CHISELVERIFY: A VERIFICATION LIBRARY FOR CHISEL

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OUTLINE

➤ Hardware Verification: Why do we need it?

➤ Current Verification Solutions for Chisel designs.

- ➤ Our solution: bringing verification to Chisel with Scala.
 - ➤ Functional Coverage
 - ➤ Constrained Random Verification
 - ➤ Bus Functional Models
 - ➤ Timed Assertions

HARDWARE VERIFICATION: WHAT AND WHY?

HARDWARE VERIFICATION: WHAT IS IT?

➤ Verification of a design = Guaranteeing that the expected features, as described in the specification, have been correctly implemented and have the correct behaviour.

- ➤ Verification = testing before tape-out.
- ➤ Validation = testing after tape-out.

HARDWARE VERIFICATION: WHY DO WE NEED IT?

➤ Enables the Computer Engineer to guarantee the working state of a design before spending a lot of money taping it out.

Saves time by allowing for automated and randomised checking, rather than writing test-benches for each possible value reached by a given port.

➤ Shouldn't take up too many resources. Too much time is currently spent on making sure that a design is correct.

CURRENT SOLUTIONS

OVERVIEW OF THE CURRENT SOLUTIONS

➤ For Chisel:

- ➤ ChiselTest: "traditional" test-benches with peek-poke-expect interfaces and forking, lacks verification features.
- ➤ **ScalaTest**: Software testing framework, not ideal for hardware, doesn't simulate the hardware, only checks the Chisel code itself.

➤ For Verilog:

- ➤ **SystemVerilog**: Extension of Verilog that enables object oriented programming and verification features inside of the test-benches.
- ➤ UVM: Verification Methodology, enables a standardised testing method that can be reused for many different DUTs.

CURRENT SOLUTIONS: WHAT'S MISSING?

- ➤ For Chisel:
 - ➤ ChiselTest: For test-benches, not really verification.
 - > ScalaTest: Not made for Hardware.

- ➤ SystemVerilog & UVM:
 - ➤ Too verbose, requires ~800 LOC for a test-bench.
 - ➤ Requires multiple languages to test a Chisel design.

OUR SOLUTION: CHISELVERIFY

CHISELVERIFY: OVERVIEW

- ➤ Hardware Verification library for Chisel, inspired by UVM.
- ➤ Powered entirely by Scala and ChiselTest.

- ➤ ChiselVerify brings the following to the Chisel ecosystem:
 - ➤ Functional Coverage
 - ➤ Constrained Random Verification
 - ➤ Bus Functional Models
 - ➤ Timed Assertions

CHISELVERIFY: FUNCTIONAL COVERAGE

FUNCTIONAL COVERAGE: WHAT IS IT?

- ➤ <u>Statement coverage</u> = **quantitative** approach to getting the verification progress.
 - ➤ How much code have we tested?

- Functional coverage = qualitative approach to getting the verification progress.
 - ➤ Which features have we tested?

FUNCTIONAL COVERAGE: WHAT IS IT?

Verification Plan

Representation of the DUT's expected features.

CoverGroup

Set of DUT ports that will be sampled together, represents a "feature".

CoverPoint

Set of values that a port is expected to have to verify a feature, represents a feature of single port.

Bins

Definition of a set of values that a port should reach during testing, done in two ways:

Range: first to last (both included)

Conditional:
Arbitrary
(portValues) =>
Boolean function

Cross

Defines a relation between two bins in a CoverPoint, i.e. how many value pairs, within the defined cross set, have these two points reached during testing.

Example:

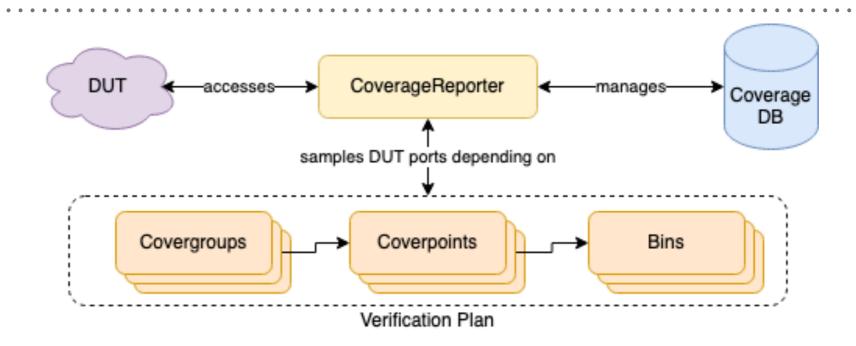
Cross(A, B, 1 to 1, 1 to 1) => Have ports A and B reached the value 1 at the same time?

