

TRANSACTIONS



Scalaris:

Users and Developers Guide

Version 0.2.0 draft

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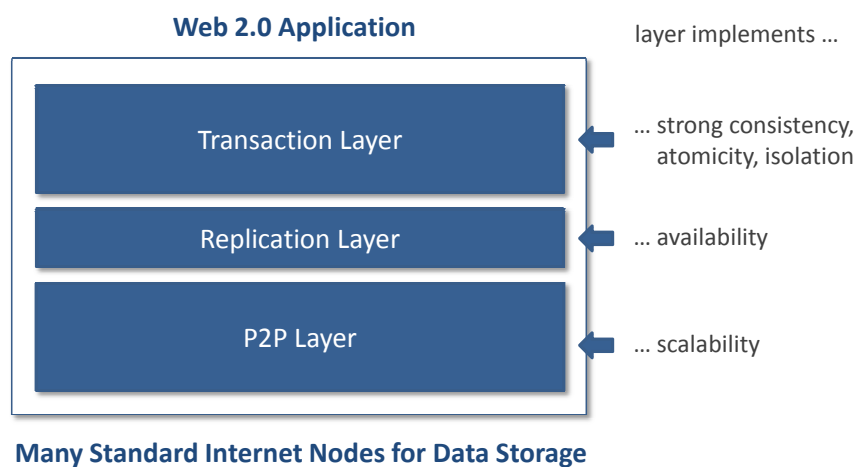
Part I.

Users Guide

1. Introduction

Scalaris is a scalable, transactional, distributed key-value store based on the peer-to-peer principle. It can be used to build scalable Web 2.0 services. The concept of Scalaris is quite simple: Its architecture consists of three layers.

It provides self-management and scalability by replicating services and data among peers. Without system interruption it scales from a few PCs to thousands of servers. Servers can be added or removed on the fly without any service downtime.



Scalaris takes care of:

- Fail-over
- Data distribution
- Replication
- Strong consistency
- Transactions

The Scalaris project was initiated by Zuse Institute Berlin and onScale solutions and was partly funded by the EU projects Selfman and XtreamOS. Additional information (papers, videos) can be found at <http://www.zib.de/CSR/Projects/scalaris> and <http://www.onscale.de/scalarix.html>.

1.1. Brewer's CAP Theorem

In distributed computing there exists the so called CAP theorem. It basically says that there are three desirable properties for distributed systems but one can only have any two of them.

Strict Consistency. Any read operation has to return the result of the latest write operation on the same data item.

Availability. Items can be read and modified at any time.

Partition Tolerance. The network on which the service is running may split into several partitions which cannot communicate with each other. Later on the networks may re-join again.

For example, a service is hosted on one machine in Seattle and one machine in Berlin. This service is partition tolerant if it can tolerate that all Internet connections over the Atlantic (and Pacific) are interrupted for a few hours and then get repaired.

The goal of Scalaris is to provide strict consistency and partition tolerance. We are willing to sacrifice availability to make sure that the stored data is always consistent. I.e. when you are running Scalaris with a replication degree of 4 and the network splits into two partitions, one partition with three replicas and one partition with one replica, you will be able to continue to use the service only in the larger partition. All requests in the smaller partition will time out until the two networks merge again. Note, most other key-value stores tend to sacrifice consistency.

2. Download and Installation

2.1. Requirements

For building and running Sclaris, some third-party modules are required which are not included in the Sclaris sources:

- Erlang R13B01 or newer
- GNU-like Make

To build the Java API (and the command-line client) the following modules are required additionally:

- Java Development Kit 6
- Apache Ant

Before building the Java API, make sure that `JAVA_HOME` and `ANT_HOME` are set. `JAVA_HOME` has to point to a JDK installation, and `ANT_HOME` has to point to an Ant installation.

2.2. Download

The sources can be obtained from <http://code.google.com/p/scalaris>. RPMs and DEBs are available from <http://download.opensuse.org/repositories/home:/tschuett/>.

2.2.1. Development Branch

You find the latest development version in the svn repository:

```
# Non-members may check out a read-only working copy anonymously over HTTP.  
svn checkout http://scalaris.googlecode.com/svn/trunk/ scalaris-read-only
```

2.2.2. Releases

Releases can be found under the 'Download' tab on the web-page.

2.3. Configuration

Sclaris reads two configuration files from the working directory: `bin/scalaris.cfg` (mandatory) and `bin/scalaris.local.cfg` (optional). The former defines default settings and is included in the release. The latter can be created by the user to alter settings. A sample file is provided as `bin/scalaris.local.cfg.example`. To run Sclaris distributed over several nodes, each node requires a `bin/scalaris.local.cfg`:

File `scalaris.local.cfg`:

```
%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
% Settings for distributed Erlang
% (see scalaris.hrl to switch)

% {boot_host, {boot, 'boot@foo.bar.com'}}}.
% {known_hosts, [{service_per_vm, 'boot@foo.bar.com'}]}.

%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
% Settings for TCP mode.
% (see scalaris.hrl to switch)

% Insert the appropriate IP-addresses for your setup
% as comma separated integers:
% IP Address, Port, and label of the boot server
{boot_host, {{127,0,0,1},14195,boot}}.

% IP Address, Port, and label of a node which is already in the system
{known_hosts, [{{{127,0,0,1},14195, service_per_vm}]}].
```

Scalaris currently distinguishes two different kinds of nodes: (a) the boot-server and (b) regular nodes. For the moment, we limit the number of boot-servers to exactly one. The remaining nodes are regular nodes. The boot-server is contacted to join the system. On all servers, the `boot_host` option defines the server where the boot server is running. In the example, it is an IP address plus a TCP port.

2.4. Build

2.4.1. Linux

Scalaris uses `autoconf` for configuring the build environment and `GNU Make` for building the code.

```
%> ./configure
%> make
%> make docs
```

For more details read `README` in the main Scalaris checkout directory.

2.4.2. Windows

We are currently not supporting Scalaris on Windows. However, we have two small bat files for building and running scalaris nodes. It seems to work but we make no guarantees.

- Install Erlang
<http://www.erlang.org/download.html>
- Install OpenSSL (for crypto module)
<http://www.slproweb.com/products/Win32OpenSSL.html>
- Checkout scalaris code from SVN
- adapt the path to your Erlang installation in `build.bat`
- start a `cmd.exe`
- go to the scalaris directory
- run `build.bat` in the cmd window

- check that there were no errors during the compilation; warnings are fine
- go to the bin sub-directory
- adapt the path to your Erlang installation in `boot.bat`, `cs_local.bat`, `cs_local2.bat` and `cs_local3.bat`
- run `boot.bat` or one of the other start scripts in the cmd window

`build.bat` will generate a `Emakefile` if there is none yet. If you have Erlang < R13B04, you will need to adapt the `Emakefile`. There will be empty lines in the first three blocks ending with “`}]`.”: add the following to these lines and try to compile again. It should work now.

```
, {d, type_forward_declarations_are_not_allowed}
, {d, forward_or_recursive_types_are_not_allowed}
```

For the most recent description please see the FAQ at <http://code.google.com/p/scalaris/wiki/FAQ>.

2.4.3. Java-API

The following commands will build the Java API for Scalaris:

```
%> make java
```

This will build `scalaris.jar`, which is the library for accessing the overlay network. Optionally, the documentation can be build:

```
%> cd java-api
%> ant doc
```

2.5. Running Scalaris

As mentioned above, in Scalaris there are two kinds of nodes:

- boot servers
- regular nodes

In every Scalaris, at least one boot server is required. It will maintain a list of nodes taking part in the system and allows other nodes to join the ring. For redundancy, it is also possible to have several boot servers. In the future, we want to eliminate this distinction, so any node is also a boot-server.

2.5.1. Running on a local machine

Open at least two shells. In the first, inside the Scalaris directory, start the boot script (`boot.bat` on Windows):

```
%> ./bin/boot.sh
```

This will start the boot server. On success <http://localhost:8000> should point to the management interface page of the boot server. The main page will show you the number of nodes currently in

the system. After a couple of seconds a first Scalaris should have started in the boot server and the number should increase to one. The main page will also allow you to store and retrieve key-value pairs but should not be used by applications to access Scalaris. See Chapter 3 on page 12 for application APIs.

In a second shell, you can now start a second Scalaris node. This will be a ‘regular server’:

```
%> ./bin/cs_local.sh
```

The second node will read the configuration file and use this information to contact the boot server and join the ring. The number of nodes on the web page should have increased to two by now.

Optionally, a third and fourth node can be started on the same machine. In a third shell:

```
%> ./bin/cs_local2.sh
```

In a fourth shell:

```
%> ./bin/cs_local3.sh
```

This will add 3 nodes to the network. The web pages at <http://localhost:8000> should show the additional nodes.

On linux you can also use the `scalarisctl` script to start boot and ‘regular’ nodes.

2.5.2. Running distributed

Scalaris can be installed on other machines in the same way as described in Section 2.6. In the default configuration, nodes will look for the boot server on `localhost` on port 14195. You should create a `scalaris.local.cfg` pointing to the node running the boot server.

```
% Insert the appropriate IP-addresses for your setup
% as comma separated integers:
% IP Address, Port, and label of the boot server
{boot_host, {{127,0,0,1},14195,boot}}.
```

If you are using the default configuration on the boot server it will listen on port 14195 and you only have to change the IP address in the configuration file. Otherwise the other nodes will not find the boot server. On the remote nodes, you only need to call `./cs_local.sh` and they will automatically contact the configured boot server.

