

A Unified View of Frequency Estimation and their Attacks on Local Differential Privacy

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Outline



Introduction:

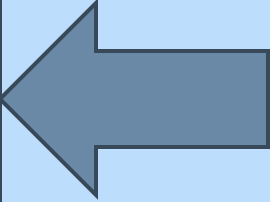
- Differential Privacy
- Local Differential Privacy
- Pure LDP Framework
- Attack Problem



Attacks:

- General Attack Formulation
- Attacking kRR
- Attacking OUE
- Attacking OLH

Frequency Estimation Techniques:

- RAPPOR
 - K Randomized Response (kRR)
 - Optimized Unary Encoding (OUE)
 - Optimized Local Hashing (OLH)
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Evaluation:

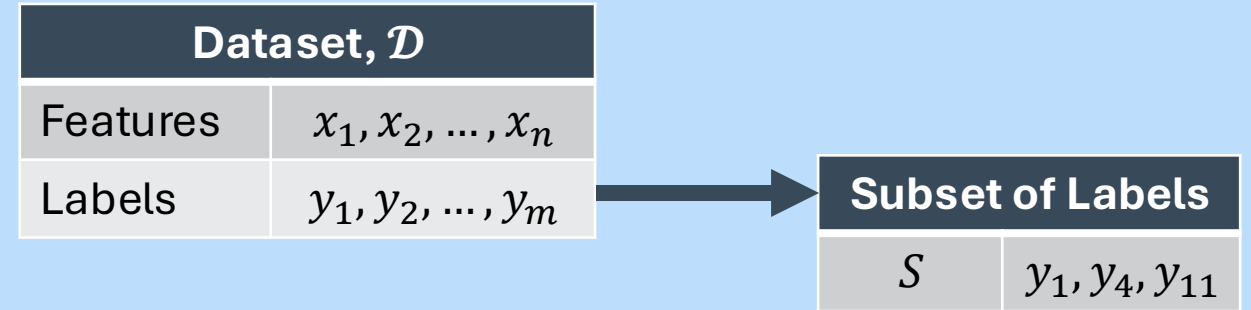
- Comparison between Estimators
- Gain from Attacks
- Impacts of Parameters on Attacks

Conclusion:



Introductory Concepts

Differential Privacy



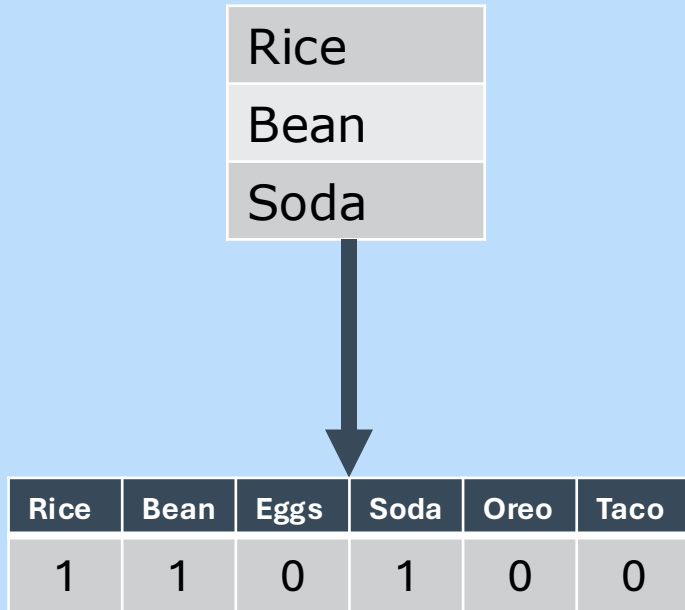
$$\Pr \left[\mathcal{M} \left(\begin{array}{c} x \\ \text{sample}_1 \\ \text{sample}_2 \\ \text{sample}_3 \\ \text{sample}_4 \\ \text{sample}_5 \end{array} \right) \in S \right] \leq e^\epsilon \Pr \left[\mathcal{M} \left(\begin{array}{c} y \\ \text{sample}_1 \\ \text{sample}_2 \\ \text{sample}_3 \\ \text{sample}_4 \\ \text{sample}_{11} \end{array} \right) \in S \right] + \delta$$

