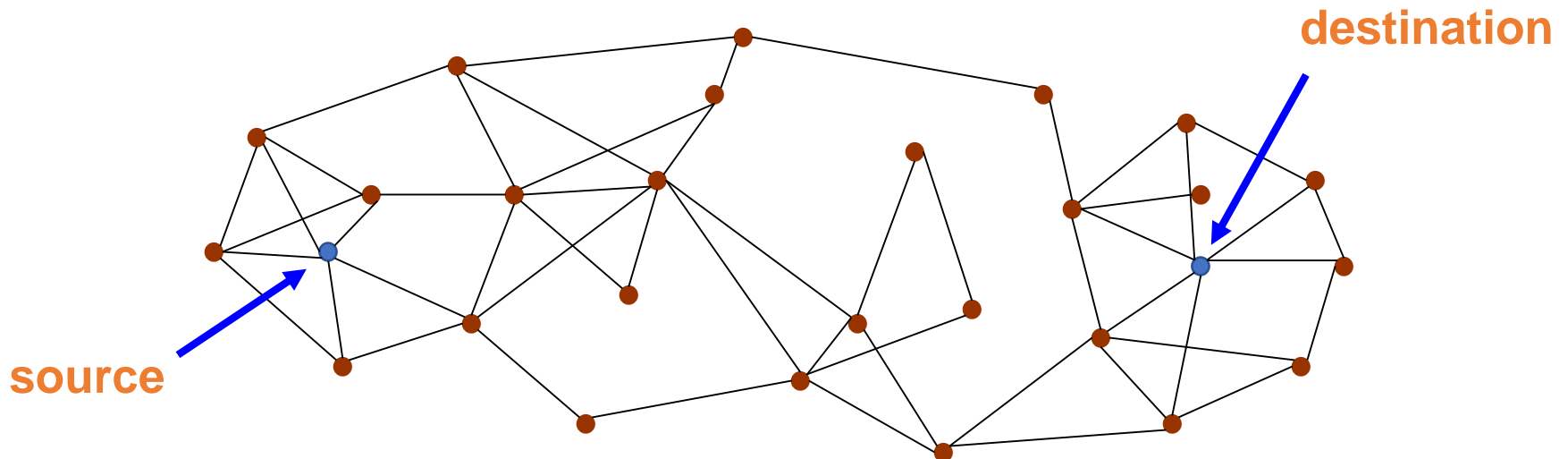


# Data Structure Programming Project #1

郭建志

# Background

- Large-scale geographic routing
- Every node only has the **local** information (i.e., its neighbors' information) and the **destination's** information



