Post-Gödel Logic and Unanswered Questions

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#blairlogicmath

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Arithmetic

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- **② First incompleteness theorem**: For any sufficiently powerful theory of arithmetic A, there is a formula G such that both $A \not\vdash G$ and $A \not\vdash \neg G$.
 - Gödel-number the formulas of arithmetic
 - Define BEW(x) as "formula with Gödel-number x is provable"
 - Apply BEW to its own negation to get a formula G with Gödel-number g meaning "The formula with Gödel-number g is not provable"
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 - A simple statement of number theory cannot be proven or disproven!
- Second incompleteness theorem: Any sufficiently powerful theory of arithmetic A cannot prove its own consistency.

$$A \not\vdash \forall x (WFF(x) \rightarrow \neg (BEW(x) \land BEW(\neg x)))$$



Primitive recursive arithmetic

- There are other important theories of arithmetic besides Peano
- Many philosophers and mathematicians accept primitive recursive arithmetic as an inherently consistent starting point for arithmetic
- Constructive and finite, so it pleases intuitionists
- Quantifier-free arithmetization of arithmetic; no quantification over infinite domains means no absolute infinity
- Infinite number of variables x, y, z, ... and a symbol for each primitive recursive function
 - Basically a recursively defined function
 - You know how long many steps it will take to compute beforehand (for but not while loops)
 - Each primitive recursive function's recursive definition is included as an axiom
 - For example, ADD(x, 0) = x, ADD(x, S(y)) = S(ADD(x, y)), MUL(x,0) = 0, MUL(x,S(y)) = ADD(MUL(x,y),y) are all axioms
- Only relies on extra axioms $S(x) \neq 0$, $S(x) = S(y) \rightarrow x = y$, and $\{\phi(0), \phi(x) \rightarrow \phi(S(x))\} \vdash \phi(y)$

Gentzen's consistency proof

- ullet Gentzen proved Peano arithmetic consistent relative to primitive recursive arithmetic when we allow induction up to ϵ_0
- Gentzen's system and Peano's system are incomparable; neither is strictly stronger
 - Gentzen's system can prove that Peano's system is consistent, whereas Peano's can't (second incompleteness theorem)
 - Peano's system can compute certain functions (e.g. Ackermann function) that aren't primitive recursive
- Many people see Gentzen's proof as "as close as we can get" to an absolute proof of the consistency of PA

Logic

Löwenheim-Skolem theorem

- Löwenheim-Skolem theorem: If a set of first-order axioms has an infinite model, then it has a model of cardinality κ for all infinite cardinals κ
- First-order theories cannot prevent arbitrarily large models
- All first-order axiomatizations with infinite models are non-categorical
 - Gödel's theorems also imply this for sufficiencly strong theories
 - Axioms are supposed to completely categorize the objects they define, but if they're not categorical, they just can't
- Leads to Skolem's paradox
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- Leads to **Skolem's paradox**
 - Löwenheim-Skolem says set theory has countable models
 - Do we know that there are uncountable objects in set theory?
 - Yes, by Cantor's theorem, so what?!??!
- Resolution: externally countable sets can appear uncountable inside of countable models because no bijection with $\mathbb N$ exists inside the model but one exists outside the model
- People started realizing that even first-order logic is really hard

Kinds of logic

- Distinguish simple syntax, semantics (meaning), and proof theory (what proofs exist)
- **Propositional logic**: Everything is either true or false
 - Connectives combine propositions to create new propositions
 - Semantics: truth tables
- First-order logic: Quantify over a domain
 - Same as propositional but variables, predicates, constants, functions, and quantifiers are added
 - Used for all our axiomatizations (Peano, ZFC, etc.)
 - Non-categorical
 - Semantics: model theory
- Second-order logic: Quantify over predicates/relations
 - Can go even higher order
 - More expressive but also much more difficult to work with
 - Incompleteness, etc. still apply, but we can bypass some Löwenheim-Skolem issues
- Note that, for example, Peano arithmetic in FOL requires infinite axioms but only five in SOL

Zermelo-Frankel set theory

- Foundational system still in use today
- We know that ZF is incomplete and cannot demonstrate consistency internally by Gödel's theorems, since we can construct a model of Peano arithmetic
- We actually know two statements that are independent of ZF (true in some models, false in others)
 - Axiom of choice (AC): can choose an element from each set in a (possibly very large) collection of sets
 - Continuum hypothesis (CH): $\aleph_1 = 2^{\aleph_0}$
- Gödel showed that ZF + AC and ZF + CH are both consistent relative to ZF
- Cohen introduced **forcing** to show that $ZF + \neg AC$ and $ZF + \neg CH$ are also consistent relative to ZF
- Thus, AC and CH are both independent of ZF-a concrete example of how ZF is incomplete

Conclusions

Difficulties

- Hilbert's program didn't really work
- Hilbert's second problem-prove arithmetic is consistent-cannot be solved outright (Gödel's incompleteness theorem)
- Hilbert's Entscheidungsproblem-find a decision procedure for formulas of FOL-cannot be solved outright (halting problem)
- Hilbert's tenth problem-find a decision procedure for Diophantine equations (polynomial equations with integer coefficients and solutions)-cannot be solved outright (Matiyasevich reduced it to the halting problem)
- Arithmetic is incomplete, undecidable, non-categorical, and possibly inconsistent
- Since our theories like ZFC and Peano arithmetic are incomplete,
 Platonism suffers a huge setback

Current work

- Keep exploring other branches of math that were never really bothered by foundational issues
- ZFC is basically the universally accepted foundation of math
 - We don't think we can prove that it's absolutely consistent, but nobody has found a contradiction so far
- Explore other foundational areas of math
 - Category theory
- Do computer science
- Keep studying logic and do really hard things, possibly hoping to discover errors or new methods that will bypass current difficulties
 - New axiomatizations of arithmetic
 - New set theories
 - New logical systems
- Incorporate philosophical arguments... some Platonists still hold out hope!

Our heroes

- Alan Turing
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Kurt Gödel

- Arguably the most pivotal logician (along with Aristotle and Frege)
- Massive contributions demolished Hilbert, Russell, and Brouwer
- Targeted by Nazis for mathematical collaborations with Jews and conscripted; fled to America in 1940
- Became besties with Einstein
 - Einstein had to coach him on how not to derail his citizenship ceremony by pointing out logical flaws in the Constitution
 - Lowkey showed that time travel is possible in Einstein's theories as a "birthday present" and caused Einstein to doubt his entire career
- Obsessive fear of poisoning and would only eat wife's cooking; starved to death in 1974 when she was hospitalized

Our heroes, continued

Bertrand Russell

- Started working outside of mathematics after the Principia
- Huge pacifist and anti-war activist
 - Arrested for speaking out in WWI
 - WWII: "War was always a great evil, but in some particularly extreme circumstances, it may be the lesser of two evils."
- Won the Nobel Prize in Literature for contributions to "freedom of thought"

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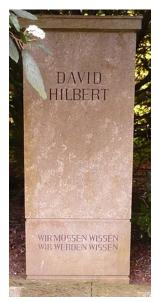
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David Hilbert

- Nazis purged the University of Göttingen in 1933, most promiment mathematicians were Jewish or had Jewish family (Weyl, Noether, Landau, etc.)
- Hilbert remained and died in obscurity 1943; funeral attended by ¡10 people

Spirit



Unanswered questions

In mathematics:

- What does it mean for an object to exist?
- Mow do we interpret mathematical results?
- Why is math useful in the real world? Should it be?
- Should we even bother doing math?
- What is truth? How can we verify it?