

ASSIGNMENT 5

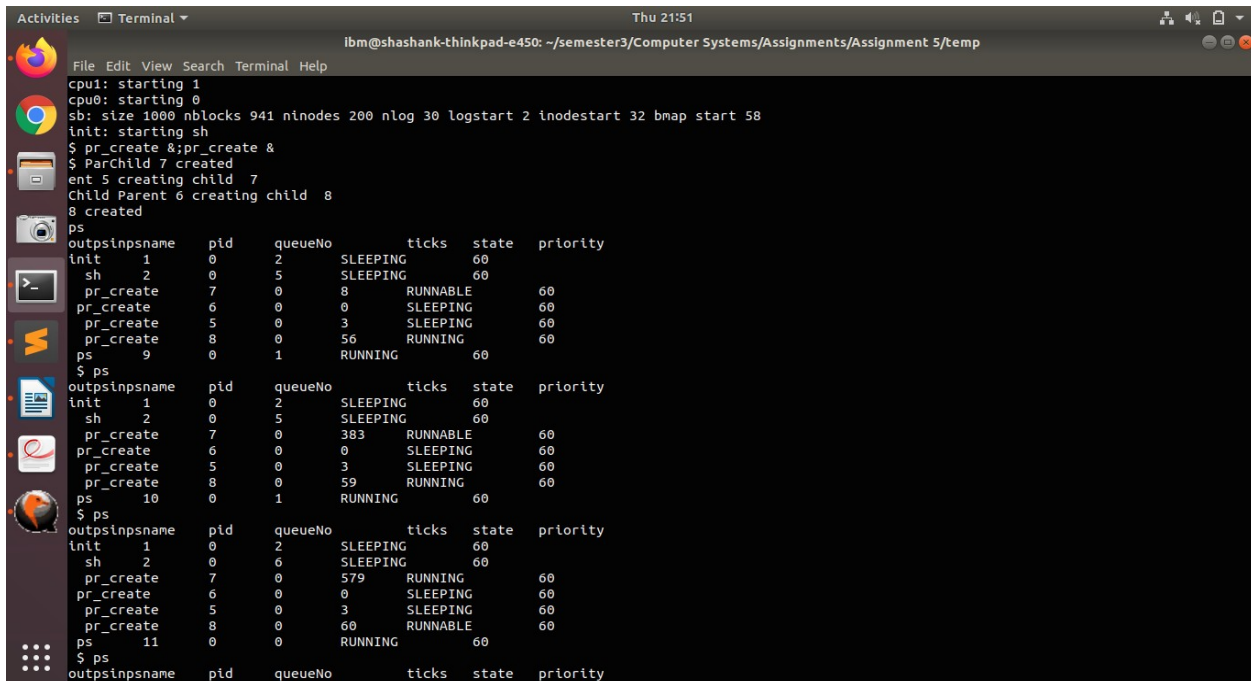
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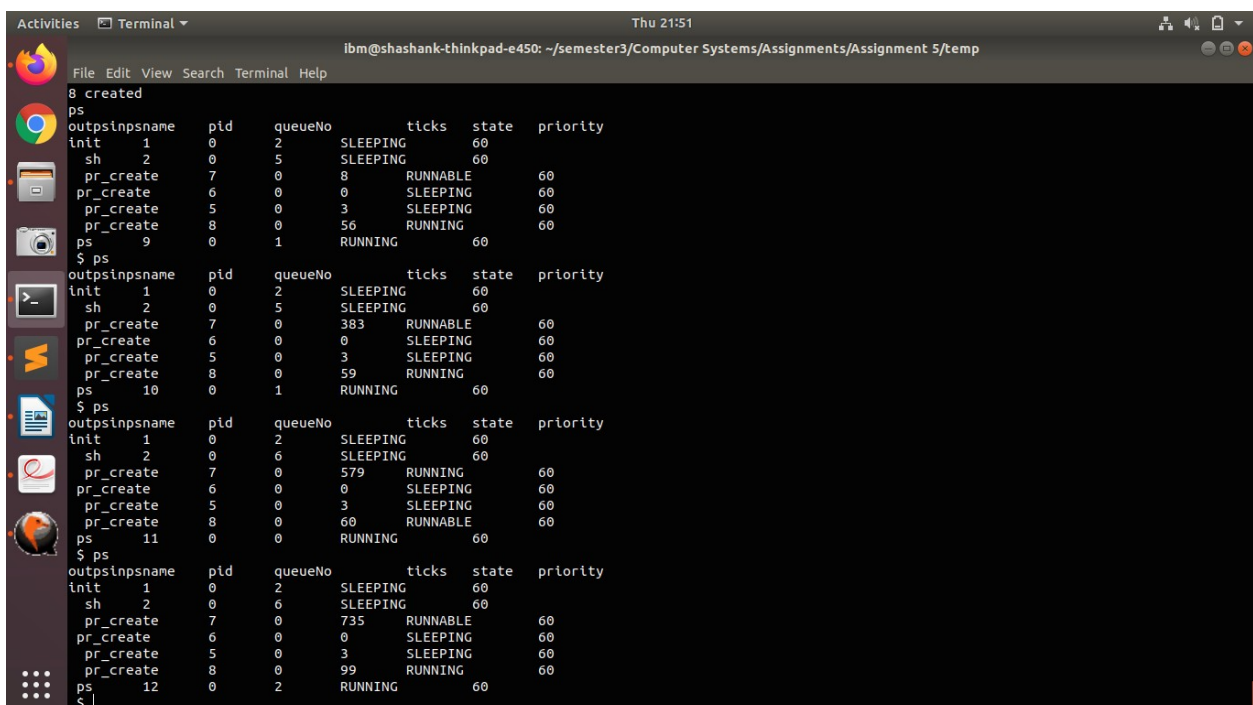
Performance Comparison of SCHEDULING POLICIES

