Automatic Language-Based Verification of Differential Privacy

by Chiké Abuah Goal: reveal trends in the population, without revealing information about individuals

Why Do We Need Differential Privacy?

 Ad-hoc privacy techniques (e.g. anonymization) can't guarantee privacy

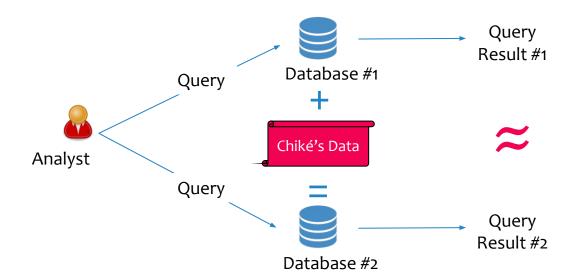


- Subject to re-identification attacks
 - Netflix recommender algorithm prize (Narayanan et al.)
 - NYC taxi data (Anthony Tockar)



Differential Privacy is a Definition of Privacy

- Adversary can't tell from analysis result whether or not Chiké participated
- Guarantee parameterized by ε (the privacy budget)



Satisfying Differential Privacy

To satisfy differential privacy, add noise to the query result

$$F(x) = f(x) +$$
Noise

How much noise? It depends on f

How many trips were taken in New York last year?

Low sensitivity

If Chiké lives in Vermont, return 100000; else return 0

High sensitivity

Why do we need Verification?

Algorithm 1 $A_{Noise-GD}$: Differentially Private Gradient Descent

Input: Data set: $\mathcal{D} = \{d_1, \dots, d_n\}$, loss function ℓ (with Lipschitz constant L), privacy parameters (ϵ, δ) , convex set \mathcal{C} , and the learning rate function $\eta \cdot [n^2] \to \mathbb{R}$.

1: Set noise variance
$$\sigma^2 \leftarrow O\left(\frac{L^2 n^2 \log(n/\delta) \log(1/\delta)}{\epsilon^2}\right)$$
.

2: $\widetilde{\theta}_1$: Choose any point from \mathcal{C} .

3: **for**
$$t = 1$$
 to $n^2 - 1$ **do**

4: Pick $d \sim_u \mathcal{D}$ with replacement.

5:
$$\widetilde{\theta}_{t+1} = \Pi_{\mathcal{C}}\left(\widetilde{\theta}_t - \eta(t) \left[n \bigtriangledown \ell(\widetilde{\theta}_t; d) + b_t \right] \right), b_t \sim \mathcal{N}\left(0, \mathbb{I}_p \sigma^2\right).$$

6: Output $\theta^{priv} = \widetilde{\theta}_{n^2}$.

How do we know it's the right amount of noise? *Manual proof*

Incorrect differentially private algorithms don't crash
They silently violate your privacy

Privacy by (programming language) Design!

- Accurate
- Accessible to non-experts
- Automatic
- Ubiquitous
- Context Matters: Support all manner of programs, systems, software, languages, modes of operation, etc.

Program Analysis Techniques

Static Analysis

- Control Flow Analysis
- Data Flow Analysis
- Abstract Interpretation
- Type Systems
- Effect Systems
- Model Checking

Dynamic Analysis

- Testing
- Monitoring (Runtime Verification)
- Program Slicing

Talk Summary: Projects

- Duet: PL, linear types for (ϵ, δ)
- DDuo: dynamic analysis system
- Solo: static analysis system

Big Picture Practical Significance

- In practice differential privacy can be difficult to analyze. This work proposes several techniques for automatic analysis and convenient implementation of differential privacy.
- Ideally, little or no domain knowledge should be required from programmers.

 Privacy violations are silent, making it an especially important field in which to apply verification techniques.

Threat Model

- Assumes an "honest but fallible" programmer.
- Does not address side-channels.
- Terminated programs can be rerun safely (while consuming the privacy budget).

