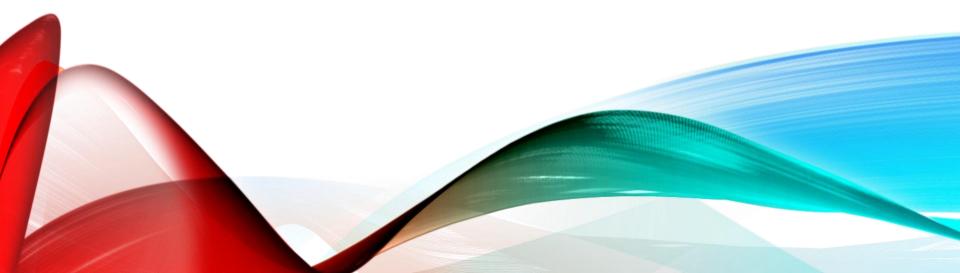


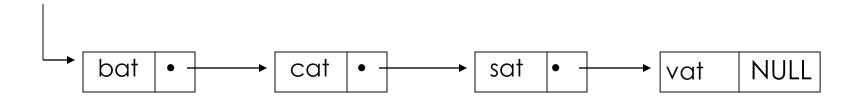
Chapter 04: Lists



INTRODUCTION

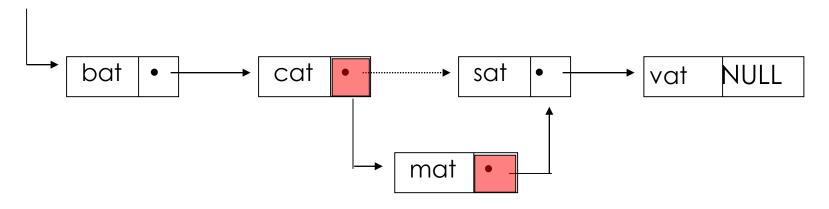
- Array successive items locate a fixed distance
- disadvantage
 - data movements during insertion and deletion
 - waste space in storing n ordered lists of varying size
- possible solution
 - linked list

SINGLY LINKED LISTS

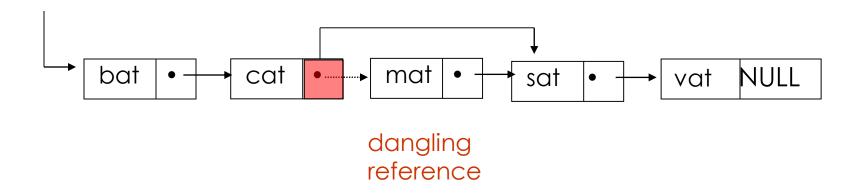


4

Insertion



Delete mat from list



EXAMPLE 4.1: CREATE A LINKED LIST OF WORDS

Declaration

7

REFERENCING THE FIELDS OF A NODE

EXAMPLE: CREATE A TWO-NODE LIST

```
Example 4.2
ptr
                                                  typedef struct list_node *list_pointer;
                                                  typedef struct list_node {
                                                            int data;
                                                            list_pointer link;
        10
                                    NULL
                              20
                                                  list_pointer ptr =NULL
            list pointer create2()
             /* create a linked list with two nodes */
               list_pointer first, second;
first = (list_pointer) malloc(sizeof(list_node));
               second = ( list_pointer) malloc(sizeof(list_node));
second -> link = NULL;
               second -> data = 20;
```

*Program 4.2:Create a two-node list

first -> data = 10; first ->link = second;

return first:

USE OF A POINTER (1)

int i, *pi;

pi = &i;

2000 pi 1000

i = 10 or *pi = 10

i 1000 *pi 10 2000 pi 1000

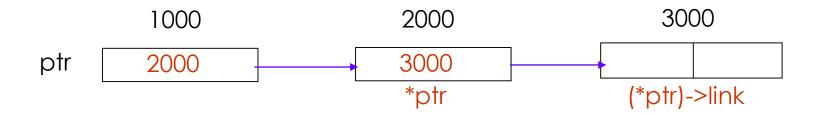
USE OF A POINTER (2)

```
typedef struct list_node *list_pointer;
typedef struct list_node {
              int data;
              list_pointer link;
list_pointer ptr = NULL;
           1000
                                        ptr->data ← (*ptr).data
ptr
           NULL
ptr = malloc(sizeof(list_node));
            1000
                                 2000
                                           *ptr
 ptr
            2000
                                  data
                                           link
```

