

Agenda

- BnF and digital preservation
- SPAR
- The Storage Abstraction Service
- Limitation of the Abstraction
- Conclusion

The National Library of France





An autonomous public establishment

2200 agents and dozens of professions, a budget of aprox. 230 M€

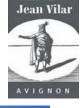
Local, national and international missions

About 1M readers per year, 300 000 visitors to the exhibitions

Large collections to manage

More than 14M books
More than 30M posters and photos
More than 250.000 manuscripts
More than 1M audiovisual material

"French" Web legal deposit











Scalable Preservation and Archiving Repository

The National Library of France (BnF)

Main missions:

- to build up the collections
- to preserve them forever
- to communicate them to the public

Legal deposit:

- legal deposit since 1537 for printed materials
- 1648: engravings and maps
- 1793: musical scores
- 1925: photos
- 1938: phonograms
- 1941: posters
- 1975: videograms and multimedia documents
- 1992: audiovisual and electronic documents
- 2006: Web legal deposit













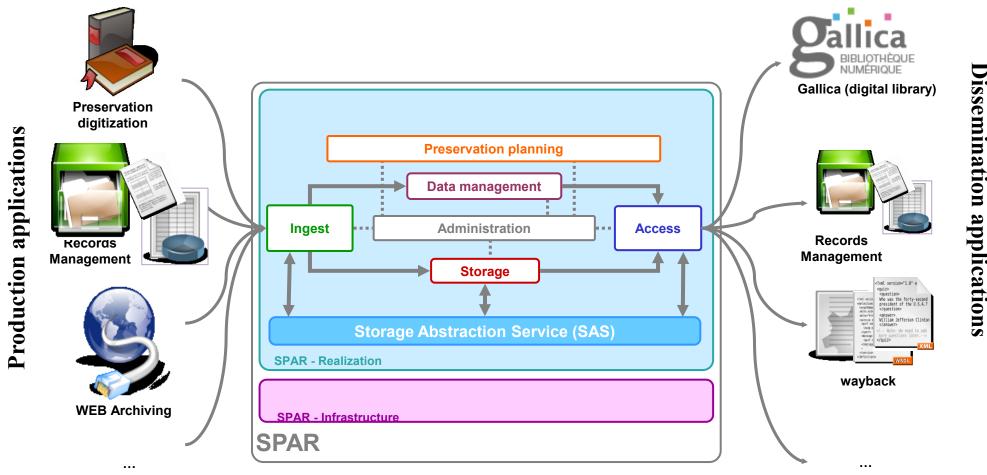
Decree n°94-3 of january 3rd of 19994 – record 2 : La Bibliothèque nationale de France a pour missions de

collecter, cataloguer, **conserver** et enrichir dans tous les champs de la connaissance, **le patrimoine national dont**

Institutional issue: missions

- Preservation is at the heart of BnF's missions
- Digital preservation is a direct extension of the preservation of BnF's collections. Two different approaches can be considered:
 - The digital form as a mean to preserve an analog document that degrades. It then enters in a digitization process (as it's the case for paper but also for analog audiovisual collections) and the digital surrogate (which may be a substitution) needs to be preserved along side with the analog
 - 2. The born-digital document (in the case of the web legal deposit since august 1st, 2006, or of the audiovisual legal deposit). This document then belonging to the heritage collections has to be preserved.

