

## 2 classes on Recursion

Recursion :

Why?

Tod {  
→ How to write recursive code  
→ Working  
→ TC/SC  
Nent Session

{  
→ Merge sort  
→ Binary Trees  
→ Dynamic Programming  
→ Backtracking

Recursion: { function calling itself }

Solve a problem using smaller version of the same problem.

↳ subproblem

$$\begin{aligned} \text{Sum}(N) &= 1 + 2 + 3 + 4 + \dots + N \\ &= 1 + 2 + 3 + 4 + \dots + N-1 + N \\ &= \text{sum}(N-1) + N \end{aligned}$$

### Recursion Code

- 1) Assumption: Decide what your function does, and assume it does exactly that
- 2) Main Logic: Solving Assumption with subproblem
- 3) Base Condition: When should code stop.  
Smallest value for which we know answer.

int sum(N) {      Calc sum of first N natural nos.

Assumption: return sum of first N natural nos.

Base Case: if (n == 1)  
                  return 1

Main logic: return (sum(n-1) + n)

}

int fact(N) {      Calculate factorial of N  
                             $1 \times 2 \times 3 \times 4 \times \dots \times N$

Assumption: return factorial value of N

Base case: if (n == 1)  
                  return 1

Main Logic: return (fact(n-1) \* n )

}

$$f(N) = 1 \times 2 \times 3 \times \dots \times N-1 \times N$$
$$f(n-1)$$

Fibonacci series  $fib(n) = fib(n-1) + fib(n-2)$

fib:      1      2      3      4      5      6      7      8      9      10      11      12  
            1      1      2      3      5      8      13      21      34      55      89      144

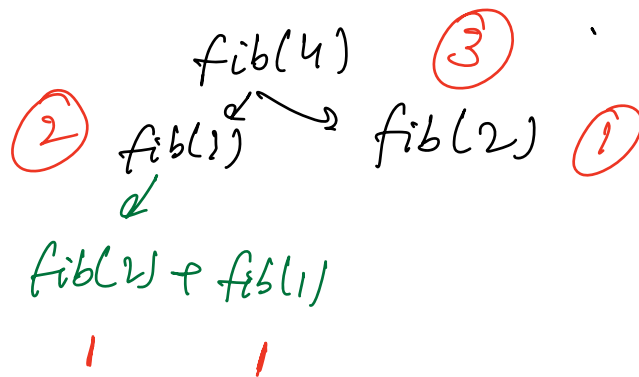
int fib(N) &  $N^{th}$  fibonacci number

Assumption: Return  $n^{th}$  fib number

Base Case: if ( $n == 1$  ||  $n == 2$ )  
                  return 1

Main logic return ( $fib(n-1) + fib(n-2)$ )

y



Example :

```
int add(x, y) {  
    return x+y  
}
```

```
int mul(x, y) {  
    return x*y  
}
```

```
int sub(x, y) {  
    return x-y  
}
```

main() {

int x=10, y=20, z=30

int a = add(x, y)     a = 30

int m = mul(a, z)     m = 900

int s = sub(m, 75)     s = 825

print(s)     825

}

main() {

int x=10, y=20, z=30

print(sub(mul(add(x, y), z), 75))

}

sub(mul(add(x, y), z), 75)

