Lecture 17 HEAPS

Outline

- The ADT Heap
- Heap Implementation of the ADT Priority Queue
- Array-based Implementation of a Heap
- Heap Sort

The ADT Heap

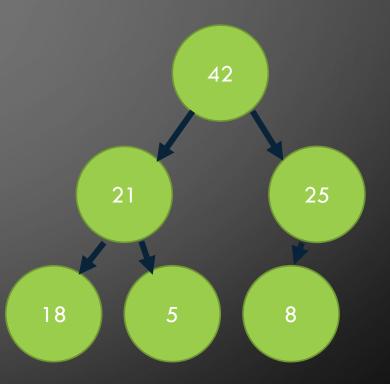
- A heap is a complete binary tree that either is
 - Empty or ...
 - Whose root contains a value ≥ each of its children and has heaps as its subtrees
- It is a special binary tree ... different in that
 - It is ordered in a weaker sense
 - it will always be a complete binary tree

```
Heap

+isEmpty(): boolean
+getNumberOfNodes(): integer
+getHeight(): integer
+peekTop(): ItemType
+add(newData: ItemType): boolean
+remove(): boolean
+clear(): void
```

Max Heap

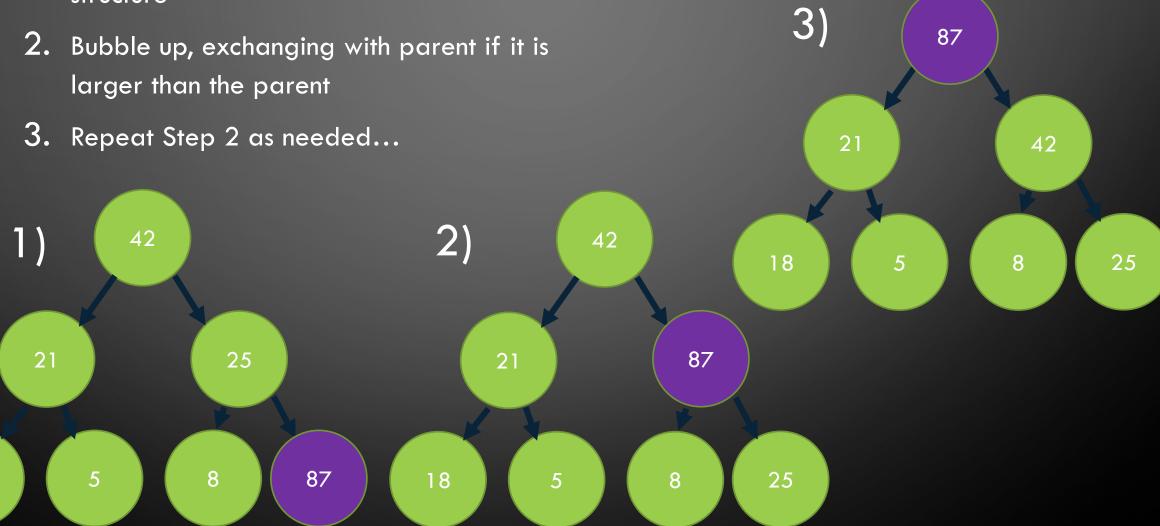
- Each parent node is larger than its children
- Only the next-in-line can be removed
- Min Heap has the opposite rule: each parent is smaller than its children



Max Heap: Push (insert)

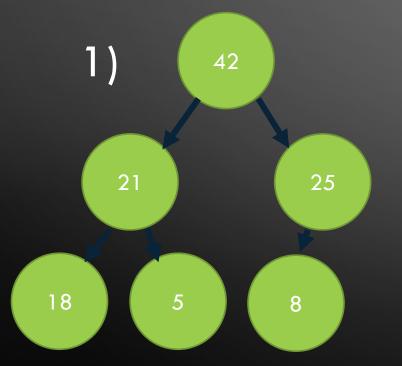
18

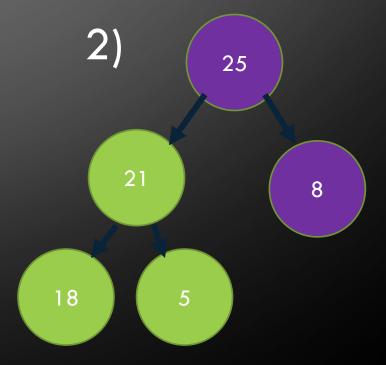
1. Add the node to the next empty spot in the structure