DEFAULT

The default scheduling is Round robin, where each process is provided a fixed time to execute called “QUANTUM”. Once a process is executed for a given time period, it is preempted and other process executes for a given time period. The screenshots of processes info when DEFAULT is used:



```
ibm@shashank-thinkpad-e450: ~/semester3/Computer Systems/Assignments/Assignment 5/temp
cpu1: starting 1
cpu0: starting 0
sb: size 1000 nblocks 941 ninodes 200 nlog 30 logstart 2 inodestart 32 bmap start 58
init: starting sh
$ pr_create & pr_create &
$ ParChild 7 created
ent 5 creating child 7
Child Parent 6 creating child 8
8 created
ps
outpsinspname pid queueNo ticks state priority
init 1 0 2 SLEEPING 60
sh 2 0 5 SLEEPING 60
pr_create 7 0 8 RUNNABLE 60
pr_create 6 0 0 SLEEPING 60
pr_create 5 0 3 SLEEPING 60
pr_create 8 0 56 RUNNING 60
ps 9 0 1 RUNNING 60
$ ps
outpsinspname pid queueNo ticks state priority
init 1 0 2 SLEEPING 60
sh 2 0 5 SLEEPING 60
pr_create 7 0 383 RUNNABLE 60
pr_create 6 0 0 SLEEPING 60
pr_create 5 0 3 SLEEPING 60
pr_create 8 0 59 RUNNING 60
ps 10 0 1 RUNNING 60
$ ps
outpsinspname pid queueNo ticks state priority
init 1 0 2 SLEEPING 60
sh 2 0 6 SLEEPING 60
pr_create 7 0 579 RUNNABLE 60
pr_create 6 0 0 SLEEPING 60
pr_create 5 0 3 SLEEPING 60
pr_create 8 0 60 RUNNABLE 60
ps 11 0 0 RUNNING 60
$ ps
outpsinspname pid queueNo ticks state priority
```



```
ibm@shashank-thinkpad-e450: ~/semester3/Computer Systems/Assignments/Assignment 5/temp
8 created
ps
outpsinspname pid queueNo ticks state priority
init 1 0 2 SLEEPING 60
sh 2 0 5 SLEEPING 60
pr_create 7 0 8 RUNNABLE 60
pr_create 6 0 0 SLEEPING 60
pr_create 5 0 3 SLEEPING 60
pr_create 8 0 56 RUNNING 60
ps 9 0 1 RUNNING 60
$ ps
outpsinspname pid queueNo ticks state priority
init 1 0 2 SLEEPING 60
sh 2 0 6 SLEEPING 60
pr_create 7 0 383 RUNNABLE 60
pr_create 6 0 0 SLEEPING 60
pr_create 5 0 3 SLEEPING 60
pr_create 8 0 59 RUNNING 60
ps 10 0 1 RUNNING 60
$ ps
outpsinspname pid queueNo ticks state priority
init 1 0 2 SLEEPING 60
sh 2 0 6 SLEEPING 60
pr_create 7 0 579 RUNNABLE 60
pr_create 6 0 0 SLEEPING 60
pr_create 5 0 3 SLEEPING 60
pr_create 8 0 60 RUNNABLE 60
ps 11 0 0 RUNNING 60
$ ps
outpsinspname pid queueNo ticks state priority
init 1 0 2 SLEEPING 60
sh 2 0 6 SLEEPING 60
pr_create 7 0 735 RUNNABLE 60
pr_create 6 0 0 SLEEPING 60
pr_create 5 0 3 SLEEPING 60
pr_create 8 0 99 RUNNING 60
ps 12 0 2 RUNNING 60
$ |
```