Digital archiving at BnF



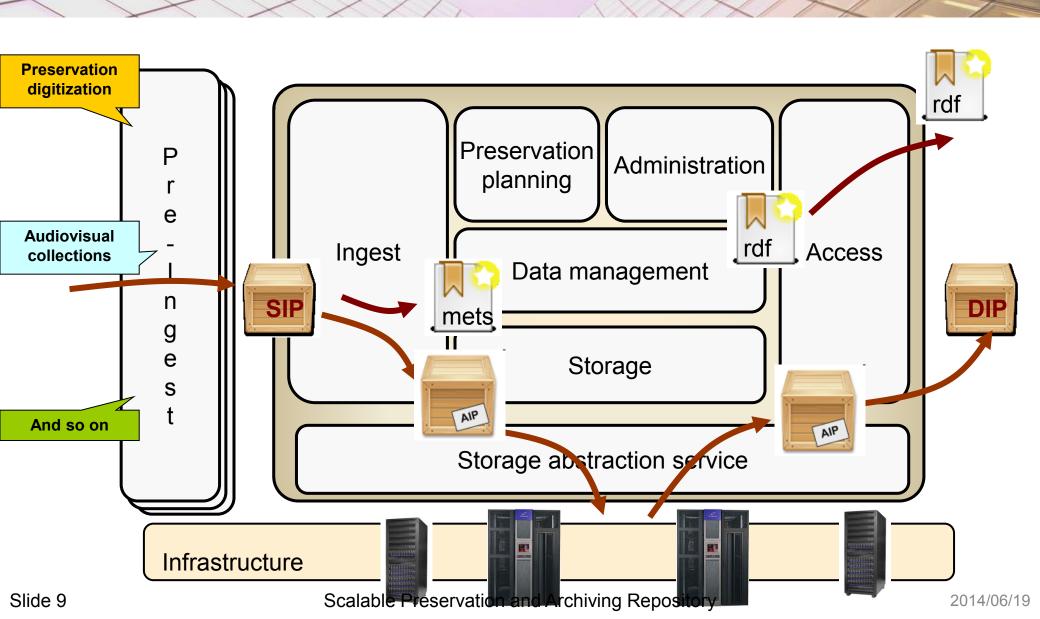
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Some facts about SPAR (at 2014/06/16)

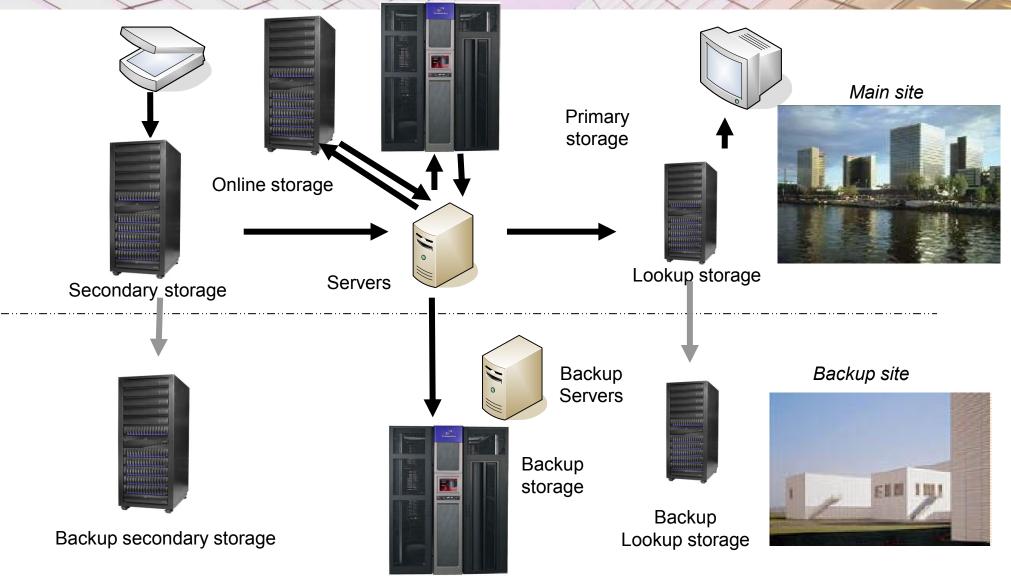
- Operation begins: may 2010
 - Current workflow of BnF's mass digitization program
 - Distributed storage over 2 sites
- 2 668 000 packages archived
- 266 703 000 data-objects (files)
- > 1,36PB (=1000TB) of raw data
- > 1 billion triples (elementary metadata)
- Very large manuscripts have been ingested (packages with size of 100GB)
- · First technological migration made: new tape generation
- Available tracks: monographs, periodicals, still images, audio, video, web archiving, office documents, third-party archiving