2.6. Installation

For simple tests, you do not need to install Scalaris. You can run it directly from the source directory. Note: `make install` will install `scalaris` into `/usr/local` and place `scalarisctl` into `/usr/local/bin`. But is more convenient to build an RPM and install it.

```
svn checkout http://scalaris.googlecode.com/svn/trunk/ scalaris-0.0.1
tar -cvjf scalaris-0.0.1.tar.bz2 scalaris-0.0.1 --exclude-vcs
cp scalaris-0.0.1.tar.bz2 /usr/src/packages/SOURCES/
rpmbuild -ba scalaris-0.0.1/contrib/scalaris.spec
```

Your source and binary RPM will be generated in `/usr/src/packages/SRPMS` and `RPMS`. We build RPMs and Debs using checkouts from svn and provide them using the openSUSE BuildService at <http://download.opensuse.org/repositories/home:/tschuett/>. Packages are available for

- Fedora 9, 10, 11, 12, 13,
- Mandriva 2008, 2009, 2009.1, 2010,
- openSUSE 11.0, 11.1, 11.2, 11.3, Factory,
- SLE 10, 11,
- CentOS 5.4,
- RHEL 5,
- Debian 5.0 and
- Ubuntu 9.04, 9.10.

Inside those repositories you will also find an erlang RPM - you don't need this if you already have a recent enough erlang version!

2.7. Logging

Scalaris uses the `log4erl` library (see `contrib/log4erl`) for logging status information and error messages. The log level can be configured in `bin/scalaris.cfg`. The default value is `error`; only errors and severe problems are logged.

```
%% @doc Loglevel: debug < info < warn < error < fatal < none
{log_level, error}.
```

In some cases, it might be necessary to get more complete logging information, e.g. for debugging. In 10.2 on page 39, we are explaining the startup process of Scalaris nodes in more detail, here the `info` level provides more detailed information.

```
%% @doc Loglevel: debug < info < warn < error < fatal < none
{log_level, info}.
```

3. Using the system

3.1. JSON API

Scalaris supports a JSON API for transactions. To minimize the necessary round trips between a client and Scalaris, it uses request lists, which contain all requests that can be done in parallel. The request list is then send to a Scalaris node with a POST message. The result is an opaque TransLog and a list containing the results of the requests. To add further requests to the transaction, the TransLog and another list of requests may be send to Scalaris. This process may be repeated as often as necessary. To finish the transaction, the request list can contain a 'commit' request as the last element, which triggers the validation phase of the transaction processing.

The JSON-API can be accessed via the Scalaris-Web-Server running on port 8000 by default and the page `jsonrpc.yaws` (For example at: <http://localhost:8000/jsonrpc.yaws>). The following example illustrates the message flow:

Client

Make a transaction, that sets two keys:

Scalaris node

→

```
{
  "method": "req_list",
  "version": "1.1",
  "params":
  [
    [
      { "write": { "keyA": "valueA" } },
      { "write": { "keyB": "valueB" } },
      { "commit": "commit" }
    ]
  ],
  "id": 0
}
```

← Scalaris sends results back

```
{ "result":
  { "results":
    [
      { "op": "commit",
        "value": "ok",
        "key": "ok" },
      { "op": "write",
        "value": "valueB",
        "key": "keyB" },
      { "op": "write",
        "value": "valueA",
        "key": "keyA" }
    ],
    "translog":
      [...]
  },
  "id" : 0
}
```

In a second transaction: Read the two keys →

```
{
  "method": "req_list",
  "version": "1.1",
  "params":
    [
      [
        { "read": "keyA" },
        { "read": "keyB" }
      ]
    ]
  "id": 0
}
```

← Scalaris sends results back

```
{ "result":
  { "results":
    [
      { "op": "read",
        "value": "valueB",
        "key": "keyB" },
      { "op": "read",
        "value": "valueA",
        "key": "keyA" }
    ],
    "translog":
      [...] // this list is the translog
              // for further operations!
              // We name it TLOG here.
  },
  "id" : 0
}
```

Calculate something with the read values →
and make further requests, here a write
and the commit for the whole transaction.
Also include the latest translog we got from
Scalaris (named TLOG here).

```
{
  "method": "req_list",
  "version": "1.1",
  "params":
  [
    TLOG, // translog from prev. result
    [
      { "write": { "keyA": "valueA2" } },
      { "commit": "commit" }
    ]
  ],
  "id" : 0
}
```

← Scalaris sends results back

```
{ "result":
  { "results":
    [ { "op": "commit",
        "value": "ok",
        "key": "ok" },
      { "op": "write",
        "value": "valueA2",
        "key": "keyA" }
    ],
    "translog":
    [...]
  },
  "id" : 0
}
```

A sample usage of the JSON API using Ruby can be found in contrib/jsonrpc.rb.

A single request list must not contain a key more than once!

The allowed requests are:

```
{ "read": "any_key" }

{ "write": { "any_key": "any_value" } }

{ "commit": "commit" }
```

The possible results are:

```
{ "op": "read", "key": "any_key", "value": "any_value" }
{ "op": "read", "key": "any_value", "fail": "reason" } // 'not_found' or 'timeout'

{ "op": "write", "key": "any_key", "value": "any_value" }
{ "op": "read", "key": "any_key", "fail": "reason" }

{ "op": "commit", "value": "ok", "key": "ok" }
{ "op": "commit", "value": "fail", "fail": "reason" }
```

3.1.1. Deleting a key

Outside transactions keys can also be deleted, but it has to be done with care, as explained in the following thread on the mailing list: http://groups.google.com/group/scalaris/browse_thread/thread/ff1d9237e218799.

```
{
  "method": "delete",
  "version": "1.1",
  "params":
    [
      { "key": "any_key" }
    ],
  "id" : 0
}
```

Two sample results

```
{ "result":
  { "ok":2, // how many replicas were deleted successssfully
    "results": [ "ok", "ok", "locks_set", "undef" ]
  }
}
```

```
{ "result":
  { "failure": "reason" }
}
```

3.2. Java command line interface

The jar file contains a small command line interface client. For convenience, we provide a wrapper script called `scalaris` which sets up the Java environment:

```
%> ./java-api/scalaris -help
Script Options:
  --help, --h          print this message and scalaris help
  --noconfig           suppress sourcing of /etc/scalaris/scalaris-java.conf
                      and $HOME/.scalaris/scalaris-java.conf config files
  --execdebug          print scalaris exec line generated by this
                      launch script

usage: scalaris [Options]
  -b,--minibench        run mini benchmark
  -d,--delete <key> <[timeout]> delete an item (default timeout: 2000ms)
                          WARNING: This function can lead to inconsistent data
                          (e.g. deleted items can re-appear). Also when
                          re-creating an item the version before the delete can
                          re-appear.
  -g,--getsubscribers <topic> get subscribers of a topic
  -h,--help            print this message
  -lh,--localhost      gets the local host's name as known to
                          Java (for debugging purposes)
  -p,--publish <topic> <message> publish a new message for the given
                          topic
  -r,--read <key>      read an item
  -s,--subscribe <topic> <url> subscribe to a topic
  -u,--unsubscribe <topic> <url> unsubscribe from a topic
  -v,--verbose         print verbose information, e.g. the
                          properties read
  -w,--write <key> <value> write an item
```

read, write and delete can be used to read, write and delete from/to the overlay, respectively. getsubscribers, publish, and subscribe are the PubSub functions. The others provide debugging and testing functionality.

```
%> ./java-api/scalaris -write foo bar
write(foo, bar)
%> ./java-api/scalaris -read foo
read(foo) == bar
```

Per default, the `scalaris` script tries to connect to a boot server at `localhost`. You can change the node it connects to (and further connection properties) by adapting the values defined in `java-api/scalaris.properties`.

3.3. Java API

The `scalaris.jar` provides the command line client as well as a library for Java programs to access Scalaris. The library provides two classes:

- `Scalaris` provides a high-level API similar to the command line client.
- `Transaction` provides a low-level API to the transaction mechanism.

For details we refer the reader to the Javadoc:

```
%> cd java-api
%> ant doc
%> firefox doc/index.html
```


4. Testing the system

4.1. Running the unit tests

There are some unit tests in the `test` directory. You can call them by running `make test` in the main directory. The results are stored in a local `index.html` file.

The tests are implemented with the `common-test` package from the Erlang system. For running the tests we rely on `run_test`, which is part of the `common-test` package, but (on `erlang < R14`) is not installed by default. `configure` will check whether `run_test` is available. If it is not installed, it will show a warning and a short description of how to install the missing file.

Note: for the unit tests, we are setting up and shutting down several overlay networks. During the shut down phase, the runtime environment will print extensive error messages. These error messages do not indicate that tests failed! Running the complete test suite takes about 3 minutes, depending on your machine. Only if the complete suite finishes, it will present statistics on failed and successful tests.

5. Troubleshooting

5.1. Network

Scalaris uses a couple of TCP ports for communication. It does not use UDP at the moment.

- 8000 HTTP Server on the boot node
- 8001 HTTP Server on the other nodes
- 14195 Port for inter-node communication (boot server)
- 14196 Port for inter-node communication (other nodes)

Please make sure that at least 14195 and 14196 are not blocked by firewalls.

Part II.

Developers Guide

6. General Hints

6.1. Coding Guidelines

- Keep the code short
- Use `gen_component` to implement additional processes
- Don't use `receive` by yourself (Exception: to implement single threaded user API calls (`cs_api`, `yaws_calls`, etc))
- Don't use `erlang:now()` , `erlang:send_after()` , `receive after` etc. in performance critical code, consider using `msg_delay` instead.
- Don't use `tc:timer()` as it catches exceptions

6.2. Testing Your Modifications and Extensions

- Run the testsuites using `make test`
- Run the java api test using `make java-test` (Scalaris output will be printed if a test fails; if you want to see it during the tests, start a `bin/boot.sh` and run the tests by `cd java; ant test`)
- Run the Ruby client by starting Scalaris and running `cd contrib; ./jsonrpc.rb`

6.3. Help with Digging into the System

- use `ets:i(t)` to get details on the local state of some processes
- consider changing `pdb.erl` to use `ets` instead of `erlang:put/get`
- Have a look at `strace -f -p PID` of beam process
- Get message statistics via the Web-interface
- enable/disable tracing for certain modules
- Use `etop` and look at the total memory size and atoms generated
- send processes `sleep` or `kill` messages to test certain behaviour (see `gen_component.erl`)
- `USE boot_server:number_of_nodes(). flush().`
- `USE admin_checkring(). flush().`

7. System Infrastructure

7.1. The Process Dictionary

- What is it? How to distinguish from Erlangs internal process dictionary?
- Joining a process group (InstanceId is a group name)
- Why do we do this... (managing several independent nodes inside a single Erlang VM)

7.2. The Communication Layer `comm`

- in general
- format of messages (tuples)
- use messages with cookies (server and client side)
- What is a message tag?