825

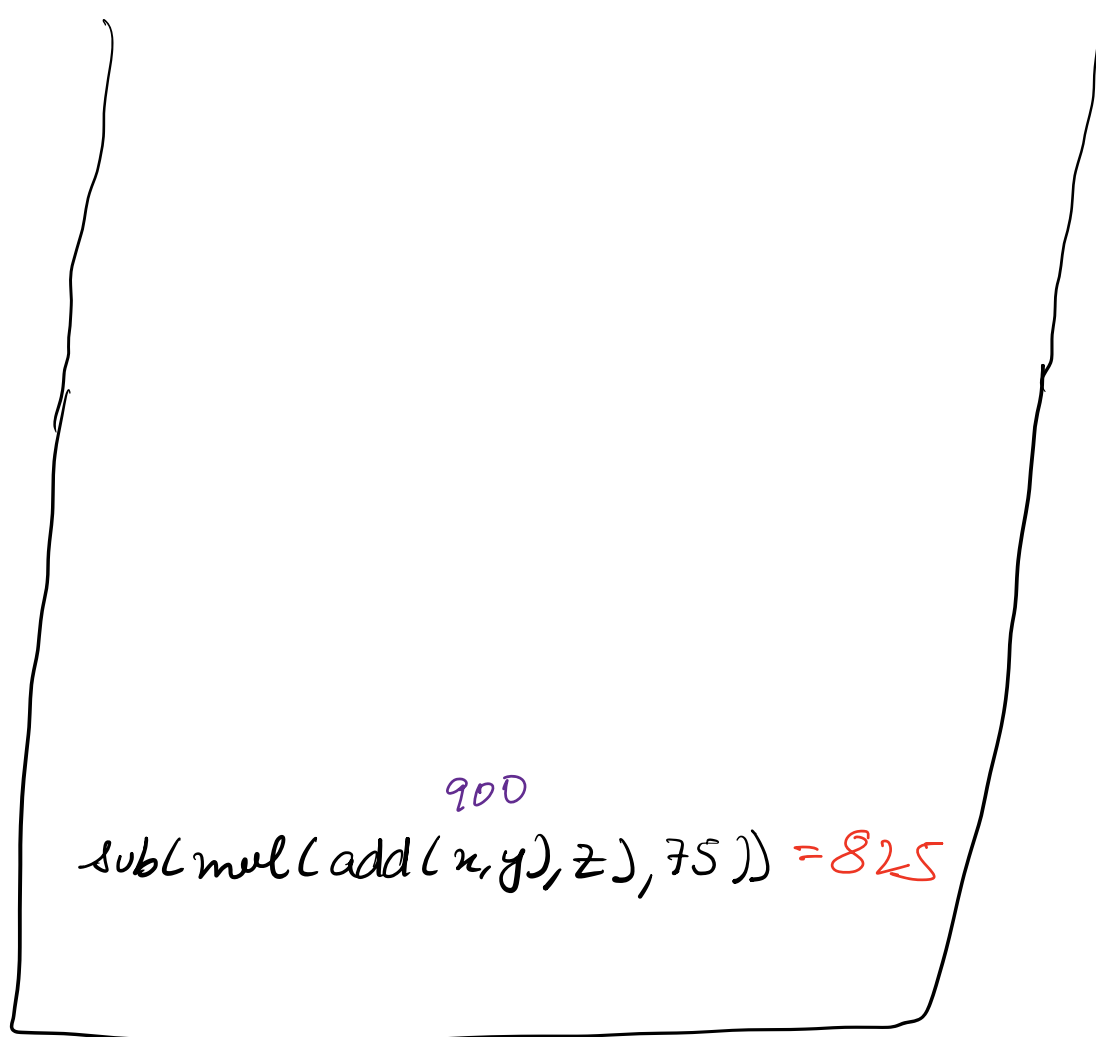
mul(add(x, y), z)

30

add(x, y) = 30

10 20

# Internal working (Call Stack)



# Works like an Idli cooker

```

int sum(N) {
    if (N==1) return 1
    return N + sum(N-1)
}

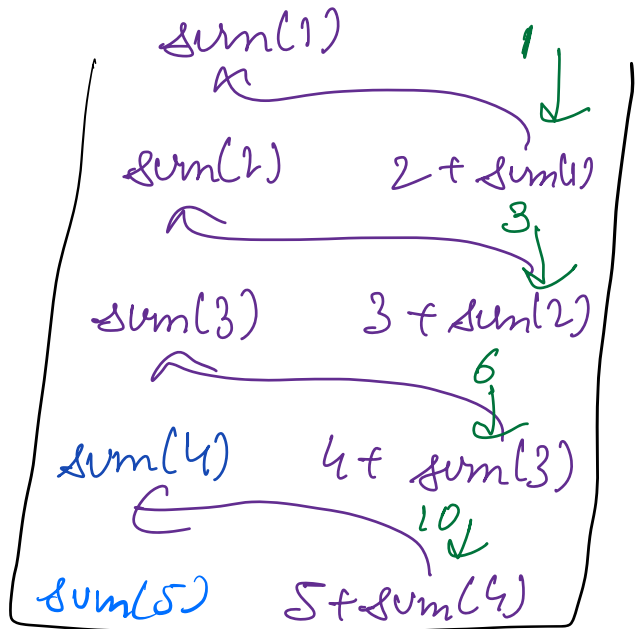
```

```

main() {
    print(sum(5))
}

```

$$1 + 2 + 3 + 4 + 5 = 15$$



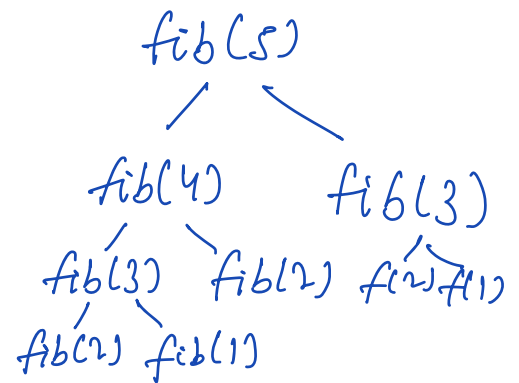
$\text{sum}(5)$   
 $5 + \text{sum}(4)$

$\text{sum}(4)$   
 $4 + \text{sum}(3)$

$\text{sum}(3)$   
 $3 + \text{sum}(2)$

$\text{sum}(2)$   
 $2 + \text{sum}(1)$

Recursion  
Tree



→ print 1 2 3 4 ..... N

void printInc(N) {

Assumption: prints 1 to N

Base Case: if (n == 1)  
print 1 return;

Main Logic: printInc(n-1)  
print(N)

}

p(N) = 1 2 3 4 ..... N-1 N

main() {

printInc(5)  
}

printInc(5)  
 ● printInc(4)  
 ● print(5)

1 2 3 4 5

printInc(4)  
 ● printInc(3)  
 ● print(4)

printInc(3)  
 ● printInc(2)  
 ● print(3)

printInc(2)  
 ● printInc(1)  
 ● print(2)

printInc(1)  
 ● print(1)

$s=0$   
 $s=1$

$e=1$   
 $e=0$

0  
 a  
 e

1  
 a  
 s



Q Given string check whether palindrome

Eg-      abacaba  $\Rightarrow$  true  
         abcd       $\Rightarrow$  false.