Timed

Same idea as the cross bin, but with an added time constraint, i.e. have the two points hit a cross set within a given number of cycles from each other?

There are 3 ways to define a time constraint: Always, Eventually and Never

FUNCTIONAL COVERAGE: HOW DID WE DO IT?



- ➤ CoverageDB: DataBase that maintains the values gathered for all of the bins across multiple tests in a test suite.
- ➤ Coverage Reporter: Handles the registration of CoverPoints and Bins to the DB, samples the bin values and creates the coverage report.
 - ➤ This is used to create the verification plan.

FUNCTIONAL COVERAGE: API

```
Rule of thumb: PointType(name, ports)(bins/conditions)
Actual API:
case class CoverPoint(pointName: String, port: Data)(bins: Bins*)
       Bins(val name: String, val range: Range)
       Bins(val name: String, val range: Range, val condition: Condition)
       case class Condition(name: String, cond: Seq[BigInt] => Boolean)
case class CrossPoint(n: String, p: Data*)(b: CrossBin*)
        case class CrossBin(name: String, ranges: Range*)
case class CoverCondition(pointName: String, ports: Data*)(conditions: Condition*)
case class TimedCross(name: String, port1: Data, port2: Data)(delay: DelayType)(bins: CrossBin*)
        case class Never(delay: Int)
        case class Eventually(delay: Int)
        case class Always(delay: Int)
        case class Exactly(delay: Int)
```

FUNCTIONAL COVERAGE: USING IT

➤ Create the coverage reporter and verification plan.

```
val cr = new CoverageReporter(dut)
cr.register(
    //Declare CoverPoints
    CoverPoint("accu", dut.io.outA)( //CoverPoint 1
        Bins("lo10", 0 until 10), Bins("First100", 0 until 100)),
    CoverPoint("test", dut.io.outB)( //CoverPoint 2
        Bins("testLo10", 0 until 10)),
    //Declare cross points
    CrossPoint("accuAndTest", dut.io.outA, dut.io.outB)(
        CrossBin("both1", 1 to 1, 1 to 1))
)
```

➤ Sample the CoverPoints inside of the test.

```
cr.sample()
```

RESULT: FUNCTIONAL COVERAGE REPORT

- Create and print the coverage report //Print coverage report cr.printReport()
- ➤ Example result:

```
======== COVERAGE REPORT =========
========= GROUP ID: 1 ==========
COVER_POINT PORT NAME: accu
BIN lo10 COVERING Range 0 until 10 HAS 10 HIT(S) = 100,00%
BIN First100 COVERING Range 0 until 100 HAS 50 HIT(S) = 50,00%
COVER_POINT PORT NAME: test
BIN testLo10 COVERING Range 0 until 10 HAS 4 HIT(S) = 40,00%
CROSS_POINT accuAndTest FOR POINTS accu AND test
BIN both1 COVERING Range 1 to 1 CROSS Range 1 to 1 HAS 1 HIT(S) = 100,00%
```

CHISELVERIFY: CONSTRAINED RANDOM VERIFICATION

CONSTRAINED RANDOM VERIFICATION: WHAT IS IT?

- ➤ Model tests using randomness:
 - ➤ Give random inputs to the DUT and expect values using a golden model.
- ➤ Guide the randomness using constraints:
 - ➤ We don't want the randomness to be uniformly distributed => use constraints to describe the random distribution.
- ➤ Idea: Add a Constraint Programming Language to Scala/Chisel
 - ➤ Create Random Objects, then define Random Variables and Constraints in it.
 - ➤ Random Objects define a **Constraint Satisfaction Problem (CSP)**. This problem is solved using a CSP Solver.
 - ➤ We use **JaCoP** as a CSP Solver.

CONSTRAINED RANDOM VERIFICATION: SYSTEMVERILOG

- > Create random objects represented by a RandObj class.
- > Store random fields as rand or randc (cyclic) variables.
- ➤ Constraints are defined for the whole object in a **constraint** block.

Example:

```
class frame_t;
rand pkt_type ptype;
rand integer len;
randc bit [1:0] no_repeat;
rand bit [7:0] payload [];
// Constraint the members
constraint legal {
  len >= 2;
  len <= 5;
  payload.size() == len;
}</pre>
```

CONSTRAINED RANDOM VERIFICATION: HOW DID WE DO IT?