7.3. The `gen_component`

Description is based on SVN revision r832.

The generic component model implemented by `gen_component` allows to add some common functionality to all the components that build up the Scalaris system. It supports:

event-handlers: message handling with a similar syntax as used in [2].

FIFO order of messages: components cannot be inadvertently locked as we do not use selective receive statements in the code.

sleep and halt: for testing components can sleep or be halted.

debugging, breakpoints, stepwise execution: to debug components execution can be steered via breakpoints, step-wise execution and continuation based on arriving events and user defined component state conditions.

basic profiling ,

state dependent message handlers: depending on its state, different message handlers can be used and switched during runtime. Thereby a kind of state-machine based message handling is supported.

prepared for process_dictionary: allows to send events to named processes inside the same group as the actual component itself (`send_to_group_member`) when just holding a reference to any group member, and

unit-testing of event-handlers: as message handling is separated from the main loop of the component, the handling of individual messages and thereby performed state manipulation can easily be tested in unit-tests by directly calling message handlers.

In Scalaris all Erlang processes should be implemented as `gen_component`. The only exception are functions interfacing to the client, where a transition from asynchronous to synchronous request handling is necessary and that are executed in the context of a client's process or a process that behaves as a proxy for a client (`cs_api`).

7.3.1. A basic `gen_component` including a message handler

To implement a `gen_component`, the component has to provide the `gen_component` behaviour:

File `gen_component.erl`:

```
45 behaviour_info(callbacks) ->
46 [
47     {init, 1},      % initialize component
48     {on, 2}         % handle a single message
49                     % on(Msg, State) -> NewState | unknown_event | kill
50 ];
```

This is illustrated by the following example:

File `idholder.erl`:

```
88 %% @doc Initialises the idholder with a random key and a counter of 0.
89 -spec init([]) -> state().
90 init(_Arg) ->
91     {get_initial_key(config:read(key_creator)), 0}.
92
93 -spec on(message(), state()) -> state() | unknown_event.
94 on({reinit, _State}) ->
95     {get_initial_key(config:read(key_creator)), 0};
96 on({get_id, PID}, {Key, Count} = State) ->
97     comm:send_local(PID, {idholder_get_id_response, Key, Count}),
98     State;
99 on({set_id, NewKey, Count}, _State) ->
100     {NewKey, Count};
101 on(_, _State) ->
102     unknown_event.
```

`your_gen_component:init/1` is called during start-up of a `gen_component` and should return the initial state to be used for this `gen_component`.

To react on messages / events, a message handler is used. The default message handler is called `your_gen_component:on/2`. This can be changed by calling `gen_component:change_handler/2` (see Section 7.3.6). When an event / message for the component arrives, this handler is called with the event itself and the current state of the component. In the handler, the state of the component may be adjusted depending upon the event. The handler itself may trigger new events / messages for itself or other components and has finally to return the updated state of the component or the atoms `unknown_event` or `kill`. It must neither call `receive` nor `timer:sleep/1` nor `erlang:exit/1`.

7.3.2. How to start a `gen_component`?

A `gen_component` can be started using one of:

```
gen_component:start(Module, Args, GenCOptions = [])
```

```
gen_component:start_link(Module, Args, GenCOptions = [])
```

Module: the the name of the module your component is implemented in

Args: List of parameters passed to `Module:init/1` for initialization

GenCOptions: optional parameter. List of options for `gen_component`

`{register, ProcessGroup, ProcessName}`: registers the new process with the given process group (also called instanceid) and name in the `process_dictionary`.
`{register_native, ProcessName}`: registers the process as a named Erlang process.
`wait_for_init`: wait for `Module:init/1` to return before returning to the caller.

These functions are compatible to the Erlang/OTP supervisors. They spawn a new process for the component which itself calls `Module:init/1` with the given `Args` to initialize the component. `Module:init/1` should return the initial state for your component. For each message sent to this component, the default message handler `Module:on(Message, State)` will be called, which should react on the message and return the updated state of your component.

`gen_component:start()` and `gen_component:start_link()` return the pid of the spawned process as `{ok, Pid}`.

7.3.3. When does a `gen_component` terminate?

A `gen_component` can be stopped using:

`gen_component:kill(Pid)` or by returning `kill` from the current message handler.

7.3.4. What happens when unexpected events / messages arrive?

Your message handler (default is `your_gen_component:on/2`) should return `unknown_event` in the final clause (`your_gen_component:on(_, _)`). `gen_component` then will nicely report on the unhandled message, the component's name, its state and currently active message handler, as shown in the following example:

```
# bin/boot.sh
[...]
(boot@localhost)10> process_dictionary ! {no_message}.
{no_message}
[error] unknown message: {no_message} in Module: process_dictionary and
handler on in State null
(boot@localhost)11>
```

The `process_dictionary` (see Section 7.1) is a `gen_component` which registers itself as named Erlang process with the `gen_component` option `register_native` and therefore can be addressed by its name in the Erlang shell. We send it a `{no_message}` and `gen_component` reports on the unhandled message. The `process_dictionary` itself continues to run and waits for further messages.

7.3.5. What if my message handler generates an exception or crashes the process?

`gen_component` catches exceptions generated by message handlers and reports them with a stack trace, the message, that generated the exception, and the current state of the component.

If a message handler terminates the process via `erlang:exit/1`, this is out of the responsibility scope of `gen_component`. As usual in Erlang, all linked processes will be informed. If for example `gen_component:start_link/2` or `/3` was used for starting the `gen_component`, the spawning process will be informed, which may be an Erlang supervisor process taking further actions.

7.3.6. Changing message handlers and implementing state dependent message responsiveness as a state-machine

Sometimes it is beneficial to handle messages depending on the state of a component. One possibility to express this is implementing different clauses depending on the state variable, another is introducing case clauses inside message handlers to distinguish between current states. Both approaches may become tedious, error prone, and may result in confusing source code.

Sometimes the use of several different message handlers for different states of the component leads to clearer arranged code, especially if the set of handled messages changes from state to state. For example, if we have a component with an initialization phase and a production phase afterwards, we can handle in the first message handler messages relevant during the initialization phase and simply queue all other requests for later processing using a common default clause.

When initialization is done, we handle the queued user requests and switch to the message handler for the production phase. The message handler for the initialization phase does not need to know about messages occurring during production phase and the message handler for the production phase does not need to care about messages used during initialization. Both handlers can be made independent and may be extended later on without any adjustments to the other.

One can also use this scheme to implement complex state-machines by changing the message handler from state to state.

To switch the message handler `gen_component:change_handler(State, new_handler)` is called as the last operation after a message in the active message handler was handled, so that the return value of `gen_component:change_handler/2` is propagated to `gen_component`. The new handler is given as an atom, which is the name of the 2-ary function in your component module to be called.

Starting with non-default message handler.

It is also possible to change the message handler right from the start in your `your_gen_component:init/1` to avoid the default message handler `your_gen_component:on/2`. Just create your initial state as usual and call `gen_component:change_handler(State, my_handler)` as the final call in your `your_gen_component:init/1`. We prepared `gen_component:change_handler/2` to return `State` itself, so this will work properly.

7.3.7. Halting and pausing a `gen_component`

Using `gen_component:kill(Pid)` and `gen_component:sleep(Pid, Time)` components can be terminated or paused.

7.3.8. Integration with `process_dictionary`: Redirecting events / messages to other `gen_components`

Each `gen_component` by itself is prepared to support `comm:send_to_group_member/3` which forwards messages inside a group of processes registered via the `process_dictionary` (see Section 7.1) by their name. So, if you hold a `Pid` of one member of a process group, you can send messages to other members of this group, if you know their registered Erlang name. You do not necessarily have to know their individual `Pid`.

In consequence, no `gen_component` can individually handle messages of the form `{send_to_group_member, _, _}` as such messages are consumed by `gen_component` itself.

7.3.9. Integration with `fd_pinger`: Replying to failure detectors

Each `gen_component` replies automatically to `{ping, Pid}` requests with a `{pong}` send to the given `Pid`. Such messages are generated, for example, by `fd_pinger` which is used by our `fd` failure detectors.

In consequence, no `gen_component` can individually handle messages of the form: `{ping, _}` as such messages are consumed by `gen_component` itself.

7.3.10. The debugging interface of `gen_component`: Breakpoints and step-wise execution

We equipped `gen_component` with a debugging interface, which especially is beneficial, when testing the interplay between several `gen_components`. It supports breakpoints which can pause the `gen_component` depending on the arriving messages or depending on user defined conditions. If a breakpoint is reached, the execution can be continued step-wise (message by message) or until the next breakpoint is reached.

We use it in our unit tests to steer protocol interleavings and to perform tests using random protocol interleavings between several processes (see `paxos_SUITE`). It allows also to reproduce given protocol interleavings for better testing.

Managing breakpoints.

Breakpoints are managed by the following functions:

`gen_component:bp_set(Pid, MsgTag, BPName)` : For the component running under `Pid` a breakpoint `BPName` is set. It is reached, when a message with a message tag `MsgTag` is next to be handled by the component (See `comm:get_msg_tag/1` and Section 7.2 for more information on message tags). The `BPName` is used as a reference for this breakpoint, for example to delete it later.

`gen_component:bp_set_cond(Pid, Cond, BPName)` : The same as `gen_component:bp_set/3` but a user defined condition implemented in `{Module, Function, Params = 2}` = `Cond` is checked by calling `Module:Function(Message, State)` to decide whether a breakpoint is reached or not. `Message` is the next message to be handled by the component and `State` is the current state of the component. `Module:Function/2` should return a boolean.

`gen_component:bp_del(Pid, BPName)` : The breakpoint `BPName` is deleted. If the component is in this breakpoint, it will not be released by this call. This has to be done separately by `gen_component:cont/1`. But the deleted breakpoint will no longer be considered for newly entering a breakpoint.

