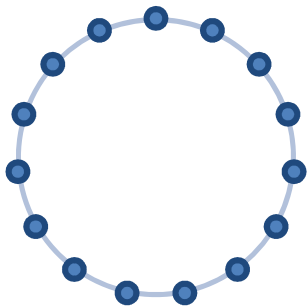


## TRANSACTIONS



## Scalaris:

### Users and Developers Guide

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Part I.

# Users Guide

# 1. Introduction

Scalaris is a scalable, transactional, distributed key-value store based on the peer-to-peer principle. It can be used to build scalable Web 2.0 services. The concept of Scalaris is quite simple: Its architecture consists of three layers.

It provides self-management and scalability by replicating services and data among peers. Without system interruption it scales from a few PCs to thousands of servers. Servers can be added or removed on the fly without any service downtime.



Scalaris takes care of:

- Fail-over
- Data distribution
- Replication
- Strong consistency
- Transactions

The Scalaris project was initiated by Zuse Institute Berlin and onScale solutions and was partly funded by the EU projects Selfman and XtreamOS. Additional information (papers, videos) can be found at <http://www.zib.de/CSR/Projects/scalaris> and <http://www.onscale.de/scalarix.html>.

## 1.1. Brewer's CAP Theorem

In distributed computing there exists the so called CAP theorem. It basically says that there are three desirable properties for distributed systems but one can only have any two of them.

**Strict Consistency.** Any read operation has to return the result of the latest write operation on the same data item.

**Availability.** Items can be read and modified at any time.

**Partition Tolerance.** The network on which the service is running may split into several partitions which cannot communicate with each other. Later on the networks may re-join again.

For example, a service is hosted on one machine in Seattle and one machine in Berlin. This service is partition tolerant if it can tolerate that all Internet connections over the Atlantic (and Pacific) are interrupted for a few hours and then get repaired.

The goal of Scalaris is to provide strict consistency and partition tolerance. We are willing to sacrifice availability to make sure that the stored data is always consistent. I.e. when you are running Scalaris with a replication degree of 4 and the network splits into two partitions, one partition with three replicas and one partition with one replica, you will be able to continue to use the service only in the larger partition. All requests in the smaller partition will time out until the two networks merge again. Note, most other key-value stores tend to sacrifice consistency.

## 1.2. Scientific Background

**Basics.** The general structure of Scalaris is modelled after Chord. The Chord paper [4] describes the ring structure, the routing algorithms, and basic ring maintenance.

The main routines of our Chord node are in `src/dht_node.erl` and the join protocol is implemented in `src/dht_node_join.erl` (see also Chap. 10). Our implementation of the routing algorithms is described in more detail in Sect. 8.3 and the actual implementation is in `src/rt_chord.erl`.

**Transactions.** The most interesting part is probably the transaction algorithms. The most current description of the algorithms and background is in [6].

The implementation consists of the paxos algorithm in `src/paxos` and the transaction algorithms itself in `src/transactions` (see also Chap. 9).

**Ring Maintenance.** We changed the ring maintenance algorithm in Scalaris. It is not the standard Chord one, but a variation of T-Man [5]. It is supposed to fix the ring structure faster. In some situations, the standard Chord algorithm is not able to fix the ring structure while T-Man can still fix it. For node sampling, our implementation relies on Cyclon [7].

The T-Man implementation can be found in `src/rm_tman.erl` and the Cyclon implementation in `src/cyclon`.

**Vivaldi Coordinates.** For some experiments, we implemented so called Vivaldi coordinates [2]. They can be used to estimate the network latency between arbitrary nodes.

The implementation can be found in `src/vivaldi.erl`.

**Gossiping.** For some algorithms, we use estimates of global information. These estimates are aggregated with the help of gossiping techniques [8].

The implementation can be found in `src/gossip.erl`.

## 2. Download and Installation

### 2.1. Requirements

For building and running Sclaris, some third-party modules are required which are not included in the Sclaris sources:

- Erlang R13B01 or newer
- GNU-like Make

To build the Java API (and the command-line client) the following modules are required additionally:

- Java Development Kit 6
- Apache Ant

Before building the Java API, make sure that `JAVA_HOME` and `ANT_HOME` are set. `JAVA_HOME` has to point to a JDK installation, and `ANT_HOME` has to point to an Ant installation.

### 2.2. Download

The sources can be obtained from <http://code.google.com/p/scalaris>. RPMs and DEBs are available from <http://download.opensuse.org/repositories/home:/tschuett/>.

#### 2.2.1. Development Branch

You find the latest development version in the svn repository:

```
# Non-members may check out a read-only working copy anonymously over HTTP.  
svn checkout http://scalaris.googlecode.com/svn/trunk/ scalaris-read-only
```

#### 2.2.2. Releases

Releases can be found under the 'Download' tab on the web-page.

### 2.3. Configuration

Sclaris reads two configuration files from the working directory: `bin/scalaris.cfg` (mandatory) and `bin/scalaris.local.cfg` (optional). The former defines default settings and is included in the release. The latter can be created by the user to alter settings. A sample file is provided as `bin/scalaris.local.cfg.example`. To run Sclaris distributed over several nodes, each node requires a `bin/scalaris.local.cfg`:



File `scalaris.local.cfg`:

```
%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
% Settings for distributed Erlang
% (see scalaris.hrl to switch)

% {boot_host, {boot, 'boot@foo.bar.com'}}.
% {known_hosts, [{service_per_vm, 'boot@foo.bar.com'}]}.

%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
% Settings for TCP mode.
% (see scalaris.hrl to switch)

% Insert the appropriate IP-addresses for your setup
% as comma separated integers:
% IP Address, Port, and label of the boot server
{boot_host, {{127,0,0,1},14195,boot}}.

% IP Address, Port, and label of a node which is already in the system
{known_hosts, [{{{127,0,0,1},14195, service_per_vm}]}].
```

Scalaris currently distinguishes two different kinds of nodes: (a) the boot-server and (b) regular nodes. For the moment, we limit the number of boot-servers to exactly one. The remaining nodes are regular nodes. The boot-server is contacted to join the system. On all servers, the `boot_host` option defines the server where the boot server is running. In the example, it is an IP address plus a TCP port.

## 2.4. Build

### 2.4.1. Linux

Scalaris uses `autoconf` for configuring the build environment and `GNU Make` for building the code.

```
%> ./configure
%> make
%> make docs
```

For more details read `README` in the main Scalaris checkout directory.

### 2.4.2. Windows

We are currently not supporting Scalaris on Windows. However, we have two small bat files for building and running scalaris nodes. It seems to work but we make no guarantees.

- Install Erlang  
<http://www.erlang.org/download.html>
- Install OpenSSL (for crypto module)  
<http://www.slproweb.com/products/Win32OpenSSL.html>
- Checkout scalaris code from SVN
- adapt the path to your Erlang installation in `build.bat`
- start a `cmd.exe`
- go to the scalaris directory
- run `build.bat` in the cmd window

- check that there were no errors during the compilation; warnings are fine
- go to the bin sub-directory
- adapt the path to your Erlang installation in `boot.bat`, `cs_local.bat`, `cs_local2.bat` and `cs_local3.bat`
- run `boot.bat` or one of the other start scripts in the cmd window

`build.bat` will generate a `Emakefile` if there is none yet. If you have Erlang < R13B04, you will need to adapt the `Emakefile`. There will be empty lines in the first three blocks ending with “`}]`.”: add the following to these lines and try to compile again. It should work now.

```
, {d, type_forward_declarations_are_not_allowed}
, {d, forward_or_recursive_types_are_not_allowed}
```

For the most recent description please see the FAQ at <http://code.google.com/p/scalaris/wiki/FAQ>.

### 2.4.3. Java-API

The following commands will build the Java API for Scalaris:

```
%> make java
```

This will build `scalaris.jar`, which is the library for accessing the overlay network. Optionally, the documentation can be build:

```
%> cd java-api
%> ant doc
```

## 2.5. Running Scalaris

As mentioned above, in Scalaris there are two kinds of nodes:

- boot servers
- regular nodes

In every Scalaris, at least one boot server is required. It will maintain a list of nodes taking part in the system and allows other nodes to join the ring. For redundancy, it is also possible to have several boot servers. In the future, we want to eliminate this distinction, so any node is also a boot-server.

### 2.5.1. Running on a local machine

Open at least two shells. In the first, inside the Scalaris directory, start the boot script (`boot.bat` on Windows):

```
%> ./bin/boot.sh
```

This will start the boot server. On success <http://localhost:8000> should point to the management interface page of the boot server. The main page will show you the number of nodes currently in

the system. After a couple of seconds a first Scalaris should have started in the boot server and the number should increase to one. The main page will also allow you to store and retrieve key-value pairs but should not be used by applications to access Scalaris. See Chapter 3 on page 13 for application APIs.

In a second shell, you can now start a second Scalaris node. This will be a ‘regular server’:

```
%> ./bin/cs_local.sh
```

The second node will read the configuration file and use this information to contact the boot server and join the ring. The number of nodes on the web page should have increased to two by now.

Optionally, a third and fourth node can be started on the same machine. In a third shell:

```
%> ./bin/cs_local2.sh
```

In a fourth shell:

```
%> ./bin/cs_local3.sh
```

This will add 3 nodes to the network. The web pages at <http://localhost:8000> should show the additional nodes.

On linux you can also use the `scalarisctl` script to start boot and ‘regular’ nodes.

### 2.5.2. Running distributed

Scalaris can be installed on other machines in the same way as described in Section 2.6. In the default configuration, nodes will look for the boot server on localhost on port 14195. You should create a `scalaris.local.cfg` pointing to the node running the boot server.

```
% Insert the appropriate IP-addresses for your setup
% as comma separated integers:
% IP Address, Port, and label of the boot server
{boot_host, {{127,0,0,1},14195,boot}}.
```

If you are using the default configuration on the boot server it will listen on port 14195 and you only have to change the IP address in the configuration file. Otherwise the other nodes will not find the boot server. On the remote nodes, you only need to call `./cs_local.sh` and they will automatically contact the configured boot server.

## 2.6. Installation

For simple tests, you do not need to install Scalaris. You can run it directly from the source directory. Note: make `install` will install scalaris into `/usr/local` and place `scalarisctl` into `/usr/local/bin`. But is more convenient to build an RPM and install it.

```
svn checkout http://scalaris.googlecode.com/svn/trunk/ scalaris-0.0.1
tar -cvjf scalaris-0.0.1.tar.bz2 scalaris-0.0.1 --exclude-vcs
cp scalaris-0.0.1.tar.bz2 /usr/src/packages/SOURCES/
rpmbuild -ba scalaris-0.0.1/contrib/scalaris.spec
```

Your source and binary RPM will be generated in `/usr/src/packages/SRPMS` and `RPMS`. We build RPMs and Debs using checkouts from svn and provide them using the openSUSE BuildService at <http://download.opensuse.org/repositories/home:/tschuett/>. Packages are available for

- Fedora 9, 10, 11, 12, 13,
- Mandriva 2008, 2009, 2009.1, 2010,
- openSUSE 11.0, 11.1, 11.2, 11.3, Factory,
- SLE 10, 11,
- CentOS 5.4,
- RHEL 5,
- Debian 5.0 and
- Ubuntu 9.04, 9.10, 10.04.

Inside those repositories you will also find an erlang package - you don't need this if you already have a recent enough erlang version!

## 2.7. Logging

*Description is based on SVN revision r1083.*

Scalaris uses the `log4erl` library (see `contrib/log4erl`) for logging status information and error messages. The log level can be configured in `bin/scalaris.cfg` for both the stdout and file logger. The default value is `warn`; only warnings, errors and severe problems are logged.

```
%% @doc Loglevel: debug < info < warn < error < fatal < none
{log_level, warn}.
{log_level_file, warn}.
```

In some cases, it might be necessary to get more complete logging information, e.g. for debugging. In Chapter 10 on page 44, we are explaining the startup process of Scalaris nodes in more detail, here the `info` level provides more detailed information.

```
%% @doc Loglevel: debug < info < warn < error < fatal < none
{log_level, info}.
{log_level_file, info}.
```

## 3. Using the system

### 3.1. JSON API

Scalaris supports a JSON API for transactions. To minimize the necessary round trips between a client and Scalaris, it uses request lists, which contain all requests that can be done in parallel. The request list is then send to a Scalaris node with a POST message. The result is an opaque TransLog and a list containing the results of the requests. To add further requests to the transaction, the TransLog and another list of requests may be send to Scalaris. This process may be repeated as often as necessary. To finish the transaction, the request list can contain a 'commit' request as the last element, which triggers the validation phase of the transaction processing.