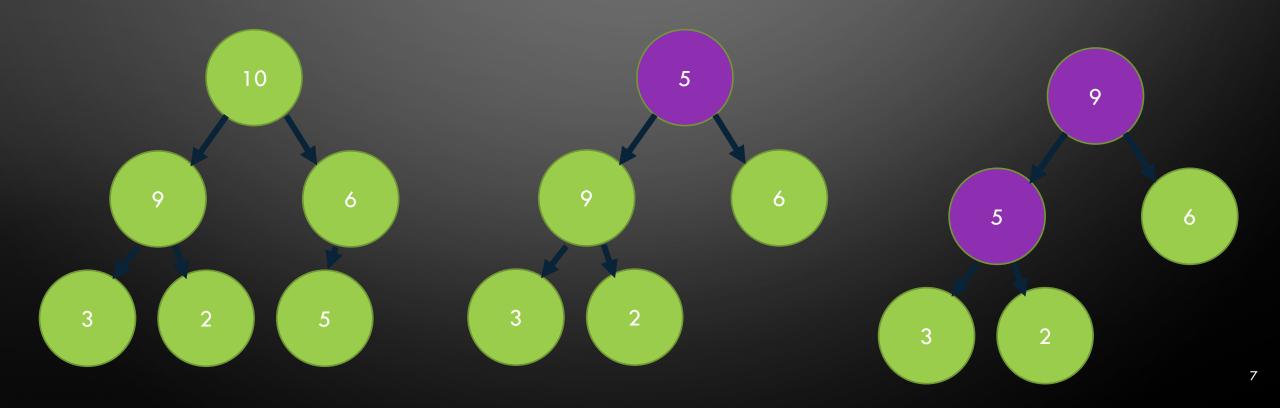
Max Heap: Pop

- Remove root
- Replace root with "last node"
- Exchange nodes with largest child as much as possible





Second Pop Example



Heap ADT

- A Heap is an ordered tree-based structure where the max (or min) object is stored at the root
 - Is this like anything we have seen?

- A Heap is not a priority queue...but
- A heap can represent the most efficient implementation of a priority queue
- Sometimes priority queues are simply called heaps

Heap Complexity

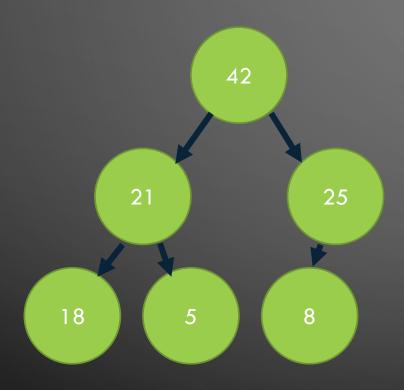
Heap top() and isEmpty() are O(1)

What is the complexity of push(item)?

What is the complexity of pop()?

Representing a Heap: Array

Because a heap is a complete binary
 tree, it can be represented using an array





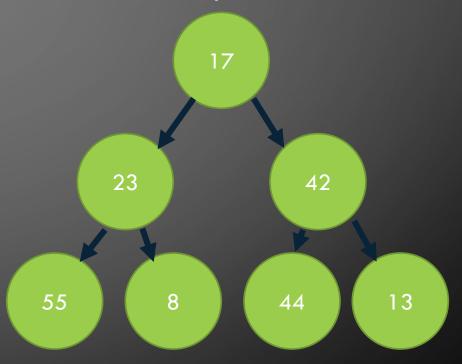
Insert Psuedocode

• The parent of position k can be represented by k >> 1

```
insert(ItemType item)
     k = heapsize + 1;
     j = k \gg 1; //j is the parent
     //while the item is larger than its parent
     while ((j \ge 1) \text{ and } (array[j] < item))
        //replace current position with its parent
        array[k] = array[j];
        k = j;
        j = j >> 1; //move to next parent
     array[k] = item;
     heapsize = heapsize + 1;
```

