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## MAJORITY ELEMENT

Given an array of size  $N$ , Return if there exists a majority element.

$\rightarrow \text{freq} > \frac{N}{2}$

SC:  $O(1)$

Ex:-

A: 1, 6, 1, 1, 2, 1

$N = 6$

ME = 1

$\text{freq} > \frac{6}{2}$

Minfreq = 4

Quiz

A: 3, 4, 3, 6, 1, 3, 2, 5, 3, 3, 3

ME = 3

$N = 11$

Minfreq = 6

Quiz

A: 4, 6, 5, 3, 4, 5, 6, 4, 4, 4

$\text{freq}(4) = 5 > 5$  X

$N = 10$

minfreq = 6

NO M.E

A: 4, 6, 5, 3, 4, <sup>i</sup>5, 6, 4, 4, 4  
↑  
i

```
int majority (int A[], int N) {  
    for (i = 0; i < N; i++) {  
        freq = 0  
        for (j = 0; j < N; j++) {  
            if (A[i] == A[j]) {  
                freq++;  
            }  
            if (freq > N/2)  
                return A[i];  
        }  
        return -1;  
    }  
}
```

TC:  $O(N^2)$

SC:  $O(1)$

Quiz At max how many M.E can be there in an Array?



$N=10$

$$\min \text{freq} = 6$$



$\leftarrow \frac{N}{2} \rightarrow \leftarrow \frac{N}{2} \rightarrow$

$\Rightarrow$  There can be a single M.E only.

Green :  $\begin{matrix} \cancel{0} & \cancel{0} & \cancel{0} & \cancel{0} & \cancel{0} & \cancel{0} \end{matrix}$

Purple :  $\cancel{0}$   
 Red :  $\cancel{0}$   
 Blue :  $\cancel{0}$

$\begin{matrix} \cancel{0} & \cancel{0} & \cancel{0} & \cancel{0} \end{matrix}$

$N=10$	M.E
$\downarrow$	6
8	$\downarrow$ 5
$\downarrow$	$\downarrow$ 4
6	$\downarrow$ 3
$\downarrow$	$\downarrow$ 2
4	
$\downarrow$	
2	

Obs:- If 1 majority & 1 non majority elements are fighting then M.E won't change.

Obs: If 2 non ME's are fighting then ME won't change.

A: <sup>0</sup>3, <sup>1</sup>4, <sup>2</sup>3, <sup>3</sup>6, <sup>4</sup>1, <sup>5</sup>3, <sup>6</sup>2, <sup>7</sup>5, <sup>8</sup>3, <sup>9</sup>3, <sup>10</sup>3

$$N = 11$$

$$\text{freq} > \frac{11}{2} = 5$$

$$ME = 3$$

N	ME
11	6
↓	↓
9	5
↓	↓
7	4
↓	↓
5	4

A: <sup>0</sup>3, <sup>1</sup>4, <sup>2</sup>3, <sup>3</sup>6, <sup>4</sup>1, <sup>5</sup>3, <sup>6</sup>2, <sup>7</sup>5, <sup>8</sup>3, <sup>9</sup>3, <sup>10</sup>3 ↑

$$ME = \cancel{3} \cancel{3} \cancel{4} \cancel{2} 3$$

$$\text{freq} = \cancel{4} \cancel{0} \cancel{4} \cancel{0} \cancel{4} \cancel{0} \cancel{4} \cancel{0} \cancel{4} \cancel{2} 3$$

Ex :-

A: 3 3 4 2 4 4 2 4 4  $N=9$

$$ME = 3 4$$

$$freq = \cancel{x} \cancel{x} \cancel{x} \cancel{0} \cancel{x} \cancel{x} \cancel{x} \cancel{x} 3$$

$$freq(4) = 5 > \frac{9}{2}$$

Ex

1 2 3 4 5 6 7

$$ME = \cancel{x} \cancel{x} \cancel{x} 7$$

$$freq = \cancel{x} \cancel{0} \cancel{x} \cancel{0} \cancel{x} \cancel{0} 1$$

Ex

1 2 3 4

$$ME = \cancel{x} 3$$

$$f = \cancel{x} \cancel{0} \cancel{x} 0$$

Ex

1 1 2 3

$$ME = 1$$

$$f = \cancel{x} \cancel{x} \cancel{x} 0$$

### MOORE'S VOTING ALGORITHM

$$TC: O(N)$$

$$SC: O(1)$$

Code

```
int me = a[0]
freq = 1
for( i = 1; i < N; i++) {
    if(a[i] == me){
        freq++
    }
    else {
        if(freq == 0){
            me = a[i];
            freq = 1;
        }
        else {
            freq--;
        }
    }
}

count = 0
for( i = 0; i < N; i++) {
    if(a[i] == me)
        count++
}

if(count > N/2)
    return me;
return -1;
```

