

Question 01

② Recursion involves solving a problem by breaking it down into smaller sub-problems and solving each subproblem in a recursive manner. It relies on the function calling itself to perform the repetitive task. Iteration involves using loops to repeatedly execute a block of code until a certain condition is met.

Recursion uses more memory compared to iteration. Iteration uses fixed amount of memory to store loop variables more efficiently while recursion call adds new frames to stack.

b) `for (int i=0; i<count; i++) {` (count → arr.length)
`int minLoc = i;`
`for (int j=i+1; j<count; j++) {`
`if (count arr[j] < arr[minLoc]) {`
`minLoc = j;`
`}`
`}`
`int temp = arr[i];`
`arr[i] = arr[minLoc];`
`arr[minLoc] = temp;`
`}` P R S E

② Advantages:-

1. Simplicity:- Selection sort is relatively simple to understand and implement.
2. Space Efficiency:- It operates directly on the input array without requiring additional data structures.

Disadvantage:-

Inefficient:- It has time complexity $O(n^2)$. As a result, it becomes inefficient for large data sets.

③

Original Array: 8 2 5 9 7 8

$\underbrace{8}_{\leftarrow} \ 2 \ 5 \ 9 \ 7 \ 8$ set to small

$2 \ \underbrace{8}_{\leftarrow} \ 5 \ 9 \ 7 \ 8$ = initial val.

$(2 + 5 + 8) > 9 + 7 = 8$ true

$2 \ 5 \ 7 \ 8 \ 9 \ 8$

$2 \ 5 \ 7 \ 8 \ 9 \ 8$

$2 \ 5 \ 7 \ 8 \ 9 \ 8$

$2 \ 5 \ 7 \ 8 \ 9 \ 8$

Sorted Array : 2 5 7 8 8 9.

- ③ 1 : Starting from 0 th index, compare the first and the second elements.

2 : If the first element is greater than the second element, they are swapped

3 : Now, compare the second and the third elements, swap them if they are not in order.

4 : The above process goes on until the last

⑧ C SNSP is at Best in Worst in Average

Insertion Sort $O(n)$ partition $O(n^2)$ merge $O(n^2)$

Insertion Sort and $O(n)$ partitioning $O(n^2)$ Selection Sort $O(n^2)$

Selection Sort: $O(n^2)$ avg. $O(n^2)$ best $O(n^2)$

Bubble sort is $O(n)$, $O(n^2)$, $O(n^2)$

2 hours as abd. damage to body

I am going to buy bag

③ original Array: 8 2 5 7 9 2 7 8

 $\theta_1 = \text{constant}$

2 2 5 2 9 1 7 1 8 0 5 7 0 1 6 1

2 8 5 9 7 8

18 the 1st edition 1998 2nd edn

2 5 8 9 7-8 2 1

answ. ansatz zweiter Satz null bsp für

2 5 8 9 7 8

→ as in max. plausibility at available data

and 25 in 83 days 9 8 is it

2 5 11 17 23 19 15

2 5 7 8 9 10

Question 02

(any two with box 4x4)

- a) i) Stack is a linear data structure that follows the Last-In-First-Out (LIFO) principle. where the last item added is the first one to be removed. It has two main operators - 'Push' to add elements to the top end of stack and 'pop' to remove the top most element.
- ii) Queue is a linear data structure that follows the First-In-First-Out (FIFO) principle. It is similar to a queue of people waiting in a line where the first person to join the queue is the first one to be served. Has two primary operations 'Enqueue' adds an elements to the end and 'Dequeue' removes the elements from the front.
- iii) Linked list is a dynamic data structure composed of nodes where each node contains data and a reference to the next node. Unlike arrays, linked lists does not store its elements in contiguous memory locations. Linked lists has various forms such as singly linked list, Doubly and circular. This allow to efficiently insertion and deletion of elements at any position.

18 J F E S

P D S F C

(b) ① `#define MAX-SIZE 100`

int top = -1;

② int STACK-DATA[MAX-SIZE];

③ int isEmpty() {
return top == -1;
}

④ int isFull() {
return top == MAX-SIZE - 1;
}

⑤ void push(int value) {
if (isFull()) {
printf("stack overflow");
return;
}
STACK-DATA[++top] = value;
}

⑥ int pop() {
if (isEmpty()) {
printf("stack underflow");
return;
}
return STACK-DATA[top--];
}

```

③ void empty_stack() {
    if (top == -1)
        return;
    pop();
    empty_stack();
}

```

Question 03

a)

→ Array \times $m = \text{get}$ Linked lists.

- Size is fixed → Size is not fixed
- Occupies less → Occupies more
- Deletions is not possible → Deletion of elements is possible
- Memory allocation → Memory Allocation
From stacks = $\text{[get]} \times$ From heap

b) A counter A scenario where customers visit a service counter, such as a bank. It helps organize and serve customers in the order they arrive, ensuring and minimize wait time.

