### Erlang

### Concurrent Programming Language and Runtime System

Lucian Pârvu

lucian@erlang.ro

June 22nd, 2012

Functional Programming Summer School Politehnica University of Bucharest Computer Science Department

# Agenda

- History
- Why Erlang was created?
- Who is using Erlang
- Sequential Erlang
- Concurrent Programming in Erlang
- OTP Open Telecom Platform
- Distributed Erlang Applications
- Debugging in Erlang
- Advantages of using Erlang
- Issues & Risks
- OpenSource Projects developed in Erlang

# Who was Erlang?

#### **Agner Krarup Erlang**

1878 - 1929

Danish mathematician, statistician and engineer, who invented the fields of traffic engineering and queueing theory.

He developed his theory of telephone traffic over several years.

His significant publications include:

1909 - "The Theory of Probabilities and Telephone Conversations" - which proves that the Poisson distribution applies to random telephone traffic.

1917 - "Solution of some Problems in the Theory of Probabilities of Significance in Automatic Telephone Exchanges" - which contains his classic formulae for loss and waiting time.



# Erlang Unit

An Erlang is a unit of telecommunications traffic measurement.

An Erlang represents the continuous use of one voice path. In practice, it is used to describe the total traffic volume of one hour.

```
Minutes of traffic in the hour = number of calls x duration = 30 \times 5 = 150
```

Hours of traffic in the hour = 150 / 60 = 2.5

Traffic figure = 2.5 Erlangs

# History

#### 1982 - 1985

Experiments with programming of telecom using > 20 different languages. Conclusion: The language must be a very high level symbolic language in order to achive productivity gains! (Leaves us with: Lisp, Prolog, Parlog...)

#### 1985 - 86

Experiments with Lisp, Prolog, Parlog etc. Conclusion: The language must contain primitives for concurrency and error recovery, and the execution model must not have back-tracking. (Rules out Lisp and Prolog.) It must also have a granularity of concurrency such that one asyncronous telephony process is represented by one process in the language. (Rules out Parlog.) We must therefore develop our own language with the desirable features of Lisp, Prolog and Parlog, but with concurrency and error recovery built into the language.

#### 1987

The first experiments with Erlang.

#### 1988

ACS/Dunder Phase I. Prototype construction of PABX functionality by external users Erlang escapes from the lab!

#### 1989

ACS/Dunder Phase 2. Reconstruction of 1/10 of the complete MD-110 system. Results: >> 10 times greater gains in efficiency at construction compared with construction in PLEX!

Further experiments with a fast implementation of Erlang.

#### 1990

Erlang is presented at ISS'90, which results in several new users, e.g Bellcore.

# History

1991 Fast implementation of Erlang is released to users. Erlang is represented at Telecom'91. More functionality such as ASN1 - Compiler, graphical interface etc.

1992 A lot of new users, e.g several RACE projects. Erlang is ported to VxWorks, PC, Macintosh etc. Three applications using Erlang are presented at ISS'92. The two first product projects using Erlang are started.

1993 Distribution is added to Erlang, which makes it possible to run a homogeneous Erlang system on a heterogeneous hardware. Decision to sell implementations Erlang externally. Separate organization in Ericsson started to maintain and support Erlang implementations and Erlang Tools.

1995 OTP Unit is formed - 60 engineers 1997 OTP Unit responsible for the distribution of Erlang

1998 (Feb) The rst demo of GPRS developed in Erlang

(Feb) Erlang was banned inside Ericsson Radio System.

(Mar) AXD301 was announced. This was possibly the largest ever program in a functional language.

(Dec) Open Source Erlang was released

(Dec) Most of the group that created Erlang resigned from Ericsson and started a new company called Bluetail AB.

2000 Bluetail was acquired by Alteon Web and six days later Alteon was acquired by Nortel Network

2002-present Joe Armstrong joins SICS

2007 "Programming Erlang - Software for a Concurrent World" - Joe Armstrong - is published

2009 Erlang Code moving to GitHub

Source: http://www.erlang.org/course/history.html

# Why Erlang was created

What if I need to build a product that needs to follow these characteristics:

- highly concurrent an distributed systems
- thousands of simultaneous transactions
- support cluster architectures and changes in the cluster topology at runtime
- support may OS's Solaris, VxWorks, Windows, Linux, Mac OS X
- support 32bit, 64bit with SMP support
- no down time (actually uptime of 99.999% which means maximum downtime of 5 minutes / year)
- highly "expressive" programming language
- recovery from software errors
- recovery from hardware errors
- trace & debug code at runtime
- update code at runtime

# Who is using Erlang?

Who uses Erlang for product development?

http://www.erlang.org/faq/introduction.html#id49610

Aptela (VoIP Service Provider)

Bluetail/Alteon/Nortel/Avaya (distributed, fault tolerant VPN Gateway)

Corelatus (SS7 monitoring).

dqdp.net (in Latvian) (Web Services).

Facebook (Facebook chat backend)

Finnish Meteorological Institute (Data acquisition and real-time monitoring)

IDT corp. (Real-time least-cost routing expert systems)

IEISS. (Electronic financial instrument exchange software)

Klarna (Electronic payment systems)

Lindenbaum (Large scale voice conferencing)

M5 Networks (VoIP Services)

Mobilearts (GSM and UMTS services)

Netkit Solutions (Network Equipment Monitoring and Operations Support Systems)

Process-one (Jabber Messaging)

# Who is using Erlang?

#### Quviq (Software Test Tool)

RabbitMQ (AMQP Enterprise Messaging)

Schlund + Partner (Messaging and Interactive Voice Response services)

Smarkets (Betting exchange and prediction market)

T-Mobile (previously one2one) (advanced call control services)

Telia (a telecomms operator)

Textendo (Innovative text messaging services)

Vail Systems (Interactive Voice Response systems)

Wavenet (SS7 and IVR applications)

#### and also:

Amazon (SimpleDB)

Basho (Riak)