SPAR: modular, OAIS compliant repository



Infrastructure





Infrastructure - Kind of hardware

Tapes

- Less costly
- High capacity
- Energy efficient
- Slow (reach the tape, mount it in the drive, load the information...)
 - ⇒ For archival documents and those that don't need quick access

Disks

- Fast access
- Most expensive
 - ⇒ For documents in the workflow
 - ⇒For documents ready for quick dissemination

(BnF

Decomposition in tracks

- To deal with the variability and heterogeneity of the data, definition of tracks
- build on the relation between the digital objects and the archival system, independently of any given organization:
 - Preservation digitization
 - Audiovisual legal deposit
 - Negotiated legal deposit (e-books, large posters,...)
 - Automatic legal deposit (surface Web)
 - Administrative production
 - Third party archiving
 - Acquisition / Donation
 - + reference track

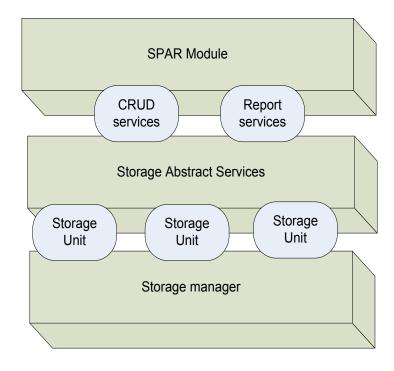
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Requirements for the Storage abstraction service



- Abstract the infrastructure with:
 - storage unit
 - aggregates storage element to defines an abstract storage defined by a class of service
 - mean time (read/write)
 - number of copies …
 - record
 - simple bit stream (no semantics)
 - some properties : checksum, logical name, lastAuditDate







Main concepts for the SAS





• A record is a information managed by the SAS. Each record is hosted on one and only one storage unit.



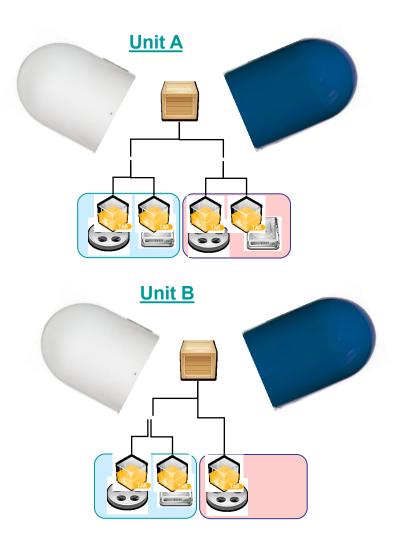
• A copy (or replica): Depending on the needs, a record can have multiple copies on the SAS. All copies are strictly identical (integrity control). The SAS only exposes one record with multiple copies. Each copy is written physically in a storage element.



• A storage unit: It handles an entity to capture the characteristics of a particular storage. Each unit is declared by an administrator who defines which elements of storage are linked as well as the number of copies. Every record in one storage unit has the same characteristics of storage.

