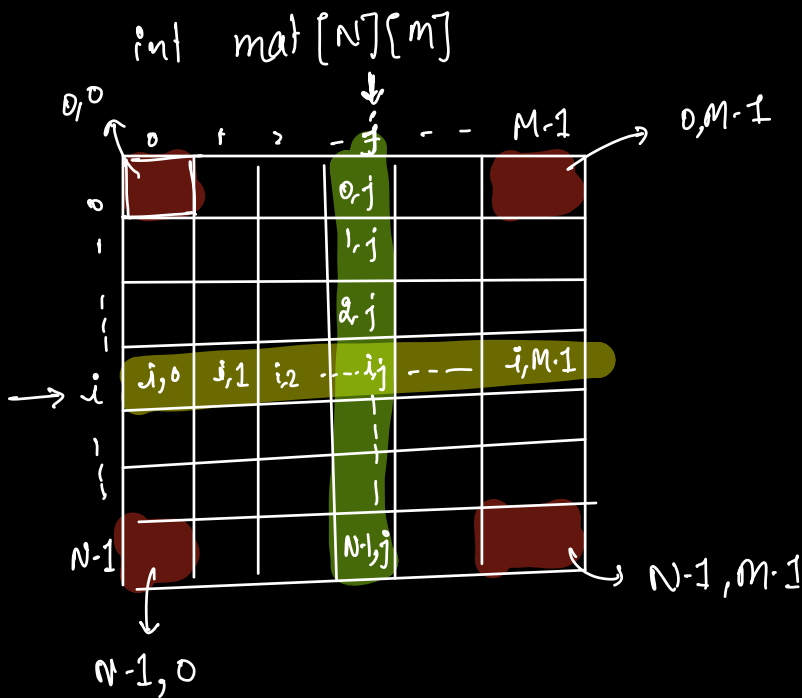
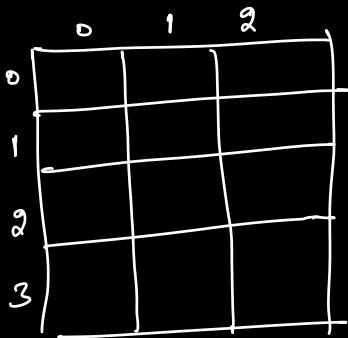


Today's content → 2D Arrays
 → Problems on 2D Arrays.

✓ How to declare 2-D array?

int mat[4][3] no. of columns.
 ↓
 no. of rows



observation 1 : If we move in j^{th} -row
col will change $[0, M-1]$

observation 2 : If we move in j^{th} -col
row will change $[0, N-1]$

Q.1 Given $\text{mat}[N][M]$. Print row-wise sum.

	0	1	2	3
0	4	3	1	7
1	2	5	8	6
2	9	1	3	4

o/p.

15

21

17

pseudo code.

```
void printSum(arr, N, M){  
    for(i=0; i < N; i++){  
        sum = 0  
        for(j=0; j < M; j++){  
            sum += arr[i][j]  
        }  
        // print(sum)  
    }  
}
```

T.C $\rightarrow O(N \times M)$
S.C $\rightarrow O(1)$

Print column-wise sum.

[To-do].

Q.2

	0	1	2	3
0	4	3	1	7
1	2	5	8	6
2	9	1	3	4

o/p \rightarrow 15, 9, 12, 17

Q2) Given $arr[N][N]$. Print Diagonals. — $\left\{ \begin{array}{l} \text{from } l \text{ to } r \\ \text{from } r \text{ to } l \end{array} \right.$

	0	1	2	3
0	(0,0)			
1		(1,1)		
2			(2,2)	
3				(3,3)

$T.C \rightarrow O(N), S.C \rightarrow O(1)$

```

i = 0, j = 0
while (i < N && j < N) {
    print(arr[i][j]);
    i++;
    j++;
}

```

	0	1	2	3
0				(0,3)
1			(1,2)	
2		(2,1)		
3	(3,0)			

$0, 3$
 \downarrow
 $1, 2$
 \downarrow
 $2, 1$
 \downarrow
 $3, 0$
 \downarrow
 $4, -1$

```

i = 0, j = N-1
while (i < N && j >= 0) {
    print(arr[i][j]);
    i++;
    j--;
}

```

$T.C \rightarrow O(N), S.C \rightarrow O(1)$

Q: Given a mat[N][M]. Print all diagonals from R to L.

diagonals starting from 0th row or M-1th column.