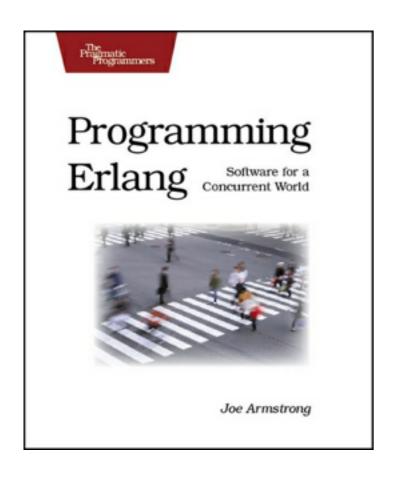
**Delicious** 

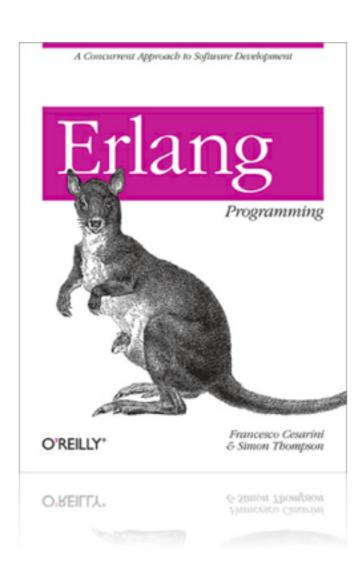
Ericsson

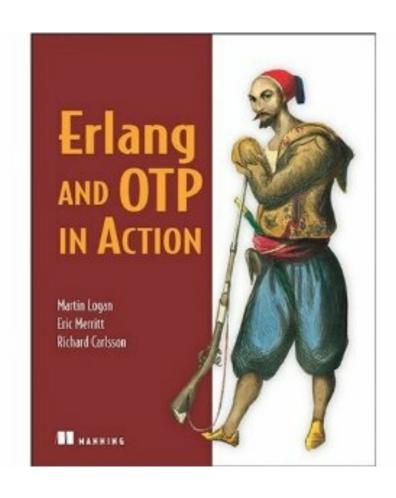
**MochiMedia** 

• • •

#### Books







### Tutorial



Learn You Some Erlang for Great Good! http://learnyousomeerlang.com/

### Basic Erlang

- The Erlang Shell
- Variables
- Data Types
- Pattern Matching
- Functions
- Modules

# The Erlang Shell

```
Last login: Wed Jun 20 16:52:16 on ttys002
dna:~ lucian$ erl
Erlang R13B02 (erts-5.7.3) [source] [64-bit] [smp:2:2] [rq:2] [async-threads:0] [kernel-poll:false]
Eshell V5.7.3 (abort with ^G)
1> pwd().
/Users/lucian
2> regs().
** Registered procs on node nonode@nohost **
                     Pid
Name
                                  Initial Call
                                                                    Reds Msgs
application controlle <0.5.0>
                                  erlang:apply/2
                                                                     444
                                                                            0
code server
                                  erlang:apply/2
                     <0.18.0>
                                                                            0
                                                                  110165
                                  erlang:apply/2
erl prim loader
                   <0.2.0>
                                                                            0
                                                                  171231
                  <0.4.0>
error logger
                                  gen event:init it/6
                                                                     220
                                                                            0
file server 2
                                  file server:init/1
                    <0.17.0>
                                                                     202
                                                                            0
global group
                     <0.16.0>
                                  global group:init/1
                                                                      60
global name server
                     <0.11.0>
                                  global:init/1
                                                                      53
                                                                            0
inet db
                     <0.15.0>
                                  inet db:init/1
                                                                    2873
                                                                            0
                                  otp_ring0:start/2
init
                     <0.0.0>
                                                                    2566
                                                                            0
kernel_safe_sup
                     <0.27.0>
                                  supervisor:kernel/1
                                                                      57
                                                                            0
kernel sup
                     <0.9.0>
                                  supervisor:kernel/1
                                                                    1169
                                                                            0
                     <0.10.0>
                                  rpc:init/1
                                                                      36
                                                                            0
rex
                                  standard error:server/2
standard error
                     <0.20.0>
                                                                            0
standard_error_sup
                                  supervisor_bridge:standar
                                                                            0
                     <0.19.0>
                                                                      40
user
                      <0.23.0>
                                  group:server/3
                                                                      38
                                  user_drv:server/2
user_drv
                     <0.22.0>
                                                                     574
                                                                            0
** Registered ports on node nonode@nohost **
Name
                      Id
                                     Command
ok
```

#### Variables

- All variable names must start with an uppercase letter.
- Erlang has single assignment variables can be given a value only once. If you
  try to change the value of a variable once it has been set, then you'll get an
  error.
- A variable that had a value assigned to it is called a bound variable;
   otherwise, it is called an unbound variable.

```
1> X = 1. 12345
2> X*2. 2
3> X = 2.
** exception error: no match of right hand side value 2
```

# Numbers, Atoms, Tuples

Floating-Point Numbers - must have a decimal point followed by at least one decimal digit. When you divide two integers with "/", the result is automatically converted to a floating-point number.

```
1> 5/3.
1.66667
2> 4/2.
2.00000
```

Atoms are global. Atoms start with lowercase letters, followed by a sequence of alphanumeric characters or the underscore () or at () sign.

```
router, device, dna@itcnetworks, nmos_domain
```

#### Tuples:

```
1> Workstation = {workstation, {dns, dna}, {domain, erlang}, {uptime, 200}, {os, macosx}}.
```

### Pattern Matching

#### Functions

```
%% Code from
    Erlang Programming
    Francecso Cesarini and Simon Thompson
    O'Reilly, 2008
    http://oreilly.com/catalog/9780596518189/
    http://www.erlangprogramming.org/
    (c) Francesco Cesarini and Simon Thompson
%% Chapter 2
-module(chapter2).
-export([area/1, area2/1, flatten/1, factorial/1]).
-import(math, [sqrt/1]).
% To calculate the area of a shape (section 2.12)
area({square, Side}) ->
  Side * Side :
area({circle, Radius}) ->
 math:pi() * Radius * Radius;
area({triangle, A, B, C}) ->
 S = (A + B + C)/2,
 math:sqrt(S*(S-A)*(S-B)*(S-C));
area(Other) ->
 {error, invalid object}.
```

```
% Variant of area using the sqrt/1 imported from the math
module.
area2({square, Side}) ->
  Side * Side ;
area2({circle, Radius}) ->
  math:pi() * Radius * Radius;
area2({triangle, A, B, C}) ->
  S = (A + B + C)/2,
  sqrt(S*(S-A)*(S-B)*(S-C));
area2(Other) ->
  {error, invalid_object}.
% To flatten a 3D object into a 2D object.
flatten(Other) -> {error, unknown shape};
flatten(cube)
                -> square;
flatten(sphere) -> circle.
% The factorial function.
factorial(0) -> 1;
factorial(N) ->
  N * factorial(N-1).
```

#### Modules

```
%% Code from
%% Erlang Programming
%% Francecso Cesarini and Simon Thompson
    O'Reilly, 2008
    http://oreilly.com/catalog/9780596518189/
    http://www.erlangprogramming.org/
%% Demo module from Chapter 2
-module(demo).
-export([double/1]).
% This is a comment.
% Everything on a line after % is ignored.
double(Value) ->
  times(Value, 2).
times(X,Y) ->
  X*Y.
%% Note: times/2 is deliberately omitted from the export list.
```

# Sequential Erlang

- Conditional Evaluations
- Guards
- Built-in Functions
- Recursion
- Handling Errors