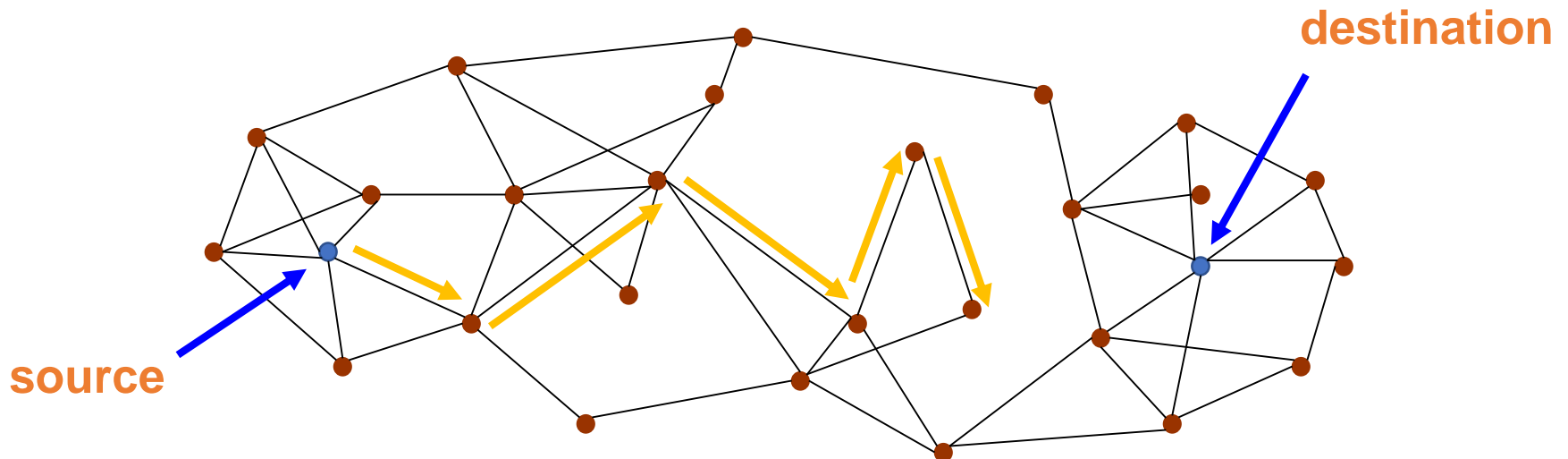
# Background

- Large-scale geographic routing
- Every node only has the **local** information (i.e., its neighbors' information) and the **destination's** information
- So, how should a source route a packet to a designated destination?



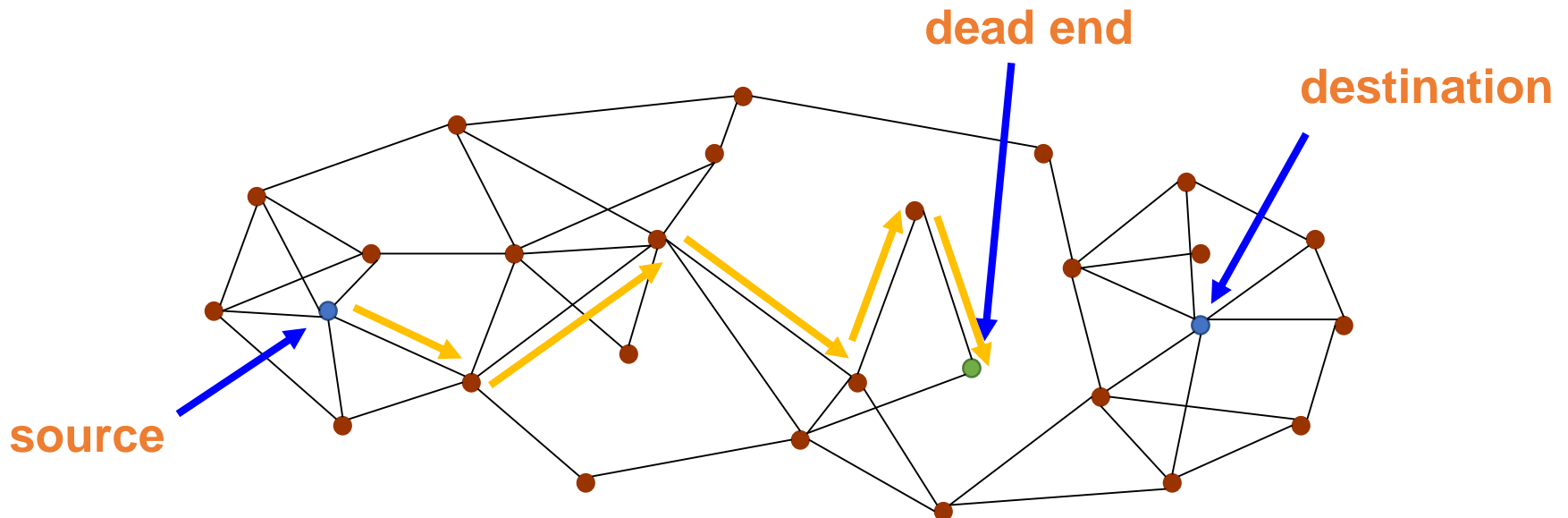
# Naïve Routing Scheme: Greedy

- Forward the packet towards the neighboring node that is **closer** to the destination



