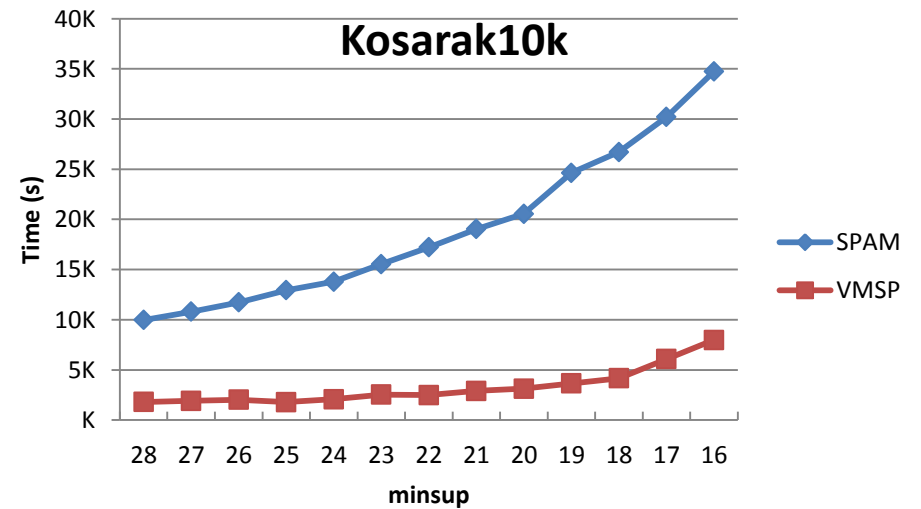
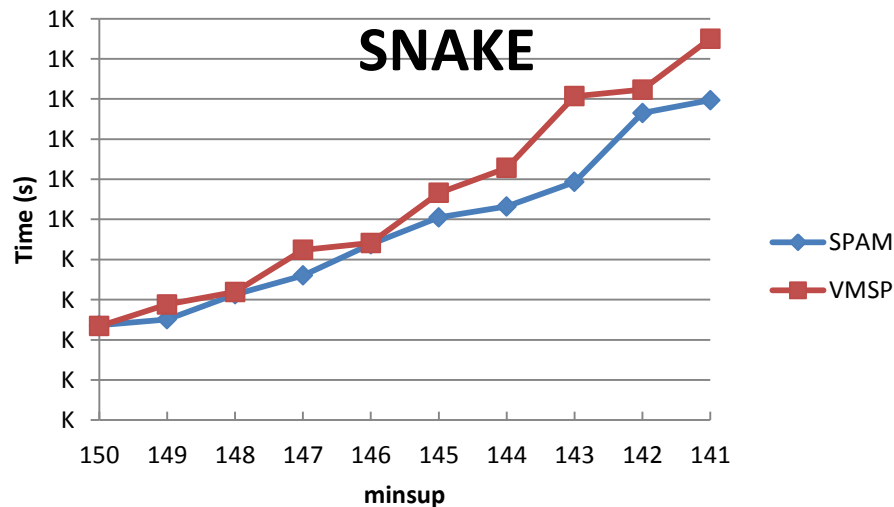
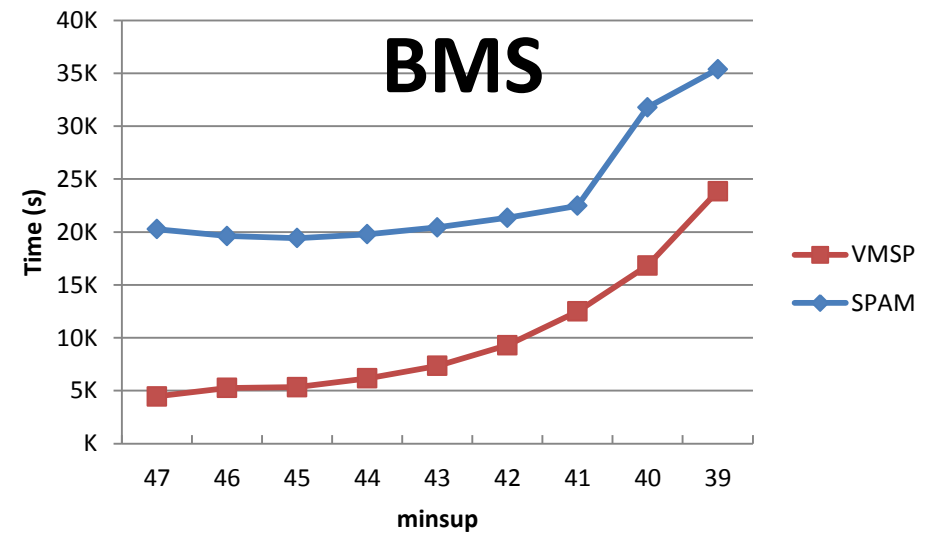
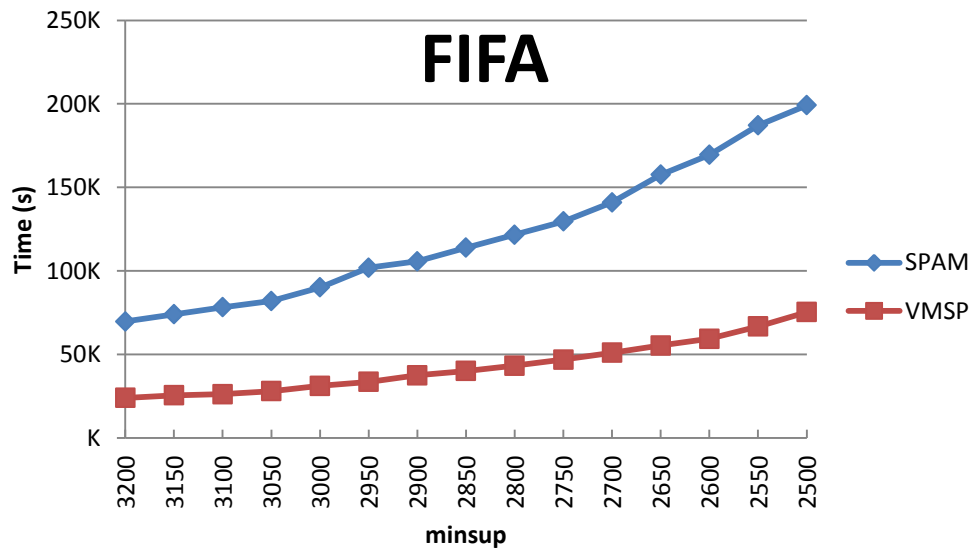
VMSP_W3 : without CPC strategy