#### Conditional Evaluations

```
%% Conditional Evaluations
                                                                    % Defining an index function using nested case expressions.
% To convert an atom representing a day into an integer.
                                                                    index3(X,Y) \rightarrow
convert(Day) ->
                                                                      case X of
  case Day of
                                                                        0 ->
    monday
            -> 1;
                                                                          case Y of
    tuesday -> 2;
                                                                           [Z|_] -> Z
    wednesday -> 3;
                                                                           end:
   thursday -> 4;
                                                                        N when N>0 ->
    friday -> 5;
                                                                           case Y of
    saturday -> 6;
                                                                            [Zs] \rightarrow index3(N-1,Zs)
    sunday -> 7;
                                                                           end
              -> {error, unknown day}
    Other
                                                                      end.
  end.
                                                                    % Examples to illustrate the scope of variables.
% Calculating the length of a list.
                                                                    f(X)
                                                                               -> Y=X+1,Y*X.
listlen([]) -> 0;
listlen([ |Xs]) -> 1 + listlen(Xs).
                                                                    g([0|Xs]) \rightarrow g(Xs);
                                                                    g([Y|Xs]) \rightarrow Y+g(Xs);
% Rewrite this directly using a case expression:
                                                                    g([])
                                                                           -> 0.
listlen2(Y) ->
                                                                    % Binding a variable in both arms of a case expression is
 case Y of
                                                                    % possible ...
   [] -> 0;
   [ | Xs ] \rightarrow 1 + listlen2(Xs)
                                                                    safe(X) ->
  end.
                                                                      case X of
                                                                        one -> Y = 12;
% Indexing into a list, i.e. looking for the nth element of a
                                                                        _ -> Y = 196
% list (with numbering from zero).
                                                                      end,
                                                                      X+Y.
index(0,[X| ])
                -> X;
index(N, [Xs]) when N>0 -> index(N-1, Xs).
                                                                    % ... but the preferred style is to assign a value to the variable.
                                                                    % where the value is defined using a case
% Defining an index function using a case expression.
                                                                    preferred(X) ->
index2(X,Y) ->
                                                                      Y = case X of
   index({X,Y}).
                                                                            one -> 12;
                                                                               -> 196
index(Z) ->
                                                                           end,
 case Z of
                                                                      X+Y.
   \{0, [X|_]\}
                         -> X;
   \{N,[\_|Xs]\}\ when N>0 \rightarrow index2(N-1,Xs)
```

end.

#### Guards

```
%% Guards
% Defining factorial using a guard
factorial(N) when N > 0 ->
  N * factorial(N - 1);
factorial(0) -> 1.
% An example of a complex guard ...
guard(X,Y) when not(((X>Y) or not(is_atom(X))) and (is_atom(Y) or (X==3.4))) ->
 X+Y.
% ... and the same guard rewritten using ; and ,
guard2(X,Y) when not(X>Y), is_atom(X); not(is_atom(Y)), X=/=3.4 ->
  X+Y.
% Examples from the examples.erl module.
even(Int) when Int rem 2 == 0 -> true;
even(Int) when Int rem 2 == 1 -> false.
number(Num) when is_integer(Num) -> integer;
number(Num) when is_float(Num) -> float;
number(_Other)
                                 -> false.
```

#### **Built-in Functions**

- Object Access and Examination
- Type Conversion
- Process Dictionary
- Meta Programming
- Process, Port, Distribution, and System Information
- Input and Output

#### Recursion

```
%% Recursion
% The bump function: add one to each element of a list.
% A first example of recursion.
bump([]) -> [];
bump([Head | Tail]) -> [Head + 1 | bump(Tail)].
% Finding the average value in a numeric list.
average(List) -> sum(List) / len(List).
sum([]) -> 0;
sum([Head | Tail]) -> Head + sum(Tail).
len([]) -> 0;
len([ | Tail]) -> 1 + len(Tail).
% Is the first argument a memeer of the second argument (a list)?
member(_, [])
                 -> false;
member(H, [H | _]) -> true;
member(H, [ | T]) -> member(H, T).
% Summing a list using tail recursion.
sum acc([],Sum) -> Sum;
sum acc([Head|Tail], Sum) -> sum acc(Tail, Head+Sum).
sum2(List) -> sum acc(List,0).
% Bumping every element in a list using an accumulator.
bump2(List) -> bump_acc(List, []).
bump acc([], Acc)
                            -> reverse(Acc);
bump acc([Head | Tail], Acc) -> bump_acc(Tail, [Head + 1 | Acc]).
```

```
% Reversing a list.
reverse(List) -> reverse_acc(List, []).
reverse_acc([], Acc) -> Acc;
reverse acc([H | T], Acc) -> reverse acc(T, [H | Acc]).
% Merging the elements of two lists.
merge(Xs,Ys) ->
lists:reverse(mergeL(Xs,Ys,[])).
mergeL([X|Xs],Ys,Zs) ->
  mergeR(Xs,Ys,[X|Zs]);
mergeL([],[],Zs) ->
  Zs.
mergeR(Xs,[Y|Ys],Zs) ->
  mergeL(Xs,Ys,[Y|Zs]);
mergeR([],[],Zs) ->
  Zs.
% Average revisited, this time using two accumulators.
average2(List) -> average acc(List, 0,0).
average acc([], Sum, Length) ->
  Sum / Length;
average acc([H | T], Sum, Length) ->
  average acc(T, Sum + H, Length + 1).
% Iterative version of sum
sum3(Boundary) -> sum acc(1, Boundary, 0).
sum acc(Index, Boundary, Sum) when Index =< Boundary ->
  sum_acc(Index + 1, Boundary, Sum + Index);
sum_acc(_I, _B, Sum)->
   Sum.
```

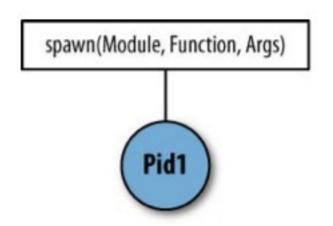
# Handling Errors

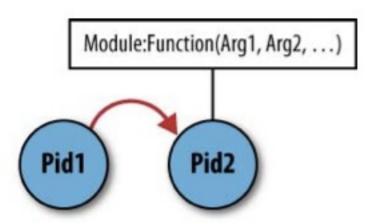
```
try_return(X) when is_integer(X) ->
try Exprs of
                                                                try return_error(X) of
    Pattern1 [when Guard1] ->
                                                                  Val -> {normal, Val}
         ExpressionBody1;
                                                                catch
    Pattern2 [when Guard2] ->
                                                                  exit:Reason -> {exit, Reason};
         ExpressionBody2
                                                                  throw:Throw -> {throw, Throw};
catch
                                                                  error:Error -> {error, Error}
    [Class1:]ExceptionPattern1 [when ExceptionGuardSeq1] ->
                                                                end.
         ExceptionBody1;
    [ClassN:]ExceptionPatternN [when ExceptionGuardSeqN] ->
                                                              try_wildcard(X) when is_integer(X) ->
         ExceptionBodyN
                                                                try return error(X)
end
                                                                  catch
                                                                  throw:Throw -> {throw, Throw};
                                                                  error: -> error;
                                                                  Type:Error -> {Type, Error};
                                                                                                        %% Will never be returned
                                                                              -> other;
                                                                                                        %% Will never be returned
                                                                              -> other
                                                                end.
                                                              try return2(X) when is integer(X) ->
                                                                try return_error(X) of
                                                                  Val -> {normal, Val}
                                                                  catch
                                                                  exit:_ -> 34;
                                                                  throw: -> 99;
                                                                  error: -> 678
                                                                end.
                                                              return(X) when is integer(X) ->
                                                                catch return_error(X).
```