`gen_component:bp_barrier(Pid)` : Delay all further handling of breakpoint requests until a breakpoint is actually entered.

Note, that the following call sequence may not catch the breakpoint at all, as during the sleep the component not necessarily consumes a ping message and the set breakpoint may already be deleted before a ping arrives.

```

gen_component:bp_set(Pid, ping, bp_ping),
timer:sleep(10),
gen_component:bp_del(Pid, bp_ping),
gen_component:cont(Pid).

```

This is where `gen_component:bp_barrier/1` can be used:

```

gen_component:bp_set(Pid, ping, bp_ping),
gen_component:bp_barrier(Pid),
%% the following breakpoint requests will not be handled before a
%% breakpoint is reached.
%% the gen_component itself is still active and handles messages as usual
%% up to the next breakpoint
gen_component:bp_del(Pid, bp_ping),
% the breakpoint was entered once, so we delete.
% next we leave the breakpoint and continue the gen_component
gen_component:cont(Pid).

```

None of the calls in the sample listing above is blocking. It just schedules all the operations, including the `bp_barrier`, for the `gen_component` and immediately finishes. The actual events of entering and continuing the breakpoint in the `gen_component` can happen independently later on, when the next ping message arrives.

Managing execution.

The execution of a `gen_component` can be managed by the following functions:

`gen_component:bp_step(Pid)` : This is the only blocking breakpoint function. It waits until the `gen_component` is in a breakpoint and has handled a single message. It returns the module, the active message handler, and the handled message as a tuple `{Module, On, Message}`. This function does not actually finish the breakpoint, but just lets a single message pass through. For further messages, no breakpoint condition has to be valid, the original breakpoint is still active. To leave a breakpoint, use `gen_component:bp_cont/1`.

`gen_component:bp_cont(Pid)` : Leaves a breakpoint. `gen_component` runs as usual until the next breakpoint is reached.

If no further breakpoints should be entered after continuation, you should delete the registered breakpoint using `gen_component:bp_del/2` before continuing the execution with `gen_component:bp_cont/1`. To ensure, that the breakpoint is entered at least once, `gen_component:bp_barrier/1` should be used before deleting the breakpoint (see the example above). Otherwise it could happen, that the delete request arrives at your `gen_component` before it was actually triggered. The following continuation request would then unintentional apply to an unrelated breakpoint that may be entered later on.

`gen_component:runnable(Pid)` : Returns whether a `gen_component` has messages to handle and is runnable. If you know, that a `gen_component` is in a breakpoint, you can use this to check, whether a `gen_component:bp_step/1` or `gen_component:bp_cont/1` is applicable to the component.

Tracing handled messages – getting a message interleaving protocol.

We use the debugging interface of `gen_component` to test protocols with random interleaving. First we start all the components involved, set breakpoints on the initialization messages for a new

Paxos consensus and then start a single Paxos instance on all of them. The outcome of the Paxos consensus is a learner_decide message. So, in paxos_SUITE:step_until_decide/3 we look for runnable processes and select randomly one of them to perform a single step until the protocol finishes with a decision.

File paxos_SUITE.erl:

```

223 -spec(prop_rnd_interleave/3 :: (1..4, 4..16, {pos_integer(), pos_integer(), pos_integer()}
224     -> boolean()).
225 prop_rnd_interleave(NumProposers, NumAcceptors, Seed) ->
226     ct:pal("Called with: paxos_SUITE:prop_rnd_interleave(~p, ~p, ~p).~n",
227         [NumProposers, NumAcceptors, Seed]),
228     Majority = NumAcceptors div 2 + 1,
229     {Proposers, Acceptors, Learners} =
230         make(NumProposers, NumAcceptors, 1, rnd_interleave),
231     %% set bp on all processes
232     [ gen_component:bp_set(element(3, X), proposer_initialize, bp)
233       || X <- Proposers ],
234     [ gen_component:bp_set(element(3, X), acceptor_initialize, bp)
235       || X <- Acceptors ],
236     [ gen_component:bp_set(element(3, X), learner_initialize, bp)
237       || X <- Learners ],
238     %% start paxos instances
239     [ proposer:start_paxosid(X, paxidrndinterl, Acceptors,
240         proposal, Majority, NumProposers, Y)
241       || {X,Y} <- lists:zip(Proposers, lists:seq(1, NumProposers)) ],
242     [ acceptor:start_paxosid(X, paxidrndinterl, Learners)
243       || X <- Acceptors ],
244     [ learner:start_paxosid(X, paxidrndinterl, Majority,
245         comm:this(), cpaxidrndinterl)
246       || X <- Learners ],
247     %% randomly step through protocol
248     OldSeed = random:seed(Seed),
249     Steps = step_until_decide(Proposers ++ Acceptors ++ Learners, cpaxidrndinterl, 0),
250     ct:pal("Needed ~p steps~n", [Steps]),
251     case OldSeed of
252         undefined -> ok;
253         _ -> random:seed(OldSeed)
254     end,
255     true.
256
257 step_until_decide(Processes, PaxId, SumSteps) ->
258     %% io:format("Step ~p~n", [SumSteps]),
259     Runnable = [ X || X <- Processes, gen_component:runnable(element(3,X)) ],
260     case Runnable of
261         [] ->
262             ct:pal("No runnable processes of ~p~n", [length(Processes)]),
263             timer:sleep(5), step_until_decide(Processes, PaxId, SumSteps);
264         _ -> ok
265     end,
266     Num = random:uniform(length(Runnable)),
267     gen_component:bp_step(element(3,lists:nth(Num, Runnable))),
268     receive
269         {learner_decide, cpaxidrndinterl, _, _Res} = _Any ->
270             %% io:format("Received ~p~n", [_Any]),
271             SumSteps
272     after 0 -> step_until_decide(Processes, PaxId, SumSteps + 1)
273     end.

```

To get a message interleaving protocol, we either can output the results of each gen_component:-bp_step/1 call together with the Pid we selected for stepping, or alter the definition of the macro TRACE_BP_STEPS in gen_component, when we execute all gen_components locally in the same Erlang virtual machine.

File gen_component.erl:

```

30 %-define(TRACE_BP_STEPS(X,Y), io:format(X,Y)). %% output on console
31 %-define(TRACE_BP_STEPS(X,Y), ct:pal(X,Y)).    %% output even if called by unittest

```

32 `-define(TRACE_BP_STEPS(X,Y), ok).`

7.3.11. Future use and planned extensions for `gen_component`

`gen_component` could be further extended. For example it could support hot-code upgrade or could be used to implement algorithms that have to be run across several components of *Scalaris* like snapshot algorithms or similar extensions.

7.4. The Process' Database (pdb)

- How to use it and how to switch from `erlang:put/set` to `ets` and implied limitations.

7.5. Writing Unittests

7.5.1. Plain unittests

7.5.2. Randomized Testing using `tester.erl`

8. Basic Structured Overlay

8.1. Ring Maintenance

8.2. T-Man

8.3. Routing Tables

Description is based on SVN revision r915.

Each node of the ring can perform searches in the overlay.

A search is done by a lookup in the overlay, but there are several other demands for communication between peers. Scalaris provides a general interface to route a message to the (other) peer, which is currently responsible for a given key.

File `lookup.erl`:

```
30 -spec unreliable_lookup(Key::?RT:key(), Msg::comm:message()) -> ok.
31 unreliable_lookup(Key, Msg) ->
32     comm:send_local(process_dictionary:get_group_member(dht_node),
33                     {lookup_aux, Key, 0, Msg}).
34
35 -spec unreliable_get_key(Key::?RT:key()) -> ok.
36 unreliable_get_key(Key) ->
37     unreliable_lookup(Key, {get_key, comm:this(), Key}).
38
39 -spec unreliable_get_key(CollectorPid::comm:mypid(),
40                         ReqId::{rdht_req_id, {pos_integer(), comm:mypid()}},
41                         Key::?RT:key()) -> ok.
42 unreliable_get_key(CollectorPid, ReqId, Key) ->
43     unreliable_lookup(Key, {get_key, CollectorPid, ReqId, Key}).
```

The message `Msg` could be a `get_key` which retrieves content from the responsible node or a `get_node` message, which returns a pointer to the node.

All currently supported messages are listed in the file `dht_node.erl`.

The message routing is implemented in `dht_node_lookup.erl`

File `dht_node_lookup.erl`:

```
27 %% @doc Find the node responsible for Key and send him the message Msg.
28 -spec lookup_aux(State::dht_node_state:state(), Key::intervals:key(),
29                 Hops::non_neg_integer(), Msg::comm:message()) -> ok.
30 lookup_aux(State, Key, Hops, Msg) ->
31     case intervals:in(Key, dht_node_state:get(State, succ_range)) of
32     true -> % found node -> terminate
33         comm:send(dht_node_state:get(State, succ_pid),
34                 {lookup_fin, Hops + 1, Msg});
35     _ ->
36         P = ?RT:next_hop(State, Key),
37         comm:send(P, {lookup_aux, Key, Hops + 1, Msg})
38     end.
```

Each node is responsible for a certain key interval. The function `intervals:in/2` is used to decide, whether the key is between the current node and its successor. If that is the case, the final step is delivers a `lookup_fin` message to the local node. Otherwise, the message is forwarded to the next nearest known peer (listed in the routing table) determined by `?RT:next_hop/2`.

`rt_beh.erl` is a generic interface for routing tables. It can be compared to interfaces in Java. In Erlang interfaces can be defined using a so called 'behaviour'. The files `rt_simple` and `rt_chord` implement the behaviour 'rt_beh'.

The macro `?RT` is used to select the current implementation of routing tables. It is defined in `include/scalaris.hrl`.

