

# A Unified View of Frequency Estimation and their Attacks on Local Differential Privacy

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# Outline

## Introduction:

- Differential Privacy
- Local Differential Privacy
- Pure LDP Framework
- Attack Problem

## Frequency Estimation Techniques:

- RAPPOR
- K Randomized Response (kRR)
- Optimized Unary Encoding (OUE)
- Optimized Local Hashing (OLH)

## Attacks:

- General Attack Formulation
- Attacking kRR
- Attacking OUE
- Attacking OLH

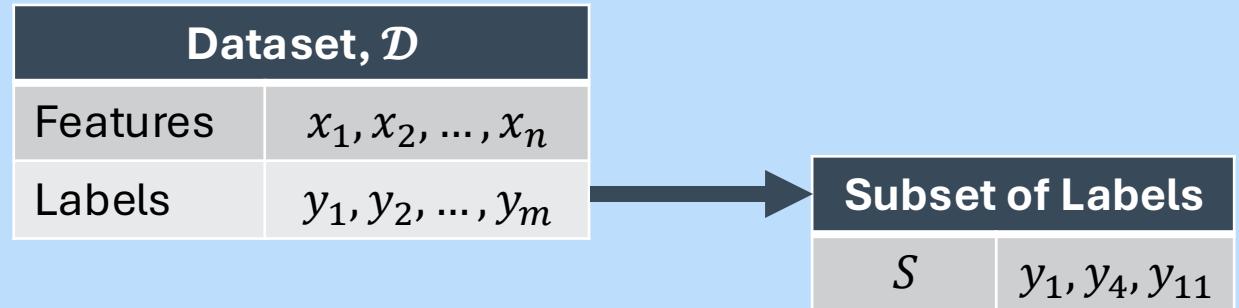
## Evaluation:

- Comparison between Estimators
- Gain from Attacks
- Impacts of Parameters on Attacks

## Conclusion:

# Introductory Concepts

# Differential Privacy



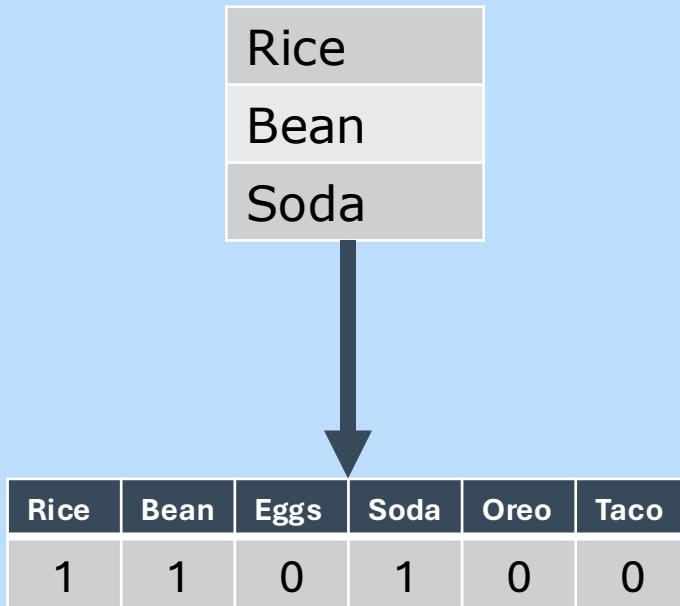
$$\Pr[M(\begin{matrix} x \\ sample_1 \\ sample_2 \\ sample_3 \\ sample_4 \\ sample_5 \end{matrix}) \in S] \leq e^\epsilon \Pr[M(\begin{matrix} y \\ sample_1 \\ sample_2 \\ sample_3 \\ sample_4 \\ sample_{11} \end{matrix}) \in S] + \delta$$

# Local Differential Privacy

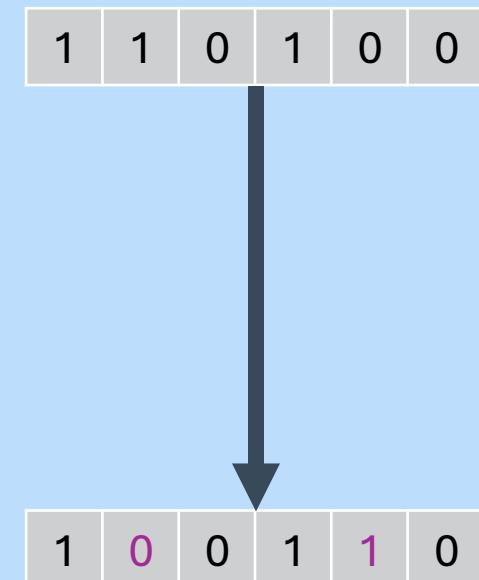
$$\Pr \left[ M\left( \begin{array}{c} x \\ sample_1 \end{array} \right) \in y \right] \leq e^\epsilon \Pr \left[ M\left( \begin{array}{c} y \\ sample_{11} \end{array} \right) \in y \right]$$