DUET:

An Expressive Language for Statically Verifying Differential Privacy

(Published at OOPSLA 2019)

ACM SIGPLAN Distinguished Paper Award Winner

Duet is the first linear typed language to support verification of advanced variants of Differential Privacy

Duet automatically proves that:

```
\begin{array}{l} \text{noisy-gradient-descent}(X,y,k,\epsilon,\delta) \triangleq \\ \text{let } \theta_0 = \text{zeros } (\text{cols } X_1) \text{ in} \\ \text{loop}[\delta \ ] \ k \text{ on } \theta_0 < X_1,y > \{t,\theta \Rightarrow \\ g_p \leftarrow \text{mgauss}[\frac{1}{m},\epsilon,\delta] < X,y > \{\text{gradient } \theta \ X \ y\} \ ; \\ \text{return } \theta - g_p \ \} \end{array}
```

satisfies

$$(2\epsilon\sqrt{2k\log(1/\delta)}, k\delta + \delta)$$

differential privacy

Duet's Contributions

- Support for (ε, δ) -DP and other variants in a linear type system
- Expressive matrix types and API for machine learning
- Ability to mix privacy variants within a single program
- Advanced privacy variants are more accurate in most machine learning scenarios

Duet gets its name...

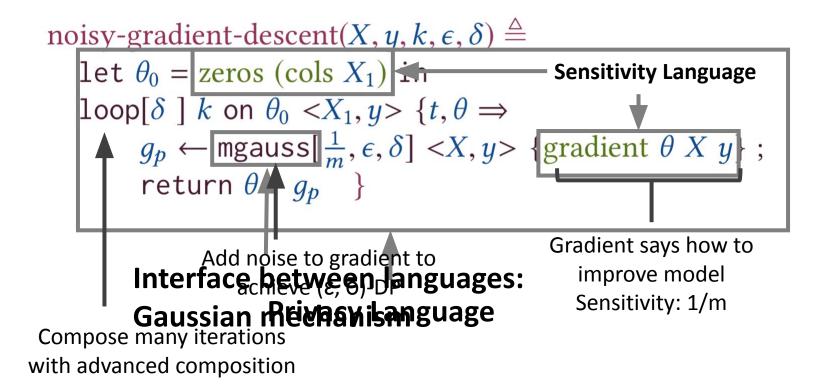
Duet is a co-design of two distinct, mutually embedded type systems:

- one for sensitivity which leverages linear typing with scaling a la Fuzz, and
- one for privacy which leverages linear typing without scaling and is novel in this work

However, to the user, it's all just one language!



Duet's Two Languages



Linear Types for Sensitivity

$$x + x$$

2-sensitive in *x*

$$\{x:_2\mathbb{R}\}\vdash x+x:\mathbb{R}$$

Context encodes sensitivity of *x*

$$\{x:_1 \tau\} \vdash x:\tau$$

$$\frac{\Gamma_1 \vdash e_1 : \mathbb{R} \qquad \Gamma_2 \vdash e_2 : \mathbb{R}}{\Gamma_1 + \Gamma_2 \vdash e_1 + e_2 : \mathbb{R}}$$

Add up sensitivities for each variable

Fuzz: Jason Reed and Benjamin C Pierce. 2010. Distance makes the types grow stronger: a calculus for differential privacy. ICFP 45, 9 (2010).

Sensitivity in Functions

$$\emptyset \vdash \lambda x : \mathbb{R} \implies x + x : \mathbb{R} \multimap_2 \mathbb{R}$$
2-sensitive function in x

Type of a 2-sensitive function

Scaling Sensitivity

$$\{y:_4\mathbb{R}\}\vdash(\lambda x:\mathbb{R}\Rightarrow x+x)\;(y+y):\mathbb{R}$$
 2-sensitive function in x

Program is 4-sensitive in y

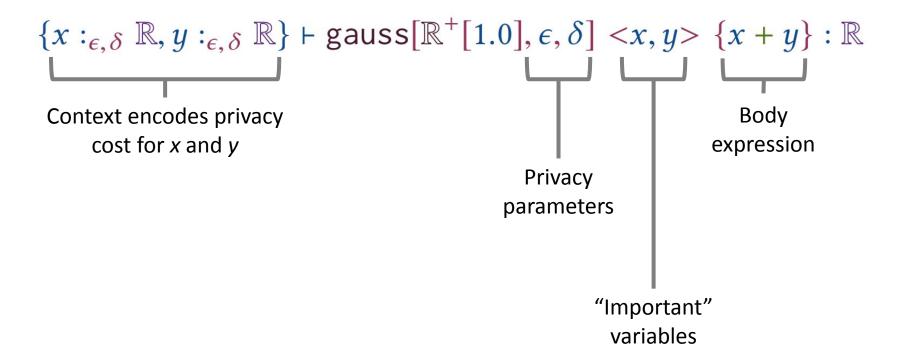
2-sensitive expression in *y*

Application rule *scales* context of the argument by the function's sensitivity

$$\frac{\neg \circ \text{-E}}{\Gamma_1 \vdash e_1 : \tau_1 \multimap_s \tau_2} \qquad \Gamma_2 \vdash e_2 : \tau_1$$

$$\frac{\Gamma_1 + s\Gamma_2 \vdash e_1 \ e_2 : \tau_2}{\Gamma_1 + s\Gamma_2 \vdash e_1 \ e_2 : \tau_2}$$