Solve using recursion.

bool ispal(str, <sup>start</sup>s, <sup>end</sup>e) {

Assumption: Return if the substring  
                 str[s:e] is palindrome

Base case: if (s > e) return true

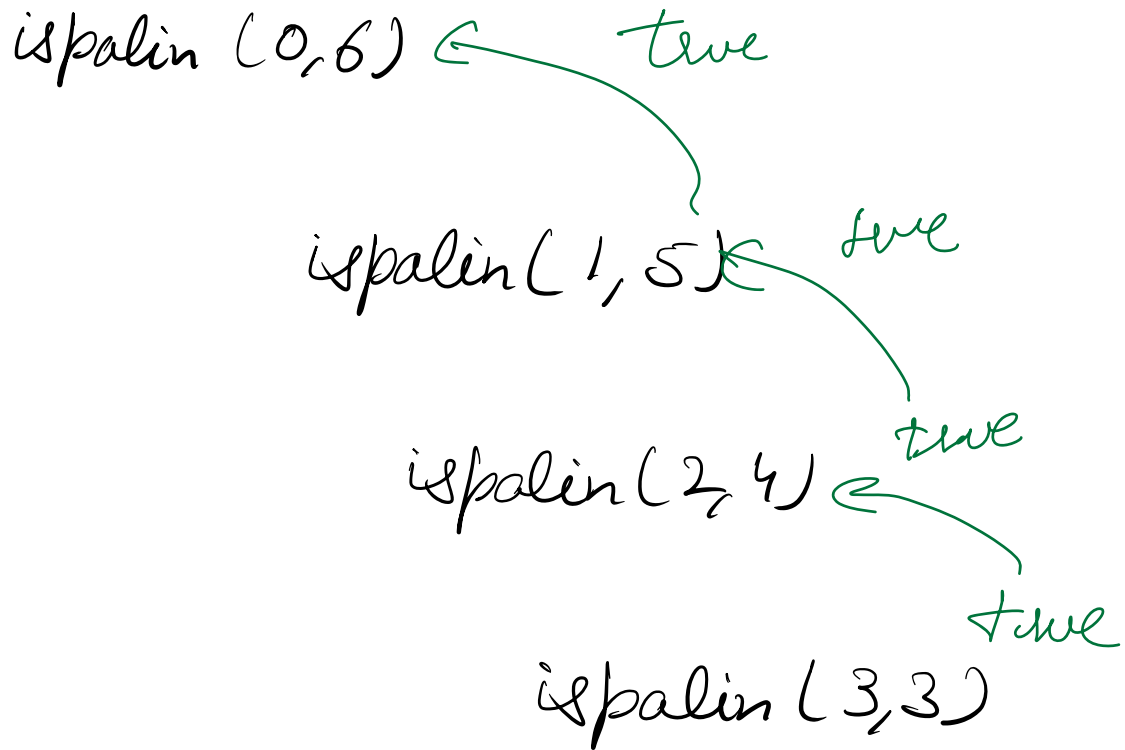
Main logic:

if (str[s] != str[e])  
    return false

else return ispalin(str, s+1, e-1)

}

0 1 2 3 4 5 6  
a b a c a b a



0 1 2 3  
a b c a

isPal(0,3)  $\leftarrow$  false  
isPal(1,2) false

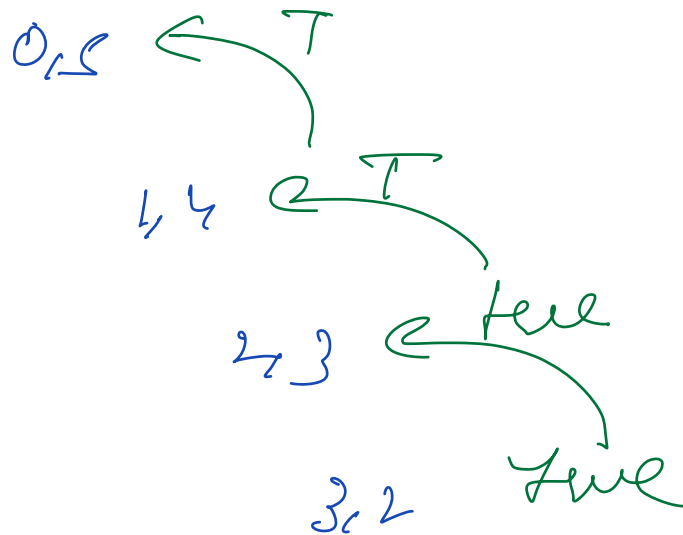
0 1 2 3  
a c c a

isPal(0,3)

isPal(1,2)

isPal(2,1)

0 1 2 3 4 5  
a d p p d a



done?

M — O — T

DP