# Protocols for Local Differential Privacy

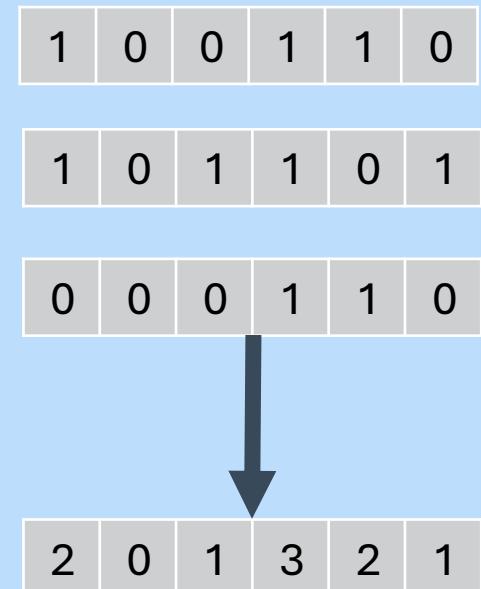
Encode



Perturb



Aggregate



# Pure Differential Privacy

- A unified framework --- Can be used to define all frequency protocols.
- Defined based on two fixed probabilities  $p^*$  and  $q^*$

$$\Pr[PE(v_1) \in \{y \mid v_1 \in \text{Support}(y)\}] = p^*$$
$$\Pr[PE(v_2) \in \{y \mid v_1 \in \text{Support}(y)\}] = q^*$$

- Based on this framework, we will describe 4 protocols – RAPPOR, kRR, OUE, OLH.

# Attacking Estimators – Problem Formulation

- Assume:  $n$  real users
- Inject:  $m$  fake users
- To increase frequency of  $r$  items:  $T = \{t_1, t_2, \dots, t_r\}$

- Goal:

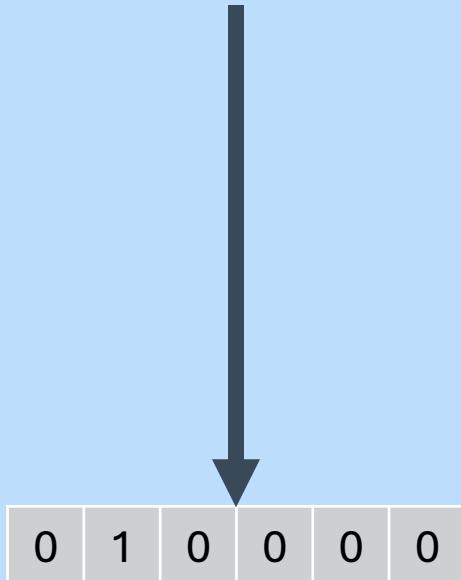
$$\Delta \tilde{f}_t = \tilde{f}_{t, \text{after}} - \tilde{f}_{t, \text{before}} \quad \forall t \in T$$
$$\max_Y \sum_{\{t \in T\}} \mathbb{E}[\Delta \tilde{f}_t]$$

# Foundational Work

# Google's RAPPOR

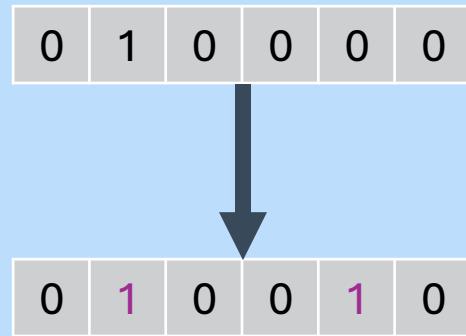
## Encode

$$\mathcal{H} = \{H_1, H_2, \dots, H_m\}$$



$$B_0[i] = 1 \quad \text{if } \exists H \in \mathcal{H}, \text{s.t., } H(v) = i$$

## Perturb



$$\Pr[B_1[i]] = 1 = \begin{cases} p = 1 - \frac{f}{2} & \text{if } B_0[i] = 1 \\ q = \frac{f}{2} & \text{if } B_0[i] = 0 \end{cases}$$

## Aggregate

- Using Linear Regression
- Using LASSO Regression

# Limitations of RAPPOR

- Use of Bloom filter reduced communication cost

However,

- Accuracy decreased significantly
- Computation cost of the aggregation step increased significantly