Q. Given an array & Q queries  
Directly

$s, e, o \Rightarrow$  Sum of all **odd** indexed elements from  $[s, e]$

$s, e, e \Rightarrow$  Sum of all **even** indexed elements from  $[s, e]$

A: 2, 3, 1, -1, 0, 8, 5, 4

Q: 2

s	e	o/e	
3	6	o	$\Rightarrow -1 + 8 = 7$
1	5	e	$\Rightarrow 1 + 0 = 1$

Ex A: 2, 3, 1, 6, 4, 5

$PS:$  2, 5, 6, 12, 16, 21

$PS_E:$  2, 2, 3, 3, 7, 7

$PS_O:$  0, 3, 3, 9, 9, 14

$$PS_E[0] = A[0]$$

$$PS_E[i] = \begin{cases} PS_E[i-1] + A[i] & i \% 2 == 0 \\ PS_E[i-1] & \text{else} \end{cases}$$

$$PS_0[0] = 0$$

$$PS_0[i] = \begin{cases} PS_0[i-1] + A[i] & i \% 2 \neq 0 \\ PS_0[i-1] & \text{else} \end{cases}$$

Quiz

A:     0     1     2     3     4  
         2     4     3     1     5

PS<sub>0</sub>:   0     4     4     5     5

Quiz

A:     0     1     2     3     4     5     6  
         4     1     0     -2     3     2     5

PS<sub>E</sub>:   4     4     4     4     7     7     12

Sum of **ODD** indexed elements from s to e  $\Rightarrow$   
 $PS_0[e] - PS_0[s-1]$

Sum of **EVEN** indexed elements from s to e  $\Rightarrow$   
 $PS_E[e] - PS_E[s-1]$

TC:  $O(N) + O(N) + O(1) : O(N+1)$

SC:  $O(N)$



Q. Given an Array, count the no. of Special Indices in the Array.

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Special Index: An index after removing which

$$\text{Sum of all odd indexed elements} = \text{Sum of all even indexed elements}$$

in the resultant array.

A :      <sup>0</sup>4    <sup>1</sup>3    <sup>2</sup>2    <sup>3</sup>7    <sup>4</sup>6    <sup>5</sup>-2

i	[ ]	SE	So	
0	<sup>0</sup> 3 <sup>1</sup> 2 <sup>2</sup> 7 <sup>3</sup> 6 <sup>4</sup> -2	8	8	✓
1	<sup>0</sup> 4 <sup>1</sup> 2 <sup>2</sup> 7 <sup>3</sup> 6 <sup>4</sup> -2	9	8	✗
2	<sup>0</sup> 4 <sup>1</sup> 3 <sup>2</sup> 7 <sup>3</sup> 6 <sup>4</sup> -2	9	9	✓
3	<sup>0</sup> 4 <sup>1</sup> 3 <sup>2</sup> 2 <sup>3</sup> 6 <sup>4</sup> -2	4	9	✗
⋮	⋮	⋮	⋮	

Quiz

	0	1	2	3	4	5
A:	4	1	5	3	7	10
	↓	↓	↖	↖	↖	↖
	0	1	2	3	4	
	4	1	3	7	10	

$$\begin{aligned} S_0 &= 8 = S_0[0-1] + S_E[3-5] \\ &= 1 + 7 = 8 \end{aligned}$$

Quiz

	0	1	2	3	4	5	6	7	8	9
A:	2	3	1	4	0	-1	2	-2	10	8
	↓	↓	↓	↖	↖	↖	↖	↖	↖	↖
	0	1	2	3	4	5	6	7	8	
	2	3	1	0	-1	2	-2	10	8	

$$S_0 = S_0[0-2] + S_E[4-9]$$

Quiz

	0	1	2	3	4	5	6	7	8	9
A:	2	3	1	4	0	-1	2	-2	10	8

$$\begin{aligned} S_E &= S_E[0-2] + S_0[4-9] \\ &= 3 + 5 = 8 \end{aligned}$$

# Sum of ODD indexed elements after removing  
index = i

$$\begin{aligned} S_o &= S_o[0, i-1] + S_e[i+1, N-1] \\ &= PS_o[i-1] + (PS_e[N-1] - PS_e[i]) \end{aligned}$$

# Sum of Even indexed elements after removing  
index = i

$$\begin{aligned} S_e &= S_e[0, i-1] + S_o[i+1, N-1] \\ &= PS_e[i-1] + PS_o[N-1] - PS_o[i] \end{aligned}$$

$$\boxed{\text{sum}[s, e] = PS[e] - PS[s-1]}$$

# i = 0

$$\begin{aligned} S_o &= S_e[1, N-1] \\ \Rightarrow S_o &= PS_e[N-1] - PS_e[0] \end{aligned}$$

$$\begin{aligned} \Rightarrow S_e &= S_o[1, N-1] \\ S_e &= PS_o[N-1] - PS_o[0] \end{aligned}$$

# i = N-1

$$S_o = S_o[0, N-2] = PS_o[N-2]$$

$$S_e = S_e[0, N-2] = PS_e[N-2]$$

1. Build  $PS_0$

2. Build  $PS_E$

3. for ( $i=0$ ;  $i < N$ ;  $i++$ ) {

    if ( $S_E == S_0$ )

        count++

}

return count;

TC:  $O(N)$

SC:  $O(N)$

———— \* ————