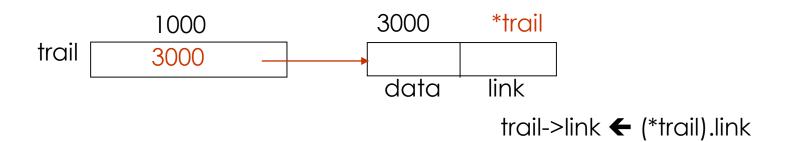
USE OF A POINTER (3)

void delete(list_pointer *ptr, list_pointer trail, list_pinter node)

ptr: a pointer point to a pointer point to a list node

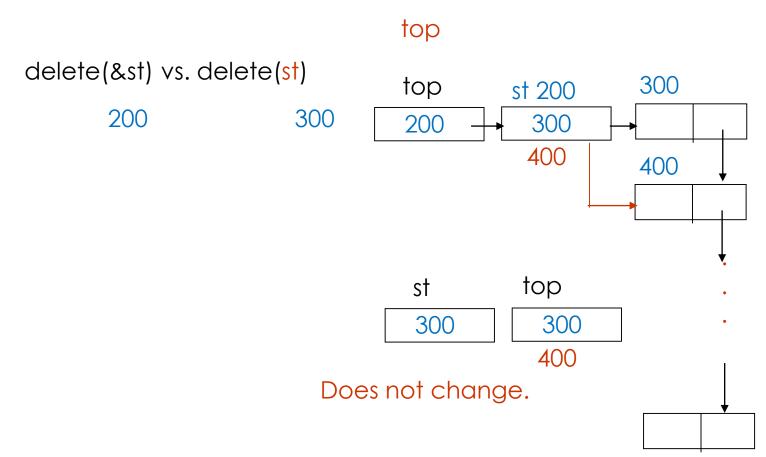


trail (node): a pointer point to a list node



USE OF A POINTER (4)

element delete(stack_pointer *top)

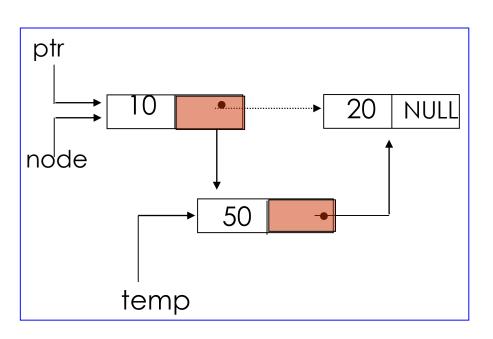


LIST INSERTION

Insert a node after a specific node

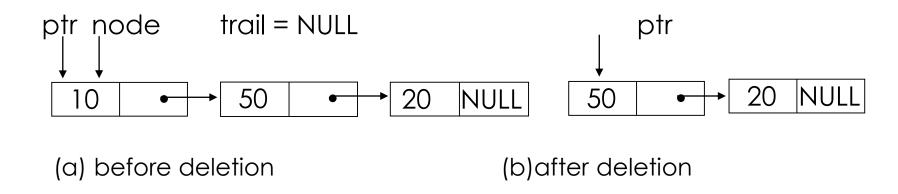
```
void insert(list_pointer *ptr, list_pointer node)
{
/* insert a new node with data = 50 into the list ptr after node */
    list_pointer temp;
    temp = (list_pointer) malloc(sizeof(list_node));
    if (IS_FULL(temp)){
        fprintf(stderr, "The memory is full\n");
        exit (1);
    }
```

LIST INSERTION



*Program 4.3:Simple insert into front of list

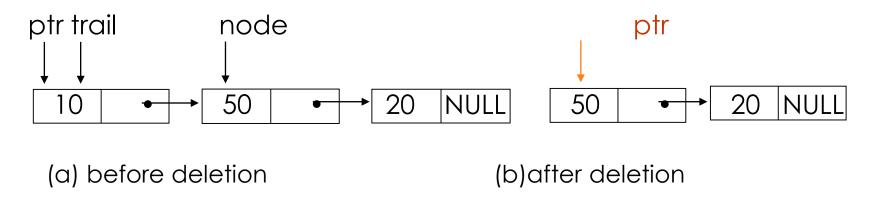




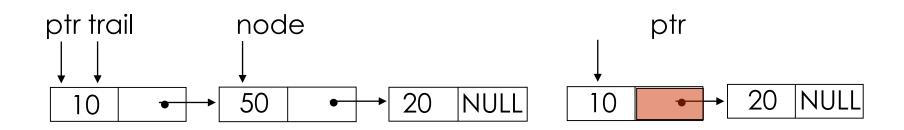
List after the function call Delete (&ptr, NULL.ptr)

List Deletion

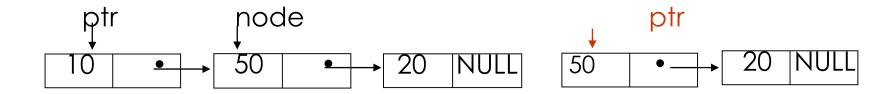
Delete the first node.



Delete node other than the first node.



