CS61B Lecture #14: Integers

Announcement:

Programming contest SATURDAY! You can still sign up.

Today: Integer Types; Readings: A Java Reference, §6.3-4. Head First Java, Chapter 10.

Readings for Upcoming Topics: Data Structures (Into Java), Chapter 1.

Integer Types and Literals

Type	Bits	Signed?	Literals
byte	8	Yes	
short	16	Yes	
char	16	No	'a' // (char) 97 '\n' // newline ((char) 10) '\t' // tab ((char) 8) '\\' // backslash 'A', '\101', '\u0041' // == (char) 65
int	32	Yes	123 0100 // Octal for 64 0x3f, 0xffffffff // Hexadecimal 63, -1 (!)
long	64	Yes	123L, 01000L, 0x3fL 1234567891011L

- "N bits" means that there are 2^N integers in the domain of the type.
- If signed, range of values is $-2^{N-1} cdots 2^{N-1} 1$.
- \bullet If unsigned, only non-negative numbers, and range is $0..2^N-1.$
- Negative numerals are just negated (positive) literals.
- Use casting for byte and short: (byte) 12, (short) 2000.

Modular Arithmetic

- Problem: How do we handle overflow, such as occurs in 10000*10000*10000?
- Some languages throw an exception (Ada), some give undefined results (C, C++)
- Java defines the result of any arithmetic operation or conversion on integer types to "wrap around"—modular arithmetic.
- That is, the "next number" after the largest in an integer type is the smallest (like "clock arithmetic").
- E.g., (byte) 128 == (byte) (127+1) == (byte) -128
- In general,
 - If the result of some arithmetic subexpression is supposed to have type T, an n-bit integer type,
 - then we compute the real (mathematical) value, x,
 - and yield a number, x', that is in the range of T, and that is equivalent to x modulo 2^n .
 - (That means that x x' is a multiple of 2^n .)

Modular Arithmetic II

- (byte) (64*8) yields 0, since $512 0 = 2 \cdot 2^8$.
- (byte) (64*2) and (byte) (127+1) yield -128, since $128 (-128) = 1 \cdot 2^8$.
- (byte) (345*6) yields 22, since $2070 22 = 8 \cdot 2^8$.
- (byte) (-30*13) yields 122, since $-390 122 = -2 \cdot 2^8$.
- (char) (-1) yields $2^{16} 1$, since $-1 (2^{16} 1) = -1 \cdot 2^{16}$.
- Natural definition for a machine that uses binary arithmetic:

Type char	Type byte
0 = 0000000000000000000000000000000000	$0 = 00000000_{2}$ $1 = 00000001_{2}$ $127 = 01111111_{2}$ $-128 = 10000000_{2}$ $-1 = 11111111_{2}$

- Terminology: rightmost (units) bit is bit 0, 2s bit is bit 1.
- ullet Hence, changing bit n modifies value by 2^n ; truncating on left to n bits computes modulo 2^n .

Negative numbers

Why this representation for -1?

$$\begin{array}{c|cccc}
 & 1 & 00000001_2 \\
+ & -1 & 11111111_2 \\
= & 0 & 1 & | 000000000_2
\end{array}$$

Only 8 bits in a byte, so bit 8 falls off, leaving 0.

- \bullet The truncated bit is in the 2^8 place, so throwing it away gives an equal number modulo 2^8 . All bits to the left of it are also divisible by 2^{8} .
- On unsigned types (char), arithmetic is the same, but we choose to represent only non-negative numbers modulo 2^{16} :

$$\begin{array}{c|ccccc}
 & 1 & 00000000000000012 \\
+ & 2^{16} - 1 & 11111111111111112 \\
= & 2^{16} + 0 & 1 | 000000000000000002
\end{array}$$

Conversion

- In general Java will silently convert from one type to another if this makes sense and no information is lost from value.
- Otherwise, cast explicitly, as in (byte) x.
- Hence, given

```
byte aByte; char aChar; short aShort; int anInt; long aLong;
// OK:
aShort = aByte; anInt = aByte; anInt = aShort; anInt = aChar;
aLong = anInt;
// Not OK, might lose information:
anInt = aLong; aByte = anInt; aChar = anInt; aShort = anInt;
aShort = aChar; aChar = aShort; aChar = aByte;
// OK by special dispensation:
aByte = 13; // 13 is compile-time constant
aByte = 12+100 // 112 is compile-time constant
```

Promotion

- Arithmetic operations (+, *, ...) promote operands as needed.
- Promotion is just implicit conversion.
- For integer operations,
 - if any operand is long, promote both to long.
 - otherwise promote both to int.
- So,

```
aByte + 3 == (int) aByte + 3 // Type int

aLong + 3 == aLong + (long) 3 // Type long

'A' + 2 == (int) 'A' + 2 // Type int

aByte = aByte + 1 // ILLEGAL (why?)
```

• But fortunately,

```
aByte += 1; // Defined as aByte = (byte) (aByte+1)
```

• Common example:

```
// Assume aChar is an upper-case letter
char lowerCaseChar = (char) ('a' + aChar - 'A'); // why cast?
```

- \bullet Java (and C, C++) allow for handling integer types as sequences of bits. No "conversion to bits" needed: they already are.
- Operations and their uses:

Mask	Set	Flip	Flip all
00101100	00101100	00101100	
& 10100111	10100111	10100111	~ 10100111
00100100	10101111	10001011	01011000

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