

Diffie - Hellman Key - Exchange / Agreement algorithm

- whitefield diffie and martin Hellman devised solution to the problem of key no agreement or key exchange, in 1976
- this solution is called the Diffie - Hellman key exchange / agreement algorithm
- The two parties , who want to communicate securely , can agree on a symmetric key using this Techniques
- used only for key agreement but not for encryption or decryption of messages.

* Alice and bob want to agree upon a key to be used.

I] Firstly , Alice and bob agree on two large prime numbers , n & g these two integers need not be kept secrete . Alice & bob use an insecure

* man in the middle attack

* Elliptic curve cryptography

- this is asymmetric public crypto system
- it provides equal security with small key size as compared with other non Ecc Algorithm
- it makes use of elliptic curve.

$$\text{eqn. : } y^2 = x^3 + ax + b$$

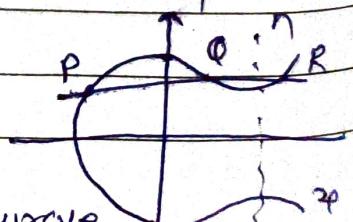


Fig Elliptic curve

* Properties

Elliptic curve is symmetric to x-axis

- Trapdoor function

$$E_g(a, b)$$

$$Q = k \times P$$

\vdots a, b, q elements of elliptic curve

- Key generation at Alice

Select private key n_A $\vdots n_A < n$

- calculate public key P_A

$$P_A = n_A \times G \rightarrow \text{generator point on EC}$$

Key generation at Bob

Select private key n_B $\vdots n_B < n$

calculate public key

$$P_B = n_B \times G$$

$$\text{secret key } K = n_B \times P_A$$

$$\text{secret key } K = n_A \times P_B$$

* encryption :- $M \rightarrow \text{Message}$

$$cm \rightarrow \{ K \cdot G, P_m + K \cdot P_B \}$$

\vdash Public key of Bob

* decryption :- $K \cdot G + n_B$

$$P_m + K \cdot P_B - K \cdot G \times n_B$$

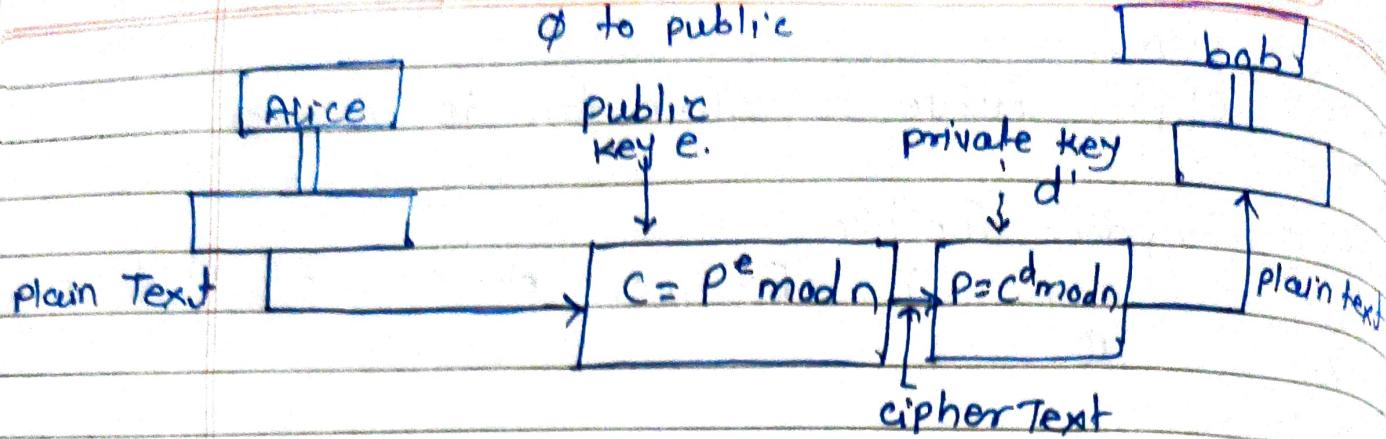
$$P_m + K \cdot P_B - K \cdot P_B = P_m$$

$$\vdash P_m$$

* RSA algorithm (R-Rivest, S-Shamir, A-Adleman)

- most common public key algorithm

- it uses two numbers e & d
public key & private key



* selection of keys

1) bob chooses two very large prime number

$$p \& q$$

2) bob multiplies two prime number ($p * q$) to find n ($n = p * q$)

