

CMPS 130 – Spring Quarter 2017 – Homework 1

Christopher Hsiao – chhsiao@ucsc.edu – 1398305

1 Exercises from pages 25, 26, and 27 of the book: 0.1 through 0.9

0.1

Examine the following formal descriptions of sets so that you understand which members they contain. Write a short informal English description of each set.

- The infinite set of all positive odd integers, or the set of all odd natural numbers.
- The infinite set of all even integers.
- The infinite set of all even natural numbers.
- The infinite set of all even natural numbers, and all natural numbers which are multiples of 3.
- The infinite set containing all palindromic bit strings.
- The finite set containing any integer n and $n + 1$. =====

0.2

Write formal descriptions of the following sets.

- $\{1, 10, 100\}$
- $\{n | n > 5 \text{ for some } n \in \mathbb{Z}\}$
- $\{1, 2, 3, 4\}$
- $\{aba\}$
- $\{''''\}$
- $\{\}$

0.3

Let A be the set $\{x, y, z\}$ and B be the set $\{x, y\}$.

- No.
- Yes.
- $\{x, y, z\}$
- $\{x, y\}$
- $\{\{x, x\}, \{x, y\}, \{y, x\}, \{y, y\}, \{z, x\}, \{z, y\}\}$
- $\{\{\}, \{x\}, \{y\}, \{x, y\}\}$

0.4

If A has a elements, and B has b elements, how many elements are in $A \times B$? Explain your answer.

Claim. $|A \times B| = a \cdot b$

Proof. This is because we assign every element of the set A to every element in the set B . That means, each element in the set A , such as A_i , has b pairings made. Thus, if there are b pairings made with every element in A , then there are a pairings, each of size b , which yields $a \cdot b$ number of pairings made. \square

0.5

If C is a set with c elements, how many elements are in the power set of C ? Explain your answer.

First, we will define the power set of C as P_C .

Claim. $|P_C| = 2^c$

Proof. \square

0.6

Let X be the set $\{1, 2, 3, 4, 5\}$ and Y be the set $\{6, 7, 8, 9, 10\}$. The unary function $f : X \rightarrow Y$ and the binary function $g : X \times Y \rightarrow Y$ are described.

- a. 7
- b. $D : [6, 7], R : [1, 5]$
- c. 6
- d. $D : [6, 10], R : [1, 5]$
- e. 8

0.7

For each part, give a relation that satisfies the condition. Let $A = \{x, y, z\}$.

- a. Reflexive, Symmetric, but not Transitive.

Let $R = \{(x, x), (y, y), (z, z), (x, y), (y, x), (y, z), (z, y)\}$

R is reflexive, since $(x, x), (y, y), (z, z) \in R$.

R is symmetric, since $(x, y), (y, x), (y, z), (z, y) \in R$.

R is not transitive, since $(x, y), (y, z) \in R$, but $(x, z) \notin R$.

- b. Reflexive, Transitive, but not Symmetric.

Let $R = \{(x, x), (y, y), (z, z), (x, y), (y, z), (x, z)\}$

R is reflexive, since $(x, x), (y, y), (z, z) \in R$.

R is transitive, since $(x, y), (y, z), (x, z) \in R$, where (x, y) and $(y, z) \implies (x, z)$

R is not symmetric, since $(x, y), (y, z), (x, z) \in R$, but $(y, x), (z, y), (z, x) \notin R$.

c. Symmetric, Transitive, but not Reflexive.

Let $R = \{(x, y), (y, x), (y, z), (z, y), (x, z), (z, x)\}$

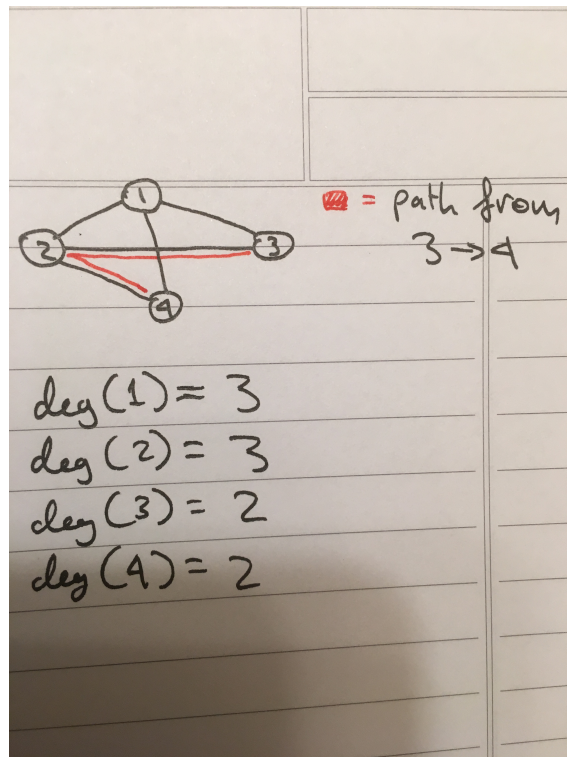
R is symmetric, since $(x, y), (y, x), (y, z), (z, y), (x, z), (z, x) \in R$.

R is transitive, since $(x, y), (y, z), (z, x) \in R$, where (x, y) and $(y, z) \implies (z, x)$.

R is not reflexive, since $(x, x), (y, y), (z, z) \notin R$.

