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CS610 Assignment 2

1 Problem 1

Performance Bugs Identified:

False Sharing

The struct tracker has word_count array whose size is equal to the number of threads. In the execution of the statement tracker.word_count[thread_id]++; (Line 142 in reference.cpp), false sharing is present which is clear from the code as well as from the generated PERF report.

To remove this false sharing, padding was done in the word_count array and its size was made equal to NUM_THREADS × 8 and instead of accessing the index thread_id, thread_id*8 was accessed by thread having tid equal to thread_id. This multiplication by 8 ensured that the memory location accessed by each thread is on separate cache line (8 uint_64 of 8 Bytes each take 64 Bytes, which is equal to the cache line size) and thus false sharing is minimised. The source file padded.cpp has these modifications.

True Sharing

Taking and releasing lock for every word and line of the input file caused true sharing and was the cause of large number of HITMs (Lines 144-146 and 133-135 in reference.cpp). This was reduced by keeping thread-private variables for the words and line counts of each thread, and then later reducing them with the count in tracker struct, thus taking lock only once. This reduced the HITMs to nearly zero. The source file improved.cpp has these modifications. Also we have removed the redundant mutex line_count_mutex since the mutex in struct tracker is sufficient for the synchronisation needed.

HITM Shared	S	tore Re	fs	Data addr	ess				- cycles -		Total	сри	
Num RmtHitm LclH Object	Source	:Line	Node{cpu list	t}		PA cnt	Code address	rmt hitm	lcl hitm	load	records	cnt	Symb
Θ Θ 3	954 0	0	Θ	0x5d23dcb17380									
0.00% 10. eference.out thread			0.00% 0{4.8-10.15	0x0	0	1	0x5d23dcb134d1	Θ	190	169	473	5	[.] thread_runner(void*)
	43% 0.00%	0.00%		0x8	0	1	0x5d23dcb134d1	0	207	163	459	6	[.] thread_runner(void*)
	0.00%	0.00%		0x10	Θ	1	0x5d23dcb134d1	Θ	199	159	482	5	[.] thread_runner(void*)
	40% 0.00%	0.00%		0x18	0	1	0x5d23dcb134d1	0	214	178	468	6	[.] thread_runner(void*)
	38% 0.00%	0.00%		0x20	0	1	0x5d23dcb134d1	0	190	161	434	4	[.] thread_runner(void*)
	40% 0.00%	0.00%		0x28	0	1	0x5d23dcb1332a	0	187	131	257	11	[.] thread_runner(void*)
0.00% 10. Ference.out thread	12% 0.00%	0.00%	0.00%	0x30	0	1	0x5d23dcb134de	Θ	225	154	1886	11	[.] thread_runner(void*)
0.00% 16.		0.00%		0x38	Θ	1	0x75eaf4e99aa4	0	462	92	2905	11	[.] pthread_mutex_unlock@@GL
0.00% 13.	52% 0.00%	0.00%	0.00%	0x38	0	1	0x75eaf4e97f40	Θ	389	192	2746	11	[.] pthread_mutex_lock@@GLIE
0.00% 5.			0{0,3-4,6,8 0.00%	0x38	Θ	1	0x75eaf4e91294	Θ	263	102	331	11	[.]GIlll_lock_wait
	ellock.c:42		0{0,3-4,6,8·	-12,14-15}			0x75eaf4e912a3		346	138	1716		[.]GIlll_lock_wait

Figure 1: PERF Report of Reference Version - Large number of HITMs in the thread_runner function