17

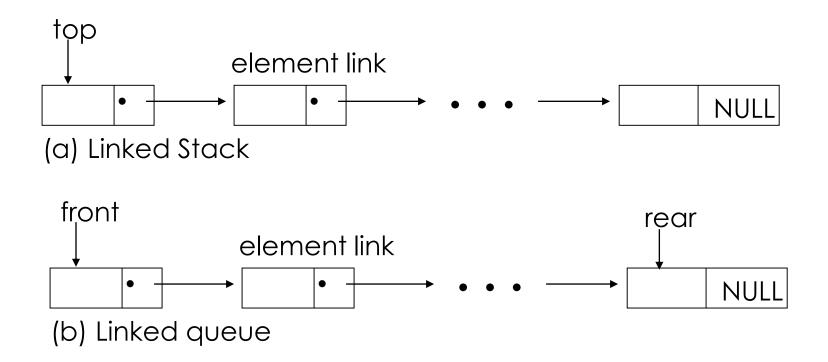


Print out a list (traverse a list)

```
void print_list(list_pointer ptr)
{
    printf("The list ocntains: ");
    for (; ptr; ptr = ptr->link)
        printf("%4d", ptr->data);
    printf("\n");
}
```

*Program 4.5: Printing a list

DYNAMICALLY LINKED STACKS AND QUEUES



*Figure 4.10: Linked Stack and queue

Represent n stacks

```
#define MAX_STACKS 10 /* maximum number of stacks */
typedef struct {
    int key;
    /* other fields */
    } element;
typedef struct stack *stack_pointer;
typedef struct stack {
    element item;
    stack_pointer link;
    };
stack_pointer top[MAX_STACKS];
```

Represent n queues

Push in the linked stack

pop from the linked stack

```
element delete(stack_pointer *top) {
/* delete an element from the stack */
  stack_pointer temp = *top;
  element item;
  if (IS_EMPTY(temp)) {
    fprintf(stderr, "The stack is empty\n");
    exit(1);
  item = temp->item;
  *top = temp->link;
  free(temp);
  return item;
*Program 4.7: Delete from a linked stack
```

ENQUEUE IN THE LINKED QUEUE

```
void addq(queue_pointer *front, queue_pointer *rear, element
item)
{ /* add an element to the rear of the queue */
 queue_pointer temp =
              (queue_pointer) malloc(sizeof (queue));
 if (IS_FULL(temp)) {
   fprintf(stderr, "The memory is full\n");
   exit(1);
   temp->item = item;
   temp->link = NULL;
   if (*front) (*rear) -> link = temp;
   else *front = temp;
   *rear = temp; }
```

DEQUEUE FROM THE LINKED QUEUE (SIMILAR TO PUSH)

```
element deleteq(queue_pointer *front) {
/* delete an element from the queue */
  queue_pointer temp = *front;
  element item:
  if (IS_EMPTY(*front)) {
    fprintf(stderr, "The queue is empty\n");
    exit(1);
  item = temp->item;
  *front = temp->link;
  free(temp);
  return item;
```

Polynomials

$$A(x) = a_{m-1}x^{e_{m-1}} + a_{m-2}x^{e_{m-2}} + ... + a_0x^{e_0}$$

Representation

```
typedef struct poly_node *poly_pointer;
typedef struct poly_node {
    int coef;
    int expon;
    poly_pointer link;
};
poly pointer a, b, c;
```

