

# Solutions of "counting problem"

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## 1 Solution via matrix multiplication

The problem from YouTube video "Olympiad level counting - How many subsets of  $\{1, \dots, 2000\}$  have a sum divisible by 5" by 3Blue1Brown (url: <https://www.youtube.com/watch?v=bOXCLR3Wric>) can be described via matrix multiplication.

We want to name the counts of sums congruent to  $(0, 1, \dots, d-1)$  modulo  $d$  ( $d=5$  in the video) as

$$c_0, c_1, \dots, c_{d-1}.$$

As per definition in video, an empty set has a sum of 0 and is a multiple of  $d$ .

If we know the counts  $c_0, c_1, \dots, c_{d-1}$  in a set of given numbers (set C), and we've got another set with counts  $a_0, a_1, \dots, a_{d-1}$  (set A), the total counts  $t_0, t_1, \dots, t_{d-1}$  for the union (set T) of these two sets will be (e. g.  $d=5$ ):

$$\begin{pmatrix} t_0 \\ t_1 \\ t_2 \\ t_3 \\ t_4 \end{pmatrix} = \begin{pmatrix} c_0 & c_4 & c_3 & c_2 & c_1 \\ c_1 & c_0 & c_4 & c_3 & c_2 \\ c_2 & c_1 & c_0 & c_4 & c_3 \\ c_3 & c_2 & c_1 & c_0 & c_4 \\ c_4 & c_3 & c_2 & c_1 & c_0 \end{pmatrix} \cdot \begin{pmatrix} a_0 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{pmatrix}.$$

For example, for the count of sums congruent to 2 in the new set T, we can combine:

- sums congruent to 2 from set C and sums congruent to 0 from set A,
- sums congruent to 1 from set C and sums congruent to 1 from set A,
- sums congruent to 0 from set C and sums congruent to 2 from set A,
- sums congruent to 4 from set C and sums congruent to 3 from set A, and
- sums congruent to 3 from set C and sums congruent to 4 from set A.

In total, it will be  $c_2 \cdot a_0 + c_1 \cdot a_1 + c_0 \cdot a_2 + c_4 \cdot a_3 + c_3 \cdot a_4$ , which is third row ( $t_2$ ) from the result of matrix multiplication.

### 1.1 Calculation for any number

First of all, we know that for an empty set, there is one sum ( $= 0$ ), so  $c_0$  is 1 and  $c_1 = c_2 = c_3 = c_4 = 0$ .

Then we combine this set with a set with only the number "1". We've got two possibilities here to:

- either don't use the new number (congruent 0 mod 5,  $a_0=1$ ), or
- or use the new number (congruent 1 mod 5,  $a_1=1$ ).

So for the new set as union of empty set and number "1", we've got the following new counts in a solution vector.

$$\begin{pmatrix} t_0 \\ t_1 \\ t_2 \\ t_3 \\ t_4 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix} \cdot \begin{pmatrix} 1 \\ 1 \\ 0 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \\ 0 \\ 0 \\ 0 \end{pmatrix} .$$

It works similarly when combining this set with number "2":

$$\begin{pmatrix} t_0 \\ t_1 \\ t_2 \\ t_3 \\ t_4 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 \end{pmatrix} \cdot \begin{pmatrix} 1 \\ 0 \\ 1 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \\ 1 \\ 1 \\ 0 \end{pmatrix} .$$

For sets with 3, 4, and 5 numbers:

$$n=3: \begin{pmatrix} t_0 \\ t_1 \\ t_2 \\ t_3 \\ t_4 \end{pmatrix} = \begin{pmatrix} 2 \\ 2 \\ 1 \\ 2 \\ 1 \end{pmatrix}, n=4: \begin{pmatrix} t_0 \\ t_1 \\ t_2 \\ t_3 \\ t_4 \end{pmatrix} = \begin{pmatrix} 4 \\ 3 \\ 3 \\ 3 \\ 3 \end{pmatrix}, n=5: \begin{pmatrix} t_0 \\ t_1 \\ t_2 \\ t_3 \\ t_4 \end{pmatrix} = \begin{pmatrix} 8 \\ 6 \\ 6 \\ 6 \\ 6 \end{pmatrix} .$$

For sets with numbers  $n > 5$ , we can repeat the matrix generation and matrix multiplication with the solution vector for  $n=5$ . If there is remainder for  $n$  divided by 5, we will apply another matrix multiplication with solution vector of  $n=1, \dots, 4$ .

