

# Log geometry (Amand P.)

05-06-14

(1)

Essentially, generalization of scheme theory.

Basic objects:  $\underline{X}$  scheme underlying +  $\alpha: M_X \rightarrow \mathcal{O}_X^*$

PRELOG STRUCTURE

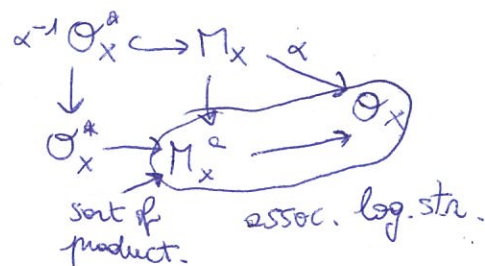
sheaf of monoids  
mult. monoid  
monoid map.  
D. Abr.: "special dust over  $X$ "

comm. mon. id.

def. log-structures are pre-log structures s.t.  $\alpha^{-1}(\mathcal{O}_X^*) \rightarrow \mathcal{O}_X^*$

(idea: things that have more than 1 preimage have to vanish; sort of the branches of the logarithm)

Fact pre-log str. gives an associated log-str.



Example  $\underline{X}$  smooth,  $D \subset X$  divisor

$$M_X(U) = \{g \in \mathcal{O}_X(U) \mid g|_{U \setminus D} \in \mathcal{O}_X^*(U \setminus D)\} \subset \mathcal{O}_X(U) \quad (\text{the case } D \text{ normal crossing is going to be "log smooth"})$$

This is very related to logarithmic differentials, that in fact is the motivation for the theory.

Example (building blocks)  $P$  monoid,  $R$  a ring  $\Rightarrow R[P]$  ring, then

$$\underline{X} = \text{spec}(R[P]), \alpha: P \rightarrow R[P]$$

e.g.  $P = \mathbb{N}, R = \mathbb{Z}, \underline{X} = \text{spec} \mathbb{Z}[t] = \mathbb{A}_{\mathbb{Z}}^1$

$$\alpha: P \rightarrow R[P] \quad \begin{matrix} e_1 \mapsto t \end{matrix} \quad \left( \begin{matrix} \text{is as before } D \subset X \\ \text{where } D = \{0\} \\ X = \mathbb{A}_{\mathbb{Z}}^1 \end{matrix} \right)$$

Morphisms  $f: \underline{X} \rightarrow \underline{Y}$  morph. of Schemes,  $\alpha_Y: M_Y \rightarrow \mathcal{O}_Y^*$  log str. on  $Y$ .

we can pull it back to  $f^*(M_Y) \rightarrow f^*(\mathcal{O}_Y^*) \rightarrow \mathcal{O}_X^* \rightsquigarrow$  call it  $f^* M_Y \rightarrow \mathcal{O}_X^*$

flat "inverse image"

A map between  $(\underline{X}, M_X)$  and  $(\underline{Y}, M_Y)$  is  $f: \underline{X} \rightarrow \underline{Y}$  and  $f^* M_Y \rightarrow M_X$

$\downarrow \quad \downarrow$   
 $f^*(\mathcal{O}_Y^*) \rightarrow \mathcal{O}_X^*$

Building blocks of morphisms two monoids  $\partial: Q \rightarrow P$ ,  $R$  ring, (2)  
 we get  $R[Q] \rightarrow R[P] \rightsquigarrow$  we get morph. of log schemes  $(\text{spec } R[P], P) \rightarrow (\text{spec } R[Q], Q)$

There is a lot of tower stuff going on.

"The log structures-morphisms that are made (etale locally) by these building blocks, are locally tower varieties/morphisms."