3) bob calculate another no. which is ϕ
 $\phi = (p-1)*(q-1)$

4) bob chooses random integer 'e' hidden neither calculates d ($dxe \bmod \phi = 1$)
 bob Announces e & n to the public & keeps

$$p = 3, q = 11$$

$$n = 3 \times 11 \\ = 33$$

$$\phi = (p-1)*(q-1) \\ = (3-1)*(11-1)$$

$$\phi = 20$$

$1 < e < \phi$ $e = 3$ [$e \& \phi$ does not have any common factor]

$P = 5$ plain text

$$C = P^e \bmod n$$

$$= 5^3 \bmod 33$$

$$C = 26$$

$$d = 7$$

Data

* digital encryption standard (DES)

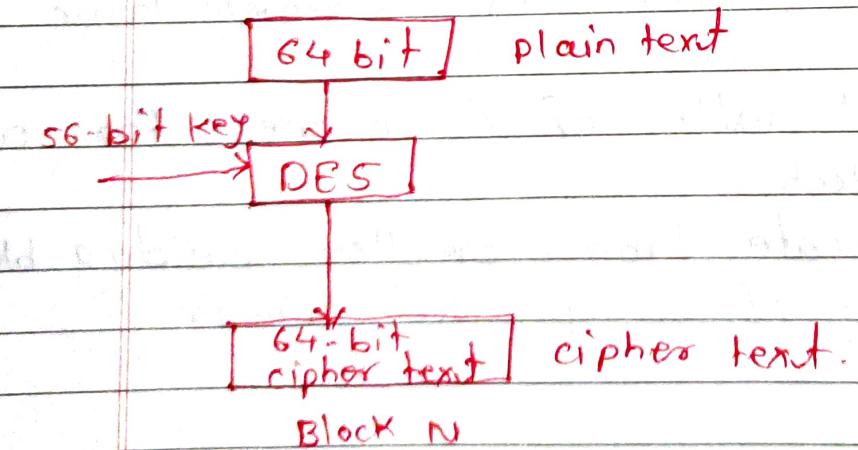
- Most widely used block cipher in world
- adopted in 1977 by NBS (now NIST)
 - as FIPS 46
- encrypt 64-bit data using 56-bit key
- has widespread use
- has been considerable controversy over its security

IBM developed Lucifer cipher

- by team led by Feistel in late 60's
- used 64-bit data blocks with 128-bit key

18-04-22

- DES is a block cipher.
- The algorithm and key are used for encryption with minor difference.
- The key length is 56 bit



→ these bits
discarded

- bit positions are 8, 16, 24, 32, 40 ↑ 56 & 64
- check parity of every bit they will be next contains any error.

Data

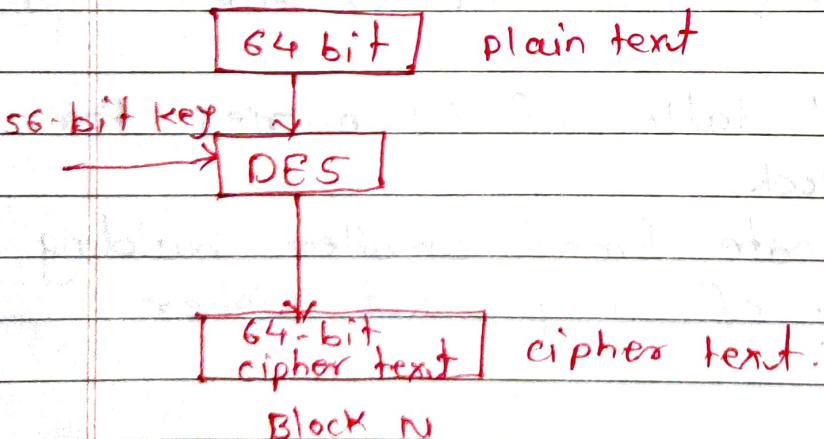
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- bit positions 1's, 8, 16, 24, 32, 40 ↑ 56 of 64
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→ these bit discarded

* Initial permutation (IP)

- First step of the data computation.
- IP records the i/p data bit.
- Even no bits in LH & odd bits in RH

Step 1: Key Transformation

- From the 56-bit key, a different 48 bit subkey is generated during each round using a process called Key Transformation
- The 56-key is divided into two halves, each of 28-bit
- These halves are circuitly shifted^{left} by one or two bit position.

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Step 2] expansion permutation

→ Right plain text

- RPT is expanded from 32 bit to 48 bit
 - Besides increasing the bit size from 32 to 48 bit are permuted as well, hence the name expansion permutation.
- the RPT is divided into 8 bit
 - Next, each 4-bit block of the above step is then expanded to corresponding 6-bit block.

1) draw DSF flowgraph to process a plain text block of 64 bit.

- The data encryption standard is also called as data encryption algorithm by ANSI

* Basic principles.