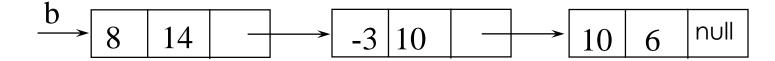
coef expon	link
------------	------

Examples

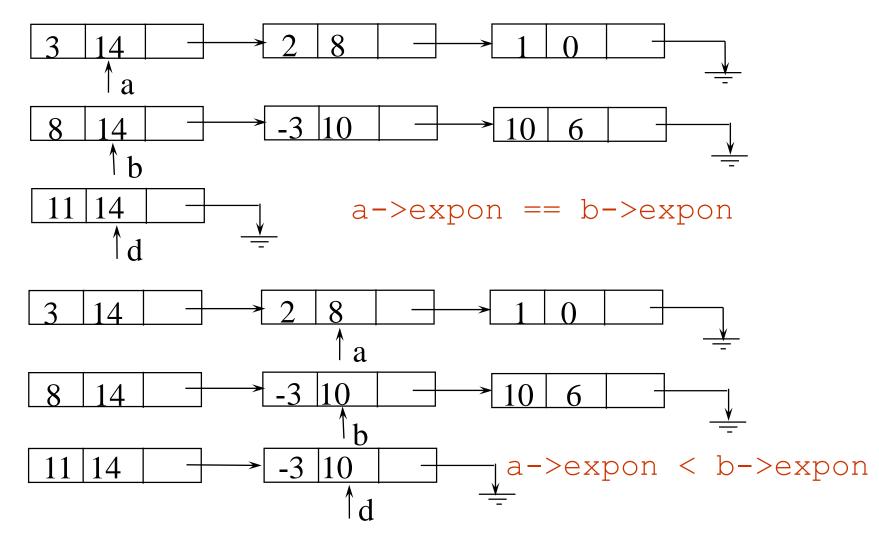
$$a = 3x^{14} + 2x^8 + 1$$



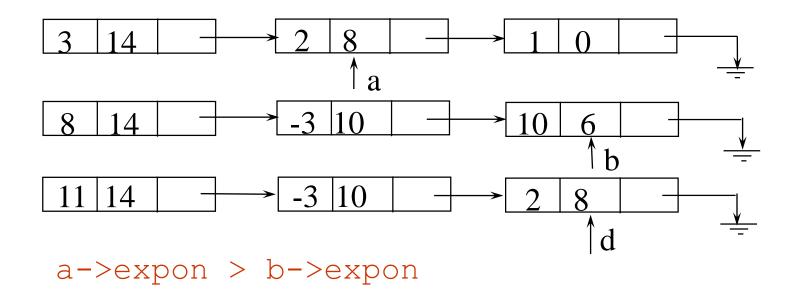
$$b = 8x^{14} - 3x^{10} + 10x^6$$



Adding Polynomials



ADDING POLYNOMIALS (CONTINUED)



ALOGRITHM FOR ADDING POLYNOMIALS

```
poly_pointer padd(poly_pointer a, poly_pointer b)
{
    poly_pointer front, rear, temp;
    int sum;
    rear = (poly_pointer)malloc(sizeof(poly_node));
    if (IS_FULL(rear)) {
        fprintf(stderr, "The memory is full\n");
        exit(1);
    }
    front = rear;
    while (a && b) {
        switch (COMPARE(a->expon, b->expon)) {
```

```
case -1: /* a->expon < b->expon */
             attach(b->coef, b->expon, &rear);
             b= b->link;
             break;
        case 0: /* a->expon == b->expon */
             sum = a->coef + b->coef;
             if (sum) attach(sum,a->expon,&rear);
             a = a \rightarrow link; b = b \rightarrow link;
             break;
        case 1: /* a->expon > b->expon */
             attach(a->coef, a->expon, &rear);
             a = a - > link;
for (; a; a = a \rightarrow link)
    attach(a->coef, a->expon, &rear);
for (; b; b=b->link)
    attach(b->coef, b->expon, &rear);
rear->link = NULL;
temp = front; front = front->link; free(temp);
return front;
```

Delete extra initial node.

ANALYSIS

- (1) coefficient additions
 0 ≤ additions ≤ min(m, n)
 where m (n) denotes the number of terms in A (B).
- (2) exponent comparisons extreme case $e_{m-1} > f_{m-1} > e_{m-2} > f_{m-2} > ... > e_0 > f_0$ m+n-1 comparisons
- (3) creation of new nodes extreme case

 m + n new nodes
 summary O(m+n)

ATTACH A TERM

```
void attach (float coefficient, int exponent,
            poly pointer *ptr)
/* create a new node attaching to the node pointed to
  by ptr. ptr is updated to point to this new node. */
    poly pointer temp;
    temp = (poly pointer) malloc(sizeof(poly node));
    if (IS FULL(temp)) {
        fprintf(stderr, "The memory is full \n");
        exit(1);
    temp->coef = coefficient;
    temp->expon = exponent;
    (*ptr) - > link = temp;
    *ptr = temp;
```

A SUITE FOR POLYNOMIALS

```
e(x) = a(x) * b(x) + d(x)

poly_pointer a, b, d, e;

a = read_poly();

b = read_poly();

d = read_poly();

temp = pmult(a, b);

e = padd(temp, d);

print poly()

pmult()
```

temp is used to hold a partial result. By returning the nodes of temp, we may use it to hold other polynomials

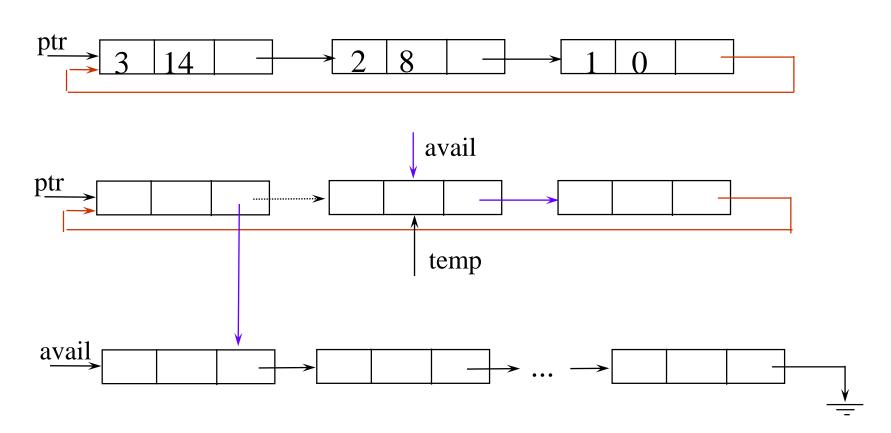
ERASE POLYNOMIALS

```
void earse(poly_pointer *ptr)
{
/* erase the polynomial pointed to by ptr */
    poly_pointer temp;
    while (*ptr) {
        temp = *ptr;
        *ptr = (*ptr)->link;
        free(temp);
    }
}
```