Ex.  $\underline{X} = \text{Spec } \frac{k[x_1, \dots, x_n]}{(x_1 \cdot x_2 \cdot \dots \cdot x_n)} \quad \mathbb{N}^n \rightarrow \mathcal{O}_X$   
 $e_i \mapsto x_i$

$p = \text{spec}(k) \xrightarrow[\underline{P}]{\text{origin}} \underline{X} \quad f^*(M_X) = k^* \oplus \mathbb{N}^n \rightarrow \mathcal{O}_p = k$   
 $(a, v) \mapsto a \text{ if } v=0$   
 $0 \text{ if } v \neq 0$   
 this  $k^*$  comes from "logification" of the prelog  $(\mathbb{N}^n \rightarrow \underline{X})$   
 this is called a log pt of  $\text{rk } n$ .  
standard when  $n=1$

Let's restrict  $\alpha: M_X \rightarrow \mathcal{O}_X$ , because otherwise there are too many pathologies.

Restrictions coming from geometry.

•  $P$  monoid is integral if  $P \rightarrow P^{gp}$  is injective

•  $P$  is fine if it is fin. gen.

•  $P$  saturated if  $m \cdot p \in P (p \in P^{gp}) \Rightarrow p \in P$ .

Ex.  $P = \mathbb{N} \setminus \{1\}, \text{spec}(P \rightarrow k[P]) = \text{spec } k[x^2, x^3]$   
 cusp.  
 here  $1 \in P^{gp} = \mathbb{Z}$   
 so this is not saturated.

Chart  $(\underline{X}, M_X)$  a log scheme,  $P$  monoid. A chart for  $M_X$  is

$\xrightarrow[\text{chart}]{\text{chart}} P_X \rightarrow M_X \xrightarrow{\alpha} \mathcal{O}_X$  s.t.  $P_X^a \rightarrow \mathcal{O}_X$  is isom. to  $M_X$

Similarly  $f: \underline{X} \rightarrow \underline{Y}$  a map. A chart is  $Q, P, Q \xrightarrow{\partial} P$  where  $\begin{matrix} P_X \rightarrow M_X \\ \uparrow \partial \quad \uparrow f^* \\ Q_X \rightarrow M_Y \end{matrix}$

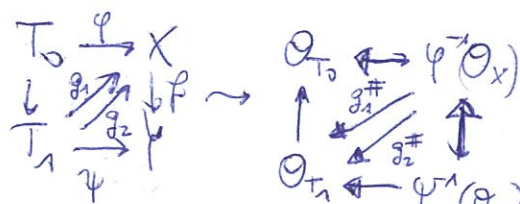
Differentials (Log smooth curves)

(Grothendieck point of view on derivations)

$T_1$  nilpot. thickening of  $T_0$  with  $T_1^2 = 0$   
 $T_0 \xrightarrow{\varphi} X \quad \downarrow f$   
 $\downarrow \rightsquigarrow \quad \downarrow$   
 $T_1 \rightarrow Y$   
 we want to define  $\Omega_{X/Y}^1$  (enhancement of  $\Omega_{X/Y}^1$ , still an  $\mathcal{O}_X$ -module)  
 $(X, Y \text{ log schemes } T_0, T_1 \text{ schemes})$



in the case of  $f$  we consider two liftings  $g_1, g_2$



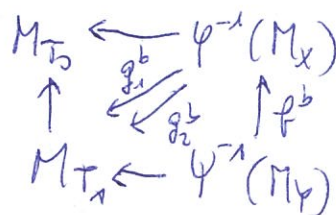
(3)

$\leadsto g_1^\# - g_2^\# : \varphi^{-1}(\mathcal{O}_X) \rightarrow J$  derivation in  $\text{Der}_Y(X, \mathcal{O}_X)$   
 $\nearrow \frac{1}{f}$   $\searrow \frac{1}{\psi}$  "d<sub>g<sub>1</sub>-g<sub>2</sub></sub>" we call it here  $J$  in general but often  $J \cong \mathcal{O}_X$  over  $X$ .