- DES is a block cipher

- it encrypts data in blocks of 64 bit each

- 64 input and 64 are output.

* Basic workflow of DES algorithm

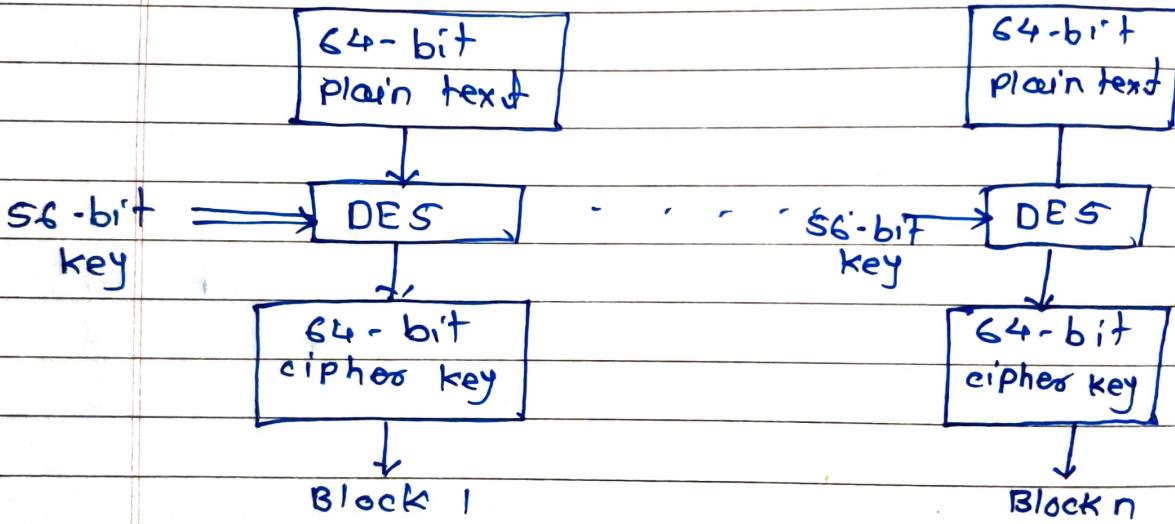


Fig. conceptual working of DES

- 8, 16, 24, 32, 40, 48, 56, 64 discarded key
- before discarding, these bits can be used for parity checking to ensure that the key does not contain any errors.
- substitution & transposition.

* flowgraph of plain text to cipher text.

step 1 :-

[plain text (64 bit)]

step 2 :-

[Initial permutation
CP]

step 3 :-

[LPT] [RPT]

step 4 :- Key →

[10 rounds]