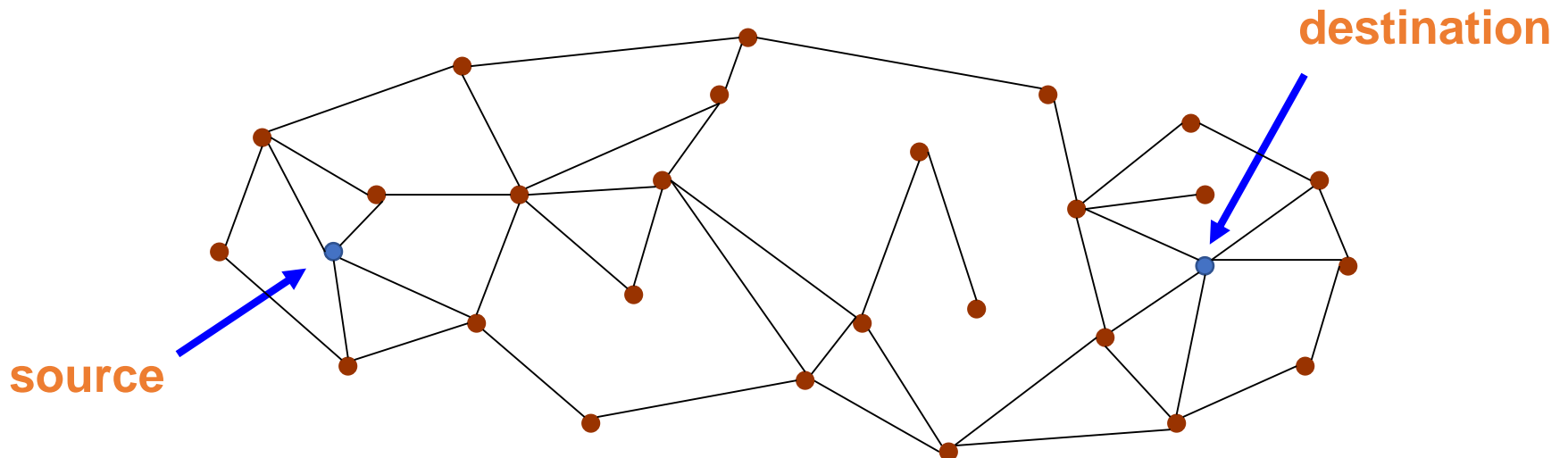
# Naïve Routing Scheme: Greedy

- Forward the packet towards the neighboring node that is **closer** to the destination
- However, the packet may be **stuck** at the dead end