Linear Types for Privacy



Privacy Scaling is Not Allowed

$$\{y:_{2\delta} \mathcal{L}_{\delta} \mathbb{R}\} \vdash (p\lambda(x:\mathbb{R}) \Rightarrow ...) (y+y):\mathbb{R}$$

Scaling is **allowed** in *E*-differential privacy ("group privacy")

Used in Fuzz (Reed & Pierce 2010) and DFuzz (Gaboardi et al. 2013)

Scaling is **not allowed** in (ε, δ) -differential privacy

Duet's Two Languages

$$\frac{\Gamma_1 \vdash e_1 : \tau_1 \multimap_s \tau_2 \qquad \Gamma_2 \vdash e_2 : \tau_1}{\Gamma_1 + s\Gamma_2 \vdash e_1 \ e_2 : \tau_2}$$

Sensitivity language: scaling allowed

Privacy language: no scaling allowed

 $\frac{-\circ^* - E}{\Gamma \vdash e : (\tau_1 @ p_1, ..., \tau_n @ p_n) \multimap^* \tau} \qquad \frac{1}{\Gamma_1} \Gamma^1 \vdash e_1 : \tau_1 \qquad \cdots \qquad \frac{1}{\Gamma_n} \Gamma^1 \vdash e_n : \tau_n} \Gamma^1 \Gamma^n \Gamma^n \vdash e(e_1, ..., e_n) : \tau$

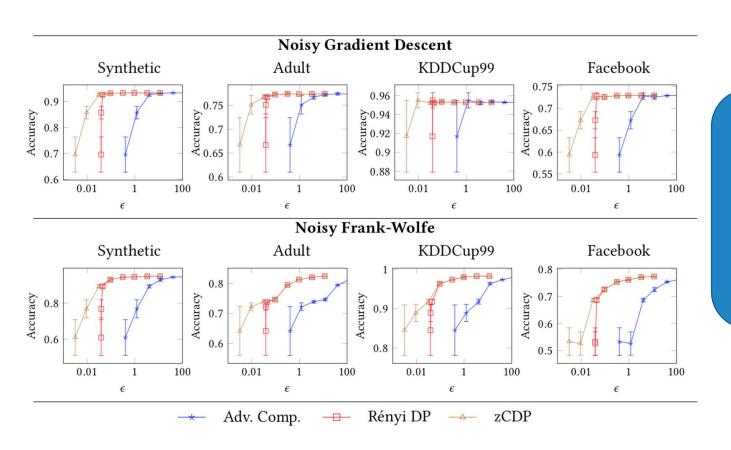
Theorem (Fundamental Property/Metric Preservation: Sensitivity)

$$\Gamma \vdash e : \tau ; \Sigma$$
 , $(\gamma_1, \gamma_2) \in \mathcal{G}_{\Sigma'} \llbracket \Gamma \rrbracket \Rightarrow (\gamma_1 \vdash e, \gamma_2 \vdash e) \in \mathcal{E}_{\Sigma' \cdot \Sigma} \llbracket \Sigma'(\tau) \rrbracket$

```
(r_{1}, r_{2}) \in \mathcal{V}_{s}[\mathbb{R}] \iff |r_{1} - r_{2}| \leq s
(\gamma_{1} \vdash e_{1}, \gamma_{2} \vdash e_{2}) \in \mathcal{E}_{s}[\tau] \iff \forall v_{1}, v_{2},
\gamma_{1} \vdash e_{1} \Downarrow v_{1} \land \gamma_{2} \vdash e_{2} \Downarrow v_{2}
\Rightarrow (v_{1}, v_{2}) \in \mathcal{V}_{s}[\tau]
(\gamma_{1}, \gamma_{2}) \in \mathcal{G}_{\Sigma}[\Gamma] \iff \forall x \in dom(\Gamma). \ (\gamma_{1}(x), \gamma_{2}(x)) \in \mathcal{V}_{\Sigma(x)}[\Gamma(x)]
```