# Concurrent Programming

- Creating Processes
- Message Passing
- Receiving Messages
- Registered Processes

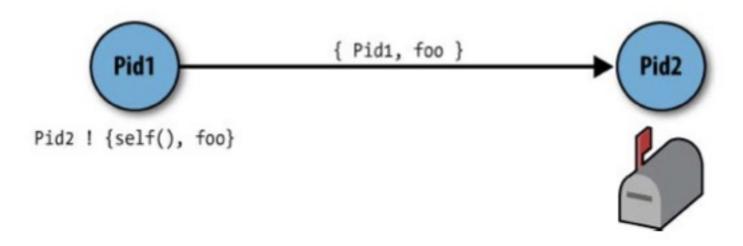
### Creating Processes



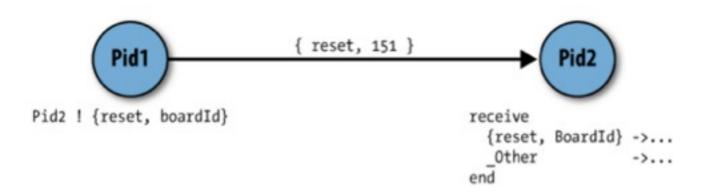


Pid2 = spawn(Module, Function, Arguments).

# Message Passing

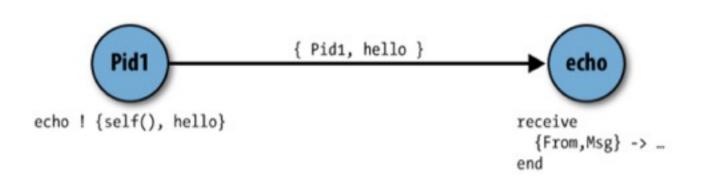


# Receiving Messages



```
%% Code from
     Erlang Programming
%% Francecso Cesarini and Simon Thompson
%% O'Reilly, 2008
%% http://oreilly.com/catalog/9780596518189/
%% http://www.erlangprogramming.org/
%% (c) Francesco Cesarini and Simon Thompson
-module(echo).
-export([go/0, loop/0]).
go() ->
   Pid = spawn(echo, loop, []),
   Pid ! {self(), hello},
   receive
     {Pid, Msg} ->
       io:format("~w~n",[Msg])
   end,
   Pid! stop.
loop() ->
   receive
     {From, Msg} ->
        From ! {self(), Msg},
        loop();
     stop ->
       true
   end.
```

### Registered Processes

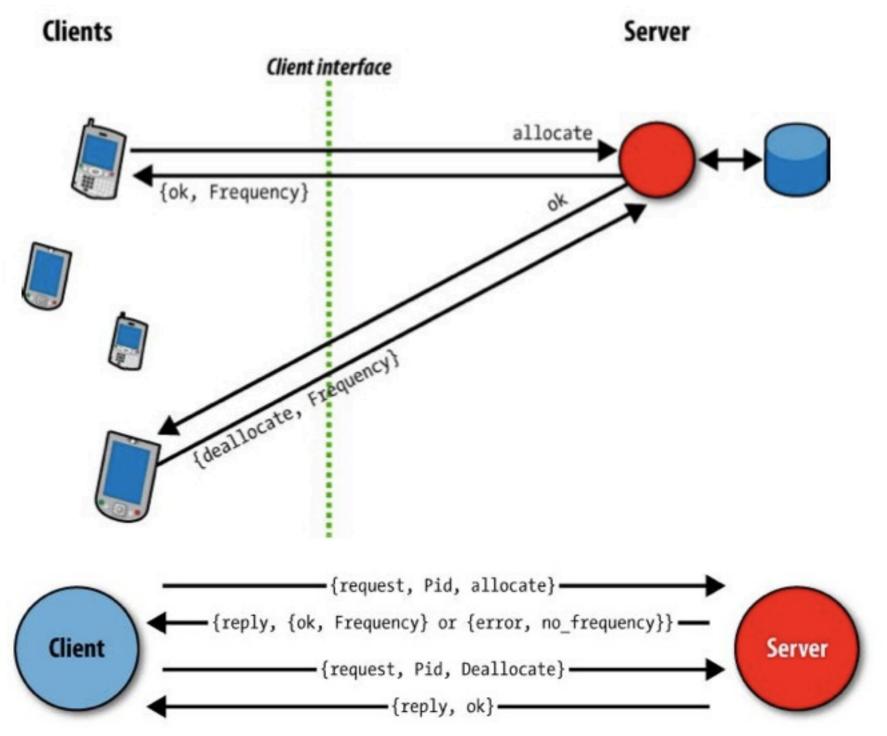


```
%% Code from
    Erlang Programming
    Francecso Cesarini and Simon Thompson
    O'Reilly, 2008
    http://oreilly.com/catalog/9780596518189/
    http://www.erlangprogramming.org/
    (c) Francesco Cesarini and Simon Thompson
-module(echo2).
-export([go/0, loop/0]).
go() ->
 register(echo, spawn(echo2, loop, [])),
 echo ! {self(), hello},
 receive
   { Pid, Msg} ->
     io:format("~w~n",[Msg])
 end.
loop() ->
 receive
   {From, Msg} ->
     From ! {self(), Msg},
     loop();
   stop ->
     true
 end.
```

# Process Design Patterns

- Client/Server Models
- Finite State Machines
- Event Managers and Handlers

#### Client/Server Models



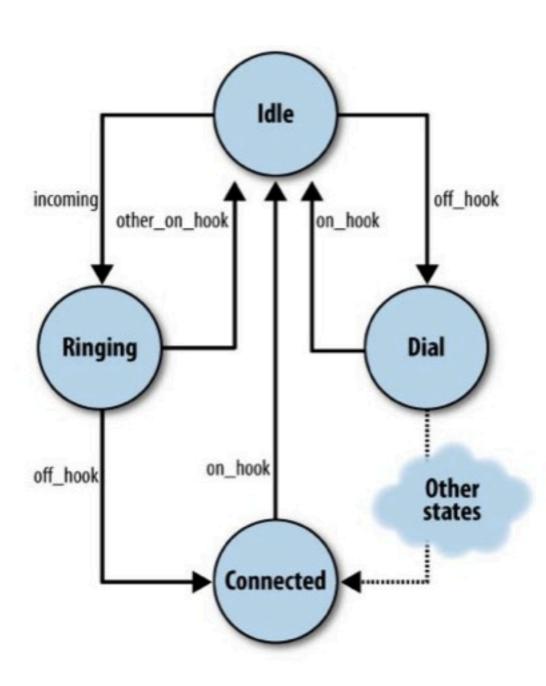
#### A Process Pattern

```
%% Code from
%% Erlang Programming
%% Francecso Cesarini and Simon Thompson
%% O'Reilly, 2008
%% http://oreilly.com/catalog/9780596518189/
%% http://www.erlangprogramming.org/
   (c) Francesco Cesarini and Simon Thompson
-module(server).
-export([start/2, stop/1, call/2]).
-export([init/1]).
start(Name, Data) ->
  Pid = spawn(generic handler, init,[Data]),
  register(Name, Pid), ok.
stop(Name) ->
 Name ! {stop, self()},
  receive {reply, Reply} -> Reply end.
call(Name, Msg) ->
  Name ! {request, self(), Msg},
  receive {reply, Reply} -> Reply end.
reply(To, Msg) ->
  To ! {reply, Msg}.
init(Data) ->
  loop(initialize(Data)).
```

```
loop(State) ->
    receive
    {request, From, Msg} ->
        {Reply,NewState} = handle_msg(Msg, State),
        reply(From, Reply),
        loop(NewState);
    {stop, From} ->
        reply(From, terminate(State))
    end.

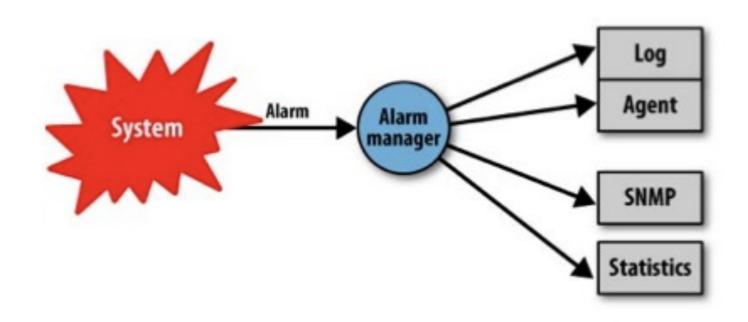
%% initialize(...) -> ...
%% handle_msg(...,...) -> ...
%% terminate(...) -> ...
```