Local Differential Privacy

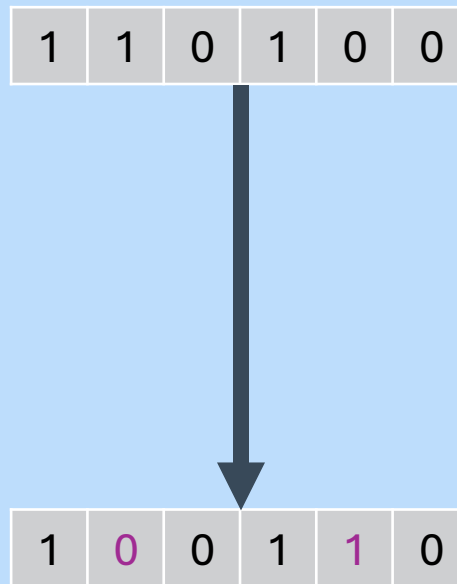
$$\Pr \left[\mathcal{M} \left(\begin{array}{c} x \\ \text{sample}_1 \end{array} \right) \in y \right] \leq e^\epsilon \Pr \left[\mathcal{M} \left(\begin{array}{c} y \\ \text{sample}_{11} \end{array} \right) \in y \right]$$

Protocols for Local Differential Privacy

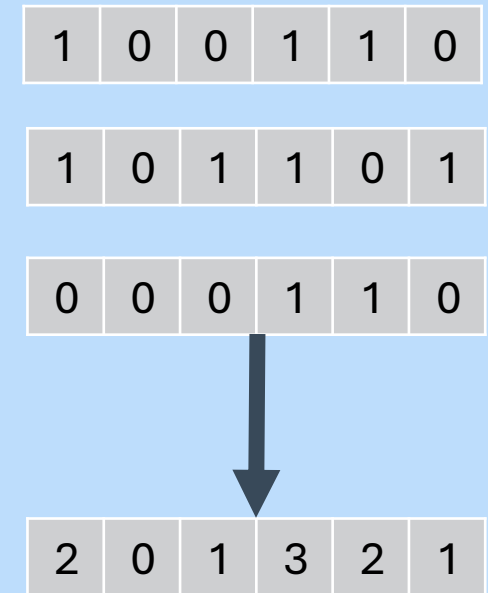
Encode



Perturb



Aggregate



Frequency Estimation Problem

Pure Differential Privacy

- A unified framework --- Can be used to define all frequency protocols.
- Defined based on two fixed probabilities p^* and q^*

$$\begin{aligned}\Pr[PE(v_1) \in \{y \mid v_1 \in \text{Support}(y)\}] &= p^* \\ \Pr[PE(v_2) \in \{y \mid v_1 \in \text{Support}(y)\}] &= q^*\end{aligned}$$

- Based on this framework, we will describe 4 protocols – RAPPOR, kRR, OUE, OLH.

Attacking Estimators – Problem Formulation

- Assume: n real users
- Inject: m fake users
- To increase frequency of r items: $T = \{t_1, t_2, \dots, t_r\}$

- Goal:


$$\Delta \tilde{f}_t = \tilde{f}_{t,after} - \tilde{f}_{t,before} \quad \forall t \in T$$
$$\max_Y \sum_{\{t \in T\}} \mathbb{E}[\Delta \tilde{f}_t]$$

Foundational Work

Google's RAPPOR

Encode

$$\mathcal{H} = \{H_1, H_2, \dots, H_m\}$$




0	1	0	0	0	0
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$$B_0[i] = 1 \quad \text{if } \exists H \in \mathcal{H}, s.t., H(v) = i$$