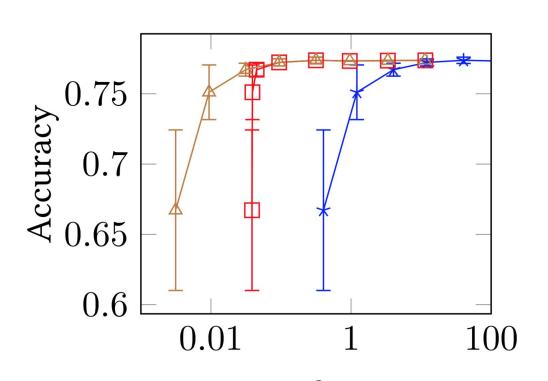
when given related initial configurations and evaluation outputs, then those outputs are related

Empirical Results (Accuracy of Trained Models)



Duet programs
produce
state-of-the-art
accuracy in
private linear
models

Benefit of Differential Privacy Variants



Recent privacy variants yield better accuracy at the same level of privacy

Adv. Comp.





Empirical Results (Typechecker Performance)

Technique	LOC	Time (ms)
Noisy G.D.	23	0.51ms
G.D. + Output Pert.	25	0.39ms
Noisy Frank-Wolfe	31	0.59ms
Minibatching	26	0.51ms
Parallel minibatching	42	0.65ms
Gradient clipping	21	0.40ms
Hyperparameter tuning	125	3.87ms
Adaptive clipping	68	1.01ms
Z-Score normalization	104	1.51ms

Typechecker computes privacy cost of complex programs in milliseconds

DDUO:

General-Purpose Dynamic Analysis for Differential Privacy

(Published at CSF 2021)

Contributions!

- includes all base types
- general language operations
- various notions of sensitivity
- advanced privacy variants
- generalizes to all future possible arbitrary inputs

DDUO gets its name from...



Product: DDUO makes it easy to automatically enforce privacy

```
def dp_gradient_descent(iterations, alpha, eps):
    eps_i = eps/iterations
    theta = np.zeros(X_train.shape[1])
    noisy_count = duet.renyi_gauss(X_train.shape[0], α = alpha, ε = eps)
    for i in range(iterations):
        grad_sum = gradient_sum(theta, X_train, y_train, sensitivity)
        noisy_grad_sum = gaussian_mech_vec(grad_sum, alpha, eps_i)
        noisy_avg_grad = noisy_grad_sum / noisy_count
        theta = np.subtract(theta, noisy_avg_grad)
    return theta
```

- Usable
- For non-experts
- Capable of complex algorithms

DDUO is prototyped in Python

- Does not require the programmer to write any additional type annotations.
- In many cases, DDUO can verify differential privacy for essentially unmodified Python programs.
- easily integrated with popular libraries like Pandas and NumPy.

```
from dduo import pandas as pd
df = pd.read_csv("data.csv")
df

# no change to sensitivity environment
df + 5

# doubles the sensitivity
df + df

( df * 5, df * df)
```

```
with dduo.Eps0dometer() as odo:
   _ = dduo.laplace(df.shape[0], ε = 1.0)
   _ = dduo.laplace(df.shape[0], ε = 1.0)
   print(odo)
```

DDUO Overview

- DDUO data sources are wrappers around sensitive data.
- DDUO tracks the sensitivity of a value to changes in the program's inputs using a sensitivity environment mapping input data sources to sensitivities.
- Sensitivity environments of Sensitive objects update as operations are applied to them (f(x) = x + x).
- DDUO tracks total privacy cost using objects called **privacy odometers**.
- DDUO also allows the analyst to place upper bounds on total privacy cost (i.e. a privacy budget) using **privacy filters**.

```
from dduo import pandas as pd
df = pd.read_csv("data.csv")
df

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   print(odo)
```

Sensitivity Analysis: Object Proxies

```
from dduo import pandas as pd
df = pd.read csv("data.csv")
df
# no change to sensitivity environment
df + 5
# doubles the sensitivity
df + df
(df * 5, df * df)
```

```
>>> Sensitive(<'DataFrame'>, data.csv → 1, L∞)
>>> Sensitive(<'DataFrame'>, data.csv → 1, L∞)
>>> Sensitive(<'DataFrame'>, data.csv → 2, L∞)
>>> (Sensitive(<'DataFrame'>, data.csv → 5,
L\infty), Sensitive(<'DataFrame'>, data.csv \Rightarrow \infty, L\infty))
```