O(n)

CIRCULARLY LINKED LISTS

circular list vs. chain



MAINTAIN AN AVAILABLE LIST

```
poly pointer get node (void)
  poly_pointer node;
  if (avail)
      node = avail;
      avail = avail->link:
  else {
      node = (poly pointer) malloc(sizeof(poly node));
      if (IS FULL (node))
          prīntf(stderr, "The memory is full\n");
          exit(1);
  return node;
```

MAINTAIN AN AVAILABLE LIST

(CONTINUED)

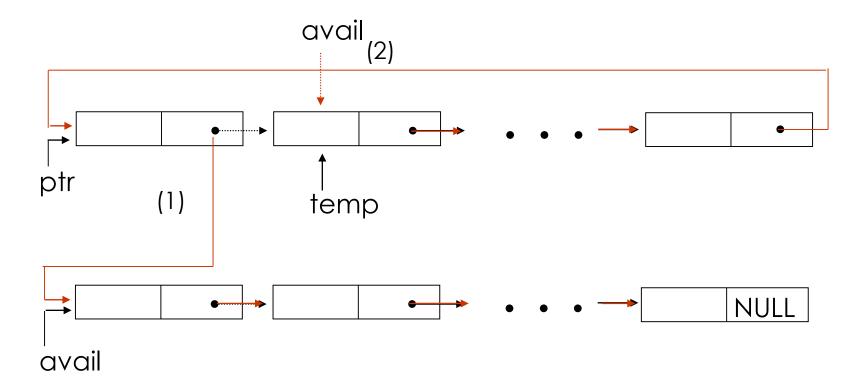
Insert ptr to the front of this list

```
void ret_node(poly_pointer ptr)
{
  ptr->link = avail;
  avail = ptr;
}

void cerase(poly_pointer *ptr)
{
  poly pointer temp;
  if (*ptr) {
    temp = (*ptr)->link;
        (*ptr)->link = avail;
        avail = temp;
        *ptr = NULL;
  }
}
(1)
```

Independent of # of nodes in a list O(1) constant time

REPRESENTING POLYNOMIALS AS CIRCULARLY LINKED LISTS



*Figure 4.14: Returning a circular list to the avail list (p.159)

Head Node

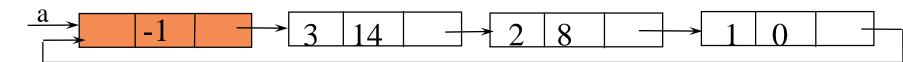
Represent polynomial as circular list.

(1) zero



Zero polynomial

(2) others



$$a = 3x^{14} + 2x^8 + 1$$

Another Padd

```
poly_pointer cpadd(poly_pointer a, poly_pointer b)
{
   poly_pointer starta, d, lastd;
   int sum, done = FALSE;
   starta = a;
   a = a->link;
   b = b->link;
   d = get_node();
   d->expon = -1;   lastd = d;
   do {
      switch (COMPARE(a->expon, b->expon)) {
        case -1: attach(b->coef, b->expon, &lastd);
        b = b->link;
        break;
```

Another Padd (Continued)

```
case 0: if (starta == a) done = TRUE;
    else {
        sum = a->coef + b->coef;
        if (sum) attach(sum,a->expon,&lastd);
        a = a->link; b = b->link;
        break;
    case 1: attach(a->coef,a->expon,&lastd);
        a = a->link;
    }
} while (!done);
lastd->link = d;
return d;
Link last node to first
```

Additional List Operations

```
typedef struct list_node *list_pointer;
typedef struct list_node {
    char data;
    list_pointer link;
};
```

Invert single linked lists
Concatenate two linked lists

INVERT SINGLE LINKED LISTS

Use two extra pointers: middle and trail.

```
list pointer invert(list pointer lead)
    list pointer middle, trail;
    middTe = NULL;
    while (lead)
        trail = middle;
        middle = lead;
        lead = lead->link;
        middle->link = trail;
    return middle;
                    0: null
                    1: lead
                    ≥2: lead
```

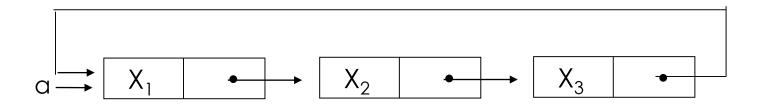
Concatenate Two Lists

```
list pointer concatenate(list pointer
             ptr1, list pointer ptr2)
  list pointer temp;
  if (IS EMPTY(ptr1)) return ptr2;
  else {
    if (!IS EMPTY(ptr2)) {
      for (temp=ptr1;temp->link;temp=temp->link);
      temp->link = ptr2;
    return ptr1;
```

O(m) where m is # of elements in the first list

OPERATIONS FOR CIRCULARLY LINKED LIST

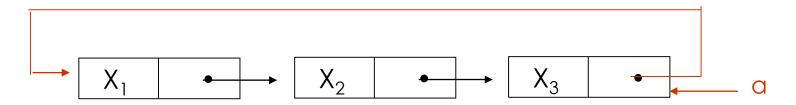
What happens when we insert a node to the front of a circular linked list?