- Uses our new heap structure to sort an array
- 1. Start with an array, and then consider it an unsorted heap





 55
 17
 44
 23
 8
 42
 13

55

44

42

17

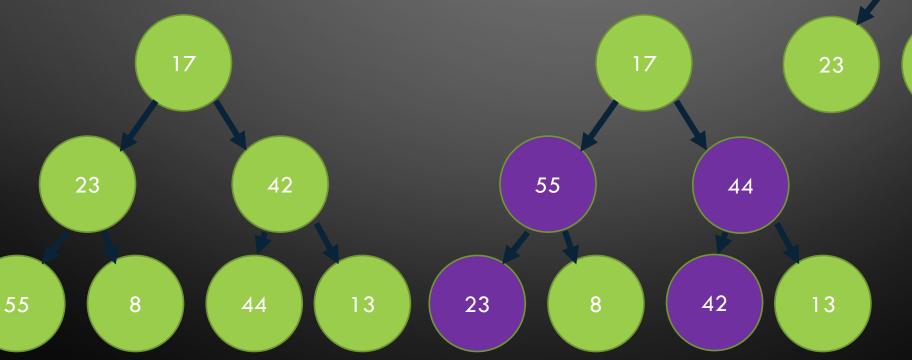
8

• Uses our new heap structure to sort an array



This is done by swapping each element with its children if they are out of place.

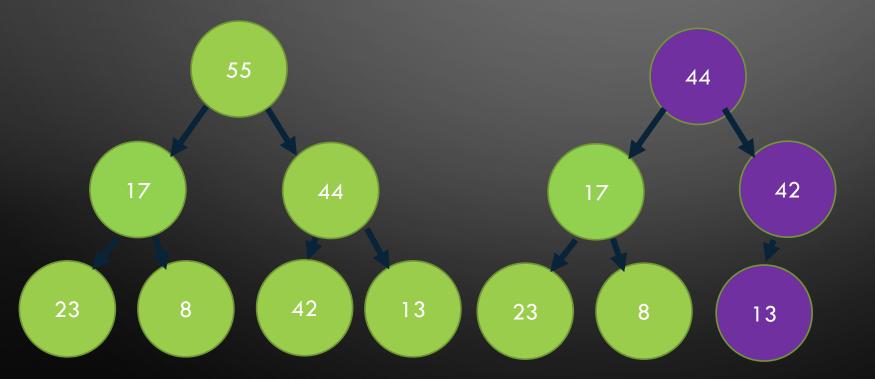
Note that at the end, 17 is not sorted correctly



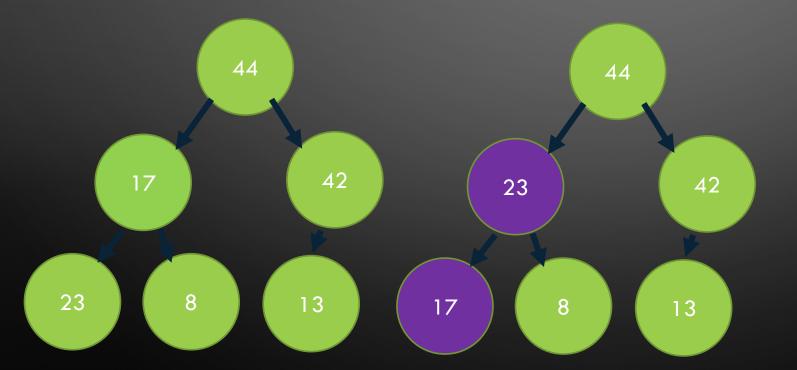
44 17 42 23 8 13 55

- Uses our new heap structure to sort an array
- 3. Remove root, keep in separate sorted array