File `scalaris.hrl`:

```
25 %%The RT macro determines which kind of routingtable is used. Uncomment the
26 %%one that is desired.
27
28 %%Standard Chord routingtable
29 -define(RT, rt_chord).
30
31 %%Simple routingtable
32 %-define(RT, rt_simple).
```

The functions, that have to be implemented for a routing mechanism are defined in the following file:

File `rt_beh.erl`:

```
35 behaviour_info(callbacks) ->
36 [
37     % create a default routing table
38     {empty, 1}, {empty_ext, 1},
39     % mapping: key space -> identifier space
40     {hash_key, 1}, {get_random_node_id, 0},
41     % routing
42     {next_hop, 2},
43     % trigger for new stabilization round
44     {init_stabilize, 3},
45     % dead nodes filtering
46     {filter_dead_node, 2},
47     % statistics
48     {to_pid_list, 1}, {get_size, 1},
49     % for symmetric replication
50     {get_keys_for_replicas, 1},
51     % address space size (throws 'throw:not_supported' if unsupported by the RT)
52     {n, 0},
53     % for debugging and web interface
54     {dump, 1},
55     % for bulkowner
56     {to_list, 1},
57     % convert from internal representation to version for dht_node
58     {export_rt_to_dht_node, 4},
59     % update pred/succ in routing stored in dht_node
60     {update_pred_succ_in_dht_node, 3},
61     % handle messages specific to a certain routing-table implementation
62     {handle_custom_message, 2},
63     % common methods, e.g. from rt_generic.hrl
64     {check, 5}, {check, 6}, {check_fd, 2},
65     {check_config, 0}
66 ];
```

`empty/1` gets a successor and generates an empty routing table for use inside the routing table implementation. The data structure of the routing table is undefined. It can be a list, a tree, a matrix ...

`empty_ext/1` similarly creates an empty external routing table for use by the `dht_node`. This process might not need all the information a routing table implementation requires and can thus work with less data.

`hash_key/1` gets a key and maps it into the overlay's identifier space.

`get_random_node_id/0` returns a random node id from the overlay's identifier space. This is used for example when a new node joins the system.

`next_hop/2` gets a `dht_node`'s state (including the external routing table representation) and a key and returns the node, that should be contacted next when searching for the key, i.e. the known node nearest to the id.

`init_stabilize/3` is called periodically to rebuild the routing table. The parameters are the identifier of the node, its successor and the old (internal) routing table state.

`filter_dead_node/2` is called by the failure detector and tells the routing table about dead nodes. This function gets the (internal) routing table and a node to remove from it. A new routing table state is returned.

`to_pid_list/1` get the PIDs of all (internal) routing table entries.

`get_size/1` get the (internal or external) routing table's size.

`get_keys_for_replicas/1` Returns for a given (un-hashed) `Key` the (hashed) keys of its replicas. This used for implementing symmetric replication.

`n/0` gets the number of available keys. An implementation may throw `throw:not_supported` if the operation is unsupported by the routing table.

`dump/1` dump the (internal) routing table state for debugging, e.g. by using the web interface. Returns a list of `{Index, Node_as_String}` tuples which may just as well be empty.

`to_list/1` convert the (external) representation of the routing table inside a given `dht_node_state` to a sorted list of known nodes from the routing table, i.e. `first=succ`, `second=next` known node on the ring, ... This is used by bulk-operations to create a broadcast tree.

`export_rt_to_dht_node/4` convert the internal routing table state to an external state. Gets the internal state, the node's ID, the predecessor and the successor for doing so.

`update_pred_succ_in_dht_node/3` update predecessor and successor node stored in an external routing table state.

`handle_custom_message/2` handle messages specific to the routing table implementation. `rt_loop` will forward unknown messages to this function.

`check/5`, `check/6`, `check_fd/2` are implemented in `rt_generic.hrl`, check for routing table changes and send an updated (external) routing table to the `dht_node` process.

`check_config/0` check that all required configuration parameters exist and satisfy certain restrictions.

8.3.1. Simple routing table

One implementation of a routing table is the `rt_simple`, which routes via the successor. Note that this is inefficient as it needs a linear number of hops to reach its goal. A more robust implementation, would use a successor list. This implementation is also not very efficient in the presence of churn.

Data types

First, the data structure of the routing table is defined:

File `rt_simple.erl`:

```
35 % @type key(). Identifier.
```

```

36 -opaque(key()::non_neg_integer()).
37 % @type rt(). Routing Table.
38 -opaque(rt()::Succ::node:node_type()).
39 -type(external_rt()::rt()).
40 -opaque(custom_message() :: any()).

```

The routing table only consists of a node (the successor). Keys in the overlay are identified by integers ≥ 0 .

A simple rm_beh behaviour

File rt_simple.erl:

```

48 %% @doc Creates an "empty" routing table containing the successor.
49 -spec empty(node:node_type()) -> rt().
50 empty(Succ) -> Succ.

```

File rt_simple.erl:

```

147 -spec empty_ext(node:node_type()) -> external_rt().
148 empty_ext(Succ) -> empty(Succ).

```

The empty routing table (internal or external) consists of the successor.

File rt_simple.erl:

```

54 %% @doc Hashes the key to the identifier space.
55 -spec hash_key(iodata() | integer()) -> key().
56 hash_key(Key) -> hash_key_(Key).
57
58 %% @doc Hashes the key to the identifier space (internal).
59 %%     Note: Needed for dialyzer to cope with the opaque key() type and the
60 %%     use of key() in get_keys_for_replicas/1.
61 -spec hash_key_(iodata() | integer()) -> non_neg_integer().
62 hash_key_(Key) when is_integer(Key) ->
63     <<N:128>> = erlang:md5(erlang:term_to_binary(Key)),
64     N;
65 hash_key_(Key) ->
66     <<N:128>> = erlang:md5(Key),
67     N.

```

Keys are hashed using MD5 and have a length of 128 bits.

File rt_simple.erl:

```

71 %% @doc Generates a random nodeid, i.e. a random 128-bit string.
72 -spec get_random_node_id() -> key().
73 get_random_node_id() -> hash_key(randoms:getRandomId()).

```

Random node id generation uses the helpers provided by the randoms module.

File rt_simple.erl:

```

152 %% @doc Returns the next hop to contact for a lookup.
153 -spec next_hop(dht_node_state:state(), key()) -> comm:mypid().
154 next_hop(State, _Key) -> node:pidX(dht_node_state:get(State, rt)).

```

Next hop is always the successor.

File `rt_simple.erl`:

```
81 %% @doc Triggered by a new stabilization round, renews the routing table.
82 -spec init_stabilize(key(), node:node_type(), rt()) -> rt().
83 init_stabilize(_Id, Succ, _RT) -> empty(Succ).
```

init_stabilize/3 resets its routing table to the current successor.

File `rt_simple.erl`:

```
87  %% @doc Removes dead nodes from the routing table (rely on periodic
88  %%      stabilization here).
89  -spec filter_dead_node(rt(), comm:mypid()) -> rt().
90  filter_dead_node(RT, _DeadPid) -> RT.
```

`filter_dead_node/2` does nothing, as only the successor is listed in the routing table and that is reset periodically in `init_stabilize/3`.

File rt_simple.erl:

```
94  %% @doc Returns the pids of the routing table entries.
95  -spec to_pid_list(rt() | external_rt()) -> [comm:mypid()].
96  to_pid_list(Succ) -> [node:pidX(Succ)].
```

to_pid_list/1 returns the pid of the successor.

File `rt_simple.erl`:

```
100 % @doc Returns the size of the routing table.
101 -spec get_size(rt() | external_rt()) -> 1.
102 get_size(_RT) -> 1.
```

The size of the routing table is always 1.

File rt_simple.erl:

```

112 % @doc Returns the replicas of the given key.
113 -spec get_keys_for_replicas(iodata() | integer()) -> [key()].
114 get_keys_for_replicas(Key) ->
115     HashedKey = hash_key_(Key),
116     [HashedKey,
117      HashedKey bxor 16#40000000000000000000000000000000,
118      HashedKey bxor 16#80000000000000000000000000000000,
119      HashedKey bxor 16#C0000000000000000000000000000000
120     ].

```

This `get_keys_for_replicas/1` implements symmetric replication.

File rt_simple.erl:

```
106 %% @doc Returns the size of the address space.  
107 -spec n() -> non_neg_integer().  
108 n() -> 16#10000000000000000000000000000000.
```

There are 2^{128} available keys.

File `rt_simple.erl`:

```
124 % @doc Dumps the RT state for output in the web interface.
125 -spec dump(RT::rt()) -> KeyValueCollection::[{Index::non_neg_integer(), Node::string()}].
126 dump(Succ) -> [{0, lists:flatten(io_lib:format("~p", [Succ]))}].
```

dump/1 lists the successor.

File `rt_simple.erl`:

```
173 %% @doc Converts the (external) representation of the routing table to a list
174 %%      in the order of the fingers, i.e. first=succ, second=shortest finger,
175 %%      third=next longer finger,...
176 -spec to_list(dht_node_state:state()) -> nodelist:nodelist().
177 to_list(State) -> [dht_node_state:get(State, rt)].
```

`to_list/1` lists the successor from the external routing table state.

File `rt_simple.erl`:

```
158 %% @doc Converts the internal RT to the external RT used by the dht_node. Both
159 %%      are the same here.
160 -spec export_rt_to_dht_node(rt(), ID::key(), Pred::node:node_type(),
161                             Succ::node:node_type()) -> external_rt().
162 export_rt_to_dht_node(RT, _Id, _Pred, _Succ) -> RT.
```

`export_rt_to_dht_node/1` states that the external routing table is the same as the internal table.

File `rt_simple.erl`:

```
166 %% @doc Updates the successor in the (external) routing table state.
167 -spec update_pred_succ_in_dht_node(Pred::node:node_type(), Succ::node:node_type(),
168                                   RT::external_rt()) -> external_rt().
169 update_pred_succ_in_dht_node(_Pred, Succ, Succ) -> Succ.
```

Updates the external routing table state with the given successor.

File `rt_simple.erl`:

```
137 %% @doc There are no custom messages here.
138 -spec handle_custom_message(custom_message(), rt_loop:state_init()) -> unknown_event.
139 handle_custom_message(_Message, _State) -> unknown_event.
```

Custom messages could be send from a routing table process on one node to the routing table process on another node and are independent from any other implementation.