In above images there are two processes with pids 7 and 8 which are RUNNING alternatively ,the “RUNNING” process gets changed once its assigned timequantum is done.

FIRST COME – FIRST SERVED (FCFS)

Processes are executed on first come, first serve basis. Other processes need to wait until the processes which are created before it are done.

The screenshots of processes info when FCFS is used:

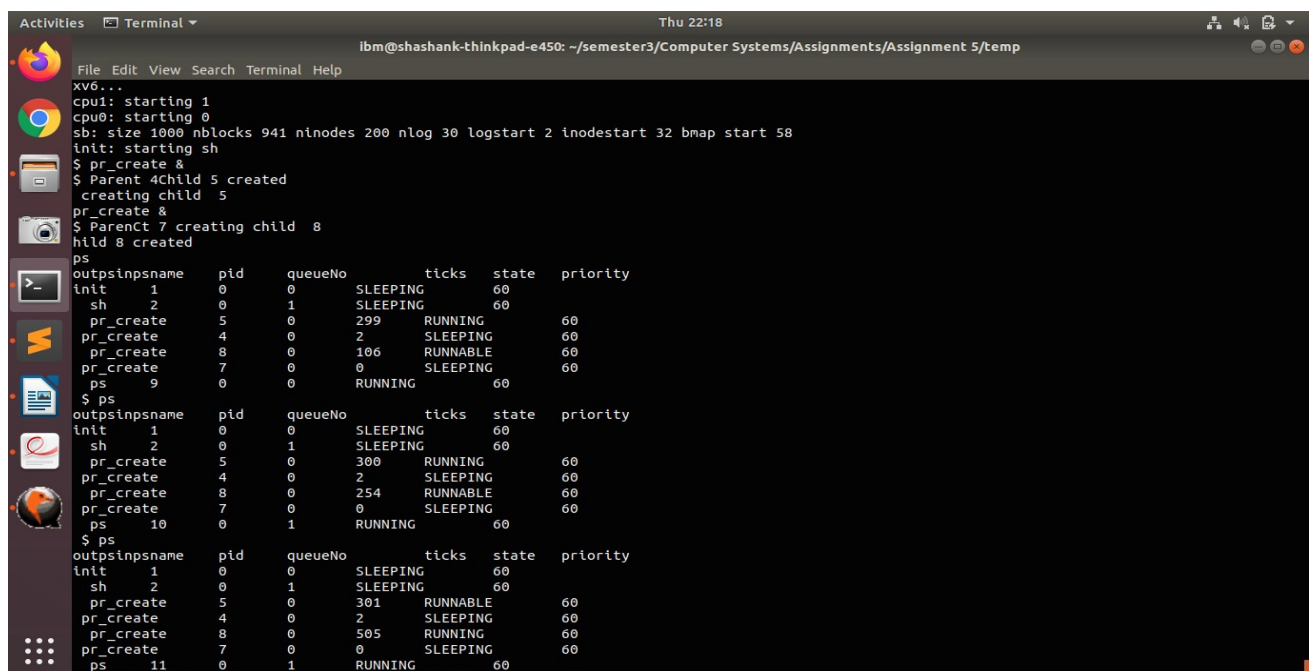
```
368+1 records in
368+1 records out
188672 bytes (189 kB, 184 KiB) copied, 0.00110625 s, 171 MB/s
qemu-system-i386 -serial mon:stdio -drive file=fs.img,index=1,media=disk,format=raw -drive file=xv6.img,index=0,media=disk,format=raw -smp 2 -m 512
xv6...
cpu1: starting 1
cpu0: starting 0
sb: size 1000 nblocks 941 ninodes 200 nlog 30 logstart 2 inodestart 32 bmap start 58
init: starting sh
$ pr_create &
$ Parent 4 creating child 5
Child 5 created
pr_create &
$ Parent 7 creating child 8
Child 8 created
ps
ps
ps
```