Side-effects/Mutation

- Since DDUO attaches sensitivity environments to values (instead of variables), the use of side effects does not affect the soundness of the analysis
- When a program variable is updated to reference a new value, that value's sensitivity environment remains attached.
- more capable than traditional type-based static analysis, due to the additional challenges there (e.g. aliasing).

```
total = 0
for i in range(20):
  total = total + df.shape[0]
return total
```

```
>>> Sensitive(<'DataFrame'>,
data.csv → 20, L∞)
```

Conditionals

```
if df.shape[0] == 10:
    return df.shape[0]
else:
    return df.shape[0] * 10000

if dduo.gauss(ε=1.0, δ=1e-5, x) > 5:
    print(dduo.gauss(ε=1.0, δ=1e-5, y))
else:
    print(dduo.gauss(ε=1000000000000, δ=1e-5, y))
```

- Branching on sensitive values is disallowed
- We don't want this anyway

- Static Analysis: Take the max!
- Dynamic Analysis: ...?
- **Adaptive** privacy analysis requires use of privacy odometers/filters

Privacy Analysis: Filters and Odometers

```
dduo.laplace(df.shape[0], \(\epsilon=1.0\)
with dduo.EpsOdometer() as odo:
   _ = dduo.laplace(df.shape[0], \(\epsilon=1.0\))
   _ = dduo.laplace(df.shape[0], \(\epsilon=1.0\))
print(odo)

with dduo.EdFilter(\(\epsilon=1.0\), \(\delta=10e-6\)) as odo:
   print('1:', dduo.gauss(df.shape[0], \(\epsilon=1.0\), \(\delta=10e-6\))
print('2:', dduo.gauss(df.shape[0], \(\epsilon=1.0\), \(\delta=10e-6\))
```

```
>>> 9.963971319623278
>>> Odometer \varepsilon(data.csv \rightarrow 2.0)
>>> 1 · 10 5627
Traceback (most recent call last):
 ...
 dduo.PrivacyFilterException
```

Loops, Composition, Variants

```
with dduo.EpsOdometer() as odo:
  for i in range(20):
    dduo.laplace(df.shape[0], \varepsilon = 1.0)
  print(odo)
with dduo.AdvEdOdometer() as odo:
  for i in range(20):
    dduo.gauss(df.shape[0], \varepsilon = 0.01, \delta = 0.001)
with dduo.EdOdometer(max delta = 1e-4) as odo:
  with dduo.RenyiDP(1e-5):
    for x in range(200):
      noisy_count = dduo.renyi_gauss(\alpha = 10,
         \varepsilon=0.2, df.shape[0])
  print(odo)
```

- sequential composition

- advanced composition

variant mixing

Gradient Descent in DDUO

- Bound sensitivity of gradient calculation.
- Add random noise.
- Descend model with noisy gradient.

```
def dp_gradient_descent(iterations, alpha, eps):
    eps i = eps/iterations
    theta = np.zeros(X_train.shape[1])
    noisy_count = duet.renyi_gauss(X_train.shape[0], \alpha = alpha, \epsilon = eps)
    for i in range(iterations):
                        = gradient_sum(theta, X_train, y_train, sensitivity)
        grad sum
        noisy grad sum = gaussian mech vec(grad sum, alpha, eps i)
        noisy_avg_grad = noisy_grad_sum / noisy_count
                        = np.subtract(theta, noisy avg grad)
        theta
    return theta
```

Theorem (Metric Preservation)

```
If: \rho_1 \sim_n^{\Sigma} \rho_2

And: \sigma_1 \sim_n^{\Sigma} \sigma_2

Then: \rho_1, \sigma_1, e \sim_n^{\Sigma} \rho_2, \sigma_2, e
```

That is, either n=0 or n=n'+1 and...

If:
$$n_1 \leq n$$

And: $\rho_1 \vdash \begin{bmatrix} \sigma_1, e \end{bmatrix} \Downarrow_{n_1} \begin{bmatrix} \sigma'_1, v_1 \end{bmatrix}$

And: $\rho_2 \vdash \begin{bmatrix} \sigma_2, e \end{bmatrix} \Downarrow_{n_2} \begin{bmatrix} \sigma'_2, v_2 \end{bmatrix}$

Then: $n_1 = n_2$

And: $\sigma'_1 \sim_{n-n_1^{\Sigma}} \sigma'_2$

And: $v_1 \sim_{n-n_1^{\Sigma}} v_2$

- the true sensitivity of a program is guaranteed to be equal to or less than the sensitivity reported by DDUO's dynamic monitor.
- accurate even for inputs which differ entirely from those used in the dynamic analysis!