Problem: move down the whole list.

*Figure 4.16: Example circular list (p.165)

A possible solution:



Note a pointer points to the last node.

*Figure 4.17: Pointing to the last node of a circular list

ptr

OPERATIONS FOR CIRCULAR LINKED LISTS

 X_3

```
void insert front (list pointer *ptr, list pointer
node)
    if (IS EMPTY(*ptr)) {
       *ptr= node;
       node->link = node;
    else {
        node - > link = (*ptr) - > link; (1)
         (*ptr) - > link = node;
                                        (2)
```

 X_2

(1)

(2)

Length of Linked List

```
int length(list_pointer ptr)
{
    list_pointer temp;
    int count = 0;
    if (ptr) {
        temp = ptr;
        do {
            count++;
            temp = temp->link;
        } while (temp!=ptr);
    }
    return count;
}
```

Equivalence Relations

A relation over a set, S, is said to be an *equivalence relation* over S *iff* it is symmetric, reflexive, and transitive over S.

```
reflexive, x=x
symmetric, if x=y, then y=x
transitive, if x=y and y=z, then x=z
```

Examples

three equivalent classes {0,2,4,7,11}; {1,3,5}; {6,8,9,10}

A ROUGH ALGORITHM TO FIND EQUIVALENCE CLASSES

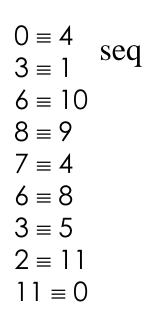
```
void equivalenec()
    initialize;
    while (there are more pairs) {
Phase
        read the next pair <i, j>;
        process this pair;
    initialize the output;
    do
        output a new equivalence class;
    } while (not done);
2
```

What kinds of data structures are adopted?

First Refinement

```
#include <stdio.h>
#include <alloc.h>
#define MAX SIZE 24
#define IS FULL(ptr) (!(ptr))
#define FATSE 0
#define TRUE 1
yoid equivalence()
          initialize seq to NULL and out to TRUE
while (there are more pairs) {
    read the next pair, <i,j>;
    put j on the seq[i] list;
    put i on the seq[j] list;
                      (i=0; i<n; i++)
if (out[i]) {
    out[i]= FALSE;
                                                                                            direct equivalence
                                output this equivalence class;
                                                                         Compute indirect equivalence
                                                                         using transitivity
```

Lists After Pairs are input



```
[0]
                       [4] [5]
                                        [7] [8] [9] [10] [11]
                [3]
                                   [6]
      [1]
            [2]
                                                   NULL
                                                         NULL
      NULL
           NULL
                             NULL
                                        NULL
                                   10
NULL
                 NULL
                       NULL
                                                               NULL
                                  NULL
                                              NULL
```

```
typedef struct node *node_pointer;
typedef struct node {
    int data;
    node_pointer link;
};
```

FINAL VERSION FOR FINDING EQUIVALENCE CLASSES

```
void main(void)
  short int out[MAX SIZE];
  node pointer seq[MAX SIZE];
  node pointer x, y, top;
  int \overline{i}, j, n;
  printf ("Enter the size (<= %d) ", MAX SIZE);
  scanf("%d", &n);
  for (i=0; i< n; i++) {
      out[i] = TRUE; seq[i] = NULL;
  printf("Enter a pair of numbers (-1 -1 to quit): ");
  scanf("%d%d", &i, &j);
      Phase 1: input the equivalence pairs:
```

```
while (i \ge 0) {
    x = (node pointer) malloc(sizeof(node));
    if (IS FULL(x))
Insert x to the top of lists seq[i] full\n");
        exit(1);
    x->data= j; x->link= seq[i]; seq[i]= x;
    if (IS FULL(x))
      finsert * to the top of lists sequils full \n");
        exit(1);
    x->data= i; x->link= seq[j]; seq[j]= x;
    printf("Enter a pair of numbers (-1 -1 to \
           quit): ");
    scanf("%d%d", &i, &j);
```