For  $n = 2000$ , we will need  $2000 / 5 = 400$  steps.

But we can also calculate solution vectors for  $n=5^1, 5^2, 5^4$ , etc. by multiplying the generated matrix with the new solution vector coming out from the matrix multiplication.

For  $n = 2000 = 5 * 400$  and  $400 = 1\ 1001\ 0000_b$ , we need to multiply with solution vectors of  $5^{16}$ ,  $5^{128}$ , and  $5^{256}$ .

## 2 Solution via generating function

In the YouTube video "Olympiad level counting - How many subsets of  $\{1, \dots, 2000\}$  have a sum divisible by 5" by 3Blue1Brown (url: <https://www.youtube.com/watch?v=bOXCLR3Wric>), it shows a solution by using a generating function:

$$f(x) = (1+x)(1+x^2)(1+x^3)\dots(1+x^{2000})$$

$$f(x) = \sum_{n=0}^N c_n x^n = 1 + x + x^2 + 2x^3 + 2x^4 + \dots$$

With  $\zeta = e^{\frac{2\pi i}{5}}$ , the sum  $f(\zeta^0) + f(\zeta^1) + f(\zeta^2) + f(\zeta^3) + f(\zeta^4)$  results in 5 times of the sum of coefficients which is multiple of 5:

$$f(\zeta^0) + f(\zeta^1) + f(\zeta^2) + f(\zeta^3) + f(\zeta^4) = 5 \cdot (c_0 + c_5 + c_{10} + \dots)$$

For numbers n which is multiple of 5, the video gives the answer that there are  $\frac{1}{5} \left( 2^n + 4 \cdot 2^{\frac{n}{5}} \right)$  sums which are divisible by 5. E. g. for n=2000, there are  $\frac{1}{5} \left( 2^{2000} + 4 \cdot 2^{400} \right)$  of them.

The video doesn't show the case if n is not a multiple of d, but this approach is still valid then. However, the generating function f will have a remaining product  $p_r$  (r is remainder of n divided by d), e. g.  $p_3(x) = (1+x)(1+x^2)(1+x^3)$  for r is 3, and the formula will look like this:

$$\frac{1}{5} \left( 2^{n-r} \cdot p_r(\zeta^0) + 2^{\frac{n-r}{5}} \cdot [p_r(\zeta^1) + p_r(\zeta^2) + p_r(\zeta^3) + p_r(\zeta^4)] \right).$$

As  $p_r(\zeta^0) = 2^r$ , the formula becomes  $\frac{1}{5} \left( 2^n + 2^{\frac{n-r}{5}} \cdot [p_r(\zeta^1) + p_r(\zeta^2) + p_r(\zeta^3) + p_r(\zeta^4)] \right)$ .

This approach works similarly with every divisor  $d \geq 2$ .

## 2.1 Calculation for remainder of sums other than 0

I've have modified this approach by adding factors while adding the generating functions. With

$\alpha_{l,k} = \zeta^{k(d-l)}$  and k is the remainder of sums when dividing by d (e. g. 5), the sum  $\alpha_{0,k}f(\zeta^0) + \alpha_{1,k}f(\zeta^1) + \alpha_{2,k}f(\zeta^2) + \alpha_{3,k}f(\zeta^3) + \alpha_{4,k}f(\zeta^4)$  results in 5 times of the sum of coefficients with remainders is k = (0, ..., 4) when divided by 5:

$$\begin{aligned} k=0: & \alpha_{0,0}f(\zeta^0) + \alpha_{1,0}f(\zeta^1) + \alpha_{2,0}f(\zeta^2) + \alpha_{3,0}f(\zeta^3) + \alpha_{4,0}f(\zeta^4) = 5 \cdot (c_0 + c_5 + c_{10} + \dots) \\ k=1: & \alpha_{0,1}f(\zeta^0) + \alpha_{1,1}f(\zeta^1) + \alpha_{2,1}f(\zeta^2) + \alpha_{3,1}f(\zeta^3) + \alpha_{4,1}f(\zeta^4) = 5 \cdot (c_1 + c_6 + c_{11} + \dots) \\ k=2: & \alpha_{0,2}f(\zeta^0) + \alpha_{1,2}f(\zeta^1) + \alpha_{2,2}f(\zeta^2) + \alpha_{3,2}f(\zeta^3) + \alpha_{4,2}f(\zeta^4) = 5 \cdot (c_2 + c_7 + c_{12} + \dots) \\ k=3: & \alpha_{0,3}f(\zeta^0) + \alpha_{1,3}f(\zeta^1) + \alpha_{2,3}f(\zeta^2) + \alpha_{3,3}f(\zeta^3) + \alpha_{4,3}f(\zeta^4) = 5 \cdot (c_3 + c_8 + c_{13} + \dots) \\ k=4: & \alpha_{0,4}f(\zeta^0) + \alpha_{1,4}f(\zeta^1) + \alpha_{2,4}f(\zeta^2) + \alpha_{3,4}f(\zeta^3) + \alpha_{4,4}f(\zeta^4) = 5 \cdot (c_4 + c_9 + c_{14} + \dots) \end{aligned}$$

For k=0, the formula  $\frac{1}{5} \left( 2^n + 4 \cdot 2^{\frac{n}{5}} \right)$  remains the same since all factors  $\alpha_{l,0}$  equal to 1. For

k=1..4, we've got the formula  $\frac{1}{5} \left( 2^n - 2^{\frac{n}{5}} \right)$  for number of sums with remainder k when divided by 5.

These formulas applies for numbers n which is multiple of 5.

If n is not a multiple of d, the generating function f will have a remaining product  $p_r$  and the formulas must be modified accordingly (see Chapter 2). For k=1..4, we'll get formulas which are

very unhandy:  $\frac{1}{5} \left( 2^n + 2^{\frac{n-r}{5}} \cdot [\alpha_{1,k} \cdot p_r(\zeta^1) + \alpha_{2,k} \cdot p_r(\zeta^2) + \alpha_{3,k} \cdot p_r(\zeta^3) + \alpha_{4,k} \cdot p_r(\zeta^4)] \right)$ . Functions in

my program will apply matrix multiplication with solution vector of n=1, ..., d-1 instead (see Section 1.1).

This approach works similarly with every divisor  $d \geq 2$ .

## 2.2 Calculation of generating function using gcd(k, d)

Given the generating function for divisor  $d$  and numbers from 1 to  $n$  (with  $n$  is a multiple of  $d$ ):

$$f(x) = \prod_{r=1}^n (1+x^r) \quad , \text{ the video by 3Blue1Brown shows that when we plug in } \zeta^k = e^{\frac{k \cdot 2\pi i}{d}} \text{ for } x,$$

with  $k=0, \dots, d-1$ , the function can be expressed as  $f(\zeta^k) = \left( \prod_{r=1}^d (1+\zeta^{rk}) \right)^{\frac{n}{d}}$ , since  $\zeta^{(r+d)k} = \zeta^{rk}$ .

Let call the product  $g(\zeta^k) = \prod_{r=1}^d (1+\zeta^{rk})$ . The larger  $d$ , the more terms need to be multiplied in  $g(\zeta^k)$ .

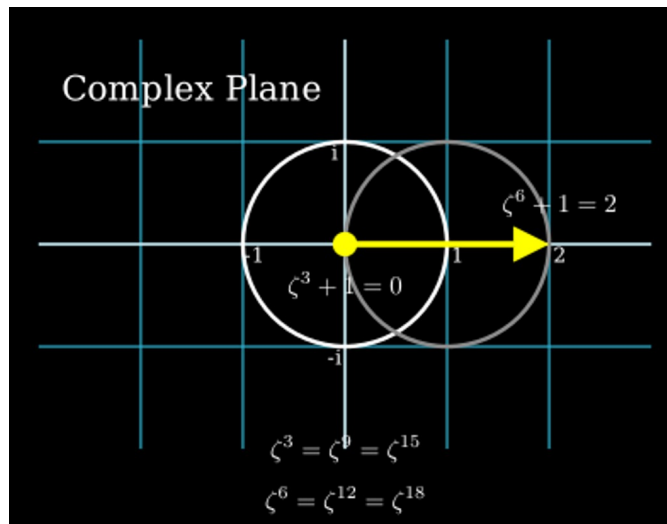
However,  $g(\zeta^k)$  can be calculated without actual multiplying the terms. Instead we can apply the following three rules.