Perturb

0	1	0	0	0	0
---	---	---	---	---	---



0	1	0	0	1	0
---	---	---	---	---	---

$$\Pr[B_1[i]] = 1 = \begin{cases} p = 1 - \frac{f}{2} & \text{if } B_0[i] = 1 \\ q = \frac{f}{2} & \text{if } B_0[i] = 0 \end{cases}$$

Aggregate

- Using Linear Regression
- Using LASSO Regression

Limitations of RAPPOR

- Use of Bloom filter reduced communication cost

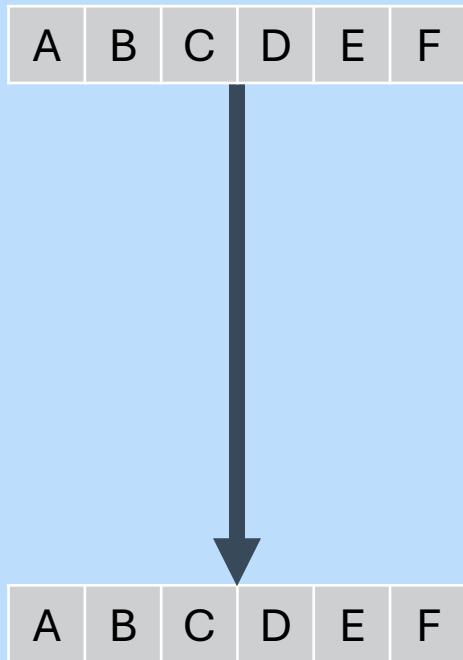
However,

- Accuracy decreased significantly
- Computation cost of the aggregation step increased significantly

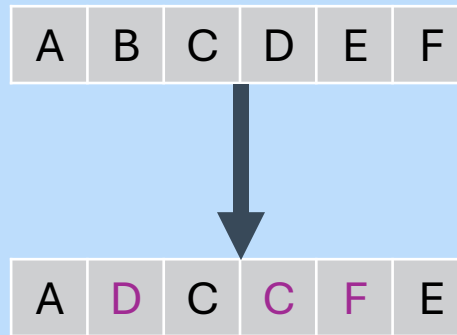
SOTA Frequency Estimators

K Randomized Response

Encode



Perturb



$$\Pr[PE(v) = i] = \begin{cases} p = \frac{e^\epsilon}{e^\epsilon + d - 1} & \text{if } i = v \\ q = \frac{1}{e^\epsilon + d - 1} & \text{otherwise} \end{cases}$$

Aggregate

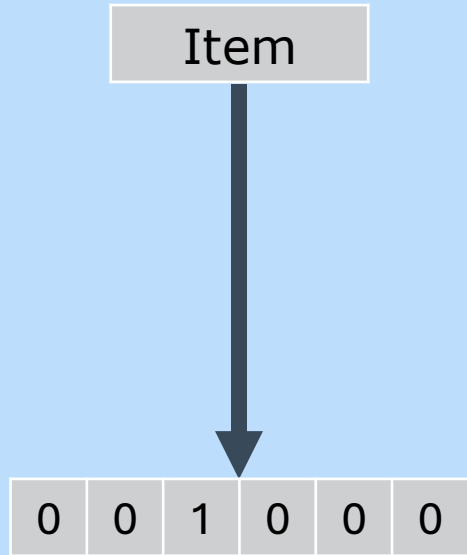
$$\tilde{f}_v = \frac{\frac{1}{n} \sum_{i=1}^n \mathbb{I}_{\{Support(y_i)\}}(v) - q^*}{p^* - q^*}$$