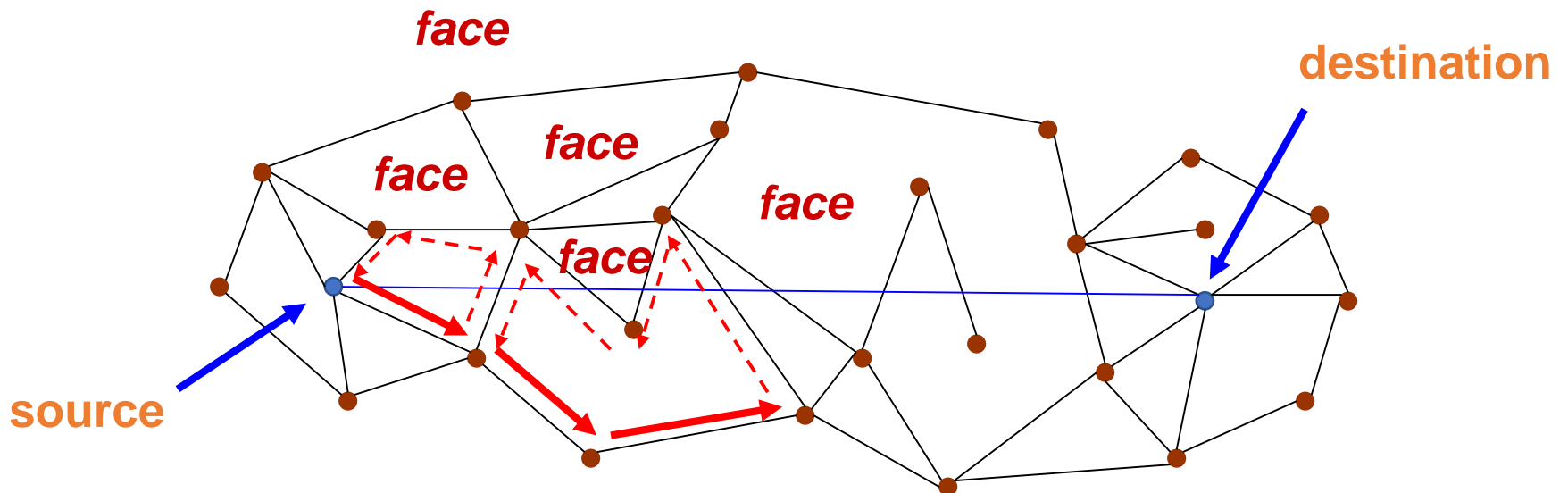
# Face Routing

- Generate a planar subgraph: no crossing links
- Right-hand rule
- Change faces



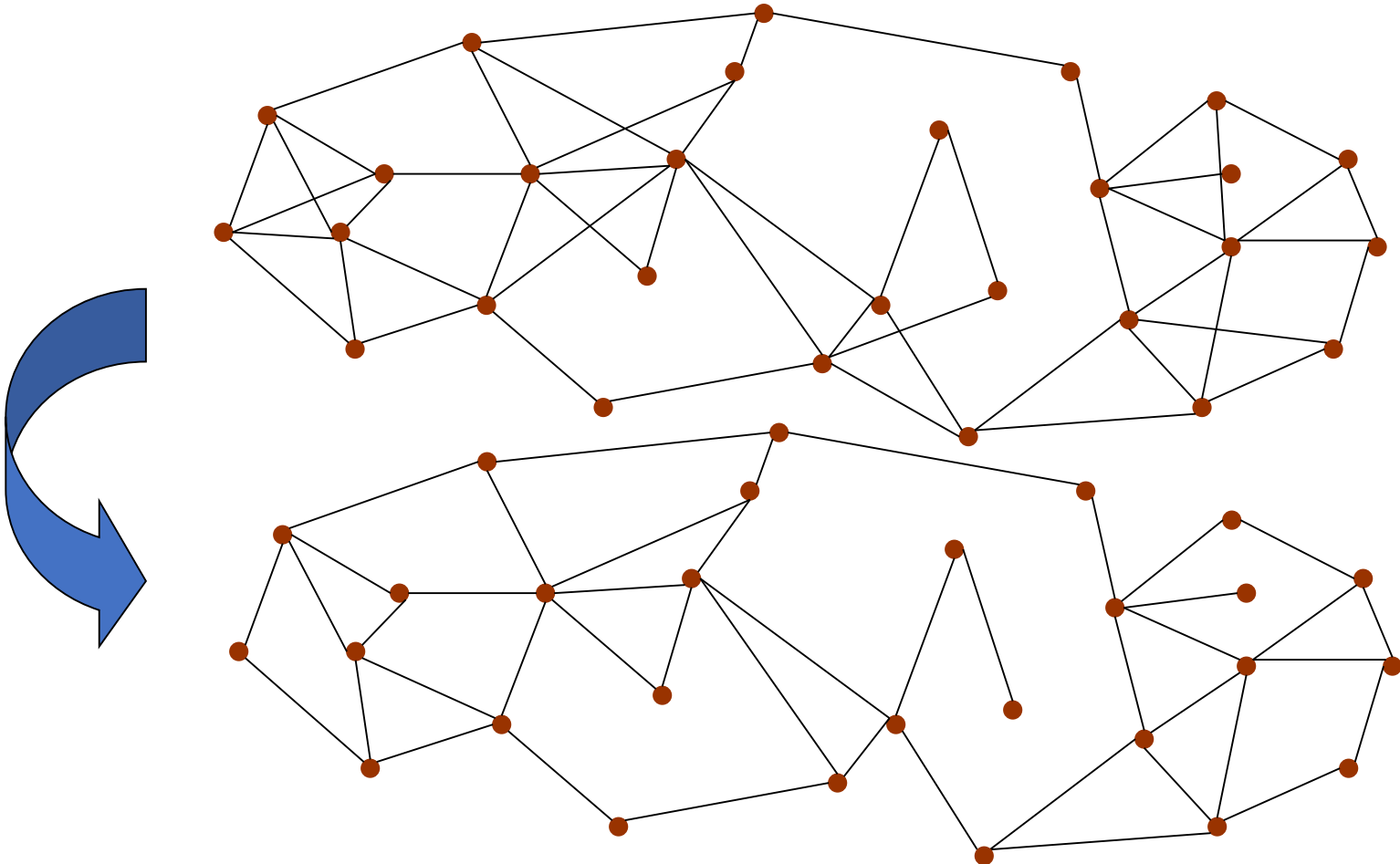
# Face Routing

- Generate a planar subgraph: no crossing links
- Right-hand rule
- Change faces



# Programming Project #1: Generate a Planar Subgraph

- Generate a planar subgraph: no crossing links





# Programming Project #1: Generate a Planar Subgraph

- Input:

- Number of nodes
- Nodes with non-negative coordinates (x, y) (the input graph is connected when we add links if  $\text{dist}(u,v) \leq 1$ )

- Procedure:

- Add a link between any two nodes u, v as  $\text{dist}(u,v) \leq 1$
- Remove some links to generate a planar graph

- Output:

- The edges before and after planarization

- The grade is proportional to **the number of remaining edges**

Input file:

```
5
0 1.5 2.3
1 1.2 2.6
2 1.3 1.2
...
```

Output file:

```
12
0 0 1
1 1 2
2 1 3
...
8
0 0 1
2 1 3
4 3 4
```

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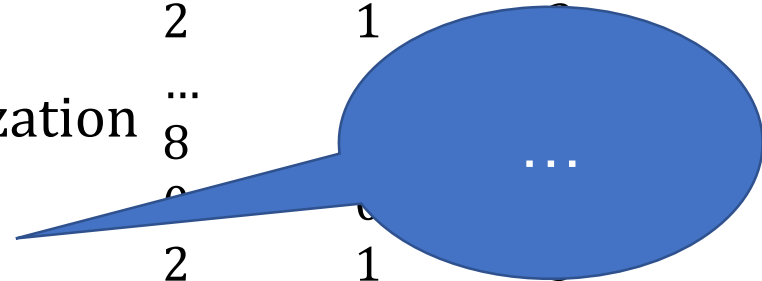
- The grade is proportional to **the number of remaining edges**

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0  1.5  2.3
1  1.2  2.6
2  1.3  1.2
...
```

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0  0  1
1  1  2
2  1
...
8
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棄選了

# Programming Project #1: Generate a Planar Subgraph

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- Procedure:

- Add a link between any two nodes u, v as  $\text{dist}(u,v) \leq 1$
- Remove some links to generate a planar graph

- Output:

- The edges before and after planarization

- Implement a **designated algorithm** to remove the links

Input file:

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0 1.5 2.3
1 1.2 2.6
2 1.3 1.2
...
```

Output file:

```
12
0 0 1
1 1 2
2 1 3
...
8
0 0 1
2 1 3
4 3 4
```

# Programming Project #1: Generate a Planar Subgraph

- Output:
  - The edges before and after planarization
- **Notice that:**
  - The first node ID should be **smaller** than the second one for each link
  - The links should be sequentially indexed in **ascending** order of the first node ID
  - If there is a tie, the two links are indexed in **ascending** order of the second node ID
  - The IDs of remaining links after planarization should be **identical** to the ones before the planarization

Input file: node.txt

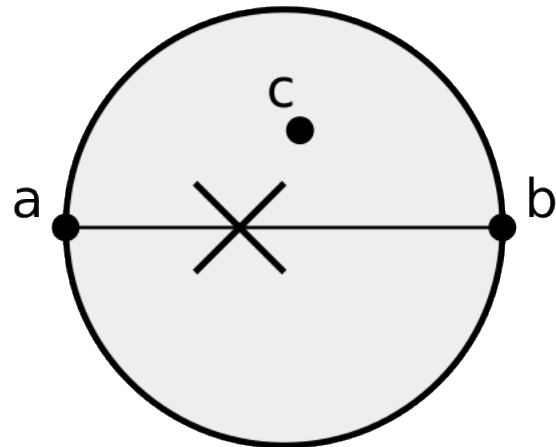
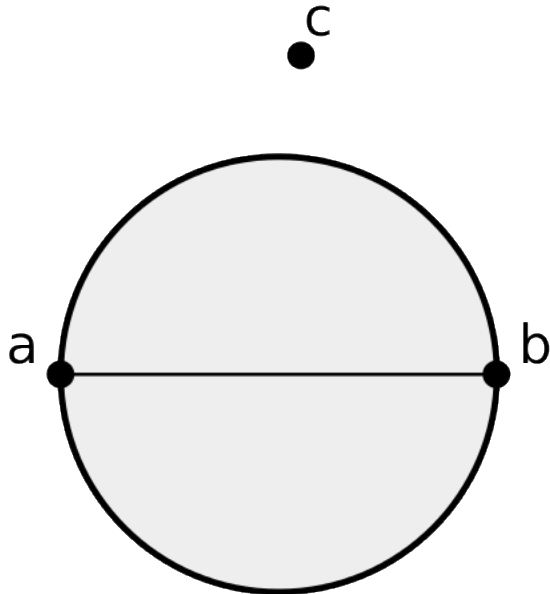
5		
0	1.5	2.3
1	1.2	2.6
2	1.3	1.2
...		