Here two processes are created with pids 5 and 8, and later no other process is getting executed as the previous started processes are not finished and the # of CPU's used are 2 so atmost 2 processes can execute at a given time.

PRIORITY BASED SCHEDULER (PBS)

A priority-based scheduler selects the process with the highest priority for execution. Until the processes of higher priority are done lower priority are not done.

The screenshots of processes info when PBS is used:



```
ibm@shashank-thinkpad-e450: ~/semester3/Computer Systems/Assignments/Assignment 5/temp
xv6...
cpu1: starting 1
cpu0: starting 0
sb: size 1000 nblocks 941 ninodes 200 nlog 30 logstart 2 inodestart 32 bmap start 58
init: starting sh
$ pr_create &
$ Parent 4child 5 created
$ creating child 5
pr_create &
$ Parent 7 creating child 8
child 8 created
ps
$ ps
  outpsname  pid  queueNo  ticks  state  priority
init         1    0         0  SLEEPING  60
sh           2    0         1  SLEEPING  60
pr_create    5    0        299  RUNNING  60
pr_create    4    0         2  SLEEPING  60
pr_create    8    0        106  RUNNABLE  60
pr_create    7    0         0  SLEEPING  60
ps           9    0         0  RUNNING  60
$ ps
  outpsname  pid  queueNo  ticks  state  priority
init         1    0         0  SLEEPING  60
sh           2    0         1  SLEEPING  60
pr_create    5    0        300  RUNNING  60
pr_create    4    0         2  SLEEPING  60
pr_create    8    0        254  RUNNABLE  60
pr_create    7    0         0  SLEEPING  60
ps          10    0         1  RUNNING  60
$ ps
  outpsname  pid  queueNo  ticks  state  priority
init         1    0         0  SLEEPING  60
sh           2    0         1  SLEEPING  60
pr_create    5    0        301  RUNNABLE  60
pr_create    4    0         2  SLEEPING  60
pr_create    8    0        505  RUNNING  60
pr_create    7    0         0  SLEEPING  60
ps          11    0         1  RUNNING  60
```

Above two processes with pids 5 and 8 are created and they both run alternatively following Round robin as their priorities are same (i.e 60)

Now set_pr command is used to update the priority of process with pid 5 to 30.