Case Studies: Dynamic Enforcement of Privacy

Algorithm	Libraries Used	Baseline	Instrumented Version	Overhead (% increase)
Noisy Gradient Descent	NumPy	5.922s	6.302s	6.42%
Multiplicative Weights (MWEM)	Pandas	0.725s	0.833s	14.90%
Private Naive Bayes Classification	DiffPrivLib	2.155s	2.423s	12.44%
Private Logistic Regression	DiffPrivLib	2.022s	3.161s	56.33%

Solo:

Lightweight Static Analysis for Differential Privacy

(Published at OOPSLA 2022)

Why Solo?

- + **Linear types** provide a strategy rooted in type theory and linear logic for tracking resources throughout the semantics of a core lambda calculus.
- + However, **not commonly available** in mainstream programming languages, and when available are usually **not adequately sophisticated**.

What is Solo?

- + Static type checking for differential privacy in Haskell
- + Doesn't use linear types!
- + Static analysis via (in tandem with) Haskell typechecker
- + Regular Haskell type errors when privacy is violated
- + **Privacy cost** along with types of functions & values
- + Encapsulation of sensitive values via constructor hiding
- + Automatic type inference

How does it work?

- + Without linear types, where do we attach sensitivities?
- + Previous dynamic sensitivity analyses have attached sensitivities to values.
- + We embed sensitivities in **base types**—the static equivalent of the dynamic strategy of attaching sensitivities to values.
- + **Type-level parameters** to represent sensitivities symbolically.
- + **Type-level computation** to compute symbolic sensitivity expressions.
- + Parametric polymorphism to generalize types over sensitivity parameters.
- + Possible in any language with these e.g. Scala, OCaml and Haskell.

Departure from Linear types

- + **case statements** are *restricted* to have the same sensitivity (or privacy cost) in each case alternative, rather than max. Solved by weaken operation.
- + No verified/typechecked **map** or **loop**. However trusted primitives are available.
- + No **general recursive datatypes**. However, sensitive collection data types are available.
- + In practice, **none of these are typically barriers** to writing differentially private programs.

Solo gets its name from ...



Basics: Tracking Sensitivity & Privacy

- Data source based approach
- Type-level sensitivity parameters on "sensitive" types
- Parametric effect monad: unifies privacy effect system with a monadic-style semantics

```
readDoubleFromIO :: ∀ m o. IO (SDouble m '[ '(o, 1) ])
(<+>) :: SDouble 'Diff senv1
 -> SDouble 'Diff senv2
 -> SDouble 'Diff (Plus senv1 senv2)
dbl :: SDouble 'Diff senv -> SDouble 'Diff (Plus senv senv)
dbl x = x < +> x
simplePrivacyFunction :: SDouble 'Diff'['(o, 1)]
 -> EpsPrivacyMonad '[ '(o, 2) ] Double
simplePrivacyFunction x = laplace @2 Proxy $ dbl x
```

Functions

Function types are polymorphic over sensitivity environments

```
-- An s-sensitive function
s sensitive :: SDouble senv m -> SDouble (ScaleSens senv s) m
-- A 1-sensitive function
one sensitive :: SDouble senv m -> SDouble senv m
-- тар
map :: \forall m s s<sub>1</sub> a b. (\forall s'. a s' \rightarrow b (s * s'))
  -> SList m a si
  \rightarrow SList m b (s * s<sub>1</sub>)
```

