

Q1) Find first missing natural number

Ex 1 3, -2, 1, 2, 7 → 4

 -9, 2, 6, 4, -8, 1, 3 → 5

 1, 2, 5, 6, 4, 3 → 7

size of array = n
max possible answer $\Rightarrow n+1$

answer is b/w 1 & $n+1$

- Idea \Rightarrow Store elements in hashset and start checking from 1 till the time no is not found.

```
for (i=0; i<N; i++)
```

```
    hs.insert(arr[i])
```

```
for (i=1; i<=(N+1); i++) {
```

```
    if (i not in hs) {
```

```
        return i
```

```
    }
```

```
}
```

TC: $O(N)$ SC: $O(N)$

Constraint

Expected TC: $O(N)$

SC: $O(1)$

Idea : Keep element in correct position.

$N = 5$

Generalize

val	idx
1	0
2	1
3	2
...	...
N	$n-1$

val	idx
1	0
2	1
...	...
x	$x-1$
...	...
N	$n-1$

any val $> N$
any val < 1
ignore

	0	1	2	3	4	5	6	7
arr[8]	1	2	3	4	9	6	7	8

i

0

1

2

3

4

5

6

7

Final array \Rightarrow

0	1	2	3	4	5	6	7
1	2	3	4	9	6	7	8

Start from 0 idn & find missing no.

Ex 2

0	1	2	3	4	5	6	7	8	9
1	2	3	-14	5	6	7	8	9	10

Ans = 4

0	1	2	3
-3	4	-2	4

Code

```
for (i=0; i < N; i++) {  
    while (ar[i] > 1 && ar[i] ≤ n && ar[i] != (i+1)) {  
        int val = ar[i]  
        if (ar[i] == ar[val-1]) break  
        else swap (ar[i], ar[val-1])  
    }  
}
```

// iterate to get missing no.

```
for (i=0; i < N; i++) {  
    if (ar[i] != (i+1))  
        return i+1  
}
```

return n+1

TC: $O(N)$

SC: $O(1)$

How TC $O(N)$

i	swap-count
0	S_0
1	S_1
2	S_2
\vdots	\vdots
$n-1$	S_{n-1}

$S_0 + S_1 + \dots + S_{n-1}$ tot no of swaps

Obs: Each swap places atleast 1 elem in correct place.

max no of swaps = N

Hence $O(N)$

difference $\Rightarrow a - b$

absolute value

$x \geq 0 \Rightarrow x$

$x < 0 \Rightarrow -x$

Q2 Max absolute diff.

Given array A of size, find max of $|A_i - A_j| + |i - j|$

Eg-

0 1 2
1, 3, -2

$$\Rightarrow 3 - (-2) + |1 - 2| = 6$$

Brute: nested for loops TC: $O(n^2)$

Note: If I ask max of $A_i - A_j$
ans = max val - min val.

Idea $f(i, j) = f(j, i)$

$\Rightarrow f(i, j) = |A_i - A_j| + (i - j)$ such that $i > j$

$$\begin{array}{l} A_i > A_j \\ A_i - A_j + i - j \\ A_i + i - (A_j + j) \\ \cancel{x_i} - \cancel{x_j} \\ x_k = A_k + k \end{array}$$

$$\begin{array}{l} A_i < A_j \\ A_j - A_i + i - j \\ A_j - j - (A_i - i) \\ \cancel{x_j} - \cancel{x_i} \\ y_k = A_k - k \end{array}$$

$$\begin{array}{c}
 \text{0} \quad \text{1} \quad \text{2} \\
 A = 1 \quad 3 \quad -2 \\
 X = 1 \quad 4 \quad 0
 \end{array}$$

$$\max - \min = 4$$

$$\begin{aligned}
 & X_i - X_j \\
 \Rightarrow & \max(X) \\
 & - \min(X)
 \end{aligned}$$