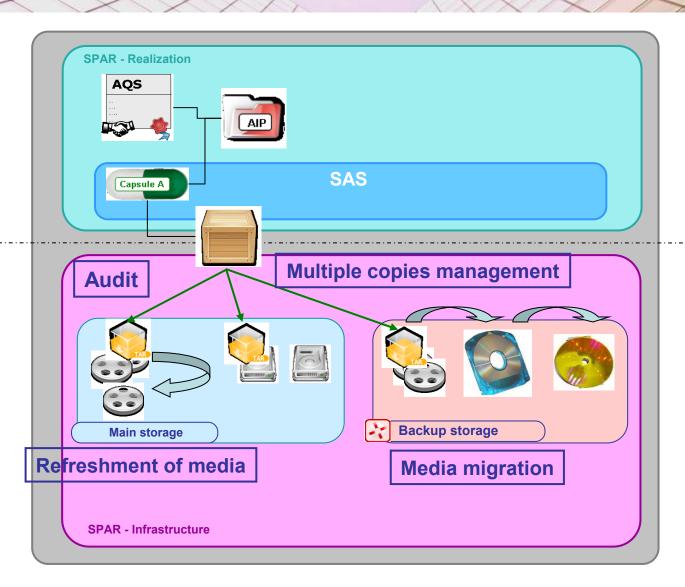


• A storage element: it is an entity very closed to the hardware (media, disk, tapes), except that it represents a part of a physical media. For example: a partition of a disk array or a pool of tapes...



Storage abstraction



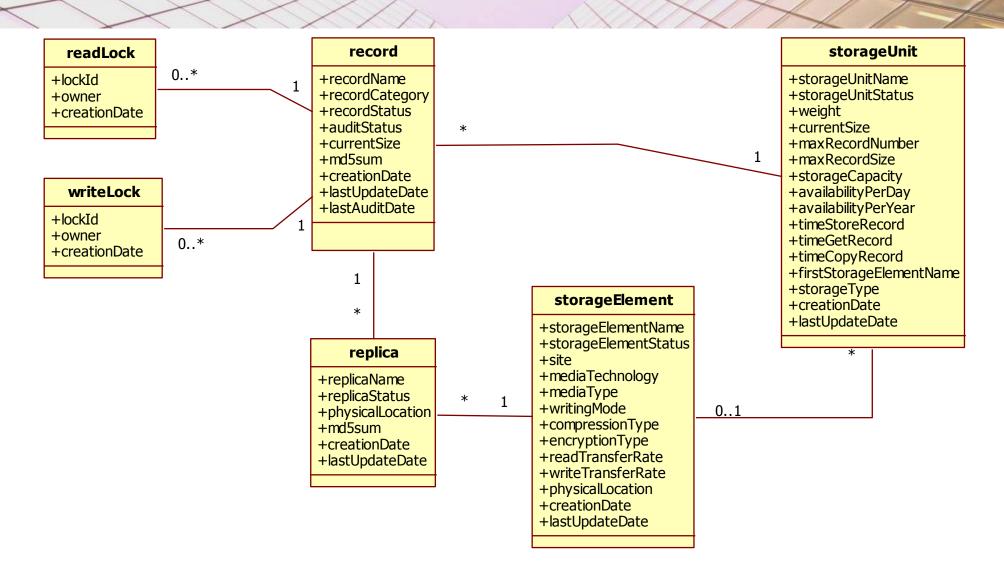


Application viewpoint

Infrastructure viewpoint

Logical Data Model





Implementation choice



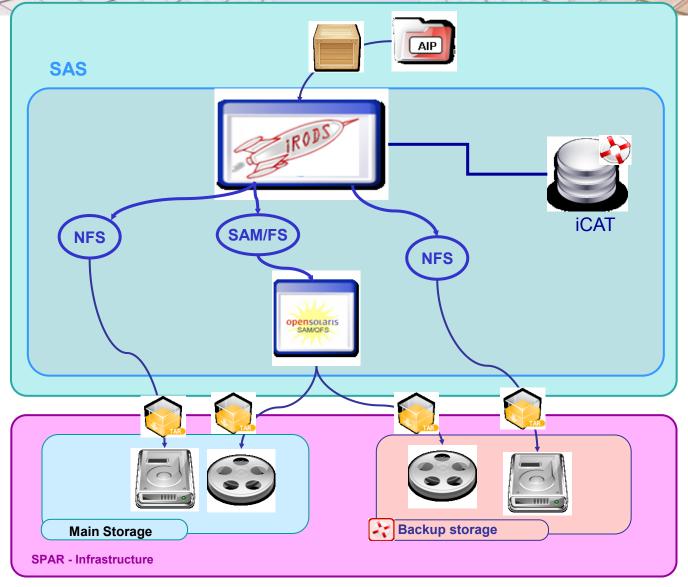
- iRods
 - scalable and proven
 - rules based