# SOTA Frequency Estimators

# K Randomized Response

Encode

A	B	C	D	E	F
---	---	---	---	---	---



A	B	C	D	E	F
---	---	---	---	---	---

Perturb

A	B	C	D	E	F
---	---	---	---	---	---



A	D	C	C	F	E
---	---	---	---	---	---

$$\Pr[PE(v) = i] = \begin{cases} p = \frac{e^\epsilon}{e^\epsilon + d - 1} & \text{if } i = v \\ q = \frac{1}{e^\epsilon + d - 1} & \text{otherwise} \end{cases}$$

Aggregate

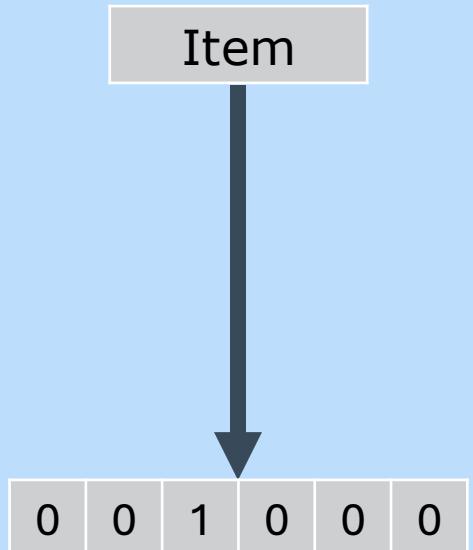
$$\tilde{f}_v = \frac{\frac{1}{n} \sum_{i=1}^n \mathbb{I}_{\{Support(y_i)\}}(v) - q^*}{p^* - q^*}$$

For kRR,

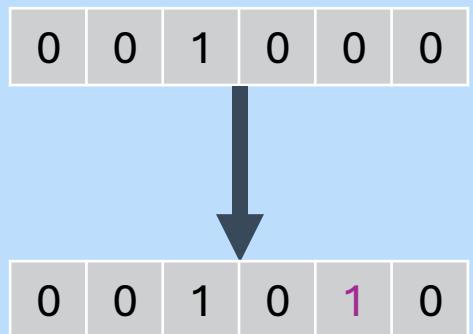
$$Support(y_i) = \{y\}$$

# Optimal Unary Encoding

Encode



Perturb



Aggregate

$$\tilde{f}_v = \frac{\frac{1}{n} \sum_{i=1}^n \mathbb{I}_{\{Support(y_i)\}}(v) - q^*}{p^* - q^*}$$

$$\Pr[PE(v) = 1] = \begin{cases} p = \frac{1}{2} & \text{if } i = v \\ q = \frac{1}{e^\epsilon + 1} & \text{otherwise} \end{cases}$$

For OUE,  
 $Support(y_i) = \{v \mid v \in [d] \text{ and } y_v = 1\}$

# Optimal Local Hashing

Encode

$$\mathcal{H} = \{H_1, H_2, \dots, H_m\}$$



$H_3$



0	0	1	0	0	0
---	---	---	---	---	---

Perturb

0	0	1	0	0	0
---	---	---	---	---	---



0	0	0	0	0	0
---	---	---	---	---	---

$$\Pr[y = \langle H, x \rangle] = \begin{cases} p = \frac{e^\epsilon}{e^\epsilon + d - 1} & \text{if } x = i \\ q = \frac{1}{e^\epsilon + d - 1} & \text{otherwise} \end{cases}$$