Phase 2: output the equivalence classes

```
for (i=0; i< n; i++) {
     if (out[i]) {
           printf("\nNew class: %5d", i);
           out[i] = FALSE;
           x = seq[i]; top = NULL;
           for (;; ) {
                 while (x) {
                                                          push
                      j = x->data;
if (out[j]) {
    printf("%5d", j);
    out[j] = FALSE;
    y = x->link; x->link = top;
    top = x; x = y;
    pop
                                                         pop
                       else x = x - \sinh x;
                 if (!top) break;
                 x = seq[top->data]; top = top->link;
```

SPARSE MATRICES

$$\begin{bmatrix} 0 & 0 & 11 & 0 \\ 12 & 0 & 0 & 0 \\ 0 & -4 & 0 & 0 \\ 0 & 0 & 0 & -15 \end{bmatrix}$$

inadequates of sequential schemes

- (1) # of nonzero terms will vary after some matrix computation
- (2) matrix just represents intermediate results

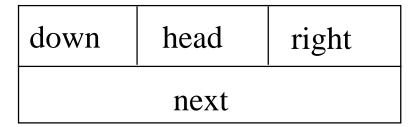
new scheme

Each column (row): a circular linked list with a head node

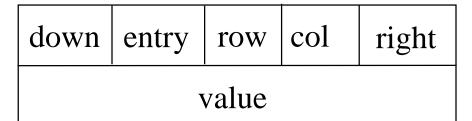
Revisit Sparse Matrices

of head nodes = max{# of rows, # of columns}

head node

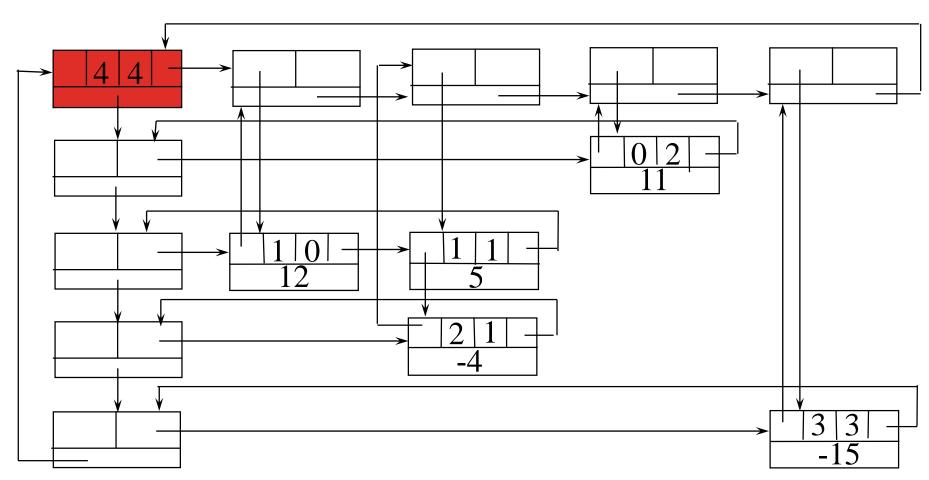


entry node



aij entry i j — aij

Linked Representation for Matrix



```
#define MAX_SIZE 50 /* size of largest matrix */
typedef enum {head, entry} tagfield;
typedef struct matrix_node *matrix_pointer;
typedef struct entry_node {
    int row;
    int col;
    int value;
    };
typedef struct matrix_node {
    matrix_pointer down;
    matrix_pointer right;
    tagfield tag;
```

6

	[0]	[1]	[2]
[0]	4	4	4
[1]	0	2	11
[2]	1	0	12
[3]	2	1	-4
[4]	3	3	-15

*Figure 4.22: Sample input for sparse matrix

Read in a Matrix

```
scanf("%d%d%d", &num rows, &num cols,
     &num terms);
num heads =
 (num cols>num rows)? num cols : num rows;
/* set up head node for the list of head
   nodes */
node->u.entry.row = num rows;
node->u.entry.col = num_cols;
if (!num heads) node->right = node;
else { /* initialize the head nodes */
  for (i=0; i<num heads; i++) {
    term= new nod\overline{e}();
    hdnode[i] = temp;
    hdnode[i]->tag = head;
    hdnode[i]->right = temp;
                                        O(max(n,m))
    hdnode[i]->u.next = temp;
```

```
current row= 0; last= hdnode[0];
for (i=0; i< num terms; i++) {
 printf("Enter row, column and value:");
  scanf("%d%d%d", &row, &col, &value);
  if (row>current row) {
    last->right= The Todo [current row];
    current row = row; last=hdnode[row];
  temp = new node();
  temp->tag=entry; temp->u.entry.row=row;
  temp->u.entry.col = col;
 temp->u.entry.value = value;
  last->right = temp; /*link to row list */
  last= temp;
 /* link to column list */
 hdnode[col]->u.next->down = temp;
 hdnode[col]=>u.next = temp;
    利用next field 存放column的last node
```

```
/*close last row */
  last->right = hdnode[current row];
  /* close all column lists */
  for (i=0; i < num cols; i++)
    hdnode[i]->u.next->down = hdnode[i];
  /* link all head nodes together */
  for (i=0; i < num heads-1; i++)
    hdnode[i]->u.\overline{n}ext = hdnode[i+1];
  hdnode[num heads-1]->u.next= node;
  node - > right = hdnode[0];
return node;
        O(max\{\#_rows, \#_cols\}+\#_terms)
```