② i) int isEmpty() {
 if (front == NULL && rear == NULL) {
 return 1;
} else {
return 0;
}
}
} 

ii) int peek() {
if (front == NULL) {
printf("Empty Queue");
return -1;
}
return front->data;
}


iii) void enqueue() {
struct node *temp = (struct node*) malloc((sizeof(struct node));
temp->data = data;
temp->next = NULL;

if (front == NULL && rear == NULL) {
front = rear = temp;
}
else {
rear->next = temp;
rear = temp;
}
}


②

S

Q

S.push(4)

4

S.push(10)

4 10

Q.enqueue(6)

4 10

Q.enqueue(8)

4 10

6 8

(a) x = S.pop()

4

6 8

(10)

S.push(2)

4 2

6 8

Q.enqueue(x)

4 2

6 8 10

Y = Q.dequeue()

4 2

8 10

(6)

S.push(x)

4 2 10

8 10

S.push(y)

4 2 10 6

8 10

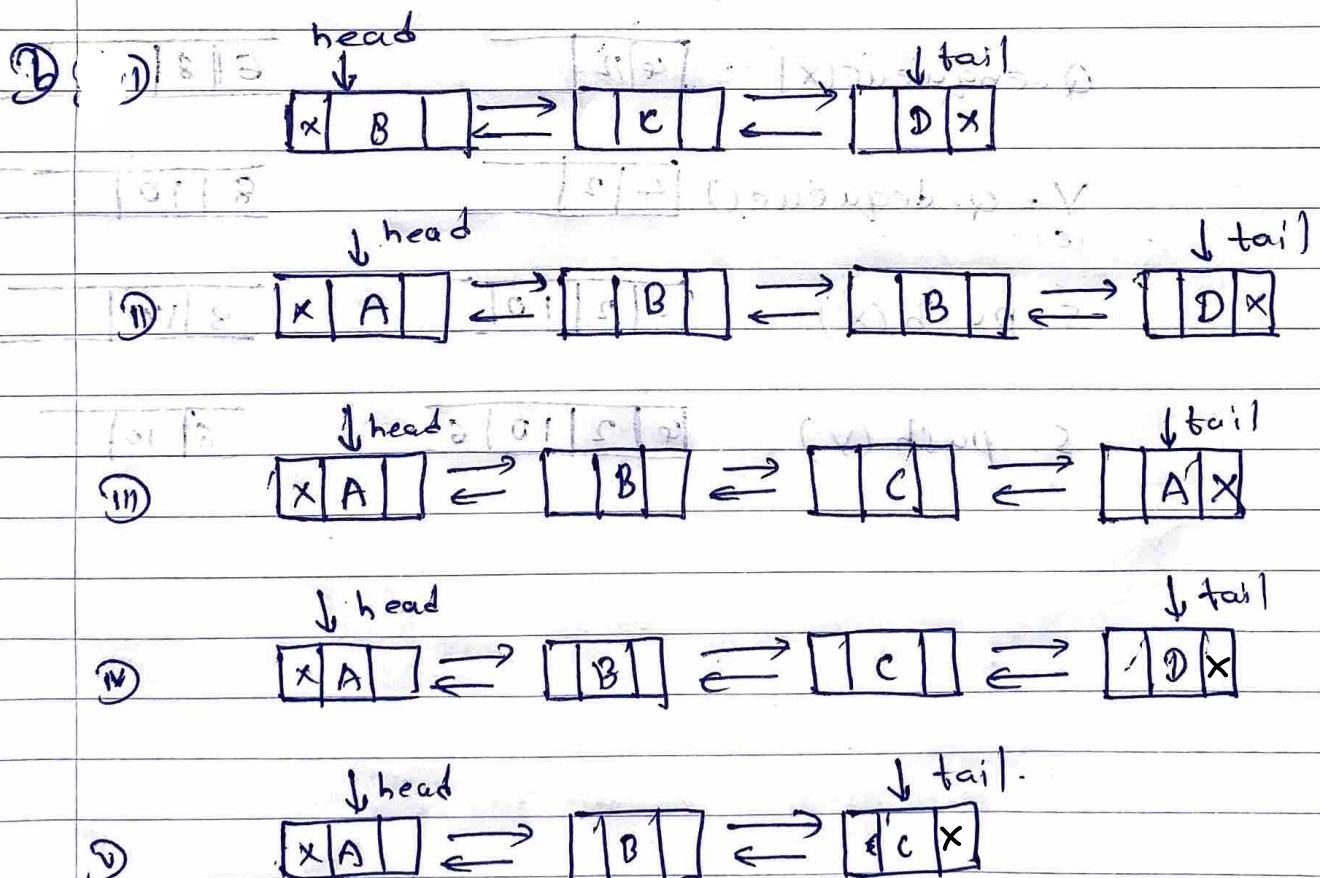
Question 4

④ Static Memory Allocation:

Memory is allocated for variables during the compilation time, and the allocation remains fixed for the lifetime of program. It is suitable for variable with known sizes and lifetimes.

Dynamic Memory Allocation:

Memory is allocated at runtime using functions like 'malloc'. It allows for variable-sized data structures and flexible memory management.



Q) i) Time complexity refer to the estimation of the amount of time an algorithm takes to run, based on the size of the input. It measure how the runtime of an algorithm increases as the input size grows.

ii) Space complexity refers to the estimation of the amount of memory an algorithm requires to run , based on the size of the input. It measure how the memory usage of an algorithm increases as the input size grows.

③ void testFunction(int n) {
 for (n) → n
 for (10) → 10
 for (n) → n
 for (20) → 20}

Total Time

$$\text{complexity} = n \times 10 \times n \times 20 \\ = 200 n^2$$

In Big'O $\Rightarrow O(n^2)$

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