- > Same ideas as in SV:
 - ➤ Random objects created by extending the **RandObj** trait using a given **Model** with a seed.
 - ➤ Random fields are added as **Rand/Randc** values inside the class.
 - Constraints are defined using either:
 - ➤ Single constraints: using "#" (e.g. val lenConstraint = len #> 2)
 - ConstraintGroups which are equivalent to SV

```
Example: class Frame extends RandObj(new Model) {
   val pkType: Rand = new Rand(0, 3)
   val len: Rand = new Rand(0, 10)
   val noRepeat: Randc = new Randc(0, 1)

   val legal: ConstraintGroup = new ConstraintGroup {
      len #>= 2
      len #<= 5
   }
}</pre>
```

CONSTRAINED RANDOM VERIFICATION: WHAT DO WE ADD?

- ➤ The list of operator used to construct constraint is the following: #<, #<=, #>, #>=,#=, div, #*, mod, #+, -, #\=, #^, in, inside.
- ➤ We also added conditional constraints using:
 - ➤ IfCon: If a condition is met, then use the constraint
 - ➤ **ElseC**: If a condition is met, then use the constraint, else use an other.

```
Example: val constraint1: crv.Constraint = IfCon(len #= 1) {
          payload.size #= 3
     } ElseC {
          payload.size #= 10
     }
```

CONSTRAINED RANDOM VERIFICATION: USING IT

➤ Each random class exposes a randomize() method, which automatically solves the constraints specified in the class and assign to each random filed a random value.

The method returns true only if the CSP Solver found a set of values that satisfy the current constraints.

Example:

```
val myPacket = new Frame(new Model)
assert(myPacket.randomize)
```

CONSTRAINED RANDOM VERIFICATION: USING IT

Create Random object:

Instantiate and use it:

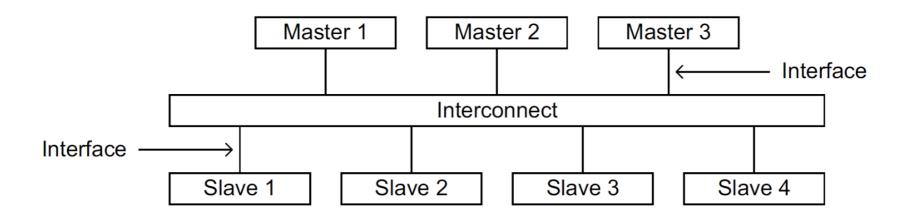
```
//Instantiate RandObj
class Packet extends RandObj(new Model(3)) {
                                                     val pckt = new Packet
    val idx = new Randc("idx", 0, 10)
   val size = new Rand("size", 1, 100)
                                                     //Check that the constraints were solvable
   val len = new Rand("len", 1, 100)
                                                     assert(pckt.randomize)
    val payload: Array[Rand] = Array.tabulate(11)(
        new Rand(s"byte[$_]", 1, 100)
                                                     // [...] ChiselTest boilerplate [...]
                                                     //Example use of random variables in a DUT
   //Example Constraint with operations
                                                     while(pckt.randomize && /*Coverage isn't 80%*/) {
    val single = payload(0) #= (len #- size)
                                                         dut.portA.poke(
                                                             pckt.payload(pckt.idx.value()).value()
   //Example conditional constraint
   val conditional = IfCon(len #= 1) {
        payload.size #= 3
                                                         // [...] Sample the coverage and update constraints
    } ElseC {
                                                     }
        payload.size #= 10
    val idxConst = idx #< payload.size</pre>
```

CHISELVERIFY: BUS FUNCTIONAL MODELS

BUS FUNCTIONAL MODEL: WHAT IS IT?

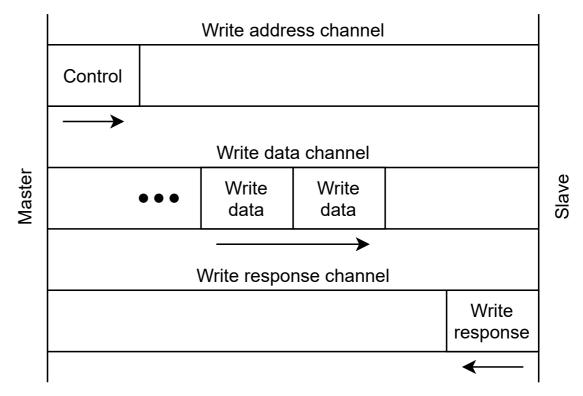
- ➤ Abstraction of the inner workings of a standardised interface.
- ➤ Allows for the use of data transfer via **Transactions**, rather than having to deal with the inner wiring manually.
- ➤ Software abstraction, is useful for faster verification.