	0	1	2	3	4
0		0,1		0,3	0,4
1	1,0		1,2	1,3	1,4
2		2,1	2,2	2,3	
3	3,0	3,1	3,2		

r, c

0, 3

↓

1, 2

↓

2, 1

↓

3, 0

↓

4, -1 → stop.

r, c

1, 4.

↓

2, 3

↓

3, 2

↓

4, 1

stop.

arr[4][5].

	0	1	2	3	4
0	5	7	3	2	6
1	11	9	20	21	13
2	14	-2	3	11	17
3	27	37	47	50	60

o/p.

5
7 11
3 9 14
2 20 -2 27
6 21 3 37
13 11 47
17 50
60

Final obs. for every s.p, print diagonal from l to r.

pseudo-code-

```
void printAllDiagonals ( arr, N, M) {  
    // print all Diagonals starting from 0th row
```

```
    for( j = 0 ; j < M ; j++) {  
        r = 0 , c = j  
        while (r < N && c >= 0) {  
            print (arr[r][c])  
            r++, c--  
        }  
    }
```

	0	1	2	3
0	0,0	0,1	0,2	0,3
1	1,0	1,1	1,2	
2	2,0	2,1	2,2	

```
    // print all Diagonals starting from last col.
```

```
    for( i = 1 ; i < N ; i++) {  
        r = i , c = M-1  
        while (r < N && c >= 0) {  
            print (arr[r][c])  
            r++, c--  
        }  
    }
```

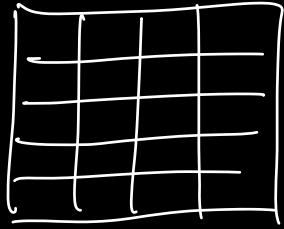
	0	1	2	3
0	0,0	0,1	0,2	0,3
1	1,0	1,1	1,2	1,3
2	2,0	2,1	2,2	2,3

T.C $\rightarrow O(N \times M)$, S.C $\rightarrow O(1)$

\Downarrow
[we are touching all the
elements once.]

arr[n] \rightarrow  $\rightarrow O(\underline{n})$.

mat [n][m]



T.C $\rightarrow O(n * m)$

[printing all elements in a fancy way doesn't]
impact T.C.

Breaks 10:10 \rightarrow 10:20

Q: Given $\text{mat}[N][N]$. Calculate the transpose of the given matrix. $[S.C \rightarrow O(1)]$

Note → take transpose in given matrix only.

	0	1	2	3	4
0	1	2	3	4	5
1	6	7	8	9	10
2	11	12	13	14	15
3	16	17	18	19	20
4	21	22	23	24	25

row 0 → col 0

row 1 → col 1

row 2 → col 2

row 3 → col 3

row 4 → col 4

	0	1	2	3	4
0	1	6	11	16	21
1	2	7	12	17	22
2	3	8	13	18	23
3	4	9	14	19	24
4	5	10	15	20	25

$i=0 \rightarrow j=[1,4]$

$i=1 \rightarrow j=[2,4]$

$i=2 \rightarrow j=[3,4]$

$i \quad j \rightarrow [i+1, N-1]$

idea: swap elements of upper half with lower half

```
void takeTranspose (arr, N) {
    for (i = 0; i < N; i++) {
        for (j = i + 1; j < N; j++) {
            // swap arr[i][j] with arr[j][i]
            temp = arr[i][j]
            arr[i][j] = arr[j][i]
            arr[j][i] = temp
        }
    }
}
```

$T.C \rightarrow O(N^2)$
 $S.C \rightarrow O(1)$

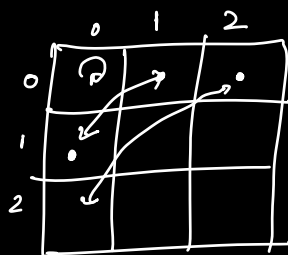
```

void takeTranspose ( arr, N) {
    for ( i = 0 ; i < N ; i++) {
        for ( j = 0 ; j < N ; j++) {
            // swap arr[i][j] with arr[j][i]
            temp = arr[i][j]
            arr[i][j] = arr[j][i]
            arr[j][i] = temp
        }
    }
}

```

$arr[0][0] \leftrightarrow arr[0][0]$
 $arr[0][1] \leftrightarrow arr[1][0]$
 $arr[1][0] \leftrightarrow arr[0][1]$

$arr[0][2] \leftrightarrow arr[2][0]$



\therefore Matrix is going to remain as it is.