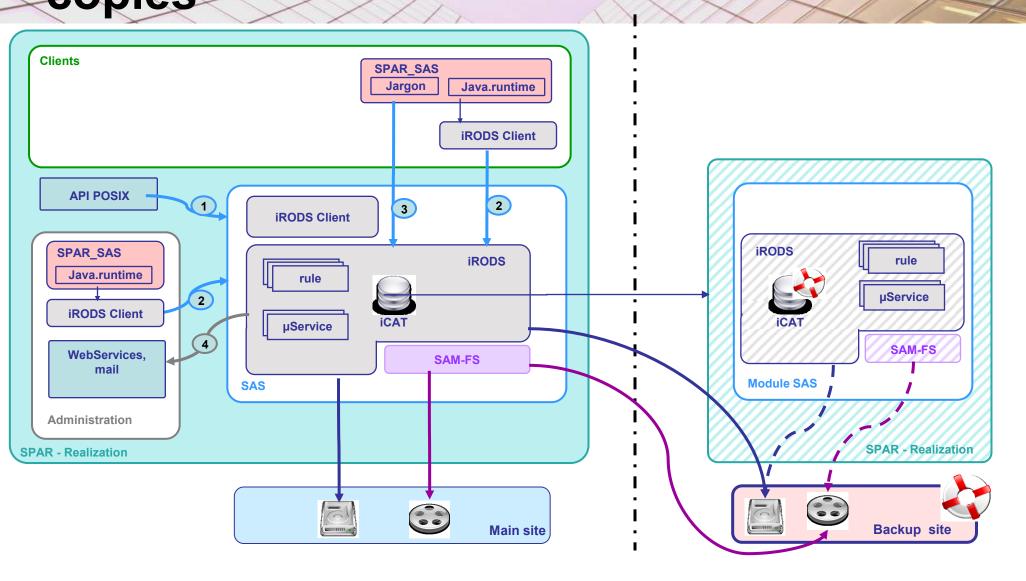
- Use of the rules to configure a storage unit
 - rule for put
 - rule for audit
 - rule for get
- General rules for
 - refreshment migration
 - duplication migration

(BnF

Technical implementation: iRODS



iRODS: Logical objects ⇔ Physical {BNF copies





- To the classic CRUD verbs, we add a Audit verb
- Implemented through rules:
 - sparPut
 - sparGet
 - sparAudit
 - sparUpdate
 - sparDElete

Dialog between Storage module and SAS

Storing an AIP

- corresponding SLA search (storage requirements)
- retrieval of available storage units with compatible class of service
- choose of the least expensive storage unit
- store the AIP as a record

Audit of an AIP

- if audit time has expired,
- asked the SAS for an audit of the copies
- retrieved the package itself for internal audit