➤ We chose to create a first BFM for the AXI4 Bus:

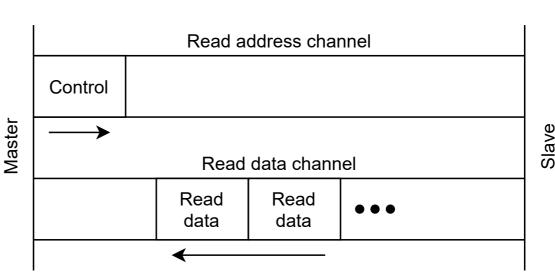


BUS FUNCTIONAL MODEL: TRANSACTIONS

- > Rather than setting every wire manually, use either:
 - ➤ WriteTransaction



➤ ReadTransaction



BUS FUNCTIONAL MODEL: HOW TO USE IT

➤ Create a **FunctionalMaster** for your AXI interfaced DUT.

```
val master = new FunctionalMaster(dut)
```

- ➤ Create transactions:
 - ➤ Write:

```
master.createWriteTrx(0, Seq(42), size = 2)

var resp = master.checkResponse()
while (resp == None) {
    resp = master.checkResponse()
    dut.clock.step()
}
```

➤ Read:

```
master.createReadTrx(0, size = 2)

var data = master.checkReadData()
while (data == None) {
    data = master.checkReadData()
    dut.clock.step()
}
```

BUS FUNCTIONAL MODEL: API

➤ WriteTransaction API:

```
createWriteTrx(
    addr: BigInt, data: Seq[BigInt], id: BigInt, len: Int,
    size: Int, burst: UInt, lock: Bool, cache: UInt,
    prot: UInt, qos: UInt, region: UInt, user: UInt
)
```

➤ ReadTransaction API:

```
createReadTrx(
    addr: BigInt, id: BigInt, len: Int, size: Int,
    burst: UInt, lock: Bool, cache: UInt, prot: UInt,
    qos: UInt, region: UInt, user: UInt
)
```

- ➤ addr: Start write address, must fit in the slave's address width.
- ➤ data: List of data to write, defaults to random data, entries in data must fit within the slave DUT's write data width, and the list can have at most len entries.
- ➤ id: Transaction id, defaults to 0, id must fit within DUT's ID width, likewise size cannot be greater than the DUT's write data width.
- ➤ len: Burst length, defaults to 0 (i.e. 1 beat).
- ➤ size: Beat size, defaults to 1B.
- **burst**: Burst type, defaults to FIXED.

CHISELVERIFY: TIMED ASSERTIONS

TIMED ASSERTIONS: QUICK OVERVIEW

- ➤ Allows to create predicated assertions that take into account certain timing delays.
- > Types of delays: Given a delay of x cycles:
 - **Exactly**: Assertion is true exactly in x cycles.
 - **Eventually**: Assertion is true at least one in the next x cycles.
 - ➤ **Always:** Assertion is true every cycle for the next x cycles.
 - ➤ **Never:** Assertion isn't true at any cycle for the next x cycles.
- ➤ Two types of Timed Assertions:
 - ➤ ExpectTimed: Uses *ChiselTest*'s expect for the assertion.
 - ➤ **AssertTimed:** Uses a software assert() for the assertion.

TIMED ASSERTIONS: USING IT

ExpectTimed[T <: Module](dut: T, port: Data, expectedVal: UInt, message: String)(delayType: DelayType)</pre> AssertTimed[T <: Module](dut: T, cond: () => Boolean, errorMsg: String)(delayType: DelayType) Example: AssertTimed(dut, () => dut.io.aEvEqC.peek().litValue() == 1, "a eventually isn't c")(Eventually(11)).join() //And the same thing but with expect ExpectTimed(dut, dut.io.aEvEqC, 1.U, "aEqb expected timing is wrong")(Exactly(6)).join()

CONCLUSION

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➤ ChiselVerify brings verification to the Chisel ecosystem.

➤ High-Level Functional backend (i.e. Scala) allows for much more efficiency during verification process (in comparison to SystemVerilog with UVM).

➤ Can be used to verify non-Chisel designs as well thanks to Chisel Blackboxes.

REFERENCES

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- ➤ K. Kuchcinski and R. Szymanek, "Jacop java constraint programming solver," 2013, cP Solvers: Modeling, Applications, Integration, and Standardization, co-located with the 19th International Conference on Principles and Practice of Constraint Programming; Conference date: 16-09-2013.
- ➤ Xillinx AXI reference guide: https://www.xilinx.com/support/documentation/ip_documentation/
 https://www.xilinx.com/support/documentation/ip_documentation/
 axi_ref_guide/latest/ug1037-vivado-axi-reference-guide.pdf

Current Project repository:

https://github.com/chiselverify/chiselverify/

➤ Project Wiki (Good way to get started):

https://github.com/chiselverify/chiselverify/wiki/

QUESTIONS?