The JSON-API can be accessed via the Scalaris-Web-Server running on port 8000 by default and the page `jsonrpc.yaws` (For example at: <http://localhost:8000/jsonrpc.yaws>). The following example illustrates the message flow:

#### Client

Make a transaction, that sets two keys:

#### Scalaris node

→

```
{
  "method": "req_list",
  "version": "1.1",
  "params":
  [
    [
      { "write": { "keyA": "valueA" } },
      { "write": { "keyB": "valueB" } },
      { "commit": "commit" }
    ]
  ],
  "id": 0
}
```

← Scalaris sends results back

```
{ "result":
  { "results":
    [
      { "op": "commit",
        "value": "ok",
        "key": "ok" },
      { "op": "write",
        "value": "valueB",
        "key": "keyB" },
      { "op": "write",
        "value": "valueA",
        "key": "keyA" }
    ],
    "translog":
      [...]
  },
  "id" : 0
}
```

In a second transaction: Read the two keys →

```
{
  "method": "req_list",
  "version": "1.1",
  "params":
    [
      [
        { "read": "keyA" },
        { "read": "keyB" }
      ]
    ]
  "id": 0
}
```

← Scalaris sends results back

```
{ "result":
  { "results":
    [
      { "op": "read",
        "value": "valueB",
        "key": "keyB" },
      { "op": "read",
        "value": "valueA",
        "key": "keyA" }
    ],
    "translog":
      [...] // this list is the translog
              // for further operations!
              // We name it TLOG here.
  },
  "id" : 0
}
```

Calculate something with the read values →  
and make further requests, here a write  
and the commit for the whole transaction.  
Also include the latest translog we got from  
Scalaris (named TLOG here).

```
{
  "method": "req_list",
  "version": "1.1",
  "params":
  [
    TLOG, // translog from prev. result
    [
      { "write": { "keyA": "valueA2" } },
      { "commit": "commit" }
    ]
  ],
  "id" : 0
}
```

← Scalaris sends results back

```
{ "result":
  { "results":
    [ { "op": "commit",
        "value": "ok",
        "key": "ok" },
      { "op": "write",
        "value": "valueA2",
        "key": "keyA" }
    ],
    "translog":
    [...]
  },
  "id" : 0
}
```

A sample usage of the JSON API using Ruby can be found in contrib/jsonrpc.rb.

A single request list must not contain a key more than once!

The allowed requests are:

```
{ "read": "any_key" }
{ "write": { "any_key": "any_value" } }
{ "commit": "commit" }
```

The possible results are:

```
{ "op": "read", "key": "any_key", "value": "any_value" }
{ "op": "read", "key": "any_value", "fail": "reason" } // 'not_found' or 'timeout'

{ "op": "write", "key": "any_key", "value": "any_value" }
{ "op": "read", "key": "any_key", "fail": "reason" }

{ "op": "commit", "value": "ok", "key": "ok" }
{ "op": "commit", "value": "fail", "fail": "reason" }
```

### 3.1.1. Deleting a key

Outside transactions keys can also be deleted, but it has to be done with care, as explained in the following thread on the mailing list: [http://groups.google.com/group/scalaris/browse\\_thread/thread/ff1d9237e218799](http://groups.google.com/group/scalaris/browse_thread/thread/ff1d9237e218799).

```
{
  "method": "delete",
  "version": "1.1",
  "params":
    [
      { "key": "any_key" }
    ],
  "id" : 0
}
```

#### Two sample results

```
{ "result":
  { "ok":2, // how many replicas were deleted successssfully
    "results": [ "ok", "ok", "locks_set", "undef" ]
  }
}
```

```
{ "result":
  { "failure": "reason" }
}
```

## 3.2. Java command line interface

The jar file contains a small command line interface client. For convenience, we provide a wrapper script called `scalaris` which sets up the Java environment:

```
%> ./java-api/scalaris -help
Script Options:
  --help, --h          print this message and scalaris help
  --noconfig           suppress sourcing of /etc/scalaris/scalaris-java.conf
                      and $HOME/.scalaris/scalaris-java.conf config files
  --execdebug          print scalaris exec line generated by this
                      launch script

usage: scalaris [Options]
  -b,--minibench        run mini benchmark
  -d,--delete <key> <[timeout]> delete an item (default timeout: 2000ms)
                          WARNING: This function can lead to inconsistent data
                          (e.g. deleted items can re-appear). Also when
                          re-creating an item the version before the delete can
                          re-appear.
  -g,--getsubscribers <topic> get subscribers of a topic
  -h,--help            print this message
  -lh,--localhost      gets the local host's name as known to
                          Java (for debugging purposes)
  -p,--publish <topic> <message> publish a new message for the given
                          topic
  -r,--read <key>      read an item
  -s,--subscribe <topic> <url> subscribe to a topic
  -u,--unsubscribe <topic> <url> unsubscribe from a topic
  -v,--verbose         print verbose information, e.g. the
                          properties read
  -w,--write <key> <value> write an item
```



read, write and delete can be used to read, write and delete from/to the overlay, respectively. getsubscribers, publish, and subscribe are the PubSub functions. The others provide debugging and testing functionality.

```
%> ./java-api/scalaris -write foo bar
write(foo, bar)
%> ./java-api/scalaris -read foo
read(foo) == bar
```

Per default, the scalaris script tries to connect to a boot server at localhost. You can change the node it connects to (and further connection properties) by adapting the values defined in java-api/scalaris.properties.

### 3.3. Java API

The scalaris.jar provides the command line client as well as a library for Java programs to access Scalaris. The library provides two classes:

- Scalaris provides a high-level API similar to the command line client.
- Transaction provides a low-level API to the transaction mechanism.

For details we refer the reader to the Javadoc:

```
%> cd java-api
%> ant doc
%> firefox doc/index.html
```

## 4. Testing the system

### 4.1. Running the unit tests

There are some unit tests in the `test` directory. You can call them by running `make test` in the main directory. The results are stored in a local `index.html` file.

The tests are implemented with the `common-test` package from the Erlang system. For running the tests we rely on `run_test`, which is part of the `common-test` package, but (on `erlang < R14`) is not installed by default. `configure` will check whether `run_test` is available. If it is not installed, it will show a warning and a short description of how to install the missing file.

Note: for the unit tests, we are setting up and shutting down several overlay networks. During the shut down phase, the runtime environment will print extensive error messages. These error messages do not indicate that tests failed! Running the complete test suite takes about 3 minutes, depending on your machine. Only if the complete suite finishes, it will present statistics on failed and successful tests.

# 5. Troubleshooting

## 5.1. Network

Scalaris uses a couple of TCP ports for communication. It does not use UDP at the moment.

- 8000 HTTP Server on the boot node
- 8001 HTTP Server on the other nodes
- 14195 Port for inter-node communication (boot server)
- 14196 Port for inter-node communication (other nodes)

Please make sure that at least 14195 and 14196 are not blocked by firewalls.

## Part II.

# Developers Guide

## 6. General Hints

### 6.1. Coding Guidelines

- Keep the code short
- Use `gen_component` to implement additional processes
- Don't use `receive` by yourself (Exception: to implement single threaded user API calls (`cs_api`, `yaws_calls`, etc))
- Don't use `erlang:now/0`, `erlang:send_after/3`, `receive after` etc. in performance critical code, consider using `msg_delay` instead.
- Don't use `timer:tc/3` as it catches exceptions. Use `util:tc/3` instead.

### 6.2. Testing Your Modifications and Extensions

- Run the testsuites using `make test`
- Run the java api test using `make java-test` (Scalaris output will be printed if a test fails; if you want to see it during the tests, start a `bin/boot.sh` and run the tests by `cd java; ant test`)
- Run the Ruby client by starting Scalaris and running `cd contrib; ./jsonrpc.rb`

### 6.3. Help with Digging into the System

- use `ets:i/0,1` to get details on the local state of some processes
- consider changing `pdb.erl` to use `ets` instead of `erlang:put/get`
- Have a look at `strace -f -p PID` of beam process
- Get message statistics via the Web-interface
- enable/disable tracing for certain modules
- Use `etop` and look at the total memory size and atoms generated
- send processes `sleep` or `kill` messages to test certain behaviour (see `gen_component.erl`)
- use `boot_server:number_of_nodes(). flush().`
- use `admin_checkring(). flush().`

### 6.4. General Erlang server loop

Servers in Erlang often use the following structure to maintain a state while processing received messages:

```
loop(State) ->  
  receive  
    Message ->  
      State1 = f(State),  
      loop(State1)  
  end.
```

The server runs an endless loop, that waits for a message, processes it and calls itself using tail-recursion in each branch. The loop works on a State, which can be modified when a message is handled.

## 7. System Infrastructure

### 7.1. Groups of Processes

- What is it? How to distinguish from Erlangs internal named processes?
- Joining a process group
- Why do we do this... (managing several independent nodes inside a single Erlang VM for testing)

### 7.2. The Communication Layer `comm`

- in general
- format of messages (tuples)
- use messages with cookies (server and client side)
- What is a message tag?

### 7.3. The `gen_component`

*Description is based on SVN revision r993.*

The generic component model implemented by `gen_component` allows to add some common functionality to all the components that build up the Scalaris system. It supports:

**event-handlers:** message handling with a similar syntax as used in [3].

**FIFO order of messages:** components cannot be inadvertently locked as we do not use selective receive statements in the code.

**sleep and halt:** for testing components can sleep or be halted.

**debugging, breakpoints, stepwise execution:** to debug components execution can be steered via breakpoints, step-wise execution and continuation based on arriving events and user defined component state conditions.

**basic profiling** ,

**state dependent message handlers:** depending on its state, different message handlers can be used and switched during runtime. Thereby a kind of state-machine based message handling is supported.

**prepared for `pid_groups`:** allows to send events to named processes inside the same group as the actual component itself (`send_to_group_member`) when just holding a reference to any group member, and

**unit-testing of event-handlers:** as message handling is separated from the main loop of the component, the handling of individual messages and thereby performed state manipulation can easily be tested in unit-tests by directly calling message handlers.

In *Scalaris* all Erlang processes should be implemented as `gen_component`. The only exception are functions interfacing to the client, where a transition from asynchronous to synchronous request handling is necessary and that are executed in the context of a client's process or a process that behaves as a proxy for a client (`cs_api`).

### 7.3.1. A basic `gen_component` including a message handler

To implement a `gen_component`, the component has to provide the `gen_component` behaviour:

File `gen_component.erl`:

```

48 -spec behaviour_info(atom()) -> [{atom(), arity()}] | undefined.
49 behaviour_info(callbacks) ->
50 [
51     {init, 1}      % initialize component
52     % note: can use arbitrary on-handler, but by default on/2 is used:
53     {on, 2}        % handle a single message
54     % on(Msg, State) -> NewState | unknown_event | kill
55 ];

```

This is illustrated by the following example:

File `msg_delay.erl`:

```

70 %% initialize: return initial state.
71 -spec init([]) -> state().
72 init([]) ->
73     MyGroup = pid_groups:my_groupname(),
74     ?TRACE("msg_delay:init for pid group ~p~n", [MyGroup]),
75     TimeTableName = list_to_atom(MyGroup ++ "_msg_delay"),
76     %% use random table name provided by ets to *not* generate an atom
77     %% TableName = pdb:new(?MODULE, [set, private]),
78     TimeTable = pdb:new(TimeTableName, [set, protected, named_table]),
79     comm:send_local(self(), {msg_delay_periodic}),
80     _State = {TimeTable, _Round = 0}.
81
82 -spec on(message(), state()) -> state().
83 on({msg_delay_req, Seconds, Dest, Msg} = _FullMsg,
84     {TimeTable, Counter} = State) ->
85     ?TRACE("msg_delay:on(~.0p, ~.0p)~n", [_FullMsg, State]),
86     Future = trunc(Counter + Seconds),
87     case pdb:get(Future, TimeTable) of
88         undefined ->
89             pdb:set({Future, [{Dest, Msg}]}, TimeTable);
90         {_, MsgQueue} ->
91             pdb:set({Future, [{Dest, Msg} | MsgQueue]}, TimeTable)
92     end,
93     State;
94
95 %% periodic trigger
96 on({msg_delay_periodic} = Trigger, {TimeTable, Counter} = _State) ->
97     ?TRACE("msg_delay:on(~.0p, ~.0p)~n", [Trigger, State]),
98     case pdb:get(Counter, TimeTable) of
99         undefined -> ok;
100         {_, MsgQueue} ->
101             _ = [ comm:send_local(Dest, Msg) || {Dest, Msg} <- MsgQueue ],
102             pdb:delete(Counter, TimeTable)
103     end,
104     comm:send_local_after(1000, self(), Trigger),
105     {TimeTable, Counter + 1};
106
107 on({web_debug_info, Requestor}, {TimeTable, Counter} = State) ->
108     KeyValueCollection =
109         [{"queued messages (in 0-10s, messages):", ""}] |
110         [begin
111             Future = trunc(Counter + Seconds),

```



```

112         Queue = case pdb:get(Future, TimeTable) of
113             undefined -> none;
114             {_, Q}    -> Q
115         end,
116         {lists:flatten(io_lib:format("~p", [Seconds])),
117          lists:flatten(io_lib:format("~p", [Queue]))},
118     end || Seconds <- lists:seq(0, 10)]],
119     comm:send_local(Requestor, {web_debug_info_reply, KeyValueType}),
120     State.

```

`your_gen_component:init/1` is called during start-up of a `gen_component` and should return the initial state to be used for this `gen_component`. Later, the current state of the component can be retrieved using `gen_component:get_state/1`.

To react on messages / events, a message handler is used. The default message handler is called `your_gen_component:on/2`. This can be changed by calling `gen_component:change_handler/2` (see Section 7.3.6). When an event / message for the component arrives, this handler is called with the event itself and the current state of the component. In the handler, the state of the component may be adjusted depending upon the event. The handler itself may trigger new events / messages for itself or other components and has finally to return the updated state of the component or the atoms `unknown_event` or `kill`. It must neither call `receive` nor `timer:sleep/1` nor `erlang:exit/1`.