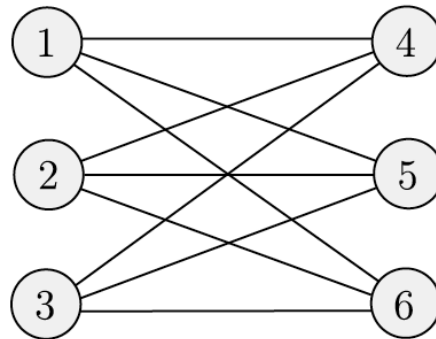
0.8

Consider the undirected graph $G = (V, E)$ where V , the set of nodes, is $\{1, 2, 3, 4\}$ and E , the set of edges, is $\{\{1, 2\}, \{2, 3\}, \{1, 3\}, \{2, 4\}, \{1, 4\}\}$. Draw the graph G . What are the degrees of each node? Indicate a path from node 3 to node 4 on your drawing of G .



0.9

Write a formal description of the following graph.



$\{\{1, 2, 3, 4, 5, 6\}, \{(1, 4), (1, 5), (1, 6), (2, 4), (2, 5), (2, 6), (3, 4), (3, 5), (3, 6)\}\}$

2 Problems page 27 in 3rd ed 0.10, 0.12 and 0.13.

0.10

1. $a = b$
2. $a^2 = ab$
3. $a^2 - a^b = ab - b^2$
4. $\frac{(a+b)(a-b)}{(a-b)} = \frac{b(a-b)}{(a-b)}$
5. $a + b = b$

Claim. *Step 4 is invalid in the case where $a = b = 1$.*

Proof. If $a = b = 1$, that means the term $(a - b) = 0$. This means that (4) is dividing by 0, an impossible operation. Thus, this is not a proof that $2 = 1$. \square

0.12

Claim. *This proof falls apart because it assumes the truth of $k = 2$.*

Proof. This proof relies heavily on the "middle" set H_1 , which contains all k -horses except the first and the last. However, at just $k = 2$, where the only horses are the first and the last, this "middle" argument falls apart completely. The proof must first establish this base case of $k = 2$ first, before proceeding to larger k values. \square

0.13

Show that every graph with two or more nodes contains two nodes that have equal degrees.
We will prove this via Proof by Contradiction.

Proof. Assume some graph $G = (V, E)$, such that all degrees of nodes in V are distinct. Thus, the set of all possible degrees is $D = \{0, 1, 2, \dots, n-1\}$. This means that there is some node V_i with degree $D_i = 0$, and some node V_j with degree $D_j = n-1$. However, $D_i = 0$ implies that there are no edges leaving this node, but D_j implies that node V_j is connected to every node in G , other than itself. This is a contradiction of our initial assumption. \therefore there are at least two nodes $\{V_i, V_j\}$ s.t. $D_i = D_j$. \square

3 Prove from the definitions of set union, intersection, complement and equality that:

$$\overline{(A \cap B)} = (\bar{A} \cup \bar{B}) \quad (1)$$

Claim. $\overline{(A \cap B)} \implies (\bar{A} \cup \bar{B})$

Proof. Consider some $x \in \overline{(A \cap B)}$.

- (1) $x \notin (A \cap B)$
- (2) $x \in \bar{A} \wedge x \in \bar{B}$
- (3) $x \notin A \wedge x \notin B$
- (4) $x \in (\bar{A} \cup \bar{B})$

\square

Claim. $(\bar{A} \cup \bar{B}) \implies \overline{(A \cap B)}$

Proof. Consider some $x \in (\bar{A} \cup \bar{B})$.

- (1) $x \in \bar{A} \vee x \in \bar{B}$
- (2) $x \notin A \vee x \notin B$
- (3) $x \in \overline{(A \cap B)}$

\square

Thus, we can see that $\overline{(A \cap B)} = (\bar{A} \cup \bar{B})$.

4 Show the set of odd numbers is countable.

The set of odd integers is a subset of the integers (which are countable), it is as most countable.

5 Prove by induction on n .

Base Case:

Show that $n = 1$ is true.

1. $i^2 = \frac{n(n+1)(2n+1)}{6}$
2. $1^2 = \frac{1(1+1)(2(1)+1)}{6}$
3. $1 = \frac{(1)(2)(3)}{6}$
4. $1 = 1$

Inductive Hypothesis:

$$\text{Assume that } \sum_{i=1}^k i^2 = \frac{k(k+1)(2k+1)}{6} \text{ for some } k \in \mathbb{N} \quad (2)$$

Inductive Step:

Claim 1. *If $n = k$ is true, then $n = k + 1$ is also true.*

Proof. For $n = k + 1$, the summation is:

$$\sum_{i=1}^{k+1} i^2 \quad (3)$$

Using our inductive hypothesis, we can rewrite this as:

$$\sum_{i=1}^k i^2 + (k+1)^2 = \frac{k(k+1)(2k+1)}{6} + (k+1)^2 \quad (4)$$

- $\frac{1}{6}(k+1)k(2k+1) + (k+1)(k+1)$
- $\frac{1}{6}(k+1)[k(2k+1) + 6k + 6]$
- $\frac{1}{6}(k+1)[2k^2 + 7k + 6]$
- $\frac{1}{6}(k+1)[2k(k+2) + 3(k+2)]$
- $\frac{1}{6}(k+1)(k+2)[2k+3]$
- $\frac{(k+1)((k+1)+1)(2(k+1)+1)}{6}$

Which is proof for $n = k + 1$ being true assuming $n = k$. Thus, by proof of induction the statement is true. □