Output file: link.txt

12		
0	0	1
1	1	2
2	1	3
...		
8		
0	0	1
2	1	3
4	3	4

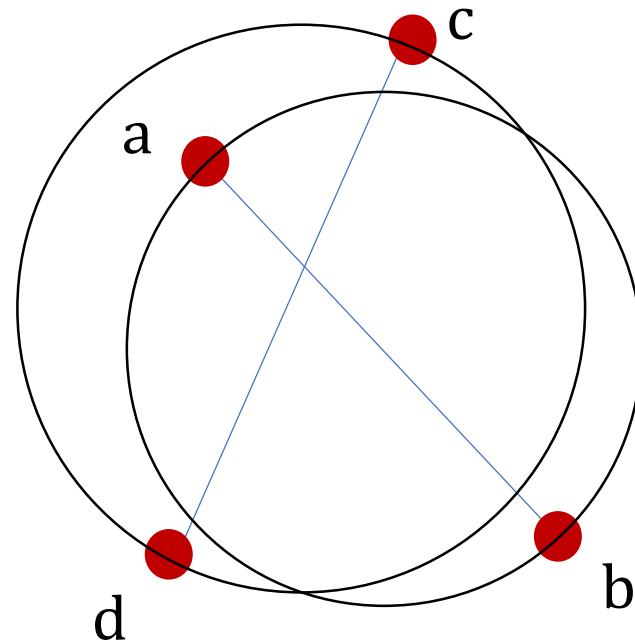
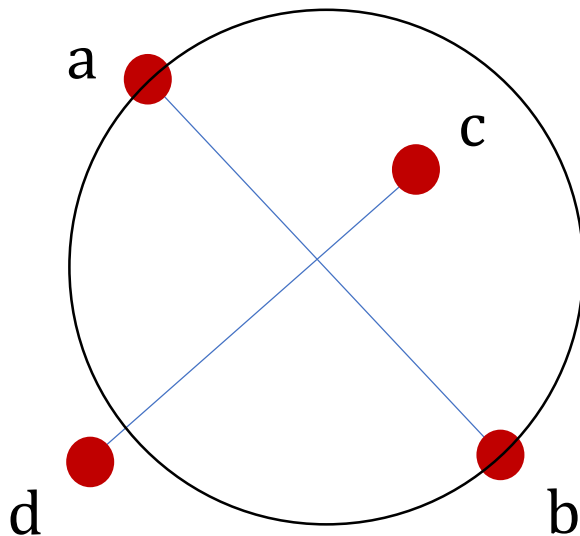
# Programming Project #1: Generate a Planar Subgraph

- Implement a **designated algorithm** to remove the links
- **Cut a link** between two neighboring nodes  $a$  and  $b$ , if there is other node  $c$  is **within their the closed disc**, where line segment  $\overline{ab}$  is a diameter



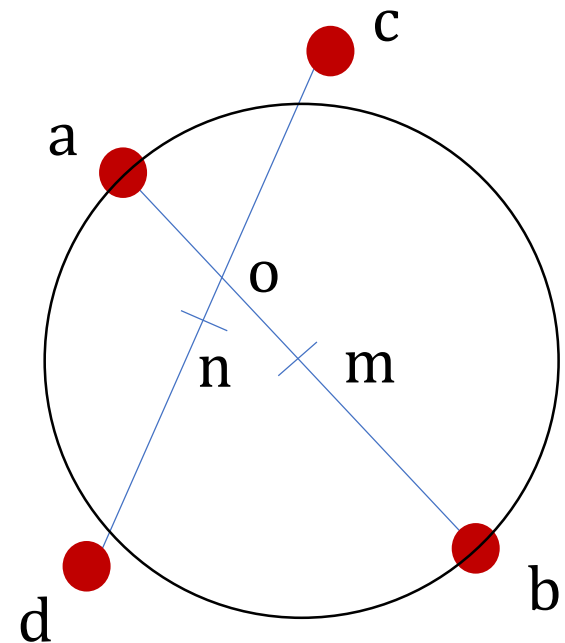
# Why does the algorithm work?