// Rectangular matrix \Rightarrow We need to have extra space.

mat[2][5]

	0	1	2	3	4
0	1	2	3	4	5
1	6	7	8	9	10

(2*5)

\Rightarrow

1	6
2	7
3	8
4	9
5	10

5*2

// Now query is simplified.

Q) Given mat[N][N]. Rotate matrix by 90° in clockwise direction.
 { S.C \rightarrow 0(1) }

	0	1	2	3	4
0	1	2	3	4	5
1	6	7	8	9	10
2	11	12	13	14	15
3	16	17	18	19	20
4	21	22	23	24	25

0th row \rightarrow 4th col
 1st row \rightarrow 3rd col
 2nd row \rightarrow 2nd col
 3rd row \rightarrow 1st col
 4th row \rightarrow 0th col

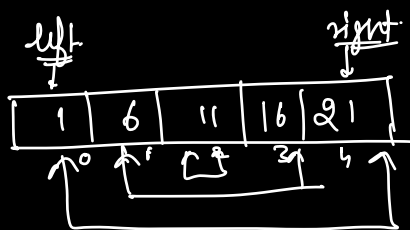
	0	1	2	3	4
0	21	16	11	6	1
1	22	17	12	7	2
2	23	18	13	8	3
3	24	19	14	9	4
4	25	20	15	10	5

\Downarrow transpose.

	0	1	2	3	4
0	1	6	11	16	21
1	2	7	12	17	22
2	3	8	13	18	23
3	4	9	14	19	24
4	5	10	15	20	25

reverse 0th row
 " 1st row
 reverse 2nd row
 reverse 3rd row
 reverse 4th row

	0	1	2	3	4
0	21	16	11	6	1
1	22	17	12	7	2
2	23	18	13	8	3
3	24	19	14	9	4
4	25	20	15	10	5



[Reverse an array]
 row.

// 1. take transpose of the given matrix. $\rightarrow N^2$

// 2. Reverse every row.

```
for( i = 0 ; i < N ; i++) {  
    left = 0 , right = N-1  
    while ( left < right ) {  
        // swap arr[i][left] with arr[i][right].  $\rightarrow N^2$   
        temp = arr[i][left]  
        arr[i][left] = arr[i][right]  
        arr[i][right] = temp.  
        left++ , right--  
    }  
}
```

T.C $\rightarrow O(N^2)$
S.C $\rightarrow O(1)$

for 360° \rightarrow

	0	1	2	3	4
0	1	2	3	4	5
1	6	7	8	9	10
2	11	12	13	14	15
3	16	17	18	19	20
4	21	22	23	24	25

(Remain same)

\rightarrow [Rotate 90° anti-clockwise.] \rightarrow [180°, 270°, ...]

// For Rectangular Matrix. // we have to take extra space.

	0	1	2	3	4
0	1	2	3	4	5
1	6	7	8	9	10

(2x5)

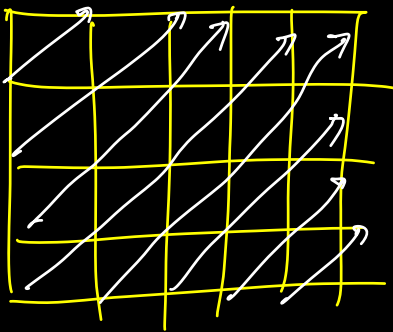
→

	0	1
0	6	1
1	7	2
2	8	3
3	9	4
4	10	5

{ 5x2 }

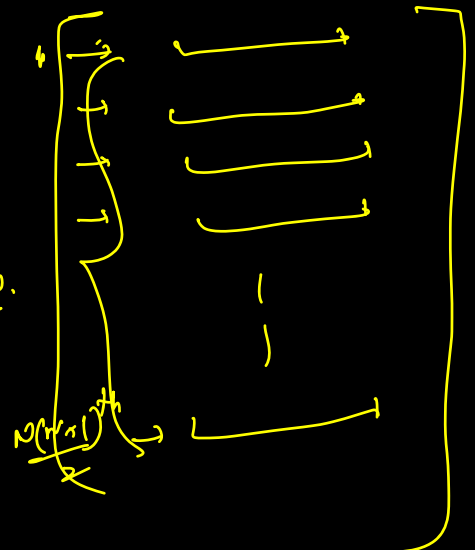
T.C → $O(N \times m)$, S.C → $O(N \times m)$

Extra Todo → Print Diagonals from 1 to n



total no. of subarrays → $\frac{N(N+1)}{2}$

length of every subarray → n.



T.C $\rightarrow O(N^2 \times N)$

\rightarrow Sliding window + Spiral matrix. \rightarrow next session.