

Bézier Curves – de Casteljau Algorithm

Examples: $n=3$, "Cubic"

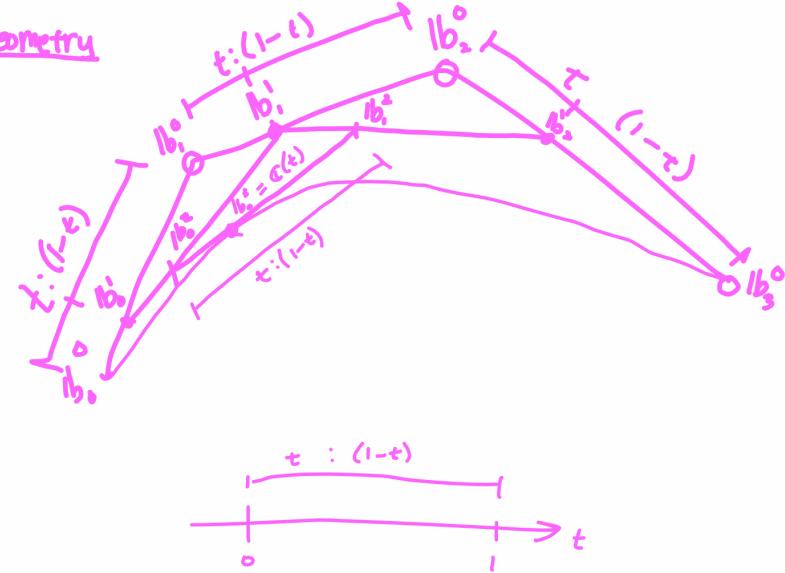


Cubic Bézier Curve

$$\begin{aligned}
 C(t) &= \sum_{i=0}^3 b_i \tilde{B}_i(t) = \sum_{i=0}^3 b_i \binom{n}{i} (1-t)^{n-i} t^i \\
 &= (1-t^3)b_0 + 3(1-t)^2 t b_1 + 3(1-t)t^2 b_2 + t^3 b_3 \\
 &= (1-t) \left\{ (1-t)^3 b_0 + 2(1-t)t b_1 + t^2 b_2 \right\} \\
 &\quad + t \left\{ (1-t)^2 b_1 + 2(1-t)t b_2 + t^2 b_3 \right\} \\
 &= (1-t) \left\{ (1-t) \langle (1-t) b_0 + t b_1 \rangle \right. \\
 &\quad \left. + t \langle (1-t) b_1 + t b_2 \rangle \right\} \\
 &\quad + t \left\{ (1-t) \langle (1-t) b_0 + t b_1 \rangle \right. \\
 &\quad \left. + t \langle (1-t) b_1 + t b_2 \rangle \right\} \\
 &\quad + t \left\{ \langle (1-t) b_0 + t b_1 \rangle \right.
 \end{aligned}$$

$$\begin{aligned}
 &= -(1-t) \{ (1-t) \mathbf{l} \mathbf{b}_0 \\
 &\quad + t \mathbf{l} \mathbf{b}_1 \} \\
 &+ t \{ (1-t) \mathbf{l} \mathbf{b}_1 \\
 &\quad + t \mathbf{l} \mathbf{b}_2 \} \\
 &= (1-t) \mathbf{l} \mathbf{b}_0^2 \\
 &\quad + t \mathbf{l} \mathbf{b}_1^2 \\
 &= \mathbf{l} \mathbf{b}_0^3 = \mathcal{C}(t)
 \end{aligned}$$

Geometry



resulting "de Casteljau scheme"

$$\begin{array}{ccccccc}
 |b_e^0 & \xrightarrow{\cdot(1-t)} & |b_o^1 & & |b_o^2 & \xrightarrow{\cdot(1-t)} & +|b_o^3 = C(t) \\
 |b_i^0 & \xrightarrow{\cdot t} & |b_i^1 & & |b_i^2 & \xrightarrow{\cdot t} & \\
 |b_e^0 & \xrightarrow{\cdot(1-t)} & |b_e^1 & & |b_e^2 & \xrightarrow{\cdot t} & \\
 |b_e^0 & \xrightarrow{\cdot t} & |b_e^1 & & & & \\
 |b_e^0 & & & & & & \text{general}
 \end{array}$$

"de Casteljau Algorithm"