Write out a Matrix

```
void mwrite(matrix pointer node)
{ /* print out the matrix in row major form */
  int i;
  matrix pointer temp, head = node->right;
  printf("\n num rows = %d, num cols= %d\n",
         node->u.entry.row, node->u.entry.col);
  printf ("The matrix by row, column, and
         value:\n\n");
O(#_rows+#_terms)
  for (i=0; i< node->u.entry.row; i++) {
    for (temp=head->right; temp!=head; temp=temp->right)
      printf("%5d%5d%5d\n", temp->u.entry.row,
           temp->u.entry.col, temp->u.entry.value);
    head= head->u.next; /* next row */
```

Free the entry and head nodes by row.

Erase a Matrix

```
void merase(matrix pointer *node)
  int i, num heads;
  matrix poi\overline{n}ter x, y, head = (*node)->right; for (i\equiv 0; i<(*node)->u.entry.row; i++) {
    y=head->right;
    while (y!=head) {
       x = y; y = y - > right; free (x);
    x= head; head= head->u.next; free(x);
  y = head;
  while (y!=*node) {
    x = y; y = y->u.next; free(x);
  free(*node); *node = NULL;
  O(#_rows+#_cols+#_terms)
```

Doubly Linked List

Move in forward and backward direction.

Singly linked list (in one direction only)
How to get the preceding node during deletion or insertion?
Using 2 pointers

Node in doubly linked list left link field (llink) data field (item) right link field (rlink)

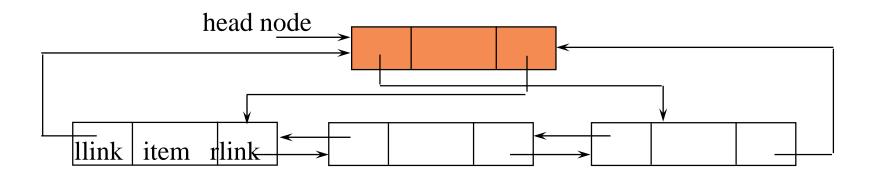
Doubly Linked Lists

```
typedef struct node *node_pointer;

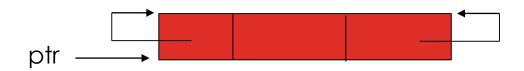
typedef struct node {
    node_pointer llink;
    element item;
    node_pointer rlink;

    rode_pointer rlink;

    rode_pointer rlink;
```

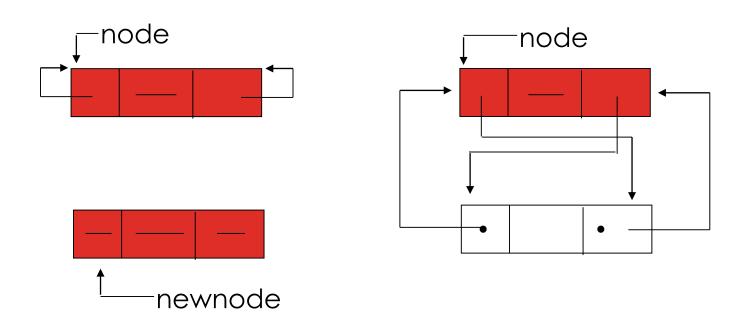






*Figure 4.24:Empty doubly linked circular list with head node (p.1



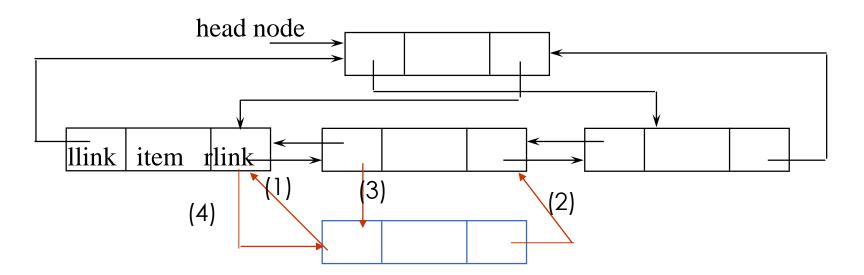


*Figure 4.25: Insertion into an empty doubly linked circular list (p.18

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Insert

```
void dinsert(node_pointer node, node_pointer newnode)
{
    (1) newnode->llink = node;
    (2) newnode->rlink = node->rlink;
    (3) node->rlink->llink = newnode;
    (4) node->rlink = newnode;
}
```



Delete