Privacy Monad

```
data EpsPrivacyCost = InfEps | EpsCost TLRat
type EpsPrivEnv = [(Source, EpsPrivacyCost)]
return :: a -> EpsPrivacyMonad '[] a
(>>=) :: EpsPrivacyMonad p1 a
  -> (a -> EpsPrivacyMonad p<sub>2</sub> b)
  -> EpsPrivacyMonad (EpsSeqComp p<sub>1</sub> p<sub>2</sub>) b
laplace :: Proxy ε
  -> SDouble s Diff
  -> EpsPrivacyMonad (TruncateSens ε s) Double
listLaplace :: Proxy ε
  -> L1List (SDouble Diff) s
  -> EpsPrivacyMonad (TruncateSens ε s) [Double]
addNoiseTwice :: TL.KnownNat (MaxSens s) =>
  SDouble s Diff -> EpsPrivacyMonad '[ '(0,5) ] Double
addNoiseTwice x = do
  y<sub>1</sub> <- laplace @2 Proxy x
  y<sub>2</sub> <- laplace @3 Proxy x
  return y_1 + y_2
```

- The return operation accepts some value and embeds it in the PrivacyMonad without causing any side-effects.
- The (>>=) (bind) operation allows us to sequence private computations using differential privacy's sequential composition property

Case study: gradient descent

```
ad :: NatS k
  -> NatS ε
  -> SMatrix σ LInf m n SDouble
  -> SMatrix o LInf m 1 SDouble
  \rightarrow EpsPrivacyMonad (ScalePriv k (TruncateSens \epsilon \sigma)) (Matrix 1 n Double)
ad k t xs vs = do
  let m_0 = \text{matrix} (\text{sn32 @ 1}) (\text{sn32 @ n}) \$ \setminus i j \rightarrow 0
       cxs = mclip xs (natS @ 1)
  let f :: SMatrix o<sub>1</sub> LInf 1 n SDouble
     -> EpsPrivacyMonad (TruncateSens \varepsilon \sigma_1) (Matrix 1 n Double)
       f = \theta \rightarrow t g = mlaplace @e Proxy (natS @5) $ xgradient $\theta$ cxs ys
       in msubM (return \theta) g
       z = mloop @(TruncateNat t 1) k (sourceM $ xbp m<sub>0</sub>) f
  Z
```

Proof: similar to DDuo but with type info

Theorem (Metric Preservation)

```
If: \gamma_1 \sim \gamma_2 \in \mathcal{G}_n^{\Sigma}[\![\Gamma]\!] And: \Gamma \vdash e : \tau@_m^{\Sigma_1}
```

Then: $\gamma_1, e \sim \gamma_2, e \in \mathcal{E}_n^{\Sigma \cdot \Sigma_1} \llbracket \tau \rrbracket$

That is, either n=0, or n=n'+1 and...

If:
$$n'' \leq n$$

And:
$$\gamma_1 \vdash e \Downarrow_{n''} v_1$$

Then:
$$\exists ! v_2. \ \gamma_2 \vdash e \Downarrow_{n''} v_2$$

And:
$$v_1 \sim v_2 \in \mathcal{V}_{n-n''}^{\Sigma \cdot \Sigma_1} \llbracket \tau \rrbracket$$

Solo: Case Studies

Solo reference implementation is a (~600 loc) Haskell library

- K-means clustering
- Cumulative Distribution Function
- Gradient Descent
- Multiplicative-Weights Exponential Mechanism (MWEM)

Related Work

Related Work: Statically Typed

Linear Type Systems

Fuzz, DFuzz (only ε-differential privacy)

Indexed Monadic Types

HOARe² (lacks multi-argument support)

Relational Type Systems

LightDP (lacks sensitivity analysis)

<u>Type Systems Enriched With Program Logics</u>

Fuzzi (less support for higher-order/type-checking automation)

Related Work: Dynamically Typed

- PINQ (lacks support for general-purpose programming)
- <u>ProPer</u> (per-user budget)
- <u>UniTrax</u> (per-user w/abstract db)
- <u>Testing Methods</u> (counter-example search)

Related Work: Other Privacy Analysis Software Libraries

- <u>DiffPrivLib</u> (no language-based sensitivity/composition)
- Google's Privacy Lib (no language-based sensitivity/composition)
- <u>Ektelo</u> (plans over library of operators, no general PL model)
- <u>DPella</u> (AST analysis, symbolic interpretation)

Future Work & Conclusions

Open Problem/Future Work

- Property-based testing, sensitivity analysis for large software libraries
- Gradual Differential Privacy

Takeaways!

- Analysis for differential privacy is important: buggy programs silently violate your privacy
- Automated enforcement of privacy can be practical in any context!

https://github.com/uvm-plaid/duet

https://github.com/uvm-plaid/dduo-python

https://github.com/uvm-plaid/solo-haskell

Papers: https://chikeabuah.github.io/



Thank You!

Questions?