[10 rounds] ← Key

step 5 :-

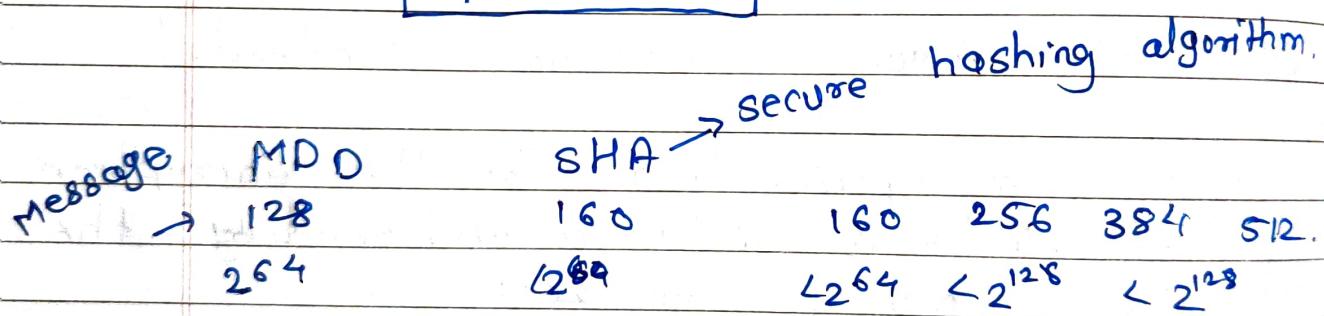
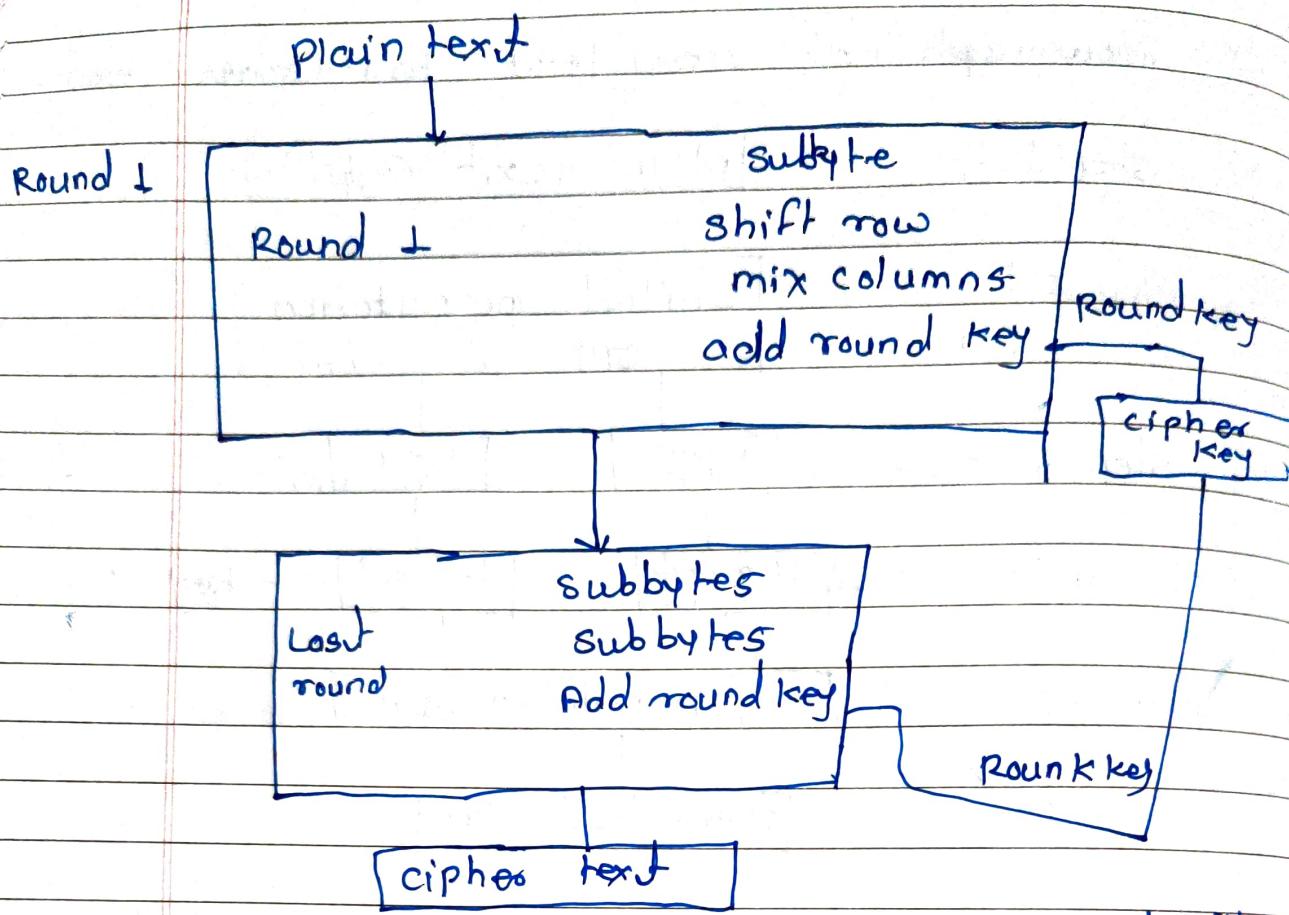
[Final permutation F_P]

step 6 :-

[cipher text (64 bit)]

Q1 AES [Advance encryption standard.]

- AES is specification for the encryption of electronic data established by the U.S (NIST) in 2001
- stronger than triple des.
-
- AES is block cipher.
- the key size can be 128 / 192 / 256 bit
- encrypts data in blocks of 128 bits each
- i/p 128 bits & output cipher text 128 bits.
- no of rounds depend on key length
 - 128 key → 10 rounds.
 - 192 key → 12 rounds
 - 256 key → 14 rounds



Q1 compare DES with AES algorithm.

SR NO	AES	DES
1.	Advance encryption standard	Data encryption standard
2.	byte oriented key length can be 128, 192, 256 bits	The key length is 56 bit DES
3.	AES can encrypt 128 bit of plain text	DES can encrypt 64 bit of plain text
4.	more secure than triple DES	less secure than AES
5.	structure based on a substitution NW	the structure based on a feistel NW
6.	No. of rounds depends on key length eg. 10 rounds (128bit)	DES involves 16 rounds.