### 7.3.2. How to start a `gen_component`?

A `gen_component` can be started using one of:

```
gen_component:start(Module, Args, GenCOptions = [])
```

```
gen_component:start_link(Module, Args, GenCOptions = [])
```

Module: the name of the module your component is implemented in

Args: List of parameters passed to `Module:init/1` for initialization

GenCOptions: optional parameter. List of options for `gen_component`

`{pid_groups_join_as, ProcessGroup, ProcessName}`: registers the new process with the given process group (also called `instanceid`) and name using `pid_groups`.

`{erlang_register, ProcessName}`: registers the process as a named Erlang process.

`wait_for_init`: wait for `Module:init/1` to return before returning to the caller.

These functions are compatible to the Erlang/OTP supervisors. They spawn a new process for the component which itself calls `Module:init/1` with the given `Args` to initialize the component. `Module:init/1` should return the initial state for your component. For each message sent to this component, the default message handler `Module:on(Message, State)` will be called, which should react on the message and return the updated state of your component.

`gen_component:start()` and `gen_component:start_link()` return the pid of the spawned process as `{ok, Pid}`.

### 7.3.3. When does a `gen_component` terminate?

A `gen_component` can be stopped using:

`gen_component:kill(Pid)` or by returning `kill` from the current message handler.

### 7.3.4. What happens when unexpected events / messages arrive?

Your message handler (default is `your_gen_component:on/2`) should return `unknown_event` in the final clause (`your_gen_component:on(_, _)`). `gen_component` then will nicely report on the unhandled message, the component's name, its state and currently active message handler, as shown in the following example:

```
# bin/boot.sh
[...]
(boot@localhost)10> pid_groups ! {no_message}.
{no_message}
[error] unknown message: {no_message} in Module: pid_groups and
handler on in State null
(boot@localhost)11>
```

The `pid_groups` (see Section 7.1) is a `gen_component` which registers itself as named Erlang process with the `gen_component` option `erlang_register` and therefore can be addressed by its name in the Erlang shell. We send it a `{no_message}` and `gen_component` reports on the unhandled message. The `pid_groups` module itself continues to run and waits for further messages.

### 7.3.5. What if my message handler generates an exception or crashes the process?

`gen_component` catches exceptions generated by message handlers and reports them with a stack trace, the message, that generated the exception, and the current state of the component.

If a message handler terminates the process via `erlang:exit/1`, this is out of the responsibility scope of `gen_component`. As usual in Erlang, all linked processes will be informed. If for example `gen_component:start_link/2` or `/3` was used for starting the `gen_component`, the spawning process will be informed, which may be an Erlang supervisor process taking further actions.

### 7.3.6. Changing message handlers and implementing state dependent message responsiveness as a state-machine

Sometimes it is beneficial to handle messages depending on the state of a component. One possibility to express this is implementing different clauses depending on the state variable, another is introducing case clauses inside message handlers to distinguish between current states. Both approaches may become tedious, error prone, and may result in confusing source code.

Sometimes the use of several different message handlers for different states of the component leads to clearer arranged code, especially if the set of handled messages changes from state to state. For example, if we have a component with an initialization phase and a production phase afterwards, we can handle in the first message handler messages relevant during the initialization phase and simply queue all other requests for later processing using a common default clause.

When initialization is done, we handle the queued user requests and switch to the message handler for the production phase. The message handler for the initialization phase does not need to know about messages occurring during production phase and the message handler for the production phase does not need to care about messages used during initialization. Both handlers can be made independent and may be extended later on without any adjustments to the other.

One can also use this scheme to implement complex state-machines by changing the message handler from state to state.

To switch the message handler `gen_component:change_handler(State, new_handler)` is called as

the last operation after a message in the active message handler was handled, so that the return value of `gen_component:change_handler/2` is propagated to `gen_component`. The new handler is given as an atom, which is the name of the 2-ary function in your component module to be called.

### Starting with non-default message handler.

It is also possible to change the message handler right from the start in your `your_gen_component:init/1` to avoid the default message handler `your_gen_component:on/2`. Just create your initial state as usual and call `gen_component:change_handler(State, my_handler)` as the final call in your `your_gen_component:init/1`. We prepared `gen_component:change_handler/2` to return `State` itself, so this will work properly.

### 7.3.7. Halting and pausing a gen\_component

Using `gen_component:kill(Pid)` and `gen_component:sleep(Pid, Time)` components can be terminated or paused.

### 7.3.8. Integration with pid\_groups: Redirecting events / messages to other gen\_components

Each `gen_component` by itself is prepared to support `comm:send_to_group_member/3` which forwards messages inside a group of processes registered via `pid_groups` (see Section 7.1) by their name. So, if you hold a `Pid` of one member of a process group, you can send messages to other members of this group, if you know their registered Erlang name. You do not necessarily have to know their individual `Pid`.

*In consequence, no `gen_component` can individually handle messages of the form `{send_to_group_member, _, _}` as such messages are consumed by `gen_component` itself.*

### 7.3.9. Replying to ping messages

Each `gen_component` replies automatically to `{ping, Pid}` requests with a `{pong}` send to the given `Pid`. Such messages are generated, for example, by `vivaldi_latency` which is used by our `vivaldi` module.

*In consequence, no `gen_component` can individually handle messages of the form: `{ping, _}` as such messages are consumed by `gen_component` itself.*

### 7.3.10. The debugging interface of gen\_component: Breakpoints and step-wise execution

We equipped `gen_component` with a debugging interface, which especially is beneficial, when testing the interplay between several `gen_components`. It supports breakpoints (bp) which can pause the `gen_component` depending on the arriving messages or depending on user defined conditions. If a breakpoint is reached, the execution can be continued step-wise (message by message) or until the next breakpoint is reached.

We use it in our unit tests to steer protocol interleavings and to perform tests using random protocol interleavings between several processes (see `paxos_SUITE`). It allows also to reproduce given protocol interleavings for better testing.

### Managing breakpoints.

Breakpoints are managed by the following functions:

`gen_component:bp_set(Pid, MsgTag, BPName)`: For the component running under `Pid` a breakpoint `BPName` is set. It is reached, when a message with a message tag `MsgTag` is next to be handled by the component (See `comm:get_msg_tag/1` and Section 7.2 for more information on message tags). The `BPName` is used as a reference for this breakpoint, for example to delete it later.

`gen_component:bp_set_cond(Pid, Cond, BPName)`: The same as `gen_component:bp_set/3` but a user defined condition implemented in `{Module, Function, Params = 2} = Cond` is checked by calling `Module:Function(Message, State)` to decide whether a breakpoint is reached or not. `Message` is the next message to be handled by the component and `State` is the current state of the component. `Module:Function/2` should return a boolean.

`gen_component:bp_del(Pid, BPName)`: The breakpoint `BPName` is deleted. If the component is in this breakpoint, it will not be released by this call. This has to be done separately by `gen_component:bp_cont/1`. But the deleted breakpoint will no longer be considered for newly entering a breakpoint.

`gen_component:bp_barrier(Pid)`: Delay all further handling of breakpoint requests until a breakpoint is actually entered.

*Note, that the following call sequence may not catch the breakpoint at all, as during the sleep the component not necessarily consumes a ping message and the set breakpoint 'sample\_bp' may already be deleted before a ping message arrives.*

```
gen_component:bp_set(Pid, ping, sample_bp),
timer:sleep(10),
gen_component:bp_del(Pid, sample_bp),
gen_component:bp_cont(Pid).
```

*To overcome this, `gen_component:bp_barrier/1` can be used:*

```
gen_component:bp_set(Pid, ping, sample_bp),
gen_component:bp_barrier(Pid),
%% After the bp_barrier request, following breakpoint requests
%% will not be handled before a breakpoint is actually entered.
%% The gen_component itself is still active and handles messages as usual
%% until it enters a breakpoint.
gen_component:bp_del(Pid, sample_bp),
% Delete the breakpoint after it was entered once (ensured by bp_barrier).
% Release the gen_component from the breakpoint and continue.
gen_component:bp_cont(Pid).
```

None of the calls in the sample listing above is blocking. It just schedules all the operations, including the `bp_barrier`, for the `gen_component` and immediately finishes. The actual events of entering and continuing the breakpoint in the `gen_component` happens independently later on, when the next ping message arrives.

## Managing execution.

The execution of a `gen_component` can be managed by the following functions:

`gen_component:bp_step(Pid)`: This is the only blocking breakpoint function. It waits until the `gen_component` is in a breakpoint and has handled a single message. It returns the module, the active message handler, and the handled message as a tuple `{Module, On, Message}`. This function does not actually finish the breakpoint, but just lets a single message pass through. For further messages, no breakpoint condition has to be valid, the original breakpoint is still active. To leave a breakpoint, use `gen_component:bp_cont/1`.

`gen_component:bp_cont(Pid)`: Leaves a breakpoint. `gen_component` runs as usual until the next breakpoint is reached.

If no further breakpoints should be entered after continuation, you should delete the registered breakpoint using `gen_component:bp_del/2` before continuing the execution with `gen_component:bp_cont/1`. To ensure, that the breakpoint is entered at least once, `gen_component:bp_barrier/1` should be used before deleting the breakpoint (see the example above). Otherwise it could happen, that the delete request arrives at your `gen_component` before it was actually triggered. The following continuation request would then unintentional apply to an unrelated breakpoint that may be entered later on.

`gen_component:runnable(Pid)`: Returns whether a `gen_component` has messages to handle and is runnable. If you know, that a `gen_component` is in a breakpoint, you can use this to check, whether a `gen_component:bp_step/1` or `gen_component:bp_cont/1` is applicable to the component.

## Tracing handled messages – getting a message interleaving protocol.

We use the debugging interface of `gen_component` to test protocols with random interleaving. First we start all the components involved, set breakpoints on the initialization messages for a new Paxos consensus and then start a single Paxos instance on all of them. The outcome of the Paxos consensus is a `learner_decide` message. So, in `paxos_SUITE:step_until_decide/3` we look for runnable processes and select randomly one of them to perform a single step until the protocol finishes with a decision.

File `paxos_SUITE.erl`:

```
228 -spec prop_rnd_interleave(1..4, 4..16, {pos_integer(), pos_integer(), pos_integer()})
229     -> true.
230 prop_rnd_interleave(NumProposers, NumAcceptors, Seed) ->
231   ct:pal("Called with: paxos_SUITE:prop_rnd_interleave(~p, ~p, ~p).~n",
232     [NumProposers, NumAcceptors, Seed]),
233   Majority = NumAcceptors div 2 + 1,
234   {Proposers, Acceptors, Learners} =
235     make(NumProposers, NumAcceptors, 1, "rnd_interleave"),
236   %% set bp on all processes
237   _ = [ gen_component:bp_set(comm:make_local(X), proposer_initialize, bp)
238     || X <- Proposers ],
239   _ = [ gen_component:bp_set(comm:make_local(X), acceptor_initialize, bp)
240     || X <- Acceptors ],
241   _ = [ gen_component:bp_set(comm:make_local(X), learner_initialize, bp)
242     || X <- Learners ],
243   %% start paxos instances
244   _ = [ proposer:start_paxosid(X, paxidrndinterl, Acceptors,
245     proposal, Majority, NumProposers, Y)
246     || {X,Y} <- lists:zip(Proposers, lists:seq(1, NumProposers)) ],
247   _ = [ acceptor:start_paxosid(X, paxidrndinterl, Learners)
248     || X <- Acceptors ],
```

```

249     _ = [ learner:start_paxosid(X, paxidrndinterl, Majority,
250                               comm:this(), cpaxidrndinterl)
251           || X <- Learners ],
252     %% randomly step through protocol
253     OldSeed = random:seed(Seed),
254     Steps = step_until_decide(Proposers ++ Acceptors ++ Learners, cpaxidrndinterl, 0),
255     ct:pal("Needed ~p steps~n", [Steps]),
256     _ = case OldSeed of
257           undefined -> ok;
258           _ -> random:seed(OldSeed)
259     end,
260     true.
261
262 step_until_decide(Processes, PaxId, SumSteps) ->
263     %% io:format("Step ~p~n", [SumSteps]),
264     Runnable = [ X || X <- Processes, gen_component:runnable(comm:make_local(X)) ],
265     case Runnable of
266     [] ->
267         ct:pal("No runnable processes of ~p~n", [length(Processes)]),
268         timer:sleep(5), step_until_decide(Processes, PaxId, SumSteps);
269     _ -> ok
270     end,
271     Num = random:uniform(length(Runnable)),
272     _ = gen_component:bp_step(comm:make_local(lists:nth(Num, Runnable))),
273     receive
274         {learner_decide, cpaxidrndinterl, _, _Res} = _Any ->
275             %% io:format("Received ~p~n", [_Any]),
276             SumSteps
277     after 0 -> step_until_decide(Processes, PaxId, SumSteps + 1)
278     end.

```

To get a message interleaving protocol, we either can output the results of each `gen_component:bp_step/1` call together with the `Pid` we selected for stepping, or alter the definition of the macro `TRACE_BP_STEPS` in `gen_component`, when we execute all `gen_components` locally in the same Erlang virtual machine.

File `gen_component.erl`:

```

31  %%-define(TRACE_BP_STEPS(X,Y), io:format(X,Y)). %% output on console
32  %%-define(TRACE_BP_STEPS(X,Y), ct:pal(X,Y)).    %% output even if called by unittest
33  -define(TRACE_BP_STEPS(X,Y), ok).

```

### 7.3.11. Future use and planned extensions for `gen_component`

`gen_component` could be further extended. For example it could support hot-code upgrade or could be used to implement algorithms that have to be run across several components of *Scalaris* like snapshot algorithms or similar extensions.