VMSP_W2W3: without FME and CPC

VMSP W1W2W3: without FME, CPC and EFN

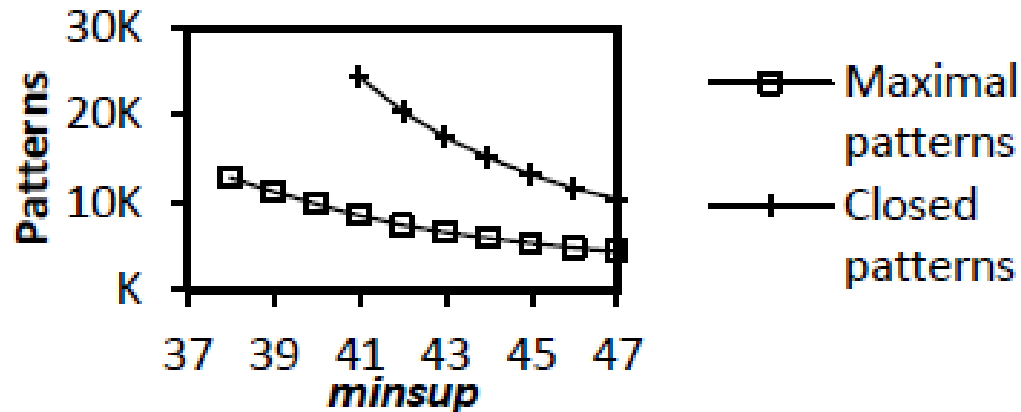
- Strategies improves the speed by up to 8 times.
- CPC is the most effective strategy