The terminal window shows the output of the 'ps' command at four different stages. The first two outputs show processes with priority 60. The third output shows the process with pid 5 having its priority changed to 30. The fourth output shows the process with pid 8 in a 'RUNNABLE' state, indicating it is waiting for the CPU.

```

$ ps
  outps\nsname  pid  queueNo  ticks  state  priority
  ----
  init          1    0          0  SLEEPING  60
  sh            2    0          0  SLEEPING  60
  pr_create     5    0        301  RUNNABLE  60
  pr_create     4    0          2  SLEEPING  60
  pr_create     8    0       505  RUNNING  60
  pr_create     7    0          0  SLEEPING  60
  ps           11    0          1  RUNNING  60

$ set_pr 5 30
here
$ ps
  outps\nsname  pid  queueNo  ticks  state  priority
  ----
  init          1    0          0  SLEEPING  60
  sh            2    0          3  SLEEPING  60
  pr_create     5    0        392  RUNNING  30
  pr_create     4    0          2  SLEEPING  60
  pr_create     8    0      1351  RUNNABLE  60
  pr_create     7    0          0  SLEEPING  60
  ps           13    0          1  RUNNING  60

$ ps
  outps\nsname  pid  queueNo  ticks  state  priority
  ----
  init          1    0          0  SLEEPING  60
  sh            2    0          3  SLEEPING  60
  pr_create     5    0      1425  RUNNING  30
  pr_create     4    0          2  SLEEPING  60
  pr_create     8    0      1933  RUNNABLE  60
  pr_create     7    0          0  SLEEPING  60
  ps           14    0          0  RUNNING  60

$ ps
  outps\nsname  pid  queueNo  ticks  state  priority
  ----
  init          1    0          0  SLEEPING  60
  sh            2    0          3  SLEEPING  60
  pr_create     5    0      1720  RUNNING  30
  pr_create     4    0          2  SLEEPING  60
  pr_create     8    0      1933  RUNNABLE  60
  pr_create     7    0          0  SLEEPING  60
  ps           15    0          0  RUNNING  60
  
```

Now only process of pid 5 is RUNNING and pid 8 process is not RUNNING as one CPU is assigned to pid 5 and other is for present command ,so pid 8 remains in RUNNABLE until process of pid 5 is done its execution.

MULTI-LEVEL FEEDBACK QUEUE SCHEDULING

MLFQ scheduler allows processes to move between different priority queues based

on their behavior and CPU bursts. If a process uses too much CPU time, it is pushed

to a lower priority queue, leaving I/O bound and interactive processes for higher priority queues. Also, to prevent starvation, it implements aging.

The screenshots of processes info when MLFQ is used:

Here initially only one process is created with pid 5, and can observe that it changes its “CURRENT QUEUE” based on the ticks performed.

```
ibm@shashank-thinkpad-e450: ~/semester3/Computer Systems/Assignments/Assignment 5/temp
pr_create 13 4 1576 RUNNING 60
$ QEMU 2.11.1 monitor - type 'help' for more information
(qemu) q
(base) ibm@shashank-thinkpad-e450:~/semester3/Computer Systems/Assignments/Assignment 5/temp$ make qemu SCHEDULER=MLFQ
qemu-system-i386 -serial mon:stdio -drive file=fs.img,index=1,media=disk,format=raw -drive file=xv6.img,index=0,media=disk,format=raw -smp 2 -m 512
xv6...
cpu1: starting 1
cpu0: starting 0
sb: size 1000 nblocks 941 ninodes 200 nlog 30 logstart 2 inodestart 32 bmap start 58
init: starting sh
$ pr_create &
$ Parent 4 creating child 5
child 5 created
ps
  outpsinsname  pid  queueNo  ticks  state  priority
init 1 1 1 SLEEPING 60
sh 2 0 4 SLEEPING 60
pr_create 5 2 3 RUNNING 60
pr_create 4 0 0 SLEEPING 60
ps 6 0 3 RUNNING 60
$ ps
  outpsinsname  pid  queueNo  ticks  state  priority
init 1 1 1 SLEEPING 60
sh 2 0 4 SLEEPING 60
pr_create 5 2 7 RUNNING 60
pr_create 4 0 0 SLEEPING 60
ps 7 0 1 RUNNING 60
$ ps
  outpsinsname  pid  queueNo  ticks  state  priority
init 1 1 1 SLEEPING 60
sh 2 0 4 SLEEPING 60
pr_create 5 3 29 RUNNING 60
pr_create 4 0 0 SLEEPING 60
ps 8 0 1 RUNNING 60
$ pr_create &
Parent 10 creating child 11
$ Child 11 created
```

