

Time complexity of Question
No. 2 and 4.

2. NO. Q. Ans:

Implementation-2
def fibonacci_2(n):

if n <= 0:

print(" ")

elif n <= 2:

return n-1

else:

return fibonacci_2(n-1) + fibonacci_2(n-2)

$\hookrightarrow O(T(n-1) + T(n-2))$

\therefore Time complexity,

$$T(n) = T(n-1) + T(n-2) + 1$$

Assume the recurrence is,

$$\begin{aligned} T(n) &= T(n-1) + T(n-1) + 1 \\ &= 2T(n-1) + 1 \end{aligned}$$

Base case,
 $T(0) = 1$

P.T.O

$$= 2(2(n-2) + 1) + 1$$

$$= (2^2 T(n-2) + 2 + 1)$$

$$= 2^3 T(n-3) + 2^2 + 2 + 1$$

$$= 2^{n+2} \vdots T(n-n+2) + 2^{n+1} + 2^n + \dots$$

$$= 2^{n+2} \cdot 1 + 2^{n+1} + 2^n + \dots \quad [\text{using base case}]$$

$$= 2^{(n+2)+1} - 1$$

$$= 2^{n+3} - 1$$

$$= 2^n \cdot 2^3 - 1$$

$$= 2^n$$

\therefore Time complexity $= 2^n$

implementation-2

```
def fibonacci_2(n):  
    fibonacci_array = [0, 1]  $\longrightarrow O(1)$   
    if n < 0:  $\longrightarrow O(1)$   
        print(---)  
    elif n <= 2:  $\longrightarrow O(1)$   
        return fibonacci[n-1]  
    else:  
        for i in range(2, n):  
            fibonacci_array.append(fibonacci_array[i-2]  
                                    + fibonacci_array[i-1])  
             $\longrightarrow O(n)$ 
```

\therefore Time complexity becomes: ~~$O(1) + O(1) + O(1)$~~

$$\begin{aligned} &= O(1) + O(1) + O(1) + O(n) \\ &= O(n) \text{ (Ans)} \end{aligned}$$

4. No. of Ans:

Part B : Time complexity:

```
def Multiply-matrix(A, B):  
    n = len(A)  —————  $O(1)$   
    c = []  —————  $O(1)$   
    for m in range(n):  —————  $O(n)$   
        s = [0] * 3  —————  $O(1)$   
        c.append(s)  —————  $O(1)$   
        for i in range(n):  —————  $O(n)$   
            for j in range(n):  —————  $O(n)$   
                for k in range(n):  —————  $O(n)$   
                    c[i][j] = int(A[i][k]) * int(B[k][j])  —————  $O(1)$   
    return c
```

$O(n \times n \times n)$
 $= O(n^3)$

Time complexity:

$$(1 + (O(1) + O(1)) + O(n) + O(n^3))$$
$$= O(n^3)$$