Aggregate

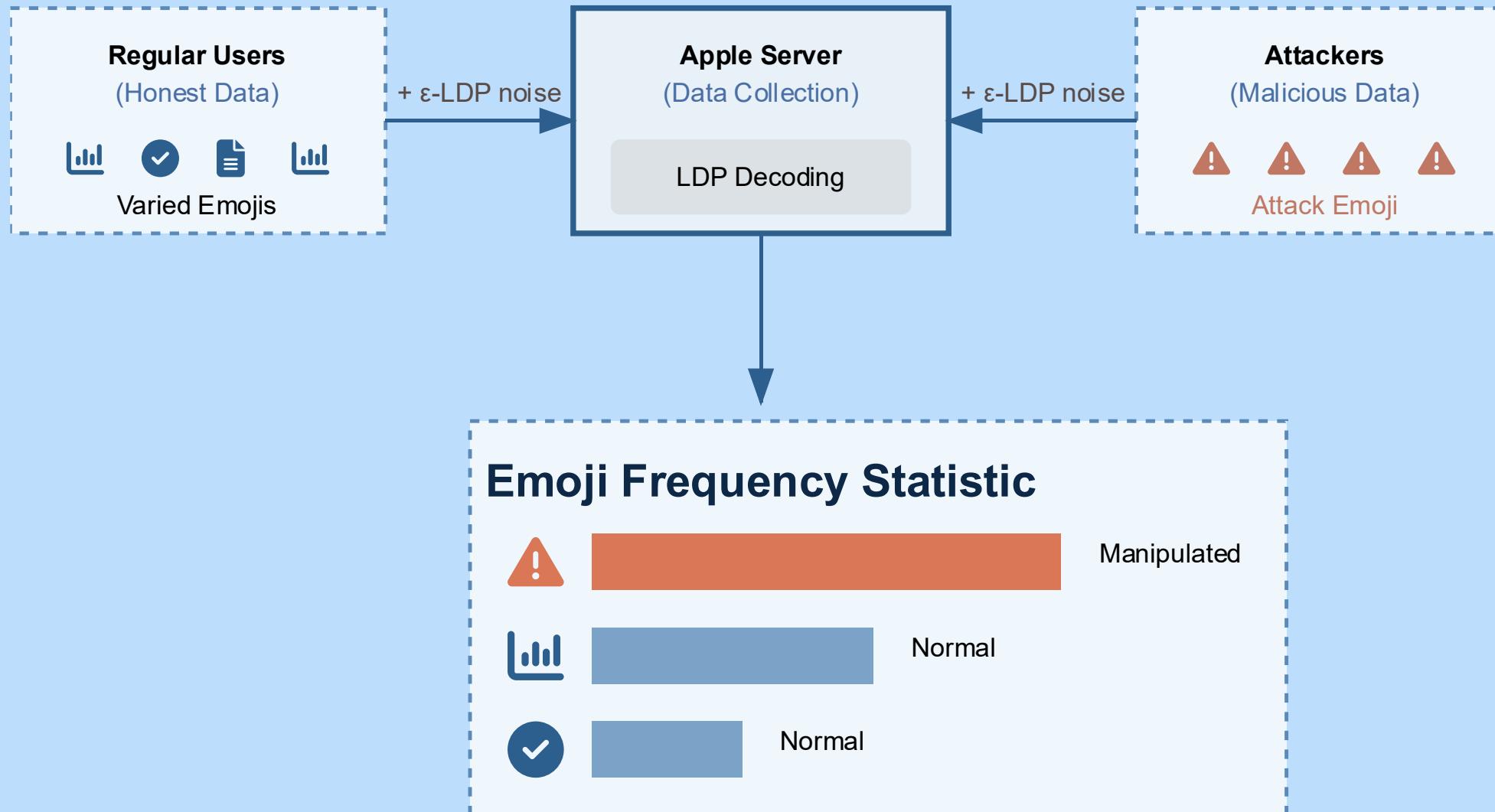
$$\tilde{f}_v = \frac{\frac{1}{n} \sum_{i=1}^n \mathbb{I}_{\{Support(y_i)\}}(v) - q^*}{p^* - q^*}$$

For OLH,

$$Support(y_i) = \{v \mid v \in [d] \text{ and } H(v) = x\}$$

# Maximum Gain Attacks on Frequency Estimators

# LDP Protocol Attack Skewing the Frequently Used Emoji Statistic



# General Formula for Maximum Gain

$$G = \frac{\sum_{\{i=n+1\}}^{\{n+m\}} \mathbb{E}[\mathbb{I}_{Support(y_i)}(t)]}{(n+m)(p^* - q^*)} - \frac{m \sum_{\{i=1\}}^{\{n\}} \mathbb{E}[\mathbb{I}_{Support(y_i)}(t)]}{n(n+m)(p^* - q^*)}$$

# Attacking kRR

- The support set is:

$$\sum_{\{t \in T\}} \mathbb{I}_{Support(y_i)}(t) = 1$$

- Gain becomes:

$$G = \frac{m}{(n + m)(p^* - q^*)} - c$$

- Plugging in  $p^*$  and  $q^*$  gives:

$$G = \beta(1 - f_T) + \frac{\beta(d - r)}{e^\epsilon - 1}$$

# Attacking OUE

- The support set is:

$$\sum_{\{t \in T\}} \mathbb{I}_{Support(y_i)}(t) = r$$

- Gain becomes:

$$G = \frac{rm}{(n+m)(p^* - q^*)} - c$$

- Plugging in  $p^*$  and  $q^*$  gives:

$$G = \beta(2r - f_T) + \frac{2\beta r}{e^\epsilon - 1}$$

# Attacking OLH

- The support set is:

$$\sum_{\{t \in T\}} \mathbb{I}_{Support(y_i)}(t) = r$$

- Gain becomes:

$$G = \frac{rm}{(n+m)(p^* - q^*)} - c$$

- Plugging in  $p^*$  and  $q^*$  gives:

$$G = \beta(2r - f_T) + \frac{2\beta r}{e^\epsilon - 1}$$

# Evaluation

# Estimators: Utility vs. Security

	<b>kRR</b>	<b>OUE</b>	<b>OLH</b>
Communication Cost	$\theta(\log d)$	$\theta(d)$	$\theta(\log n)$
Variance	$n \cdot \frac{d - 2 + e^\epsilon}{(e^\epsilon - 1)^2}$	$n \cdot \frac{4e^\epsilon}{(e^\epsilon - 1)^2}$	$n \cdot \frac{4e^\epsilon}{(e^\epsilon - 1)^2}$
Gain of MGA	$\beta(1 - f_T) + \frac{\beta(d - r)}{e^\epsilon - 1}$	$\beta(2r - f_T) + \frac{2\beta r}{e^\epsilon - 1}$	$\beta(2r - f_T) + \frac{2\beta r}{e^\epsilon - 1}$

$n$  = number of real users

$m$  = number of fake users

$\beta = \frac{m}{n+m}$  = fraction of fake users

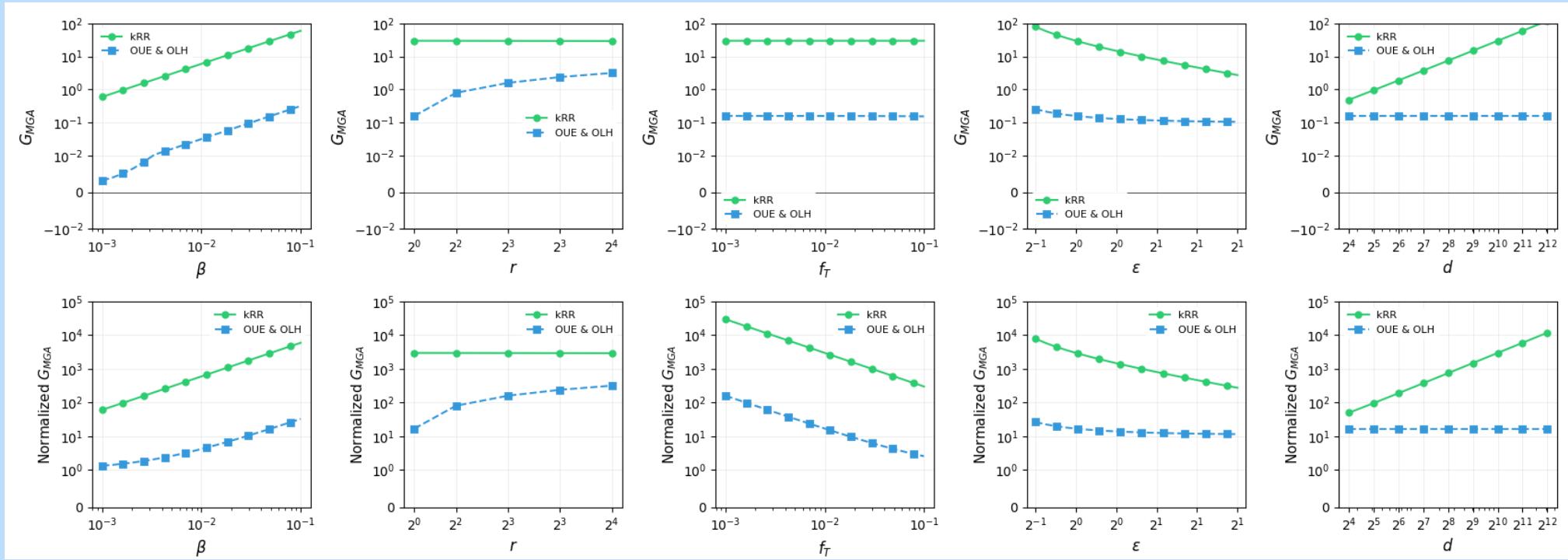
$d$  = domain size

$T$  = set of target items

$r = |T|$

$$f_T = \sum_{t \in T} f_t$$

# MGA Comparison



# Conclusion

# Summary and Open Questions

- Different Frequency Estimation Protocols under LDP
- Attacks on Estimators
  - kRR performs poorly when the domain size increases
  - OUE, OLH perform poorly when the target items increases
- Open Question:
  - Large domain, Large target items
  - How to handle complex data types
  - Utility – Privacy trade-offs

Thank You