8.3.2. Chord routing table

The file `rt_chord.erl` implements Chord's routing.

Data types

File `rt_chord.erl`:

```
42 -opaque(key()::non_neg_integer()).
43 -opaque(rt()::gb_tree()).
44 -type(external_rt()::gb_tree()).           %% @todo: make opaque
45 -type(index() :: {pos_integer(), pos_integer()}).
46 -opaque(custom_message() ::
47         {rt_get_node_response, Index::pos_integer(), Node::node:node_type()}).
```

The routing table is a `gb_tree`. Identifiers in the ring are integers. Note that in Erlang integer can be of arbitrary precision. For Chord, the identifiers are in $[0, 2^{128})$, i.e. 128-bit strings.

The `rt_beh` behaviour for Chord (excerpt)

File `rt_chord.erl`:

```
55 %% @doc Creates an empty routing table.
56 -spec empty(node:node_type()) -> rt().
57 empty(_Succ) -> gb_trees:empty().
```

File `rt_chord.erl`:

```
210 -spec empty_ext(node:node_type()) -> external_rt().
211 empty_ext(_Succ) -> gb_trees:empty().
```

`empty/1` returns an empty `gb_tree`, same for `empty_ext/1`.

`rt_chord:hash_key/1`, `rt_chord:get_random_node_id/0`, `rt_chord:get_keys_for_replicas/1` and `rt_chord:n/0` are implemented like their counterparts in `rt_simple.erl`.

File `rt_chord.erl`:

```
215 %% @doc Returns the next hop to contact for a lookup.
216 %%     Note, that this code will be called from the dht_node process and
217 %%     it will thus have an external_rt!
218 -spec next_hop(dht_node_state:state(), key()) -> comm:mypid().
219 next_hop(State, Id) ->
220     RT = dht_node_state:get(State, rt),
221     case get_size(RT) == 0 orelse
222         intervals:in(Id, dht_node_state:get(State, succ_range)) of
223     true -> dht_node_state:get(State, succ_pid); % -> succ
224     _ -> % check routing table:
225         case util:gb_trees_largest_smaller_than(Id, RT) of
226             {value, _Key, Node} -> node:pidX(Node);
227             nil -> % forward to largest finger
228                 {_Key, Node} = gb_trees:largest(RT),
229                 node:pidX(Node)
230         end
231     end.
```

If the (external) routing table contains at least one item, the next hop is retrieved from the `gb_tree`. It will be the node with the largest id that is smaller than the id we are looking for. If the routing table is empty, the successor is chosen. However, if we haven't found the key in our routing table, the next hop will be our largest finger, i.e. entry.

File `rt_chord.erl`:

```
86 %% @doc Starts the stabilization routine.
87 -spec init_stabilize(key(), node:node_type(), rt()) -> rt().
88 init_stabilize(Id, _Succ, RT) ->
89     % calculate the longest finger
90     Key = calculateKey(Id, first_index()),
91     % trigger a lookup for Key
92     lookup:unreliable_lookup(Key, {rt_get_node, comm:this(), first_index()}),
93     RT.
```

The routing table stabilization is triggered for the first index and then runs asynchronously, as we do not want to block the `rt_loop` to perform other request while recalculating the routing table.

We have to find the node responsible for the calculated finger and therefore perform a lookup for the node with a `rt_get_node` message, including a reference to ourselves as the reply-to address and the index to be set.

The lookup performs an overlay routing by passing the message until the responsible node is found. There, the message is delivered to the `dht_node`. At the destination the message is handled in `dht_node.erl`:

File `dht_node.erl`:

```
110 on({rt_get_node, Source_PID, Index}, State) ->
111     comm:send(Source_PID, {rt_get_node_response, Index, dht_node_state:get(State, node)}),
112     State;
```

The remote node sends the requested information back directly. It includes a reference to itself in a `rt_get_node_response` message which will be handled by `rt_chord:handle_custom_message/2` that calls `rt_chord:stabilize/5`:

File `rt_chord.erl`:

```
192 %% @doc Chord reacts on 'rt_get_node_response' messages in response to its
193 %% 'rt_get_node' messages.
194 -spec handle_custom_message(custom_message(), rt_loop:state_init()) -> rt_loop:state_init().
195 handle_custom_message({rt_get_node_response, Index, Node}, State) ->
196     OldRT = rt_loop:get_rt(State),
197     Id = rt_loop:get_id(State),
198     Succ = rt_loop:get_succ(State),
199     Pred = rt_loop:get_pred(State),
200     NewRT = stabilize(Id, Succ, OldRT, Index, Node),
201     check(OldRT, NewRT, Id, Pred, Succ),
202     rt_loop:set_rt(State, NewRT).
```

File `rt_chord.erl`:

```
142 %% @doc Updates one entry in the routing table and triggers the next update.
143 -spec stabilize(key(), node:node_type(), rt(), pos_integer(), node:node_type())
144     -> rt().
145 stabilize(Id, Succ, RT, Index, Node) ->
146     case node:is_valid(Node)
147     andalso (node:id(Succ) /= node:id(Node))
148     andalso (not intervals:in(
149         node:id(Node),
150         intervals:mk_from_node_ids(Id, node:id(Succ))) of
151     true ->
152         NewRT = gb_trees:enter(Index, Node, RT),
153         Key = calculateKey(Id, next_index(Index)),
154         lookup:unreliable_lookup(Key, {rt_get_node, comm:this(),
155             next_index(Index)}),
156         NewRT;
157     _ -> RT
158 end.
```

`stabilize/5` assigns the received routing table entry and triggers the routing table stabilization for the the next shorter entry using the same mechanisms as described above.

If the shortest finger is the successor, then filling the routing table is stopped, as no further new entries would occur. It is not necessary, that `Index` reaches 1 to make that happen. If less than 2^{128} nodes participate in the system, it may happen earlier.

File `rt_chord.erl`:

```
97 %% @doc Removes dead nodes from the routing table.
98 -spec filter_dead_node(rt(), comm:mypid()) -> rt().
99 filter_dead_node(RT, DeadPid) ->
100     DeadIndices = [Index || {Index, Node} <- gb_trees:to_list(RT),
101         node>equals(Node, DeadPid)],
102     lists:foldl(fun(Index, Tree) -> gb_trees:delete(Index, Tree) end,
103         RT, DeadIndices).
```

`filter_dead_node` removes dead entries from the `gb_tree`.

File `rt_chord.erl`:

```
235 -spec export_rt_to_dht_node(rt(), key(), node:node_type(), node:node_type())
236     -> external_rt().
237 export_rt_to_dht_node(RT, Id, Pred, Succ) ->
238     Tree = gb_trees:enter(node:id(Succ), Succ,
239         gb_trees:enter(node:id(Pred), Pred, gb_trees:empty())),
240     util:gb_trees_foldl(fun (_K, V, Acc) ->
241         % only store the ring id and the according node structure
242         case node:id(V) == Id of
243             true  -> Acc;
244             false -> gb_trees:enter(node:id(V), V, Acc)
245         end
246     end, Tree, RT).
```

`export_rt_to_dht_node` converts the internal `gb_tree` structure based on indices into the external representation optimised for look-ups, i.e. a `gb_tree` with node ids and the nodes themselves.

8.4. Local Datastore

8.5. Cyclon

8.6. Vivaldi Coordinates

8.7. Estimated Global Information (Gossiping)

8.8. Load Balancing

8.9. Broadcast Trees

9. Transactions in Scalaris

9.1. The Paxos Module

9.2. Transactions using Paxos Commit

9.3. Applying the Tx-Modules to replicated DHTs

Introduces transaction processing on top of a Overlay

10. How a node joins the system

10.1. General Erlang server loop

Servers in Erlang often use the following structure to maintain a state while processing received messages:

```
loop(State) ->
  receive
    Message ->
      State1 = f(State),
      loop(State1)
  end.
```

The server runs an endless loop, that waits for a message, processes it and calls itself using tail-recursion in each branch. The loop works on a `State`, which can be modified when a message is handled.

10.2. Starting additional local nodes after boot

After booting a new Scalaris-System as described in Section 2.5.1 on page 9, ten additional local nodes can be started by typing `admin:add_nodes(10)` in the Erlang-Shell that the boot process opened¹.

```
scalaris/bin> ./boot.sh
[...]
```

```
=INFO REPORT==== 12-May-2009::16:24:18 ===
Yaws: Listening to 0.0.0.0:8000 for servers
- http://localhost:8000 under ../docroot
[info] [ CC ] this() == {{127,0,0,1},14195}
[info] [ DNC <0.96.0> ] starting DeadNodeCache
[info] [ DNC <0.96.0> ] starting Dead Node Cache
[info] [ RM <0.97.0> ] starting ring maintainer

[info] [ RT <0.99.0> ] starting routingtable
[info] [ Node <0.101.0> ] joining 315238232250031455306327244779560426902
[info] [ Node <0.101.0> ] join as first 315238232250031455306327244779560426902
[info] [ FD <0.74.0> ] starting pinger for {{127,0,0,1},14195,<0.101.0>}
[info] [ Node <0.101.0> ] joined
[info] [ CY ] Cyclon spawn: {{127,0,0,1},14195,<0.102.0>}
(boot@csr-pc9)1> admin:add_nodes(10)
```

In the following we will trace what this function does in order to add additional nodes to the system.

The function `admin:add_nodes(int)` is defined as follows.