#### Finite State Machines



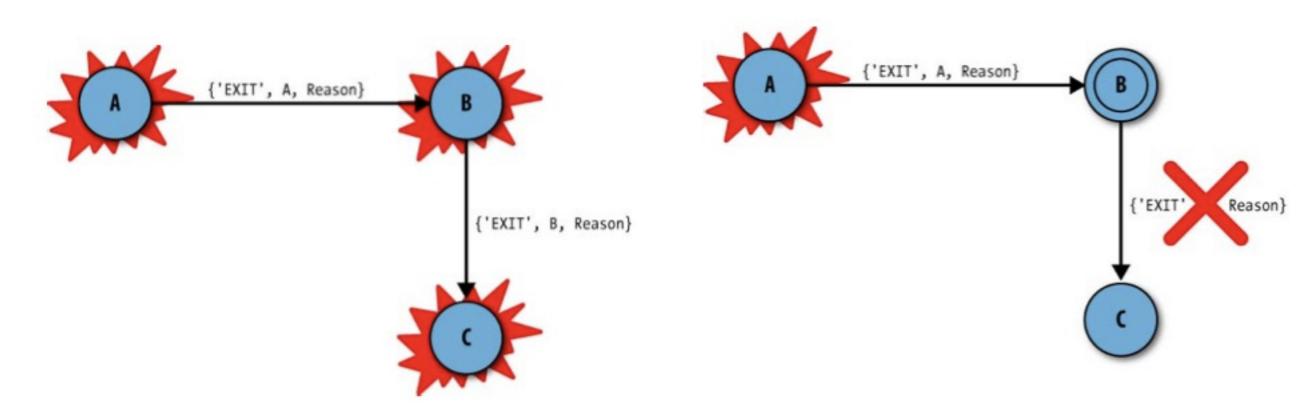
```
%% Code from
     Erlang Programming
%% Francecso Cesarini and Simon Thompson
%% O'Reilly, 2008
    http://oreilly.com/catalog/9780596518189/
    http://www.erlangprogramming.org/
    (c) Francesco Cesarini and Simon Thompson
-module(fsm).
-export([idle/0,ringing/1]).
idle() ->
  receive
    {Number, incoming} ->
      start_ringing(),
      ringing(Number);
    off hook ->
      start_tone(),
      dial()
  end.
ringing(Number) ->
  receive
    {Number, other_on_hook} ->
      stop_ringing(),
      idle();
    {Number, off_hook} ->
      stop_ringing(),
      connected(Number)
  end.
```

# Event Managers and Handlers



### Process Error Handling

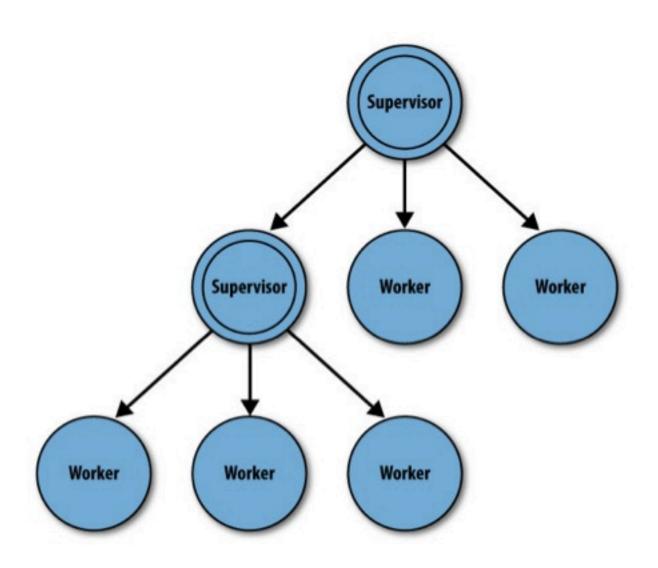
#### Process Links and Exit Signals



Propagation of exit signals

Trapping Exits

# Robust Systems



## Records & Macros

```
Simple Macros
-define(TIMEOUT, 1000).

receive
    after ?TIMEOUT -> ok
end.

Parameterized Macros
-define(DBG(Str, Args), io:format(Str, Args)).
```

# Live Code Upgrade

The ability to load new and updated modules during runtime allows systems to run without interruption, not only when errors are being fixed but also when new func- tionality is added.

It also reduces the turnaround time of bugs and facilitates testing, as in most cases, systems do not have to be restarted for patches to be validated and deployed.

The software upgrade mechanism relies on a set of simple but powerful constructs on which more powerful tools are built.

These upgrade tools are used in pretty much every Erlang-based system where downtime has to be reduced to a minimum.

# Funs and Higher-Order Functions

```
1> Bump = fun(Int) -> Int + 1 end.
#Fun<erl_eval.6.13229925>
2> Bump(10). 11
3> (fun(Int) -> Int + 1 end)(9).
10
```

#### **Functions as Results**

```
times(X) ->
    fun (Y) -> X*Y end.

This is a function that takes one argument: X,

times(X) -> ...

and whose result is this expression:

fun (Y) -> X*Y end.
```

### **Functions as Arguments**

```
doubleAll([]) ->
    [];
doubleAll([X|Xs]) ->
    [X*2 | doubleAll(Xs)].
```

### **Predefined Higher-Order Functions**

```
all(Predicate, List)
any(Predicate, List)
dropwhile(Predicate, List)
filter(Predicate, List)
foldl(Fun, Accumulator, List)
map(Fun, List)
partition(Predicate, List)
```

# List Comprehensions

### **General List Comprehensions**

[Expression || Generators, Guards, Generators, ...]