**Rule 1:**  $k=0 \Rightarrow g(\zeta^k) = \prod_{r=1}^d (1+1) = 2^d$

**Rule 2:** If  $d$  has 2 as factor more often than  $k$  (in other words:  $\frac{d}{\gcd(k,d)}$  is even, where  $\gcd$  means the greatest common divisor), then at least one of the  $(1+\zeta^{rk})$  is 0 and  $g(\zeta^k)=0$ .

*Example 1:*  $d=6, k=3$

Since  $\gcd(3, 6) = 3$  and  $6 / 3 = 2$  is even, so are  $(1+\zeta^{rk})=0$  for  $r \in \{1, 3, 5\}$  and  $g(\zeta^k)=0$ .

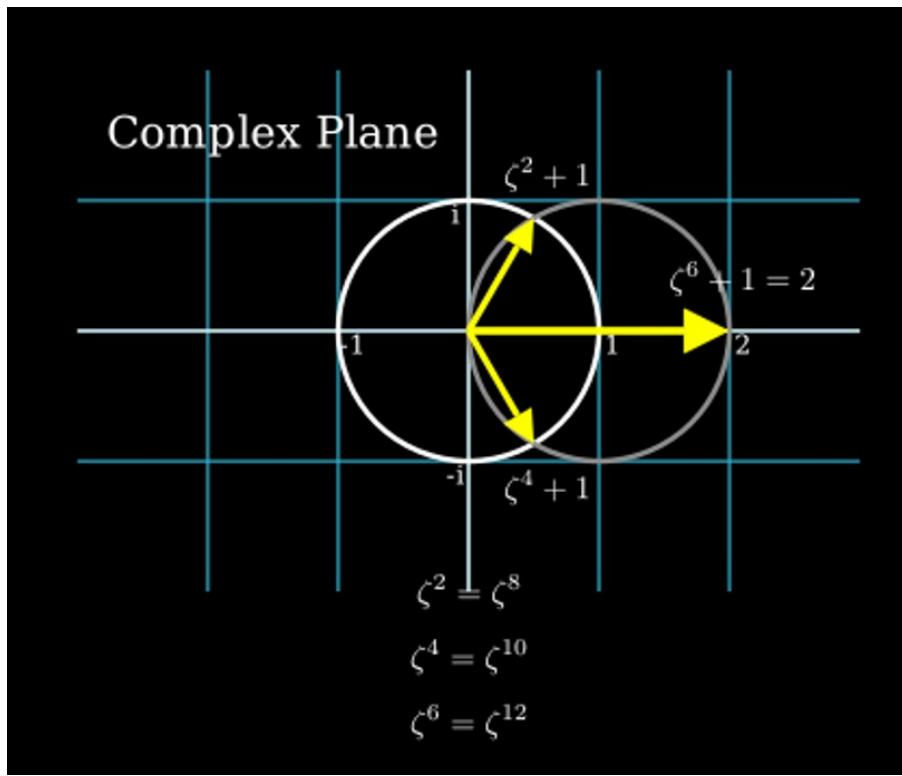


*Example 2:* Rules 2 can be applied if  $d=2^s, s \in \mathbb{Z}^+$ , since  $d/\gcd(k, d)$  is even for  $k = 1, \dots, d-1$ , and so is  $g(\zeta^k)=0$ .

**Rule 3:** In the remaining cases, the following applies:  $g(\zeta^k) = 2^{gcd(k, d)}$  .

*Example 3:*  $d = 6, k = 2$

Since  $gcd(2, 6) = 2$  and  $6 / 2 = 3$  is not even, so is  $g(\zeta^k) = 2^{gcd(2, 6)} = 2^2 = 4$  .



*Example 4:* Rule 3 can be applied if  $d$  is an odd prime number, since  $gcd(k, d)$  is 1 for  $k = 1, \dots, d-1$ . So is  $g(\zeta^k) = 2^{gcd(k, d)} = 2^1 = 2$  .