* advantages of ~~block~~ blowfish algorithm

1. it is faster and much better than DES encryption.
2. it uses only primitive operations such as addition, XOR and table look-up, making its design and implementation simple.
3. The blowfish algorithm also has a lesser amount of operations to complete compared to other encryption algorithm.
4. Blowfish is fast block cipher except when changing key.

* disadvantages.

1. blowfish uses a 64-bit block size which makes vulnerable to attacks.
2. A reduced round variant of blowfish is known to be susceptible to known plain text.

* Application.

1. Bulk encryption
2. Packet encryption (ATM Packet)
3. password hashing

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* RSA algorithm (Rivest Shamir Adleman)

- RSA algorithm is asymmetric cryptography algorithm
- Asymmetric actually means that two different keys i.e. public key and private key

RSA encryption algorithm.

- choose any 2 (large) prime numbers
- compute $n = p * q$
- compute $\phi(n) = (p-1) * (q-1)$
- choose e such that $1 < e < \phi(n)$ and $\phi(n)$ are coprime.
- compute a value for d such that $(d * e) \% \phi(n) = 1$
- public key is (e, n)
- private key is (d, n)
- the encryption of message m is $c = m^e \% n$
- the decryption of cipher text c is $m = c^d \% n$

example.

choose $p = 3$ & $q = 11$

compute $n = p * q = 33$

compute $\phi(n) = (p-1) * (q-1) = (3-1) * (11-1) = 20$

$$e = 7$$

$$(d * e) \% \phi(n) = 1 \quad [d = 3 \quad (3 * 7) \% 20 = 1]$$

public key is $(e, n) = (7, 33)$

private key is $(d, n) = (3, 33)$

the encryption of $m = 2$ is $c = 2^7 \% 33 = 29$

$$c = 29 \quad \text{is } m = 29^3 \% 33 = 2$$

Q. compare public key encryption with conventional encryption.

Sr. NO. conventional.
public key en

1. it is type of crypto-graphic system which uses a single key both encrypt the Message and decrypt it

public
conventional key
encryption.
it is a type of encryption scheme which instead of a single key, uses a pair of keys to encrypt the message and decrypt it

2. The same secrete key is shared by the sender and the recipient and must be kept secrete at all times.

The public can be shared freely to anyone while the private key kept secrete & is known only the recipient

- Faster

- Slower.

3. it is ~~more~~ less secure

it is more secure.

* principle of public key crypto systems :-

- Public key cryptography is cryptographic technique that involves 'two distinct key' for encryption and decryption.
- that's why it's also known as asymmetric key cryptography
- it is totally based on mathematical function.

* There are two basic principles of any cryptosystem .

- ① confidentiality
- ② authenticity

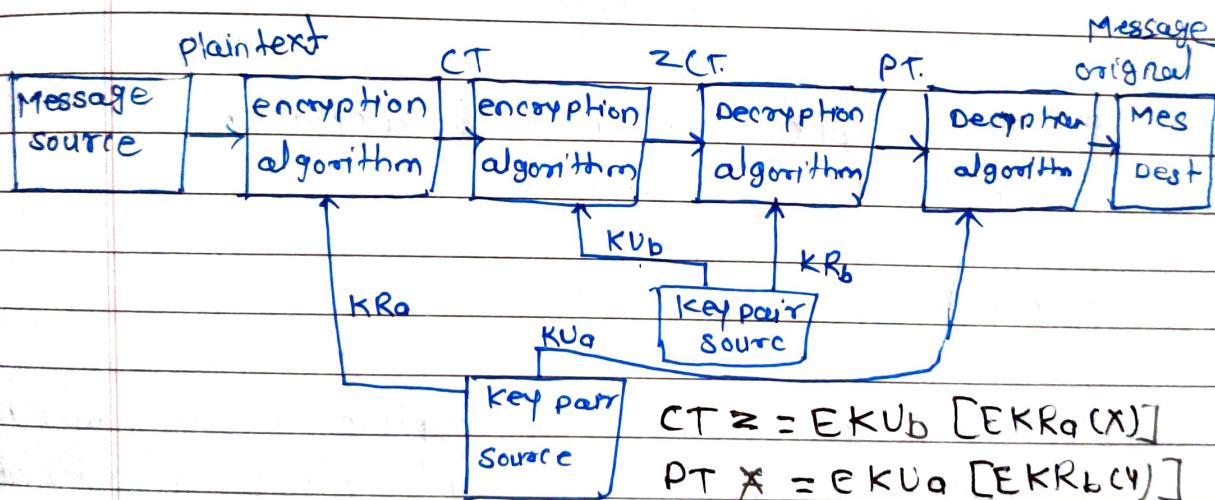


fig public key crypto system.

explanation points:-

- ① plaintext
- ② encryption algorithm.
- ③ public and private keys.
- ④ cipher text
- ⑤ Decryption algorithm.