For kRR,

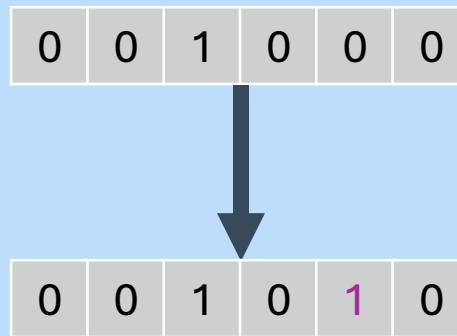
$$Support(y_i) = \{y\}$$

Optimal Unary Encoding

Encode



Perturb



$$\Pr[PE(v) = 1] = \begin{cases} p = \frac{1}{2} & \text{if } i = v \\ q = \frac{1}{e^\epsilon + 1} & \text{otherwise} \end{cases}$$

Aggregate

$$\tilde{f}_v = \frac{\frac{1}{n} \sum_{i=1}^n \mathbb{I}_{\{Support(y_i)\}}(v) - q^*}{p^* - q^*}$$

For OUE,

$$Support(y_i) = \{v \mid v \in [d] \text{ and } y_v = 1\}$$

Optimal Local Hashing

Encode

$$\mathcal{H} = \{H_1, H_2, \dots, H_m\}$$



H_3



0	0	1	0	0	0
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Perturb

0	0	1	0	0	0
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0	0	0	0	0	0
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$$\Pr[y = \langle H, x \rangle] = \begin{cases} p = \frac{e^\epsilon}{e^\epsilon + d - 1} & \text{if } x = i \\ q = \frac{1}{e^\epsilon + d - 1} & \text{otherwise} \end{cases}$$

Aggregate

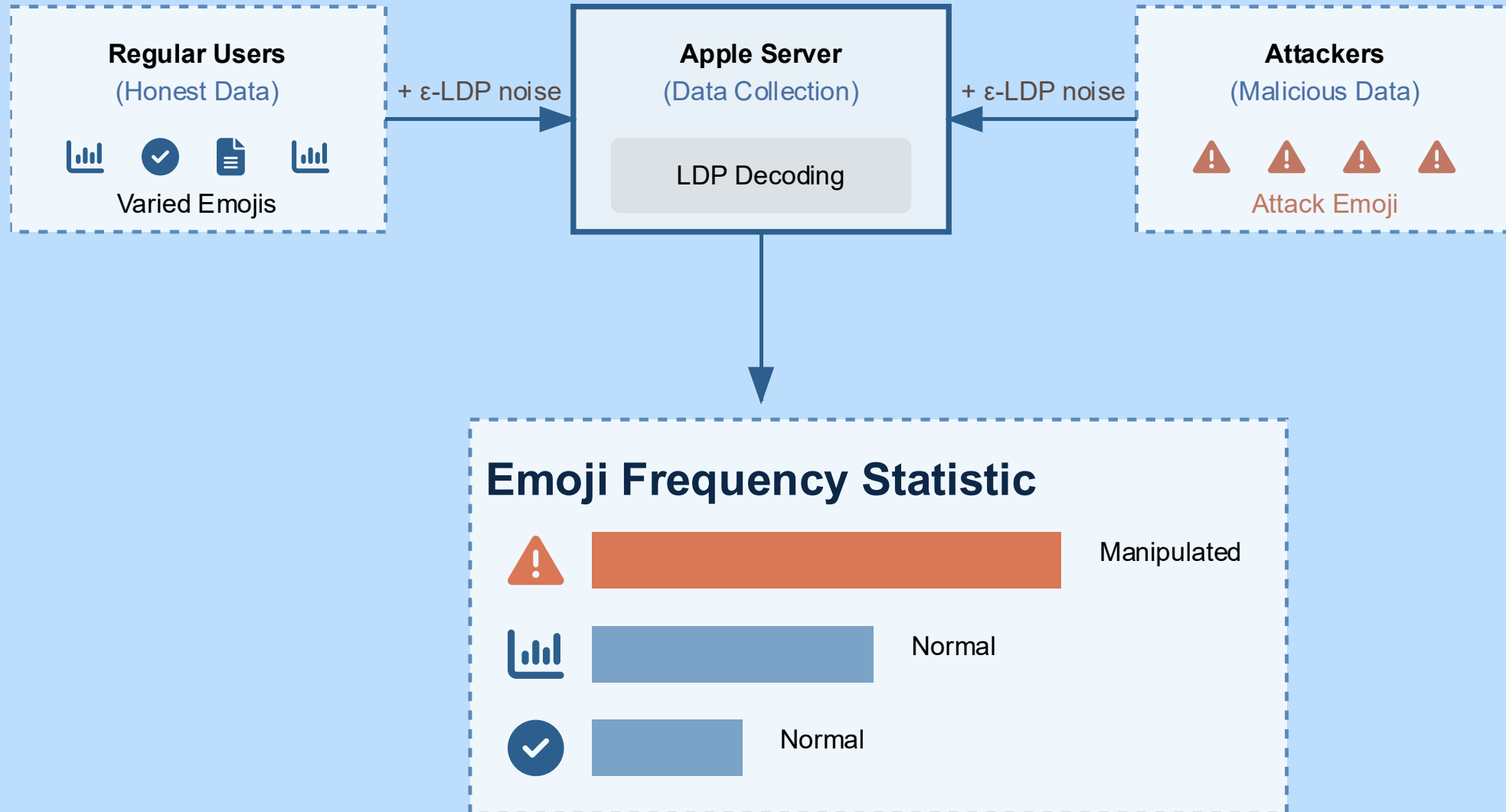
$$\tilde{f}_v = \frac{\frac{1}{n} \sum_{i=1}^n \mathbb{I}_{\{\text{Support}(y_i)\}}(v) - q^*}{p^* - q^*}$$