File `admin.erl`:

```
33 % @doc add new Scalaris nodes on the local node
```

¹Increase the log level to `info` to get the detailed startup logs. See Section 2.7 on page 11

```

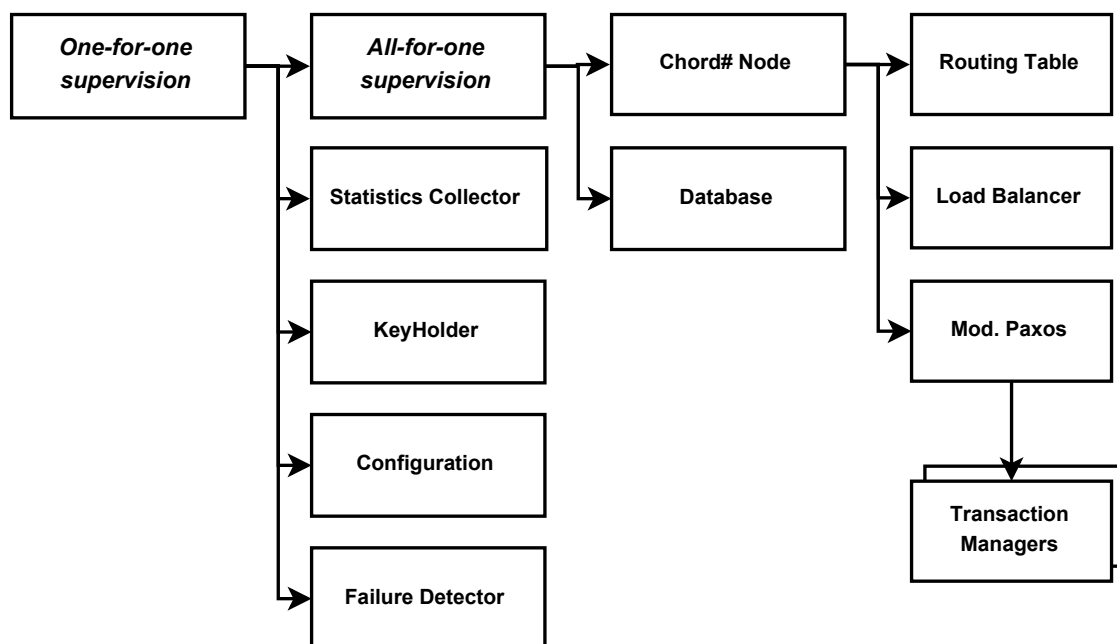
34 -spec add_nodes(non_neg_integer()) -> ok.
35 add_nodes(0) ->
36     ok;
37 add_nodes(Count) ->
38     [ begin
39         Desc = util:sup_supervisor_desc(randoms:getRandomId(),
40                                         sup_dht_node, start_link),
41         supervisor:start_child(main_sup, Desc)
42     end || _ <- lists:seq(1, Count) ],
43     ok.

```

It calls `add_nodes_loop(Count, Delay)` with a delay of 0. This function starts a new child for the main supervisor `main_sup`. As defined by the parameters, to actually perform the start, the function `sup_dht_node:start_link(i)` is called by the Erlang supervisor mechanism. For more details on the OTP supervisor mechanism see Chapter 18 of the Erlang book [1] or the online documentation at <http://www.erlang.org/doc/man/supervisor.html>.

10.2.1. Supervisor-tree of a Scalaris node

When starting a new node in the system, the following supervisor tree is created:



10.2.2. Starting the or-supervisor and general processes of a node

Starting supervisors is a two step process: the supervisor mechanism first calls the `init()` function of the defined module (`sup_dht_node:init(i)` in this case) and then calls the start function (`start_link` here).

So, let's have a look at `sup_dht_node:init(,)` the 'Scalaris or supervisor'.

File `sup_dht_node.erl`:

```

40 -spec init([any()]) -> {ok, {{one_for_one, MaxRetries::pos_integer(),
41                               PeriodInSeconds::pos_integer()},
42                               [ProcessDescr::any()]} }.

```



```

43 init([Options]) ->
44     InstanceId = string:concat("dht_node_", randoms:getRandomId()),
45     process_dictionary:register_process(InstanceId, sup_dht_node, self()),
46     boot_server:connect(),
47     KeyHolder =
48         util:sup_worker_desc(idholder, idholder, start_link,
49                               [InstanceId]),
50     Supervisor_AND =
51         util:sup_supervisor_desc(cs_supervisor_and, sup_dht_node_core, start_link,
52                                   [InstanceId, Options]),
53     RingMaintenance =
54         util:sup_worker_desc(?RM, ?RM, start_link, [InstanceId]),
55     RoutingTable =
56         util:sup_worker_desc(routingtable, rt_loop, start_link, [InstanceId]),
57     DeadNodeCache =
58         util:sup_worker_desc(deadnodecache, dn_cache, start_link, [InstanceId]),
59     Vivaldi =
60         util:sup_worker_desc(vivaldi, vivaldi, start_link, [InstanceId]),
61     Reregister =
62         util:sup_worker_desc(dht_node_reregister, dht_node_reregister, start_link,
63                               [InstanceId]),
64     DC_Clustering =
65         util:sup_worker_desc(dc_clustering, dc_clustering, start_link,
66                               [InstanceId]),
67     Cyclon =
68         util:sup_worker_desc(cyclon, cyclon, start_link, [InstanceId]),
69     Gossip =
70         util:sup_worker_desc(gossip, gossip, start_link, [InstanceId]),
71     {ok, {{one_for_one, 10, 1},
72         [
73             Reregister,
74             KeyHolder,
75             RoutingTable,
76             Supervisor_AND,
77             Cyclon,
78             DeadNodeCache,
79             RingMaintenance,
80             Vivaldi,
81             DC_Clustering,
82             Gossip
83             %% _RSE
84         ]}}}.

```

The return value of the `init()` function specifies the child processes of the supervisor and how to start them. Here, we define a list of processes to be observed by a `one_for_one` supervisor. The processes are: `KeyHolder`, `DeadNodeCache`, `RingMaintenance`, `RoutingTable`, and a `Supervisor_AND` process.

The term `{one_for_one, 10, 1}` specifies that the supervisor should try 10 times to restart each process before giving up. `one_for_one` supervision means, that if a single process stops, only that process is restarted. The other processes run independently.

The `sup_dht_node:init()` is finished and the supervisor module, starts all the defined processes by calling the functions that were defined in the list of the `sup_dht_node:init()`.

For a join of a new node, we are only interested in the starting of the `Supervisor_AND` process here. At that point in time, all other defined processes are already started and running.

10.2.3. Starting the and-supervisor with a peer and its local database

Again, the OTP will first call the `init()` function of the corresponding module:

File `sup_dht_node_core.erl`:

```

40 -spec init([instanceid() | [any()]]) -> {ok, {{one_for_all, MaxRetries::pos_integer(),
41                                     PeriodInSeconds::pos_integer()},
42                                     [ProcessDescr::any()]} }.
43 init([InstanceId, Options]) ->
44     process_dictionary:register_process(InstanceId, sup_dht_node_core, self()),
45     Proposer =
46         util:sup_worker_desc(proposer, proposer, start_link, [InstanceId]),
47     Acceptor =
48         util:sup_worker_desc(acceptor, acceptor, start_link, [InstanceId]),
49     Learner =
50         util:sup_worker_desc(learner, learner, start_link, [InstanceId]),
51     Node =
52         util:sup_worker_desc(dht_node, dht_node, start_link,
53                             [InstanceId, Options]),
54     Delayer =
55         util:sup_worker_desc(msg_delay, msg_delay, start_link,
56                             [InstanceId]),
57     TX =
58         util:sup_supervisor_desc(sup_dht_node_core_tx, sup_dht_node_core_tx, start_link,
59                                 [InstanceId]),
60     {ok, {{one_for_all, 10, 1},
61         [
62             Proposer, Acceptor, Learner,
63             Node,
64             Delayer,
65             TX
66         ]}}.

```

It defines three processes, that have to be observed using an `one_for_all`-supervisor, which means, that if one fails, all have to be restarted. Passed to the `init` function is the `InstanceId`, a random number to make nodes unique. It was calculated a bit earlier in the code. Exercise: Try to find where.

As you can see from the list, the DB is started before the `Node`. This is intended and important, because `dht_node` uses the database, but not vice versa. The supervisor first completely initializes the DB process and afterwards calls `dht_node:start_link()`. We only go into details here, for the latter.

File `dht_node.erl`:

```

384 %% @doc spawns a scalaris node, called by the scalaris supervisor process
385 -spec start_link(instanceid()) -> {ok, pid()}.
386 start_link(InstanceId) ->
387     start_link(InstanceId, []).
388
389 -spec start_link(instanceid(), [any()]) -> {ok, pid()}.
390 start_link(InstanceId, Options) ->
391     gen_component:start_link(?MODULE, [InstanceId, Options],
392                             [{register, InstanceId, dht_node}, wait_for_init]).

```

`dht_node` implements the `gen_component` behaviour. This component was developed by us to enable us to write code which is similar in syntax and semantics to the examples in [2]. Similar to the supervisor behaviour, the component has to provide an `init` function, but here it is used to initialize the state of the component. This function is described in the next section.

Note: `?MODULE` is a predefined Erlang macro, which expands to the module name, the code belongs to (here: `dht_node`).

10.2.4. Initializing a `dht_node`-process

File `dht_node.erl`:

```

364 %% @doc joins this node in the ring and calls the main loop
365 -spec init([instanceid() | [any()]] -> {join, {as_first}, msg_queue:msg_queue()} | {join, {phase1},
366 init([_InstanceId, Options]) ->
367     %io:format("~p~n", [Options]),
368     % first node in this vm and also vm is marked as first
369     % or unit-test
370     case lists:member(first, Options) andalso
371         (is_unittest() orelse
372         application:get_env(boot_cs, first) == {ok, true}) of
373     true ->
374         trigger_known_nodes(),
375         idholder:get_id(),
376         {join, {as_first}, msg_queue:new()};
377     _ ->
378         idholder:get_id(),
379         {join, {phase1}, msg_queue:new()}
380     end.

```

The `gen_component` behaviour registers the `dht_node` in the process dictionary. Formerly, the process had to do this itself, but we moved this code into the behaviour. If the `dht_node` is the first node, it will start immediately by triggering all known nodes (to initialize the comm layer) and entering the join process accordingly. The node also retrieves its `Id` from the `idholder`: `idholder:get_id()`. In the first call, a random identifier is returned, otherwise the latest set value. If the `dht_node-process` failed and is restarted by its supervisor, this call to the `idholder` ensures, that the node still keeps its `Id`, assuming that the `idholder` process is not failing. This is important for the load-balancing and for consistent responsibility of nodes to ensure consistent lookup in the structured overlay.