## 7.4. The Process' Database (pdb)

- How to use it and how to switch from `erlang:put/set` to `ets` and implied limitations.

## 7.5. Writing Unittests

### 7.5.1. Plain unittests

### 7.5.2. Randomized Testing using `tester.erl`

## 8. Basic Structured Overlay

### 8.1. Ring Maintenance

### 8.2. T-Man

### 8.3. Routing Tables

*Description is based on SVN revision r1453.*

Each node of the ring can perform searches in the overlay.

A search is done by a lookup in the overlay, but there are several other demands for communication between peers. Scalaris provides a general interface to route a message to the (other) peer, which is currently responsible for a given key.

File `api_dht_raw.erl`:

```
31 -spec unreliable_lookup(Key::?RT:key(), Msg::comm:message()) -> ok.
32 unreliable_lookup(Key, Msg) ->
33     comm:send_local(pid_groups:find_a(dht_node),
34                     {lookup_aux, Key, 0, Msg}).
35
36 -spec unreliable_get_key(Key::?RT:key()) -> ok.
37 unreliable_get_key(Key) ->
38     unreliable_lookup(Key, {get_key, comm:this(), Key}).
39
40 -spec unreliable_get_key(CollectorPid::comm:myid(),
41                         ReqId::{rdht_req_id, pos_integer()},
42                         Key::?RT:key()) -> ok.
43 unreliable_get_key(CollectorPid, ReqId, Key) ->
44     unreliable_lookup(Key, {get_key, CollectorPid, ReqId, Key}).
```

The message `Msg` could be a `get_key` which retrieves content from the responsible node or a `get_node` message, which returns a pointer to the node.

All currently supported messages are listed in the file `dht_node.erl`.

The message routing is implemented in `dht_node_lookup.erl`

File `dht_node_lookup.erl`:

```
27 %% @doc Find the node responsible for Key and send him the message Msg.
28 -spec lookup_aux(State::dht_node_state:state(), Key::intervals:key(),
29                 Hops::non_neg_integer(), Msg::comm:message()) -> ok.
30 lookup_aux(State, Key, Hops, Msg) ->
31     Neighbors = dht_node_state:get(State, neighbors),
32     case intervals:in(Key, nodelist:succ_range(Neighbors)) of
33     true -> % found node -> terminate
34         P = node:pidX(nodelist:succ(Neighbors)),
35         comm:send(P, {lookup_fin, Key, Hops + 1, Msg});
36     ->
37         P = ?RT:next_hop(State, Key),
38         comm:send(P, {lookup_aux, Key, Hops + 1, Msg})
```





```

64     {check, 4}, {check, 5},
65     {check_config, 0}
66 ];

```

empty/1 gets a successor and generates an empty routing table for use inside the routing table implementation. The data structure of the routing table is undefined. It can be a list, a tree, a matrix ...

empty\_ext/1 similarly creates an empty external routing table for use by the dht\_node. This process might not need all the information a routing table implementation requires and can thus work with less data.

hash\_key/1 gets a key and maps it into the overlay's identifier space.

get\_random\_node\_id/0 returns a random node id from the overlay's identifier space. This is used for example when a new node joins the system.

next\_hop/2 gets a dht\_node's state (including the external routing table representation) and a key and returns the node, that should be contacted next when searching for the key, i.e. the known node nearest to the id.

init\_stabilize/2 is called periodically to rebuild the routing table. The parameters are the identifier of the node, its successor and the old (internal) routing table state. This method may send messages to the routing\_table process which need to be handled by the handle\_custom\_message/handler since they are implementation-specific.

update/7 is called when the node's ID, predecessor and/or successor changes. It updates the (internal) routing table with the (new) information.

filter\_dead\_node/2 is called by the failure detector and tells the routing table about dead nodes. This function gets the (internal) routing table and a node to remove from it. A new routing table state is returned.

to\_pid\_list/1 get the PIDs of all (internal) routing table entries.

get\_size/1 get the (internal or external) routing table's size.

get\_replica\_keys/1 Returns for a given (hashed) Key the (hashed) keys of its replicas. This used for implementing symmetric replication.

n/0 gets the number of available keys. An implementation may throw `throw:not_supported` if the operation is unsupported by the routing table.

dump/1 dump the (internal) routing table state for debugging, e.g. by using the web interface. Returns a list of `{Index, Node_as_String}` tuples which may just as well be empty.

to\_list/1 convert the (external) representation of the routing table inside a given dht\_node\_state to a sorted list of known nodes from the routing table, i.e. first=succ, second=next known node on the ring, ... This is used by bulk-operations to create a broadcast tree.

export\_rt\_to\_dht\_node/2 convert the internal routing table state to an external state. Gets the internal state and the node's neighborhood for doing so.

handle\_custom\_message/2 handle messages specific to the routing table implementation. rt\_loop will forward unknown messages to this function.

check/5, check/6 check for routing table changes and send an updated (external) routing table to the dht\_node process.

check\_config/0 check that all required configuration parameters exist and satisfy certain restrictions.

### 8.3.1. The routing table process (rt\_loop)

The rt\_loop module implements the process for all routing tables. It processes messages and calls the appropriate methods in the specific routing table implementations.

File `rt_loop.erl`:

```
40 -opaque(state_active() :: {Neighbors      :: nodelist:neighborhood(),
41                               RTState      :: ?RT:rt(),
42                               TriggerState  :: trigger:state()}).
43 -type(state_inactive() :: {inactive,
44                               MessageQueue::msg_queue:msg_queue(),
45                               TriggerState::trigger:state()}).
46 %% -type(state() :: state_active() | state_inactive()).
```

If initialized, the node's id, its predecessor, successor and the routing table state of the selected implementation (the macro `RT` refers to).

File `rt_loop.erl`:

```
153 on_active({trigger_rt}, {Neighbors, OldRT, TriggerState}) ->
154     % start periodic stabilization
155     % log:log(debug, "[ RT ] stabilize"),
156     NewRT = ?RT:init_stabilize(Neighbors, OldRT),
157     ?RT:check(OldRT, NewRT, Neighbors, true),
158     % trigger next stabilization
159     NewTriggerState = trigger:next(TriggerState),
160     new_state(Neighbors, NewRT, NewTriggerState);
```

Periodically (see `routingtable_trigger` and `pointer_base_stabilization_interval` config parameters) a trigger message is sent to the `rt_loop` process that starts the periodic stabilization implemented by each routing table.

File `rt_loop.erl`:

```
138 % update routing table with changed ID, pred and/or succ
139 on_active({update_rt, OldNeighbors, NewNeighbors}, {_Neighbors, OldRT, TriggerState}) ->
140     case ?RT:update(OldRT, OldNeighbors, NewNeighbors) of
141         {trigger_rebuild, NewRT} ->
142             % trigger immediate rebuild
143             NewTriggerState = trigger:now(TriggerState),
144             ?RT:check(OldRT, NewRT, OldNeighbors, NewNeighbors, true),
145             new_state(NewNeighbors, NewRT, NewTriggerState);
146         {ok, NewRT} ->
147             ?RT:check(OldRT, NewRT, OldNeighbors, NewNeighbors, true),
148             new_state(NewNeighbors, NewRT, TriggerState)
149     end;
```

Every time a node's neighborhood changes, the `dht_node` sends an `update_rt` message to the routing table which will call `?RT:update/7` that decides whether the routing table should be rebuild. If so, it will stop any waiting trigger and schedule an immediate (periodic) stabilization.

### 8.3.2. Simple routing table (`rt_simple`)

One implementation of a routing table is the `rt_simple`, which routes via the successor. Note that this is inefficient as it needs a linear number of hops to reach its goal. A more robust implementation, would use a successor list. This implementation is also not very efficient in the presence of churn.

#### Data types

First, the data structure of the routing table is defined:

File `rt_simple.erl`:

```
26 -type key_t() :: 0..16#FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF. % 128 bit numbers
27 -type rt_t() :: Succ::node:node_type().
28 -type external_rt_t() :: Succ::node:node_type().
29 -type custom_message() :: none().
```

The routing table only consists of a node (the successor). Keys in the overlay are identified by integers  $\geq 0$ .

### A simple `rm_beh` behaviour

File `rt_simple.erl`:

```
41 %% @doc Creates an "empty" routing table containing the successor.
42 empty(Neighbors) -> nodelist:succ(Neighbors).
```

File `rt_simple.erl`:

```
207 empty_ext(Neighbors) -> empty(Neighbors).
```

The empty routing table (internal or external) consists of the successor.

File `rt_simple.erl`:

Keys are hashed using MD5 and have a length of 128 bits.

File `rt_simple.erl`:

```
61 %% @doc Generates a random node id, i.e. a random 128-bit number.
62 get_random_node_id() ->
63     case config:read(key_creator) of
64         random -> hash_key_(randoms:getRandomId());
65         random_with_bit_mask ->
66             {Mask1, Mask2} = config:read(key_creator_bitmask),
67             (hash_key_(randoms:getRandomId()) band Mask2) bor Mask1
68     end.
```

Random node id generation uses the helpers provided by the `randoms` module.

File `rt_simple.erl`:

```
211 %% @doc Returns the next hop to contact for a lookup.
212 next_hop(State, _Key) -> node:pidX(dht_node_state:get(State, rt)).
```

Next hop is always the successor.

File `rt_simple.erl`:

```
76 %% @doc Triggered by a new stabilization round, renews the routing table.
77 init_stabilize(Neighbors, _RT) -> empty(Neighbors).
```

`init_stabilize/2` resets its routing table to the current successor.

File `rt_simple.erl`:

```
81 %% @doc Updates the routing table due to a changed node ID, pred and/or succ.
82 -spec update(OldRT::rt(), OldNeighbors::nodelist:neighborhood(),
83             NewNeighbors::nodelist:neighborhood()) -> {ok, rt()}.
84 update(_OldRT, _OldNeighbors, NewNeighbors) ->
85     {ok, nodelist:succ(NewNeighbors)}.
```

update/7 updates the routing table with the new successor.

File `rt_simple.erl`:

```

89  %% @doc Removes dead nodes from the routing table (rely on periodic
90  %%      stabilization here).
91  filter_dead_node(RT, _DeadPid) -> RT.

```

filter\_dead\_node/2 does nothing, as only the successor is listed in the routing table and that is reset periodically in init\_stabilize/2.

File rt\_simple.erl:

```
95 %% @doc Returns the pids of the routing table entries.
96 to_pid_list(Succ) -> [node:pidX(Succ)].
```

to\_pid\_list/1 returns the pid of the successor.

File `rt_simple.erl`:

```
100 %% @doc Returns the size of the routing table.
101 get_size(_RT) -> 1.
```

The size of the routing table is always 1.

File `rt_simple.erl`:

```

139 %% @doc Returns the replicas of the given key.
140 get_replica_keys(Key) ->
141     [Key,
142      Key bxor 16#40000000000000000000000000000000,
143      Key bxor 16#80000000000000000000000000000000,
144      Key bxor 16#C0000000000000000000000000000000]
145 ]

```

This `get_replica_keys/1` implements symmetric replication.

File rt\_simple.erl:

[illegible]

There are  $2^{128}$  available keys.

File rt\_simple.erl:

```
149 %% @doc Dumps the RT state for output in the web interface.
150 dump(Succ) -> [{"0", lists:flatten(io_lib:format("~p", [Succ]))}].
```

dump/1 lists the successor.

File rt\_simple.erl:

```

222 %% @doc Converts the (external) representation of the routing table to a list
223 %%      in the order of the fingers, i.e. first=succ, second=shortest finger,
224 %%      third=next longer finger,...
225 to_list(State) -> [dht_node_state:get(State, rt)].

```

to\_list/1 lists the successor from the external routing table state.

File `rt_simple.erl`:

```
216 %% @doc Converts the internal RT to the external RT used by the dht_node. Both
217 %%      are the same here.
218 export_rt_to_dht_node(RT, _Neighbors) -> RT.
```

`export_rt_to_dht_node/2` states that the external routing table is the same as the internal table.

File `rt_simple.erl`:

```
168 %% @doc There are no custom messages here.
169 -spec handle_custom_message
170       (custom_message() | any(), rt_loop:state_active()) -> unknown_event.
171 handle_custom_message(_Message, _State) -> unknown_event.
```

Custom messages could be send from a routing table process on one node to the routing table process on another node and are independent from any other implementation.

File `rt_simple.hrl`:

```
175 %% @doc Notifies the dht_node and failure detector if the routing table changed.
176 %%      Provided for convenience (see check/5).
177 check(OldRT, NewRT, Neighbors, ReportToFD) ->
178     check(OldRT, NewRT, Neighbors, Neighbors, ReportToFD).
179
180 %% @doc Notifies the dht_node if the (external) routing table changed.
181 %%      Also updates the failure detector if ReportToFD is set.
182 %%      Note: the external routing table only changes the internal RT has
183 %%      changed.
184 check(OldRT, NewRT, _OldNeighbors, NewNeighbors, ReportToFD) ->
185     case OldRT == NewRT of
186     true -> ok;
187     _ ->
188         Pid = pid_groups:get_my(dht_node),
189         RT_ext = export_rt_to_dht_node(NewRT, NewNeighbors),
190         comm:send_local(Pid, {rt_update, RT_ext}),
191         % update failure detector:
192         case ReportToFD of
193         true ->
194             NewPids = to_pid_list(NewRT),
195             OldPids = to_pid_list(OldRT),
196             fd:update_subscriptions(OldPids, NewPids);
197         _ -> ok
198         end
199     end.
```

Checks whether the routing table changed and in this case sends the `dht_node` an updated (external) routing table state. Optionally the failure detector is updated. This may not be necessary, e.g. if `check` is called after a crashed node has been reported by the failure detector (the failure detector already unsubscribes the node in this case).