For OLH,

$$\begin{aligned} \text{Support}(y_i) \\ = \{v \mid v \in [d] \text{ and } H(v) = x\} \end{aligned}$$

Maximum Gain Attacks on Frequency Estimators

LDP Protocol Attack Skewing the Frequently Used Emoji Statistic



General Formula for Maximum Gain

$$G = \frac{\sum_{\{i=n+1\}}^{\{n+m\}} \mathbb{E}[\mathbb{I}_{\text{Support}(y_i)}(t)]}{(n+m)(p^* - q^*)} - \frac{m \sum_{\{i=1\}}^{\{n\}} \mathbb{E}[\mathbb{I}_{\text{Support}(y_i)}(t)]}{n(n+m)(p^* - q^*)}$$

Attacking kRR

- The support set is:

$$\sum_{\{t \in T\}} \mathbb{I}_{\text{Support}(y_i)}(t) = 1$$

- Gain becomes:

$$G = \frac{m}{(n + m)(p^* - q^*)} - c$$

- Plugging in p^* and q^* gives:

$$G = \beta(1 - f_T) + \frac{\beta(d - r)}{e^\epsilon - 1}$$

Attacking OUE

- The support set is:

$$\sum_{\{t \in T\}} \mathbb{I}_{\text{Support}(y_i)}(t) = r$$

- Gain becomes:

$$G = \frac{rm}{(n+m)(p^* - q^*)} - c$$

- Plugging in p^* and q^* gives:

$$G = \beta(2r - f_T) + \frac{2\beta r}{e^\epsilon - 1}$$

Attacking OLH

- The support set is:

$$\sum_{\{t \in T\}} \mathbb{I}_{\text{Support}(y_i)}(t) = r$$

- Gain becomes:

$$G = \frac{rm}{(n+m)(p^* - q^*)} - c$$

- Plugging in p^* and q^* gives:

$$G = \beta(2r - f_T) + \frac{2\beta r}{e^\epsilon - 1}$$

Evaluation

Estimators: Utility vs. Security

	kRR	OUE	OLH
Communication Cost	$\theta(\log d)$	$\theta(d)$	$\theta(\log n)$
Variance	$n \cdot \frac{d - 2 + e^\epsilon}{(e^\epsilon - 1)^2}$	$n \cdot \frac{4e^\epsilon}{(e^\epsilon - 1)^2}$	$n \cdot \frac{4e^\epsilon}{(e^\epsilon - 1)^2}$
Gain of MGA	$\beta(1 - f_T) + \frac{\beta(d - r)}{e^\epsilon - 1}$	$\beta(2r - f_T) + \frac{2\beta r}{e^\epsilon - 1}$	$\beta(2r - f_T) + \frac{2\beta r}{e^\epsilon - 1}$

n = number of real users

m = number of fake users

$\beta = \frac{m}{n+m}$ = fraction of fake users

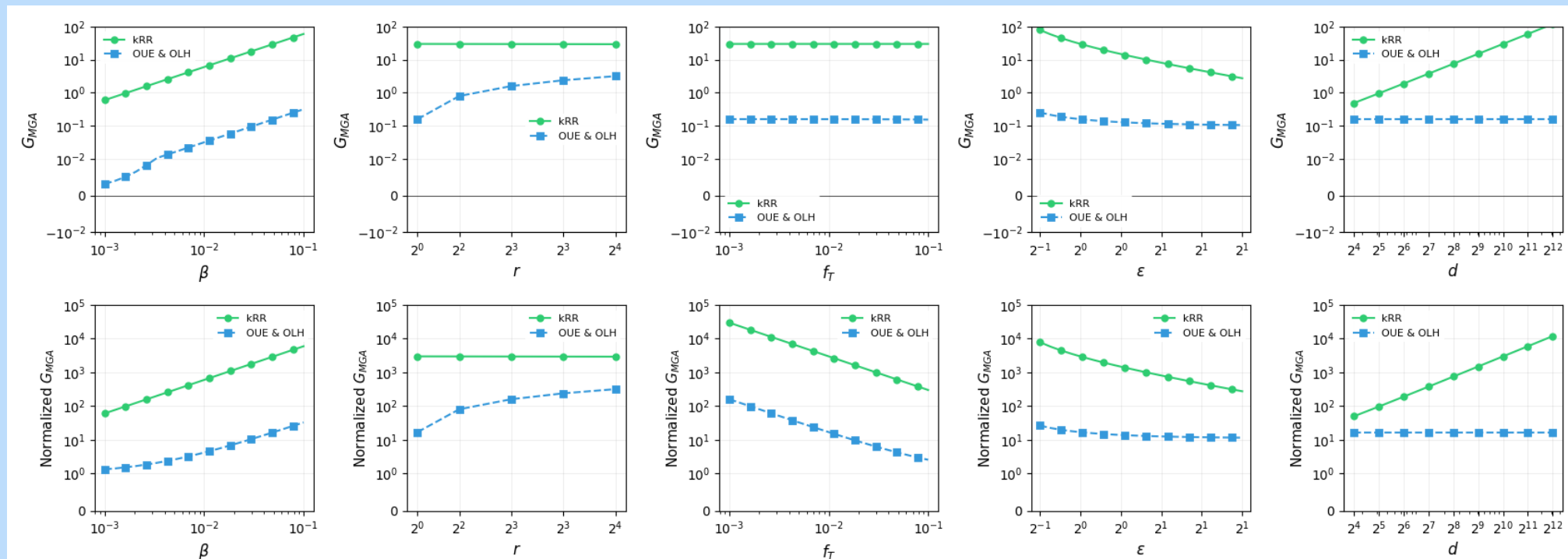
d = domain size

T = set of target items

$r = |T|$

$$f_T = \sum_{t \in T} f_t$$

MGA Comparison



Conclusion

Summary and Open Questions

- Different Frequency Estimation Protocols under LDP
- Attacks on Estimators
 - kRR performs poorly when the domain size increases
 - OUE, OLH perform poorly when the target items increases
- Open Question:
 - Large domain, Large target items
 - How to handle complex data types
 - Utility – Privacy trade-offs

Thank You