#### **Generators**

A generator has the form Pattern <- List, where Pattern is a pattern that is matched with elements from the List expression. You can read the symbol <- as "comes from"; it's also like the mathematical symbol  $\in$ , meaning "is an element of."

#### **Guards**

Guards are just like guards in function definitions, giving a true or false result. The variables in the guards are those which appear in generators to the left of the guard (and any other variables defined at the outer level).

### **Expression**

The expression specifies what the elements of the result will look like.

```
1> [X || X <- [1,2,3], X rem 2 == 0].
[2]

2> Database = [ {francesco, harryPotter}, {simon, jamesBond}, {marcus, jamesBond}, {francesco, daVinciCode} ].
...
3> [Person || {Person,_} <- Database].
[francesco, simon, marcus, francesco]

4> [ {X,Y} || X <- lists:seq(1,3), Y <- lists:seq(X,3) ].
[{1,1},{1,2},{1,3},{2,2},{2,3},{3,3}]</pre>
```

## ETS and DETS Tables

To handle fast searches, Erlang provides two mechanisms. This chapter introduces Erlang Term Storage (ETS) and Disk Erlang Term Storage (Dets), two mechanisms for memory- and disk-efficient storage and retrieval of large collections of data.

### Set

In a set, each key can occur only once. So, using this kind of table for the index example will mean there can be only one element in the table for each word.

### **Ordered** set

An ordered set has the same property as the set, but it is stored so that the elements can be traversed following the lexicographical order on the keys.

### Bag

A bag allows multiple entries for the same key. The elements have to be distinct: in the index example, this means there can be only one entry for a particular word on a particular line.

### **Duplicate bag**

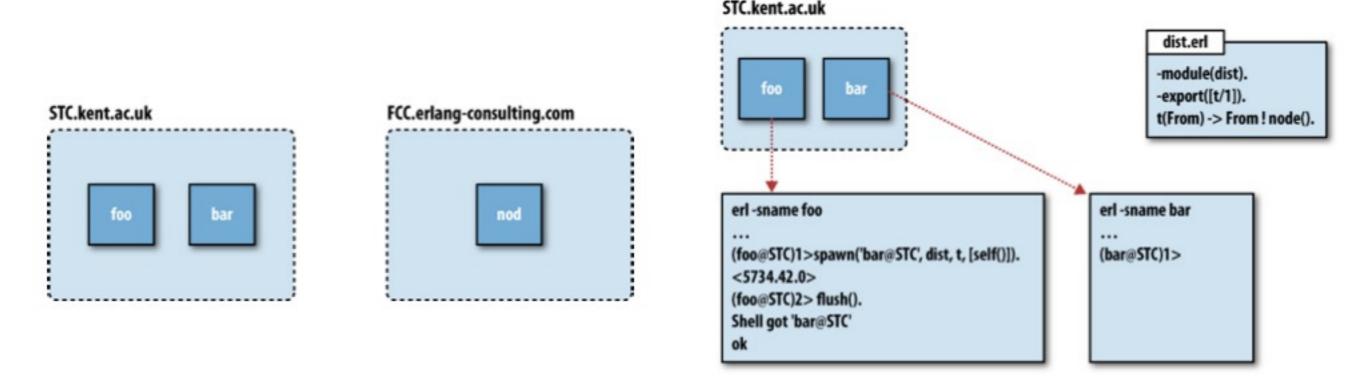
A duplicate bag allows duplicated elements, as well as duplicated keys.

# Distributed Programming

A distributed Erlang system consists of a number of Erlang runtime systems communicating with each other. Each such runtime system is called a node.

Message passing between processes at different nodes, as well as links and monitors, are transparent when pids are used.

Registered names, however, are local to each node. This means the node must be specified as well when sending messages etc. using registered names.



## **OTP Behaviours**

The standard Erlang/OTP behaviours are:

### gen\_server

For implementing the server of a client-server relation.

### gen\_fsm

For implementing finite state machines.

### gen\_event

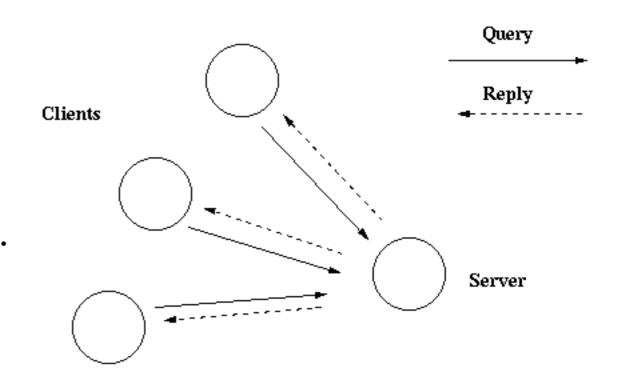
For implementing event handling functionality.

### supervisor

For implementing a supervisor in a supervision tree.

## Generic Servers

```
-module(ch3).
-behaviour(gen_server).
-export([start_link/0]).
-export([alloc/0, free/1]).
-export([init/1, handle_call/3, handle_cast/2]).
start link() ->
    gen_server:start_link({local, ch3}, ch3, [], []).
alloc() ->
    gen_server:call(ch3, alloc).
free(Ch) ->
    gen_server:cast(ch3, {free, Ch}).
init(_Args) ->
    {ok, channels()}.
handle_call(alloc, _From, Chs) ->
    {Ch, Chs2} = alloc(Chs),
    {reply, Ch, Chs2}.
handle_cast({free, Ch}, Chs) ->
    Chs2 = free(Ch, Chs),
    {noreply, Chs2}.
```

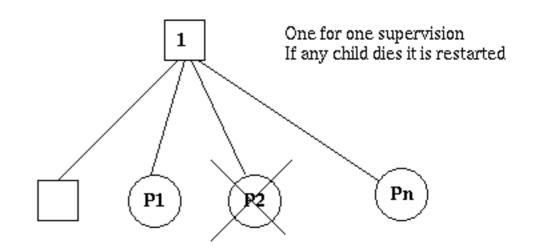


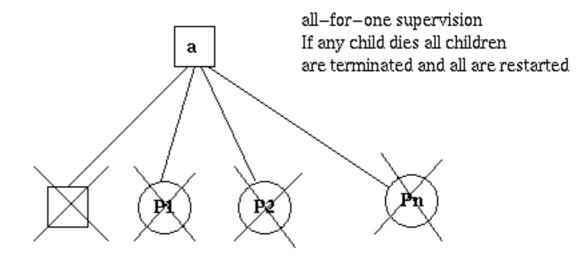
The Client-server model

# Supervisors

```
-module(ch_sup).
-behaviour(supervisor).
-export([start link/0]).
-export([init/1]).
start_link() ->
    supervisor:start_link(ch_sup, []).
init( Args) ->
    {ok, {{one_for_one, 1, 60}},
          [{ch3, {ch3, start_link, []},
            permanent, brutal_kill, worker, [ch3]}]}}.
one_for_one is the restart strategy.
1 and 60 defines the maximum restart
frequency.
The tuple {ch3, ...} is a child
```

specification.