Shar	ed Cach	e Line Di	stributio	n Pareto												
# # Share		TM		Store Ref	s		Data addı	ess				cycles -		Total	сри	
	RmtHitm	LclHitm	L1 Hit Source:L	L1 Miss	N/A cou lis	t}	0ffset	Node	PA cnt	Code address	rmt hitm	lcl hitm	load	records	cnt	Symbol
#																
0	0	695	0	0	0	0x580b	ff9154c0									
naddod ou			0.00%			.11-12.14}	0x0	Θ	1	0x580bff911325	0	391	277	228	8	[.] thread_runner(void*)
•	0.00%	0.58%	0.00%	0.00%	´ ′0.00%		0x8	Θ	1	0x580bff9114d6	0	356	373	1830	8	[.] thread_runner(void*)
	0.00%	10.65%	0.00%	0.00%	0.00%		0x10	Θ	1	0x77236b097 f 40	Θ	794	399	2078	8	[.] pthread_mutex_lock@@GLIBC_
libc.so.6	0.00%	5.04%	0.00%	0.00%	0.00%		0x10	Θ	1	0x77236b099aa4	0	404	23	1987	8	[.] pthread_mutex_unlock@@GLIB
libc.so.6	0.00%	2.88%	0.00%	0.00%	0.00%		0x10	Θ	1	0x77236b0912a3	0	575	221	306	8	[.]GIlll_lock_wait
libc.so.6	0.00%	1.29%	0.00%	0.00%	0.00%	,11-12,14} ; 1-12,14}	0x10	Θ	1	0x77236b091294	0	302	405	22	7	[.]GIlll_lock_wait
libc.so.6	0.00%	0.29%	0.00%	0.00%	΄ ΄Θ.ΘΘ%		0x18	0	1	0x77236b097f4a	0	345	369	1851	8	[.] pthread_mutex_lock@@GLIBC_
	0.00%	0.43%	0.00%	0.00%	0.00%		0x1c	Θ	1	0x77236b097f60	0	388	367	1807	8	[.] pthread_mutex_lock@@GLIBC_
	0.00%	0.14%	0.00%	0.00%	0.00%		0x1c	Θ	1	0x77236b099a8c	9	735	7	852	8	[.] pthread_mutex_unlock@@GLIB
libc.so.6	0.00%	65.18%	0.00%	0.00%	0.00%		0x20	Θ	1	0x77236b097ef4	0	402	398	1990	8	[.] pthread_mutex_lock@@GLIBC_
	0.00%	0.43%		0.00%	0.00%		0x20	0	1	0x77236b099a74	0	465	365	1836	8	[.] pthread_mutex_unlock@@GLIB
	0.00%	0.43%	0.00%	0.00%	0.00%		0x20	0	1	0x77236b099a97	0	702	383	1648	8	[.] pthread_mutex_unlock@@GLIB
libc.so.6	0.00%	0.14%	0.00%	0.00%	0.00%	,11-12,14} ; ,11-12,14}	0x20	Θ	1	0x77236b097f22	0	166	393	2037	8	[.] pthread_mutex_lock@@GLIBC_
CIDC.SO.C	peni	eau_mutex		4 010	74,0,0-9	,11-12,143										

Figure 2: PERF Report of Padded Version - Small number of HITMs in the thread_runner function, majority of them are due to locking



Figure 3: PERF Report of Improved Version - No HITMs

Version	HITMs	Time (s)
Reference	5904	1.651
Padded	1291	1.009
Improved	0	0.070

Table 1: Performance Comparison of three versions (Data taken on csews172 with 11 MB files)

Running the Code

There is a Makefile and a bash script in the code. make generates the executables for three versions. run.sh generates the PERF report for all three versions.

make

./run.sh <path to the file which has the names of input files>

The three PERF reports will be generated as perf_report_reference.out, perf_report_padded.out and perf_report_improved.out.

Also, there is a file <code>generate_input.cpp</code> which takes an integer <code>i</code> as CLI and generates a large random text file <code>inp<i>.txt</code> in the <code>test1</code> directory. <code>generate_input.out</code> is also generated by the <code>make</code> command. Ensure that <code>test1</code> directory is present before running <code>generate_input.out</code>.

2 Problem 2

The directory problem2-dir contains the source files. Below is the compilation and execution commands -

g++ problem2.cpp -lpthread
./a.out -inp=<input_path> -thr=<num_producers> -lns=<lines_per_thread> -buf=<shared_buffer_size>
-out=<output_path>

3 Problem 3

for
$$i = 1$$
, $N - 2$
for $j = i + 1$, N
A (i, $j - i$) = A (i, $j - i - 1$) - A (i + 1, $j - i$) + A (i - 1, $i + j - 1$)

We consider the following pairs and type of dependencies:

1. A(i, j - i) and A(i, j - i - 1)

$$i_0 = i_0 + \Delta i \implies \boxed{\Delta i = 0}$$

$$j_0 - i_0 = j_0 + \Delta j - i_0 - \Delta i - 1 \implies \Delta j = \Delta i + 1 \implies \boxed{\Delta j = 1}$$

Thus the flow dependency is (0,1) or (0,+), which is a valid flow dependency. Since, we have a valid flow dependency, the corresponding anti-dependency would be invalid.

2. A(i, j - i) and A(i + 1, j - i)

$$i_0 = i_0 + \Delta i + 1 \implies \boxed{\Delta i = -1}$$

$$j_0 - i_0 = j_0 + \Delta j - i_0 - \Delta i \implies \Delta j = \Delta i \implies \boxed{\Delta j = -1}$$

The flow dependency is (-1, -1), but it is an invalid dependency since the first non-zero component of the vector is negative. However, there will be an anti-dependency with the vector equal to (1, 1).

3. A(i, j - i) and A(i - 1, i + j - 1)

$$i_0 = i_0 + \Delta i - 1 \implies \boxed{\Delta i = 1}$$

$$j_0 - i_0 = i_0 + \Delta i + j_0 + \Delta j - 1 \implies \boxed{\Delta j = -2i_0}$$

In this case, Δi is positive and $\Delta j = -2i_0$. Again since i_0 is always positive, Δj will be always negative. Thus the dependency vector will be (+,-). It is a valid flow dependency and so there is no anti-dependency.

There is no output dependency for any of the memory access pairs.