$$\begin{array}{c}
 \text{0} \quad \text{1} \quad \text{2} \\
 A = 1 \quad 3 \quad -2 \\
 Y = 1 \quad 2 \quad -4
 \end{array}$$

$$\max - \min = 6$$

$$\begin{aligned}
 & Y_i - Y_j \\
 & \max(Y) \\
 & - \min(Y)
 \end{aligned}$$

$$\text{ans} = \max(\text{ans}_x, \text{ans}_y)$$

$$TC: O(N)$$

$$SC: O(N)$$

x, y

for $i \geq 0; i < n$ iff

$$x(i) = a(i) + i$$

$$y(i) = a(i) - i$$

$$\text{ans}_x = \max(x) - \min(x)$$

$$\text{ans}_y = \max(y) - \min(y)$$

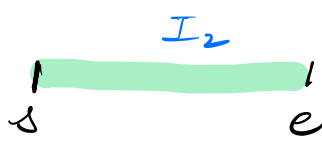
return $\max(\text{ans}_x, \text{ans}_y)$

Q3 Merge intervals.

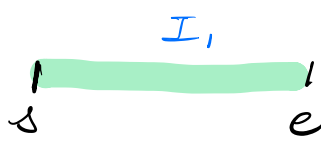
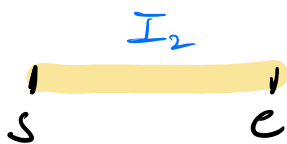
Interval is $[a, b]$

overlap \Rightarrow If intersecting at 1 or more elem,
then they overlap.

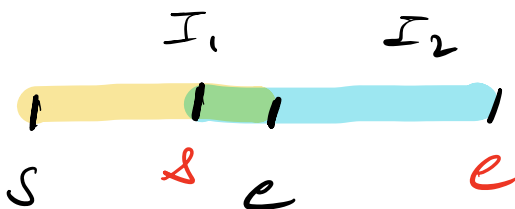
<u>Ex</u>	2, 6	3, 7	2, 7
	2, 8	4, 6	2, 8
	3, 7	4, 10	3, 10
	3, 6	6, 10	3, 10
	2, 5	8, 10	no overlap



$\left. \begin{array}{l} I_1.e < \\ I_2.s \end{array} \right\}$



$\left. \begin{array}{l} I_2.e < \\ I_1.s \end{array} \right\}$



merged interval
 $s \swarrow \searrow e$

$\min(I_1.s, I_2.s)$ $\max(I_1.e, I_2.e)$

Given N ^{non overlap} intervals, sorted based on start
 new interval I comes, merge all

$[1, 3]$	$I = $	$12, 22$	$1, 3$
$[4, 7]$		$10, 22$	$4, 7$
$[10, 14]$		$10, 24$	$10, 24$
$[16, 19]$			$27, 30$
$[21, 24]$			$32, 35$
$[27, 30]$			return.
$[32, 35]$			

$N=5$	$[1, 5]$	$I = $	$[12, 22]$	$1, 5$
	$[8, 10]$		$14, 22$	$8, 10$
	$[11, 14]$		$11, 24$	$11, 24$
	$[15, 20]$			
	$[20, 24]$			

Code

Say $ar[N]$ Intervals, Interval I comes

Intervals $[]$ merge (Intervals $ar[]$, Interval I) {

for ($i=0; i < N; i++$) { arraylist<Interval> ans

 // i^{th} Interval = $ar[i]$

 if ($ar[i].e < I.s$)

 ans.insert($ar[i]$)

 else if ($I.e < ar[i].s$) {

 ans.insert(I)

 for ($j=i; j < n; j++$)

 ans.insert($ar[j]$)

 return ans,

}

else {

//

 // update Interval I after ?

$I.s = \min(I.s, ar[i].s)$

$I.e = \max(I.e, ar[i].e)$

}

Y

ans.insert(I)

return ans

Y

TC : $O(n)$

SC: $O(1)$

{done}

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