# (a)mnesia

Mnesia is a distributed DataBase Management System (DBMS), appropriate for telecommunications applications and other Erlang applications which require continuous operation and exhibit soft real-time properties.

Listed below are some of the most important and attractive capabilities, Mnesia provides:

- A relational/object hybrid data model which is suitable for telecommunications applications.
- A specifically designed DBMS query language, QLC (as an add-on library).
- Persistence. Tables may be coherently kept on disc as well as in main memory.
- Replication. Tables may be replicated at several nodes.
- Atomic transactions. A series of table manipulation operations can be grouped into a single atomic transaction.
- Location transparency. Programs can be written without knowledge of the actual location of data.
- Extremely fast real time data searches.
- Schema manipulation routines. It is possible to reconfigure the DBMS at runtime without stopping the system.

# Socket Programming

```
%% File : tcp.erl
%%% Description : Example from Chapter 15,
-module(tcp).
-export([client/2, send/2, server/0, wait connect/2,
      get request/3, handle/2]).
client(Host, Data) ->
    {ok, Socket} = gen_tcp:connect(Host, 1234, [binary, {packet, 0}]),
    send(Socket, Data),
    ok = gen tcp:close(Socket).
send(Socket, <<Chunk:100/binary, Rest/binary>>) ->
    gen_tcp:send(Socket, Chunk),
    send(Socket, Rest);
send(Socket, Rest) ->
    gen_tcp:send(Socket, Rest).
server() ->
    {ok, ListenSocket} = gen tcp:listen(1234, [binary, {active, false}]),
    wait connect(ListenSocket,0).
wait connect(ListenSocket, Count) ->
    {ok, Socket} = gen_tcp:accept(ListenSocket),
    spawn(?MODULE, wait connect, [ListenSocket, Count+1]),
    get request(Socket, [], Count).
```

# Interfaces with other Programming Languages

- JInterface
- C Node
- Ports
- NIFs

# Java - JInterface

```
OtpNode self = new OtpNode("gurka");
  OtpMbox mbox = self.createMbox("echo");
  OtpErlangObject o;
  OtpErlangTuple msg;
  OtpErlangPid from;

while (true) {
    try {
        o = mbox.receive();
        if (o instanceof OtpErlangTuple) {
            msg = (OtpErlangTuple)o;
            from = (OtpErlangPid)(msg.elementAt(0));
            mbox.send(from,msg.elementAt(1));
    }
    catch (Exception e) {
        System.out.println("" + e);
    }
}
```

## **C** Nodes

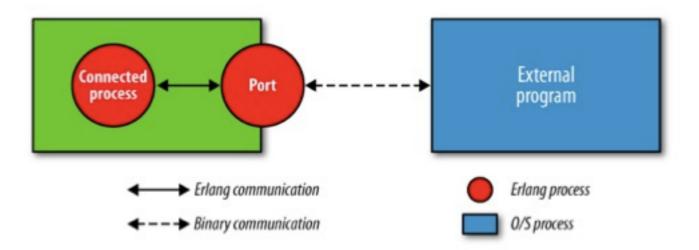
From Erlang's point of view, the C node is treated like a normal Erlang node.

Therefore, calling the functions foo and bar only involves sending a message to the C node asking for the function to be called, and receiving the result.

Sending a message requires a recipient; a process which can be defined using either a pid or a tuple consisting of a registered name and a node name.

## Ports

An Erlang port allows communication between an Erlang node and an external program through binary messages sent to and from an Erlang process running in the node— known as the connected process of the port—and the external program, running in a separate operating system thread.



## **NIFs**

A NIF (Native Implemented Function) is a function that is implemented in C instead of Erlang. NIFs appear as any other functions to the callers.

```
#include "erl_nif.h"
extern int foo(int x);
extern int bar(int y);
static ERL_NIF_TERM foo_nif(ErlNifEnv* env, int argc, const ERL_NIF_TERM argv[])
    int x, ret;
   if (!enif_get_int(env, argv[0], &x)) {
     return enif_make_badarg(env);
    ret = foo(x);
    return enif_make_int(env, ret);
}
static ERL_NIF_TERM bar_nif(ErlNifEnv* env, int argc, const ERL_NIF_TERM argv[])
    int y, ret;
    if (!enif_get_int(env, argv[0], &y)) {
     return enif_make_badarg(env);
    ret = bar(y);
    return enif_make_int(env, ret);
}
static ErlNifFunc nif_funcs[] = {
    {"foo", 1, foo_nif},
    {"bar", 1, bar_nif}
};
ERL_NIF_INIT(complex6, nif_funcs, NULL, NULL, NULL, NULL)
```

```
-module(complex6).
-export([foo/1, bar/1]).
-on_load(init/0).

init() ->
    ok = erlang:load_nif("./complex6_nif", 0).

foo(_X) ->
    exit(nif_library_not_loaded).
bar(_Y) ->
    exit(nif_library_not_loaded).
```

# the dbg Tracer

The basic steps of tracing for function calls are on a non-live node:

```
> dbg:start().
                                         % start dbg
                                         % start a simple tracer process
> dbg:tracer().
> dbg:tp(Module, Function, Arity, []). % specify MFA you are interested in
                                         % trace calls (c) of that MFA for all processes.
> dbg:p(all, c).
... trace here
> dbg:stop_clear().
                                         % stop tracer and clear effect of tp and p calls.
```

```
You can trace for multiple functions at the same time. > dbg:tpl(Module, '_', []). %
> dbg:tpl(Module, '_', '[]).
> dbg:tpl(Module, Function, '_', []).
> dbg:tpl(Module, Function, Arity, []).
> dbg:tpl(Module, Function, Arity, []).

% all calls in Module:Function with any arity.
% all calls to Module:Function/Arity.
> dbg:tpl(M, F, A, [{'_', [], [{return_trace}]}]). % same as before, but also show return value.
```

You can select the processes to trace on with the call to p().

```
> dbg:p(all, c). % trace calls to selected functions by all functions
> dbg:p(new, c). % trace calls by processes spawned from now on
> dbg:p(Pid, [c, m]). % trace calls and messages of a given process
```