### 8.3.3. Chord routing table (`rt_chord`)

The file `rt_chord.erl` implements Chord's routing.

#### Data types

File `rt_chord.erl`:

```
26 -type key_t() :: 0..16#FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF. % 128 bit numbers
27 -type rt_t() :: gb_tree().
```

```

28 -type external_rt_t() :: gb_tree().
29 -type index() :: {pos_integer(), non_neg_integer()}.
30 -opaque custom_message() ::
31     {rt_get_node, Source_PID::comm:mypid(), Index::index()} |
32     {rt_get_node_response, Index::index(), Node::node:node_type()}.

```

The routing table is a `gb_tree`. Identifiers in the ring are integers. Note that in Erlang integer can be of arbitrary precision. For Chord, the identifiers are in  $[0, 2^{128})$ , i.e. 128-bit strings.

### The `rm_beh` behaviour for Chord (excerpt)

File `rt_chord.erl`:

```

44 %% @doc Creates an empty routing table.
45 empty(_Neighbors) -> gb_trees:empty().

```

File `rt_chord.erl`:

```

274 empty_ext(_Neighbors) -> gb_trees:empty().

```

`empty/1` returns an empty `gb_tree`, same for `empty_ext/1`.

`rt_chord:hash_key/1`, `rt_chord:get_random_node_id/0`, `rt_chord:get_replica_keys/1` and `rt_chord:-n/0` are implemented like their counterparts in `rt_simple.erl`.

File `rt_chord.erl`:

```

278 %% @doc Returns the next hop to contact for a lookup.
279 %%     If the routing table has less entries than the rt_size_use_neighbors
280 %%     config parameter, the neighborhood is also searched in order to find a
281 %%     proper next hop.
282 %%     Note, that this code will be called from the dht_node process and
283 %%     it will thus have an external_rt!
284 next_hop(State, Id) ->
285     Neighbors = dht_node_state:get(State, neighbors),
286     case intervals:in(Id, nodelist:succ_range(Neighbors)) of
287         true -> node:pidX(nodelist:succ(Neighbors));
288         _ ->
289             % check routing table:
290             RT = dht_node_state:get(State, rt),
291             RTSize = get_size(RT),
292             NodeRT = case util:gb_trees_largest_smaller_than(Id, RT) of
293                 {value, _Key, N} ->
294                     N;
295                 nil when RTSize == 0 ->
296                     nodelist:succ(Neighbors);
297                 nil -> % forward to largest finger
298                     {_Key, N} = gb_trees:largest(RT),
299                     N
300             end,
301             FinalNode =
302                 case RTSize < config:read(rt_size_use_neighbors) of
303                     false -> NodeRT;
304                     _ ->
305                         % check neighborhood:
306                         nodelist:largest_smaller_than(Neighbors, Id, NodeRT)
307                 end,
308             node:pidX(FinalNode)
309     end.

```

If the (external) routing table contains at least one item, the next hop is retrieved from the `gb_tree`. It will be the node with the largest id that is smaller than the id we are looking for. If the routing

table is empty, the successor is chosen. However, if we haven't found the key in our routing table, the next hop will be our largest finger, i.e. entry.

File `rt_chord.erl`:

```

77 %% @doc Starts the stabilization routine.
78 init_stabilize(Neighbors, RT) ->
79     % calculate the longest finger
80     Id = nodelist:nodeid(Neighbors),
81     Key = calculateKey(Id, first_index()),
82     % trigger a lookup for Key
83     api_dht_raw:unreliable_lookup(Key, {send_to_group_member, routing_table,
84                                         {rt_get_node, comm:this(), first_index()}}),
85     RT.

```

The routing table stabilization is triggered for the first index and then runs asynchronously, as we do not want to block the `rt_loop` to perform other request while recalculating the routing table.

We have to find the node responsible for the calculated finger and therefore perform a lookup for the node with a `rt_get_node` message, including a reference to ourselves as the reply-to address and the index to be set.

The lookup performs an overlay routing by passing the message until the responsible node is found. There, the message is delivered to the `routing_table` process. The remote node sends the requested information back directly. It includes a reference to itself in a `rt_get_node_response` message. Both messages are handled by `rt_chord:handle_custom_message/2`:

File `rt_chord.erl`:

```

219 %% @doc Chord reacts on 'rt_get_node_response' messages in response to its
220 %%      'rt_get_node' messages.
221 -spec handle_custom_message
222       (custom_message(), rt_loop:state_active()) -> rt_loop:state_active();
223       (any(), rt_loop:state_active()) -> unknown_event.
224 handle_custom_message({rt_get_node, Source_PID, Index}, State) ->
225     MyNode = nodelist:node(rt_loop:get_neighb(State)),
226     comm:send(Source_PID, {rt_get_node_response, Index, MyNode}),
227     State;
228 handle_custom_message({rt_get_node_response, Index, Node}, State) ->
229     OldRT = rt_loop:get_rt(State),
230     Id = rt_loop:get_id(State),
231     Succ = rt_loop:get_succ(State),
232     NewRT = stabilize(Id, Succ, OldRT, Index, Node),
233     check(OldRT, NewRT, rt_loop:get_neighb(State), true),
234     rt_loop:set_rt(State, NewRT);
235 handle_custom_message(_Message, _State) ->
236     unknown_event.

```

File `rt_chord.erl`:

```

151 %% @doc Updates one entry in the routing table and triggers the next update.
152 -spec stabilize(MyId::key() | key_t(), Succ::node:node_type(), OldRT::rt(),
153               Index::index(), Node::node:node_type()) -> NewRT::rt().
154 stabilize(Id, Succ, RT, Index, Node) ->
155     case (node:id(Succ) /= node:id(Node)) % reached succ?
156     andalso (not intervals:in(           % there should be nothing shorter
157                 node:id(Node),           % than succ
158                 node:mk_interval_between_ids(Id, node:id(Succ)))) of
159     true ->
160         NewRT = gb_trees:enter(Index, Node, RT),
161         Key = calculateKey(Id, next_index(Index)),
162         Msg = {rt_get_node, comm:this(), next_index(Index)},
163         api_dht_raw:unreliable_lookup(
164             Key, {send_to_group_member, routing_table, Msg}),
165         NewRT;
166     _ -> RT

```



167       end.

stabilize/5 assigns the received routing table entry and triggers the routing table stabilization for the the next shorter entry using the same mechanisms as described above.

If the shortest finger is the successor, then filling the routing table is stopped, as no further new entries would occur. It is not necessary, that Index reaches 1 to make that happen. If less than  $2^{128}$  nodes participate in the system, it may happen earlier.

File `rt_chord.erl`:

```
171 %% @doc Updates the routing table due to a changed node ID, pred and/or succ.
172 -spec update(OldRT::rt(), OldNeighbors::nodelist:neighborhood(),
173             NewNeighbors::nodelist:neighborhood()) -> {trigger_rebuild, rt()}.
174 update(_OldRT, _OldNeighbors, NewNeighbors) ->
175     % to be on the safe side ...
176     {trigger_rebuild, empty(NewNeighbors)}.
```

Tells the `rt_loop` process to rebuild the routing table starting with an empty (internal) routing table state.

File `rt_chord.erl`:

```
89 %% @doc Removes dead nodes from the routing table.
90 filter_dead_node(RT, DeadPid) ->
91     DeadIndices = [Index || {Index, Node} <- gb_trees:to_list(RT),
92                          node:same_process(Node, DeadPid)],
93     lists:foldl(fun(Index, Tree) -> gb_trees:delete(Index, Tree) end,
94                 RT, DeadIndices).
```

`filter_dead_node` removes dead entries from the `gb_tree`.

File `rt_chord.erl`:

```
313 export_rt_to_dht_node(RT, Neighbors) ->
314     Id = nodelist:nodeid(Neighbors),
315     Pred = nodelist:pred(Neighbors),
316     Succ = nodelist:succ(Neighbors),
317     Tree = gb_trees:enter(node:id(Succ), Succ,
318                          gb_trees:enter(node:id(Pred), Pred, gb_trees:empty())),
319     util:gb_trees_foldl(fun (_K, V, Acc) ->
320                         % only store the ring id and the according node structure
321                         case node:id(V) == Id of
322                             true -> Acc;
323                             false -> gb_trees:enter(node:id(V), V, Acc)
324                         end
325                     end, Tree, RT).
```

`export_rt_to_dht_node` converts the internal `gb_tree` structure based on indices into the external representation optimised for look-ups, i.e. a `gb_tree` with node ids and the nodes themselves.

File `rt_chord.hrl`:

```
240 %% @doc Notifies the dht_node and failure detector if the routing table changed.
241 %%      Provided for convenience (see check/5).
242 check(OldRT, NewRT, Neighbors, ReportToFD) ->
243     check(OldRT, NewRT, Neighbors, Neighbors, ReportToFD).
244
245 %% @doc Notifies the dht_node if the (external) routing table changed.
246 %%      Also updates the failure detector if ReportToFD is set.
247 %%      Note: the external routing table also changes if the Pred or Succ
248 %%      change.
249 check(OldRT, NewRT, OldNeighbors, NewNeighbors, ReportToFD) ->
250     case OldRT == NewRT andalso
```

```

251         nodelist:pred(OldNeighbors) == nodelist:pred(NewNeighbors) andalso
252         nodelist:succ(OldNeighbors) == nodelist:succ(NewNeighbors) of
253     true -> ok;
254 - ->
255     Pid = pid_groups:get_my(dht_node),
256     RT_ext = export_rt_to_dht_node(NewRT, NewNeighbors),
257     comm:send_local(Pid, {rt_update, RT_ext}),
258     % update failure detector:
259     case ReportToFD of
260     true ->
261         NewPids = to_pid_list(NewRT),
262         OldPids = to_pid_list(OldRT),
263         fd:update_subscriptions(OldPids, NewPids);
264     - -> ok
265     end
266 end.

```

Checks whether the routing table changed and in this case sends the `dht_node` an updated (external) routing table state. Optionally the failure detector is updated. This may not be necessary, e.g. if check is called after a crashed node has been reported by the failure detector (the failure detector already unsubscribes the node in this case).

## 8.4. Local Datastore

## 8.5. Cyclon

## 8.6. Vivaldi Coordinates

## 8.7. Estimated Global Information (Gossiping)

## 8.8. Load Balancing

## 8.9. Broadcast Trees

## 9. Transactions in Scalaris

### 9.1. The Paxos Module

### 9.2. Transactions using Paxos Commit

### 9.3. Applying the Tx-Modules to replicated DHTs

Introduces transaction processing on top of a Overlay

## 10. How a node joins the system

*Description is based on SVN revision r1370.*

After booting a new Scalaris-System as described in Section 2.5.1 on page 10, ten additional local nodes can be started by typing `admin:add_nodes(10)` in the Erlang-Shell that the boot process opened <sup>1</sup>.

```
scalaris> ./bin/boot.sh
[...]  
(boot@csr-pc9)1> admin:add_nodes(10)
```

In the following we will trace what this function does in order to add additional nodes to the system. The function `admin:add_nodes(pos_integer())` is defined as follows.