SPAM vs VMSP

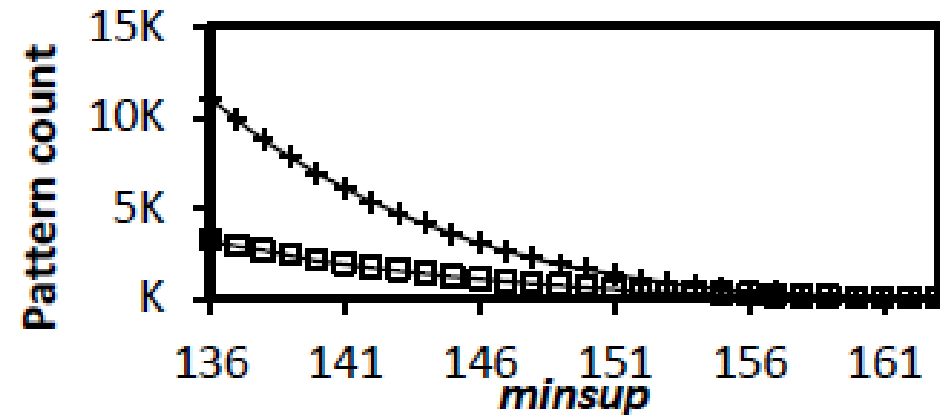


Pattern count

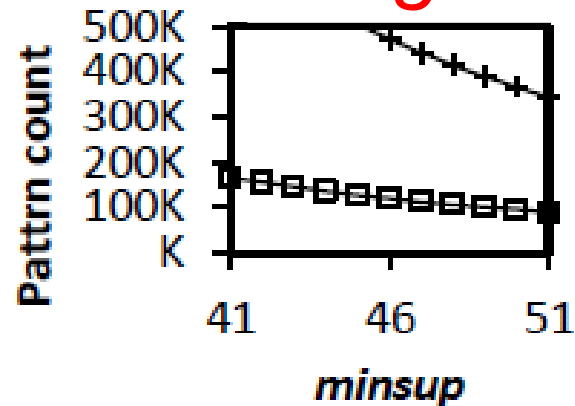
BMS



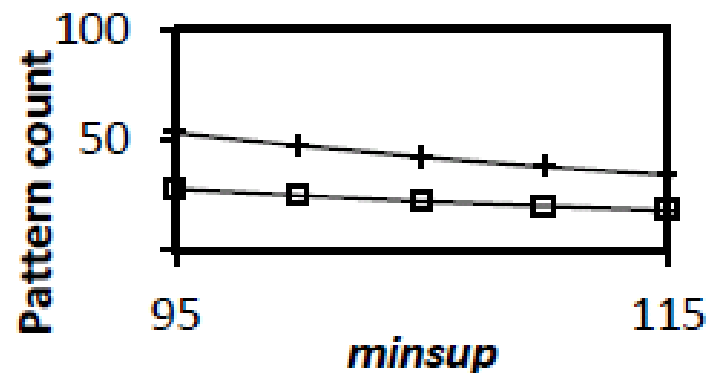
Snake



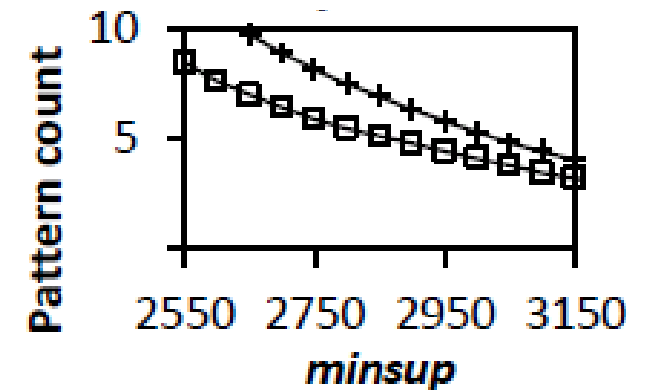
Sign



Leviathan



FIFA



Much less maximal sequential patterns than closed patterns

eg.: Snake – 28 %, Sign = 25 %

Conclusion

- **VMSP**

- a new vertical algorithm to discover **maximal sequential patterns**

- includes three novel strategies:

- **EFN**: Efficient Filtering of Non maximal patterns

- **FME**: Forward-Maximal Extension checking

- **CPC**: Candidate pruning with Co-occurrence map

- up to 100 times faster than **MaxSP**

- Source code and datasets available as part of the **SPMF data mining library** (GPL 3).



Open source Java data mining software, 66 algorithms
<http://www.philippe-fournier-viger.com/spmf/>

Thank you. Questions?



Open source Java data mining software, 55 algorithms
<http://www.philippe-fournier-viger.com/spmf/>

Applications of SPMF

- Web usage mining
- Stream mining
- Optimizing join indexes in data warehouses
- E-learning
- Smartphone usage log mining
- Opinion mining on the web
- Insider thread detection on the cloud
- Classifying edits on Wikipedia
- Linguistics
- Library recommendation,
- restaurant recommendation,
- web page recommendation
- Analyzing DOS attack in network data
- Anomaly detection in medical treatment
- Text retrieval
- Predicting location in social networks
- Manufacturing simulations
- Retail sale forecasting
- Mining source code
- Forecasting crime incidents
- Analyzing medical pathways
- Intelligent and cognitive agents
- Chemistry