Retrieve of an AIP

- get back the first ok copy
- if one is found bad, launch an audit

Storing an AIP

```
sasPUT(*rodsName, *cPath, *localFilePath, *mainRsrc, *inChksum)
  *putStatus = -1;
                                                   Save replica 1 with checksum test
  *rodsPath = "*cPath/*rodsName";
  acObjPutWithDateAndChksumAsAVUs(
    *rodsPath, *mainRsrc, *localFilePath,
                                                             Query for replica
    *inChksum, *putStatus);
                                                                resources
  if (*putStatus == 0) {
   acGetResourceMetaAttribute(*mainRsrc, "replicaResources", *replList);
    if (*replList != "" && *replList != "null") {
      acSetDataObjAttribute(
                                                                      Get the list
        *rodsName, *cPath, "recordStatus", "WRITE SCHEDULED");
      delay("<PLUSET>1m</PLUSET>") {
        acReplicateFile(*replList, *rodsPath, *cPath, *rodsName);
    } else {
                                                      Plan the creation of each replica
      acSetDataObjAttribute(
        *rodsName, *cPath, "recordStatus", "OK");
```





```
sasAUDIT(*rodsName, *parentColl, *stageDir,
 *replicasToReplace = list("");
  *rodsPath = "";
 acGetFullPathFromParentCollection(*rodsName, *parentColl, *rodsPath);
 *goodReplicaEncountered = 0;
 acGetDataObjMetaAttribute(*rodsName, *parentColl, "MD5SUM", *objStoredChksum);
 acGetDataObjLocations(*rodsName, *parentColl, *matchingObjects);
                                                                                                Retrieve all the
 foreach (*obj in *matchingObjects) {
                                                                                                     replicas
   msiGetValByKey(*obj, "DATA_PATH", *objPhysicalPath);
msiGetValByKey(*obj, "RESC_NAME", *currRescName);
   *objPhysicalMD5 = "";
   msiExecCmd("rodsMD5sum", "*objPhysicalPath", "null", "null", "null", *cmdOut);
   msiGetStdoutInExecCmdOut(*cmdOut, *objPhysicalMD5);
   if (*objStoredChksum == *objPhysicalMD5) {
                                                                                                    Calculate the
      *goodReplicaEncountered = 1;
    } else {
      *replicasToReplace = cons(*replicasToReplace,"*currRescName");
                                                                                                     checksum
 if (*goodReplicaEncountered == 1) {
   if (*replicasToReplace != "") {
      acSetDataObjAttribute(*rodsName, *parentColl, "recordStatus", "DIRTY");
      *allReplicaReplacementOk = 1;
      foreach (*replica in *replicasToReplace) {
        acReplaceStaleReplica(*rodsPath, "*replica", *replStatus);
        if (*replStatus != 0) {
                                                                                      Replace the bad replica
          *allReplicaReplacementOk = 0;
      if (*allReplicaReplacementOk == 1) {
        acSetDataObjAttribute(*rodsName, *parentColl, "recordStatus", "OK");
    } else {
      acSetDataObjAttribute(*rodsName, *parentColl, "recordStatus", "OK");
                                                                                              Update the record
  } else {
                                                                                                      status
   acSetDataObjAttribute(*rodsName, *parentColl, "recordStatus", "ERROR");
```

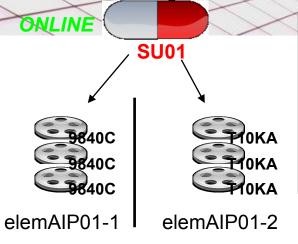
Retrieving an AIP

Slide 25

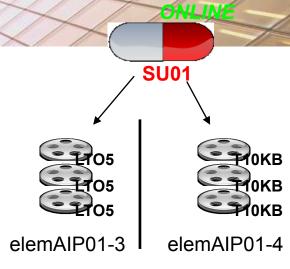
```
sasGET(*rodsName, *parentColl, *stageDir, *getStatus) {
  *replicasToReplace = list("");
                                                                                       No good replica yet
  *rodsPath = "";
                                                                                               found
 acGetFullPathFromParentCollection(*rodsName, *parentColl, *rodsPath);
  *goodReplicaEncountered = 0;
 acGetDataObjMetaAttribute(*rodsName, *parentColl, "MD5SUM", *objStoredChksum);
 acGetDataObjLocations(*rodsName, *parentColl, *matchingObjects);
 foreach (*obj in *matchingObjects) {
   msiGetValByKey(*obj, "DATA_PATH", *objPhysicalPath);
msiGetValByKey(*obj, "RESC_NAME", *currRescName);
    *objPhysicalMD5 = "";
   msiExecCmd("rodsMD5sum", "*objPhysicalPath", "null", "null", "null", *cmdOut);
   msiGetStdoutInExecCmdOut(*cmdOut, *objPhysicalMD5);
    if (*objStoredChksum == *objPhysicalMD5) {
      *goodReplicaEncountered = 1;
      break:
                                                                                 Calculate the checksum
    } else {
      *replicasToReplace = cons(*replicasToReplace,"*currRescName");
 if (*goodReplicaEncountered == 1) {
    if (*replicasToReplace != "") {
      acSetDataObjAttribute(*rodsName, *parentColl, "recordStatus", "DIRTY");
      foreach(*replica in *replicasToReplace) {
        delay("<PLUSET>1m</PLUSET>") {
          acReplaceStaleReplica(*rodsPath, *replica, *replStatus);
                                                          Plan the replacement of the bad replica
    *localFilePath = "*stageDir/*rodsName";
    *qetContext = "qetting *rodsName to *localFilePath from resource : *currRescName ...";
    msiDataObjGet(*rodsPath, "localPath=*localFilePath++++rescName=*currRescName++++forceFlag=",
    *qetStatuš);
   else {
    acSetDataObjAttribute(*rodsName, *parentColl, "recordStatus", "ERROR");
                                                                                       Get the first good
                                                                                              replica
```