Input: $\{b_0^0, \dots, b_n^0, t \in [0, 1]\}$

Output: $P(t)$

Algo:

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for (j=1) to n do
    for (i=0) to (n-j) do
         $b_i^j = (1-t)b_i^{j-1} + t b_{i+1}^{j-1}$ 
return  $\{b_0^n\}$ 

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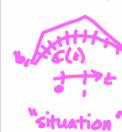
Complexity - time

→ count # of multiplications

$O(n^2)$
in terms of
multiplication

BUT:
Can use $O(n)$ -time algorithm
for evaluation,
e.g. HORNER's method

Summary

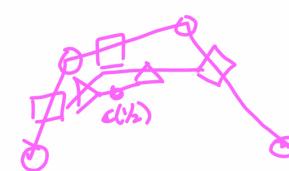


"situation"

$$c(t) = \sum_{i=0}^n b_i^i B_i^n(t)$$

algebra

geometry



scheme

o	□	△	•
o	□	△	•
o	□	△	•
o	□	△	•

Properties of Bézier curves

1. End point interpolation: $c(0) = \dots = b_0^n$
 $c(1) = \dots = b_n^n$

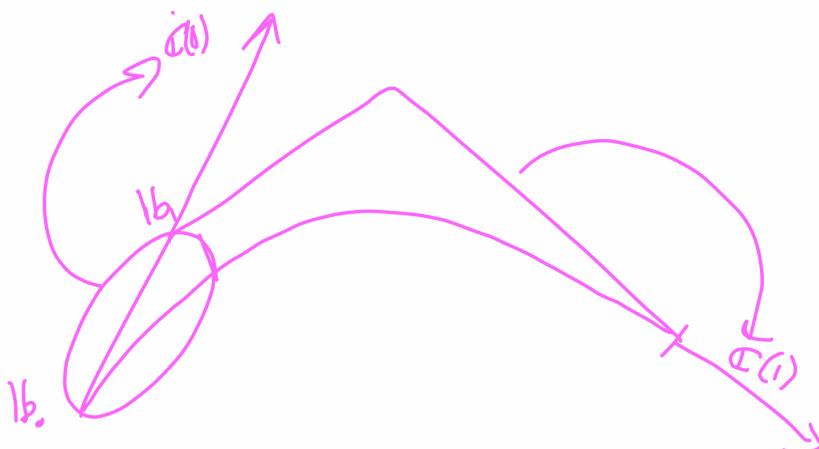
2. Derivatives/velocities @ end points:

→ differentiate curve:

$$\begin{aligned} \frac{d}{dt} c(t) = \dot{c} &= \left(\sum_{i=0}^n b_i^i B_i^n(t) \right)' = \dots \\ &= n \sum_{i=0}^{n-1} (b_{i+1} - b_i) B_i^{n-1}(t) \end{aligned}$$

$$\Rightarrow \dot{c}(0) = n(l_b - l_{b_0})$$

$$\dot{c}(1) = n(l_{b_n} - l_{b_{n-1}})$$



② Convex hull properties

" $C(t)$ is a subset of / C_i is inside the convex hull of its defining control points."



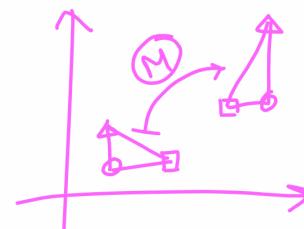
Convex hull of $\{l_{b_0}, \dots, l_{b_n}\}$
is
 $CH \{l_{b_0}, \dots, l_{b_n}\}$
 $\{IP | IP = \sum_{i=0}^n w_i l_{b_i} \text{ where } w_i \geq 0\}$

Midterm Topics

- 1) DDA, Bresenham...
- 2) Line & Polygon clipping - know names
- 3) 2D transforms
 - translate
 - scale
 - reflect
 - rotate

} apply sequence such operation to point, polygon,

Usually
"homogeneous" Matrices



- 4) Projections Parallel, perspective

Oblique projection onto ≥ 0 plane

7) Coord./Coord. system transformations



6) Decomposing Difficult

2D/3D transformations into a sequence of SIMPLE one,

7) Phong

special cases



$$I = I_a$$

$$I = I_a + I_d$$

$$I = I_a + I_d + I_s$$

8) Gouraud