File `admin.erl`:

```
37 % @doc add new Scalaris nodes on the local node  
38 -spec add_node_at_id(?RT:key()) ->  
39     ok | {error, already_present} | {already_started, pid() | undefined} | term().  
40 add_node_at_id(Id) ->  
41     add_node([{{dht_node, id}, Id}, {skip_psv_lb}]).  
42  
43 -spec add_node([tuple()]) ->  
44     ok | {error, already_present} | {already_started, pid() | undefined} | term().  
45 add_node(Options) ->  
46     DhtNodeId = randoms:getRandomId(),  
47     Desc = util:sup_supervisor_desc(  
48         DhtNodeId, config:read(dht_node_sup), start_link,  
49         [{my_sup_dht_node_id, DhtNodeId} | Options]),  
50     case supervisor:start_child(main_sup, Desc) of  
51         {ok, _Child} -> ok;  
52         {ok, _Child, _Info} -> ok;  
53         {error, _Error} = X -> X  
54     end.  
55  
56 -spec add_nodes(non_neg_integer()) ->  
57     nothing_to_do | [ok | {error, already_present} |  
58         {already_started, pid() | undefined} | term()],...].  
59 add_nodes(0) -> nothing_to_do;  
60 add_nodes(Count) ->  
61     [add_node([]) || _X <- lists:seq(1, Count)].
```

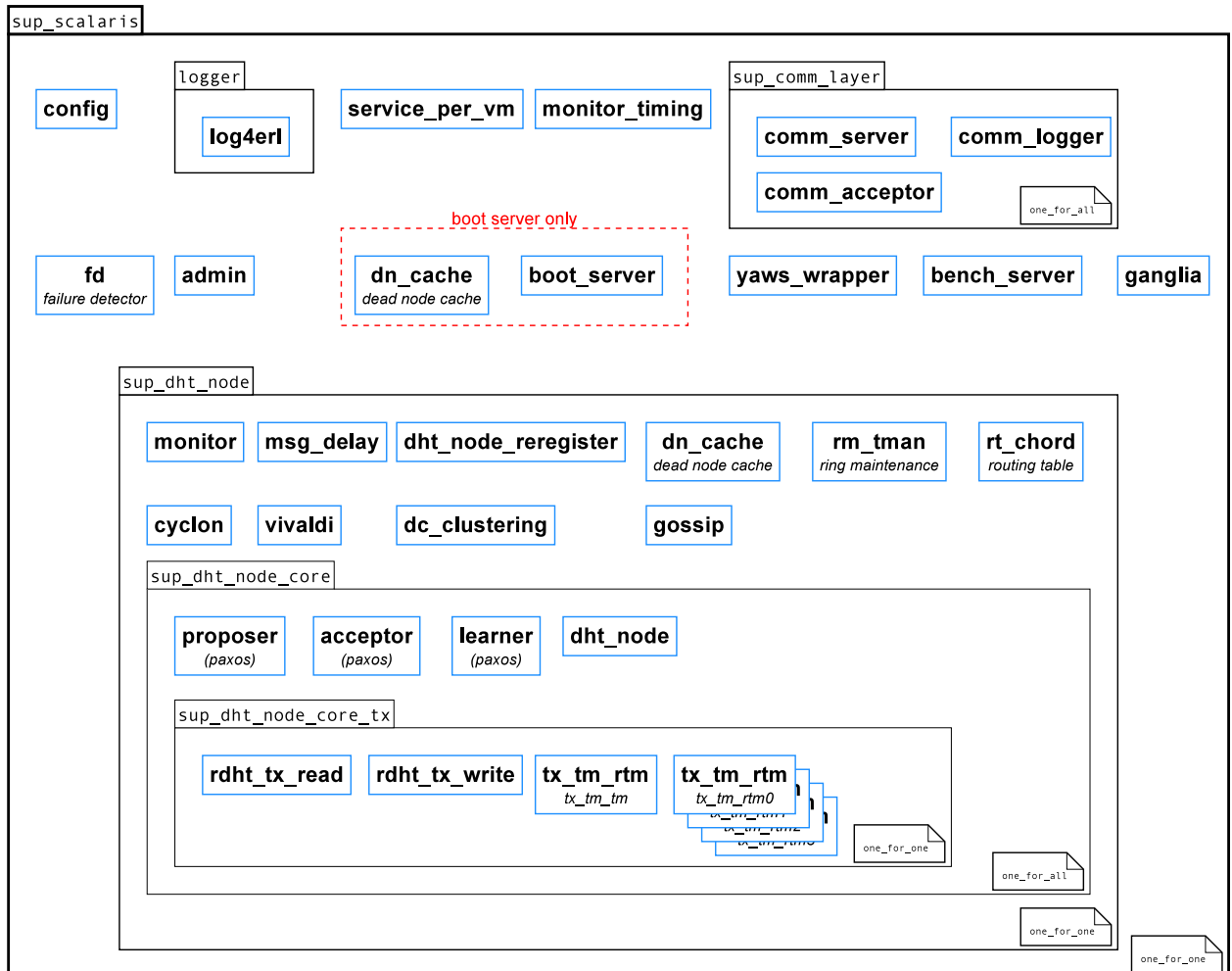
It calls `admin:add_node([])` Count times. This function starts a new child with the given options for the main supervisor `main_sup`. In particular, it sets a random ID that is passed to the new node as its suggested ID to join at. To actually perform the start, the function `sup_dht_node:start_link/1` is called by the Erlang supervisor mechanism. For more details on the OTP supervisor mechanism see Chapter 18 of the Erlang book [1] or the online documentation at <http://www.erlang.org/doc/man/supervisor.html>.

---

<sup>1</sup>Increase the log level to `info` to get more detailed startup logs. See Section 2.7 on page 12

## 10.1. Supervisor-tree of a Scalaris node

When a new Erlang VM with a Scalaris node is started, a `sup_scalaris` supervisor is started that creates further workers and supervisors according to the following scheme (processes starting order: left to right, top to bottom):



When new nodes are started using `admin:add_node/1`, only new `sup_dht_node` supervisors are started.

## 10.2. Starting the `sup_dht_node` supervisor and general processes of a node

Starting supervisors is a two step process: a call to `supervisor:start_link/2,3`, e.g. from a custom supervisor's own `start_link` method, will start the supervisor process. It will then call `Module:init/1` to find out about the restart strategy, maximum restart frequency and child processes. Note that `supervisor:start_link/2,3` will not return until `Module:init/1` has returned and all child processes have been started.

Let's have a look at `sup_dht_node:init/1`, the 'DHT node supervisor'.

File `sup_dht_node.erl`:

```
44 -spec init([tuple()]) -> {ok, {{one_for_one, MaxRetries::pos_integer(),
45                               PeriodInSeconds::pos_integer()},
46                               [ProcessDescr::any()]} }.
47 init(Options) ->
48     DHTNodeGroup = pid_groups:new("dht_node_"),
49     pid_groups:join_as(DHTNodeGroup, ?MODULE),
50     boot_server:connect(),
51
52     Cyclon = util:sup_worker_desc(cyclon, cyclon, start_link, [DHTNodeGroup]),
53     DC_Clustering =
54         util:sup_worker_desc(dc_clustering, dc_clustering, start_link,
55                             [DHTNodeGroup]),
56     DeadNodeCache =
57         util:sup_worker_desc(deadnodecache, dn_cache, start_link,
58                             [DHTNodeGroup]),
59     Delayer =
60         util:sup_worker_desc(msg_delay, msg_delay, start_link,
61                             [DHTNodeGroup]),
62     Gossip =
63         util:sup_worker_desc(gossip, gossip, start_link, [DHTNodeGroup]),
64     Reregister =
65         util:sup_worker_desc(dht_node_reregister, dht_node_reregister,
66                             start_link, [DHTNodeGroup]),
67     RoutingTable =
68         util:sup_worker_desc(routing_table, rt_loop, start_link,
69                             [DHTNodeGroup]),
70     SupDHTNodeCore_AND =
71         util:sup_supervisor_desc(sup_dht_node_core, sup_dht_node_core,
72                                 start_link, [DHTNodeGroup, Options]),
73     Vivaldi =
74         util:sup_worker_desc(vivaldi, vivaldi, start_link, [DHTNodeGroup]),
75     Monitor =
76         util:sup_worker_desc(monitor, monitor, start_link, [DHTNodeGroup]),
77
78     %% order in the following list is the start order
79     {ok, {{one_for_one, 10, 1},
80         [
81             Monitor,
82             Delayer,
83             Reregister,
84             DeadNodeCache,
85             RoutingTable,
86             Cyclon,
87             Vivaldi,
88             DC_Clustering,
89             Gossip,
90             SupDHTNodeCore_AND
91         ]}}.
```

The return value of the `init/1` function specifies the child processes of the supervisor and how to start them. Here, we define a list of processes to be observed by a `one_for_one` supervisor. The processes are: `Monitor`, `Delayer`, `Reregister`, `DeadNodeCache`, `RingMaintenance`, `RoutingTable`, `Cyclon`, `Vivaldi`, `DC_Clustering`, `Gossip` and a `SupDHTNodeCore_AND` process in this order.

The term `{one_for_one, 10, 1}` specifies that the supervisor should try 10 times to restart each process before giving up. `one_for_one` supervision means, that if a single process stops, only that process is restarted. The other processes run independently.

When the `sup_dht_node:init/1` is finished the supervisor module starts all the defined processes by calling the functions that were defined in the returned list.

For a join of a new node, we are only interested in the starting of the `SupDHTNodeCore_AND` process here. At that point in time, all other defined processes are already started and running.

## 10.3. Starting the `sup_dht_node_core` supervisor with a peer and some paxos processes

Like any other supervisor the `sup_dht_node_core` supervisor calls its `sup_dht_node_core:init/1` function:

File `sup_dht_node_core.erl`:

```
40 -spec init({pid_groups:groupname(), Options::[tuple()]}) ->
41     {ok, {{one_for_all, MaxRetries::pos_integer(),
42         PeriodInSeconds::pos_integer()},
43         [ProcessDescr::any()]}}.
44 init({DHTNodeGroup, Options}) ->
45     pid_groups:join_as(DHTNodeGroup, ?MODULE),
46     PaxosProcesses = util:sup_supervisor_desc(sup_paxos, sup_paxos,
47         start_link, [DHTNodeGroup, []]),
48     DHTNodeModule = config:read(dht_node),
49     DHTNode = util:sup_worker_desc(dht_node, DHTNodeModule, start_link,
50         [DHTNodeGroup, Options]),
51     TX =
52         util:sup_supervisor_desc(sup_dht_node_core_tx, sup_dht_node_core_tx, start_link,
53             [DHTNodeGroup]),
54     {ok, {{one_for_all, 10, 1},
55         [
56             PaxosProcesses,
57             DHTNode,
58             TX
59         ]}}.
```

It defines five processes, that have to be observed using a `one_for_all`-supervisor, which means, that if one fails, all have to be restarted. The `dht_node` module implements the main component of a full Scalaris node which glues together all the other processes. Its `dht_node:start_link/2` function will get the following parameters: (a) the processes' group that is used with the `pid_groups` module and (b) a list of options for the `dht_node`. The process group name was calculated a bit earlier in the code. *Exercise: Try to find where.*

File `dht_node.erl`:

```
386 %% @doc spawns a scalaris node, called by the scalaris supervisor process
387 -spec start_link(pid_groups:groupname(), [tuple()]) -> {ok, pid()}.
388 start_link(DHTNodeGroup, Options) ->
389     gen_component:start_link(?MODULE, Options,
390         [{pid_groups_join_as, DHTNodeGroup, dht_node}, wait_for_init]).
```

Like many other modules, the `dht_node` module implements the `gen_component` behaviour. This behaviour was developed by us to enable us to write code which is similar in syntax and semantics to the examples in [3]. Similar to the supervisor behaviour, a module implementing this behaviour has to provide an `init/1` function, but here it is used to initialize the state of the component. This function is described in the next section.

Note: `?MODULE` is a predefined Erlang macro, which expands to the module name, the code belongs to (here: `dht_node`).

## 10.4. Initializing a `dht_node`-process

File `dht_node.erl`:

```
369 %% @doc joins this node in the ring and calls the main loop
```

```

370 -spec init(Options::[tuple()]) -> dht_node_state:state().
371 init(Options) ->
372     {my_sup_dht_node_id, MySupDhtNode} = lists:keyfind(my_sup_dht_node_id, 1, Options),
373     erlang:put(my_sup_dht_node_id, MySupDhtNode),
374     % get my ID (if set, otherwise chose a random ID):
375     Id = case lists:keyfind({dht_node, id}, 1, Options) of
376         {{dht_node, id}, IdX} -> IdX;
377         _ -> ?RT:get_random_node_id()
378     end,
379     case is_first(Options) of
380         true -> dht_node_join:join_as_first(Id, 0, Options);
381         _ -> dht_node_join:join_as_other(Id, 0, Options)
382     end.

```

The `gen_component` behaviour registers the `dht_node` in the process dictionary. Formerly, the process had to do this itself, but we moved this code into the behaviour. If an ID was given to `dht_node:init/1` function as a `{{dht_node, id}, KEY}` tuple, the given `Id` will be used. Otherwise a random key is generated. Depending on whether the node is the first inside a VM marked as first or not, the according function in `dht_node_join` is called. Also the pid of the node's supervisor is kept for future reference.

## 10.5. Actually joining the ring

After retrieving its identifier, the node starts the join protocol which processes the appropriate messages calling `dht_node_join:process_join_state(Message, State)`. On the existing node, join messages will be processed by `dht_node_join:process_join_msg(Message, State)`.

### 10.5.1. A single node joining an empty ring

File `dht_node_join.erl`:

```

98 -spec join_as_first(Id::?RT:key(), IdVersion::non_neg_integer(), Options::[tuple()])
99     -> dht_node_state:state().
100 join_as_first(Id, IdVersion, _Options) ->
101     % ugly hack to get a valid ip-address into the comm-layer
102     dht_node:trigger_known_nodes(),
103     log:log(info, "[ Node ~w ] joining as first: (~.0p, ~.0p)",
104         [self(), Id, IdVersion]),
105     Me = node:new(comm:this(), Id, IdVersion),
106     % join complete, State is the first "State"
107     finish_join(Me, Me, Me, ?DB:new(), msg_queue:new()).

```

If the ring is empty, the joining node will be the only node in the ring and will thus be responsible for the whole key space. It will trigger all known nodes to initialize the comm layer and then finish the join. `dht_node_join:finish_join/5` just creates a new state for a Scalaris node consisting of the given parameters (the node as itself, its predecessor and successor, an empty database and the queued messages that arrived during the join). It then activates all dependent processes and creates a routing table from this information.