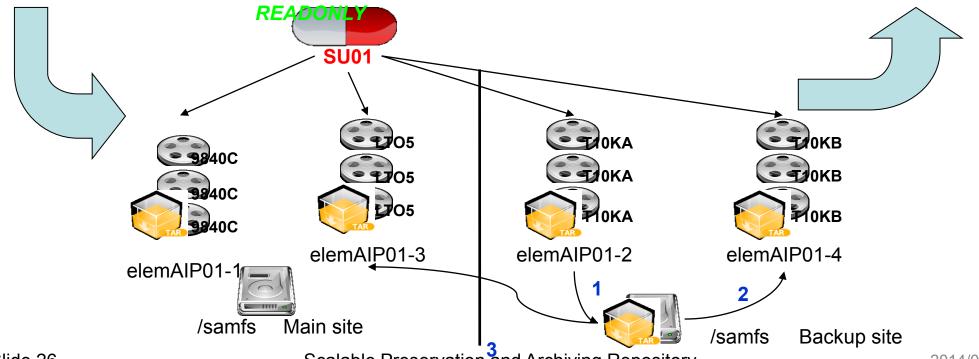
Scalable Preservation and Archiving Repository

Transparent media migration



Main site: from 9840C (40GB) to LTO5 (1.5TB)
Backup site: from T10000A (500GB) to T10000B (1TB)
3 months operation





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Issues of storage and SAS

- Provide a intelligent storage system:
 - Driven by the services it should supplied
 - Auto-healing
 - Able to make regular audits
 - Able to take into account new hardware, using replication migration
 - Able to migrate from equivalent hardware (tape change), using refreshment migration

Storage resources management

Accès

Connecté(e) en tant que : spar_adr **SPAR Réalisation** Administration

Administration > Surveillance matérielle > Suivre les capsules

Versement

Suivre les caps	sules			?
Capsule	\$ Statut	% occupation		Nb enregistrements
aip-ref	ONLINE	50	100	93
capsAIPO1	READONLY	50	100	149463
capsAIPO2	OFFLINE	50	100	2783
capsAIPO3	OFFLINE	50	100	717
capsAIPO4	ONLINE	50	100	31041
capsAIP05	OFFLINE	50	100	1

Limits on the abstraction

- With non-random access ressources, need to know where the records are
- Batch 'vsn2irods' to inform on where are located the record
- => allow for optimization of retrieval and auditing

Technical metadata

- Additionnal metadata linked to the archival package
 - sparprovenance:recordLastUpdateDate: date of the last update of a record
 - sparprovenance:recordLastAuditDate: date of the last audit of a record
 - sparreference:replicaLocations: serial numbers of the tapes where the replicas are located