Now another process with pid 11 is created and is in RUNNING state as it lies in queue 0 and can observe the change of queue of that process based on the ticks performed.

```
ibm@shashank-thinkpad-e450: ~/semester3/Computer Systems/Assignments/Assignment 5/temp
  outpsinsname  pid  queueNo  ticks  state  priority
init 1 1 1 SLEEPING 60
sh 2 0 19 SLEEPING 60
pr_create 5 4 3475 RUNNABLE 60
pr_create 4 0 0 SLEEPING 60
pr_create 11 0 9 RUNNING 60
pr_create 10 0 0 SLEEPING 60
ps 26 0 0 RUNNING 60
$ ps
  outpsinsname  pid  queueNo  ticks  state  priority
init 1 1 1 SLEEPING 60
sh 2 1 4 SLEEPING 60
pr_create 5 4 3475 RUNNABLE 60
pr_create 4 0 0 SLEEPING 60
pr_create 11 2 21 RUNNING 60
pr_create 10 0 0 SLEEPING 60
ps 27 0 1 RUNNING 60
$ pps
exec: fail
exec pps failed
$ ps
  outpsinsname  pid  queueNo  ticks  state  priority
init 1 1 1 SLEEPING 60
sh 2 0 25 SLEEPING 60
pr_create 5 4 4079 RUNNABLE 60
pr_create 4 0 0 SLEEPING 60
pr_create 11 2 21 RUNNING 60
pr_create 10 0 0 SLEEPING 60
ps 29 0 1 RUNNING 60
$ ps
  outpsinsname  pid  queueNo  ticks  state  priority
init 1 1 1 SLEEPING 60
sh 2 0 26 SLEEPING 60
pr_create 5 4 4384 RUNNABLE 60
pr_create 4 0 0 SLEEPING 60
pr_create 11 2 21 RUNNING 60
pr_create 10 0 0 SLEEPING 60
ps 30 1 0 RUNNING 60
$ |
```

CONCLUSION:

In FCFS no preemption is allowed so the process created should WAIT until all the processes created before it are executed. Not suitable for REAL TIME COMPUTING and NOT INTERACTIVE.

But it is easy to implement...

Avg waiting time of processes can be large if the first created process runs for a long time.

In PBS higher priority processes are executed so the processes which need to meet a deadline can be given a higher priority and can be done.

But a process with low priority keeps WAITING until all other higher priority processes are done which leads to STARVATION.....

In RR, STARVATION doesn't occur because for each round robin cycle, every process is given a fixed time to execute. No process is left behind.

But if time quantum is longer than needed, it tends to exhibit the same behavior as FCFS and if time quantum is shorter than needed, the number of times that CPU switches from one process to another process, increases. This leads to decrease in CPU efficiency as number of CONTEXT SWITCHES are increased.

In MLFQ, processes are RESPONSIVE but STARVATION of processes can't be avoided as processes in lower priority queues need to wait indefinitely until higher priority queues are done.

Additional overhead due to MULTIPLE QUEUES as process needs to be switched between the queues.