- Prove by contradiction
- Assume that there is an intersection node on the segments  $\overline{ab}$  and  $\overline{cd}$



# Why does the algorithm work?

- Let  $m$  and  $n$  denote the middle points of  $\overline{ab}$  and  $\overline{cd}$ , respectively
- Without loss of generality, we assume that  $\overline{am}$  and  $\overline{cn}$  intersect at node  $o$
- Since node  $a$  is not in the circle of which  $\overline{cd}$  is the diameter,  $\overline{cn} < \overline{an} \dots (1)$
- Similarly, since node  $c$  is not in the circle of which  $\overline{ab}$  is the diameter,  $\overline{am} < \overline{cm} \dots (2)$



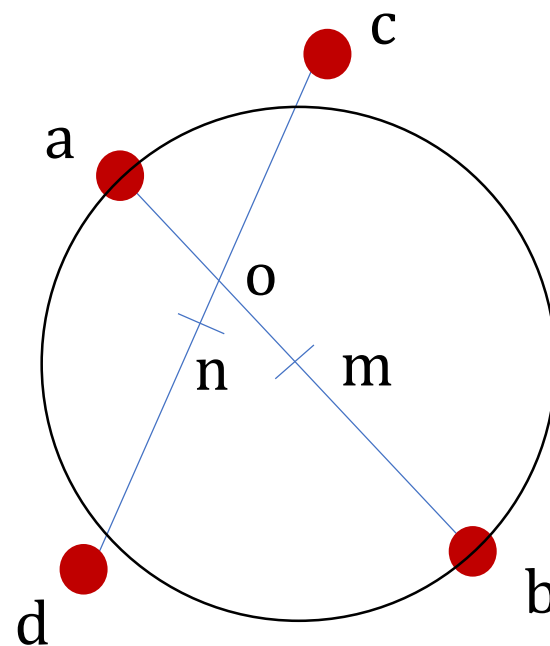


# Why does the algorithm work?

- $\overline{cn} < \overline{an} \dots (1)$
- $\overline{am} < \overline{cm} \dots (2)$
- $(1) + (2) = \overline{cn} + \overline{am} < \overline{an} + \overline{cm} \dots (3)$

- Due to triangular inequality,  
 $\overline{ao} + \overline{on} \geq \overline{an} \dots (4)$
- Similarly,  $\overline{co} + \overline{om} \geq \overline{cm} \dots (5)$
- $(4) + (5) = \overline{ao} + \overline{om} + \overline{co} + \overline{on}$   
 $= \overline{am} + \overline{cn} \geq \overline{an} + \overline{cm} \dots (6)$

- $(3) \rightarrow \leftarrow (6)$



## Input Sample: node.txt

11

0	1.4142	1.1534
1	1.97	1.85
2	0.8823	1.3926
3	1.9996	1.9484
4	0.6301	1.9145
5	0.3329	1.1756
6	1.952	0.5866
7	1.8181	0.4043
8	0.4949	1.946
9	0.2079	1.4758
10	0.1364	1.4084

# Output Sample: link.txt

22			11		
0	0	1	0	0	1
1	0	2	1	0	2
2	0	3	3	0	6
3	0	6	5	1	3
4	0	7	6	2	4
5	1	3	7	2	5
6	2	4	12	4	8
7	2	5	17	5	10
8	2	8	18	6	7
9	2	9	19	8	9
10	2	10	21	9	10
11	4	5			
12	4	8			
13	4	9			
14	4	10			
15	5	8			
16	5	9			
17	5	10			
18	6	7			
19	8	9			
20	8	10			
21	9	10			

# Note

- Deadline:  
10/11 Thu
- E-course
- C Source code
- Readme including:  
*brief description* and  
*time complexity*  
analysis of you  
implementation