* Requirements of public key cryptosystem

- it is computationally easy for the Sender A knowing the public key and message to be encrypted M , to generate the corresponding ciphertext $C = EK_{Ub}(M)$
- it is computationally easy for Receiver B to decrypt the resulting ciphertext using the private key to recover the original message $M = DK_{Rb}(C) = DK_{Rb}[EK_{Ub}(M)]$

* public key cryptanalysis.

public key encryption scheme is vulnerable to a brute force attack. The counter measure is to use large keys.

* Block cipher principles & algorithm (Blowfish, AES, DES)

* Block cipher principles.

- Block cipher is an encryption method which divides the plain text into blocks of fixed size block has an equal number bits.
- at a time block cipher operates only one block of plaintext and applies key on it to produce the corresponding block of ciphertext.

eg. Data encryption standard.

* Block cipher principles.

① Number of rounds :-

The number of rounds judges the strength of block cipher algorithm. It is considered that more it is the number of round difficult is cryptoanalysis to break the algorithm. eg. DES has 16 rounds.

② Design of Function F

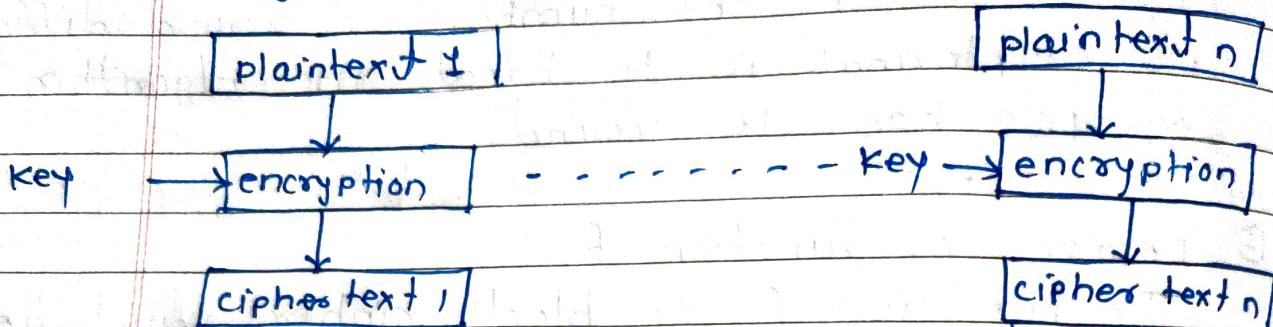
- The function F of block cipher must be designed such that it must be impossible for cryptoanalysis to uncrumble the substitution.
- The criterion that strengthens the function F is its non-linearity.
- More the function F is nonlinear, more it would be difficult to crack it.
- While designing the function F, it should be confirmed that it has good avalanche property which states that a change in one bit of input must reflect the change in one bit of output. In other words, if a change in one bit of input reflects the change in many bits of output.

③ Key schedule algorithm

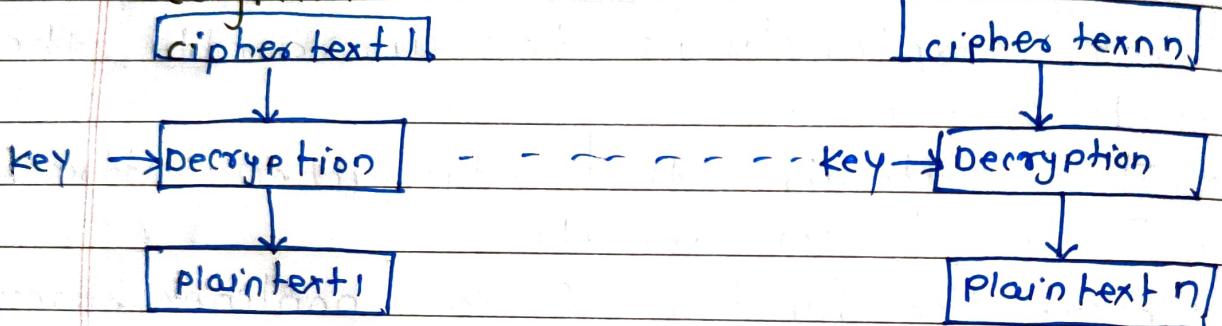
- In the Feistel block cipher structure, each round would generate a sub-key for increasing the complexity of cryptoanalysis.
- The avalanche effect makes more complex in deriving sub-key.

* Block cipher modes of operation.

1) ^{book} electronic code book (ECB) :-
encryption.



* Decryption.