The `dht_node_state:state()` type is defined in

File `dht_node_state.erl`:

```

50 -record(state, {rt          = ?required(state, rt)          :: ?RT:external_rt(),
51                rm_state    = ?required(state, rm_state)    :: rm_loop:state(),
52                join_time    = ?required(state, join_time)   :: util:time(),
53                db           = ?required(state, db)           :: ?DB:db(),

```



```

54         tx_tp_db = ?required(state, tx_tp_db) :: any(),
55         proposer = ?required(state, proposer) :: pid(),
56         % slide with pred (must not overlap with 'slide with succ'):
57         slide_pred = null :: slide_op:slide_op() | null,
58         % slide with succ (must not overlap with 'slide with pred'):
59         slide_succ = null :: slide_op:slide_op() | null,
60         msg_fwd = [] :: [{intervals:interval(), comm:mypid()}],
61         % additional range to respond to during a move:
62         db_range = [] :: [{intervals:interval(), slide_op:id()}]
63     }).
64 -opaque state() :: #state{}.

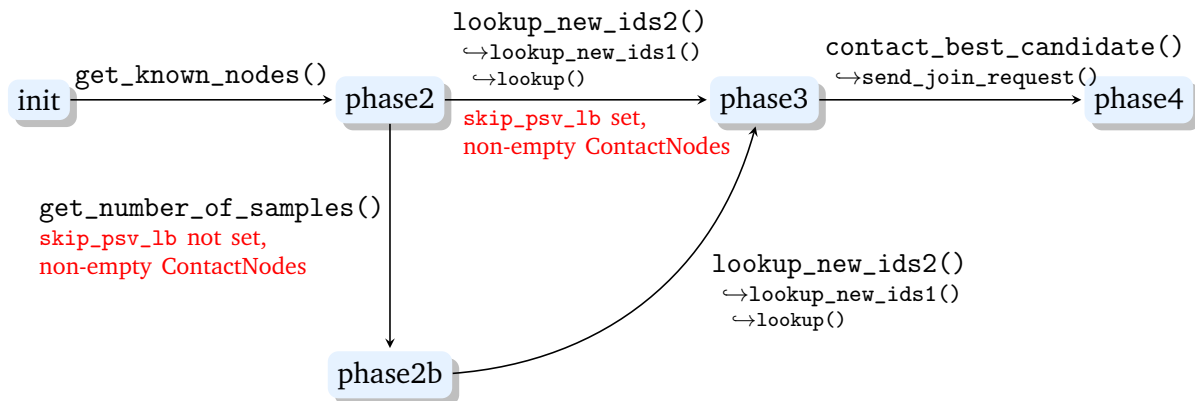
```

### 10.5.2. A single node joining an existing (non-empty) ring

If a node joins an existing ring, its join protocol will step through the following four phases:

- **phase2** finding nodes to contact with the help of the configured `known_hosts`
- **phase2b** getting the number of Ids to sample (may be skipped)
- **phase3** lookup nodes responsible for all sampled Ids
- **phase4** joining a selected node and setting up item movements

The following figure shows a (non-exhaustive) overview of the transitions between the phases in the normal case. We will go through these step by step and discuss what happens if errors occur.



At first all nodes set in the `known_hosts` configuration parameter are contacted. Their responses are then handled in phase 2. In order to separate the join state from the ordinary `dht_node` state, the `gen_component` is instructed to use the `dht_node:on_join/2` message handler which delegates every message to `dht_node_join:process_join_state/2`.

File `dht_node_join.erl`:

```

111 -spec join_as_other(Id::?RT:key(), IdVersion::non_neg_integer(), Options::tuple())
112     -> {'$gen_component', [{on_handler, Handler::on_join}],
113         State::{join, phase2(), msg_queue:msg_queue()}}.
114 join_as_other(Id, IdVersion, Options) ->
115     log:log(info, "[ Node ~w ] joining, trying ID: (~.0p, ~.0p)",
116         [self(), Id, IdVersion]),
117     get_known_nodes(util:get_pids_uid()),
118     JoinUUID = util:get_pids_uid(),
119     msg_delay:send_local(get_join_timeout() div 1000, self(),
120         {join, timeout, JoinUUID}),
121     gen_component:change_handler(
122         {join, {phase2, JoinUUID, Options, IdVersion, [], [Id], []},
123             msg_queue:new()},
124         on_join).

```

## Phase 2 and 2b

Phase 2 collects all `dht_node` processes inside the contacted VMs. It therefore mainly processes `get_dht_nodes_response` messages and integrates all received nodes into the list of available connections. The next step depends on whether the `{skip_psv_lb}` option for skipping any passive load balancing algorithm has been given to the `dht_node` or not. If it is present, the node will only use the ID that has been initially passed to `dht_node_join:join_as_other/3`, issue a lookup for the responsible node and move to phase 3. Otherwise, the passive load balancing's `lb_psv_*:-get_number_of_samples/1` method will be called asking for the number of IDs to sample. Its answer will be processed in phase 2b.

`get_dht_nodes_response` messages arriving in phase 2b or later will be processed anyway and received `dht_node` processes will be integrated into the connections. These phases' operations will not be interrupted and nothing else is changed though.

File `dht_node_join.erl`:

```
152 % in phase 2 add the nodes and do lookups with them / get number of samples
153 process_join_state({get_dht_nodes_response, Nodes} = _Msg,
154                   {join, JoinState, QueuedMessages})
155   when element(1, JoinState) == phase2 ->
156     ?TRACE_JOIN1(_Msg, JoinState),
157     Connections = [{null, Node} || Node <- Nodes, Node /= comm:this()],
158     JoinState1 = add_connections(Connections, JoinState, back),
159     NewJoinState = phase2_next_step(JoinState1, Connections),
160     ?TRACE_JOIN_STATE(NewJoinState),
161     {join, NewJoinState, QueuedMessages};
162
163 % in all other phases, just add the provided nodes:
164 process_join_state({get_dht_nodes_response, Nodes} = _Msg,
165                   {join, JoinState, QueuedMessages})
166   when element(1, JoinState) == phase2b orelse
167     element(1, JoinState) == phase3 orelse
168     element(1, JoinState) == phase4 ->
169     ?TRACE_JOIN1(_Msg, JoinState),
170     Connections = [{null, Node} || Node <- Nodes, Node /= comm:this()],
171     JoinState1 = add_connections(Connections, JoinState, back),
172     ?TRACE_JOIN_STATE(JoinState1),
173     {join, JoinState1, QueuedMessages};
```

Phase 2b will handle `get_number_of_samples` messages from the passive load balance algorithm. Once received, new (unique) IDs will be sampled randomly so that the total number of join candidates (selected IDs together with fully processed candidates from further phases) is at least as high as the given number of samples. Afterwards, lookups will be created for all previous IDs as well as the new ones and the node will move to phase 3.

File `dht_node_join.erl`:

```
199 % note: although this message was send in phase2, also accept message in
200 % phase2, e.g. messages arriving from previous calls
201 process_join_state({join, get_number_of_samples, Samples, Conn} = _Msg,
202                   {join, JoinState, QueuedMessages})
203   when element(1, JoinState) == phase2 orelse
204     element(1, JoinState) == phase2b ->
205     ?TRACE_JOIN1(_Msg, JoinState),
206     % prefer node that send get_number_of_samples as first contact node
207     JoinState1 = reset_connection(Conn, JoinState),
208     % (re-)issue lookups for all existing IDs and
209     % create additional samples, if required
210     NewJoinState = lookup_new_ids2(Samples, JoinState1),
211     ?TRACE_JOIN_STATE(NewJoinState),
212     {join, NewJoinState, QueuedMessages};
213
214 % ignore message arriving in other phases:
```

```

215 process_join_state({join, get_number_of_samples, _Samples, Conn} = _Msg,
216                    {join, JoinState, QueuedMessages}) ->
217     ?TRACE_JOIN1(_Msg, JoinState),
218     NewJoinState = reset_connection(Conn, JoinState),
219     ?TRACE_JOIN_STATE(NewJoinState),
220     {join, NewJoinState, QueuedMessages};

```

Lookups will make Scalaris find the node currently responsible for a given ID and send a request to simulate a join to this node, i.e. a `get_candidate` message. Note that during such an operation, the joining node would become the existing node's predecessor. The simulation will be delegated to the passive load balance algorithm the joining node requested, as set by the `join_lb_psv` configuration parameter.

File `dht_node_join.erl`:

```

481 process_join_msg({join, get_candidate, Source_PID, Key, LbPsv, Conn} = _Msg, State) ->
482     ?TRACE1(_Msg, State),
483     LbPsv:create_join(State, Key, Source_PID, Conn);

```

### Phase 3

The result of the simulation will be send in a `get_candidate_response` message and will be processed in phase 3 of the joining node. It will be integrated into the list of processed candidates. If there are no more IDs left to process, the best among them will be contacted. Otherwise further `get_candidate_response` messages will be awaited. Such messages will also be processed in the other phases where the candidate will be simply added to the list.

File `dht_node_join.erl`:

```

252 process_join_state({join, get_candidate_response, OrigJoinId, Candidate, Conn} = _Msg,
253                    {join, JoinState, QueuedMessages})
254     when element(1, JoinState) == phase3 ->
255     ?TRACE_JOIN1(_Msg, JoinState),
256     JoinState0 = reset_connection(Conn, JoinState),
257     JoinState1 = remove_join_id(OrigJoinId, JoinState0),
258     JoinState2 = integrate_candidate(Candidate, JoinState1, front),
259     NewJoinState =
260         case get_join_ids(JoinState2) of
261             [] -> % no more join ids to look up -> join with the best:
262                 contact_best_candidate(JoinState2);
263             [_|_] -> % still some unprocessed join ids -> wait
264                 JoinState2
265         end,
266     ?TRACE_JOIN_STATE(NewJoinState),
267     {join, NewJoinState, QueuedMessages};
268
269 % In phase 2 or 2b, also add the candidate but do not continue.
270 % In phase 4, add the candidate to the end of the candidates as they are sorted
271 % and the join with the first has already started (use this candidate as backup
272 % if the join fails). Do not start a new join.
273 process_join_state({join, get_candidate_response, OrigJoinId, Candidate, Conn} = _Msg,
274                    {join, JoinState, QueuedMessages})
275     when element(1, JoinState) == phase2 orelse
276         element(1, JoinState) == phase2b orelse
277         element(1, JoinState) == phase4 ->
278     ?TRACE_JOIN1(_Msg, JoinState),
279     JoinState0 = reset_connection(Conn, JoinState),
280     JoinState1 = remove_join_id(OrigJoinId, JoinState0),
281     JoinState2 = case get_phase(JoinState1) of
282                     phase4 -> integrate_candidate(Candidate, JoinState1, back);
283                     _ -> integrate_candidate(Candidate, JoinState1, front)
284                 end,
285     ?TRACE_JOIN_STATE(JoinState2),

```

```
286 {join, JoinState2, QueuedMessages};
```

If `dht_node_join:contact_best_candidate/1` is called and candidates are available (there should be at this stage!), it will sort the candidates by using the passive load balance algorithm, send a `join_request` message and continue with phase 4.

File `dht_node_join.erl`:

```
784 %% @doc Contacts the best candidate among all stored candidates and sends a
785 %%      join_request (Timeouts = 0).
786 -spec contact_best_candidate(JoinState::phase_2_4())
787     -> phase2() | phase2b() | phase4().
788 contact_best_candidate(JoinState) ->
789     contact_best_candidate(JoinState, 0).
790 %% @doc Contacts the best candidate among all stored candidates and sends a
791 %%      join_request. Timeouts is the number of join_request_timeout messages
792 %%      previously received.
793 -spec contact_best_candidate(JoinState::phase_2_4(), Timeouts::non_neg_integer())
794     -> phase2() | phase2b() | phase4().
795 contact_best_candidate(JoinState, Timeouts) ->
796     JoinState1 = sort_candidates(JoinState),
797     send_join_request(JoinState1, Timeouts).
```

File `dht_node_join.erl`:

```
801 %% @doc Sends a join request to the first candidate. Timeouts is the number of
802 %%      join_request_timeout messages previously received.
803 %%      PreCond: the id has been set to the ID to join at and has been updated
804 %%      in JoinState.
805 -spec send_join_request(JoinState::phase_2_4(), Timeouts::non_neg_integer())
806     -> phase2() | phase2b() | phase4().
807 send_join_request(JoinState, Timeouts) ->
808     case get_candidates(JoinState) of
809     [] -> % no candidates -> start over (should not happen):
810         start_over(JoinState);
811     [BestCand | _] ->
812         Id = node_details:get(lb_op:get(BestCand, n1_new), new_key),
813         IdVersion = get_id_version(JoinState),
814         NewSucc = node_details:get(lb_op:get(BestCand, n1succ_new), node),
815         Me = node:new(comm:this(), Id, IdVersion),
816         CandId = lb_op:get(BestCand, id),
817         ?TRACE_SEND(node:pidX(NewSucc), {join, join_request, Me, CandId}),
818         comm:send(node:pidX(NewSucc), {join, join_request, Me, CandId}),
819         msg_delay:send_local(
820             get_join_request_timeout() div 1000, self(),
821             {join, join_request_timeout, Timeouts, CandId, get_join_uuid(JoinState)}),
822         set_phase(phase4, JoinState)
823     end.
```

The `join_request` message will be received by the existing node which will set up a slide operation with the new node. If it is not responsible for the key (anymore), it will deny the request and reply with a `{join, join_response, not_responsible, Node}` message.