Now consider the log situation, we want  $j^* M_{T_1} \cong M_{T_0}$  so that

$$1 \rightarrow 1+J \rightarrow M_{T_1} \rightarrow M_{T_0} \rightarrow 1$$

$$g_1^b - g_2^b : \varphi^{-1} M_X \rightarrow J \text{ ADDITIVE, such that}$$



(1)  $D(ab) = D(a) + D(b)$

(2)  $\alpha(a) \cdot D(a) = D(\alpha(a))$

(3)  $D(a) = 0 \quad \forall a \in \text{Im}(f^{-1}(M_Y) \rightarrow M_X)$

(2) gives  $D(a) = \frac{D(\alpha(a))}{\alpha(a)}$   
 or D acts as "dlog"

So  $g_1, g_2$  gives  $(D, D)$   
 $\nearrow$  classical derivation  $\nwarrow$  log derivatives-like with inputs in  $M_X$

Point  $(D, D)$  can be constructed locally, they define a sheaf  $\text{Der}_Y(X, J)$   
 and when  $J \cong \mathcal{O}_X$  this is dual to  $\Omega_{X/Y}^1$  (cotangent bundle)

Log smooth  $f: X \rightarrow Y$  is smooth if  $\forall$   $\exists g: T_1 \rightarrow X$  making it commute (if  $\exists!$ , then is étale)  
 sort of evaluative criterion  
 we follow these definitions in the log case.

Kato's thm  $f: X \rightarrow Y$  is log smooth (étale) if  $f$  étale locally on  $X$  there are charts  $(P, Q, \theta)$  s.t.  
 (a) kernel + torsion of coker of  $Q^{\text{gp}} \rightarrow P^{\text{gp}}$  are finite with invertible order (char  $p$  order (to define étale we need the whole cokernel)  
 (b)  $\underline{X} \rightarrow \underline{Y} \times_{\text{Spec}(\mathbb{Z}[Q])} \text{Spec}(\mathbb{Z}[P])$  needs to be étale

(b) is like asking that locally the map is a projection from a toric variety to another.

(4)

e.g.  $|N \setminus \{1\} \rightarrow P$ ,  $\text{spec } k[P]$  is log smooth!

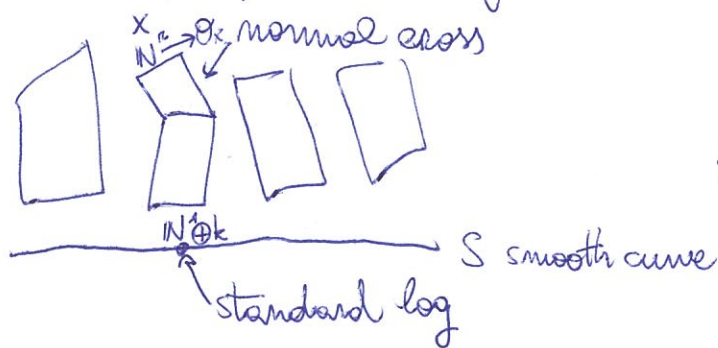
def. log curves  $f: X \rightarrow S$  gen fibers are reduced 1-dim'l connected  $X, S$  are fine, saturated, and  $f$  is log smooth.

Kato log curves are semistable

(log smooth curve is stable if it has no inf. automorphism)

$\leadsto \mathcal{M}_g$  is "collection" of stable log smooth curves.

other example

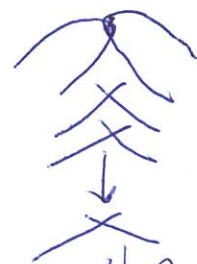


here  $X \rightarrow \text{pt}$  is still log smooth, with  $\Theta: N^1 \rightarrow N^2$   
 $e_1 \mapsto (e_1, \dots, e_1)$   
 diagonal

If  $X \rightarrow p, k^{\text{ss}} \oplus N$  does not admit a log structure of  $\Delta$  diagonal type, then  $X$  cannot be smoothed in a smooth total space.

(Classically known as d-semistability)

other example admissible covers problem:



this map is not flat!

but there are natural log structures which make  $\pi$  étale

Mochizuki's thesis is about this: using log geometry admissible covers are just log-étale maps.

What about surfaces? For K3, Ogura. Classically one tries to do semistable reduction on families  $X \rightarrow S$  and try to make fibers as less singular as possible.  
 $\leadsto$  there is mod. space of log smooth K3 surfaces