* Advantages of ECB

- Parallel encryption of blocks of bits is possible, thus the faster way of encryption.
- Simple way of the block cipher.

* Disadvantages of ECB

- prone to cryptanalysis since there is a direct relationship between plaintext & ciphertext.

② cipher Block chaining

Initial vector
↓
encryption.

IV
[plaintext 1]
↓
XOR

Key → encryption

↓
ciphertext 1

Plaintext n

XOR

Key → encrypt

ciphertext n

Decryption

ciphertext 1 | C2

key → Decrypt

XOR

plaintext 1

C2

DE

XOR

P2

Cn

DE

XOR

Pn

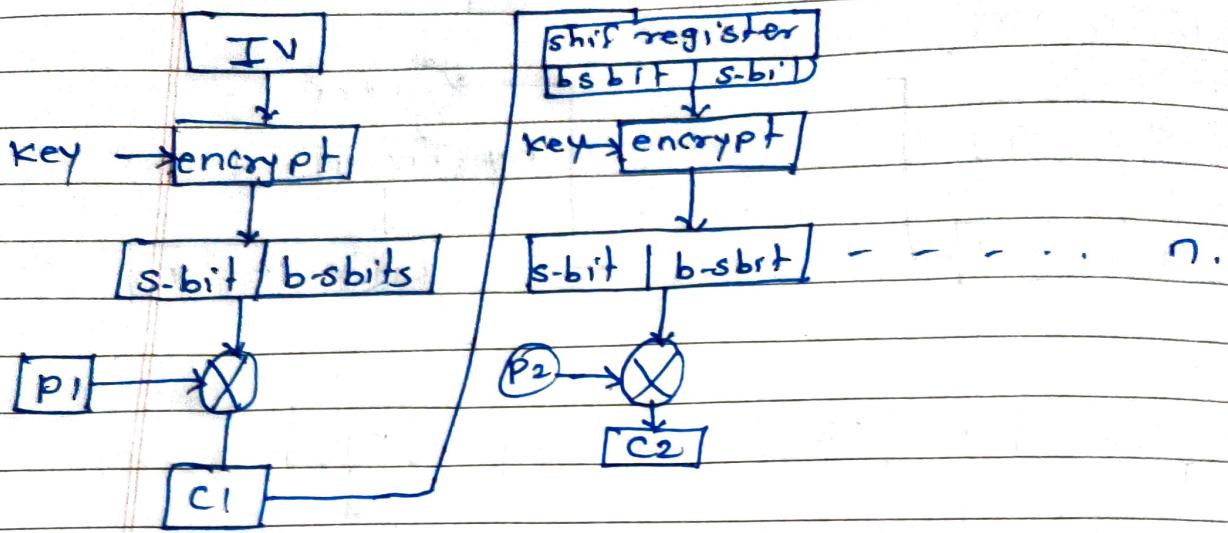
* Advantages of CBC :-

- CBC is a good authentication mechanism
- better cryptoanalysis than ECB
- CBC works well for input greater than b bits

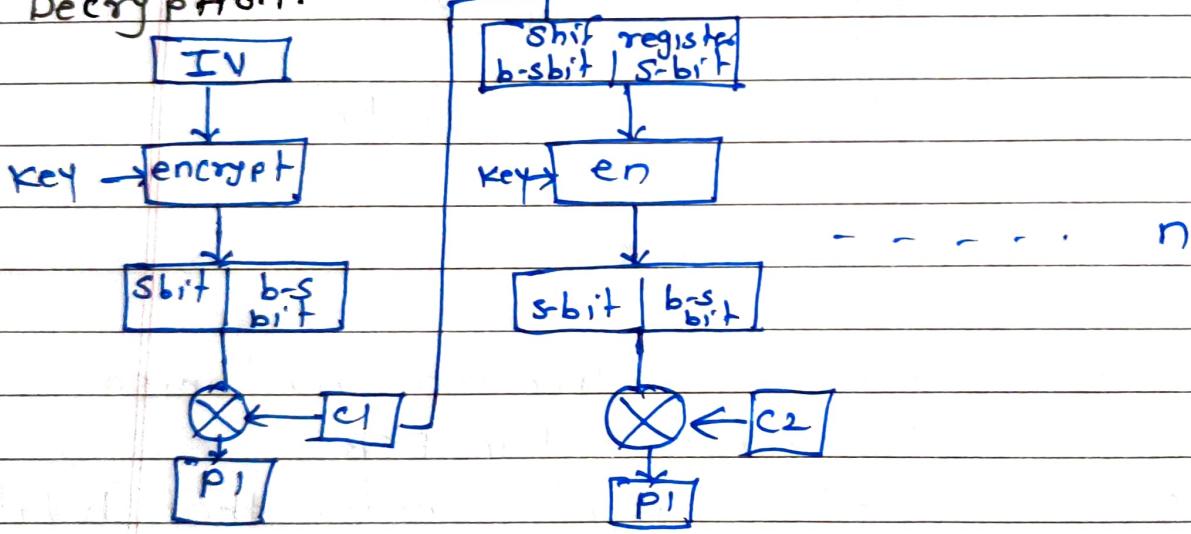
* Disadvantages.