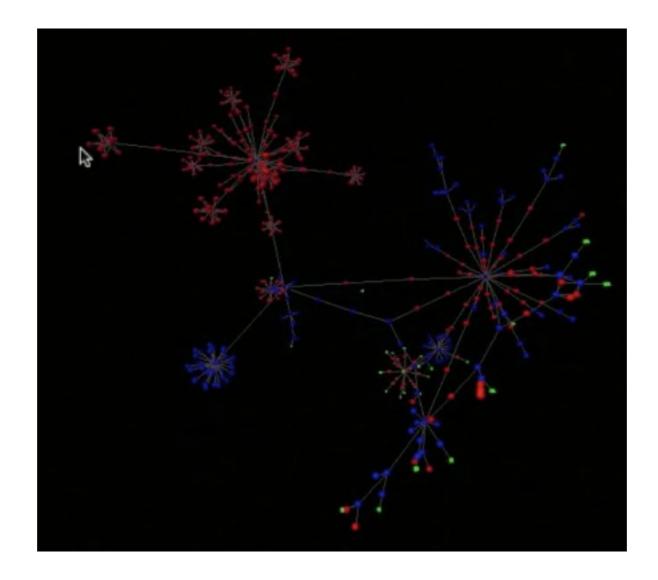
# Example Trace App

### A Visual Tour of Erlang

http://www.youtube.com/watch?v=IHoWfeNuAN8

### **Erlang Process Visualizer.**

http://github.com/krestenkrab/erlubi



# Types and Documentation

```
%% @doc Create the ETS and DETS tables which implement the database. The
%% argument gives the filename which is used to hold the DETS table.
%% If the table can be created, an `ok' tuple containing a
%% reference to the created table is returned; if not, it returns an `error'
%% tuple with a string describing the error.
%% @spec create tables(string()) -> {ok, reference()} | {error, string()}
-spec(create_tables(string()) -> {ok, ref()} | {error, string()}).
create tables(FileName) ->
    ets:new(subRam, [named_table, {keypos, #usr.msisdn}]),
    ets:new(subIndex, [named table]),
    dets:open file(subDisk, [{file, FileName}, {keypos, #usr.msisdn}]).
%% @doc Close the ETS and DETS tables implementing the database.
%% Returns either `ok' or and `error'
%% tuple with the reason for the failure to close the DETS table.
%% @spec close tables() -> ok | {error, string()}
-spec(close tables() -> ok | {error, string()}).
close tables() ->
    ets:delete(subRam),
    ets:delete(subIndex),
    dets:close(subDisk).
```

# Dialyzer

### Dialyzer: A Discrepancy AnaLYZer for ERlang Programs

```
Step I
$ dialyzer --build_plt -r <erl-lib>/kernel-2.12.5/ebin <erl-lib>/ stdlib-1.15.5/ebin <erl-lib>/mnesia-4.4.7/ebin
Creating PLT /Users/simonthompson/.dialyzer plt ...
re.erl:41: Call to missing or unexported function unicode:characters to binary/2
re.erl:134: Call to missing or unexported function unicode:characters to list/2
re.erl:200: Call to missing or unexported function re:compile/2
re.erl:226: Call to missing or unexported function unicode:characters to binary/2
re.erl:245: Call to missing or unexported function unicode:characters_to_list/2
re.erl:505: Call to missing or unexported function unicode:characters to list/2
re.erl:545: Call to missing or unexported function unicode:characters to binary/2
Unknown functions:
compile:file/2
compile:forms/2
compile:noenv forms/2
compile:output_generated/1 c
rypto:des3 cbc decrypt/5
crypto:start/0
done in 16m43.44s
done (warnings were emitted)
Step 2
$ dialyzer -c usr.erl usr_db.erl
Checking whether the PLT /Users/simonthompson/.dialyzer plt is up-to-date... yes
Proceeding with analysis...
usr.erl:110: The pattern [] can never match the type {'error', 'instance'}
usr db.erl:69: Call to missing or unexported function ets:safefixtable/2
done in 0m0.33s
done (warnings were emitted)
```

### **EUnit**

```
-module(fib).
-export([fib/1]).
-include_lib("eunit/include/eunit.hrl").
fib(0) -> 1;
fib(1) \rightarrow 1;
fib(N) when N > 1 -> fib(N-1) + fib(N-2).
fib_test_() ->
   [?_assert(fib(0) =:= 1),
    ?_assert(fib(1) =:= 1),
    ?_assert(fib(2) =:= 2),
    ?_assert(fib(3) =:= 3),
    ?_assert(fib(4) =:= 5),
    ?_assert(fib(5) =:= 8),
    ?_assertException(error, function_clause, fib(-1)),
   ?_assert(fib(31) =:= 2178309)
   ].
```

## QuickCheck

```
%% Encode/Decode bool
%%----
prop_bool() ->
    ?FORALL({Id, Bool},
        {non_neg_integer(), oneof([boolean(), 0, 1])},
        begin
          Fun = fun (B) when B = := 1; B = := true \rightarrow true;
           (B) when B = := 0; B = := false \rightarrow false
            end,
          {{Id, Fun(Bool)}, <<>>} =:=
        (?DECODE((?ENCODE(Id, Bool, bool)), bool))
        end).
enclode bool test () ->
    [?_assertMatch(<<8, 1>>, (?ENCODE(1, true, bool))),
     ?_assertMatch(<<8, 0>>, (?ENCODE(1, false, bool))),
     ?_assertMatch(<<40, 1>>, (?ENCODE(5, 1, bool))),
     ?_assertMatch(<<40, 0>>, (?ENCODE(5, 0, bool)))].
decode_bool_test_() ->
    [? assertMatch({{1, true}, <<>>},
           (?DECODE(<<8, 1>>, bool))),
     ? assertMatch({{1, false}, <<>>},
           (?DECODE(<<8, 0>>, bool)))].
```

# Applications

lost Watched This Month	Most Forked This Month
ericmoritz / wsdemo	ericmoritz / wsdemo
erlang / otp	erlang / otp
krestenkrab / erlubi	basho / rebar
extend / cowboy	99s extend / cowboy
apache / couchdb	apache / couchdb
ost Watched Overall	Most Forked Overall
erlang / otp	erlang / otp
apache / couchdb	apache / couchdb
mochi / mochiweb	basho / rebar
basho / rebar	mochi / mochiweb
arcieri / reia	processone / ejabberd

## Conferences

Erlang User Conference 2012 Erlang Factory SF Bay Area 2012 Krakow Erlang Factory Lite 2012 Brussels Erlang Factory Lite Brisbane Factory Lite Erlang User Conference 2011 !!! > 320 Amsterdam Factory Lite Paris Erlang Factory Lite Edinburgh Factory Lite Erlang Factory London 2011 Erlang Factory Lite Munich Erlang Factory SF Bay Area 2011 !!! > 170 Erlang Factory Lite Krakow 2010 Erlang User Conference 2010 Tutorial Workshop 2010, Stockholm Erlang Factory Lite LA Erlang Factory London 2010 Erlang Factory SF Bay Area 2010



Erlang Factory Lite Krakow 2009
Property-based Testing Tutorial Workshop
2009, Stockholm
Erlang User Conference 2009, Stockholm
SIGPLAN Erlang Workshop 2009, Edinburgh
Erlang Factory London 2009
Erlang Factory SF Bay Area 2009

# Advantages

- Nice architectures can be desinged
- Different way of thinking applications
- Debugging live is really nice
- Extensive Documentation

## Issues & Risks

- Sometimes I miss Types
- Concurrency debugging for large applications is hard, even in Erlang
- Compilation time on large projects is big
- Memory/Process Leaks can crash your system if not handled correctly

## Get Started

- Books
- GitHub
- Learn You some Erlang
- Mailing List(s)
- Erlang Factory Conferences

# SpawnFest 2012

http://spawnfest.com/ July 7-8, 2012

### SPAWNFEST 2012

An annual 48 hour development competition in which teams of skilled developers get exactly one weekend to create the best Erlang applications they can

### Erlang is quite hot today

### Thank You!