File `dht_node_join.erl`:

```
493 process_join_msg({join, join_request, NewPred, CandId} = _Msg, State)
494     when (not is_atom(NewPred)) -> % avoid confusion with not_responsible message
495     ?TRACE1(_Msg, State),
496     TargetId = node:id(NewPred),
497     case dht_node_move:can_slide_pred(State, TargetId, {join, 'rcv'}) of
498     true ->
499         try
500             % TODO: implement step-wise join
501             MoveFullId = util:get_global_uid(),
502             Neighbors = dht_node_state:get(State, neighbors),
503             SlideOp = slide_op:new_sending_slide_join(
```

```

504             MoveFullId, NewPred, join, Neighbors),
505             SlideOp1 = slide_op:set_phase(SlideOp, wait_for_pred_update_join),
506             RMSubscrTag = {move, slide_op:get_id(SlideOp1)},
507             rm_loop:subscribe(self(), RMSubscrTag,
508                             fun(_OldNeighbors, NewNeighbors) ->
509                                 NewPred := nodelist:pred(NewNeighbors)
510                             end,
511                             fun dht_node_move:rm_notify_new_pred/4, 1),
512             State1 = dht_node_state:add_db_range(
513                 State, slide_op:get_interval(SlideOp1),
514                 slide_op:get_id(SlideOp1)),
515             send_join_response(State1, SlideOp1, NewPred, CandId)
516         catch throw:not_responsible ->
517             ?TRACE_SEND(node:pidX(NewPred),
518                         {join, join_response, not_responsible, CandId}),
519             comm:send(node:pidX(NewPred),
520                     {join, join_response, not_responsible, CandId}),
521             State
522         end;
523     - ->
524         ?TRACE([" ~.Op ]~n ignoring join_request from ~.Op due to a running slide~n",
525               [self(), NewPred]),
526         State
527     end;

```

If it is responsible for the ID and is not participating in a slide with its current predecessor, it will set up a slide with the joining node:

File dht\_node\_join.erl:

```

860 -spec send_join_response(State::dht_node_state:state(),
861                         NewSlideOp::slide_op:slide_op(),
862                         NewPred::node:node_type(), CandId::lb_op:id())
863     -> dht_node_state:state().
864 send_join_response(State, SlideOp, NewPred, CandId) ->
865     MoveFullId = slide_op:get_id(SlideOp),
866     NewSlideOp =
867         slide_op:set_timer(SlideOp, get_join_response_timeout(),
868                             {join, join_response_timeout, NewPred, MoveFullId, CandId}),
869     MyOldPred = dht_node_state:get(State, pred),
870     MyNode = dht_node_state:get(State, node),
871     ?TRACE_SEND(node:pidX(NewPred),
872                 {join, join_response, MyNode, MyOldPred, MoveFullId, CandId}),
873     comm:send(node:pidX(NewPred),
874               {join, join_response, MyNode, MyOldPred, MoveFullId, CandId}),
875     % no need to tell the ring maintenance -> the other node will trigger an update
876     % also this is better in case the other node dies during the join
877     %% rm_loop:notify_new_pred(comm:this(), NewPred),
878     dht_node_state:set_slide(State, pred, NewSlideOp).

```

## Phase 4

The joining node will receive the join\_response message in phase 4 of the join protocol. If everything is ok, it will notify its ring maintenance process that it enters the ring, start all required processes and join the slide operation set up by the existing node in order to receive some of its data.

If the join candidate's node is not responsible for the candidate's ID anymore or the candidate's ID already exists, the next candidate is contacted until no further candidates are available and the join protocol starts over using dht\_node\_join:start\_over/1.

Note that the join\_response message will actually be processed in any phase. Therefore, if messages arrive late, the join can be processed immediately and the rest of the join protocol does not

need to be executed again.

File dht\_node\_join.erl:

```
325 process_join_state({join, join_response, not_responsible, CandId} = _Msg,
326                   {join, JoinState, QueuedMessages} = State)
327   when element(1, JoinState) == phase4 ->
328     ?TRACE_JOIN1(_Msg, JoinState),
329     % the node we contacted is not responsible for the selected key anymore
330     % -> try the next candidate, if the message is related to the current candidate
331     case get_candidates(JoinState) of
332       [] -> % no candidates -> should not happen in phase4!
333         log:log(error, "[ Node ~w ] empty candidate list in join phase 4, "
334                      "starting over", [self()]),
335         NewJoinState = start_over(JoinState),
336         ?TRACE_JOIN_STATE(NewJoinState),
337         {join, NewJoinState, QueuedMessages};
338       [Candidate | _Rest] ->
339         case lb_op:get(Candidate, id) == CandId of
340           false -> State; % unrelated/old message
341           _ ->
342             log:log(info,
343                    "[ Node ~w ] node contacted for join is not responsible "
344                    "for the selected ID (anymore), trying next candidate",
345                    [self()]),
346             NewJoinState = try_next_candidate(JoinState),
347             ?TRACE_JOIN_STATE(NewJoinState),
348             {join, NewJoinState, QueuedMessages}
349         end
350     end;
351
352 % in other phases remove the candidate from the list (if it still exists):
353 process_join_state({join, join_response, not_responsible, CandId} = _Msg,
354                   {join, JoinState, QueuedMessages}) ->
355   ?TRACE_JOIN1(_Msg, JoinState),
356   {join, remove_candidate(CandId, JoinState), QueuedMessages};
357
358 % note: accept (delayed) join_response messages in any phase
359 process_join_state({join, join_response, Succ, Pred, MoveId, CandId} = _Msg,
360                   {join, JoinState, QueuedMessages} = State) ->
361   ?TRACE_JOIN1(_Msg, JoinState),
362   % only act on related messages, i.e. messages from the current candidate
363   Phase = get_phase(JoinState),
364   State1 = case get_candidates(JoinState) of
365     [] when Phase == phase4 -> % no candidates -> should not happen in phase4!
366       log:log(error, "[ Node ~w ] empty candidate list in join phase 4, "
367                    "starting over", [self()]),
368       NewJoinState = start_over(JoinState),
369       ?TRACE_JOIN_STATE(NewJoinState),
370       {join, NewJoinState, QueuedMessages};
371     [Candidate | _Rest] ->
372       % in all other phases, ignore the delayed join_response
373       % if no candidates exist
374       [CandidateNode = node_details:get(lb_op:get(Candidate, n1succ_new), node),
375        CandidateNodeSame = node:same_process(CandidateNode, Succ),
376        case lb_op:get(Candidate, id) == CandId of
377          false ->
378            log:log(warn, "[ Node ~w ] ignoring old or unrelated "
379                       "join_response message", [self()]),
380            State; % ignore old/unrelated message
381          _ when not CandidateNodeSame ->
382            % id is correct but the node is not (should never happen!)
383            log:log(error, "[ Node ~w ] got join_response but the node "
384                       "changed, trying next candidate", [self()]),
385            NewJoinState = try_next_candidate(JoinState),
386            ?TRACE_JOIN_STATE(NewJoinState),
387            {join, NewJoinState, QueuedMessages};
388          _ ->
389            MyId = node_details:get(lb_op:get(Candidate, n1_new), new_key),
390            MyIdVersion = get_id_version(JoinState),
391            case MyId == node:id(Succ) orelse MyId == node:id(Pred) of
```

```

392         true ->
393             log:log(warn, "[ Node ~w ] chosen ID already exists, "
394                 "trying next candidate", [self()]),
395             % note: can not keep Id, even if skip_psv_lb is set
396             JoinState1 = remove_candidate_front(JoinState),
397             NewJoinState = contact_best_candidate(JoinState1),
398             ?TRACE_JOIN_STATE(NewJoinState),
399             {join, NewJoinState, QueuedMessages};
400         ->
401             ?TRACE("[ ~.0p ]~n joined MyId:~.0p, MyIdVersion:~.0p~n "
402                 "Succ: ~.0p~n Pred: ~.0p~n",
403                 [self(), MyId, MyIdVersion, Succ, Pred]),
404             Me = node:new(comm:this(), MyId, MyIdVersion),
405             log:log(info, "[ Node ~w ] joined between ~w and ~w",
406                 [self(), Pred, Succ]),
407             rm_loop:notify_new_succ(node:pidX(Pred), Me),
408             rm_loop:notify_new_pred(node:pidX(Succ), Me),
409
410             finish_join_and_slide(Me, Pred, Succ, ?DB:new(),
411                 QueuedMessages, MoveId)
412         end
413     end
414 end,
415 State1;

```

File dht\_node\_join.erl:

```

882 %% @doc Finishes the join and sends all queued messages.
883 -spec finish_join(Me::node:node_type(), Pred::node:node_type(),
884     Succ::node:node_type(), DB::?DB:db(),
885     QueuedMessages::msg_queue:msg_queue())
886     -> dht_node_state:state().
887 finish_join(Me, Pred, Succ, DB, QueuedMessages) ->
888     RMState = rm_loop:init(Me, Pred, Succ),
889     Neighbors = rm_loop:get_neighbors(RMState),
890     % wait for the ring maintenance to initialize and tell us its table ID
891     rt_loop:activate(Neighbors),
892     cyclon:activate(),
893     vivaldi:activate(),
894     dc_clustering:activate(),
895     gossip:activate(node:mk_interval_between_nodes(Pred, Me)),
896     dht_node_reregister:activate(),
897     msg_queue:send(QueuedMessages),
898     NewRT_ext = ?RT:empty_ext(Neighbors),
899     dht_node_state:new(NewRT_ext, RMState, DB).
900
901 %% @doc Finishes the join by setting up a slide operation to get the data from
902 %% the other node and sends all queued messages.
903 -spec finish_join_and_slide(Me::node:node_type(), Pred::node:node_type(),
904     Succ::node:node_type(), DB::?DB:db(),
905     QueuedMessages::msg_queue:msg_queue(), MoveId::slide_op:id())
906     -> {'$gen_component', [{on_handler, Handler::on}],
907         State::dht_node_state:state()}.
908 finish_join_and_slide(Me, Pred, Succ, DB, QueuedMessages, MoveId) ->
909     State = finish_join(Me, Pred, Succ, DB, QueuedMessages),
910     SlideOp = slide_op:new_receiving_slide_join(MoveId, Pred, Succ, node:id(Me), join),
911     SlideOp1 = slide_op:set_phase(SlideOp, wait_for_node_update),
912     State1 = dht_node_state:set_slide(State, succ, SlideOp1),
913     State2 = dht_node_state:add_msg_fwd(
914         State1, slide_op:get_interval(SlideOp1),
915         node:pidX(slide_op:get_node(SlideOp1))),
916     RMSubscrTag = {move, slide_op:get_id(SlideOp1)},
917     comm:send_local(self(), {move, node_update, RMSubscrTag}),
918     gen_component:change_handler(State2, on).

```

The macro ?RT maps to the configured routing algorithm. It is defined in include/scalaris.hrl. For further details on the routing see Chapter 8.3 on page 32.



## Timeouts and other errors

The following table summarizes the timeout messages send during the join protocol on the joining node. It shows in which of the phases each of the messages is processed and describes (in short) what actions are taken. All of these messages are influenced by their respective config parameters, e.g. `join_timeout` parameter in the config files defines an overall timeout for the whole join operation. If it takes longer than `join_timeout` ms, a `{join, timeout}` will be send and processed as given in this table.

	<code>known_hosts_timeout</code>	<code>get_number_of_samples_timeout</code>	<code>lookup_timeout</code>	<code>join_request_timeout</code>	<code>timeout</code>
<b>phase2</b>	get known nodes from configured VMs	ignore	ignore	ignore	
<b>phase2b</b>	ignore	remove contact node, re-start join → phase 2 or 2b	ignore	ignore	
<b>phase3</b>	ignore	ignore	remove contact node, lookup remaining IDs → phase 2 or 3	ignore	
<b>phase3b</b>	ignore	ignore	ignore	ignore	re-start join → phase 2 or 2b
<b>phase4</b>	ignore	ignore	ignore	timeouts < 3? <sup>2</sup> → contact candidate otherwise: remove candidate no candidates left? → phase 2 or 2b otherwise: → contact next one → phase 3b or 4	

On the existing node, there is only one timeout message which is part of the join protocol: the `join_response_timeout`. It will be send when a slide operation is set up and if the timeout hits before the next message exchange, it will increase the slide operation's number of timeouts. The slide will be aborted if at least `join_response_timeouts` timeouts have been received. This parameter is set in the config file.

### Misc. (all phases)

Note that join-related messages arriving in other phases than those handling them will be ignored. Any other messages during a `dht_node`'s join will be queued and re-send when the join is complete.

<sup>2</sup>set by the `join_request_timeouts` config parameter



# 11. Directory Structure of the Source Code

The directory tree of Scalaris is structured as follows:

<code>bin</code>	contains shell scripts needed to work with Scalaris (e.g. start the boot services, start a node, ...)
<code>contrib</code>	necessary third party packages (yaws and log4erl)
<code>doc</code>	generated Erlang documentation
<code>docroot</code>	root directory of the node's webserver
<code>ebin</code>	the compiled Erlang code (beam files)
<code>java-api</code>	a Java API to Scalaris
<code>log</code>	log files
<code>src</code>	contains the Scalaris source code
<code>test</code>	unit tests for Scalaris
<code>user-dev-guide</code>	contains the sources for this document

## 12. Java API

For the Java API documentation, we refer the reader to the documentation generated by javadoc or doxygen. The following commands create the documentation:

```
%> cd java-api  
%> ant doc  
%> doxygen
```

The documentation can then be found in `java-api/doc/index.html` (javadoc) and `java-api/doc-doxygen/html/index.html` (doxygen).

The API is divided into four classes:

- `de.zib.scalarisc.Transaction` for (multiple) operations inside a transaction
- `de.zib.scalarisc.TransactionSingleOp` for single transactional operations
- `de.zib.scalarisc.ReplicatedDHT` for non-transactional (inconsistent) access to the replicated DHT items, e.g. deleting items
- `de.zib.scalarisc.PubSub` for topic-based publish/subscribe operations

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