- parallel encryption is not possible since every encryption requires a previous cipher.

③ cipher feedback mode.
 * encryption.



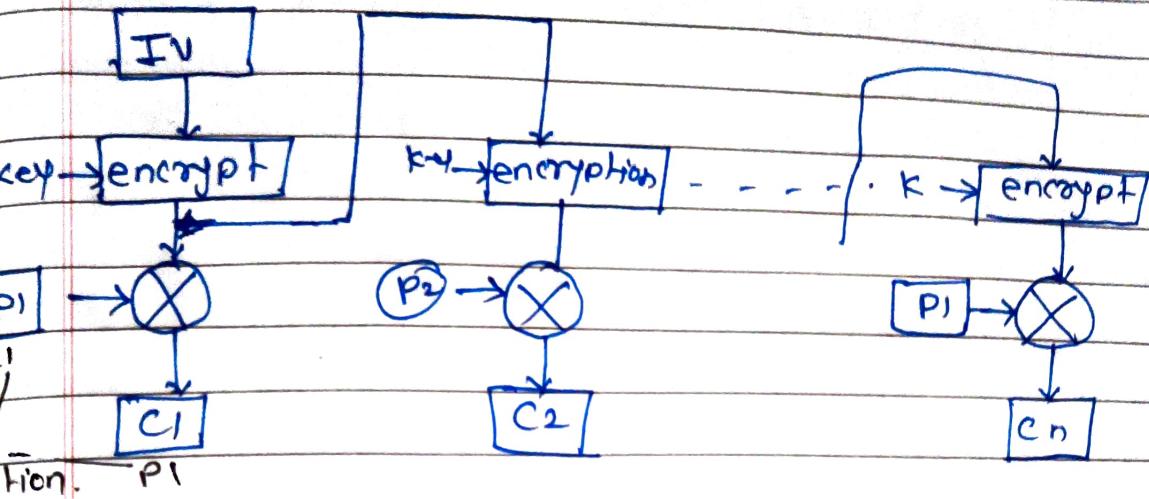
Decryption.



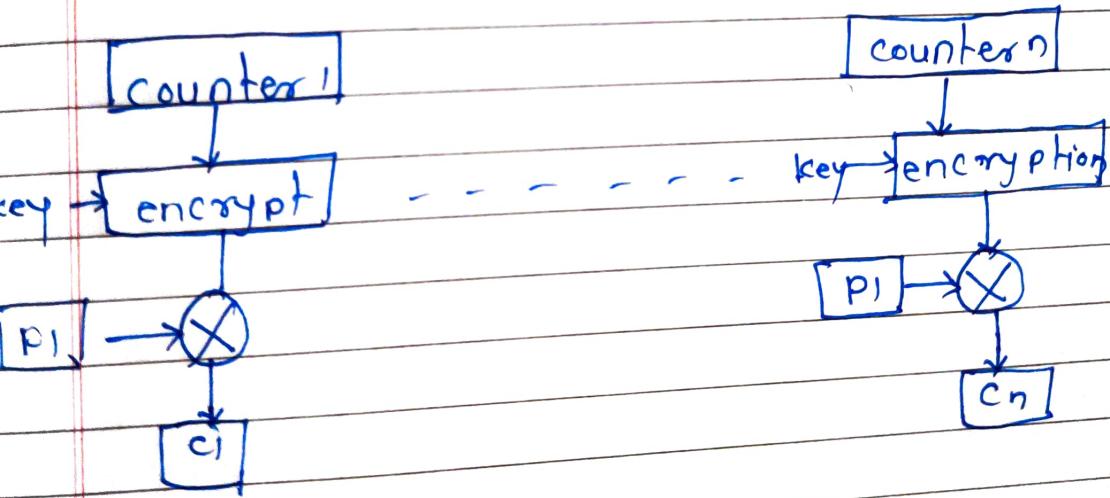
* Advantages-

- since there is some data loss due to the use of shift register, thus it is difficult for applying cryptanalysis.

④ output feedback mode.
encryption and decryption.



⑤ counter mode :-
encryption and decryption.



Advantages:

- i) plaintext can be many different ciphertexts.