Notion of plans

- The [Administration] module executes plans over a large number of archival packages
- In order for the plans to be efficient, need to know where some of the copies are located
 register the location in the data management index (triplestore)

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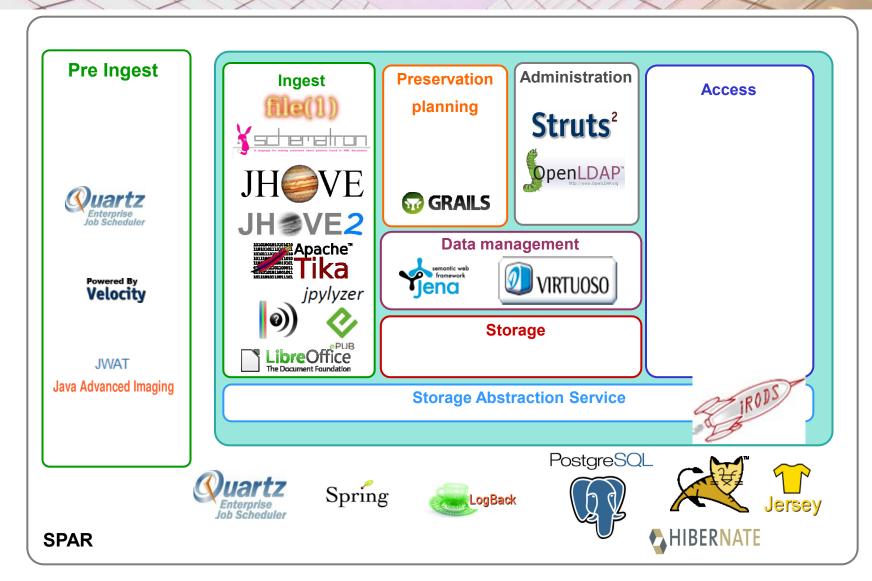


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Conclusion

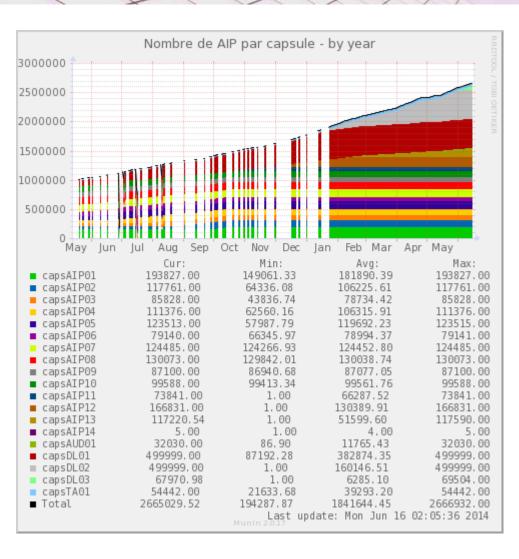
- iRods provides a great way to abstract the application from the work of the administrators
- Work to do:
 - Use of Jargon
 - Migration to iRods 4.0 and use of composable resources
 - Better integration with "tapes" resources

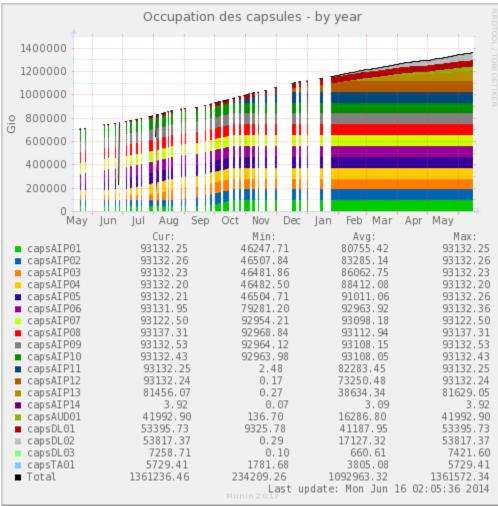




Volume







Main computer room