If a node changes its position in the ring for load-balancing, the `idholder` will be informed and the `dht_node` finishes itself. This triggers a restart of the corresponding database process via the `and-supervisor`. When the supervisor restarts both processes, they will retrieve the new position in the ring from the `idholder` and join the ring there.

TODO: The supervisor was configured to restart a node at most 10 times. Does that mean, that a node can only change its position in the ring 10 times (caused by load-balancing)?

10.2.5. Actually joining the ring

After retrieving its identifier, the node starts the join protocol which processes the appropriate messages calling `dht_node_join:process_join_msg(Message, State)`.

File `dht_node_join.erl`:

```

71 process_join_msg({idholder_get_id_response, Id, IdVersion},
72                 {join, {as_first}, QueuedMessages}) ->
73     log:log(info, "[ Node ~w ] joining as first: ~p", [self(), Id]),
74     Me = node:new(comm:this(), Id, IdVersion),
75     rt_beh:initialize(Id, Me, Me),
76     NewState = dht_node_state:new(?RT:empty_ext(Me),
77                                   nodelist:new_neighborhood(Me),
78                                   ?DB:new(Id)),
79     comm:send_local(get_local_dht_node_reregister_pid(), {go}),
80     msg_queue:send(QueuedMessages),
81     %log:log(info, "[ Node ~w ] joined", [self()]),
82     NewState; % join complete, State is the first "State"

```

If the ring is empty, the joining node is the only node in the ring and will be responsible for the whole key space. `join_first` just creates a new state for a `Scalaris` node consisting of an empty routing table, a `successorlist` containing itself, itself as its predecessor, a reference to itself, its responsibility area from `Id` to `Id` (the full ring), and a load balancing schema.

The state is defined in

File dht_node_state.erl:

```
56 -spec new(?RT:external_rt(), Neighbors::nodelist:neighborhood(),
57         ?DB:db()) -> state().
58 new(RT, Neighbors, DB) ->
59     #state{rt = RT,
60           neighbors = Neighbors,
61           join_time = now(),
62           trans_log = #translog{tid_tm_mapping = dict:new(),
63                                decided         = gb_trees:empty(),
64                                undecided        = gb_trees:empty()
65                                },
66           db = DB,
67           tx_tp_db = tx_tp:init(),
68           proposer = process_dictionary:get_group_member(paxos_proposer)
69     }.
```

If a node joins an existing ring, it will at first try to contact all dht_node processes in any VM configured in known_hosts.

File dht_node_join.erl:

```
89 % 1. get my key
90 process_join_msg({idholder_get_id_response, Id, IdVersion},
91                 {join, {phase1}, QueuedMessages}) ->
92     %io:format("p1: got key~n"),
93     log:log(info, "[ Node ~w ] joining", [self()]),
94     % send message to avoid code duplication
95     comm:send_local(self(), {known_hosts_timeout}),
96     {join, {phase2, Id, IdVersion}, QueuedMessages};
97
98 % 2. Find known hosts
99 process_join_msg({known_hosts_timeout},
100                {join, {phase2, _Id, _IdVersion}, _QueuedMessages} = State) ->
101     %io:format("p2: known hosts timeout~n"),
102     KnownHosts = config:read(known_hosts),
103     % contact all known VMs
104     _Res = [comm:send(KnownHost, {get_dht_nodes, comm:this()})
105            || KnownHost <- KnownHosts],
106     %io:format("~p~n", [_Res]),
107     % timeout just in case
108     comm:send_local_after(1000, self(), {known_hosts_timeout}),
109     State;
110
111 process_join_msg({get_dht_nodes_response, []},
112                 {join, {phase2, _Id, _IdVersion}, _QueuedMessages} = State) ->
113     %io:format("p2: got empty dht_nodes_response~n"),
114     % there is a VM with no nodes
115     State;
116
117 process_join_msg({get_dht_nodes_response, Nodes = [_|_]},
118                 {join, {phase2, Id, IdVersion}, QueuedMessages} = State) ->
119     %io:format("p2: got dht_nodes_response ~p~n", [lists:delete(comm:this(), Nodes)]),
120     case lists:delete(comm:this(), Nodes) of
121     [] ->
122         State;
123     [First | Rest] ->
124         comm:send(First, {lookup_aux, Id, 0, {get_node, comm:this(), Id}}),
125         comm:send_local_after(3000, self(), {lookup_timeout}),
126         {join, {phase3, Rest, Id, IdVersion}, QueuedMessages}
127     end;
```

These nodes will be send a lookup request for the node currently responsible for the new node's id – the successor for the joining node. If this lookup fails for some reason, it is tried again.

File dht_node_join.erl:

```
131 % 3. lookup my position
132 process_join_msg({lookup_timeout},
133                 {join, {phase3, [], Id, IdVersion}, QueuedMessages}) ->
134     %io:format("p3: lookup_timeout~n"),
135     % no more nodes left, go back to step 2
136     comm:send_local(self(), {known_hosts_timeout}),
137     {join, {phase2, Id, IdVersion}, QueuedMessages};
138
139 process_join_msg({get_node_response, Id, Succ},
140                 {join, {phase3, _DHTNodes, Id, IdVersion}, QueuedMessages}) ->
141     %io:format("p3: lookup success~n"),
142     % got my successor
143     Me = node:new(comm:this(), Id, IdVersion),
144     % announce join request
145     comm:send(node:pidX(Succ), {join, Me}),
146     {join, {phase4, Succ, Me}, QueuedMessages};
```

If its (future) successor is found, this new node will send a join message including a reference to itself and the chosen Id. This message is received by the old node in dht_node.erl:

File dht_node.erl:

```
341 on({join, NewPred}, State) ->
342     dht_node_join:join_request(State, NewPred);
```

This triggers a call to dht_node_join:join_request/2 on the old node.

File dht_node_join.erl:

```
47 -spec join_request(dht_node_state:state(), NewPred::node:node_type()) ->
48     dht_node_state:state().
49 join_request(State, NewPred) ->
50     MyNewInterval =
51         intervals:mk_from_nodes(NewPred, dht_node_state:get(State, node)),
52     {DB, HisData} = ?DB:split_data(dht_node_state:get(State, db), MyNewInterval),
53
54     %TODO: split data [{Key, Value, Version}], schedule transfer
55
56     comm:send(node:pidX(NewPred),
57               {join_response, dht_node_state:get(State, pred), HisData}),
58     % TODO: better already update our range here directly than waiting for an
59     % updated state from the ring_maintenance?
60     rm_beh:notify_new_pred(comm:this(), NewPred),
61     dht_node_state:set_db(State, DB).
```

The dht_node will update the interval it is responsible for and notify the ring maintenance of its new predecessor. It will also remove all key-value pairs from its database which are now in the responsibility of the joining node and send a join_response message to the new node with its former predecessor and the data the new node has to host.

File dht_node_join.erl:

```
150 % 4. joining my neighbors
151 process_join_msg({join_response, Pred, Data},
152                 {join, {phase4, Succ, Me}, QueuedMessages}) ->
153     Id = node:id(Me),
154     %io:format("p4: join_response~n"),
155     % @TODO data shouldn't be moved here, might be large
156     log:log(info, "[ Node ~w ] got pred ~w", [self(), Pred]),
157     DB = ?DB:add_data(?DB:new(Id), Data),
158     rt_beh:initialize(Id, Pred, Succ),
159     rm_beh:notify_new_succ(node:pidX(Pred), Me),
160     State = dht_node_state:new(?RT:empty_ext(Succ),
161                                nodelist:new_neighborhood(Pred, Me, Succ), DB),
```

```
162     cs_replica_stabilization:recreate_replicas(dht_node_state:get(State, my_range)),
163     comm:send_local(get_local_dht_node_reregister_pid(), {go}),
164     msg_queue:send(QueuedMessages),
165     State;
```

Back on the joining node: it waits for the `join_response` message in phase 4 of the join protocol. The next steps after the message is received from the old node are to initialize the maintenance components for the ring and routing table, the database and the state of the `dht_node`. The `cs_replica_stabilization:recreate_replicas()` function is called, which is not yet implemented. It would recreate necessary replicas that were lost due to load-balancing, node failures and lost updates during the data transfer.

The macro `?RT` maps to the configured routing algorithm and `?RM` to the configured ring maintenance algorithm. It is defined in `include/scalaris.hrl`. For further details on the routing see Chapter 8.3 on page 29.

Note that join-related messages arriving in other phases than those handling them will be ignored. Any other messages during a `dht_node`'s join will be queued and re-send when the join is complete.

TODO: What, if the `Id` is exactly the same as that of the existing node? This could lead to lookup and responsibility inconsistency? Can this be triggered by the load-balancing? This is a bug, that should be fixed!!!

11. Directory Structure of the Source Code

The directory tree of Scalaris is structured as follows:

bin	contains shell scripts needed to work with Scalaris (e.g. start the boot services, start a node, ...)
contrib	necessary third party packages (yaws and log4erl)
doc	generated Erlang documentation
docroot	root directory of the boot server's webserver
docroot_node	root directory of the normal node's webserver
ebin	the compiled Erlang code (beam files)
java-api	a Java API to Scalaris
log	log files
src	contains the Scalaris source code
test	unit tests for Scalaris
user-dev-guide	contains the sources for this document

12. Java API

For the Java API documentation, we refer the reader to the documentation generated by javadoc or doxygen. The following commands create the documentation:

```
%> cd java-api  
%> ant doc  
%> doxygen
```

The documentation can then be found in `java-api/doc/index.html` (javadoc) and `java-api/doc-doxygen/html/index.html` (doxygen).

We provide two kinds of APIs:

- high-level access with `de.zib.scalarisc.Scalaris`
- low-level access with `de.zib.scalarisc.Transaction`

The former provides general functions for reading, writing and deleting single key-value pairs and an API for the built-in PubSub-service. The latter allows the user to write custom transactions which can modify an arbitrary number of key-value pairs within one transaction.

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