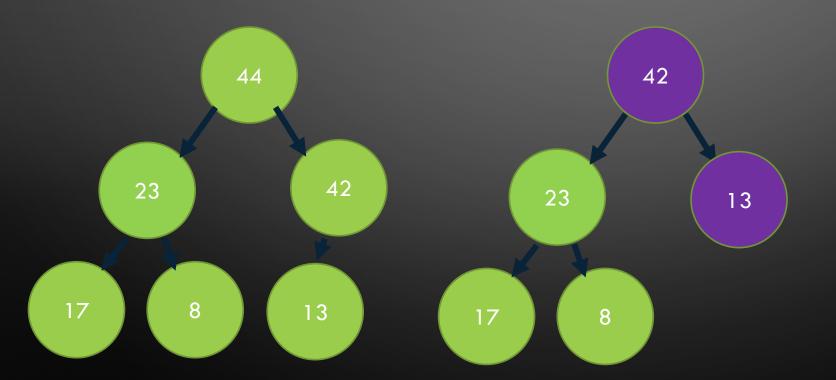
Promote largest child



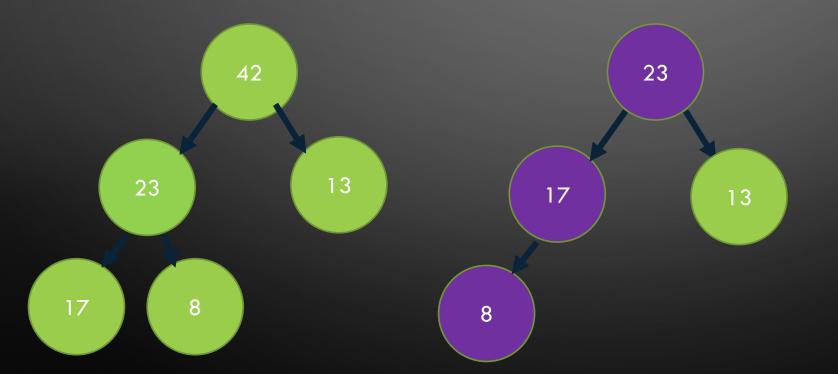
- 44 23 42 17 8 13 55
- Uses our new heap structure to sort an array
- 4. Repeat steps 2-3 until sorted



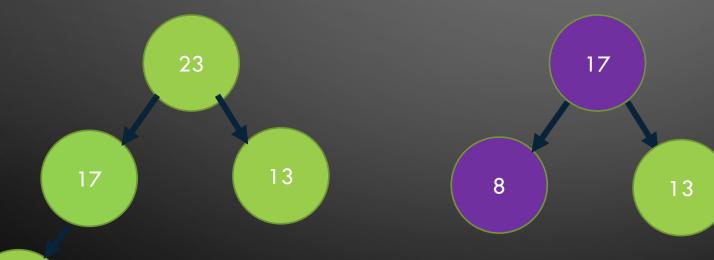
42 23 13 17 8 55 44



23 17 13 8 55 44 42



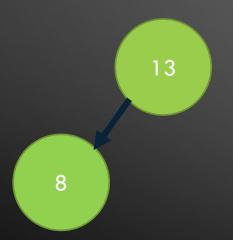
 17
 8
 13
 55
 44
 42
 23



 13
 8
 55
 44
 42
 23
 17



8 55 44 42 23 17 13





55 44 42 23 17 13 8



Heap Sort Complexity

- What is the time complexity of heap sort?
- Average Case: O(n log n)
- Worst Case: O(n log n)
- What is the space overhead of heap sort?
- O(1)

Assignment/Homework

- Reading pp. 537 -545, 547-587
- P4 (Courseweb) due on Today.
- ICE 11: BST (Courseweb) due on Thursday.
- Homework 8: BST (Gradescope) due on Thursday.